Graph and Dijkstra Algorithm

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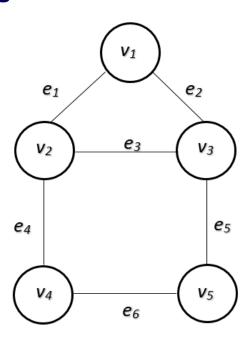
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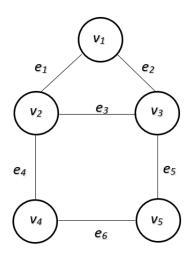
Graph

- \square A graph G is a data structure that consists of a set of vertices V and edges E
 - An edge e is a pair of vertices (v_i, v_j) , where $v_i, v_j \in V$
 - For example, the following picture shows a graph with 5 vertices and 6 edges



Adjacency Matrix

- □ If two vertices in a graph are connected by an edge, we say the vertices are adjacent
 - In our graph example, vertex v_1 has two adjacent vertices, v_2 and v_3
 - Based on this property, we can use an adjacency matrix or adjacency list to represent a graph
 - We can also use an adjacency list to represent a graph



	v_1	v_2	v_3	v_4	v_5
v_1	0	1	1	0	0
v_2	1	0	1	1	0
v_3	1	1	0	0	1
v_4	0	1	0	0	1
v_5	0	0	1	1	0

Adjacency Matrix

v_1	v_2, v_3			
v_2	v_1, v_3, v_4			
v_3	v_1, v_2, v_5			
v_4	v_2, v_5			
v_5	v_3, v_4			

Adjacency List

Dijkstra Algorithm for Shortest Path

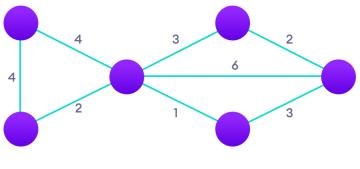
☐ The algorithm exists in many variants

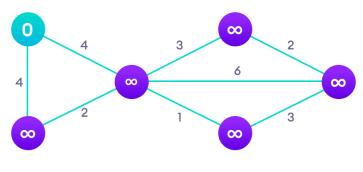
- Dijkstra's original algorithm finds the shortest path between two given nodes
- But a more common variant fixes a single node as the "source" node and finds shortest paths from the source to all other nodes in the graph, producing a shortest-path tree

Dijkstra Algorithm (Source to All)

- 1. Mark all nodes unvisited
- □ 2. Assign to every node a tentative distance value
 - Zero for the initial node and infinity for all other nodes
- □ 3. Select an unvisited node that is marked with the smallest tentative distance
- □ 4. Update the tentative distances of all the unvisited neighbors through the selected node and mark the selected node as visited
- ☐ 5. If all the nodes have been marked visited, then stop
 - Otherwise, go back to step 3

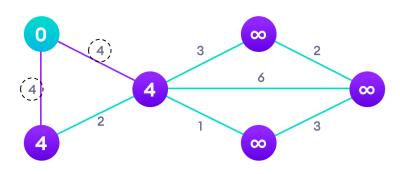
Example (1/2)

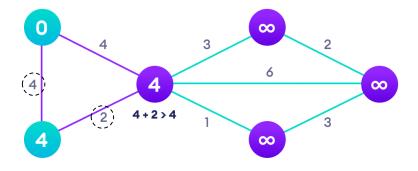




Step: 1

Step: 2

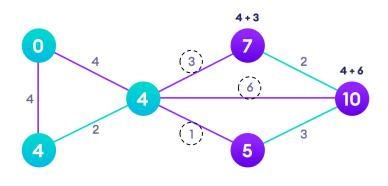




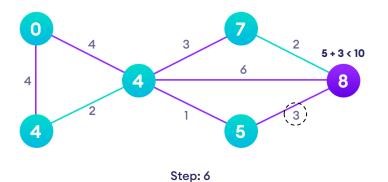
Step: 3

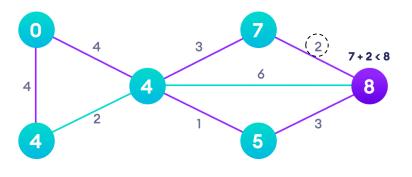
Step: 4

Example (2/2)



Step: 5





Step: 7

