# Hourai's Template

## 哈尔滨工业大学 蓬莱人形

## 2024年11月7日

# 目录

1	动态规划 4						
	1.1	多重背包	4				
	1.2	树形背包	4				
	1.3	动态动态规划 1	5				
	1.4	插头 dp	8				
	1.5	斜率优化	11				
2	数据结构						
	2.1	 平衡树					
	2.2	珂朵莉树					
	2.3	可并堆					
	2.4	线性基					
	2.5	Link Cut 树					
	2.6	线段树					
	2.7	根号数据结构					
3	树论		35				
3	がりに 3.1	。 点分树					
	3.1	<ul><li>長链剖分</li><li></li></ul>					
	3.3	重链剖分					
	3.4	<b>樹哈希</b>					
	$\frac{3.4}{3.5}$	Prufer 序列					
		Eruter 分列					
	3.6	<u>版</u>	44				
4	图论		<b>46</b>				
	4.1	仙人掌	46				
	4.2	三元环计数	48				
	4.3	四元环计数	48				
	4.4	基环树	49				
	4.5	2-SAT	52				
	4.6	割点	53				
	4.7	边双连通分量	54				
	4.8	点双连通分量	55				
	4.9	强连通分量	55				

目录 2

5	网络	流		
	5.1	费用流		
	5.2	最小割树		
	5.3	最大流		
	5.4	上下界费用流		
	5.5	上下界最大流		
6	数学	67		
	6.1	线性代数 67		
	6.2	大步小步 76		
	6.3	中国剩余定理		
	6.4	狄利克雷前缀和		
	6.5	万能欧几里得		
	6.6	扩展欧几里得		
	6.7	快速离散对数		
	6.8	快速最大公约数		
	6.9	原根		
	6.10	快速乘法逆元(离线)		
	6.11	快速乘法逆元(在线)		
	6.12	拉格朗日插值		
	6.13	min-max 容斥		
		Barrett 取模		
	6.15	Pollard's Rho		
	6.16	polya 定理		
		min25 筛		
		杜教筛		
		PN 筛		
	-	ヴ田粉主 		
		二次剩余		
		单位根反演		
	0.22	<b>丰恒恨及俱</b>		
7	7 多项式			
	7.1	NTT 全家桶		
	7.2	FWT 全家桶		
	7.3	任意模数 NTT		
8	字符			
	8.1	AC 自动机		
	8.2	扩展 KMP		
	8.3	Manacher		
	8.4	回文自动机		
	8.5	后缀平衡树		
	8.6	后缀数组(倍增)		
	8.7	后缀数组(SAIS)		
	8.8	广义后缀自动机(离线)112		
	8.9	广义后缀自动机(在线)113		
	8.10	后缀自动机		
	8.11	字典树		

目录

3

9	计算几何					
	9.1	二维凸包	117			
	9.2	最小圆覆盖	120			
	9.3	最左转线	121			
	9.4	二维基础	126			
10	10 其他					
	10.1	笛卡尔树	129			
		CDQ 分治				
	10.3	自适应辛普森	131			
	10.4	模拟退火	132			
	10.5	伪随机生成	133			

## 1 动态规划

#### 1.1 多重背包

#### 1.1.1 用法

n 个物品,m 容量背包,第 i 个物品重量为  $w_i$  价值为  $v_i$  共有  $c_i$  个,计算不超过容量的情况下最多拿多少价值的物品。

```
#include <bits/stdc++.h>
   using namespace std;
   using i64 = long long;
   const int INF = 1e9;
   const i64 INFL = 1e18;
 6 | int qread();
   const int MAXN = 4e4 + 3;
 7
   int F[MAXN];
   int main(){
9
10
        int n, m;
11
        cin >> n >> m;
12
        for(int i = 1; i \leq n; ++ i){
13
            int w, v, c;
            cin >> w >> v >> c;
14
15
            // w: value, v: volume, c: count
            for(int j = 0; j < v; ++ j){}
16
                deque <tuple<int, int> > Q;
17
                for(int k = 0; j + k * v \leq m; ++ k){
18
                     int x = j + k * v;
19
                     int f = F[x] - (x / v) * w;
20
21
                    while(!Q.empty() & get<0>(Q.back ()) \leq f)
22
                         Q.pop back ();
23
                     Q.push_back(\{f, x\});
24
                    while(!Q.empty() & get<1>(Q.front()) < x - c * v)
25
                         Q.pop front();
                     F[x] = get<0>(Q.front()) + (x / v) * w;
26
                }
27
            }
28
29
30
        cout << F[m] << endl;</pre>
31
        return 0;
32
```

#### 1.2 树形背包

```
1 #include<bits/stdc++.h>
 2 using namespace std;
 3 typedef long long i64;
   const int MAXN = 2e3 + 3;
 5 | vector<int> E[MAXN];
   int W[MAXN];
 6
   int F[MAXN][MAXN], S[MAXN];
 7
   void dfs(int u, int f){
 8
 9
       F[u][1] = W[u];
10
       S[u]
                = 1;
       for(auto &v : E[u]) if(v \neq f){
11
12
            dfs(v, u);
            for(int i = S[u]; i \ge 1; -- i)
13
```

```
for(int j = S[v]; j \ge 1; -- j)
14
                     F[u][i + j] = max(F[u][i + j], F[u][i] + F[v][j]);
15
16
            S[u] += S[v];
        }
17
18
19
   int main(){
20
        int n, m;
21
        cin >> n >> m;
        for(int i = 1; i \leq n; ++ i){
22
23
            int f;
24
            cin >> f >> W[i];
25
            E[f].push_back(i);
26
        }
27
        dfs(0, 0);
        cout << F[0][m + 1] << endl;
28
29
        return 0;
30
   }
```

## 1.3 动态动态规划 1

#### 1.3.1 例题

给定一棵 n 个点的树, 点有点权, 求最大独立集。m 次修改, 每次把 x 的权值修改成 y。

```
#include<bits/stdc++.h>
   #define up(l, r, i) for(int i = l, END##i = r; i \leq END##i; ++ i)
   #define dn(r, l, i) for(int i = r, END##i = l;i \ge END##i;-- i)
3
4 using namespace std;
5
   typedef long long i64;
   const int INF = 1e9;
7
   const int MAXN = 1e5 + 3;
   int W[MAXN];
8
   struct Mat{ int M[2][2]; };
9
   struct Vec{ int V[2];
10
                             };
11
   Mat operator *(const Mat &a, const Mat &b){
12
       Mat c;
       c.M[0][0] = max(a.M[0][0] + b.M[0][0], a.M[0][1] + b.M[1][0]);
13
       c.M[0][1] = max(a.M[0][0] + b.M[0][1], a.M[0][1] + b.M[1][1]);
14
       c.M[1][0] = max(a.M[1][0] + b.M[0][0], a.M[1][1] + b.M[1][0]);
15
16
       c.M[1][1] = max(a.M[1][0] + b.M[0][1], a.M[1][1] + b.M[1][1]);
17
       return c;
18
   Vec operator *(const Mat &a, const Vec &v){
19
20
       Vec r;
       r.V[0] = max(a.M[0][0] + v.V[0], a.M[0][1] + v.V[1]);
21
       r.V[1] = max(a.M[1][0] + v.V[0], a.M[1][1] + v.V[1]);
22
23
       return r;
24
25
   namespace Gra{
       vector<int> E[MAXN];
26
       int G[MAXN], S[MAXN], D[MAXN], T[MAXN], F[MAXN];
27
       int X[MAXN], Y[MAXN];
28
       int H[MAXN][2];
29
30
       int K[MAXN][2];
       struct Mat M[MAXN];
31
32
       void dfs1(int u, int f){
           S[u] = 1;
33
34
           F[u] = f;
```

```
for(auto &v : E[u]) if(v \neq f){
35
                dfs1(v, u);
36
                S[u] += S[v];
37
                if(S[v] > S[G[u]]) G[u] = v;
38
39
            }
        }
40
        int o;
41
42
        void dfs2(int u, int f){
43
            if(u = G[f])
                X[u] = X[f];
44
45
            else
46
                X[u] = u;
47
            H[u][0] = H[u][1] = 0;
            K[u][0] = K[u][1] = 0;
48
49
            const int &g = G[u];
            D[u] = ++ o;
50
51
            T[o] = u;
52
            if(g){
                dfs2(g, u);
53
54
                Y[u] = Y[g];
                K[u][0] += max(K[g][0], K[g][1]);
55
56
                K[u][1] += K[g][0];
            } else {
57
                Y[u] = u;
58
59
60
            for(auto \delta v : E[u]) if(v \neq f \delta \delta v \neq g){
61
                dfs2(v, u);
                H[u][0] += max(K[v][0], K[v][1]);
62
63
                H[u][1] += K[v][0];
64
            M[u].M[0][0] = H[u][0];
65
            M[u].M[0][1] = H[u][0];
66
            M[u].M[1][0] = H[u][1] + W[u];
67
68
            M[u].M[1][1] = -INF;
69
            K[u][0] += H[u][0];
            K[u][1] += H[u][1] + W[u];
70
        }
71
72
   namespace Seg{
73
74
        const int SIZ = 4e5 + 3;
        struct Mat M[SIZ];
75
76
        #define lc(t) (t << 1)
        #define rc(t) (t \ll 1 | 1)
77
        void pushup(int t, int a, int b){
78
            M[t] = M[lc(t)] * M[rc(t)];
79
80
        void build(int t, int a, int b){
81
            if(a = b){
82
83
                M[t] = Gra :: M[Gra :: T[a]];
84
            } else {
85
                int c = a + b >> 1;
86
                build(lc(t), a, c);
87
                build(rc(t), c + 1, b);
88
                pushup(t, a, b);
            }
89
90
        void modify(int t, int a, int b, int p, const Mat &w){
91
```

```
92
            if(a = b){
                 M[t] = w;
93
            } else {
94
 95
                 int c = a + b >> 1;
96
                 if(p \le c) modify(lc(t), a, c, p, w);
                          modify(rc(t), c + 1, b, p, w);
 97
98
                 pushup(t, a, b);
            }
99
100
        Mat query(int t, int a, int b, int l, int r){
101
102
            if(l \leq a \& b \leq r){
                 return M[t];
103
104
            } else {
105
                 int c = a + b >> 1:
                 if(r \le c) return query(lc(t), a, c , l, r); else
106
                 if(l > c) return query(rc(t), c + 1, b, l, r); else
107
108
                     query(rc(t), c + 1, b, l, r);
109
            }
110
        }
111
112
113
    int qread();
    int main(){
114
115
        int n = qread(), m = qread();
        up(1, n, i)
116
            W[i] = qread();
117
118
        up(2, n, i){
            int u = qread(), v = qread();
119
120
            Gra :: E[u].push_back(v);
121
            Gra :: E[v].push_back(u);
122
        Gra :: dfs1(1, 0);
123
        Gra :: dfs2(1, 0);
124
        Seg :: build(1, 1, n);
125
        Vec v0;
126
        v0.V[0] = v0.V[1] = 0;
127
        up(1, m, i){
128
129
            using namespace Gra;
            int x = qread(), y = qread();
130
131
            W[x] = y;
            int u = x;
132
            while(u \neq 0){
133
                 const int &v = X[u];
134
                 const int &f = F[v];
135
                 M[u].M[0][0] = H[u][0];
136
137
                 M[u].M[0][1] = H[u][0];
                 M[u].M[1][0] = H[u][1] + W[u];
138
                 M[u].M[1][1] = -INF;
139
140
                 const Vec p = Seg :: query(1, 1, n, D[v], D[Y[u]]) * v0;
                 Seg :: modify(1, 1, n, D[u], M[u]);
141
                 const Vec q = Seg :: query(1, 1, n, D[v], D[Y[u]]) * v0;
142
                 if(f \neq 0){
143
                     H[f][0] = H[f][0] - max(p.V[0], p.V[1]) + max(q.V[0], q
144
                        .V[1]);
                     H[f][1] = H[f][1] - p.V[0] + q.V[0];
145
146
147
                 u = f;
```

## 1.4 插头 dp

#### 1.4.1 例题

给出  $n \times m$  的方格,有些格子不能铺线,其它格子必须铺,形成一个闭合回路。问有多少种铺法?

```
#include<bits/stdc++.h>
   #define up(l, r, i) for(int i = l, END##i = r;i \leq END##i;++ i)
   #define dn(r, l, i) for(int i = r, END##i = l;i \ge END##i;-- i)
   using namespace std;
   using i64 = long long;
   const int INF = 1e9;
6
   const i64 INFL = 1e18;
7
   const int MAXN = 20 + 3;
   const int MAXM = 67108864 + 3;
9
   namespace HashT{
10
       const int SIZ = 19999997;
11
       int H[SIZ], V[SIZ], N[SIZ], t;
12
13
       bool F[SIZ];
14
       i64 W[SIZ];
       void add(int u, int v, bool f, i64 w){
15
16
           V[++ t] = v, N[t] = H[u], F[t] = f, W[t] = w, H[u] = t;
17
       i648 find(int u, bool f){
18
19
            for(int p = H[u % SIZ];p;p = N[p])
                if(V[p] = u \& F[p] = f)
20
21
                    return W[p];
           add(u % SIZ, u, f, 0);
22
23
            return W[t];
       }
24
25
   char S[MAXN][MAXN];
26
27
   int qread();
28
   int n, m;
   vector <pair<int, bool>, i64> > M[2];
   int getp(int s, int p){
30
31
       return (s >> (2 * p - 2)) & 3;
32
   int setw(int s, int p, int w){
33
34
       return (s & \sim(3 << (2 * p - 2))) | (w << (2 * p - 2));
35
   int findr(int s, int p){
36
37
       int c = 0;
38
       for(int q = p; q \leq m + 1; ++ q){}
            if(((s >> (2 * q - 2)) & 3) = 1) ++ c;
39
40
           if(((s \gg (2 * q - 2)) \& 3) = 2) -- c;
           if(c = 0)
41
42
                return q;
       }
43
```

```
44
       return -1;
45
   int findl(int s, int p){
46
47
        int c = 0;
        for(int q = p; q \geqslant 1; -- q) \{
48
            if(((s \gg (2 * q - 2)) \& 3) = 2) + c;
49
            if(((s \gg (2 * q - 2)) \& 3) = 1) -- c;
50
51
            if(c = 0)
52
                return q;
53
54
       return -1;
55
56
   void state(int s){
57
       return ;
        up(1, m + 1, i){
58
59
            switch(getp(s, i)){
60
                case 0 : putchar('#'); break;
                case 1 : putchar('('); break;
61
                case 2 : putchar(')'); break;
62
63
                case 3 : putchar('E');
64
65
        }
        puts("");
66
67
68
   int main(){
69
        n = qread(), m = qread();
70
        up(1, n, i)
            scanf("%s", S[i] + 1);
71
72
        int o = 0;
73
        #define X M[ o]
74
        #define Y M[!o]
75
        vector <pair<int, bool> > T;
        X.push_back(\{\{0, 0\}, 1\});
76
77
        up(1, n, i){
78
            Y.clear();
79
            for(auto &u : X){
                auto [s0, c] = u;
80
81
                auto [s, f] = s0;
                if(getp(s, m + 1) = 0)
82
83
                    Y.push_back(\{\{s << 2, f\}, c\});
84
            o ~= 1;
85
86
            up(1, m, j){
                int x = j, y = j + 1;
87
                for(auto &u : X){
88
89
                    auto [s0, c] = u;
                    auto [s, f] = s0;
90
91
                    int a = getp(s, x);
92
                    int b = getp(s, y);
93
                    int t = setw(setw(s, x, 0), y, 0);
                    #define update(t, c) HashT :: find(t, f) += c, T.
94
                        push_back({t, f})
                    if(S[i][j] = '.'){
95
                                            // 经过该格
96
                         if(a = 1 & b = 1){
                             t = setw(t, findr(s, y), 1),
97
98
                             update(t, c);
99
                         } else
```

```
if(a = 2 & b = 2){
100
                             t = setw(t, findl(s, x), 2),
101
                             update(t, c);
102
103
                         } else
                         if(a = 1 \& b = 2)
104
                             if(f = false) // 还没有闭合回路
105
                                 f = true, update(t, c);
106
107
                         } else
                         if(a = 2 & b = 1){
108
                             update(t, c);
109
                         } else
110
                         if(a = 0 \& b = 0)
111
112
                             t = setw(t, x, 1);
                             t = setw(t, y, 2);
113
                             update(t, c);
114
                                     // a = 0 || b = 0
115
                         } else {
116
                             int t1 = setw(t, x, a | b);
                             int t2 = setw(t, y, a | b);
117
                             update(t1, c);
118
119
                             update(t2, c);
120
121
                     if(S[i][j] = '*'){ // 不经过该格
122
                         if(a = 0 \& b = 0)
123
                             update(t, c);
124
                     }
125
126
                 Y.clear();
127
128
                 for(auto &u : T){
                     auto [s, f] = u;
129
                     if(HashT :: find(s, f) \neq 0){
130
                         Y.push_back(\{\{s, f\}, HashT :: find(s, f)\});
131
132
                         HashT :: find(s, f) = 0;
                     }
133
134
                 T.clear(), o ^{-} 1;
135
             }
136
137
        i64 \text{ ans} = 0;
138
        for(auto &u : X){
139
            auto [s0, c] = u;
140
            auto [s, f] = s0;
141
142
            bool g = true;
            up(1, m + 1, i)
143
                 g &= getp(s, i) = 0;
144
            f &= g;
145
             if(f)
146
147
                 ans = c;
148
        }
        printf("%lld\n", ans);
149
150
        return 0;
151
```

#### 1.5 斜率优化

#### 1.5.1 形式

考虑一个经典的 dp 转移方程如下:

$$f_i = \max_{j \le i} \{ f(j) + w(j, i) \}$$

我们将式子拆成三个部分: 只跟 i 有关或者与 i,j 均不相关的部分 a(i), 只跟 j 有关的部分 b(j), 跟 i,j 均有关的部分 c(i,j):

$$f_i = a(i) + \max_{j < i} \{b(j) + c(i, j)\}$$

斜率优化可被用来解决这样一个情形:  $c(i,j) = ic_j$ 。此时 b(j) + c(i,j) 可视作关于 j 的一次函数。如果  $c_j$  随着 j 的增大而单调,那么可用单调栈维护;否则可以考虑 CDQ 分治或者在凸包上二分。在凸包上可以使用二分查询最高/最低点。

#### 1.5.2 例题

玩具装箱。原始转移方程为:

$$f_i = \max_{j < i} \{ f_j + (s_i - s_j - L')^2 \}$$
  
其中  $s_i = i + \sum_{j \le i} c_i, L' = L + 1$ 。 将其分类得到: 
$$f_i = \max_{j < i} \{ f_j + s_i^2 + s_j^2 + L'^2 - 2s_i s_j + 2s_j L' - 2s_i L' \}$$
$$= (s_i^2 - 2s_i L' + L'^2) + \max_{j < i} \{ (f_j + s_j^2 + 2s_j L') - 2s_i s_j \}$$

在原始的玩具装箱中, $s_j$  单调增加,也就是斜率单调增加。因此可以直接使用单调栈维护凸包。同时  $s_i$  也单调增加,因此可以用指针维护。

```
#include<bits/stdc++.h>
 2 using namespace std;
 3 typedef long long
                        i64;
 4 typedef long double f80;
 5 const int INF = 2147483647;
 6 const int MAXN= 5e4 + 3;
 7 | int n, L, p, e, C[MAXN], Q[MAXN];
   f80 S[MAXN], F[MAXN];
   f80 gtx(int x){ return S[x]; }
   f80 gty(int x){ return F[x] + S[x] * S[x]; }
10
   f80 gtw(int x){ return -2.0 * (L - S[x]); }
11
   f80 gtk(int x,int y){ return (gty(y) - gty(x)) / (gtx(y) - gtx(x)); }
12
   int main(){
13
14
       cin \gg n \gg L;
        for(int i = 1; i \leq n; ++ i){
15
            cin \gg C[i];
16
            S[i] = S[i - 1] + C[i];
17
18
       for(int i = 1; i \leq n; ++ i){
19
            S[i] += i;
20
21
22
       e = p = 1, L ++, Q[p] = 0;
23
        for(int i = 1; i \leq n; ++ i){
            while(e \& gtk(Q[e], Q[e + 1]) < gtw(i))
24
25
                ++ e;
```

```
int j = Q[e];
26
27
          F[i] = F[j] + pow(S[i] - S[j] - L, 2);
          while(1  gtk(Q[p], i))
28
29
             e -= (e = p), -- p;
          Q[++p]=i;
30
31
      printf("%.0Lf\n", F[n]);
32
33
      return 0;
34
```

## 2 数据结构

#### 2.1 平衡树

## 2.1.1 无旋 Treap

```
#include<bits/stdc++.h>
2 using namespace std;
3 using i64 = long long;
4 const int INF = 1e9;
   const i64 INFL = 1e18;
   typedef unsigned int
6
                               u32;
   typedef unsigned long long u64;
7
8
   mt19937_64 MT(114514);
9
   namespace Treap{
10
       const int SIZ = 1e6 + 1e5 + 3;
11
       int F[SIZ], C[SIZ], S[SIZ], W[SIZ], X[SIZ][2], sz;
12
       u64 H[SIZ];
       int newnode(int w){
13
           W[++ sz] = w, C[sz] = S[sz] = 1; H[sz] = MT();
14
15
           return sz;
16
17
       void pushup(int x){
            S[x] = C[x] + S[X[x][0]] + S[X[x][1]];
18
19
20
       pair<int, int> split(int u, int x){
21
            if(u = 0)
22
                return make_pair(0, 0);
23
           if(W[u] > x){
24
                auto [a, b] = split(X[u][0], x);
                X[u][0] = b, pushup(u);
25
                return make_pair(a, u);
26
27
            } else {
                auto [a, b] = split(X[u][1], x);
28
                X[u][1] = a, pushup(u);
29
30
                return make pair(u, b);
            }
31
32
33
       int merge(int a, int b){
            if(a = 0 || b = 0)
34
35
                return a | b;
           if(H[a] < H[b]){
36
                X[a][1] = merge(X[a][1], b), pushup(a);
37
38
                return a;
39
            } else {
                X[b][0] = merge(a, X[b][0]), pushup(b);
40
```

```
41
                return b;
            }
42
43
       void insert(int &root, int w){
44
45
            auto [p, q] = split(root, w
            auto [a, b] = split( p, w - 1);
46
            if(b \neq 0){
47
48
                ++ S[b], ++ C[b];
49
            } else b = newnode(w);
              = merge(a, b);
50
51
            root = merge(p, q);
       }
52
       void erase(int &root, int w){
53
            auto [p, q] = split(root, w
54
            auto [a, b] = split( p, w - 1);
55
            -- C[b], -- S[b];
56
57
                = C[b] = 0 ? a : merge(a, b);
            root = merge(p, q);
58
59
60
        int find_rank(int &root, int w){
            int x = root, o = x, a = 0;
61
62
            for(;x;){
                if(w < W[x])
63
                    o = x, x = X[x][0];
64
65
                else {
66
                    a += S[X[x][0]];
                    if(w = W[x])
67
                        o = x; break;
68
69
                    }
70
                    a += C[x];
71
                    o = x, x = X[x][1];
72
            }
73
74
            return a + 1;
75
       int find_kth(int &root, int w){
76
            int x = root, o = x, a = 0;
77
78
            for(;x;){
                if(w \leq S[X[x][0]])
79
80
                    o = x, x = X[x][0];
81
                else {
82
                    W -= S[X[x][0]];
83
                    if(w \leq C[x]){
84
                        o = x; break;
                    }
85
86
                    w -= C[x];
87
                    o = x, x = X[x][1];
88
89
            }
90
           return W[x];
91
92
        int find_pre(int &root, int w){
           return find_kth(root, find_rank(root, w) - 1);
93
94
       int find_suc(int &root, int w){
95
96
            return find_kth(root, find_rank(root, w + 1));
        }
97
```

```
98
 99
    // == TEST ===
100 | int qread();
    int main(){
101
        using namespace Treap;
102
         int n = qread(), m = qread(), root = 0;
103
         for(int i = 1; i \leq n; ++ i){
104
105
             int a = qread(); insert(root, a);
106
        int last_ans = 0, ans = 0;
107
         for(int i = 1; i \leq m; ++ i){
108
             int op = qread(), x = qread() ^ last_ans;
109
110
             switch(op){
111
                 case 1 : insert(root, x); break;
                 case 2 : erase (root, x); break;
112
                 case 3 : ans ^= (last_ans = find_rank(root, x)); break;
113
114
                 case 4 : ans ^= (last_ans = find_kth (root, x)); break;
                 case 5 : ans ^= (last_ans = find_pre (root, x)); break;
115
                 case 6 : ans '= (last ans = find suc (root, x)); break;
116
             }
117
118
119
         printf("%d\n", ans);
120
        return 0;
121
```

#### 2.1.2 Splay

```
1 #include < bits/stdc++.h>
   using namespace std;
 3 using i64 = long long;
   const int INF = 1e9;
 4
   const i64 INFL = 1e18;
 6
   namespace Splay{
       const int SIZ = 1e6 + 1e5 + 3;
 7
        int F[SIZ], C[SIZ], S[SIZ], X[SIZ][2], size;
 8
 9
        bool T[SIZ];
10
        bool is_root(int x){ return
                                       F[x]
       bool is_rson(int x){ return X[F[x]][1] = x;}
11
12
        void push_down(int x){
            if(!T[x]) return;
13
            int lc = X[x][0], rc = X[x][1];
14
            if(lc) T[lc] \cong 1, swap(X[lc][0], X[lc][1]);
15
            if(rc) T[rc] '= 1, swap(X[rc][0], X[rc][1]);
16
           T[x] = 0;
17
18
        void pushup(int x){
19
20
            S[x] = C[x] + S[X[x][0]] + S[X[x][1]];
21
       void rotate(int x){
22
            int y = F[x], z = F[y];
23
            bool f = is_rson(x);
24
            bool g = is_rson(y);
25
26
            int \delta t = X[x][!f];
            if(z){X[z][g] = x;}
27
28
            if(t){ F[t]
                         = y; }
            X[y][f] = t, t = y;
29
30
            F[y] = x, pushup(y);
```

```
F[x] = z, pushup(x);
31
        }
32
        void splay(int &r, int x, int g = 0){
33
34
            for(int f;f = F[x], f \neq g;rotate(x))
                if(F[f] \neq g) rotate(is rson(x) = is rson(f) ? f : x);
35
36
            if(is root(x)) r = x;
        }
37
38
        int get_kth(int &r, int w){
39
            int x = r, o = x;
40
            for(;x;){
41
                push_down(x);
                if(w \leq S[X[x][0]]) o = x, x = X[x][0]; else {
42
43
                    W -= S[X[x][0]];
44
                    if(C[x] \& w \leq C[x])\{o = x; break;\}
                    W -= C[x], o = x, x = X[x][1];
45
46
47
            }
48
            splay(r, o); return o;
49
50
        int build(int l, int r){
            if(l = r)
51
52
                C[l] = S[l] = 1; return l;
53
            int c = l + r \gg 1, a = 0, b = 0;
54
            if(l \le c - 1) a = build(l, c - 1), F[a] = c, X[c][0] = a;
55
            if(c + 1 \le r) b = build(c + 1, r), F[b] = c, X[c][1] = b;
56
            C[c] = 1, pushup(c); return c;
57
58
59
        void output(int n, int &r){
            push down(r);
60
            if(X[r][0]) output(n, X[r][0]);
61
            if(r \neq 1 & r \neq n + 2) printf("%d ", r - 1);
62
            if(X[r][1]) output(n, X[r][1]);
63
        }
64
65
66
   int gread();
   int main(){
67
68
        using namespace Splay;
        int n = qread(), m = qread();
69
70
        int root = build(1, n + 2);
        for(int i = 1; i \leq m; ++ i){
71
            int l = qread() + 1, r = qread() + 1;
72
73
            int u = get kth(root, r + 1);
            int v = get_kth(root, l - 1);
74
            splay(root, v, 0), splay(root, u, v);
75
            int t = X[u][0];
76
            T[t] \cong 1, swap(X[t][0], X[t][1]);
77
78
79
        output(n, root);
80
        return 0;
81
```

#### 2.1.3 Treap

```
1 #include<bits/stdc++.h>
2 #define up(l, r, i) for(int i = l, END##i = r;i ≤ END##i;++ i)
3 #define dn(r, l, i) for(int i = r, END##i = l;i ≥ END##i;-- i)
```

```
4 using namespace std;
 5 typedef long long i64;
 6 const int INF = 2147483647;
 7 typedef unsigned int
 8 typedef unsigned long long u64;
9 | mt19937 64 MT(114514);
10 | namespace Treap{
11
       const int SIZ = 1e6 + 1e5 + 3;
12
        int F[SIZ], C[SIZ], S[SIZ], W[SIZ], X[SIZ][2], sz;
13
       u64 H[SIZ]:
        bool is_root(int x){ return F[x]
14
       bool is_rson(int x){ return X[F[x]][1] = x;}
15
16
        int newnode(int w){
           W[++ sz] = w, C[sz] = S[sz] = 1; H[sz] = MT();
17
            return sz;
18
19
20
       void pushup(int x){
            S[x] = C[x] + S[X[x][0]] + S[X[x][1]];
21
22
23
       void rotate(int &root, int x){
            int y = F[x], z = F[y];
24
25
            bool f = is_rson(x);
            bool g = is_rson(y);
26
            int \delta t = X[x][!f];
27
            if(z){ X[z][g] = x; } else root = x;
28
           if(t){ F[t]
29
                         = y; }
           X[y][f] = t, t = y;
30
            F[y] = x, pushup(y);
31
32
            F[x] = z, pushup(x);
33
34
       void insert(int &root, int w){
            if(root = 0) {root = newnode(w); return;}
35
36
            int x = root, o = x;
            for(;x;o = x, x = X[x][w > W[x]]){
37
                ++ S[x]; if(w = W[x]) \{ ++ C[x], o = x; break; \}
38
39
            if(W[o] \neq w){
40
41
                if(w < W[o]) X[o][0] = newnode(w), F[sz] = o, o = sz;
                             X[o][1] = newnode(w), F[sz] = o, o = sz;
42
43
           while(!is_root(o) & H[o] < H[F[o]])</pre>
44
45
                rotate(root, o);
46
       void erase(int &root, int w){
47
48
            int x = root, o = x;
49
            for(;x;o = x, x = X[x][w > W[x]]){
                -- S[x]; if(w = W[x]){ -- C[x], o = x; break;}
50
51
52
            if(C[o] = \emptyset)
                while(X[o][0] || X[o][1]){
53
                    u64 wl = X[o][0] ? H[X[o][0]] : ULLONG_MAX;
54
                    u64 wr = X[o][1] ? H[X[o][1]] : ULLONG_MAX;
55
56
                    if(wl < wr){
57
                        int p = X[o][0]; rotate(root, p);
58
                    } else {
59
                        int p = X[o][1]; rotate(root, p);
60
                    }
```

```
61
 62
                 if(is_root(o)){
 63
                     root = 0;
 64
                 } else {
                     X[F[o]][is\_rson(o)] = 0;
 65
 66
             }
 67
 68
         int find rank(int &root, int w){
 69
 70
             int x = root, o = x, a = 0;
 71
             for(;x;){
                 if(w < W[x])
 72
 73
                     o = x, x = X[x][0];
 74
                      a += S[X[x][0]];
 75
                      if(w = W[x]){
 76
 77
                          o = x; break;
                      }
 78
 79
                     a += C[x];
 80
                     o = x, x = X[x][1];
 81
 82
             }
 83
             return a + 1;
 84
         int find_kth(int &root, int w){
 85
 86
             int x = root, o = x, a = 0;
             for(;x;){
 87
                 if(w \leq S[X[x][0]])
 88
 89
                     o = x, x = X[x][0];
 90
                 else {
 91
                     w -= S[X[x][0]];
 92
                     if(w \leq C[x]){
 93
                          o = x; break;
                      }
 94
 95
                     w -= C[x];
                     o = x, x = X[x][1];
 96
 97
 98
             }
 99
             return W[x];
100
         int find_pre(int &root, int w){
101
             return find_kth(root, find_rank(root, w) - 1);
102
103
         int find_suc(int &root, int w){
104
             return find_kth(root, find_rank(root, w + 1));
105
106
         }
107
    int qread();
108
109
    int main(){
110
         using namespace Treap;
         int n = qread(), m = qread(), root = 0;
111
         up(1, n, i){
112
             int a = qread(); insert(root, a);
113
114
         int last_ans = 0, ans = 0;
115
116
         up(1, m, i){
             int op = qread(), x = qread() ^ last_ans;
117
```

```
switch(op){
118
119
                 case 1 : insert(root, x); break;
                 case 2 : erase (root, x); break;
120
                 case 3 : ans ^= (last ans = find rank(root, x)); break;
121
                 case 4 : ans ^= (last_ans = find_kth (root, x)); break;
122
                 case 5 : ans ^= (last_ans = find_pre (root, x)); break;
123
                 case 6 : ans ^= (last_ans = find_suc (root, x)); break;
124
             }
125
        }
126
        printf("%d\n", ans);
127
128
        return 0;
129
```

#### 2.2 珂朵莉树

```
#include<bits/stdc++.h>
   using namespace std:
 3 |using u64 = unsigned long long;
   const int MAXN = 1e6 + 3;
   int power(int a, int b, int p){
 5
 6
        int r = 1;
       while(b){
 7
8
            if(b & 1) r = 1ll * r * a % p;
 9
            b >>= 1, a = 1ll * a * a % p;
10
11
        return r;
12
13
   namespace ODT {
14
        // <pos_type, value_type>
15
        map <int, long long> M;
16
        // 分裂为 [1, p) 和 [p, +inf), 返回后者迭代器
17
        auto split(int p) {
            auto it = prev(M.upper_bound(p));
18
            return M.insert(
19
20
                it,
                make_pair(p, it \rightarrow second)
21
22
            );
        }
23
24
        // 区间赋值
        void assign(int l, int r, int v) {
25
26
            auto it = split(l);
27
            split(r + 1);
            while (it \rightarrow first \neq r + 1) {
28
29
                it = M.erase(it);
            }
30
31
            M[l] = v;
        }
32
        // // 执行操作
33
        // void perform(int l, int r) {
34
               auto it = split(l);
35
        //
        //
               split(r + 1);
36
37
        //
               while (it \rightarrow first \neq r + 1) {
38
        //
                   // Do something...
        //
39
                   it = next(it);
40
        //
41
        // }
42
        void modify1(int l, int r, int w) {
```

```
auto it = split(l);
43
44
             split(r + 1);
             while(it \rightarrow first \neq r + 1) {
45
46
                 it \rightarrow second += w;
                 it = next(it);
47
             }
48
        }
49
50
        void modify2(int l, int r, int w) {
             assign(l, r, w);
51
52
53
        long long query1(int l, int r, int k) {
             auto it = split(l);
54
55
             split(r + 1);
             map <long long, int> T;
56
             while(it \rightarrow first \neq r + 1) {
57
                 T[it \rightarrow second] += next(it) \rightarrow first - it \rightarrow first;
58
59
                 it = next(it);
60
             for(auto &[w, c]: T){
61
62
                 if(c \ge k)
63
                      return w;
64
                 k -= c;
             }
65
             return -1;
66
67
68
        long long query2(int l, int r, int x, int y) {
69
             auto it = split(l);
             split(r + 1);
70
71
             int ans = 0;
72
             while(it \rightarrow first \neq r + 1) {
                 int c = next(it) \rightarrow first - it \rightarrow first;
73
                 long long a = it \rightarrow second;
74
                 ans = (ans + 1ll * c * power(a % y, x, y)) % y;
75
76
                 it = next(it);
             }
77
78
             return ans;
        }
79
80
   };
   const int MOD = 1e9 + 7;
81
   int read(int &seed){
82
83
        int ret = seed;
84
        seed = (seed * 7ll + 13) % MOD;
85
        return ret;
86
87
   int main(){
88
        ios :: sync_with_stdio(false);
89
        cin.tie(nullptr);
90
        int n, m, seed, vmax;
91
        cin >> n >> m >> seed >> vmax;
        ODT :: M[n + 1] = 0;
92
        for(int i = 1; i \leq n; ++ i){
93
             int a = read(seed) % vmax + 1;
94
95
             ODT :: M[i] = a;
96
        for(int i = 1; i \leq m; ++ i){
97
98
             int op = read(seed) % 4 + 1;
99
             int l = read(seed) % n + 1;
```

```
100
             int r = read(seed) \% n + 1;
101
             int x, y;
             if(l > r)
102
103
                 swap(l, r);
             if(op = 3){
104
                 x = (read(seed) \% (r - l + 1)) + 1;
105
106
             } else
107
                 x = read(seed) \% vmax + 1;
             if(op = 4)
108
                 y = read(seed) \% vmax + 1;
109
             if(op = 1){
110
                 ODT :: modify1(l, r, x);
111
112
             } else
             if(op = 2){
113
                 ODT :: modify2(l, r, x);
114
             } else
115
116
             if(op = 3){
                 cout \ll ODT :: query1(l, r, x) \ll "\n";
117
             } else
118
119
             if(op = 4){
                 cout \ll ODT :: query2(l, r, x, y) \ll "\n";
120
121
             }
         }
122
123
         return 0;
124
```

#### 2.3 可并堆

```
#include <bits/stdc++.h>
   using namespace std;
   using i64 = long long;
 3
   const int INF = 1e9;
   const i64 INFL = 1e18;
 5
   namespace LeftHeap{
 7
       const int SIZ = 1e5 + 3;
 8
       int W[SIZ], D[SIZ];
 9
       int L[SIZ], R[SIZ];
       int F[SIZ], s;
10
11
       bool E[SIZ];
        int merge(int u, int v){
12
            if(u = 0 || v = 0)
13
14
                return u | v;
            if(W[u] > W[v] || (W[u] = W[v] & u > v))
15
                swap(u, v);
16
            int &lc = L[u];
17
            int &rc = R[u];
18
19
            rc = merge(rc, v);
20
            if(D[lc] < D[rc])
                swap(lc, rc);
21
            D[u] = min(D[lc], D[rc]) + 1;
22
23
            if(lc \neq 0) F[lc] = u;
            if(rc \neq 0) F[rc] = u;
24
25
            return u;
        }
26
27
       void pop(int &root){
            int root0 = merge(L[root], R[root]);
28
29
            F[root0] = root0;
```

```
30
            F[root] = root0;
            E[root] = true;
31
32
            root = root0;
        }
33
34
        int top(int &root){
            return W[root];
35
36
37
        int getfa(int u){
            return u = F[u] ? u : F[u] = getfa(F[u]);
38
39
        int newnode(int w){
40
            ++ s;
41
42
            W[s] = w;
43
            F[s] = s;
            D[s] = 1;
44
45
            return s;
46
        }
47
   // == TEST ==
48
49
   int gread();
   const int MAXN = 1e5 + 3;
50
   int A[MAXN], O[MAXN];
52
   int main(){
53
        int n, m;
54
        cin >> n >> m;
55
        for(int i = 1; i \leq n; ++ i){
56
            cin \gg A[i];
57
            O[i] = LeftHeap :: newnode(A[i]);
58
        for(int i = 1; i \leq m; ++ i){}
59
60
            int op;
            cin >> op;
61
            if(op = 1){
62
                int x, y;
63
64
                cin >> x >> y;
                if(LeftHeap :: E[0[x]])
65
66
                     continue;
67
                if(LeftHeap :: E[O[y]])
                     continue;
68
69
                int fx = LeftHeap :: getfa(0[x]);
                int fy = LeftHeap :: getfa(0[y]);
70
71
                if(fx \neq fy){
72
                     LeftHeap :: merge(fx, fy);
                }
73
            } else {
74
                int x;
75
                cin >> x;
76
                if(LeftHeap :: E[0[x]]){
77
78
                     cout << -1 << endl;
79
                     continue;
80
                int fx = LeftHeap :: getfa(0[x]);
81
                cout << LeftHeap :: top(fx) << endl;</pre>
82
83
                LeftHeap :: pop(fx);
            }
84
85
        return 0;
86
```

**87** | }

## 2.4 线性基

```
#include <bits/stdc++.h>
   using namespace std;
   using i64 = long long;
   const int INF = 1e9;
   const i64 INFL = 1e18;
 5
 6
   namespace LB{
        const int SIZ = 60 + 3;
7
8
        i64 W[SIZ], h = 60;
        void insert(i64 w){
9
            for(int i = h; i \ge 0; -- i){
10
                 if(w & (1ll << i)){
11
12
                     if(!W[i]){
                         W[i] = w;
13
14
                         break;
15
                     } else {
                         w \sim W[i];
16
17
                 }
18
19
            }
20
        }
21
        i64 query(i64 x){
            for(int i = h; i \ge 0; -- i){
22
                 if(W[i]){
23
                     x = max(x, x ^ W[i]);
24
25
26
27
            return x;
        }
28
29
   namespace realLB{
30
31
        const int SIZ = 500 + 3;
        long double W[SIZ][SIZ];
32
33
        int n = 0;
34
        void init(int n0){
35
            n = n0;
36
        bool zero(long double w){
37
38
            return fabs(w) < 1e-9;</pre>
39
40
        bool insert(long double X[]){
            for(int i = 1; i \leq n; ++ i){
41
                 if(!zero(X[i])){
42
43
                     if(zero(W[i][i])){
44
                         for(int j = 1; j \leq n; ++ j)
                              W[i][j] = X[j];
45
46
                         return true;
                     } else {
47
48
                         long double t = X[i] / W[i][i];
                         for(int j = 1; j \leq n; ++ j)
49
50
                              X[j] = t * W[i][j];
51
                     }
                 }
52
53
            }
```

```
54
            return false;
        }
55
56
   // == TEST ===
57
   int gread();
58
59
   const int MAXN = 500 + 3;
   long double X[MAXN][MAXN], C[MAXN];
60
   int I[MAXN];
61
   bool cmp(int a, int b){
62
        return C[a] < C[b];</pre>
63
64
65
   int main(){
66
        int n, m;
67
        cin >> n >> m;
        realLB :: init(m);
68
69
        for(int i = 1; i \leq n; ++ i){
70
            for(int j = 1; j \leq m; ++ j){
71
                 cin >> X[i][j];
            }
72
73
        for(int i = 1; i \leq n; ++ i){
74
75
            cin \gg C[i];
            I[i] = i;
76
        }
77
        sort(I + 1, I + 1 + n, cmp);
78
79
        int ans = 0, cnt = 0;
80
        for(int i = 1; i \leq n; ++ i){}
            int x = I[i];
81
82
            if(realLB :: insert(X[x]))
83
                 ans += C[x],
                 cnt += 1;
84
85
        cout << cnt << " " << ans << endl;
86
87
        return 0;
88
   }
```

#### 2.5 Link Cut 树

```
1
   #include<bits/stdc++.h>
   using namespace std;
   typedef long long i64;
3
   namespace LinkCutTree{
4
       const int SIZ = 1e5 + 3;
5
6
       int F[SIZ], C[SIZ], S[SIZ], W[SIZ], A[SIZ], X[SIZ][2], size;
7
       bool T[SIZ]:
       bool is_root(int x){ return X[F[x]][0] \neq x \& X[F[x]][1] \neq x;}
8
9
       bool is_rson(int x){ return X[F[x]][1] = x;}
10
       int new_node(int w){
11
            ++ size;
           W[size] = w, C[size] = S[size] = 1;
12
           A[size] = w, F[size] = 0;
13
           X[size][0] = X[size][1] = 0;
14
15
           return size;
16
       void push_up(int x){
17
           S[x] = C[x] + S[X[x][0]] + S[X[x][1]];
18
           A[x] = W[x] ^ A[X[x][0]] ^ A[X[x][1]];
19
```

```
20
        void push_down(int x){
21
            if(!T[x]) return;
22
            int lc = X[x][0], rc = X[x][1];
23
            if(lc) T[lc] ^= 1, swap(X[lc][0], X[lc][1]);
24
            if(rc) T[rc] ^{\sim} 1, swap(X[rc][0], X[rc][1]);
25
            T[x] = false;
26
        }
27
        void update(int x){
28
            if(!is root(x)) update(F[x]); push down(x);
29
30
       void rotate(int x){
31
32
            int y = F[x], z = F[y];
            bool f = is rson(x);
33
            bool g = is rson(y);
34
35
            if(is_root(y)){
                F[x] = z;
36
37
                F[y] = x;
                X[y][f] = X[x][!f], F[X[x][!f]] = y;
38
39
                X[x][!f] = y;
            } else {
40
41
                F[x] = z;
                F[y] = x;
42
43
                X[z][g] = x;
                X[y][f] = X[x][!f], F[X[x][!f]] = y;
44
45
                X[x][!f] = y;
            }
46
47
            push_up(y), push_up(x);
48
        void splay(int x){
49
50
            update(x);
            for(int f = F[x]; f = F[x], !is_root(x); rotate(x))
51
                if(!is_root(f)) rotate(is_rson(x) = is_rson(f) ? f : x);
52
53
        int access(int x){
54
55
            int p;
56
            for(p = 0; x; p = x, x = F[x]){
57
                splay(x), X[x][1] = p, push_up(x);
58
59
            return p;
60
        void make_root(int x){
61
62
            x = access(x);
            T[x] \cong 1, swap(X[x][0], X[x][1]);
63
64
        int find root(int x){
65
            access(x), splay(x), push_down(x);
66
            while(X[x][0]) x = X[x][0], push_down(x);
67
68
            splay(x);
69
            return x;
70
        void link(int x, int y){
71
            make_{root}(x), splay(x), F[x] = y;
72
73
        }
        void cut(int x, int p){
74
75
            make\_root(x), access(p), splay(p), X[p][0] = F[x] = 0;
        }
76
```

```
void modify(int x, int w){
 77
 78
             splay(x), W[x] = w, push_up(x);
         }
 79
 80
    const int MAXN = 1e5 + 3;
 81
 82
    map<pair<int, int>, bool> M;
 83
    int n, m;
 84
    int main(){
 85
         cin >> n >> m;
 86
         for(int i = 1; i \leq n; ++ i){
 87
             int a; cin >> a;
 88
             LinkCutTree :: new_node(a);
 89
 90
         for(int i = 1; i \leq m; ++ i){}
             int o; cin >> o;
 91
 92
             if(0 = 0){
 93
                 int u, v; cin >> u >> v;
 94
                 LinkCutTree :: make_root(u);
 95
                 int p = LinkCutTree :: access(v);
                 printf("%d\n", LinkCutTree :: A[p]);
 96
             } else if(0 = 1){
 97
 98
                 int u, v; cin >> u >> v;
                 int a = LinkCutTree :: find_root(u);
 99
                 int b = LinkCutTree :: find_root(v);
100
                 if(a \neq b){
101
102
                     LinkCutTree :: link(u, v);
                      M[make_pair(min(u, v), max(u, v))] = true;
103
104
105
             } else if(0 = 2){
                 int u, v; cin >> u >> v;
106
                 if(M.count(make_pair(min(u, v), max(u, v)))){
107
                     M.erase(make_pair(min(u, v), max(u, v)));
108
                      LinkCutTree :: cut(u, v);
109
110
             } else {
111
                 int u, w; cin >> u >> w;
112
113
                 LinkCutTree :: modify(u, w);
114
             }
115
116
         return 0;
117
```

#### 2.6 线段树

#### 2.6.1 李超树

```
#include<bits/stdc++.h>
2 using namespace std;
3 typedef long long i64;
   struct Line{ int id; double k, b; Line() = default;};
5
   namespace LCSeg{
       const int SIZ = 2e5 + 3;
6
7
       struct Line T[SIZ];
       #define lc(t) (t << 1)
8
9
       #define rc(t) (t \ll 1 | 1)
       bool cmp(int p, Line x, Line y){
10
           double w1 = x.k * p + x.b;
11
```

```
double w2 = y.k * p + y.b;
12
           double d = w1 - w2;
13
           if(fabs(d) < 1e-8) return x.id > y.id;
14
15
           return d < 0;</pre>
16
       void merge(int t, int a, int b, Line x, Line y){
17
            int c = a + b >> 1;
18
19
            if(cmp(c, x, y)) swap(x, y);
           if(cmp(a, y, x)){
20
                T[t] = x; if(a \neq b) merge(rc(t), c + 1, b, T[rc(t)], y);
21
22
            } else {
23
                T[t] = x; if(a \neq b) merge(lc(t), a, c, T[lc(t)], y);
24
25
       void modify(int t, int a, int b, int l, int r, Line x){
26
27
           if(l ≤ a & b ≤ r) merge(t, a, b, T[t], x);
28
           else {
29
                int c = a + b >> 1;
                if(l \le c) modify(lc(t), a, c, l, r, x);
30
                if(r > c) modify(rc(t), c + 1, b, l, r, x);
31
32
33
       }
34
       void query(int t, int a, int b, int p, Line &x){
35
           if(cmp(p, x, T[t])) x = T[t];
           if(a \neq b){
36
37
                int c = a + b >> 1;
                if(p \le c) query(lc(t), a, c, p, x);
38
39
                if(p > c) query(rc(t), c + 1, b, p, x);
40
            }
       }
41
42
   const int MOD1 = 39989;
43
   const int MOD2 = 1e9;
44
   int gread();
45
   int m = 39989, o;
46
   int main(){
47
       int n = qread(), last_ans = 0;
48
49
       for(int i = 1; i \leq n; ++ i){}
            int op = gread(); if(op = \emptyset){
50
51
                int k = (qread() + last_ans - 1) % MOD1 + 1;
52
                Line x = \{0, 0, 0\}; LCSeg :: query(1, 1, m, k, x);
                printf("%d\n", last_ans = x.id);
53
            } else {
54
                int _x1 = (qread() + last_ans - 1) % MOD1 + 1;
55
                int _y1 = (qread() + last_ans - 1) % MOD2 + 1;
56
57
                int _x2 = (qread() + last_ans - 1) % MOD1 + 1;
                int _y2 = (qread() + last_ans - 1) % MOD2 + 1;
58
                if(x1 > x2) swap(x1, x2), swap(y1, y2);
59
                double k, b; int d = ++ o;
60
                if(_x1 = _x2) k = 0, b = max(_y1, _y2);
61
                    else k = 1.0 * (_y2 - _y1) / (_x2 - _x1), b = _y1 - k *
62
                        _x1;
                Line x = \{d, k, b\}; LCSeg :: modify(1, 1, m, _x1, _x2, x);
63
64
            }
       }
65
66
       return 0;
67
```

#### 2.6.2 线段树 3

```
#include<bits/stdc++.h>
 2 using namespace std;
   typedef long long i64;
 4 const int INF = 2e9;
 5 | const int MAXN= 5e5 + 3;
 6 | int A[MAXN];
7
   struct Node{
        i64 sum; int len, max1, max2, max_cnt, his_mx;
 8
 9
       Node():
            sum(0), max1(-INF), max2(-INF), max_cnt(0), his_mx(-INF), len
10
               (0) \{ \}
        Node(int w):
11
            sum(w), max1( w), max2(-INF), max_cnt(1), his_mx(
12
                                                                     w), len
               (1) \{\}
13
        bool update(int w1, int w2, int h1, int h2){
            his_mx = max(\{his_mx, max1 + h1\});
14
15
            \max 1 += w1, \max 2 += w2;
16
            sum += 1ll * w1 * max_cnt + 1ll * w2 * (len - max_cnt);
17
            return max1 > max2;
        }
18
   };
19
20
   struct Tag{
        int max_add, max_his_add, umx_add, umx_his_add; bool have;
21
22
        void update(int w1, int w2, int h1, int h2){
            \max his add = \max(\max his add, \max add + h1);
23
24
            umx_his_add = max(umx_his_add, umx_add + h2);
25
            max_add += w1, umx_add += w2, have = true;
        }
26
27
       void clear(){
28
            max add = max his add = umx add = umx his add = have = 0;
29
        }
30
   };
31
   struct Node operator +(Node a, Node b){
32
       Node t;
        t.max1 = max(a.max1, b.max1);
33
34
        if(t.max1 \neq a.max1){
            if(a.max1 > t.max2) t.max2 = a.max1;
35
36
        } else{
37
            if(a.max2 > t.max2) t.max2 = a.max2;
            t.max_cnt += a.max_cnt;
38
39
        if(t.max1 \neq b.max1){
40
            if(b.max1 > t.max2) t.max2 = b.max1;
41
42
43
            if(b.max2 > t.max2) t.max2 = b.max2;
44
            t.max_cnt += b.max_cnt;
45
        }
46
       t.sum = a.sum + b.sum, t.len = a.len + b.len;
        t.his mx = max(a.his mx, b.his mx);
47
48
       return t;
49
50
   namespace Seg{
51
       const int SIZ = 2e6 + 3;
        struct Node W[SIZ]; struct Tag T[SIZ];
52
        \#define lc(t) (t << 1)
53
54
        #define rc(t) (t \ll 1 | 1)
```

```
void push_up(int t, int a, int b){
55
56
            W[t] = W[lc(t)] + W[rc(t)];
 57
        void push down(int t, int a, int b){
 58
             if(a = b) T[t].clear();
59
             if(T[t].have){
60
                 int c = a + b \gg 1, x = lc(t), y = rc(t);
61
                 int w = max(W[x].max1, W[y].max1);
62
                 int w1 = T[t].max add, w2 = T[t].umx add, w3 = T[t].
63
                    max his add, w4 = T[t].umx his add;
64
                 if(w = W[x].max1)
                     W[x].update(w1, w2, w3, w4),
65
66
                     T[x].update(w1, w2, w3, w4);
67
                 else
                     W[x].update(w2, w2, w4, w4),
 68
69
                     T[x].update(w2, w2, w4, w4);
70
                 if(w = W[y].max1)
                     W[y].update(w1, w2, w3, w4),
71
                     T[y].update(w1, w2, w3, w4);
72
 73
                 else
74
                     W[y].update(w2, w2, w4, w4),
75
                     T[y].update(w2, w2, w4, w4);
76
                 T[t].clear();
             }
77
78
79
        void build(int t, int a, int b){
             if(a = b){W[t] = Node(A[a]), T[t].clear();} else {
80
                 int c = a + b >> 1; T[t].clear();
81
82
                 build(lc(t), a,
                                  c);
83
                 build(rc(t), c + 1, b);
                 push_up(t, a, b);
84
             }
85
86
        void modiadd(int t, int a, int b, int l, int r, int w){
87
             if(l \leq a \& b \leq r){
88
                 T[t].update(w, w, w, w);
89
90
                 W[t].update(w, w, w, w);
91
             } else {
                 int c = a + b >> 1; push down(t, a, b);
92
93
                 if(l \le c) \mod iadd(lc(t), a, c, l, r, w);
                 if(r > c) modiadd(rc(t), c + 1, b, l, r, w);
94
95
                 push_up(t, a, b);
             }
96
97
        void modimin(int t, int a, int b, int l, int r, int w){
98
99
             if(l \leq a \& b \leq r){
                 if(w > W[t].max1) return; else
100
                 if(w > W[t].max2){
101
                     int k = w - W[t].max1;
102
                     T[t].update(k, 0, k, 0);
103
                     W[t].update(k, 0, k, 0);
104
                 } else {
105
106
                     int c = a + b >> 1;
107
                     push_down(t, a, b);
                     modimin(lc(t), a,
                                          c, l, r, w);
108
109
                     modimin(rc(t), c + 1, b, l, r, w);
                     push_up(t, a, b);
110
```

```
111
             } else {
112
                 int c = a + b >> 1; push_down(t, a, b);
113
114
                 if(l \le c) modimin(lc(t), a, c, l, r, w);
                 if(r > c) modimin(rc(t), c + 1, b, l, r, w);
115
116
                 push_up(t, a, b);
             }
117
118
        Node query(int t, int a, int b, int l, int r){
119
             if(l \le a \& b \le r) return W[t];
120
121
             int c = a + b >> 1; Node ret; push_down(t, a, b);
            if(l \le c) ret = ret + query(lc(t), a, c, l, r);
122
123
            if(r > c) ret = ret + query(rc(t), c + 1, b, l, r);
124
             return ret;
        }
125
126
127
    int qread();
128
    int main(){
        int n = gread(), m = gread();
129
130
        for(int i = 1; i \leq n; ++ i)
            A[i] = qread();
131
132
        Seg :: build(1, 1, n);
        for(int i = 1; i \leq m; ++ i){
133
             int op = qread();
134
             if(op = 1){
135
136
                 int l = qread(), r = qread(), w = qread();
137
                 Seg :: modiadd(1, 1, n, l, r, w);
             } else if(op = 2){
138
139
                 int l = qread(), r = qread(), w = qread();
                 Seg :: modimin(1, 1, n, l, r, w);
140
             } else if(op = 3){
141
                 int l = qread(), r = qread();
142
                 auto p = Seg :: query(1, 1, n, l, r);
143
                 printf("%lld\n", p.sum);
144
             else if(op = 4)
145
                 int l = qread(), r = qread();
146
                 auto p = Seg :: query(1, 1, n, l, r);
147
148
                 printf("%d\n", p.max1);
             else if(op = 5)
149
150
                 int l = qread(), r = qread();
                 auto p = Seg :: query(1, 1, n, l, r);
151
                 printf("%d\n", p.his_mx);
152
153
154
155
        return 0;
156
```

## 2.6.3 扫描线

```
#include bits / stdc ++ . h>
using namespace std;
using i64 = long long;
const int INF = 1e9;
const i64 INFL = 1e18;
int qread(){
   int w = 1, c, ret;
```

```
while((c = getchar()) > '9' || c < '0') w = (c = '-' ? -1 : 1);
 8
           ret = c - '0';
       while((c = getchar()) \geq '0' & c \leq '9') ret = ret * 10 + c - '0';
9
10
       return ret * w;
11
12
   const int MAXN = 1e5 + 3;
   int X1[MAXN], Y1[MAXN];
13
14
   int X2[MAXN], Y2[MAXN];
   int n, h, H[MAXN * 2];
15
   namespace Seg{
16
17
       #define lc(t) (t \ll 1)
        #define rc(t) (t \ll 1 | 1)
18
19
       const int SIZ = 8e5 + 3;
        int T[SIZ], S[SIZ], L[SIZ];
20
        void pushup(int t, int a, int b){
21
22
            S[t] = 0;
23
            if(a \neq b){
                S[t] = S[lc(t)] + S[rc(t)];
24
                L[t] = L[lc(t)] + L[rc(t)];
25
26
27
            if(T[t]) S[t] = L[t];
28
       void modify(int t, int a, int b, int l, int r, int w){
29
30
            if(l \leq a \& b \leq r){
31
                T[t] += w, pushup(t, a, b);
32
            } else {
33
                int c = a + b >> 1;
                if(l \leq c) modify(lc(t), a, c, l, r, w);
34
35
                if(r > c) modify(rc(t), c + 1, b, l, r, w);
                pushup(t, a, b);
36
37
38
       void build(int t, int a, int b){
39
40
            if(a = b){
                L[t] = H[a] - H[a - 1];
41
42
            } else {
                int c = a + b >> 1;
43
44
                build(lc(t), a, c);
                build(rc(t), c + 1, b);
45
46
                pushup(t, a, b);
            }
47
48
       int query(int t){
49
50
            return S[t];
        }
51
52
   tuple <int, int, int> P[MAXN];
53
   tuple <int, int, int> Q[MAXN];
54
   int main(){
55
56
       n = qread();
        for(int i = 1; i \leq n; ++ i){
57
            X1[i] = qread(), Y1[i] = qread();
58
            X2[i] = qread(), Y2[i] = qread();
59
60
            if(X1[i] > X2[i]) swap(X1[i], X2[i]);
            if(Y1[i] > Y2[i]) swap(Y1[i], Y2[i]);
61
62
            H[++h] = Y1[i];
            H[++h] = Y2[i];
63
```

```
P[i] = make_tuple(X1[i], Y1[i], Y2[i]);
64
            Q[i] = make_tuple(X2[i], Y1[i], Y2[i]);
65
        }
66
67
        sort(H + 1, H + 1 + h);
        sort(P + 1, P + 1 + n);
68
69
        sort(Q + 1, Q + 1 + n);
        int o = unique(H + 1, H + 1 + h) - H - 1;
70
71
        Seg :: build(1, 1, o);
        i64 \text{ ans} = 0, \text{ last} = -1;
72
73
        int p = 1, q = 1;
74
        while (p \leq n \mid | q \leq n)
            int x = INF;
75
76
            if(p \leq n) x = min(x, get<0>(P[p]));
77
            if(q \leq n) x = min(x, get<0>(Q[q]));
            if(last \neq -1){
78
79
                ans += 1ll * Seg :: query(1) * (x - last);
80
            }
81
            last = x;
            while(q \leq n \&\& get<0>(Q[q]) = x){
82
83
                auto [x, l, r] = Q[q]; ++ q;
                l = lower_bound(H + 1, H + 1 + 0, l) - H + 1;
84
85
                r = lower_bound(H + 1, H + 1 + 0, r) - H;
                Seg :: modify(1, 1, o, l, r, 1);
86
87
            while(p \le n \& get<0>(P[p]) = x){
88
                auto [x, l, r] = P[p]; ++ p;
89
90
                l = lower_bound(H + 1, H + 1 + 0, l) - H + 1;
                r = lower_bound(H + 1, H + 1 + 0, r) - H;
91
92
                Seg :: modify(1, 1, o, l, r, -1);
            }
93
94
95
        printf("%lld\n", ans);
96
        return 0;
97
```

## 2.7 根号数据结构

#### 2.7.1 块状链表

```
#include <bits/stdc++.h>
 2
   using namespace std;
 3
   using i64 = long long;
   const int INF = 1e9;
   const i64 INFL = 1e18;
 5
 6
   namespace BLOCK{
 7
        const int SIZ = 1e6 + 1e5 + 3;
        const int BSZ = 2000;
 8
9
        list <vector<int> > block;
        void build(int n, const int A[]){
10
            for(int l = 0, r = 0; r \neq n;){
11
                l = r;
12
                r = min(l + BSZ / 2, n);
13
14
                vector \langle int \rangle V0(A + l, A + r);
                block.emplace back(V0);
15
16
            }
17
18
        int get_kth(int k){
```

```
for(auto it = block.begin();it ≠ block.end();++ it){
19
                   if(it \rightarrow size() < k)
20
                        k = it \rightarrow size();
21
22
                   else return it \rightarrow at(k - 1);
23
24
              return -1;
         }
25
26
         int get_rank(int w){
              int ans = 0;
27
              for(auto it = block.begin();it ≠ block.end();++ it){
28
29
                   if(it \rightarrow back() < w)
                        ans += it \rightarrow size();
30
31
                   else {
                        ans += lower bound(it \rightarrow begin(), it \rightarrow end(), w) - it
32
                            \rightarrow begin();
33
                        break;
34
                   }
35
36
              return ans + 1;
37
         }
         // 插入到第 k 个位置
38
39
         void insert(int k, int w){
              for(auto it = block.begin();it ≠ block.end();++ it){
40
                   if(it \rightarrow size() < k)
41
                        k = it \rightarrow size();
42
43
                   else{
                        it \rightarrow insert(it \rightarrow begin() + k - 1, w);
44
45
                        if(it \rightarrow size() > BSZ){
46
                             vector \langle int \rangle V1(it \rightarrow begin(), it \rightarrow begin() + BSZ
47
                             vector \langle int \rangle V2(it \rightarrow begin() + BSZ / 2, it \rightarrow end
                                 ());
                             *it = V2;
48
                             block.insert(it, V1);
49
                        }
50
51
                        return;
52
                   }
              }
53
         }
54
55
         // 删除第 k 个数
         void erase(int k){
56
              for(auto it = block.begin();it ≠ block.end();++ it){
57
58
                   if(it \rightarrow size() < k)
                        k = it \rightarrow size();
59
60
                   else{
                        it \rightarrow erase(it \rightarrow begin() + k - 1);
61
                        if(it \rightarrow empty())
62
                             block.erase(it);
63
64
                        return;
                   }
65
              }
66
         }
67
68
69
    int gread();
    const int MAXN = 1e5 + 3;
70
    int A[MAXN];
72 |// ≡= TEST ≡=
```

```
int main(){
 73
         ios :: sync_with_stdio(false);
 74
 75
         cin.tie(nullptr);
 76
         int n, m;
 77
         cin >> n >> m;
 78
         for(int i = 1; i \leq n; ++ i)
             cin \gg A[i];
 79
 80
         sort(A + 1, A + 1 + n);
         A[n + 1] = INT MAX;
 81
         BLOCK :: build(n + 1, A + 1);
 82
         int last = 0;
 83
         int ans = 0;
 84
 85
         for(int i = 1; i \leq m; ++ i){
 86
             int op;
             cin >> op;
 87
             if(op = 1){
 88
                  int x; cin >> x; x ^= last;
 89
 90
                  int k = BLOCK :: get_rank(x);
                  BLOCK :: insert(k, x);
 91
 92
             } else
             if(op = 2){
 93
 94
                  int x; cin \gg x; x ^- last;
                  int k = BLOCK :: get_rank(x);
 95
                  BLOCK :: erase(k);
 96
 97
             } else
 98
             if(op = 3){
 99
                  int x; cin \gg x; x ^{\sim} last;
                  int k = BLOCK :: get_rank(x);
100
101
                  last = k, ans ^= last;
             } else
102
             if(op = 4){
103
                  int x; cin \gg x; x ^{\sim} last;
104
                  int k = BLOCK :: get_kth (x);
105
106
                  last = k, ans ^{\sim} last;
             } else
107
             if(op = 5){
108
109
                  int x; cin \gg x; x ^- last;
110
                  int k = BLOCK :: get_rank(x);
                  last = BLOCK :: get_kth (k - 1), ans ^{\sim} last;
111
112
             } else
             if(op = 6){
113
                  int x; cin >> x; x ^= last;
114
115
                  int k = BLOCK :: get rank(x + 1);
                  last = BLOCK :: get_kth (k), ans ^= last;
116
             }
117
118
         cout << ans << endl;
119
120
         return 0;
121
```

#### 2.7.2 莫队二次离线

```
#include<bits/stdc++.h>
#define up(l, r, i) for(int i = l, END##i = r;i ≤ END##i;++ i)
#define dn(r, l, i) for(int i = r, END##i = l;i ≥ END##i;-- i)
using namespace std;
typedef long long i64;
```

```
const int INF = 2147483647;
 7
   const int MAXN= 1e5 + 3;
8 | const int MAXM= (1 << 14) + 3;
9 int n, m, k, maxt = 16383, X[MAXM], C[MAXM], t;
10 | const int BUF_SIZE = 1e6;
11 | char *p1, *p2, BUF[BUF SIZE];
12 | inline char readc();
13 | inline int gread();
14 int A[MAXN], bsize; i64 B[MAXN], R[MAXN];
   struct Qry1{ int l, r, id; }0[MAXN];
15
   struct Qry2{ int id, l, r; };
16
17
   struct Qry3{ int id, l, r; };
18
   bool cmp(Qry1 a, Qry1 b){
       return a.l / bsize = b.l / bsize ? a.r < b.r : a.l < b.l;</pre>
19
20
21
   vector <Qry2> P[MAXN];
22
   vector <Qry3> Q[MAXN];
23
   int main(){
        n = gread(), m = gread(), k = gread(), bsize = sgrt(m + 1);
24
25
        up(1, n, i) A[i] = qread();
26
        up(1, m, i){
27
            int l = gread(), r = gread(); 0[i] = {l, r, i};
28
29
       sort(0 + 1, 0 + 1 + m, cmp);
30
        int l = 1, r = 0;
31
        up(1, m, i){
32
            int p = 0[i].l, q = 0[i].r;
            if(r < q){
33
34
                P[r
                       ].push_back({ i, r + 1, q});
35
                Q[l - 1].push_back({-i, r + 1, q});
36
           if(r > q){
37
38
                       ].push_back(\{-i, q + 1, r\});
39
                Q[l - 1].push_back({ i, q + 1, r});
            }
40
41
           r = q;
42
           if(l > p){
                P[p].push_back({-i, p, l - 1});
43
                Q[r].push_back({i, p, l - 1});
44
45
            }
            if(l < p){
46
47
                P[l].push_back({ i, l, p - 1});
48
                Q[r].push_back({-i, l, p - 1});
49
            l = p;
50
51
        up(0, maxt, i) if(\_builtin\_popcount(i) = k) X[++ t] = i;
52
53
        up(0, n, i){
            up(1, t, j) + C[A[i] ^ X[j]];
54
55
            for(auto &o : P[i]){
                if(o.id > 0) R[ o.id] += C[A[o.l]];
56
                              R[-o.id] -= C[A[o.l]];
57
                else
58
                if(o.l < o.r)
59
                    P[i + 1].push_back({o.id, o.l + 1, o.r});
60
            for(auto &o : Q[i]){
61
                up(o.l, o.r, j){
62
```

3 树论 35

```
if(o.id > 0) R[ o.id] += C[A[j]];
63
64
                                   R[-o.id] -= C[A[j]];
                 }
65
            }
66
            P[i].clear(), Q[i].clear();
67
68
            P[i].shrink_to_fit();
            Q[i].shrink_to_fit();
69
        }
70
        i64 \text{ ans} = 0;
71
        up(1, m, i){ ans += R[i], B[O[i].id] = ans; }
72
        up(1, m, i) printf("%lld\n", B[i]);
73
74
        return 0;
75
   }
```

## 3 树论

#### 3.1 点分树

#### 3.1.1 例题

给定 n 个点组成的树, 点有点权  $v_i$ 。m 个操作, 分为两种:

- $0 \times k$  查询距离 x 不超过 k 的所有点的点权之和;
- 0 x y 将点 x 的点权修改为 y。

```
#include<bits/stdc++.h>
   #define endl "\n"
 2
   using namespace std;
   const int MAXN = 1e5 + 3;
   vector<int> E[MAXN];
   namespace LCA{
 6
 7
        const int SIZ = 1e5 + 3;
        int D[SIZ], F[SIZ];
8
9
        int P[SIZ], Q[SIZ], o;
        void dfs(int u, int f){
10
11
            P[u] = ++ o;
            Q[o] = u;
12
            F[u] = f;
13
14
            D[u] = D[f] + 1;
15
            for(auto &v : E[u]) if(v \neq f){
                dfs(v, u);
16
17
        }
18
19
        const int MAXH = 18 + 3;
20
        int h = 18;
        int ST[SIZ][MAXH];
21
        int cmp(int a, int b){
22
23
            return D[a] < D[b] ? a : b;
24
25
        int T[SIZ], n;
26
        void init(int _n){
27
            n = _n;
            dfs(1, 0);
28
29
            for(int i = 1; i \leq n; ++ i)
                ST[i][0] = Q[i];
30
31
            for(int i = 2; i \leq n; ++ i)
```

3 树论 36

```
T[i] = T[i >> 1] + 1;
32
            for(int i = 1; i \leq h; ++ i){
33
                for(int j = 1; j \le n; ++ j) if(j + (1 << i - 1) \le n){
34
                     ST[j][i] = cmp(ST[j][i - 1], ST[j + (1 << i - 1)][i -
35
                        1]);
36
                }
            }
37
38
39
        int lca(int a, int b){
            if(a = b)
40
41
                return a;
            int l = P[a];
42
43
            int r = P[b];
            if(l > r)
44
45
                swap(l, r);
            ++ l;
46
            int d = T[r - l + 1];
47
            return F[cmp(ST[l][d], ST[r - (1 << d) + 1][d])];
48
49
50
        int dis(int a, int b){
            return D[a] + D[b] - 2 * D[lca(a, b)];
51
52
        }
53
   namespace BIT{
54
        void modify(int D[], int n, int p, int w){
55
56
            ++ p;
57
            while(p \leq n)
58
                D[p] += w, p += p & -p;
59
60
        int query(int D[], int n, int p){
            if(p < 0) return 0;
61
62
            p = min(n, p + 1);
            int r = 0;
63
64
            while (p > 0)
65
                r += D[p], p -= p & -p;
66
            return r;
        }
67
68
   namespace PTree{
69
70
        const int SIZ = 1e5 + 3;
        bool V[SIZ];
71
        int S[SIZ], L[SIZ];
72
73
        vector<int> EE[MAXN];
        int *D1[MAXN];
74
        int *D2[MAXN];
75
        void dfs1(int s, int &g, int u, int f){
76
            S[u] = 1;
77
78
            int maxsize = 0;
79
            for(auto \&v : E[u]) if(v \neq f \&v !V[v]){
80
                dfs1(s, g, v, u);
                if(S[v] > maxsize)
81
                     maxsize = S[v];
82
83
                S[u] += S[v];
84
            }
            maxsize = max(maxsize, s - S[u]);
85
            if(maxsize \leq s / 2)
86
87
                g = u;
```

```
88
 89
         int n;
         void build(int s, int &g, int u, int f){
 90
 91
             dfs1(s, g, u, f);
             V[g] = true, L[g] = s;
 92
             for(auto &u : E[g]) if(!V[u]){
 93
                 int h = 0;
 94
 95
                 if(S[u] < S[g]) build(S[u], h, u, 0);
                                   build(s - S[g], h, u, 0);
 96
 97
                 EE[g].push_back(h);
 98
                 EE[h].push_back(g);
             }
 99
100
         }
101
         int F[SIZ];
         void dfs2(int u, int f){
102
             F[u] = f;
103
             for(auto &v : EE[u]) if(v \neq f){
104
105
                 dfs2(v, u);
             }
106
107
         void build(int _n){
108
109
             n = _n;
             int s = n, g = 0;
110
111
             dfs1(s, g, 1, 0);
             V[g] = true, L[g] = s;
112
113
             for(auto &u : E[g]){
114
                 int h = 0;
                 if(S[u] < S[g]) build(S[u], h, u, 0);
115
116
                                   build(s - S[g], h, u, 0);
                 EE[g].push_back(h);
117
                 EE[h].push_back(g);
118
119
             dfs2(g, 0);
120
             for(int i = 1; i \leq n; ++ i){
121
122
                 L[i] += 2;
                 D1[i] = new int[L[i] + 3];
123
                 D2[i] = new int[L[i] + 3];
124
125
                 for(int j = 0; j < L[i] + 3; ++ j)
                      D1[i][j] = D2[i][j] = 0;
126
127
             }
128
129
         void modify(int x, int w){
130
             int u = x;
             while(1){
131
                 BIT :: modify(D1[x], L[x], LCA :: dis(u, x), w);
132
                 int y = F[x];
133
                 if(y \neq 0){
134
135
                      int e = LCA :: dis(x, y);
136
                      BIT :: modify(D2[x], L[x], LCA :: dis(u, y), w);
137
                      x = y;
138
                 } else break;
             }
139
140
141
         int query(int x, int d){
142
             int ans = 0, u = x;
143
             while(1){
                 ans += BIT :: query(D1[x], L[x], d - LCA :: dis(u, x));
144
```

```
145
                  int y = F[x];
146
                  if(y \neq 0){
147
                      int e = LCA :: dis(x, y);
148
                      ans -= BIT :: query(D2[x], L[x], d - LCA :: dis(u, y));
                      x = y;
149
150
                  } else break;
             }
151
152
             return ans;
         }
153
154
155
    int W[MAXN];
    int main(){
156
157
         ios :: sync_with_stdio(false);
158
         int n, m;
159
         cin >> n >> m;
         for(int i = 1; i \leq n; ++ i){
160
161
             cin >> W[i];
162
163
         for(int i = 2; i \leq n; ++ i){}
             int u, v;
164
165
             cin \gg u \gg v;
166
             E[u].push_back(v);
167
             E[v].push_back(u);
         }
168
169
         LCA :: init(n);
170
         PTree :: build(n);
171
         for(int i = 1; i \leq n; ++ i)
             PTree :: modify(i, W[i]);
172
173
         int lastans = 0;
174
         for(int i = 1; i \leq m; ++ i){
             int op; cin >> op;
175
             if(op = 0){
176
177
                  int x, d;
178
                  cin >> x >> d;
179
                  x \cong lastans;
                  d ^= lastans;
180
                  cout << (lastans = PTree :: query(x, d)) << endl;
181
182
             } else {
                  int x, w;
183
184
                  cin >> x >> w;
                  x ~ lastans;
185
                  w ~ lastans;
186
187
                  PTree :: modify(x, -W[x]
                  PTree :: modify(x, W[x] = w);
188
             }
189
190
         }
191
         return 0;
192
```

#### 3.2 长链剖分

```
#include<bits/stdc++.h>
using namespace std;
using i64 = long long;
const int INF = 1e9;
const i64 INFL = 1e18;
const int MAXN= 5e5 + 3;
```

```
const int MAXM= 19 + 3;
 7
 8
   vector <int> P[MAXN];
9 | vector <int> Q[MAXN];
10 | vector <int> E[MAXN];
   int h = 19;
11
   int L[MAXN], F[MAXN], G[MAXN], D[MAXN], S[MAXM][MAXN];
12
   void dfs1(int u, int f){
13
14
        L[u] = 1, S[0][u] = f;
15
        F[u] = f, D[u] = D[f] + 1;
        for(int i = 1; i \leq h; ++ i)
16
17
            S[i][u] = S[i - 1][S[i - 1][u]];
        for(auto &v : E[u]) if(v \neq f){
18
19
            dfs1(v, u);
            if(L[v] > L[G[u]])
20
21
                G[u] = v;
22
            L[u] = max(L[u], L[v] + 1);
        }
23
24
25
   int T[MAXN];
26
   void dfs2(int u, int f){
27
        if(u = G[f]){
28
            T[u] = T[f];
            P[T[u]].push back(u);
29
30
            Q[T[u]].push_back(F[Q[T[u]].back()]);
31
        } else {
32
            T[u] = u;
33
            P[u].push back(u);
34
            Q[u].push_back(u);
35
36
        if(G[u]) dfs2(G[u], u);
37
        for(auto &v : E[u]) if(v \neq f \& v \neq G[u])
38
            dfs2(v, u);
39
40
   typedef unsigned int
                                u32:
41
   typedef unsigned long long u64;
42
   int n, q; u32 s;
43
   u32 get(u32 x) {
44
       x ^- x << 13;
        x ^ x > 17;
45
46
        x ^ x < 5;
47
       return s = x;
48
49
   int gread();
50
   int H[MAXN];
51
   int main(){
52
        scanf("%d%d%u", &n, &q, &s);
53
        int root = 0; H[0] = -1;
54
        for(int i = 1; i \leq n; ++ i){
55
            int f = qread();
56
            if(f = 0)
57
                root = i;
58
            else {
59
                E[f].push_back(i);
60
                E[i].push_back(f);
61
62
            H[i] = H[i >> 1] + 1;
        }
63
```

```
64
        dfs1(root, 0);
65
        dfs2(root, 0);
        int lastans = 0;
66
67
        i64 \text{ realans} = 0;
        for(int i = 1; i \leq q; ++ i){
68
             int x = (get(s) ^ lastans) % n + 1;
69
             int k = (get(s) ^ lastans) % D[x];
70
71
            if(k = 0){
72
                 lastans = x;
73
             } else {
74
                 int h = H[k];
                 k -= 1 << h;
75
76
                 x = S[h][x];
77
                 int t = T[x];
                 k \rightarrow D[x] - D[t];
78
79
                 if(k > 0){
80
                     x = Q[t][k];
                 } else {
81
82
                     x = P[t][-k];
83
84
                 lastans = x;
85
             }
            realans ~ 1ll * i * lastans;
86
87
        printf("%lld\n", realans);
88
89
        return 0;
90
```

## 3.3 重链剖分

```
#include<bits/stdc++.h>
   using namespace std;
 3 | using i64 = long long;
   const int INF = 1e9;
   const i64 INFL = 1e18;
 5
   const int MAXN= 1e5 + 3;
 6
 7
   int MOD;
   int n, m, root;
 8
9
   int A[MAXN];
   int gread();
10
   vector <int> E[MAXN];
11
   int S[MAXN], G[MAXN], D[MAXN], F[MAXN];
12
   void dfs1(int u, int f){
13
14
       S[u] = 1, G[u] = 0, D[u] = D[f] + 1, F[u] = f;
15
        for(auto \delta v : E[u]) if(v \neq f){
            dfs1(v, u);
16
            S[u] += S[v];
17
18
            if(S[v] > S[G[u]])
19
                G[u] = v;
        }
20
21
   int B[MAXN];
22
   int P[MAXN], Q[MAXN], T[MAXN], L[MAXN], R[MAXN], cnt;
23
24
   void dfs2(int u, int f){
25
        P[++ cnt] = u, B[cnt] = A[u], Q[u] = cnt;
       L[u] = cnt;
26
27
       if(u \neq G[f]) T[u] = u;
```

```
T[u] = T[f];
28
            else
        if(G[u]) dfs2(G[u], u);
29
        for(auto \&v : E[u]) if(v \neq f \&v \neq G[u]){
30
31
            dfs2(v, u);
32
33
       R[u] = cnt;
34
35
   namespace Seg{
       #define lc(t) (t << 1)
36
       #define rc(t) (t << 1 | 1)
37
       const int SIZ = 4e5 + 3;
38
        i64 S[SIZ], T[SIZ];
39
40
        void pushup(int t, int a, int b){
            S[t] = (S[lc(t)] + S[rc(t)]) % MOD;
41
42
       void pushdown(int t, int a, int b){
43
44
            if(T[t]){
                int c = a + b >> 1;
45
                T[lc(t)] = (T[lc(t)] + T[t]) \% MOD;
46
47
                T[rc(t)] = (T[rc(t)] + T[t]) \% MOD;
                S[lc(t)] = (S[lc(t)] + 1ull * (c - a + 1) * T[t]) % MOD;
48
49
                S[rc(t)] = (S[rc(t)] + 1ull * (b - c) * T[t]) % MOD;
                T[t] = 0;
50
            }
51
52
       void modify(int t, int a, int b, int l, int r, int w){
53
            if(1 \le a \& b \le r)
54
                S[t] = (S[t] + 1ll * w * (b - a + 1)) % MOD;
55
56
                T[t] = (T[t] + w) \% MOD;
57
            } else {
                int c = a + b >> 1;
58
59
                pushdown(t, a, b);
                if(l \leq c) modify(lc(t), a, c, l, r, w);
60
                if(r > c) modify(rc(t), c + 1, b, l, r, w);
61
                pushup(t, a, b);
62
            }
63
64
65
        i64 query(int t, int a, int b, int l, int r){
            if(1 \le a \& b \le r)
66
67
                return S[t];
68
            int c = a + b >> 1;
69
            i64 \text{ ans} = 0;
70
            pushdown(t, a, b);
            if(l \le c) ans = (ans + query(lc(t), a, c, l, r)) % MOD;
71
            if(r > c) ans = (ans + query(rc(t), c + 1, b, l, r)) % MOD;
72
73
            return ans;
74
       void build(int t, int a, int b){
75
76
            if(a = b){
77
                S[t] = B[a] \% MOD;
            } else {
78
                int c = a + b >> 1;
79
                build(lc(t), a, c);
80
81
                build(rc(t), c + 1, b);
82
                pushup(t, a, b);
83
            }
        }
84
```

```
85
 86
    int main(){
         n = qread(), m = qread(), root = qread(), MOD = qread();
 87
         for(int i = 1; i \leq n; ++ i)
 88
             A[i] = qread();
 89
 90
         for(int i = 2; i \leq n; ++ i){}
             int u = qread(), v = qread();
 91
 92
             E[u].push_back(v);
             E[v].push back(u);
 93
 94
 95
         dfs1(root, 0);
 96
         dfs2(root, 0);
 97
         Seg :: build(1, 1, n);
         for(int i = 1; i \leq m; ++ i){
 98
             int op = gread();
 99
             if(op = 1){
100
101
                 int u = qread(), v = qread(), k = qread();
                 while(T[u] \neq T[v]){
102
                      if(D[T[u]] < D[T[v]])
103
                          swap(u, v);
104
                      Seg :: modify(1, 1, n, Q[T[u]], Q[u], k);
105
106
                      u = F[T[u]];
                 }
107
                 if(D[u] < D[v]) swap(u, v);
108
109
                 Seg :: modify(1, 1, n, Q[v], Q[u], k);
110
             else if(op = 2)
                 int u = gread(), v = gread();
111
                 i64 \text{ ans} = 0;
112
113
                 while(T[u] \neq T[v]){
                      if(D[T[u]] < D[T[v]])
114
                          swap(u, v);
115
                      ans = (ans + Seg :: query(1, 1, n, Q[T[u]], Q[u])) %
116
                         MOD:
117
                      u = F[T[u]];
118
                 if(D[u] < D[v]) swap(u, v);
119
                 ans = (ans + Seg :: query(1, 1, n, Q[v], Q[u])) % MOD;
120
                 printf("%lld\n", ans);
121
             } else if(op = 3){
122
123
                 int x = qread(), w = qread();
                 Seg :: modify(1, 1, n, L[x], R[x], w);
124
             } else {
125
126
                 int x = qread();
                 printf("%lld\n", Seg :: query(1, 1, n, L[x], R[x]));
127
             }
128
129
130
         return 0;
131
```

#### 3.4 树哈希

#### 3.4.1 用法

给定大小为 n 的以 1 为根的树,计算  $h_i$  表示子树 i 的哈希值,计算有多少个本质不同的值。

```
#include<bits/stdc++.h>
using namespace std;
using u64 = unsigned long long;
```

```
const int MAXN = 1e6 + 3;
 5
   u64 xor_shift(u64 x){
        x ^ x < 33;
 6
        x ^ x > 7;
7
        x ^ x < 17;
 8
        return x;
 9
10
11
   u64 H[MAXN];
   vector <int> E[MAXN];
12
   void dfs(int u, int f){
13
14
       H[u] = 1;
        for(auto &v: E[u]) if(v \neq f){
15
16
            dfs(v, u);
17
            H[u] += H[v];
18
       H[u] = xor_shift(H[u]); // !important
19
20
21
   int main(){
22
        int n;
23
        cin >> n;
        for(int i = 2; i \leq n; ++ i){
24
25
            int u, v;
            cin >> u >> v;
26
27
            E[u].push_back(v);
28
            E[v].push_back(u);
        }
29
        dfs(1, 0);
30
31
        sort(H + 1, H + 1 + n);
32
        cout << (unique(H + 1, H + 1 + n) - H - 1) << endl;
33
        return 0;
34
```

### 3.5 Prufer 序列

#### 3.5.1 用法

给定大小为 n 的以 1 为根的树,计算  $h_i$  表示子树 i 的哈希值,计算有多少个本质不同的值。

```
#include<bits/stdc++.h>
1
2
  using namespace std;
  const int MAXN = 5e6 + 3;
   int D[MAXN], F[MAXN], P[MAXN];
5
   vector<int> tree2prufer(int n){
      vector <int> P(n);
6
      for(int i = 1, j = 1; i \le n - 2; ++ i, ++ j){
7
          while(D[j]) ++ j;
8
9
          P[i] = F[j];
         10
11
             P[i + 1] = F[P[i]], i ++;
      }
12
13
      return P;
14
   vector<int> prufer2tree(int n){
15
      vector <int> F(n);
16
      for(int i = 1, j = 1; i \leq n - 1; i \neq j){
17
18
          while(D[j]) ++ j;
          F[j] = P[i];
19
20
```

```
F[P[i]] = P[i + 1], i ++;
21
22
        }
23
        return F;
24
25
   int main(){
26
        ios :: sync_with_stdio(false);
27
        cin.tie(nullptr);
28
        int n, m;
29
        cin >> n >> m;
        vector <int> ANS;
30
31
        if(m = 1){
                      // tree \rightarrow prufer
32
            for(int i = 1; i \leq n - 1; ++ i){
33
                 cin \gg F[i], D[F[i]] ++;
            }
34
35
            ANS = tree2prufer(n);
36
        } else {
                         // prufer \rightarrow tree
37
            for(int i = 1; i \leq n - 2; ++ i){
38
                 cin \gg P[i], D[P[i]] ++;
39
40
            P[n - 1] = n;
            ANS = prufer2tree(n);
41
42
        long long ans = 0, cnt = 0;
43
44
        for(int i = 1; i \le n - (m = 1 ? 2 : 1); ++ i)
            ans ^= 1ll * ANS[i] * i;
45
46
        cout \ll ans \ll "\n";
47
        return 0;
48
```

## 3.6 虚树

```
1 #include < bits / stdc++.h>
   using namespace std;
 3 | const int MAXN = 5e5 + 3;
 4 | vector<pair<int, int> > E[MAXN];
   namespace LCA{
 5
 6
        const int SIZ = 5e5 + 3;
7
        int D[SIZ], H[SIZ], F[SIZ];
8
        int P[SIZ], Q[SIZ], o;
9
        void dfs(int u, int f){
10
            P[u] = ++ o;
            Q[o] = u;
11
            F[u] = f;
12
            D[u] = D[f] + 1;
13
            for(auto \delta[v, w] : E[u]) if(v \neq f){
14
15
                H[v] = H[u] + w, dfs(v, u);
16
17
        }
18
        const int MAXH = 18 + 3;
19
        int h = 18;
        int ST[SIZ][MAXH];
20
        int cmp(int a, int b){
21
22
            return D[a] < D[b] ? a : b;
23
        int T[SIZ], n;
24
        void init(int _n, int root){
25
26
            n = _n;
```

```
dfs(root, 0);
27
            for(int i = 1;i ≤ n; ++ i)
28
                ST[i][0] = Q[i];
29
30
            for(int i = 2; i \leq n; ++ i)
                T[i] = T[i >> 1] + 1;
31
32
            for(int i = 1; i \leq h; ++ i){}
                for(int j = 1; j \le n; ++ j) if(j + (1 << i - 1) \le n){
33
34
                     ST[j][i] = cmp(ST[j][i - 1], ST[j + (1 << i - 1)][i -
35
                }
            }
36
37
        }
38
        int lca(int a, int b){
39
            if(a = b)
40
                return a;
            int l = P[a];
41
42
            int r = P[b];
43
            if(l > r)
44
                swap(l, r);
45
            ++ l;
            int d = T[r - l + 1];
46
47
            return F[cmp(ST[l][d], ST[r - (1 << d) + 1][d])];
48
        int dis(int a, int b){
49
            return H[a] + H[b] - 2 * H[lca(a, b)];
50
51
52
   bool cmp(int a, int b){
53
54
        return LCA :: P[a] < LCA :: P[b];
55
   bool I[MAXN];
56
   vector <int> E1[MAXN];
57
   vector <int> V1;
58
   void solve(vector <int> &V){
59
        using LCA :: lca;
60
        using LCA :: D;
61
62
        stack <int> S;
        sort(V.begin(), V.end(), cmp);
63
        S.push(1);
64
65
        int v, l;
        for(auto &u : V) I[u] = true;
66
        for(auto u : V) if(u \neq 1)
67
68
            int f = lca(u, S.top());
69
            l = -1;
            while(D[v = S.top()] > D[f]){
70
71
                if(l \neq -1)
                     E1[v].push_back(l);
72
73
                V1.push_back(l = v), S.pop();
74
75
            if(l \neq -1)
                E1[f].push_back(l);
76
            if(f \neq S.top())
77
78
                S.push(f);
79
            S.push(u);
        }
80
        l = -1;
81
        while(!S.empty()){
82
```

```
83
            v = S.top();
            if(l \neq -1)
84
85
                E1[v].push_back(l);
86
            V1.push_back(l = v), S.pop();
87
88
        // dfs(1, 0); // SOLVE HERE !!!
        for(auto &u : V1)
89
90
            E1[u].clear(), I[u] = false;
91
        V1.clear();
92
```

# 4 图论

## 4.1 仙人掌

## 4.1.1 例题

给定一个仙人掌, 多组询问 u, v 之间最短路长度。

```
#include<bits/stdc++.h>
   #define up(l, r, i) for(int i = l, END##i = r; i \leq END##i; ++ i)
 2
   #define dn(r, l, i) for(int i = r, END##i = l; i \ge END##i; -- i)
   using namespace std;
   typedef long long i64;
 5
   const int INF = 2147483647;
 7
   const int MAXN= 2e5 + 3;
 8
   const int MAXM= 2e5 + 3;
   const int MAXD= 18 + 3;
9
10
   struct edge{int u, v, w;};
   vector <edge> V1[MAXN];
11
   vector <edge> V2[MAXN];
12
   vector <int> H[MAXN];
13
   int n, D[MAXN], W[MAXN], F[MAXD][MAXN];
14
   int o, X[MAXN], L[MAXN];
15
   bool E[MAXN];
16
   void dfs1(int u, int f){
17
       D[u] = D[f] + 1, F[0][u] = f;
18
19
        for(auto &e : V1[u]) if(e.v \neq f){
            if(D[e.v] & D[e.v] < D[u]){
20
21
                int a = e.u;
22
                int b = e.v;
23
                int c = ++ o, t = c + n;
24
                H[c].push_back(a);
                L[c] = W[a] - W[b] + e.w;
25
                while (a \neq b)
26
                    E[a] = true, a = F[0][a], H[c].push_back(a);
27
28
                for(auto &x : H[c]){
29
                    int w = min(W[x] - W[b], L[c] - W[x] + W[b]);
30
                    V2[x].push_back(edge{x, t, w});
31
                    V2[t].push_back(edge{t, x, w});
                }
32
            } else if(!D[e.v]){
33
                W[e.v] = W[u] + e.w, dfs1(e.v, u);
34
            }
35
36
37
        for(auto &e : V1[u]) if(D[e.v] > D[u]){
            if(!E[e.v]){
38
```

```
39
                V2[e.u].push_back({e.u, e.v, e.w});
                V2[e.v].push_back({e.v, e.u, e.w});
40
            }
41
       }
42
43
44
   int d = 18;
   void dfs2(int u, int f){
45
46
       D[u] = D[f] + 1, F[0][u] = f;
47
       up(1, d, i) F[i][u] = F[i - 1][F[i - 1][u]];
       for(auto &e : V2[u]) if(e.v \neq f){
48
           X[e.v] = X[e.u] + e.w;
49
50
           dfs2(e.v, u);
       }
51
52
   int lca(int u, int v){
53
       if(D[u] < D[v]) swap(u, v);
54
55
       dn(d, 0, i) if(D[F[i][u]] \ge D[v]) u = F[i][u];
       if(u = v) return u;
56
       dn(d, 0, i) if(F[i][u] \neq F[i][v]) u = F[i][u], v = F[i][v];
57
       return F[0][u];
58
59
60
   int jump(int u, int v){
       dn(d, 0, i) if(D[F[i][v]] > D[u]) v = F[i][v];
61
62
       return v;
63
   int dis(int x, int y){
64
       int t = lca(x, y);
65
66
       if(t > n){
67
           int u = jump(t, x);
           int v = jump(t, y);
68
           int w = abs(W[u] - W[v]);
69
70
            int l = min(w, L[t - n] - w);
71
           return X[x] - X[u] + X[y] - X[v] + 1;
72
       } else {
            return X[x] + X[y] - 2 * X[t];
73
74
       }
75
76
   int m, q;
   int gread();
77
78
   int main(){
       n = qread(), m = qread();
79
80
       up(1, m, i){
            int u = qread(), v = qread(), w = qread();
81
           V1[u].push_back(edge{u, v, w});
82
           V1[v].push_back(edge{v, u, w});
83
84
       dfs1(1, 0);
85
       dfs2(1, 0);
86
87
       up(1, q, i){
            int u = qread(), v = qread();
88
            printf("%d\n", dis(u, v));
89
90
91
       return 0;
92
```

#### 4.2 三元环计数

#### 4.2.1 三元环计数

无向图:考虑将所有点按度数从小往大排序,然后将每条边定向,由排在前面的指向排在后面的,得到一个有向图。然后考虑枚举一个点,再枚举一个点,暴力数,具体见代码。结论是,这样定向后,每个点的出度是  $O(\sqrt{m})$  的。复杂度  $O(m\sqrt{m})$ 。有向图:不难发现,上述方法枚举了三个点,计算有向图三元环也就只需要处理下方向的事,这个由于算法够暴力,随便改改就能做了。

```
1 // 无向图
2 | ll n, m; cin >> n >> m;
   vector<pair<ll, ll>> Edges(m);
   vector<vector<ll>>> G(n + 2);
4
5 | \text{vector} < \text{ll} > \text{deg}(n + 2);
   for (auto \delta[i, j]: Edges) cin \gg i \gg j, ++deg[i], ++deg[j];
6
   for (auto [i, j] : Edges) {
7
       if (deg[i] > deg[j] \parallel (deg[i] = deg[j] & i > j)) swap(i, j);
8
9
       G[i].emplace_back(j);
10
   vector<ll> val(n + 2);
11
   ll ans = 0;
12
14
       for (auto j : G[i]) ++val[j];
       for (auto j : G[i]) for (auto k : G[j]) ans += val[k];
15
       for (auto j : G[i]) val[j] = 0;
16
17
   // 有向图
18
19
   ll n, m; cin >> n >> m;
   vector<pair<ll, ll>>> Edges(m);
20
21
   vector<vector<pll>>> G(n + 2);
22
   vector<ll> deg(n + 2);
   for (auto δ[i, j] : Edges) cin >> i >> j, ++deg[i], ++deg[j];
   for (auto [i, j] : Edges) {
24
25
       Il flg = 0;
26
       if (deg[i] > deg[j] || (deg[i] = deg[j] & i > j)) swap(i, j), flg
27
       G[i].emplace_back(j, flg);
28
   vector<ll> in(n + 2), out(n + 2);
29
30
   ll ans = 0;
   for (ll i = 1; i ≤ n; ++i) {
31
       for (auto [j, w] : G[i]) w ? (++in[j]) : (++out[j]);
32
33
       for (auto [j, w1] : G[i]) for (auto [k, w2] : G[j]) {
           if (w1 = w2) ans += w1 ? in[k] : out[k];
34
35
       for (auto [j, w] : G[i]) in[j] = out[j] = 0;
36
37
   cout \ll ans \ll '\n';
38
```

## 4.3 四元环计数

#### 4.3.1 四元环计数

From zpk

• 无向图: 类似, 由于定向后出度结论过于强大, 可以暴力。讨论了三种情况。

• **有向图**: 缺少题目,但应当类似三元环计数有向形式记录定向边和原边的正反关系。因为此法 最强的结论是定向后出度  $O(\sqrt{m})$ ,实际上方法很暴力,应当不难数有向形式的。

```
ll n, m; cin >> n >> m;
 vector<pair<ll, ll>>> Edges(m);
 3 | vector<vector<ll>>> G(n + 2), iG(n + 2);
 4 | vector<ll> deg(n + 2);
 5 | for (auto \delta[i, j]: Edges) cin \gg i \gg j, \#deg[i], \#deg[j];
 6 | for (auto [i, j] : Edges) {
       if (deg[i] > deg[j] \parallel (deg[i] = deg[j] \& i > j)) swap(i, j);
7
       G[i].emplace back(j), iG[j].emplace back(i);
 8
9
10
   ll\ ans = 0;
   vector < ll > v1(n + 2), v2(n + 2);
11
   for (ll i = 1; i ≤ n; ++i) {
12
       for (auto j : G[i]) for (auto k : G[j]) ++v1[k];
13
       for (auto j : iG[i]) for (auto k : G[j]) ans += v1[k], ++v2[k];
14
       for (auto j : G[i]) for (auto k : G[j]) ans += v1[k] * (v1[k] - 1)
15
           / 2, v1[k] = 0;
       for (auto j : iG[i]) for (auto k : G[j]) {
16
            if (deg[k] > deg[i] \mid | (deg[k] = deg[i] & k > i)) ans += v2[k
17
               ] * (v2[k] - 1) / 2;
           v2[k] = 0;
18
        }
19
20
21
   cout \ll ans \ll '\n';
```

### 4.4 基环树

```
1 #include<bits/stdc++.h>
 2 using namespace std;
 3 typedef long long i64;
 4 const int INF = 1e9;
 5 | const i64 INFL = 1e18;
 6 const int MAXN = 1e5 + 3;
 7 | using edge = tuple<int, int, int>;
   vector <edge> E[MAXN];
 8
9 | vector <edge> W;
10 | vector <int> C;
   edge F[MAXN];
11
   bool V[MAXN];
12
   int I[MAXN], o;
13
14
   void dfs0(int u, int e){
       V[u] = true;
15
16
        I[u] = ++ o;
17
        for(auto \&[i, v, w] : E[u]) if(i \neq e){
            if(V[v]){
18
                if(I[v] < I[u]){
19
20
                    for(int p = u; p \neq v;){
                         auto \delta[j, f, x] = F[p];
21
                         C.push_back(p);
22
23
                        W.push_back({j, p, x});
24
                         p = f;
25
                    }
26
                    C.push_back(v);
                    W.push_back({i, v, w});
27
```

```
28
                }
29
            } else {
                F[v] = \{i, u, w\};
30
31
                dfs0(v, i);
32
            }
        }
33
34
35
   namespace Problem2{
   // 	≡= 删除环上第 i 条边, 求直径 ≡=
36
        i64 H[MAXN], A1[MAXN], B1[MAXN], A2[MAXN], B2[MAXN], A3[MAXN], B3[
37
           MAXN];
        i64 L[MAXN];
38
39
        i64 dis = 0;
        void dfs1(int u, int e){
40
            for(auto \&[i, v, w] : E[u]) if(i \neq e){
41
42
                if(!V[v]){
                    dfs1(v, i);
43
                    dis = max(dis, L[u] + w + L[v]);
44
                    L[u] = max(L[u], L[v] + w);
45
46
            }
47
48
49
       int main(){
            int n;
50
51
            cin >> n;
52
            for(int i = 1; i \leq n; ++ i){
53
                int u, v, w;
                cin >> u >> v >> w;
54
55
                E[u].push_back({i, v, w});
                E[v].push_back({i, u, w});
56
57
            dfs0(1, 0);
58
            memset(V, 0, sizeof(V));
59
60
            for(auto &u : C)
61
                V[u] = true;
            for(auto &u : C){
62
63
                dfs1(u, 0);
64
            int l = 0, r = C.size() - 1;
65
            for(int i = l;i ≤ r; ++ i){
66
67
                int x = C[i];
                if(i > 0)
68
69
                    H[i] = H[i - 1] + get<2>(W[i - 1]);
70
                A1[i] = L[x] + H[i];
                B1[i] = L[x] - H[i];
71
72
                A2[i] = L[x] - H[i];
                B2[i] = L[x] + H[i];
73
74
            i64 h = H[r] + get<2>(W.back());
75
            for(int i = l; i \leq r; ++ i)
76
                A1[i] = max(i = l ? -INFL : A1[i - 1], L[C[i]] + H[i]),
77
                A2[i] = max(i = l ? -INFL : A2[i - 1], L[C[i]] - H[i]);
78
79
            for(int i = r; i \ge l; -- i)
80
                B1[i] = max(i = r ? -INFL : B1[i + 1], L[C[i]] - H[i]),
                B2[i] = max(i = r ? -INFL : B2[i + 1], L[C[i]] + H[i]);
81
            A3[l] = -INFL, B3[r] = -INFL;
82
            for(int i = l + 1; i \leq r; ++ i){}
83
```

```
84
                 int x = C[i];
                 i64 \text{ w} = A2[i - 1] + L[x] + H[i];
 85
                 A3[i] = max(A3[i - 1], w);
 86
 87
             for(int i = r - 1; i \ge l; -- i){
 88
                 int x = C[i];
 89
                 i64 w = B2[i + 1] + L[x] - H[i];
 90
 91
                 B3[i] = max(B3[i + 1], w);
             }
 92
             i64 t = INFL;
 93
 94
             for(int i = l;i < r; ++ i){
                 i64 d = A1[i] + B1[i + 1] + h;
 95
 96
                 i64 g = A2[i] + B2[i + 1] + 0;
 97
                 d = max({d, dis, A3[i], B3[i + 1]});
 98
                 t = min(t, d);
 99
100
             t = min(t, max(A3[r], dis));
101
             if(t \% 2 = 0)
                 cout << t / 2 << ".0" << endl;
102
             if(t \% 2 = 1)
103
                 cout << t / 2 << ".5" << endl;
104
105
             return 0;
         }
106
107
108
    namespace Problem3{
109
    // === 求最大点权独立集 ====
         int A[MAXN];
110
         i64 X[MAXN], Y[MAXN];
111
112
         i64 P[MAXN][2], Q[MAXN][2];
         void dfs1(int u, int e){
113
             for(auto \delta[i, v, w] : E[u]) if(i \neq e){
114
                 if(!V[v]){
115
                      dfs1(v, i);
116
                      Y[u] += max(X[v], Y[v]);
117
118
                      X[u] += Y[v];
119
120
121
             X[u] += A[u];
         }
122
123
         int main(){
124
             int n;
             cin >> n;
125
126
             for(int i = 1; i \leq n; ++ i){
127
                 cin >> A[i];
128
129
             for(int i = 1; i \leq n; ++ i){
                 int u, v;
130
                 cin >> u >> v;
131
132
                 ++ u, ++ v;
133
                 E[u].push_back({i, v, 0});
                 E[v].push_back({i, u, 0});
134
             }
135
136
             double p;
137
             cin >> p;
             dfs0(1, 0);
138
             memset(V, 0, sizeof(V));
139
140
             for(auto &u : C)
```

```
141
                 V[u] = true;
142
             for(auto &u : C){
                 dfs1(u, 0);
143
144
             int l = 0, r = C.size() - 1;
145
146
             P[0][1] = X[C[0]];
             P[0][0] = -INFL;
147
148
             Q[0][0] = Y[C[0]];
             Q[0][1] = -INFL;
149
             for(int i = l + 1; i \leq r; ++ i){}
150
                 int x = C[i];
151
                 P[i][1] = X[x] + P[i - 1][0];
152
153
                 P[i][0] = Y[x] + max(P[i - 1][0], P[i - 1][1]);
154
                 Q[i][1] = X[x] + Q[i - 1][0];
                 Q[i][0] = Y[x] + max(Q[i - 1][0], Q[i - 1][1]);
155
156
157
             i64 ans = \max(\{P[r][0], Q[r][0], Q[r][1]\});
158
             cout << fixed << setprecision(1) << ans * p << endl;</pre>
             return 0;
159
         }
160
161
162
    int main(){
163
         return Problem3 :: main();
164
```

#### 4.5 2-SAT

#### 4.5.1 例题

n 个变量 m 个条件,形如若  $x_i = a$  则  $y_j = b$ ,找到任意一组可行解或者报告无解。

```
#include<bits/stdc++.h>
 2 using namespace std;
 3 | using i64 = long long;
   const int INF = 1e9;
   const i64 INFL = 1e18;
   namespace SCC{
 6
 7
       const int MAXN= 2e6 + 3;
       vector <int> V[MAXN];
 8
9
       stack <int> S;
        int D[MAXN], L[MAXN], C[MAXN], o, s;
10
        bool F[MAXN], I[MAXN];
11
       void add(int u, int v){ V[u].push_back(v); }
12
       void dfs(int u){
13
            L[u] = D[u] = ++ o, S.push(u), I[u] = F[u] = true;
14
15
            for(auto &v : V[u]){
                if(F[v]){
16
                    if(I[v]) L[u] = min(L[u], D[v]);
17
18
                } else {
                    dfs(v), L[u] = min(L[u], L[v]);
19
20
21
            if(L[u] = D[u]){
22
23
                int c = ++ s;
                while(S.top() \neq u){
24
                    int v = S.top(); S.pop();
25
                    I[v] = false;
26
27
                    C[v] = c;
```

```
28
29
                 S.pop(), I[u] = false, C[u] = c;
             }
30
        }
31
32
33
   const int MAXN = 1e6 + 3;
34
   int X[MAXN][2], o;
35
   int main(){
        ios :: sync_with_stdio(false);
36
37
        int n, m;
38
        cin >> n >> m;
39
        for(int i = 1; i \leq n; ++ i)
40
            X[i][0] = ++ o;
41
        for(int i = 1; i \leq n; ++ i)
            X[i][1] = ++ o;
42
43
        for(int i = 1; i \leq m; ++ i){
44
             int a, x, b, y;
45
            cin \gg a \gg x \gg b \gg y;
46
             SCC :: add(X[a][!x], X[b][y]);
47
            SCC :: add(X[b][!y], X[a][x]);
48
49
        for(int i = 1; i \leq 0; ++ i)
             if(!SCC :: F[i])
50
                 SCC :: dfs(i);
51
52
        bool ok = true;
53
        for(int i = 1; i \leq n; ++ i){
             if(SCC :: C[X[i][0]] = SCC :: C[X[i][1]])
54
                 ok = false;
55
56
        if(ok){
57
             cout << "POSSIBLE" << endl;</pre>
58
59
             for(int i = 1; i \leq n; ++ i){
                 int a = SCC :: C[X[i][0]];
60
61
                 int b = SCC :: C[X[i][1]];
                 if(a < b)
62
                     cout << 0 << " ";
63
64
                 else
65
                     cout << 1 << " ";
             }
66
67
            cout << endl;</pre>
68
        } else {
69
            cout << "IMPOSSIBLE" << endl;</pre>
70
71
        return 0;
72
   }
```

## 4.6 割点

```
#include<bits/stdc++.h>
using namespace std;
const int MAXN= 2e4 + 3;
const int MAXM= 1e5 + 3;
vector<int> V[MAXN];
int n, m, o, D[MAXN], L[MAXN];
bool F[MAXN], C[MAXN];
void dfs(int u, int g){
L[u] = D[u] = ++ o, F[u] = true; int s = 0;
```

```
10
        for(auto &v : V[u]){
11
            if(!F[v]){
12
                 dfs(v, g), ++ s;
13
                 L[u] = min(L[u], L[v]);
                 if(u \neq g & L[v] \geqslant D[u]) C[u] = true;
14
15
            } else {
                 L[u] = min(L[u], D[v]);
16
17
            }
18
19
        if(u = g & s > 1) C[u] = true;
20
   int main(){
21
22
        cin >> n >> m;
23
        for(int i = 1; i \leq m; ++ i){
24
            int u, v;
25
            cin >> u >> v;
26
            V[u].push_back(v);
27
            V[v].push_back(u);
28
29
        for(int i = 1; i \leq n; ++ i)
            if(!F[i]) dfs(i, i);
30
31
        vector <int> ANS;
        for(int i = 1; i \leq n; ++ i)
32
33
            if(C[i]) ANS.push_back(i);
34
        cout << ANS.size() << endl;</pre>
35
        for(auto &u : ANS)
36
            cout << u << " ";
37
        return 0;
38
   }
```

## 4.7 边双连通分量

```
1 #include < bits/stdc++.h>
2 using namespace std;
3 using i64 = long long;
   const int INF = 1e9;
4
5 | const i64 INFL = 1e18;
6 const int MAXN= 5e5 + 3;
7
   vector <vector<int>>> A;
   vector <pair<int, int>> V[MAXN];
9
   stack <int> S;
   int D[MAXN], L[MAXN], o;
10
   bool I[MAXN];
11
   void dfs(int u, int l){
12
       D[u] = L[u] = ++ o; I[u] = true, S.push(u); int s = 0;
13
       for(auto &p : V[u]) {
14
15
            int v = p.first, id = p.second;
16
           if(id \neq l){
                if(D[v]){
17
                    if(I[v])
                                L[u] = min(L[u], D[v]);
18
19
                } else {
                    dfs(v, id), L[u] = min(L[u], L[v]), ++ s;
20
21
            }
22
23
       if(D[u] = L[u]){
24
25
           vector <int> T;
```

## 4.8 点双连通分量

```
#include<bits/stdc++.h>
   using namespace std;
 2
 3 using i64 = long long;
 4 const int INF = 1e9;
 5 const i64 INFL = 1e18;
 6 | const int MAXN= 5e5 + 3;
 7
   vector <vector<int>> A;
 8
   vector <int> V[MAXN];
   stack <int> S;
9
   int D[MAXN], L[MAXN], o; bool I[MAXN];
10
   void dfs(int u, int f){
11
       D[u] = L[u] = ++ o; I[u] = true, S.push(u); int s = 0;
12
13
        for(auto &v : V[u]) if(v \neq f){
14
            if(D[v]){
                if(I[v])
                          L[u] = min(L[u], D[v]);
15
            } else {
16
17
                dfs(v, u), L[u] = min(L[u], L[v]), ++ s;
                if(L[v] \geqslant D[u]){
18
19
                    vector <int> T;
                    while(S.top() \neq v){
20
21
                        int t = S.top(); S.pop();
22
                        T.push_back(t), I[t] = false;
23
                    T.push_back(v), S.pop(), I[v] = false;
24
25
                    T.push_back(u);
26
                    A.push_back(T);
                }
27
            }
28
29
        if(f = 0 \& s = 0){
30
31
           A.push_back({u});
32
        }
33
```

### 4.9 强连通分量

```
#include<bits/stdc++.h>
using namespace std;
using i64 = long long;
const int INF = 1e9;
const i64 INFL = 1e18;
const int MAXN= 5e5 + 3;
vector <int> V[MAXN];
stack <int> S;
```

```
int D[MAXN], L[MAXN], C[MAXN], o, s;
10
   bool F[MAXN], I[MAXN];
   void add(int u, int v){ V[u].push_back(v); }
11
12
   void dfs(int u){
        L[u] = D[u] = ++ o, S.push(u), I[u] = F[u] = true;
13
        for(auto &v : V[u]){
14
            if(F[v]){
15
16
                 if(I[v]) L[u] = min(L[u], D[v]);
17
            } else {
18
                 dfs(v), L[u] = min(L[u], L[v]);
19
20
        }
21
        if(L[u] = D[u]){
22
            int c = ++ s;
            while(S.top() \neq u){
23
24
                 int v = S.top(); S.pop();
25
                 I[v] = false;
26
                 C[v] = c;
27
28
            S.pop(), I[u] = false, C[u] = c;
        }
29
30
31
   vector <int> ANS[MAXN];
32
   int main(){
33
        int n, m;
34
        cin \gg n \gg m;
35
        for(int i = 1; i \leq m; ++ i){
36
            int u, v;
37
            cin >> u >> v;
38
            V[u].push_back(v);
39
40
        for(int i = 1; i \leq n; ++ i)
            if(!F[i])
41
42
                 dfs(i);
43
        for(int i = 1; i \leq n; ++ i){
            ANS[C[i]].push_back(i);
44
        }
45
46
        cout << s << endl;</pre>
        for(int i = 1; i \leq n; ++ i) if(F[i]){
47
48
            int c = C[i];
            sort(ANS[c].begin(), ANS[c].end());
49
50
            for(auto &u : ANS[c])
                 cout << u << " ", F[u] = false;</pre>
51
52
            cout << endl;</pre>
        }
53
54
        return 0;
55
```

# 5 网络流

## 5.1 费用流

```
#include<bits/stdc++.h>
using namespace std;
namespace MCMF{
using i64 = long long;
```

```
const i64 INF = 1e18;
 5
 6
        const int MAXN = 1e5 + 3;
       const int MAXM = 2e5 + 3;
7
        int H[MAXN], V[MAXM], N[MAXM], W[MAXM], F[MAXM], o = 1, n;
8
        void add(int u, int v, int f, int c){
9
            V[++ o] = v, N[o] = H[u], H[u] = o, F[o] = f, W[o] = c;
10
            V[++ o] = u, N[o] = H[v], H[v] = o, F[o] = 0, W[o] = -c;
11
12
            n = max(n, u);
            n = max(n, v);
13
14
15
       void clear(){
            for(int i = 1; i \leq n; ++ i)
16
17
                H[i] = 0;
18
            n = 0, o = 1;
19
20
       bool I[MAXN];
21
        i64 D[MAXN];
        bool spfa(int s, int t){
22
23
            queue <int> Q;
24
            Q.push(s), I[s] = true;
25
            for(int i = 1; i \leq n; ++ i)
26
                D[i] = INF;
            D[s] = 0;
27
            while(!Q.empty()){
28
                int u = Q.front(); Q.pop(), I[u] = false;
29
30
                for(int i = H[u];i;i = N[i]){
31
                    const int &v = V[i];
32
                    const int &f = F[i];
33
                    const int &w = W[i];
34
                    if(f & D[u] + w < D[v]){
35
                         D[v] = D[u] + w;
                         if(!I[v]) Q.push(v), I[v] = true;
36
                    }
37
                }
38
            }
39
            return D[t] \neq INF;
40
41
       int C[MAXN]; bool T[MAXN];
42
        pair<i64, i64> dfs(int s, int t, int u, i64 maxf){
43
44
            if(u = t)
                return make_pair(maxf, 0);
45
            i64 totf = 0;
46
47
            i64 \text{ totc} = 0;
            T[u] = true;
48
            for(int &i = C[u];i;i = N[i]){
49
50
                const int &v = V[i];
                const int &f = F[i];
51
                const int &w = W[i];
52
53
                if(f & D[v] = D[u] + w & !T[v]){
                    auto p = dfs(s, t, v, min(1ll * F[i], maxf));
54
55
                    i64 f = p.first;
56
                    i64 c = p.second;
                    F[i
57
                           ] -= f;
                    F[i ^ 1] += f;
58
                    totf += f;
59
60
                    totc += 1ll * f * W[i] + c;
61
                    maxf -= f;
```

```
if(maxf = 0){
62
                         T[u] = false;
63
                         return make_pair(totf, totc);
64
65
                    }
                }
66
            }
67
            T[u] = false;
68
69
            return make_pair(totf, totc);
        }
70
        pair<i64, i64> mcmf(int s, int t){
71
72
            i64 \ ans1 = 0;
            i64 ans2 = 0;
73
            pair<i64, i64> r;
74
            while(spfa(s, t)){
75
                memcpy(C, H, sizeof(H));
76
                r = dfs(s, t, s, INF);
77
                ans1 += r.first;
78
79
                ans2 += r.second;
            }
80
81
            return make_pair(ans1, ans2);
        }
82
83
84
   int gread();
   int main(){
85
        int n = qread(), m = qread(), s = qread(), t = qread();
86
87
        for(int i = 1;i ≤ m;++ i){
            int u = gread(), v = gread(), f = gread(), c = gread();
88
            MCMF :: add(u, v, f, c);
89
90
91
        pair<long long, long long> ans = MCMF :: mcmf(s, t);
92
        printf("%lld %lld\n", ans.first, ans.second);
       return 0;
93
94
```

### 5.2 最小割树

## 5.2.1 用法

给定无向图求出最小割树,点 u 和 v 作为起点终点的最小割为树上 u 到 v 路径上边权的最小 信。

```
#include <bits/stdc++.h>
 1
   using namespace std;
 3 | int gread();
 4
   namespace Dinic{
       const long long INF = 1e18;
 5
       const int SIZ = 1e5 + 3;
 6
 7
       int n, m;
       int H[SIZ], V[SIZ], N[SIZ], F[SIZ], t = 1;
 8
       int add(int u, int v, int f){
 9
            V[++ t] = v, N[t] = H[u], F[t] = f, H[u] = t;
10
           V[++ t] = u, N[t] = H[v], F[t] = 0, H[v] = t;
11
12
            n = max(n, u);
            n = max(n, v);
13
14
            return t - 1;
15
       void clear(){
16
```

```
17
            for(int i = 1; i \leq n; ++ i)
18
                H[i] = 0;
19
            n = m = 0, t = 1;
        }
20
21
        int D[SIZ];
        bool bfs(int s, int t){
22
23
            queue <int> Q;
24
            for(int i = 1; i \leq n; ++ i)
25
                D[i] = 0;
            Q.push(s), D[s] = 1;
26
27
            while(!Q.empty()){
28
                int u = Q.front(); Q.pop();
29
                for(int i = H[u];i;i = N[i]){
                     const int &v = V[i];
30
                     const int &f = F[i];
31
                     if(f \neq 0 & !D[v]){
32
33
                         D[v] = D[u] + 1;
34
                         Q.push(v);
                     }
35
36
                }
37
38
            return D[t] \neq 0;
        }
39
40
        int C[SIZ];
        long long dfs(int s, int t, int u, long long maxf){
41
42
            if(u = t)
43
                return maxf;
            long long totf = 0;
44
45
            for(int &i = C[u];i;i = N[i]){
46
                const int &v = V[i];
47
                const int &f = F[i];
                if(D[v] = D[u] + 1){
48
                    long long resf = dfs(s, t, v, min(maxf, 1ll * f));
49
50
                    totf += resf;
51
                    maxf -= resf;
                           ] -= resf;
52
                     F[i
                     F[i ^ 1] += resf;
53
54
                     if(maxf = 0)
                         return totf;
55
56
                }
            }
57
58
            return totf;
59
       long long dinic(int s, int t){
60
61
            long long ans = 0;
62
            while(bfs(s, t)){
                memcpy(C, H, sizeof(H));
63
64
                ans += dfs(s, t, s, INF);
65
66
            return ans;
        }
67
68
69
   namespace GHTree{
70
        const int MAXN = 500 + 5;
71
        const int MAXM = 1500 + 5;
72
        const int INF = 1e9;
        int n, m, U[MAXM], V[MAXM], W[MAXM], A[MAXM], B[MAXM];
73
```

```
74
         void add(int u, int v, int w){
 75
             ++ m;
             U[m] = u;
 76
 77
             V[m] = v;
             W[m] = w;
 78
 79
             A[m] = Dinic :: add(u, v, w);
             B[m] = Dinic :: add(v, u, w);
 80
 81
             n = max(n, u);
             n = max(n, v);
 82
 83
         vector <pair<int, int> > E[MAXN];
 84
         void build(vector <int> N){
 85
 86
             int s = N.front();
 87
             int t = N.back();
             if(s = t) return;
 88
 89
             for(int i = 1; i \leq m; ++ i){
 90
                 int a = A[i]; Dinic :: F[a] = W[i], Dinic :: F[a ^ 1] = 0;
                 int b = B[i]; Dinic :: F[b] = W[i], Dinic :: F[b ^ 1] = 0;
 91
 92
 93
             int w = Dinic :: dinic(s, t);
 94
             E[s].push_back(make_pair(t, w));
 95
             E[t].push_back(make_pair(s, w));
 96
             vector <int> P;
             vector <int> Q;
 97
             for(auto &u : N){
 98
                 if(Dinic :: D[u] \neq 0)
 99
100
                      P.push_back(u);
                 else
101
102
                      Q.push_back(u);
103
104
             build(P), build(Q);
105
         int D[MAXN];
106
         int cut(int s, int t){
107
108
             queue <int> Q; Q.push(s);
             for(int i = 1; i \leq n; ++ i)
109
110
                 D[i] = -1;
111
             D[s] = INF;
             while(!Q.empty()){
112
113
                 int u = Q.front(); Q.pop();
                 for(auto &e : E[u]){
114
115
                      int v = e.first;
116
                      int w = e.second;
                      if(D[v] = -1){
117
                          D[v] = min(D[u], w);
118
119
                          Q.push(v);
                      }
120
                 }
121
122
             }
123
             return D[t];
         }
124
125
```

### 5.3 最大流

```
#include <bits/stdc++.h>
using namespace std;
```

```
3 using i64 = long long;
   const int INF = 1e9;
 5
   const i64 INFL = 1e18;
   namespace Dinic{
 6
 7
        const i64 INF = 1e18;
        const int SIZ = 5e5 + 3;
 8
9
        int n, m;
10
        int H[SIZ], V[SIZ], N[SIZ], F[SIZ], t = 1;
        void add(int u, int v, int f){
11
            V[++ t] = v, N[t] = H[u], F[t] = f, H[u] = t;
12
13
            V[++ t] = u, N[t] = H[v], F[t] = 0, H[v] = t;
14
            n = max(n, u);
15
            n = max(n, v);
16
        void clear(){
17
18
            for(int i = 1; i \leq n; ++ i)
19
                H[i] = 0;
20
            n = m = 0, t = 1;
21
22
        int D[SIZ];
23
        bool bfs(int s, int t){
24
            queue <int> Q;
            for(int i = 1; i \leq n; ++ i)
25
                D[i] = 0;
26
            Q.push(s), D[s] = 1;
27
28
            while(!Q.empty()){
                int u = Q.front(); Q.pop();
29
30
                for(int i = H[u];i;i = N[i]){
31
                     const int &v = V[i];
32
                     const int &f = F[i];
33
                     if(f \neq 0 & !D[v]){
34
                         D[v] = D[u] + 1;
                         Q.push(v);
35
                     }
36
                }
37
38
39
            return D[t] \neq 0;
40
        int C[SIZ];
41
42
        i64 dfs(int s, int t, int u, i64 maxf){
43
            if(u = t)
44
                return maxf;
45
            i64 \text{ totf} = 0;
            for(int &i = C[u];i;i = N[i]){
46
                const int &v = V[i];
47
48
                const int &f = F[i];
                if(D[v] = D[u] + 1){
49
                     i64 resf = dfs(s, t, v, min(maxf, 1ll * f));
50
51
                    totf += resf;
52
                    maxf -= resf;
                           ] -= resf;
53
                     F[i ^ 1] += resf;
54
55
                    if(maxf = 0)
56
                         return totf;
                }
57
            }
58
59
            return totf;
```

```
60
        i64 dinic(int s, int t){
61
62
            i64 \text{ ans} = 0;
63
            while(bfs(s, t)){
                memcpy(C, H, sizeof(H));
64
65
                ans += dfs(s, t, s, INF);
66
67
            return ans;
        }
68
69
70
   // == TEST ===
   int qread();
71
72
   int main(){
        int n = qread(), m = qread(), s = qread(), t = qread();
73
        for(int i = 1;i ≤ m; ++ i){
74
75
            int u = qread(), v = qread(), f = qread();
76
            Dinic :: add(u, v, f);
77
        printf("%lld\n", Dinic :: dinic(s, t));
78
79
        return 0;
80
```

## 5.4 上下界费用流

#### 5.4.1 用法

- add(u, v, l, r, c): 连一条容量在 [l,r] 的从 u 到 v 的费用为 c 的边;
- solve(): 计算无源汇最小费用可行流;
- solve(s, t): 计算有源汇最小费用最大流。

```
1 #include<bits/stdc++.h>
   using namespace std;
   using i64 = long long;
   const int INF = 1e9;
   const i64 INFL = 1e18;
 5
 6
   namespace MCMF{
 7
       const int MAXN = 1e5 + 3;
       const int MAXM = 2e5 + 3;
 8
 9
       int H[MAXN], V[MAXM], N[MAXM], W[MAXM], F[MAXM], o = 1, n;
        void add0(int u, int v, int f, int c){
10
            V[++ o] = v, N[o] = H[u], H[u] = o, F[o] = f, W[o] = c;
11
           V[++ o] = u, N[o] = H[v], H[v] = o, F[o] = 0, W[o] = -c;
12
            n = max(n, u);
13
            n = max(n, v);
14
15
        }
       bool I[MAXN];
16
        i64 D[MAXN];
17
       bool spfa(int s, int t){
18
19
            queue <int> Q;
            Q.push(s), I[s] = true;
20
            for(int i = 1; i \leq n; ++ i)
21
22
                D[i] = INFL;
23
            D[s] = 0;
            while(!Q.empty()){
24
25
                int u = Q.front(); Q.pop(), I[u] = false;
26
                for(int i = H[u];i;i = N[i]){
```

```
const int &v = V[i];
27
28
                     const int &f = F[i];
                     const int &w = W[i];
29
30
                     if(f & D[u] + w < D[v]){
                         D[v] = D[u] + w;
31
                         if(!I[v]) Q.push(v), I[v] = true;
32
                     }
33
                }
34
35
            }
            return D[t] \neq INFL;
36
37
38
        int C[MAXN]; bool T[MAXN];
39
        pair<i64, i64> dfs(int s, int t, int u, i64 maxf){
            if(u = t)
40
                return make pair(maxf, 0);
41
42
            i64 totf = 0;
43
            i64 \text{ totc} = 0;
44
            T[u] = true;
            for(int &i = C[u];i;i = N[i]){
45
46
                const int &v = V[i];
                const int &f = F[i];
47
48
                const int &w = W[i];
                if(f \& D[v] = D[u] + w \& !T[v]){
49
                     auto p = dfs(s, t, v, min(1ll * F[i], maxf));
50
                     i64 f = p.first;
51
52
                     i64 c = p.second;
                     F[i
53
                           ] -= f;
                     F[i ^ 1] += f;
54
55
                     totf += f;
                     totc += 1ll * f * W[i] + c;
56
57
                     maxf -= f;
                     if(maxf = 0){
58
                         T[u] = false;
59
60
                         return make_pair(totf, totc);
                     }
61
                }
62
            }
63
64
            T[u] = false;
            return make_pair(totf, totc);
65
66
        }
        pair<i64, i64> mcmf(int s, int t){
67
68
            i64 \ ans1 = 0;
69
            i64 \text{ ans } 2 = 0;
            pair<i64, i64> r;
70
            while(spfa(s, t)){
71
72
                memcpy(C, H, sizeof(H));
                r = dfs(s, t, s, INFL);
73
74
                ans1 += r.first;
75
                ans2 += r.second;
            }
76
            return make_pair(ans1, ans2);
77
        }
78
79
        i64 cost0;
80
        int G[MAXN];
        void add(int u, int v, int l, int r, int c){
81
82
            G[v] += l;
            G[u] -= l;
83
```

```
cost0 += 1ll * l * c;
 84
 85
             add0(u, v, r - l, c);
 86
 87
         i64 solve(){
             int s = ++ n;
 88
 89
             int t = ++ n;
             i64 \text{ sum} = 0;
 90
 91
             for(int i = 1; i \leq n - 2; ++ i){
                  if(G[i] < 0)
 92
                      add0(i, t, -G[i], 0);
 93
 94
                  else
                      add0(s, i, G[i], 0), sum += G[i];
 95
 96
 97
             auto res = mcmf(s, t);
 98
             if(res.first \neq sum)
 99
                  return -1;
100
             return res.second + cost0;
101
102
         i64 solve(int s0, int t0){
103
             add0(t0, s0, INF, 0);
             int s = ++ n;
104
105
             int t = ++ n;
106
             i64 \text{ sum} = 0;
             for(int i = 1; i \le n - 2; ++ i){
107
                  if(G[i] < 0)
108
109
                      add0(i, t, -G[i], 0);
110
                  else
                      add0(s, i, G[i], 0), sum += G[i];
111
             }
112
             auto res = mcmf(s, t);
113
             if(res.first \neq sum)
114
115
                  return -1;
116
             return res.second + cost0;
         }
117
118
    // == TEST ===
119
    int qread();
120
121
    int main(){
122
         return 0;
123
```

## 5.5 上下界最大流

#### 5.5.1 用法

```
• add(u, v, l, r, c): 连一条容量在 [l, r] 的从 u 到 v 的边;
```

- solve(): 检查是否存在无源汇可行流;
- solve(s, t): 计算有源汇最大流。

```
#include <bits/stdc++.h>
using namespace std;
int qread();
using i64 = long long;
const int INF = 1e9;
const i64 INFL = 1e18;
namespace MCMF{
```

```
const int MAXN = 1e5 + 3;
 8
9
        const int MAXM = 2e5 + 3;
10
        int H[MAXN], V[MAXM], N[MAXM], F[MAXM], o = 1, n;
        void add0(int u, int v, int f){
11
            V[++ o] = v, N[o] = H[u], H[u] = o, F[o] = f;
12
            V[++ o] = u, N[o] = H[v], H[v] = o, F[o] = 0;
13
            n = max(n, u);
14
15
            n = max(n, v);
16
        i64 D[MAXN]:
17
        bool bfs(int s, int t){
18
19
            queue <int> Q;
20
            for(int i = 1; i \leq n; ++ i)
                D[i] = 0;
21
22
            Q.push(s), D[s] = 1;
23
            while(!Q.empty()){
24
                int u = Q.front(); Q.pop();
                for(int i = H[u];i;i = N[i]){
25
                     const int &v = V[i];
26
27
                     const int &f = F[i];
                     if(f \neq 0 \& !D[v])
28
29
                         D[v] = D[u] + 1;
30
                         Q.push(v);
                     }
31
                }
32
33
34
            return D[t] \neq 0;
        }
35
36
        int C[MAXN];
37
        i64 dfs(int s, int t, int u, i64 maxf){
38
            if(u = t)
                return maxf;
39
            i64 \text{ totf} = 0;
40
            for(int &i = C[u];i;i = N[i]){
41
                const int &v = V[i];
42
                const int &f = F[i];
43
                if(f & D[v] = D[u] + 1){
44
45
                     i64 f = dfs(s, t, v, min(1ll * F[i], maxf));
                           ] -= f;
                     F[i
46
                     F[i ^ 1] += f;
47
                     totf += f;
48
                     maxf -= f;
49
                     if(maxf = 0){
50
                         return totf;
51
                     }
52
53
                }
54
55
            return totf;
56
        i64 mcmf(int s, int t){
57
            i64 \text{ ans} = 0;
58
            while(bfs(s, t)){
59
                memcpy(C, H, sizeof(H));
60
61
                ans += dfs(s, t, s, INFL);
62
63
            return ans;
        }
64
```

```
int G[MAXN];
 65
         void add(int u, int v, int l, int r){
 66
             G[v] += l;
 67
 68
             G[u] -= l;
 69
             add0(u, v, r - l);
 70
         void clear(){
 71
 72
             for(int i = 1; i \le 0; ++ i){
 73
                  N[i] = F[i] = V[i] = 0;
 74
 75
             for(int i = 1; i \leq n; ++ i){
                 H[i] = G[i] = C[i] = 0;
 76
 77
 78
             o = 1, n = 0;
 79
 80
         bool solve(){
             int s = ++ n;
 81
 82
             int t = ++ n;
 83
             i64 \text{ sum} = 0;
 84
             for(int i = 1; i \le n - 2; ++ i){
                  if(G[i] < 0)
 85
 86
                      add0(i, t, -G[i]);
 87
                  else
                      add0(s, i, G[i]), sum += G[i];
 88
 89
 90
             auto res = mcmf(s, t);
 91
             if(res \neq sum)
 92
                  return true;
 93
             return false;
         }
 94
 95
         i64 solve(int s0, int t0){
 96
             add0(t0, s0, INF);
 97
             int s = ++ n;
 98
             int t = ++ n;
 99
             i64 \text{ sum} = 0;
             for(int i = 1; i \le n - 2; ++ i){}
100
101
                  if(G[i] < 0)
102
                      add0(i, t, -G[i]);
                  else
103
104
                      add0(s, i, G[i]), sum += G[i];
105
106
             auto res = mcmf(s, t);
             if(res \neq sum)
107
108
                  return -1;
             return mcmf(s0, t0);
109
         }
110
111
112
    const int MAXN = 1e3 + 3;
113
    const int MAXM = 365 + 3;
114
    int G[MAXN], A[MAXN], B[MAXM];
    int main(){
115
         ios :: sync_with_stdio(false);
116
117
         cin.tie(nullptr);
118
         int n, m, o = 0;
119
         while(cin >> n >> m){
120
             int s = ++ o;
121
             int t = ++ o;
```

```
122
             for(int i = 1; i \leq m; ++ i){
123
                  cin >> G[i];
124
                  A[i] = ++ o;
125
                  MCMF :: add(A[i], t, G[i], INF);
126
              for(int i = 1; i \leq n; ++ i){
127
                  B[i] = ++ o;
128
129
                  int c, d;
                  cin \gg c \gg d;
130
                  MCMF :: add(s, B[i], 0, d);
131
132
                  for(int j = 1; j \leq c; ++ j){
133
                      int t, l, r;
134
                      cin \gg t \gg l \gg r;
135
                      t ++:
136
                      MCMF :: add(B[i], A[t], l, r);
137
138
              }
139
             cout \ll MCMF :: solve(s, t) \ll "\n\n";
140
             MCMF :: clear();
141
         }
142
         return 0;
143
    }
```

# 6 数学

## 6.1 线性代数

## 6.1.1 行列式

```
1 #include<bits/stdc++.h>
 2 using namespace std;
 3 using i64 = long long;
   const int INF = 1e9;
   const i64 INFL = 1e18;
   const int MAXN = 600 + 3;
7
   int MOD;
   struct Mat{
 8
9
        int n, m;
        int W[MAXN][MAXN];
10
11
        Mat(int _n = 0, int _m = 0){
12
            n = n;
            m = _m;
13
14
            for(int i = 1; i \leq n; ++ i)
                for(int j = 1; j \leq m; ++ j)
15
                     W[i][j] = 0;
16
        }
17
18
   };
19
   int mat_det(Mat a){
20
        int ans = 1;
21
        const int &n = a.n;
22
        for(int i = 1; i \leq n; ++ i){
23
            int f = -1;
            for(int j = i; j \leq n; ++ j) if(a.W[j][i] \neq 0){
24
25
                f = j;
26
                break;
            }
27
```

```
if(f = -1){
28
29
                 return 0;
30
            if(f \neq i){
31
                 for(int j = 1; j \leq n; ++ j)
32
33
                     swap(a.W[i][j], a.W[f][j]);
34
                 ans = MOD - ans;
35
            for(int j = i + 1; j \le n; ++ j) if(a.W[j][i]){
36
                 while(a.W[j][i]){
37
38
                     int u = a.W[i][i];
39
                     int v = a.W[j][i];
40
                     if(u > v){
                          for(int k = 1; k \leq n; ++ k)
41
                              swap(a.W[i][k], a.W[j][k]);
42
43
                          ans = MOD - ans;
44
                          swap(u, v);
                     }
45
46
                     int rate = v / u;
47
                     for(int k = 1; k \leq n; ++ k){
                          a.W[j][k] = (a.W[j][k] - 1ll * rate * a.W[i][k] %
48
                             MOD + MOD) % MOD;
                     }
49
                 }
50
            }
51
52
53
        for(int i = 1; i \leq n; ++ i)
            ans = 1ll * ans * a.W[i][i] % MOD;
54
55
        return ans;
56
   int main(){
57
        int n;
58
59
        cin >> n >> MOD;
        Mat A(n, n);
60
61
        for(int i = 1; i \leq n; ++ i)
            for(int j = 1; j \leq n; ++ j)
62
                 cin >> A.W[i][j], A.W[i][j] %= MOD;
63
64
        cout << mat_det(A) << endl;</pre>
65
        return 0;
66
```

#### 6.1.2 高斯消元与求秩(实数)

```
1 #include<bits/stdc++.h>
   using namespace std;
   using i64 = long long;
   const int INF = 1e9;
 4
   const i64 INFL = 1e18;
   const int MAXN = 100 + 3;
   const double EPS = 1e-9;
 7
   struct Mat{
 8
9
       int n, m;
       double W[MAXN][MAXN];
10
       Mat(int _n = 0, int _m = 0){
11
12
           n = _n;
           m = _m;
13
14
            for(int i = 1; i \leq n; ++ i)
```

```
15
                 for(int j = 1; j \leq m; ++ j)
16
                     W[i][j] = 0;
        }
17
18
   };
19
   bool zero(double f){
        return fabs(f) < EPS;</pre>
20
21
22
   int mat_rank(Mat &a){
23
        const int &n = a.n;
24
        const int &m = a.m;
25
        int cnt = 0;
26
        for(int i = 1; i \leq m; ++ i){
27
            int p = cnt + 1;
28
            int f = -1;
            for(int j = p; j \leq n; ++ j){
29
30
                 if(!zero(a.W[j][i])){
31
                     f = j;
32
                     break;
                 }
33
34
            if(f = -1)
35
36
                 continue;
            if(f \neq p){
37
                 for(int j = 1; j \leq m; ++ j)
38
                     swap(a.W[p][j], a.W[f][j]);
39
40
41
            ++ cnt;
            for(int j = p + 1; j \le n; ++ j){
42
                 double rate = a.W[j][i] / a.W[p][i];
43
44
                 for(int k = 1; k \leq m; ++ k){
45
                     a.W[j][k] -= rate * a.W[p][k];
46
            }
47
48
        return cnt;
49
50
   double X[MAXN];
51
52
   int main(){
53
        int n;
54
        cin >> n;
        Mat A(n, n);
55
56
        Mat T(n, n + 1);
        for(int i = 1;i ≤ n;++ i){
57
            for(int j = 1; j \leq n; ++ j)
58
                 cin >> A.W[i][j];
59
60
            for(int j = 1; j \leq n; ++ j)
                 T.W[i][j] = A.W[i][j];
61
62
            cin \gg T.W[i][n + 1];
63
        }
64
        int res1 = mat_rank(A);
65
        int res2 = mat_rank(T);
        if(res1 \neq res2)
66
67
            cout << -1 << endl;
68
        else
69
        if(res2 < n)
70
            cout << 0 << endl;</pre>
71
        else {
```

```
72
            for(int i = n; i \ge 1; -- i){
                X[i] = T.W[i][n + 1] / T.W[i][i];
73
                for(int j = i - 1; j \ge 1; -- j){
74
                     double rate = T.W[j][i] / T.W[i][i];
75
                     T.W[j][ i] -= rate * T.W[i][
76
                                                           i];
                     T.W[j][n + 1] -= rate * T.W[i][n + 1];
77
                }
78
79
            for(int i = 1; i \leq n; ++ i)
80
                cout \ll "x" \ll i \ll "=" \ll fixed \ll setprecision(2) \ll X[i]
81
                     << endl;
82
        }
83
        return 0;
84
```

#### 6.1.3 高斯消元与求秩(整数)

```
1 #include<bits/stdc++.h>
 2 using namespace std;
 3
   using i64 = long long;
 4 const int INF = 1e9;
 5 const i64 INFL = 1e18;
   const int MAXN = 100 + 3;
   const int MOD = 998244353;
 7
 8
   struct Mat{
9
        int n, m;
        int W[MAXN][MAXN];
10
11
       Mat(int _n = 0, int _m = 0){
12
            n = _n;
13
            m = _m;
14
            for(int i = 1; i \leq n; ++ i)
15
                for(int j = 1; j \leq m; ++ j)
                    W[i][j] = 0;
16
        }
17
18
   int power(int a, int b){
19
20
        int r = 1;
       while(b){
21
22
            if(b & 1) r = 1ll * r * a % MOD;
23
            b >>= 1, a = 1ll * a * a % MOD;
24
25
       return r;
26
27
   int inv(int x){
28
       return power(x, MOD - 2);
29
30
   int mat_rank(Mat &a){
       const int &n = a.n;
31
32
        const int &m = a.m;
        int cnt = 0;
33
        for(int i = 1; i \leq m; ++ i){
34
            int p = cnt + 1;
35
36
            int f = -1;
            for(int j = p; j ≤ n; ++ j){
37
38
                if(a.W[j][i] \neq 0){
39
                    f = j;
40
                    break;
```

```
41
                 }
42
            if(f = -1)
43
44
                 continue;
45
            if(f \neq p){
46
                 for(int j = 1; j \leq m; ++ j)
                     swap(a.W[p][j], a.W[f][j]);
47
            }
48
49
            ++ cnt:
            int invp = inv(a.W[p][i]);
50
            for(int j = p + 1; j \leq n; ++ j){
51
                 int rate = 1ll * a.W[j][i] * invp % MOD;
52
53
                 for(int k = 1; k \leq m; ++ k){
                     a.W[j][k] = (a.W[j][k] - 1ll * rate * a.W[p][k] % MOD +
54
                         MOD) % MOD;
55
56
            }
57
58
        return cnt;
59
   int X[MAXN];
60
61
   int main(){
        int n;
62
63
        cin >> n;
        Mat A(n, n);
64
65
        Mat T(n, n + 1);
        for(int i = 1; i \leq n; ++ i){
66
            for(int j = 1; j \leq n; ++ j)
67
                 cin >> A.W[i][j];
68
69
            for(int j = 1; j \leq n; ++ j)
70
                 T.W[i][j] = A.W[i][j];
71
            cin \gg T.W[i][n + 1];
72
        }
73
        int res1 = mat_rank(A);
74
        int res2 = mat_rank(T);
75
        if(res1 \neq res2)
76
            cout << -1 << endl;
77
        else
        if(res2 < n)
78
79
            cout << 0 << endl;</pre>
80
        else {
            for(int i = n; i \ge 1; -- i){
81
                 int invp = inv(T.W[i][i]);
82
                 X[i] = 111 * T.W[i][n + 1] * invp % MOD;
83
                 for(int j = i - 1; j \ge 1; -- j){
84
                     int rate = 1ll * T.W[j][i] * invp % MOD;
85
86
                     T.W[j][
                                 i] = (T.W[j][
                                                 i] - 1ll * rate * T.W[i][
                            i] % MOD + MOD) % MOD;
87
                     T.W[j][n + 1] = (T.W[j][n + 1] - 1ll * rate * T.W[i][n]
                        + 1] % MOD + MOD) % MOD;
88
89
90
            for(int i = 1; i \leq n; ++ i)
                 cout << "x" << i << "=" << X[i] << endl;
91
92
        }
93
        return 0;
94
```

#### 6.1.4 矩阵求逆

```
#include<bits/stdc++.h>
   using namespace std;
   using i64 = long long;
   const int INF = 1e9;
 4
   const i64 INFL = 1e18;
 5
   const int MAXN = 400 + 3;
 7
   const int MOD = 1e9 + 7;
   struct Mat{
 8
9
        int n, m;
10
        int W[MAXN][MAXN];
11
        Mat(int _n = 0, int _m = 0){
12
            n = _n;
            m = _m;
13
            for(int i = 1; i \leq n; ++ i)
14
15
                for(int j = 1; j \leq m; ++ j)
                     W[i][j] = 0;
16
        }
17
18
19
   int power(int a, int b){
20
        int r = 1;
        while(b){
21
22
            if(b & 1) r = 1ll * r * a % MOD;
23
            b >>= 1, a = 1ll * a * a % MOD;
24
25
        return r;
26
27
   int inv(int x){
28
        return power(x, MOD - 2);
29
30
   bool mat inv(Mat &a){
31
        const int &n = a.n;
        Mat b(n, n);
32
        for(int i = 1; i \leq n; ++ i)
33
34
            b.W[i][i] = 1;
35
        for(int i = 1; i \leq n; ++ i){
36
            int f = -1;
            for(int j = i; j \leq n; ++ j) if(a.W[j][i] \neq 0){
37
38
                f = j;
39
                break;
40
            if(f = -1){
41
42
                return false;
43
            if(f \neq i){
44
                for(int j = 1; j \leq n; ++ j)
45
                     swap(a.W[i][j], a.W[f][j]),
46
47
                     swap(b.W[i][j], b.W[f][j]);
48
            int invp = inv(a.W[i][i]);
49
            for(int j = i + 1; j \leq n; ++ j){}
50
                int rate = 1ll * a.W[j][i] * invp % MOD;
51
52
                for(int k = 1; k \leq n; ++ k){
                     a.W[j][k] = (a.W[j][k] - 1ll * rate * a.W[i][k] % MOD +
53
                         MOD) % MOD;
54
                     b.W[j][k] = (b.W[j][k] - 1ll * rate * b.W[i][k] % MOD +
                         MOD) % MOD;
```

```
55
                }
            }
56
57
58
        for(int i = n; i \ge 1; -- i){
59
            int invp = inv(a.W[i][i]);
            for(int j = 1; j \leq n; ++ j){
60
                 a.W[i][j] = 1ll * a.W[i][j] * invp % MOD;
61
62
                 b.W[i][j] = 1ll * b.W[i][j] * invp % MOD;
63
            for(int j = i - 1; j \ge 1; -- j){
64
                 int rate = 1ll * a.W[j][i] % MOD;
65
                 for(int k = 1; k \leq n; ++ k){
66
67
                     a.W[j][k] = (a.W[j][k] - 1ll * rate * a.W[i][k] % MOD +
                          MOD) % MOD;
                     b.W[j][k] = (b.W[j][k] - 1ll * rate * b.W[i][k] % MOD +
68
                          MOD) % MOD;
69
                 }
            }
70
71
72
        for(int i = 1; i \leq n; ++ i)
            for(int j = 1; j \leq n; ++ j)
73
74
                 a.W[i][j] = b.W[i][j];
75
        return true;
76
   int X[MAXN];
77
78
   int main(){
        int n;
79
        cin >> n;
80
81
        Mat A(n, n);
82
        for(int i = 1; i \leq n; ++ i)
83
            for(int j = 1; j \leq n; ++ j)
84
                 cin >> A.W[i][j];
        bool res = mat_inv(A);
85
        if(res = false){}
86
            cout << "No Solution" << endl;</pre>
87
88
        } else {
            for(int i = 1; i \leq n; ++ i)
89
                 for(int j = 1; j \leq n; ++ j)
90
                     cout \ll A.W[i][j] \ll " \n"[j = n];
91
92
93
        return 0;
94
```

## 6.1.5 矩阵树

**LGV** 定理叙述 设 G 是一张有向无环图,边带权,每个点的度数有限。给定起点集合  $A = \{a_1, a_2, \dots, a_n\}$ ,终点集合  $B = \{b_1, b_2, \dots, b_n\}$ 。

- 一段路径  $p:v_0\to^{w_1}v_1\to^{w_2}v_2\to\cdots\to^{w_k}v_k$  的边权被定义为  $\omega(p)=\prod w_i$ 。
- 一对顶点 (a,b) 的权值被定义为  $e(a,b) = \sum_{p:a \to b} \omega(p)$ 。

设矩阵 M 如下:

$$M = \begin{pmatrix} e(a_1, b_1) & e(a_1, b_2) & \cdots & e(a_1, b_n) \\ e(a_2, b_1) & e(a_2, b_2) & \cdots & e(a_2, b_n) \\ \vdots & \vdots & \ddots & \vdots \\ e(a_n, b_1) & e(a_n, b_2) & \cdots & e(a_n, b_n) \end{pmatrix}$$

从 A 到 B 得到一个不相交的路径组  $p=(p_1,p_2,\cdots,p_n)$ ,其中从  $a_i$  到达  $b_{\pi_i}$ , $\pi$  是一个排列。 定义  $\sigma(\pi)$  是  $\pi$  逆序对的数量。

给出 LGV 的叙述如下:

$$\det(M) = \sum_{p:A \to B} (-1)^{\sigma(\pi)} \prod_{i=1}^{n} \omega(p_i)$$

可以将边权视作边的重数, 那么 e(a,b) 就可以视为从 a 到 b 的不同路径方案数。

### 矩阵树定理 对于无向图,

- 定义度数矩阵  $D_{i,j} = [i = j] \deg(i)$ ;
- 定义邻接矩阵  $E_{i,j} = E_{j,i}$  是从 i 到 j 的边数个数;
- 定义拉普拉斯矩阵 L = D E。

对于无向图的矩阵树定理叙述如下:

$$t(G) = \det(L_i) = \frac{1}{n} \lambda_1 \lambda_2 \cdots \lambda_{n-1}$$

其中  $L_i$  是将 L 删去第 i 行和第 i 列得到的子式。

对于有向图,类似于无向图定义入度矩阵、出度矩阵、邻接矩阵  $D^{\rm in},D^{\rm out},E$ ,同时定义拉普拉斯矩阵  $L^{\rm in}=D^{\rm in}-E,L^{\rm out}-E$ 。

$$t^{\text{leaf}}(G, k) = \det(L_k^{\text{in}})$$
$$t^{\text{root}}(G, k) = \det(L_k^{\text{out}})$$

其中  $t^{\text{leaf}}(G, k)$  表示以 k 为根的叶向树, $t^{\text{root}}(G, k)$  表示以 k 为根的根向树。

**BEST 定理** 对于一个有向欧拉图 G, 记点 i 的出度为  $out_i$ , 同时 G 的根向生成树个数为 T。T 可以任意选取根。则 G 的本质不同的欧拉回路个数为:

$$T\prod_{i}(\operatorname{out}_{i}-1)!$$

```
1 #include<bits/stdc++.h>
 2 using namespace std;
 3 using i64 = long long;
 4 const int INF = 1e9;
 5 const i64 INFL = 1e18;
   const int MAXN = 300 + 3;
   const int MOD = 1e9 + 7;
 7
   struct Mat{
 8
 9
       int n, m;
10
       int W[MAXN][MAXN];
       Mat(int n = 0, int m = 0){
11
12
           n = _n;
13
           m = _m;
            for(int i = 1; i \leq n; ++ i)
14
                for(int j = 1; j \leq m; ++ j)
15
```

```
16
                     W[i][j] = 0;
        }
17
   };
18
19
   int mat_det(Mat a){
20
        int ans = 1;
21
        const int &n = a.n;
        for(int i = 1; i \leq n; ++ i){
22
23
            int f = -1;
            for(int j = i; j \leq n; ++ j) if(a.W[j][i] \neq 0){
24
25
                 f = j;
26
                 break;
27
            }
28
            if(f = -1){
29
                return 0;
30
            if(f \neq i){
31
32
                 for(int j = 1; j \leq n; ++ j)
33
                     swap(a.W[i][j], a.W[f][j]);
34
                 ans = MOD - ans;
35
            for(int j = i + 1; j \le n; ++ j) if(a.W[j][i]){
36
37
                 while(a.W[j][i]){
38
                     int u = a.W[i][i];
39
                     int v = a.W[j][i];
                     if(u > v){
40
41
                         for(int k = 1; k \leq n; ++ k)
42
                              swap(a.W[i][k], a.W[j][k]);
43
                         ans = MOD - ans;
44
                         swap(u, v);
                     }
45
46
                     int rate = v / u;
47
                     for(int k = 1; k \leq n; ++ k){
                         a.W[j][k] = (a.W[j][k] - 1ll * rate * a.W[i][k] %
48
                             MOD + MOD) % MOD;
                     }
49
                }
50
            }
51
52
        for(int i = 1; i \leq n; ++ i)
53
54
            ans = 1ll * ans * a.W[i][i] % MOD;
55
        return ans;
56
   int D[MAXN];
57
   int W[MAXN][MAXN];
58
   int main(){
59
        int n, m, t;
60
        cin \gg n \gg m \gg t;
61
62
        for(int i = 1; i \leq m; ++ i){
63
            int u, v, w;
64
            cin >> u >> v >> w;
            if(u \neq v){
65
                 if(t = 0){ // 无向图
66
                     D[u] = (D[u] + w) \% MOD;
67
68
                     D[v] = (D[v] + w) \% MOD;
                     W[u][v] = (W[u][v] + w) % MOD;
69
70
                     W[v][u] = (W[v][u] + w) % MOD;
                            // 叶向树
                 } else {
71
```

```
72
                    D[v] = (D[v] + w) \% MOD;
73
                    W[u][v] = (W[u][v] + w) \% MOD;
                }
74
            }
75
76
77
       Mat A(n - 1, n - 1);
        for(int i = 2; i \leq n; ++ i)
78
79
            for(int j = 2; j ≤ n; ++ j) // 以 1 为根的叶向树
80
                A.W[i - 1][j - 1] = MOD - W[i][j];
        for(int i = 2; i \leq n; ++ i)
81
            A.W[i - 1][i - 1] = (D[i] + A.W[i - 1][i - 1]) \% MOD;
82
        cout << mat_det(A) << endl;</pre>
83
84
        return 0;
85
```

# 6.2 大步小步

#### 6.2.1 用法

给定 a, p 求出 x 使得  $a^x = y \pmod{p}$ , 其中 p 为质数。

```
#include<bits/stdc++.h>
   using namespace std;
 2
   int power(int a, int b, int p){
 3
        int r = 1;
 4
 5
       while(b){
            if(b & 1) r = 1ll * r * a % p;
6
7
            b >>= 1, a = 1ll * a * a % p;
8
 9
       return r;
10
   namespace BSGS {
11
12
        unordered_map <int, int> M;
                                          // a ^{^{\prime}} x = y (mod p)
        int solve(int a, int y, int p){
13
            M.clear();
14
            int B = sqrt(p);
15
            int w1 = y, u1 = power(a, p - 2, p);
16
17
            int w2 = 1, u2 = power(a, B, p);
            for(int i = 0; i < B; ++ i){</pre>
18
19
                M[w1] = i;
                w1 = 111 * w1 * u1 % p;
20
21
22
            for(int i = 0;i 
                if(M.count(w2)){
23
                    return i * B + M[w2];
24
25
26
                w2 = 111 * w2 * u2 % p;
27
28
            return -1;
        }
29
30
   int main(){
31
32
        int p, b, n;
33
        cin \gg p \gg b \gg n;
34
        int ans = BSGS :: solve(b, n, p);
35
        if(ans = -1){
            cout << "no solution\n";</pre>
36
37
        } else {
```

### 6.3 中国剩余定理

#### 6.3.1 定理

对于线性方程:

$$\begin{cases} x \equiv a_1 \pmod{m_1} \\ x \equiv a_2 \pmod{m_2} \\ \dots \\ x \equiv a_n \pmod{m_n} \end{cases}$$

如果  $a_i$  两两互质, 可以得到 x 的解  $x \equiv L \pmod{M}$ , 其中  $M = \prod m_i$ , 而 L 由下式给出:

$$L = \left(\sum a_i m_i \times \left(\left(M/m_i\right)^{-1} \bmod m_i\right)\right) \bmod M$$

```
#include<bits/stdc++.h>
   using namespace std;
   const int MAXN = 100 + 3;
   long long A[MAXN], B[MAXN], M = 1;
   long long exgcd(long long a, long long b, long long &x, long long &y){
        if(a = 0){
 6
            x = 0, y = 1; return b;
 7
        } else {
 8
            long long x0 = 0, y0 = 0;
9
            long long d = exgcd(b % a, a, x0, y0);
10
            x = y0 - (b / a) * x0;
11
12
            y = x0;
13
            return d;
        }
14
15
   int main(){
16
17
        int n;
18
        cin >> n;
        for(int i = 1; i \leq n; ++ i){
19
            cin \gg B[i] \gg A[i];
20
21
            M = M * B[i];
22
23
       long long L = 0;
        for(int i = 1; i \leq n; ++ i){
24
            long long m = M / B[i], b, k;
25
            exgcd(m, B[i], b, k);
26
            L = (L + (_int128)A[i] * m * b) % M;
27
28
29
        L = (L \% M + M) \% M;
30
        cout << L << endl;
        return 0;
31
32
```

### 6.4 狄利克雷前缀和

### 6.4.1 用法

计算:

$$s(i) = \sum_{d|i} f_d$$

```
#include<bits/stdc++.h>
 2 using namespace std;
 3 const int MAXN = 2e7 + 3;
 4 unsigned A[MAXN];
   int p, P[MAXN]; bool V[MAXN];
 5
   void solve(int n){
 6
        for(int i = 2; i \leq n; ++ i){
7
            if(!V[i]){
8
9
                P[++p] = i;
                for(int j = 1; j ≤ n / i; ++ j){ // 前缀和
10
11
                     A[j * i] += A[j];
                 }
12
13
            for(int j = 1; j \le p \& P[j] \le n / i; ++ j){
14
                V[i * P[j]] = true;
15
16
                if(i \% P[j] = 0)
17
                     break;
            }
18
19
        }
20
   unsigned seed;
21
   inline unsigned read(){
22
23
        seed ~ seed << 13;
24
        seed ^- seed >> 17;
25
        seed ^= seed << 5;
26
        return seed;
27
28
   int main(){
29
        int n;
30
        cin >> n >> seed;
        for(int i = 1; i \leq n; ++ i){
31
32
            A[i] = read();
33
        }
34
        solve(n);
35
        unsigned ans = 0;
        for(int i = 1; i \leq n; ++ i){
36
            ans ^{\sim} A[i];
37
38
39
        cout << ans << endl;</pre>
40
        return 0;
41
```

### 6.5 万能欧几里得

### 6.5.1 类欧几里得(万能欧几里得)

From zpk

一种神奇递归,对  $y = \left\lfloor \frac{Ax + B}{C} \right\rfloor$  向右和向上走的每步进行压缩,做到  $O(\log V)$  复杂度。其中  $A \geq C$  就是直接压缩,向右之后必有至少  $\lfloor A/C \rfloor$  步向上。A < C 实际上切换 x,y 轴后,相当于压缩了一个上取整折线,而上取整下取整可以互化,便又可以递归。

代码中从 (0,0) 走到  $(n,\lfloor (An+B)/C \rfloor)$ ,假设了  $A,B,C \geq 0,C \neq 0$ (类欧基本都作此假设),U,R 矩阵是从右往左乘的,对列向量进行优化,和实际操作顺序恰好相反。快速幂的  $\log$  据说可以被递归过程均摊掉,实际上并不会导致变成两个  $\log$ 。

```
Matrix solve(ll n, ll A, ll B, ll C, Matrix R, Matrix U) { // (0, 0)
1
      走到 (n, (An+B)/C)
      if (A \ge C) return solve(n, A \% C, B, C, U.qpow(A / C) * R, U);
2
      ll l = B / C, r = (A * n + B) / C;
3
       if (l = r) return R.qpow(n) * U.qpow(l); // l = r \rightarrow l = r or A
4
          = 0 \text{ or } n = 0.
      ll p = (C * r - B - 1) / A + 1;
5
      return R.qpow(n - p) * U * solve(r - l - 1, C, C - B % C + A - 1, A
6
          , U, R) * U.qpow(l);
7
  }
```

### 6.6 扩展欧几里得

#### 6.6.1 内容

给定 a, b, 求出  $ax + by = \gcd(a, b)$  的一组 x, y。

```
int exgcd(int a, int b, int &x, int &y){
1
2
        if(a = 0){
            x = 0, y = 1; return b;
3
        } else {
4
5
            int x0 = 0, y0 = 0;
            int d = exgcd(b % a, a, x0, y0);
6
            x = y0 - (b / a) * x0;
7
            y = x0;
8
9
            return d;
10
        }
11
```

### 6.7 快速离散对数

# 6.7.1 用法

给定原根 g 以及模数  $\operatorname{mod}$ , T 次询问 x 的离散对数。 复杂度  $\mathcal{O}(\operatorname{mod}^{2/3} + T \log \operatorname{mod})$ 。

```
#include<bits/stdc++.h>
 1
   using namespace std;
   int power(int a, int b, int p){
 3
 4
       int r = 1;
       while(b){
 5
 6
            if(b & 1) r = 1ll * r * a % p;
 7
            b >>= 1, a = 1ll * a * a % p;
 8
 9
       return r;
10
11
   namespace BSGS {
       unordered_map <int, int> M;
12
13
       int B, U, P, g;
```

```
14
        void init(int g, int P0, int B0){
15
            M.clear();
            B = B0;
16
17
            P = P0;
            U = power(power(g, B, P), P - 2, P);
18
19
            int w = 1;
            for(int i = 0; i < B; ++ i){
20
21
                M[w] = i;
                w = 111 * w * g % P;
22
23
24
25
        int solve(int y){
26
            int w = y;
            for(int i = 0;i < P / B;++ i){</pre>
27
                if(M.count(w)){
28
29
                     return i * B + M[w];
30
                }
31
                W = 111 * W * U % P;
            }
32
33
            return -1;
        }
34
35
36
   const int MAXN = 1e5 + 3;
   int H[MAXN], P[MAXN], H0, p, h, g, mod;
37
   bool V[MAXN];
38
39
   int solve(int x){
        if(x \leq h){
40
41
            return H[x];
42
43
        int v = mod / x, r = mod % x;
44
        if(r < x - r){
            return ((H0 + solve(r)) % (mod - 1) - H[v] + mod - 1) % (mod -
45
               1);
        } else {
46
            return (solve(x - r) - H[v + 1] + mod - 1) % (mod - 1);
47
        }
48
49
   int main(){
50
        ios :: sync_with_stdio(false);
51
52
        cin.tie(nullptr);
53
        int T;
54
        cin >> g >> mod;
55
        h = sqrt(mod) + 1;
        BSGS :: init(g, mod, sqrt(1ll * mod * sqrt(mod) / log10(mod)));
56
       H0 = BSGS :: solve(mod - 1);
57
58
       H[1] = 0;
59
        for(int i = 2; i \leq h; ++ i){
60
            if(!V[i]){
                P[++ p] = i;
61
62
                H[i] = BSGS :: solve(i);
63
            for(int j = 1; j \le p \& P[j] \le h / i; ++ j){
64
65
                int &p = P[j];
66
                H[i * p] = (H[i] + H[p]) % (mod - 1);
                V[i * p] = true;
67
68
                if(i \% p = 0)
69
                     break;
```

```
70
             }
        }
71
72
        cin >> T;
        while(T --){
73
             int x, tmp = 0;
74
75
             cin >> x;
             cout \ll solve(x) \ll "\n";
76
77
78
        return 0;
79
```

# 6.8 快速最大公约数

#### 6.8.1 用法

已知小值域 m 以及 n 次询问,  $\mathcal{O}(m)$  预处理,  $\mathcal{O}(1)$  单次查询 x,y 的最大公约数。

```
#include<bits/stdc++.h>
   #define up(l, r, i) for(int i = l, END##i = r; i \leq END##i; ++ i)
 2
   #define dn(r, l, i) for(int i = r, END##i = l;i ≥ END##i;-- i)
   using namespace std;
 4
   typedef long long i64;
   const int INF = 2147483647;
 6
   const int MAXN= 5e3 + 3;
7
   const int MAXT= 1e6 + 3;
   const int MAXM= 1e3 + 3;
9
   int G[MAXM][MAXM];
10
   int T[MAXT][3];
11
12
   int A[MAXN], B[MAXN], o = 1e6, h = 1e3, V[MAXT];
   int tgcd(int a, int b){
13
14
        if(a ≤ h & b ≤ h) return G[a][b];
15
       return a = b ? a : 1;
16
   int qgcd(int a, int b){
17
18
        int ans = 1;
19
        up(0, 2, i){
            if(T[b][i] > h){
20
                if(a % T[b][i] = 0) a \neq T[b][i], ans *= T[b][i];
21
22
            } else {
23
                int d = G[a % T[b][i]][T[b][i]];
                a \not= d, ans *= d;
24
25
26
27
       return ans;
28
29
   const int MOD = 998244353;
30
   int main(){
31
        ios :: sync_with_stdio(false);
       cin.tie(nullptr);
32
        up(1, h, i) G[0][i] = G[i][0] = i;
33
34
        up(1, h, i) up(1, h, j){
            if(i ≥ j) G[i][j] = G[i - j][j];
35
                else G[i][j] = G[i][j - i];
36
37
        up(2, o, i) if(!V[i]){
38
39
           V[i] = i;
            for(int j = 2; i * j \leq o; ++ j)
40
41
                if(!V[i * j]) V[i * j] = i;
```

```
42
        T[1][0] = T[1][1] = T[1][2] = 1;
43
44
        up(2, o, i){
            int p = V[i];
45
46
            int a = T[i / p][0];
            int b = T[i / p][1];
47
            int c = T[i / p][2];
48
49
            int x, y, z;
50
            if(p \ge h){
51
                x = 1, y = i / p, z = p;
52
            } else {
53
                if(c * p \leq h){
54
                    x = a, y = b, z = c * p;
55
                else if(b * p \leq h){
56
57
                    x = a, y = b * p, z = c;
58
                    if(y > z) swap(y, z);
59
                else if(a * p \leq h){
60
                     x = a * p, y = b, z = c;
61
62
                    if(x > y) swap(x, y);
63
                     if(y > z) swap(y, z);
                } else {
64
65
                    x = a * b, y = c, z = p;
                    if(x > y) swap(x, y);
66
67
                    if(y > z) swap(y, z);
                    if(x > z) swap(x, z);
68
                }
69
            }
70
71
            T[i][0] = x;
72
            T[i][1] = y;
73
            T[i][2] = z;
        }
74
75
        int n;
76
        cin >> n;
        up(1, n, i) cin \gg A[i];
77
        up(1, n, i) cin \gg B[i];
78
79
        up(1, n, i){
            int s = 0, u = 1;
80
81
            up(1, n, j){
                int d = qgcd(A[i], B[j]);
82
                u = 111 * u * i % MOD;
83
                s = (s + 111 * d * u) % MOD;
84
85
86
            printf("%d\n", s);
        }
87
88
        return 0;
89
```

### 6.9 原根

### 6.9.1 用法

计算 P 的最小原根。

Prime	r	k	g	Prime	r	k	g
5	1	2	2	3221225473		30	5
17	1	4	3	75161927681	35	31	3
97	3	5	5	77309411329	9	33	7
193	3	6	5	206158430209	3	36	22
257	1	8	3	2061584302081	15	37	7
7681	15	9	17	2748779069441	5	39	3
12289	3	12	11	6597069766657	3	41	5
40961	5	13	3	39582418599937		42	5
65537	1	16	3	79164837199873		43	5
786433	3	18	10	263882790666241		44	7
5767169	11	19	3	1231453023109121		45	3
7340033	7	20	3	1337006139375617		46	3
23068673	11	21	3	3799912185593857	27	47	5
104857601	25	22	3	4222124650659841	15	48	19
167772161	5	25	3	7881299347898369	7	50	6
469762049	7	26	3	31525197391593473		52	3
1004535809	479	21	3	180143985094819841	5	55	6
2013265921	15	27	31	1945555039024054273	27	56	5
2281701377	17	27	3	4179340454199820289	29	57	3

```
#include<bits/stdc++.h>
   using namespace std;
   int power(int a, int b, int p){
3
       int r = 1;
4
       while(b){
5
           if(b & 1) r = 1ll * r * a % p;
6
7
           b >>= 1, a = 1ll * a * a % p;
8
9
       return r;
10
11
   int getphi(int x){
12
       int t = x, r = x;
       for(int i = 2;i ≤ x / i;++ i){
13
           if(t \% i = 0){
14
               r = r / i * (i - 1);
15
               while(t % i = 0)
16
```

```
17
                     t \neq i;
            }
18
19
        if(t \neq 1){
20
            r = r / t * (t - 1);
21
22
23
        return r;
24
25
   vector <int> getprime(int x){
26
        vector <int> p;
27
        int t = x;
28
        for(int i = 2; i \le x / i; ++ i){
29
            if(t \% i = 0){
                 p.push back(i);
30
                while(t % i = 0)
31
32
                     t \neq i;
            }
33
34
        if(t \neq 1)
35
36
            p.push_back(x);
37
        return p;
38
   bool test(int g, int m, int mm, vector<int> &P){
39
40
        for(auto &p: P){
41
            if(power(g, mm / p, m) = 1)
42
                 return false;
43
44
        return true;
45
46
   int get_genshin(int m){
47
        int mm = getphi(m);
        vector <int> P = getprime(mm);
48
        for(int i = 1;;++ i){
49
            if(test(i, m, mm, P))
50
51
                 return i;
        }
52
53
54
   int main(){
55
        cout << get_genshin(998244353) << endl;</pre>
56
        return 0;
57
```

## 6.10 快速乘法逆元(离线)

### 6.10.1 用法

离线计算  $x = [x_1, x_2, \cdots, x_n]$  在模 p 意义下的逆元。

```
#include<bits/stdc++.h>
2
   using namespace std;
   int power(int a, int b, int p){
3
4
       int r = 1;
5
       while(b){
6
            if(b & 1) r = 1ll * r * a % p;
            b >>= 1, a = 1ll * a * a % p;
7
8
9
       return r;
10 | }
```

```
const int MAXN = 5e6 + 3;
11
12
   int A[MAXN], B[MAXN];
   int P[MAXN], Q[MAXN];
13
14
   int main(){
        ios :: sync_with_stdio(false);
15
16
        cin.tie(nullptr);
        int n, p, K, S = 1;
17
18
        cin >> n >> p >> K;
        P[0] = 1;
19
        for(int i = 1; i \leq n; ++ i){
20
21
            cin \gg A[i];
22
            P[i] = 111 * P[i - 1] * A[i] % p;
23
24
        Q[n] = power(P[n], p - 2, p);
        for(int i = n; i \ge 1; -- i){
25
26
            Q[i - 1] = 1ll * Q[i] * A[i] % p;
27
            B[i] = 111 * Q[i] * P[i - 1] % p;
        }
28
29
        int ans = 0;
30
        for(int i = 1; i \leq n; ++ i){
            S = 111 * S * K % p;
31
32
            ans = (ans + 1ll * S * B[i]) % p;
        }
33
34
        cout \ll ans \ll "\n";
35
        return 0;
36
```

# 6.11 快速乘法逆元(在线)

# 6.11.1 用法

在线计算  $x = [x_1, x_2, \cdots, x_n]$  在模 p 意义下的逆元。

```
1 #include < bits/stdc++.h>
 2 | using namespace std;
   const int MAXN = 1e7 + 3;
 3
   pair<int, int> F[MAXN], G[MAXN];
 4
   int I[MAXN];
   using u32 = uint32_t;
 6
7
   u32 read(u32 &seed){
        seed '= seed << 13;
8
        seed ^{\sim} seed >> 17;
9
10
        seed ^- seed << 5;
11
        return seed;
12
13
   int main(){
14
        ios :: sync_with_stdio(false);
15
        cin.tie(nullptr);
16
        u32 seed;
        int n, p;
17
        cin >> n >> p >> seed;
18
19
        int m = pow(p, 1.0 / 3.0);
20
        I[1] = 1;
21
        for(int i = 2; i \leq p / m; ++ i){
22
            I[i] = 1ll * (p / i) * (p - I[p % i]) % p;
23
        for(int i = 1; i < m; ++ i){</pre>
24
25
            for(int j = i + 1; j \leq m; ++ j){}
```

```
if(!F[i * m * m / j].second){
26
                   F[i * m * m / j] = \{ i, j \};
27
                   G[i * m * m / j] = \{ i, j \};
28
29
           }
30
       }
31
       F[
             0] = G[ 0] = \{ 0, 1 \};
32
       F[m * m] = G[m * m] = \{ 1, 1 \};
33
       34
           F[i] = F[i - 1];
35
       for(int i = m * m - 1; i \ge 1; -- i) if(!G[i].second)
36
           G[i] = G[i + 1];
37
38
       int lastans = 0;
       for(int i = 1; i \leq n; ++ i){
39
           int a, inv;
40
           a = (read(seed) ^ lastans) % (p - 1) + 1;
41
42
           int w = 1ll * a * m * m / p;
           auto δyy1 = F[w].second; // *avoid y1 in <cmath>
43
           if(1ll * a * yy1 % p \leq p / m){
44
45
               inv = 1ll * I[1ll * a * yy1 % p] * yy1 % p;
           } else {
46
47
               auto &yy2 = G[w].second;
               inv = 1ll * I[1ll * a * (p - yy2) % p] * (p - yy2) % p;
48
49
50
           lastans = inv;
51
       cout << lastans << "\n";</pre>
52
       return 0;
53
54
```

### 6.12 拉格朗日插值

### 6.12.1 定理

给定 n 个横坐标不同的点  $(x_i, y_i)$ ,可以唯一确定一个 n-1 阶多项式如下:

$$f(x) = \sum_{i=1}^{n} \frac{\prod_{j \neq i} (x - x_j)}{\prod_{j \neq i} (x_i - x_j)} \cdot y_i$$

下面代码先求出了多项式再计算 f(k), 也可以直接带入计算。

```
#include<bits/stdc++.h>
 2 using namespace std;
 3 const int MAXN = 2e3 + 3;
   const int MOD = 998244353;
   int X[MAXN], Y[MAXN], F[MAXN], G[MAXN], H[MAXN], A[MAXN];
   int power(int a, int b){
 7
       int r = 1;
 8
       while(b){
            if(b & 1) r = 1ll * r * a % MOD;
9
10
            b >>= 1, a = 1ll * a * a % MOD;
11
12
       return r;
13
14 | int main(){
15
       int n, k;
16
       cin >> n >> k;
17
       for(int i = 1; i \leq n; ++ i){
```

```
cin \gg X[i] \gg Y[i];
18
        }
19
        F[0] = 1;
20
        for(int i = 1;i ≤ n;++ i){ // 计算 prod(x - x_i)
21
            for(int j = 0; j ≤ n; ++ j){
22
                G[j] = ((j = 0 ? 0 : F[j - 1]) - 1 | x F[j] * X[i] % MOD +
23
                     MOD) % MOD;
24
25
            for(int j = 0; j \leq n; ++ j){
                F[j] = G[j];
26
27
        }
28
        for(int i = 1; i \leq n; ++ i){
29
            for(int j = 0; j \leq n; ++ j){
30
                G[j] = F[j];
31
32
            for(int j = n;j \geq 0;-- j){ // 计算 prod(x - x_j) / (x - x_i)
33
                H[j] = G[j + 1];
34
                G[j] = (G[j] + 1ll * H[j] * X[i]) % MOD;
35
36
            int w = 1;
                                          // 计算 inv(prod(x_i - x_j))
37
38
            for(int j = 1; j \leq n; ++ j) if(j \neq i)
                w = 111 * w * (X[i] - X[j] + MOD) % MOD;
39
            w = 1ll * power(w, MOD - 2) * Y[i] % MOD;
40
            for(int j = 0; j \leq n; ++ j)
41
                A[j] = (A[j] + 1ll * w * H[j]) % MOD;
42
        }
43
        int t = 1, ans = 0;
44
45
        for(int i = 0; i \le n - 1; ++ i){
46
            ans = (ans + 1ll * A[i] * t) % MOD;
47
            t = 1ll * t * k % MOD;
48
49
        cout << ans << endl;</pre>
50
```

### 6.13 min-max 容斥

### 6.13.1 定理

$$\max_{i \in S} \{x_i\} = \sum_{T \subseteq S} (-1)^{|T|-1} \min_{j \in T} \{x_j\}$$

$$\min_{i \in S} \{x_i\} = \sum_{T \subseteq S} (-1)^{|T|-1} \max_{j \in T} \{x_j\}$$

期望意义下上式依然成立。

另外设  $\max^k$  表示第 k 大的元素,可以推广为如下式子:

$$\max_{i \in S}^{k} \{x_i\} = \sum_{T \subseteq S} (-1)^{|T|-k} \binom{|T-1|}{k-1} \min_{j \in T} \{x_j\}$$

此外在数论上可以得到:

$$\lim_{i \in S} \{x_i\} = \prod_{T \subseteq S} \left( \gcd_{j \in T} \{x_j\} \right)^{(-1)^{|T|-1}}$$

#### 6.13.2 应用

对于计算 "n 个属性都出现的期望时间"问题,设第 i 个属性第一次出现的时间是  $t_i$ ,所求即为  $\max(t_i)$ ,使用  $\min$ -max 容斥转为计算  $\min(t_i)$ 。

比如 n 个独立物品,每次抽中物品 i 的概率是  $p_i$ ,问期望抽多少次抽中所有物品。那么就可以 计算  $\min_S$  表示第一次抽中物品集合 S 内物品的时间,可以得到:

$$\max_{U} = \sum_{S \subseteq U} (-1)^{|S|-1} \min_{S} = \sum_{S \subseteq U} (-1)^{|S|-1} \cdot \frac{1}{\sum_{x \in S} p_x}$$

### 6.14 Barrett 取模

### 6.14.1 用法

调用 init 计算出 S 和 X, 得到计算  $x \mod P = (x \times X)/2^{60} + S$ 。

```
#include<bits/stdc++.h>
   using namespace std;
   long long S = 0, X = 0;
 4
   void init(int MOD){
       while((1 << (S + 1)) < MOD)
 5
 6
 7
       X = ((_int128)1 \ll 60 + S) / MOD + !!(((_int128)1 \ll 60 + S) %
           MOD):
        cerr << S << " " << X << endl;
 8
 9
10
   int power(long long x, int y, int MOD){
11
        long long r = 1;
       while(y){
12
            if(y & 1){
13
                r = r * x;
14
                r = r - MOD * ((__int128)r * X >> 60 + S);
15
16
17
            x = x * x;
            x = x - MOD * ((_int128)x * X >> 60 + S);
18
19
            y >>= 1;
        }
20
21
        return r;
22
   int main(){
23
24
        init(998244353);
        cout << power(2, 10, 998244353) << endl;</pre>
25
        cout << power(2, 20, 998244353) << endl;</pre>
26
27
        cout << power(2, 30, 998244353) << endl;
28
        cout << power(2, 40, 998244353) << endl;</pre>
29
        return 0;
30
```

### 6.15 Pollard's Rho

#### 6.15.1 用法

- 调用 test(n) 判断 *n* 是否是质数;
- 调用 rho(n) 计算 n 分解质因数后的结果,不保证结果有序。

```
#include<bits/stdc++.h>
 1
   using namespace std;
   using i64 = long long;
   const int INF = 1e9;
   const i64 INFL = 1e18;
   i64 step(i64 a, i64 c, i64 m){
 6
 7
        return ((__int128)a * a + c) % m;
 8
9
   i64 multi(i64 a, i64 b, i64 m){
        return (__int128) a * b % m;
10
11
12
   i64 power(i64 a, i64 b, i64 m){
13
        i64 r = 1;
14
        while(b){
15
            if(b & 1) r = multi(r, a, m);
            b >>= 1, a = multi(a, a, m);
16
17
18
        return r;
19
20
   mt19937_64 MT;
21
   bool test(i64 n){
22
        if(n < 3 || n % 2 = 0)
23
            return n = 2;
        i64 u = n - 1, t = 0;
24
25
        while(u % 2 = 0)
26
            u \not= 2,
27
            t += 1;
28
        int test_time = 20;
        for(int i = 1; i ≤ test_time; ++ i){
29
30
            i64 a = MT() % (n - 2) + 2;
31
            i64 v = power(a, u, n);
            if(v = 1){
32
33
                continue;
            }
34
35
            int s;
            for(s = 0; s < t; ++ s){
36
37
                if(v = n - 1)
38
                     break;
                v = multi(v, v, n);
39
40
            if(s = t)
41
42
                return false;
43
44
        return true;
45
46
   basic_string<i64> rho(i64 n){
47
        if(n = 1)
48
            return {};
        if(test(n)){
49
50
            return {n};
        }
51
        i64 a = MT() \% (n - 1) + 1;
52
        i64 \times 1 = MT() \% (n - 1);
53
54
        i64 x2 = x1;
55
        for(int i = 1;;i <<= 1){
56
            i64 \text{ tot} = 1;
            for(int j = 1; j \leq i; ++ j){
57
```

```
x2 = step(x2, a, n);
58
                 tot = multi(tot, llabs(x1 - x2), n);
59
                 if(j \% 127 = 0){
60
                     i64 d = \_gcd(tot, n);
61
                     if(d > 1)
62
                         return rho(d) + rho(n / d);
63
                 }
64
65
            i64 d = \_gcd(tot, n);
66
            if(d > 1)
67
                return rho(d) + rho(n / d);
68
            x1 = x2;
69
        }
70
71
   // == TEST ===
72
73
   int main(){
74
        int T;
        cin >> T;
75
        for(int _ = 1; _ \leq T; ++ _){
76
            i64 n, p = 0;
77
            cin >> n;
78
79
            auto res = rho(n);
            for(auto &u : res)
80
                 p = max(p, u);
81
            if(res.size() = 1)
82
                 cout << "Prime" << endl;</pre>
83
84
            else
                 cout << p << endl;</pre>
85
86
87
        return 0;
88
```

# 6.16 polya 定理

### 6.16.1 Burnside 引理

记所有染色方案的集合为 X,其中单个染色方案为 x。一种对称操作  $g \in X$  作用于染色方案  $x \in X$  上可以得到另外一种染色 x'。

将所有对称操作作为集合 G, 那么  $Gx = \{gx \mid g \in G\}$  是与 x 本质相同的染色方案的集合,形式化地称为 x 的轨道。统计本质不同染色方案数,就是统计不同轨道个数。

Burnside 引理说明如下:

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|$$

其中  $X^g$  表示在  $g \in G$  的作用下,**不动点**的集合。不动点被定义为 x = gx 的 x。

### 6.16.2 Polya 定理

对于通常的染色问题, X 可以看作一个长度为 n 的序列, 每个元素是 1 到 m 的整数。可以将 n 看作面数、m 看作颜色数。Polya 定理叙述如下:

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} \sum_{g \in G} m^{c(g)}$$

其中 c(g) 表示对一个序列做轮换操作 g 可以分解成多少个置换环。

然而,增加了限制(比如要求某种颜色必须要多少个),就无法直接应用 Polya 定理,需要利用 Burnside 引理进行具体问题具体分析。

#### 6.16.3 应用

给定 n 个点 n 条边的环,现在有 n 种颜色,给每个顶点染色,询问有多少种本质不同的染色方案。

显然 X 是全体元素在 1 到 n 之间长度为 n 的序列,G 是所有可能的单次旋转方案,共有 n 种,第 i 种方案会把 1 置换到 i。于是:

$$\operatorname{ans} = \frac{1}{|G|} \sum_{i=1}^{n} m^{c(g_i)}$$

$$= \frac{1}{n} \sum_{i=1}^{n} n^{\gcd(i,n)}$$

$$= \frac{1}{n} \sum_{d|n}^{n} n^d \sum_{i=1}^{n} [\gcd(i,n) = d]$$

$$= \frac{1}{n} \sum_{d|n}^{n} n^d \varphi(n/d)$$

```
#include<bits/stdc++.h>
 2
   using namespace std;
   const int MOD = 1e9 + 7;
   int power(int a, int b){
 4
        int r = 1;
 5
 6
       while(b){
            if(b & 1) r = 1ll * r * a % MOD;
 7
            b >>= 1, a = 1ll * a * a % MOD;
 8
        }
 9
10
       return r;
11
   vector <tuple<int, int> > P;
12
13
   void solve(int step, int n, int d, int f, int &ans){
        if(step = P.size()){}
14
            ans = (ans + 1ll * power(n, n / d) * f) % MOD;
15
        } else {
16
            auto [w, c] = P[step];
17
            int dd = 1, ff = 1;
18
            for(int i = 0; i \leq c; ++ i){
19
                solve(step + 1, n, d * dd, f * ff, ans);
20
                ff = ff * (w - (i = 0));
21
22
                dd = dd * w;
            }
23
        }
24
25
26
   int main(){
27
        int T;
        cin >> T;
28
       while(T --){
29
            int n, t;
30
31
            cin >> n;
32
            t = n;
33
            for(int i = 2; i * i \le n; ++ i) if(n % i = 0){
                int w = i, c = 0;
34
                while(t % i = 0){
35
                    t \neq i, c +;
36
```

```
37
                 P.push_back({ w, c });
38
39
            if(t \neq 1){
40
                 P.push_back({ t, 1 });
41
42
            int ans = 0;
43
44
            solve(0, n, 1, 1, ans);
            ans = 111 * ans * power(n, MOD - 2) % MOD;
45
46
             cout << ans << endl;</pre>
47
            P.clear();
        }
48
49
        return 0;
50
```

### 6.17 min25 筛

设有一个积性函数 f(n), 满足  $f(p^k)$  可以快速求,考虑搞一个在质数位置和 f(n) 相等的 g(n),满足它有完全积性,并且单点和前缀和都可以快速求,然后通过第一部分筛出 g 在质数位置的前缀和,从而相当于得到 f 在质数位置的前缀和,然后利用它,做第二部分,求出 f 的前缀和。

```
第一部分: G_k(n) = \sum_{i=1}^n [\text{mindiv}(i) > p_k \text{ or isprime}(i)] g(i) (p_0 = 1), 则有 G_k(n) = G_{k-1}(n) - g(p_k)(G_{k-1}(n/p_k) - G_{k-1}(p_{k-1})),复杂度 O(n^{3/4}/\log n)。
第二部分: F_k(n) = \sum_{i=1}^n [\text{mindiv}(i) \ge p_k] f(i),F_k(n) = \sum_{\substack{h \ge k \\ p_h^2 \le n}} \sum_{\substack{c \ge 1 \\ p_h^{c+1} \le n}} (f(p_h^c) F_{h+1}(n/p_h^c) + \sum_{\substack{k \le k \\ p_h^2 \le n}} (f(p_k^c) F_{k+1}(n/p_k^c) + \sum_{\substack{k \le k \\ p_h^2 \le n}} (f(p_k^c) F_{k+1}(n/p_k^c) + \sum_{\substack{k \le k \\ p_h^2 \le n}} (f(p_k^c) F_{k+1}(n/p_k^c) + \sum_{\substack{k \le k \\ p_h^2 \le n}} (f(p_k^c) F_{k+1}(n/p_k^c) + \sum_{\substack{k \le k \\ p_h^2 \le n}} (f(p_k^c) F_{k+1}(n/p_k^c) + \sum_{\substack{k \le k \\ p_h^2 \le n}} (f(p_k^c) F_{k+1}(n/p_k^c) + \sum_{\substack{k \le k \\ p_h^2 \le n}} (f(p_k^c) F_{k+1}(n/p_k^c) + \sum_{\substack{k \le k \\ p_h^2 \le n}} (f(p_k^c) F_{k+1}(n/p_k^c) + \sum_{\substack{k \le k \\ p_h^2 \le n}} (f(p_k^c) F_{k+1}(n/p_k^c) + \sum_{\substack{k \le k \\ p_h^2 \le n}} (f(p_k^c) F_{k+1}(n/p_k^c) + \sum_{\substack{k \le k \\ p_h^2 \le n}} (f(p_k^c) F_{k+1}(n/p_k^c) + \sum_{\substack{k \le k \\ p_h^2 \le n}} (f(p_k^c) F_{k+1}(n/p_k^c) + \sum_{\substack{k \le k \\ p_h^2 \le n}} (f(p_k^c) F_{k+1}(n/p_k^c) + \sum_{\substack{k \le k \\ p_h^2 \le n}} (f(p_k^c) F_{k+1}(n/p_k^c) + \sum_{\substack{k \le k \\ p_h^2 \le n}} (f(p_k^c) F_{k+1}(n/p_k^c) + \sum_{\substack{k \le k \\ p_h^2 \le n}} (f(p_k^c) F_{k+1}(n/p_k^c) + \sum_{\substack{k \le k \\ p_h^2 \le n}} (f(p_k^c) F_{k+1}(n/p_k^c) + \sum_{\substack{k \le k \\ p_h^2 \le n}} (f(p_k^c) F_{k+1}(n/p_k^c) + \sum_{\substack{k \le k \\ p_h^2 \le n}} (f(p_k^c) F_{k+1}(n/p_k^c) + \sum_{\substack{k \le k \\ p_h^2 \le n}} (f(p_k^c) F_{k+1}(n/p_k^c) + \sum_{\substack{k \le k \\ p_h^2 \le n}} (f(p_k^c) F_{k+1}(n/p_k^c) + \sum_{\substack{k \le k \\ p_h^2 \le n}} (f(p_k^c) F_{k+1}(n/p_k^c) + \sum_{\substack{k \le k \\ p_h^2 \le n}} (f(p_k^c) F_{k+1}(n/p_k^c) + \sum_{\substack{k \le k \\ p_h^2 \le n}} (f(p_k^c) F_{k+1}(n/p_k^c) + \sum_{\substack{k \le k \\ p_h^2 \le n}} (f(p_k^c) F_{k+1}(n/p_k^c) + \sum_{\substack{k \le k \\ p_h^2 \le n}} (f(p_k^c) F_{k+1}(n/p_k^c) + \sum_{\substack{k \le k \\ p_h^2 \le n}} (f(p_k^c) F_{k+1}(n/p_k^c) + \sum_{\substack{k \le k \\ p_h^2 \le n}} (f(p_k^c) F_{k+1}(n/p_k^c) + \sum_{\substack{k \le k \\ p_h^2 \le n}} (f(p_k^c) F_{k+1}(n/p_k^c) + \sum_{\substack{k \le k \\ p_h^2 \le n}} (f(p_k^c) F_{k+1}(n/p_k^c) + \sum_{\substack{k \le k \\ p_h^2 \le n}} (f(p_k^c) F_{k+1}(n/p_k^c) + \sum_{\substack{k \le k \\ p_h^2 \le n}} (f(p_k^c) F_{k+1}(n/p_k^
```

 $f(p_h^{c+1})) + F_{\text{prime}}(n) - F_{\text{prime}}(p_{k-1})$ ,在  $n \leq 10^{13}$  可以证明复杂度  $O(n^{3/4}/\log n)$ 。 常见细节问题:

- 由于 n 通常是  $10^{10}$  到  $10^{11}$  的数,导致 n 会爆 int,  $n^2$  会爆 long long,而且往往会用自然数幂和,更容易爆,所以要小心。
- 记  $s = \lfloor \sqrt{n} \rfloor$ ,由于 F 递归时会去找  $F_{h+1}$ ,会访问到 s 以内最大的质数往后的一个质数,而已经证明对于所有  $n \in \mathbb{N}^+$ ,[n+1,2n] 中有至少一个质数,所以只需要筛到 2s 即可。
- 注意补回 f(1)。

```
// 预处理, $1$ 所在的块也算进去了
1
2
   namespace init {
       ll init_n, sqrt_n;
3
4
       vector<ll> np, p, id1, id2, val;
5
       ll cnt:
6
       void main(ll n) {
           init_n = n, sqrt_n = sqrt(n);
7
           ll M = sqrt_n * 2; // 筛出一个 > floor(sqrt(n)) 的质数, 避免后
8
              续讨论边界
           np.resize(M + 1), p.resize(M + 1);
9
           for (ll i = 2; i ≤ M; ++i) {
10
               if (!np[i]) p[++p[0]] = i;
11
12
               for (ll j = 1; j \leq p[0]; ++j) {
                   if (i * p[j] > M) break;
13
                   np[i * p[j]] = 1;
14
                   if (i \% p[j] = 0) break;
15
16
           }
17
           p[0] = 1;
18
           id1.resize(sqrt_n + 1), id2.resize(sqrt_n + 1);
19
20
           val.resize(1);
21
           for (ll l = 1, r, v; l \le n; l = r + 1) {
```

```
v = n / l, r = n / v;
22
                if (v \leq sqrt_n) id1[v] = ++cnt;
23
                else id2[init_n / v] = ++cnt;
24
                val.emplace back(v);
25
            }
26
27
       ll id(ll n) {
28
29
            if (n ≤ sqrt_n) return id1[n];
            else return id2[init n / n];
30
        }
31
32
   using namespace init;
33
34
   // 计算 $G_k$, 两个参数分别是 $g$ 从 $2$ 开始的前缀和和 $g$
   auto calcG = [\delta] (auto\delta\delta sum, auto\delta\delta g) \rightarrow vector<ll> {
       vector<ll> G(cnt + 1);
36
       for (int i = 1; i \leq cnt; ++i) G[i] = sum(val[i]);
37
38
       ll pre = 0;
        for (int i = 1; p[i] * p[i] \le n; ++i) {
39
            for (int j = 1; j \leq cnt; ++j) {
40
41
                if (p[i] * p[i] > val[j]) break;
                ll tmp = id(val[j] / p[i]);
42
43
                G[j] = (G[j] - g(p[i]) * (G[tmp] - pre)) % MD;
            }
44
45
            pre = (pre + g(p[i])) % MD;
46
47
        for (int i = 1; i \leq cnt; ++i) G[i] = (G[i] % MD + MD) % MD;
48
       return G;
49
   };
   // 计算 $F_k$, 直接搜, 不用记忆化。`fp` 是 $F_{\text{prime}}$, `pc` 是
50
       $p^c$, 其中 `f(p[h] ^ c)` 要替换掉。
   function<ll(ll, int)> calcF = [8] (ll m, int k) {
51
       if (p[k] > m) return 0;
52
       ll ans = (fp[id(m)] - fp[id(p[k - 1])]) % MD;
53
        for (int h = k; p[h] * p[h] \le m; ++h) {
54
            ll pc = p[h], c = 1;
55
            while (pc * p[h] \leq m) {
56
                ans = (ans + calcF(m / pc, h + 1) * f(p[h] ^ c)) % MD;
57
58
                ++c, pc = pc * p[h], ans = (ans + f(p[h] ^ c)) % MD;
            }
59
60
61
       return ans;
62
   };
```

#### 6.18 杜教筛

#### 6.18.1 用法

对于积性函数 f,找到易求前缀和的积性函数 g,h 使得 h=f\*g,根据递推式计算  $S(n)=\sum_{i=1}^n f(i)$ :

$$S(n) = H(n) - \sum_{d=1}^{n} g(d) \times S(\lfloor \frac{n}{d} \rfloor)$$

#### 6.18.2 例题

• 对于  $f = \varphi$ , 寻找 g = 1, h = id;

• 对于  $f = \mu$ , 寻找  $g = 1, h = \varepsilon$ 。

```
#include<bits/stdc++.h>
 1
   using namespace std;
   const int MAXN = 1e7 + 3;
   const int H = 1e7;
 5 int P[MAXN], p; bool V[MAXN];
   long long ph[MAXN], sph[MAXN];
 7
   long long mu[MAXN], smu[MAXN];
   long long tp[MAXN];
 8
   long long solve_ph(long long N){
9
        for(int d = N / H; d \ge 1; -- d){
10
11
            long long n = N / d;
            long long wh = 111 * n * (n + 1) / 2;
12
            tp[d] = wh;
13
            for(long long l = 2, r; l \leq n; l = r + 1){
14
15
                r = n / (n / l);
                long long wg = r - l + 1;
16
                long long ws = n / l \le H ? sph[n / l] : tp[N / (n / l)];
17
18
                tp[d] -= wg * ws;
19
20
21
        return N \leq H? sph[N] : tp[1];
22
23
   long long solve_mu(long long N){
        for(int d = N / H; d \geqslant 1; -- d){
24
25
            long long n = N / d;
26
            long long wh = 1;
            tp[d] = wh;
27
            for(long long l = 2, r; l \leq n; l = r + 1){
28
29
                r = n / (n / l);
30
                long long wg = r - l + 1;
31
                long long ws = n / l \le H ? smu[n / l] : tp[N / (n / l)];
32
                tp[d] -= wg * ws;
            }
33
34
35
        return N \leq H? smu[N] : tp[1];
36
   int main(){
37
38
        ios :: sync_with_stdio(false);
39
        cin.tie(nullptr);
        ph[1] = 1;
40
        mu[1] = 1;
41
        for(int i = 2; i \leq H; ++ i){
42
            if(!V[i]){
43
44
                P[++ p] = i;
                ph[i] = i - 1;
45
                mu[i] = -1;
46
47
            for(int j = 1; j \le p \& P[j] \le H / i; ++ j){
48
                int &p = P[j];
49
                V[i * p] = true;
50
                if(i \% p = 0){
51
52
                     ph[i * p] = ph[i] * p;
                     mu[i * p] = 0;
53
54
                     break;
                } else {
55
56
                     ph[i * p] = ph[i] * (p - 1);
```

```
mu[i * p] = -mu[i];
57
                  }
58
             }
59
60
        for(int i = 1; i \leq H; ++ i){
61
62
             sph[i] = sph[i - 1] + ph[i];
             smu[i] = smu[i - 1] + mu[i];
63
64
        int T;
65
        cin >> T;
66
67
        while(T \rightarrow \emptyset){
             int n;
68
69
             cin >> n;
             cout << solve_ph(n) << " " << solve_mu(n) << "\n";
70
71
72
        return 0;
73
    }
```

### 6.19 PN 筛

### 6.19.1 用法

对于积性函数 f(x),寻找积性函数 g(x) 使得 g(p) = f(p),且 g 易求前缀和 G。 令  $h = f * g^{-1}$ ,可以证明只有 PN 处 h 的函数值非 0,PN 指每个素因子幂次都不小于 2 的数。同时可以证明 n 以内的 PN 只有  $\mathcal{O}(\sqrt{n})$  个,且可以暴力枚举质因子幂次得到所有 PN。可利用下面公式计算  $h(p^c)$ :

$$h(p^c) = f(p^c) - \sum_{i=1}^{c} g(p^i) \times h(p^{c-i})$$

#### 6.19.2 例题

定义积性函数 f(x) 满足  $f(p^k) = p^k(p^k - 1)$ , 计算  $\sum f(i)$ 。

取  $g(p) = id(p)\varphi(p) = f(p)$ ,根据  $g*id = id_2$  利用杜教筛求解。 $h(p^c)$  的值利用递推式进行计算。

```
#include<bits/stdc++.h>
   using namespace std;
3
   const int MAXN = 1e7 + 3;
  const int MAXM = 1e5 + 3;
   const int H = 1e7;
5
   const int MOD = 1e9 + 7;
   const int DIV2 = 5000000004;
7
   const int DIV6 = 166666668;
8
   int P[MAXN], p; bool V[MAXN];
   int g[MAXN], le[MAXN], ge[MAXN];
10
11
   int s1(long long n){
                          // 1^1 + 2^1 + ... + n^1
12
       n \% = MOD;
       return 1ll * n * (n + 1) % MOD * DIV2 % MOD;
13
14
   int s2(long long n){ // 1^2 + 2^2 + ... + n^2
15
16
       n \% = MOD;
       return 1ll * n * (n + 1) % MOD * (2 * n + 1) % MOD * DIV6 % MOD;
17
18
19 int sg(long long n, long long N){
```

```
return n \leq H? le[n] : ge[N / n];
20
21
22
   int sieve_du(long long N){
23
        for(int d = N / H; d \ge 1; -- d)
            long long n = N / d;
24
25
            int wh = s2(n);
            for(long long l = 2, r; l \leq n; l = r + 1){
26
27
                r = n / (n / l);
                int wg = (s1(r) - s1(l - 1) + MOD) \% MOD;
28
                int ws = sg(n / l, N);
29
30
                ge[d] = (ge[d] + 1ll * wg * ws) % MOD;
31
            ge[d] = (wh - ge[d] + MOD) % MOD;
32
33
        return N \leq H? le[N] : ge[1];
34
35
36
   vector <int> hc[MAXM], gc[MAXM];
37
   int ANS;
   void sieve pn(int last, long long x, int h, long long N){
38
39
        ANS = (ANS + 1ll * h * sg(N / x, N)) % MOD;
        for(long long i = last + 1; x \leq N / P[i] / P[i]; ++ i){
40
41
            int c = 2;
            for(long long t = x * P[i] * P[i];t \leq N;t *= P[i], c ++){
42
43
                int hh = 1ll * h * hc[i][c] % MOD;
44
                sieve_pn(i, t, hh, N);
45
        }
46
47
48
   int main(){
49
        ios :: sync_with_stdio(false);
50
        cin.tie(nullptr);
        g[1] = 1;
51
        for(int i = 2; i \leq H; ++ i){
52
53
            if(!V[i]){
                P[++ p] = i, g[i] = 111 * i * (i - 1) % MOD;
54
55
            for(int j = 1; j \le p \& P[j] \le H / i; ++ j){
56
57
                int \delta p = P[j];
                V[i * p] = true;
58
59
                if(i \% p = 0){
                     g[i * p] = 111 * g[i] * p % MOD * p % MOD;
60
61
                     break;
62
                } else {
63
                    g[i * p] = 1ll * g[i] * p % MOD * (p - 1) % MOD;
                }
64
65
            }
66
67
        for(int i = 1; i \leq H; ++ i){}
68
            le[i] = (le[i - 1] + g[i]) % MOD;
69
70
       long long N;
71
        cin >> N;
        for(int i = 1;i ≤ p & 1ll * P[i] * P[i] ≤ N;i ++){
72
73
            int &p = P[i];
74
            hc[i].push_back(1);
75
            gc[i].push_back(1);
            for(long long c = 1, t = p; t \leq N; t = t * p, ++ c){}
76
```

```
if(c = 1){
77
                    gc[i].push_back(1ll * p * (p - 1) % MOD);
78
79
                } else {
                    gc[i].push_back(1ll * gc[i].back() * p % MOD * p % MOD)
80
81
                int w = 1ll * (t % MOD) * ((t - 1) % MOD) % MOD;
82
83
                int s = 0;
84
                for(int j = 1; j \leq c; ++ j){
                    s = (s + 1ll * gc[i][j] * hc[i][c - j]) % MOD;
85
86
                hc[i].push_back((w - s + MOD) % MOD);
87
            }
88
       }
89
       sieve_du(N);
90
       sieve_pn(0, 1, 1, N);
91
92
       cout << ANS << "\n";
93
       return 0;
94
```

### 6.20 常用数表

### 6.20.1 分拆数表

$\overline{n}$	10	20	30	40	50	60	70	80	90	100
p(n)	42	627	5604	37338	204226	966467	4087968	15796476	56634173	190569292

### 6.20.2 因数个数表

$\overline{N}$	10 <sup>1</sup>	$10^{2}$	$10^{3}$	$10^{4}$	$10^{5}$	$10^{6}$	$10^{7}$	108	10 <sup>9</sup>
$\max d(n)$	4	12	32	64	128	240	448	768	1344
$\max \omega(n)$	2	3	4	5	6	7	8	8	9
N	10 <sup>10</sup>	10 <sup>11</sup>	$10^{12}$	$10^{13}$	$10^{14}$	$10^{15}$	$10^{16}$	$10^{17}$	$10^{18}$
$\max d(n)$	2304	4032	6720	10752	17280	26880	41472	64512	103680
$\max \omega(n)$	10	10	11	12	12	13	13	14	15

### 6.20.3 大质数

1018 级别:

- $P = 10^{18} + 3$ , 好记。
- P = 2924438830427668481,可以进行 NTT, $P = 174310137655 \times 2^24 + 1$ ,原根为 3。

### 6.21 二次剩余

#### 6.21.1 用法

多次询问,每次询问给定奇素数 p 以及 y,在  $\mathcal{O}(\log p)$  复杂度计算 x 使得  $x^2 \equiv 0 \pmod p$  或者无解。

```
#include<bits/stdc++.h>
 1
   using namespace std;
 3
   int power(int a, int b, int p){
 4
       int r = 1;
       while(b){
 5
 6
            if(b & 1) r = 1ll * r * a % p;
 7
            b >>= 1, a = 1ll * a * a % p;
 8
9
       return r;
10
11
   bool check(int x, int p){
12
       return power(x, (p - 1) / 2, p) = 1;
13
14
   struct Node {
15
       int real, imag;
16
   };
   Node mul(const Node a, const Node b, int p, int v){
17
        int nreal = (1ll * a.real * b.real + 1ll * a.imag * b.imag % p * v)
18
            % p;
19
        int nimag = (1ll * a.real * b.imag + 1ll * a.imag * b.real) % p;
       return { (nreal), nimag };
20
21
   Node power(Node a, int b, int p, int v){
22
23
       Node r = \{ 1, 0 \};
       while(b){
24
            if(b & 1) r = mul(r, a, p, v);
25
26
            b >>= 1, a = mul(a, a, p, v);
27
28
       return r;
29
30
   mt19937 MT;
   void solve(int n, int p, int &x1, int &x2){
31
32
       if(n = 0){
            x1 = x2 = 0;
33
34
            return;
35
        if(!check(n, p)){
36
37
            x1 = x2 = -1;
38
            return;
39
40
       int a, t;
41
       do {
            a = MT() \% p;
42
43
        }while(check(t = (1ll * a * a - n + p) % p, p));
       Node u = \{ a, 1 \};
44
45
       x1 = power(u, (p + 1) / 2, p, t).real;
46
       x2 = (p - x1) \% p;
47
        if(x1 > x2)
48
            swap(x1, x2);
49
50 | int main(){
```

```
ios :: sync_with_stdio(false);
51
52
        cin.tie(nullptr);
53
        int T;
54
        cin >> T;
        while(T --){
55
56
            int n, p, x1, x2;
            cin \gg n \gg p;
57
58
            solve(n, p, x1, x2);
            if(x1 = -1){
59
                 cout << "Hola!\n";</pre>
60
            } else {
61
                 if(x1 = x2){
62
63
                     cout << x1 << "\n";
64
                     cout << x1 << " " << x2 << "\n";
65
66
67
            }
68
69
        return 0;
70
```

### 6.22 单位根反演

#### 6.22.1 定理

给出单位根反演如下:

$$[d\mid n] = \frac{1}{d} \sum_{i=0}^{d-1} \omega_d^{ni}$$

因为题太难了不会做, 所以没有例题, 咕着。

# 7 多项式

### 7.1 NTT 全家桶

### 7.1.1 用法

多项式全家桶。

- 包含基础多项式算法: 快速傅里叶变换(FFT)及其逆变换(IFFT)、快速数论变换(NTT)及 其逆变换(INTT);
- 包含基于 NTT 的扩展多项式算法:多项式乘法(MUL)、多项式乘法逆元(INV)、多项式微分(DIF)、多项式积分(INT)、多项式对数(LN)、多项式指数(EXP)、多项式开根(SQT)、多项式平移(即计算 G(x) = F(x+c),SHF)。

```
#include<bits/stdc++.h>
#define up(l, r, i) for(int i = l, END##i = r;i ≤ END##i;++ i)
#define dn(r, l, i) for(int i = r, END##i = l;i ≥ END##i;-- i)
using namespace std;
using i64 = long long;
const int INF = 1e9;
const i64 INFL = 1e18;
const int MOD = 998244353;
```

```
int power(int a, int b){
9
10
        int r = 1;
11
       while(b){
12
            if(b & 1) r = 1ll * r * a % MOD;
13
            b >>= 1, a = 1ll * a * a % MOD;
14
15
       return r;
16
   int inv(int x){
17
18
       return power(x, MOD - 2);
19
   const int MAX_{-} = (1 << 19) + 3;
20
21
   struct cplx{
       double a, b; cplx(double a = 0, double b = 0) :a(a), b(b){}
22
        cplx operator +(cplx t){ return cplx(a + t.a, b + t.b); }
23
24
        cplx operator -(cplx t){ return cplx(a - t.a, b - t.b); }
25
       cplx operator *(cplx t){ return cplx(a * t.a - b * t.b, a * t.b + b
            * t.a): }
        cplx operator *(int t) { return cplx(a * t, b * t); }
26
27
   };
   const long double pi = acos(-1);
28
29
   namespace Poly{
       void FFT(int n, cplx Z[]){
30
            static int W[MAX ];
31
32
            int l = 1; W[0] = 0;
33
            while (n >>= 1)
34
                up(0, l - 1, i)
                    W[l++] = W[i] << 1 | 1, W[i] <<= 1;
35
36
            up(0, l - 1, i)
37
                if(W[i] > i) swap(Z[i], Z[W[i]]);
            for (n = l >> 1, l = 1;n;n >>= 1, l <<= 1){
38
                cplx*S = Z, o(cos(pi / l), sin(pi / l));
39
                up(0, n - 1, i){
40
41
                    cplx s(1, 0);
42
                    up(0, l - 1, j){
                        cplx x = S[j] + s * S[j + l];
43
                        cplx y = S[j] - s * S[j + l];
44
45
                        S[j] = x, S[j + l] = y, s = s * o;
46
47
                    S += l << 1;
                }
48
            }
49
50
       void IFFT(int n, cplx Z[]){
51
52
            FFT(n, Z); reverse(Z + 1, Z + n);
53
            up(0, n - 1, i)
                Z[i].a \neq 1.0 * n, Z[i].b \neq 1.0 * n;
54
55
       void NTT(int n, int Z[]){
56
57
            static int W[MAX_];
58
            int g = 3, l = 1;
           W[0] = 0;
59
60
           while (n >>= 1)
61
                up(0, l - 1, i)
                    W[l++] = W[i] << 1 | 1, W[i] <<= 1;
62
63
            up(0, l - 1, i)
64
                if (W[i] > i)swap(Z[i], Z[W[i]]);
```

```
for (n = l >> 1, l = 1;n;n >>= 1, l <<= 1){
 65
                 int* S = Z, o = power(g, (MOD - 1) / l / 2);
 66
 67
                 up(0, n - 1, i){
 68
                     int s = 1;
                     up(0, l - 1, j){
 69
                         int x = (S[j] + 1ll * s * S[j + l] % MOD
                                                                          ) %
 70
                         int y = (S[j] - 1ll * s * S[j + l] % MOD + MOD) %
 71
                            MOD:
                         S[j] = x, S[j + l] = y;
 72
 73
                         s = 111 * s * o % MOD;
 74
                     }
 75
                     S += l << 1;
                 }
 76
             }
 77
 78
 79
        void INTT(int n, int Z[]){
             NTT(n, Z); reverse(Z + 1, Z + n);
 80
 81
             int o = inv(n);
 82
             up(0, n - 1, i)
                 Z[i] = 111 * Z[i] * o % MOD;
 83
 84
                                                     // 乘法
        void MUL(int n, int A[], int B[]){
 85
            NTT(n, A), NTT(n, B);
 86
 87
             up(0, n - 1, i)
                 A[i] = 111 * A[i] * B[i] % MOD;
 88
 89
             INTT(n, A);
 90
        void INV(int n, int Z[], int T[]){
 91
                                                     // 乘法逆
 92
             static int A[MAX_];
 93
             up(0, n - 1, i)
 94
                 T[i] = 0;
             T[0] = power(Z[0], MOD - 2);
 95
             for (int l = 1;l < n;l <<= 1){</pre>
 96
                        0, 2 * l - 1, i) A[i] = Z[i];
 97
                 up(2 * l, 4 * l - 1, i) A[i] = 0;
 98
 99
                 NTT(4 * l, A), NTT(4 * l, T);
100
                 up(0, 4 * l - 1, i)
                     T[i] = (2ll * T[i] - 1ll * A[i] * T[i] % MOD * T[i] %
101
                        MOD + MOD) % MOD;
102
                 INTT(4 * l, T);
103
                 up(2 * l, 4 * l - 1, i)
104
                     T[i] = 0;
             }
105
106
         void DIF(int n, int Z[], int T[]){
107
             up(0, n - 2, i)
108
                 T[i] = 111 * Z[i + 1] * (i + 1) % MOD;
109
110
            T[n - 1] = 0;
111
        void INT(int n, int c, int Z[], int T[]){
112
             up(1, n - 1, i)
113
                 T[i] = 111 * Z[i - 1] * inv(i) % MOD;
114
115
             T[0] = c;
116
        void LN(int n, int* Z, int* T){
                                                   // 求对数
117
             static int A[MAX_];
118
```

```
static int B[MAX ];
119
120
             up(0, 2 * n - 1, i)
                 A[i] = B[i] = 0;
121
122
             DIF(n, Z, A);
             INV(n, Z, B);
123
124
             MUL(2 * n, A, B);
             INT(n, 0, A, T);
125
126
         }
        void EXP(int n, int* Z, int* T){
                                                 // 求指数
127
             static int A[MAX ];
128
129
             static int B[MAX_];
             up(1, 2 * n - 1, i) T[i] = 0;
130
131
             T[0] = 1;
             for (int l = 1; l < n; l <<= 1){
132
                 LN (2 * 1, T, A);
133
                        0, 2 * l - 1, i)
134
                 up(
135
                     B[i] = (-A[i] + Z[i] + MOD) \% MOD;
                 B[0] = (B[0] + 1) \% MOD;
136
                 up(2 * l, 4 * l - 1, i)
137
138
                     T[i] = B[i] = 0;
                 MUL(4 * l, T, B);
139
             }
140
141
        void SQT(int n, int* Z, int* T){
                                                   // 开根
142
143
             static int A[MAX_];
             static int B[MAX_];
144
145
             up(1, 2 * n - 1, i) T[i] = 0;
             T[0] = 1;
146
147
             int o = inv(2);
             for (int l = 1;l < n;l <<= 1){</pre>
148
                 INV(2 * l, T, A);
149
                 up(0, 2 * l - 1, i)
150
                     B[i] = Z[i];
151
                 up(2 * l, 4 * l - 1, i)
152
153
                     A[i] = B[i] = 0;
                 MUL(4 * l, A, B);
154
                 up(0, 2 * l - 1, i)
155
156
                     T[i] = 111 * (T[i] + A[i]) * o % MOD;
             }
157
158
        void SHF(int n, int c, int* Z, int* T){
159
                                                   // 平移
             static int A[MAX_];
160
161
             static int B[MAX ];
162
             static int F[MAX_];
163
             static int G[MAX_];
             int o = 1;
164
             up(1, n - 1, i)
165
                 F[i] = 111 * F[i - 1] *
166
                                               i % MOD,
                 G[i] = 1ll * G[i - 1] * inv(i) % MOD;
167
168
             up(0, n - 1, i)
                 A[i] = 111 * Z[n - 1 - i] * F[n - 1 - i] % MOD;
169
             up(0, n - 1, i){
170
171
                 B[i] = 111 * G[i] * o % MOD;
172
                 o = 111 * o * c % MOD;
173
             int l = 1; while (l < 2 * n - 1) l <<= 1;</pre>
174
             up(n, l - 1, i)
175
```

```
A[i] = B[i] = 0;

MUL(l, A, B);

up(0, n - 1, i)

T[n - 1 - i] = 1ll * G[n - 1 - i] * A[i] % MOD;

180
}
```

### 7.2 FWT 全家桶

#### 7.2.1 用法

沃尔什全家桶。

包含与卷积、或卷积、异或卷积, 定义分别为二进制与、或、异或带入下式:

$$b_k = \sum_{i \otimes j = k} a_i \times b_j$$

```
#include<bits/stdc++.h>
 1
   #define up(l, r, i) for(int i = l, END##i = r;i \leq END##i;++ i)
 2
   #define dn(r, l, i) for(int i = r, END##i = l; i \ge END##i; -- i)
   using namespace std;
   using i64 = long long;
 5
   const int INF = 1e9;
 6
   const i64 INFL = 1e18;
 7
   const int MOD = 998244353;
   namespace Solve1{
                        // and 卷积
9
        void FWT(int n, int *A){
10
            for(int l = 1 << n, u = 2, v = 1; u ≤ l; u <<= 1, v <<= 1)</pre>
11
                for(int j = 0; j < l; j += u)</pre>
12
                    for(int k = 0; k < v; ++ k)
13
                         A[j + v + k] = (A[j + v + k] + A[j + k]) \% MOD;
14
15
16
        void IFWT(int n, int *A){
            for(int l = 1 << n, u = l, v = l / 2;u > 1;u >>= 1, v >>= 1)
17
                for(int j = 0; j < l; j += u)
18
                    for(int k = 0; k < v; ++ k)
19
                         A[j + v + k] = (A[j + v + k] - A[j + k] + MOD) %
20
                            MOD:
        }
21
22
23
   namespace Solve2{ // or 卷积
        void FWT(int n, int *A){
24
            for(int l = 1 << n, u = 2, v = 1; u ≤ l; u <<= 1, v <<= 1)</pre>
25
                for(int j = 0; j < l; j += u)
26
                    for(int k = 0; k < v; ++ k)
27
                         A[j + k] = (A[j + k] + A[j + v + k]) \% MOD;
28
29
        void IFWT(int n, int *A){
30
            for(int l = 1 << n, u = l, v = l / 2;u > 1;u >>= 1, v >>= 1)
31
                for(int j = 0; j < l; j += u)
32
33
                    for(int k = 0; k < v; ++ k)
                         A[j + k] = (A[j + k] - A[j + v + k] + MOD) \% MOD;
34
        }
35
36
   namespace Solve3{ // xor 卷积
37
        void FWT(int n, int *A){
38
            for(int l = 1 << n, u = 2, v = 1;u ≤ l;u <<= 1, v <<= 1)</pre>
39
```

```
for(int j = 0; j < l; j += u)
40
41
                    for(int k = 0; k < v; ++ k){
                         int a = A[j + k];
42
43
                         int b = A[j + v + k];
                                   ] = (a + b + MOD) \% MOD;
44
                         A[j + k]
45
                         A[j + v + k] = (a - b + MOD) \% MOD;
                    }
46
47
        void IFWT(int n, int *A){
48
            int div2 = (MOD + 1) / 2;
49
50
            for(int l = 1 << n, u = l, v = l / 2;u > 1;u >>= 1, v >>= 1)
                for(int j = 0; j < l; j += u)
51
52
                    for(int k = 0; k < v; ++ k){
53
                         int a = A[j + k];
                         int b = A[j + v + k];
54
55
                                    ] = 1ll * (a + b + MOD) * div2 % MOD;
                         A[j + k]
56
                         A[j + v + k] = 111 * (a - b + MOD) * div2 % MOD;
                    }
57
        }
58
59
```

### 7.3 任意模数 NTT

```
#include<bits/stdc++.h>
   #define up(l, r, i) for(int i = l, END##i = r; i \leq END##i; ++ i)
   #define dn(r, l, i) for(int i = r, END##i = l;i \ge END##i;-- i)
3
4 using namespace std;
5
   using i64 = long long;
   const int INF = 1e9;
7
   const i64 INFL = 1e18;
   const int MAX_{-} = (1 << 19) + 3;
8
   template <typename T>
9
10
   struct cplx0{
       T a, b; cplx0(T _a = 0, T _b = 0) :a(_a), b(_b){}
11
       cplx0 operator +(cplx0 t){ return cplx0(a + t.a, b + t.b); }
12
       cplx0 operator -(cplx0 t){ return cplx0(a - t.a, b - t.b); }
13
14
       cplx0 operator *(cplx0 t){ return cplx0(a * t.a - b * t.b, a * t.b
           + b * t.a); }
15
       cplx0 operator *(int t) { return cplx0(a * t, b * t); }
16
17
   using cplx = cplx0<double>;
18
   const long double pi = acos(-1);
19
   namespace Poly{
       void FFT(int n, cplx Z[]){
20
21
            static int W[MAX ];
22
            int l = 1; W[0] = 0;
23
           while (n >>= 1)
24
                up(0, l - 1, i)
                    W[l++] = W[i] << 1 | 1, W[i] <<= 1;
25
           up(0, l - 1, i)
26
                if(W[i] > i) swap(Z[i], Z[W[i]]);
27
            for (n = l >> 1, l = 1;n;n >>= 1, l <<= 1){
28
29
                cplx*S = Z;
                cplx0<long double> o(cosl(pi / l), sinl(pi / l));
30
31
                up(0, n - 1, i){
                    cplx0<long double> s(1, 0);
32
33
                    up(0, l - 1, j){
```

```
34
                         cplx x = S[j] + cplx(s.a, s.b) * S[j + l];
                         cplx y = S[j] - cplx(s.a, s.b) * S[j + l];
35
                         S[j] = x, S[j + l] = y, s = s * o;
36
37
                    S += l << 1;
38
                }
39
            }
40
41
42
        void IFFT(int n, cplx Z[]){
43
            FFT(n, Z); reverse(Z + 1, Z + n);
44
            up(0, n - 1, i)
                Z[i].a \neq 1.0 * n, Z[i].b \neq 1.0 * n;
45
46
        }
47
   const int MAXN = (1 << 19) + 3;
48
   const int BLOCK = 32768;
49
50
   cplx A1[MAXN], A2[MAXN], B1[MAXN], B2[MAXN];
51
   int n, m, L, mod;
   cplx P[MAXN], Q[MAXN];
52
53
   void FFTFFT(int L, cplx X[], cplx Y[]){
        for(int i = 0;i < L;++ i){</pre>
54
55
            P[i].a = X[i].a;
            P[i].b = Y[i].a;
56
57
        Poly :: FFT(L, P);
58
59
        for(int i = 0;i < L;++ i){</pre>
60
            Q[i] = (i = 0 ? P[0] : P[L - i]);
            Q[i].b = -Q[i].b;
61
62
63
        for(int i = 0; i < L; ++ i){
            X[i] = (P[i] + Q[i]);
64
            Y[i] = (Q[i] - P[i]) * cplx(0, 1);
65
            X[i].a \neq 2.0, X[i].b \neq 2.0;
66
67
            Y[i].a \neq 2.0, Y[i].b \neq 2.0;
        }
68
69
70
   int main(){
71
        ios :: sync_with_stdio(false);
72
        cin.tie(nullptr);
73
        cin >> n >> m >> mod;
        for(int i = 0; i \leq n; ++ i){
74
75
            int a; cin >> a;
76
            a %= mod;
77
            A1[i].a = a / BLOCK;
78
            A2[i].a = a \% BLOCK;
79
        for(int i = 0;i ≤ m;++ i){
80
81
            int a; cin >> a;
82
            a %= mod;
83
            B1[i].a = a / BLOCK;
84
            B2[i].a = a \% BLOCK;
85
86
        for(L = 1; L \le n + m; L <<= 1);
87
        FFTFFT(L, A1, A2);
        FFTFFT(L, B1, B2);
88
89
        for(int i = 0;i < L;++ i){
            P[i] = A1[i] * B1[i] + cplx(0, 1) * A2[i] * B1[i];
90
```

8 字符串 106

```
91
             Q[i] = A1[i] * B2[i] + cplx(0, 1) * A2[i] * B2[i];
         }
 92
 93
         Poly :: IFFT(L, P);
         Poly :: IFFT(L, Q);
 94
 95
         for(int i = 0; i < L; ++ i){</pre>
             long long a1b1 = P[i].a + 0.5;
 96
 97
             long long a2b1 = P[i].b + 0.5;
 98
             long long a1b2 = Q[i].a + 0.5;
             long long a2b2 = Q[i].b + 0.5;
 99
             long long w = ((a1b1 % mod * (BLOCK * BLOCK % mod)) + ((a2b1 +
100
                a1b2) % mod) * BLOCK + a2b2) % mod;
101
             if(i \leq n + m)
102
                 cout << w << " ";
         }
103
104
         return 0;
105
```

# 8 字符串

### 8.1 AC 自动机

```
#include<bits/stdc++.h>
   using namespace std;
   using i64 = long long;
   const int MOD = 1e9 + 7;
 5
   namespace ACAM{
        const int MAXN =1e6 + 3;
 6
7
        const int MAXM = 26 + 3;
 8
        int C[MAXN][MAXM], o;
 9
        void insert(char *S){
            int p = 0, len = 0;
10
            for(int i = 0;S[i];++ i){
11
                int e = S[i] - 'a';
12
                if(C[p][e]){
13
14
                     p = C[p][e];
15
                } else {
16
                     p = C[p][e] = ++ o;
17
18
                ++ len;
            }
19
20
21
        int F[MAXN];
        void build(){
22
23
            queue <int> Q; Q.push(0);
24
            while(!Q.empty()){
                int u = Q.front(); Q.pop();
25
                for(int i = 0; i < 26; ++ i){
26
27
                     int v = C[u][i];
                     if(v = 0)
28
29
                         continue;
30
                     int p = F[u];
31
                    while(!C[p][i] & p \neq 0)
32
                         p = F[p];
33
                     if(C[p][i] & C[p][i] \neq v)
34
                         F[v] = C[p][i];
                    Q.push(v);
35
```

8 字符串 107

# 8.2 扩展 KMP

#### 8.2.1 定义

```
z_i^{(1)} = |\operatorname{lcp}(b, \operatorname{suffix}(b, i))|z_i^{(2)} = |\operatorname{lcp}(b, \operatorname{suffix}(a, i))|
```

```
#include<bits/stdc++.h>
 1
   using namespace std;
 3
   typedef long long i64;
   const int MAXN = 2e7 + 3;
 5 | char A[MAXN], B[MAXN * 2];
 6
   | int n, m, l, r, Z[MAXN * 2];
 7
   i64 ans1, ans2;
   int main(){
8
9
        scanf("%s%s", A + 1, B + 1);
        n = strlen(A + 1);
10
11
       m = strlen(B + 1);
        l = 0, r = 0; Z[1] = 0, ans1 = m + 1;
12
13
        for(int i = 2; i \leq m; ++ i){
            if(i \le r) Z[i] = min(r - i + 1, Z[i - l + 1]);
14
15
                       Z[i] = 0;
            while(B[Z[i] + 1] = B[i + Z[i]])
16
                ++ Z[i];
17
            if(i + Z[i] - 1 > r)
18
                r = i + Z[i] - 1, l = i;
19
            ans1 ^- 1ll * i * (Z[i] + 1);
20
        }
21
22
       l = 0, r = 0;
        Z[1] = 0, B[m + 1] = '#', strcat(B + 1, A + 1);
23
24
        for(int i = 2; i \le n + m + 1; ++ i){}
25
            if(i \le r) Z[i] = min(r - i + 1, Z[i - l + 1]);
26
                       Z[i] = 0;
            while(B[Z[i] + 1] = B[i + Z[i]])
27
                ++ Z[i];
28
            if(i + Z[i] - 1 > r)
29
                r = i + Z[i] - 1, l = i;
30
31
        for(int i = m + 2; i \le n + m + 1; ++ i){}
32
33
            ans2 ^{-1}1 ll * (i - m - 1) * (Z[i] + 1);
34
        printf("%lld\n%lld\n", ans1, ans2);
35
36
        return 0;
37
```

### 8.3 Manacher

```
#include<bits/stdc++.h>
using namespace std;
using i64 = long long;
```

8 字符串 108

```
const int INF = 1e9;
 5
   const i64 INFL = 1e18;
   const int MAXN= 2.2e7 + 11;
 7
   char S[MAXN], T[MAXN]; int n, R[MAXN];
   int main(){
 8
        scanf("%s", S + 1);
9
        n = strlen(S + 1);
10
11
        for(int i = 1; i \leq n; ++ i){
            T[2 * i - 1] = S[i];
12
                      ] = '#':
13
            T[2 * i]
14
       T[0] = '#';
15
16
        n = 2 * n;
17
        int p = 0, x = 0, ans = 0;
        for(int i = 1; i \leq n; ++ i){
18
            if(i \le p) R[i] = min(R[2 * x - i], p - i);
19
20
            while(i - R[i] - 1 \geq 0 & T[i + R[i] + 1] = T[i - R[i] - 1])
                ++ R[i];
21
            if(i + R[i] > p){
22
23
                p = i + R[i];
24
                x = i;
25
            }
26
            ans = max(ans, R[i]);
27
        printf("%d\n", ans);
28
29
        return 0;
30
```

## 8.4 回文自动机

```
#include<bits/stdc++.h>
   using namespace std;
   const int MAXM = 26 + 3;
 4
   namespace PAM{
        const int SIZ = 5e5 + 3;
 5
        int n, s, F[SIZ], L[SIZ], D[SIZ];
6
7
        int M[SIZ][MAXM];
        char S[SIZ];
8
9
        void init(){
            S[0] = '\$', n = 1;
10
            F[s = 0] = -1, L[0] = -1, D[0] = 0;
11
            F[s = 1] = 0, L[1] = 0, D[1] = 0;
12
        }
13
        void extend(int &last, char c){
14
15
            S[++ n] = c;
            int e = c - 'a';
16
            int a = last;
17
18
            while (c \neq S[n - 1 - L[a]])
                a = F[a];
19
            if(M[a][e]){
20
                last = M[a][e];
21
            } else {
22
23
                int cur = M[a][e] = ++ s;
                L[cur] = L[a] + 2;
24
25
                if(a = 0){
                    F[cur] = 1;
26
27
                } else {
```

```
int b = F[a];
28
29
                     while (c \neq S[n - 1 - L[b]])
                         b = F[b];
30
31
                     F[cur] = M[b][e];
32
                D[cur] = D[F[cur]] + 1;
33
                last = cur;
34
35
            }
        }
36
37
38
   const int MAXN = 5e5 + 3;
   char T[MAXN];
39
40
   int main(){
41
        PAM :: init();
        int m = 0, last = 0, lastans = 0;
42
43
        for(char c = getchar();isalpha(c);c = getchar()){
44
            char d = (c - 97 + lastans) % 26 + 97;
45
            PAM :: extend(last, d);
            cout << (lastans = PAM :: D[last]) << " ";</pre>
46
47
        }
48
       return 0;
49
   }
   /*
50
51
   azzzyyzyyx
52 1 2 1 2 3 2 2 2 3 3
53
   */
```

## 8.5 后缀平衡树

#### 8.5.1 本代码尚未完成

## 8.6 后缀数组(倍增)

```
#include<bits/stdc++.h>
 2 using namespace std;
 3 using i64 = long long;
 4 const int INF = 1e9;
 5 | const i64 INFL = 1e18;
   const int MAXN = 1e6 + 3;
 6
 7
   int n, m;
   int A[MAXN], B[MAXN];
 8
   int C[MAXN], R[MAXN], P[MAXN], Q[MAXN];
9
   char S[MAXN];
10
   int main(){
11
12
        scanf("%s", S), n = strlen(S), m = 256;
13
        for(int i = 0;i < n;++ i) R[i] = S[i];</pre>
        for (int k = 1; k \le n; k <<= 1){
14
            for(int i = 0; i < n; ++ i){}
15
                Q[i] = ((i + k > n - 1) ? 0 : R[i + k]);
16
                P[i] = R[i];
17
18
                m = max(m, R[i]);
            }
19
20 | #define fun(a, b, c) \
            memset(C, 0, sizeof(int) * (m + 1));
21
22
            for(int i = 0;i < n;++ i) C[a] +=</pre>
                                                         1;
```

```
for(int i = 1; i \leq m; ++ i) C[i] += C[i - 1];
23
            for(int i = n - 1; i \ge 0; -- i) c[-- C[a]] = b;
24
25
            fun(Q[ i ], i , B)
26
            fun(P[B[i]], B[i], A)
   #undef fun
27
28
            int p = 1; R[A[0]] = 1;
            for(int i = 1; i \le n - 1; ++ i){
29
30
                bool f1 = P[A[i]] = P[A[i - 1]];
                bool f2 = Q[A[i]] = Q[A[i - 1]];
31
                R[A[i]] = f1 \& f2 ? R[A[i-1]] : ++ p;
32
33
            if (m = n) break;
34
35
36
       for(int i = 0; i < n; ++ i)
            printf("%u ", A[i] + 1);
37
38
       return 0;
39
   }
```

# 8.7 后缀数组(SAIS)

```
1 #include < bits/stdc++.h>
   using namespace std;
3 using i64 = long long;
   const int INF = 1e9;
4
   const i64 INFL = 1e18;
6
   const int MAXN = 1e6 + 3;
   const int MAXM = 256 + 3;
7
   #define LTYPE 0
8
   #define STYPE 1
9
10
   void induce_sort(int n, int S[], int T[], int m, int LM[], int SA[],
      int C[]){
       vector <int> BL(n);
11
       vector <int> BS(n);
12
13
       vector <int> BM(n);
14
       fill(SA, SA + n, -1);
       for(int i = 0; i < n; ++ i){
                                     // 预处理桶
15
16
           BM[i] = BS[i] = C[i] - 1;
           BL[i] = i = 0 ? 0 : C[i - 1];
17
18
19
       for(int i = m - 1;i ≥ 0;-- i) // 放置 LMS 后缀
           SA[BM[S[LM[i]]] --] = LM[i];
20
21
       for(int i =
                   0, p;i < n; ++ i) // 计算 L 类型后缀的位置
           if(SA[i] > 0 \& T[p = SA[i] - 1] = LTYPE)
22
               SA[BL[S[p]] ++] = p;
23
       for(int i = n - 1, p;i ≥ 0;-- i) // 计算 S 类型后缀的位置
24
25
           if(SA[i] > 0 \& T[p = SA[i] - 1] = STYPE)
26
               SA[BS[S[p]] --] = p;
27
   // 长度 n,字符集 [0, n),要求最后一个元素为 0
28
   // 例如输入 ababa 传入 n = 6, S = [1 2 1 2 1 0]
29
30
   void sais(int n, int S[], int SA[]){
       vector <int> T(n);
31
32
       vector <int> C(n);
33
       vector \langle int \rangle I(n, -1);
34
       T[n - 1] = STYPE;
       for(int i = n - 2;i ≥ 0;-- i){ // 递推类型
35
```

```
T[i] = S[i] = S[i + 1] ? T[i + 1] : (S[i] < S[i + 1] ? STYPE :
36
                LTYPE);
37
                                    // 统计个数
        for(int i = 0;i < n;++ i){</pre>
38
            C[S[i]] ++;
39
40
                                        // 前缀累加
        for(int i = 1; i < n; ++ i){</pre>
41
42
            C[i] += C[i - 1];
43
        vector <int> P:
44
45
        for(int i = 0;i < n;++ i){
                                     // 统计 LMS 后缀
            if(T[i] = STYPE & (i = 0 || T[i - 1] = LTYPE))
46
47
                I[i] = P.size(), P.push_back(i);
48
49
50
        int m = P.size(), tot = 0, cnt = 0;
51
        induce_sort(n, S, T.data(), m, P.data(), SA, C.data());
        vector <int> S0(m), SA0(m);
52
        for(int i = 0, x, y = -1; i < n; ++ i){
53
            if((x = I[SA[i]]) \neq -1){
54
                if(tot = 0 || P[x + 1] - P[x] \neq P[y + 1] - P[y])
55
56
                    tot ++;
                else for(int p1 = P[x], p2 = P[y]; p2 \leq P[y + 1]; ++ p1, ++
57
                    p2){
                     if((S[p1] \ll 1 \mid T[p1]) \neq (S[p2] \ll 1 \mid T[p2]))
58
59
                         tot ++; break;
60
61
62
                S0[y = x] = tot - 1;
            }
63
64
        if(tot = m){
65
            for(int i = 0;i < m;++ i)</pre>
66
67
                SA0[S0[i]] = i;
        } else {
68
            sais(m, S0.data(), SA0.data());
69
70
71
        for(int i = 0; i < m; ++ i)
            S0[i] = P[SA0[i]];
72
73
        induce_sort(n, S, T.data(), m, S0.data(), SA, C.data());
74
   int S[MAXN], SA[MAXN], H[MAXM], G[MAXM];
75
76
   int main(){
        int n = 0, t = 0, m = 256;
77
        for(char c = cin.get();isgraph(c);c = cin.get()){
78
79
            S[n \leftrightarrow] = c;
            H[c] ++;
80
81
82
        for(int i = 0; i < m; ++ i){</pre>
83
            t += !!H[i], G[i] = t;
84
        for(int i = 0;i < n;++ i){</pre>
85
86
            S[i] = G[S[i]];
87
        sais(n + 1, S, SA);
88
89
        for(int i = 1; i \leq n; ++ i){
            cout << SA[i] + 1 << " ";
90
```

```
91 }
92 return 0;
93 }
```

## 8.8 广义后缀自动机(离线)

```
#include<bits/stdc++.h>
   using namespace std;
   using i64 = long long;
 4
   const int INF = 1e9;
   const i64 INFL = 1e18;
 6
   const int MAXM= 26 + 3;
   namespace SAM{
 7
        const int SIZ = 2e6 + 3;
 8
9
        int M[SIZ][MAXM];
10
        int L[SIZ], F[SIZ], S[SIZ];
        int s = 0, h = 25;
11
        void init(){
12
13
            F[0] = -1, s = 0;
14
15
        void extend(int &last, char c){
            int e = c - 'a';
16
17
            int cur = ++ s;
            L[cur] = L[last] + 1;
18
19
            int p = last;
20
            while (p \neq -1 \& H[p][e])
                M[p][e] = cur, p = F[p];
21
22
            if(p = -1){
                F[cur] = 0;
23
24
            } else {
25
                int q = M[p][e];
26
                if(L[p] + 1 = L[q]){
                     F[cur] = q;
27
                } else {
28
29
                     int clone = ++ s;
30
                     L[clone] = L[p] + 1;
31
                     F[clone] = F[q];
                     for(int i = 0; i \leq h; ++ i)
32
33
                         M[clone][i] = M[q][i];
34
                     while(p \neq -1 & M[p][e] = q)
                         M[p][e] = clone, p = F[p];
35
                     F[cur] = F[q] = clone;
36
                }
37
38
39
            last = cur;
40
41
        void solve(){
42
            i64 \text{ ans} = 0;
            for(int i = 1; i \leq s; ++ i)
43
                ans += L[i] - L[F[i]];
44
45
            cout << ans << endl;</pre>
        }
46
47
   namespace Trie{
48
        const int SIZ = 1e6 + 3;
49
50
        int M[SIZ][MAXM], s, h = 25;
51
        void insert(char *S){
```

```
int p = 0;
52
            for(int i = 0;S[i];++ i){
53
54
                 int e = S[i] - 'a';
55
                 if(M[p][e]){
56
                     p = M[p][e];
57
                 } else
                     p = M[p][e] = ++ s;
58
            }
59
        }
60
        int 0[SIZ];
61
62
        void build_sam(){
            queue <int> Q;
63
64
            Q.push(0);
65
            while(!Q.empty()){
                 int u = Q.front(); Q.pop();
66
                 for(int i = 0; i \leq h; ++ i){
67
68
                     char c = i + 'a';
                     if(M[u][i]){
69
                          int v = M[u][i];
70
71
                          O[v] = O[u];
72
                          SAM :: extend(0[v], c);
73
                          Q.push(v);
                     }
74
75
                 }
            }
76
        }
77
78
79
   const int MAXN = 1e6 + 3;
80
   char S[MAXN];
81
   int main(){
82
        SAM :: init();
83
        int n, last = 0;
84
        cin >> n;
85
        for(int i = 1; i \leq n; ++ i){
            scanf("%s", S);
86
            Trie :: insert(S);
87
88
89
        Trie :: build_sam();
90
        SAM :: solve();
91
        cout << SAM :: s + 1 << endl;</pre>
92
        return 0;
93
```

### 8.9 广义后缀自动机(在线)

```
#include<bits/stdc++.h>
   using namespace std;
2
3 using i64 = long long;
   const int INF = 1e9;
   const i64 INFL = 1e18;
   const int MAXM= 26 + 3;
6
7
   namespace SAM{
8
       const int SIZ = 2e6 + 3;
9
       int M[SIZ][MAXM];
10
       int L[SIZ], F[SIZ], S[SIZ];
       int s = 0, h = 25;
11
12
       void init(){
```

```
13
            F[0] = -1, s = 0;
        }
14
15
        void extend(int &last, char c){
            int e = c - 'a';
16
17
            if(M[last][e]){
18
                 int p = last;
                 int q = M[last][e];
19
20
                 if(L[q] = L[last] + 1){
                     last = q;
21
22
                 } else {
23
                     int clone = ++ s;
24
                     L[clone] = L[p] + 1;
25
                     F[clone] = F[q];
                     for(int i = 0; i \leq h; ++ i)
26
27
                         M[clone][i] = M[q][i];
28
                     while(p \neq -1 & M[p][e] = q)
29
                         M[p][e] = clone, p = F[p];
30
                     F[q] = clone;
                     last = clone;
31
32
33
            } else {
34
                 int cur = ++ s;
35
                 L[cur] = L[last] + 1;
                 int p = last;
36
                 while (p \neq -1 \& \{M[p][e])
37
                     M[p][e] = cur, p = F[p];
38
                 if(p = -1){
39
                     F[cur] = 0;
40
                 } else {
41
42
                     int q = M[p][e];
43
                     if(L[p] + 1 = L[q]){
44
                          F[cur] = q;
                     } else {
45
46
                         int clone = ++ s;
                         L[clone] = L[p] + 1;
47
                          F[clone] = F[q];
48
                          for(int i = 0; i \leq h; ++ i)
49
50
                              M[clone][i] = M[q][i];
                         while (p \neq -1 \& M[p][e] = q)
51
52
                              M[p][e] = clone, p = F[p];
                          F[cur] = F[q] = clone;
53
                     }
54
55
56
                 last = cur;
            }
57
58
59
        void solve(){
60
            i64 \text{ ans} = 0;
61
            for(int i = 1; i \leq s; ++ i)
62
                 ans += L[i] - L[F[i]];
63
            cout << ans << endl;</pre>
        }
64
65
66
   const int MAXN = 1e6 + 3;
   char S[MAXN];
67
   int main(){
68
69
        SAM :: init();
```

```
70
        int n, last = 0;
71
        cin >> n;
72
        for(int i = 1; i \leq n; ++ i){
             scanf("%s", S);
73
74
            int m = strlen(S);
75
            last = 0;
            for(int j = 0; j < m; ++ j){
76
77
                 SAM :: extend(last, S[j]);
78
             }
79
80
        SAM :: solve();
81
        cout << SAM :: s + 1 << endl;</pre>
82
        return 0;
83
```

## 8.10 后缀自动机

```
#include<bits/stdc++.h>
   using namespace std;
 3
   using i64 = long long;
   const int INF = 1e9;
 5 const i64 INFL = 1e18;
   const int MAXM= 26 + 3;
 7
   namespace SAM{
        const int SIZ = 2e6 + 3;
 8
9
        int M[SIZ][MAXM];
        int L[SIZ], F[SIZ], S[SIZ];
10
11
        int last = 0, s = 0, h = 25;
12
        void init(){
13
            F[0] = -1, last = s = 0;
14
15
        void extend(char c){
            int cur = ++ s, e = c - 'a';
16
            L[cur] = L[last] + 1;
17
            S[cur] = 1;
18
19
            int p = last;
20
            while (p \neq -1 \& \{M[p][e]\})
                M[p][e] = cur, p = F[p];
21
22
            if(p = -1){
23
                F[cur] = 0;
            } else {
24
25
                int q = M[p][e];
26
                if(L[p] + 1 = L[q]){
27
                    F[cur] = q;
28
                } else {
29
                    int clone = ++ s;
30
                    L[clone] = L[p] + 1;
31
                    F[clone] = F[q];
32
                    S[clone] = 0;
                    for(int i = 0; i \leq h; ++ i)
33
                        M[clone][i] = M[q][i];
34
35
                    while(p \neq -1 & M[p][e] = q)
36
                        M[p][e] = clone, p = F[p];
                    F[cur] = F[q] = clone;
37
                }
38
            }
39
40
            last = cur;
```

```
41
        }
42
        vector <int> E[SIZ];
43
        void build(){
44
            for(int i = 1; i \leq s; ++ i){
45
                 E[F[i]].push_back(i);
46
        }
47
48
        i64 \text{ ans} = 0;
49
        void dfs(int u){
            for(auto &v : E[u]){
50
51
                 dfs(v), S[u] += S[v];
52
53
            if(S[u] > 1)
                 ans = max(ans, 1ll * S[u] * L[u]);
54
        }
55
56
57
   const int MAXN = 1e6 + 3;
58
   char S[MAXN];
   int main(){
59
60
        SAM :: init();
        scanf("%s", S); int n = strlen(S);
61
62
        for(int i = 0;i < n;++ i)</pre>
            SAM :: extend(S[i]);
63
64
        SAM :: build();
        SAM :: dfs (0);
65
        printf("%lld\n", SAM :: ans);
66
67
        return 0;
68
```

### 8.11 字典树

```
1 #include<bits/stdc++.h>
   using namespace std;
 3 using i64 = long long;
   const int INF = 1e9;
 5
   const i64 INFL = 1e18;
   const int MAXM= 10 + 3;
   namespace Trie{
 7
8
       const int SIZ = 1e6 + 3;
9
       int M[SIZ][MAXM], s, h = 10;
       void extend(int &last, char c){
10
            int e = c - 'a';
11
            if(M[last][e]){
12
13
                last = M[last][e];
14
15
                last = M[last][e] = ++ s;
16
17
        }
       void insert(char *S){
18
19
            int p = 0;
            for(int i = 0;S[i];++ i){
20
                int e = S[i] - 'a';
21
                if(M[p][e]){
22
23
                    p = M[p][e];
24
                } else
25
                    p = M[p][e] = ++ s;
26
            }
```

```
27 }
28 }
```

# 9 计算几何

## 9.1 二维凸包

#### 9.1.1 例题

给定 n 个点,保证每三点不共线。要求找到一个简单多边形满足它不是凸包,使得该多边形面积最大。

```
#include<bits/stdc++.h>
 1
   using namespace std;
   using i64 = long long;
 3
   const int MAXN = 2e5 + 3;
   int X[MAXN], Y[MAXN];
 5
   struct Frac {
 6
 7
        int a, b;
        Frac (int _a, int _b){
 8
 9
            if(_b < 0){
                a = -_a, b = -_b;
10
11
            } else {
12
                a = _a, b = _b;
13
14
        }
15
   };
16
   struct Node {
17
        int x, y;
   }P[MAXN];
18
19
   bool operator < (const Frac A, const Frac B){</pre>
20
        return 1ll * A.a * B.b - 1ll * A.b * B.a < 0;
21
   bool operator < (const Node A, const Node B){</pre>
22
23
        return A.x = B.x ? A.y > B.y : A.x < B.x;
24
25
   const Frac intersect(Node A, Node B){
        int a = B.y - A.y;
26
27
        int b = A.x - B.x;
        assert(b \neq 0);
28
        if(b < 0){
29
30
            a = -a, b = -b;
31
32
        return Frac(a, b);
33
   bool F[MAXN];
34
   int main(){
35
36
        int TT;
37
        cin >> TT;
        while(TT -- ){
38
39
            int n;
40
            cin >> n;
            int maxx = -1e9, minx = 1e9;
41
            for(int i = 1; i \leq n; ++ i){
42
43
                auto \delta[x, y] = P[i];
44
                cin \gg x \gg y;
45
                F[i] = false;
```

```
46
            sort(P + 1, P + 1 + n);
47
            vector <int> Q1, Q2, Q;
48
            // Q1 计算上凸壳, Q2 计算下凸壳
49
            for(int i = 1; i \leq n; ++ i){
50
51
                 auto \delta[x, y] = P[i];
                 if(Q1.size() \leq 1){
52
                     Q1.push_back(i);
53
54
                 } else {
                     while(Q1.size() \geq 2){
55
56
                         auto \delta[x1, y1] = P[Q1[Q1.size() - 1]];
                         auto \&[x2, y2] = P[Q1[Q1.size() - 2]];
57
58
                         long long cmp = 1ll * (y - y1) * (x1 - x2) - 1ll *
                             (x - x1) * (y1 - y2);
                          if(cmp > 0){
59
60
                              Q1.pop_back();
61
                          } else break;
62
63
                     Q1.push_back(i);
64
                 if(Q2.size() \leq 1){
65
66
                     Q2.push_back(i);
                 } else {
67
                     while(Q2.size() \geq 2){
68
                         auto \delta[x1, y1] = P[Q2[Q2.size() - 1]];
69
70
                         auto \delta[x2, y2] = P[Q2[Q2.size() - 2]];
                         long long cmp = 1ll * (y - y1) * (x1 - x2) - 1ll *
71
                             (x - x1) * (y1 - y2);
72
                          if(cmp < 0){
73
                              Q2.pop_back();
74
                          } else break;
75
76
                     Q2.push_back(i);
                 }
77
            }
78
79
            Q = Q1;
80
            for(int i = Q2.size(); i \neq 0; i \rightarrow 0
81
                 if(i \neq Q2.size())
82
                     Q.push_back(Q2[i - 1]);
83
84
            long long area = 0;
85
            int x0 = P[Q[0]].x;
86
            int y0 = P[Q[0]].y;
            for(int i = 1;i + 1 < Q.size(); ++ i){</pre>
87
                 auto \&[x1, y1] = P[Q[
88
                                            i]];
89
                 auto \delta[x2, y2] = P[Q[i + 1]];
                 area += 1ll * (x1 - x0) * (y2 - y0) - 1 ll * (x2 - x0) * (y1)
90
                     - y0);
91
            }
92
            area = -area;
            for(auto &i: Q1) F[i] = true;
93
            for(auto &i: Q2) F[i] = true;
94
95
            bool ok = false;
96
            for(int i = 1; i \le n; ++ i) if(!F[i]){
97
                 ok = true;
98
                 maxx = max(maxx, P[i].x);
99
                 minx = min(minx, P[i].x);
```

```
100
             if(!ok){
101
                 cout \ll -1 \ll "\n";
102
103
                 continue;
             }
104
105
             vector <int> L1;
             vector <int> L2;
106
             // L1 插入 kx + b 维护下凸壳
107
             for(int i = 1; i \le n; ++ i) if(!F[i]){
108
                 auto \delta[k, b] = P[i];
109
                 if(!L1.empty() & k = P[L1.back()].x)
110
111
                      continue;
112
                 while(L1.size() \geq 2){
                      auto &P1 = P[L1[L1.size() - 1]];
113
                      auto &P2 = P[L1[L1.size() - 2]];
114
                      Frac i1 = intersect(P1, P[i]);
115
116
                      Frac i2 = intersect(P2, P[i]);
117
                      if(i1 < i2){
                          L1.pop back();
118
119
                      } else break;
120
121
                 L1.push_back(i);
             }
122
             // L2 插入 kx + b 维护上凸壳
123
             for(int i = n; i \ge 1; -- i) if(!F[i]){
124
125
                 auto \delta[k, b] = P[i];
                 if(!L2.empty() & k = P[L2.back()].x)
126
127
                      continue;
                 while(L2.size() \geq 2){
128
129
                      auto &P1 = P[L2[L2.size() - 1]];
                      auto &P2 = P[L2[L2.size() - 2]];
130
                      Frac i1 = intersect(P1, P[i]);
131
132
                      Frac i2 = intersect(P2, P[i]);
133
                      if(i1 < i2){
134
                          L2.pop_back();
135
                      } else break;
136
137
                 L2.push_back(i);
138
139
             vector <Frac> E1;
             E1.push_back(Frac( -2e9, 1 ));
140
             for(int i = 0;i + 1 < L1.size();++ i){</pre>
141
142
                 auto &P1 = P[L1[i
                                        ]];
143
                 auto \&P2 = P[L1[i + 1]];
144
                 E1.push_back(intersect(P1, P2));
145
             vector <Frac> E2;
146
             E2.push_back(Frac( -2e9, 1 ));
147
             for(int i = 0; i + 1 < L2.size(); ++ i){</pre>
148
149
                 auto &P1 = P[L2[i
                                        ]];
                 auto \&P2 = P[L2[i + 1]];
150
151
                 E2.push_back(intersect(P1, P2));
152
153
             long long ans = 0;
             for(int i = 0; i + 1 < Q.size(); ++ i){</pre>
154
                 auto \delta[x1, y1] = P[Q[i]
155
                 auto \&[x2, y2] = P[Q[i + 1]];
156
```

```
157
                  long long w = 1ll * x2 * y1 - 1ll * x1 * y2;
158
                  int A = y2 - y1;
                  int B = x1 - x2;
159
160
                  int x = 0, y = 0;
                  if(B = 0){
161
162
                      if(A > 0){
                          x = minx, y = 0;
163
164
                      } else {
165
                          x = maxx, y = 0;
166
                  } else
167
                  if(B < 0){
168
169
                      Frac K = Frac(-A, -B);
170
                      int p = 0;
                      for(int k = 20; k \ge 0; -- k){
171
172
                           int pp = p | 1 << k;</pre>
173
                          if(pp < E1.size() & E1[pp] < K){
174
                               p = pp;
                           }
175
                      }
176
                      x = P[L1[p]].x;
177
178
                      y = P[L1[p]].y;
                  } else {
179
180
                      Frac K = Frac( A, B);
181
                      int p = 0;
182
                      for(int k = 20; k \ge 0; -- k){
183
                           int pp = p | 1 << k;
184
                           if(pp < E2.size() & E2[pp] < K){
185
                               p = pp;
                           }
186
187
                      }
                      x = P[L2[p]].x;
188
                      y = P[L2[p]].y;
189
190
                  ans = \max(\text{ans, area - }(w + 1ll * A * x + 1ll * B * y));
191
192
             // cerr << "ans = " << ans << endl;
193
194
             cout \ll ans \ll "\n";
         }
195
196
         return 0;
197
```

### 9.2 最小圆覆盖

```
#include "2d.cpp"
   point geto(point a, point b, point c) {
2
3
       double a1, a2, b1, b2, c1, c2;
4
       point ans(0, 0);
       a1 = 2 * (b.x - a.x), b1 = 2 * (b.y - a.y),
5
       c1 = sqr(b.x) - sqr(a.x) + sqr(b.y) - sqr(a.y);
6
       a2 = 2 * (c.x - a.x), b2 = 2 * (c.y - a.y),
7
       c2 = sqr(c.x) - sqr(a.x) + sqr(c.y) - sqr(a.y);
8
       if (equal(a1, 0)) {
9
            ans.y = c1 / b1;
10
11
            ans.x = (c2 - ans.y * b2) / a2;
       } else if (equal(b1, 0)) {
12
13
            ans.x = c1 / a1;
```

```
ans.y = (c2 - ans.x * a2) / b2;
14
15
        } else {
            ans.x = (c2 * b1 - c1 * b2) / (a2 * b1 - a1 * b2);
16
            ans.y = (c2 * a1 - c1 * a2) / (b2 * a1 - b1 * a2);
17
18
19
        return ans;
20
21
   mt19937 MT;
   circ minimal(vector <point> V){
22
23
        shuffle(V.begin(), V.end(), MT);
24
        point o = V[0];
25
        double r = 0;
26
        for(int i = 0;i < V.size();++ i) {</pre>
            if (sign(dis(o, V[i]) - r) \neq 1) continue;
27
            o.x = (V[i].x + V[0].x) / 2;
28
            o.y = (V[i].y + V[0].y) / 2;
29
30
            r = dis(V[i], V[0]) / 2;
            for(int j = 0; j < i; ++ j) {</pre>
31
                if (sign(dis(o, V[j]) - r) \neq 1) continue;
32
33
                o.x = (V[i].x + V[j].x) / 2;
                o.y = (V[i].y + V[j].y) / 2;
34
35
                r = dis(V[i], V[j]) / 2;
                for(int k = 0;k < j;++ k) {</pre>
36
                     if (sign(dis(o, V[k]) - r) \neq 1) continue;
37
                     o = geto(V[i], V[j], V[k]);
38
39
                     r = dis(o, V[i]);
40
            }
41
42
        }
43
        circ res;
44
        res.o = o;
45
        res.r = r;
46
        return res;
47
```

### 9.3 最左转线

```
#include "2d.cpp"
 1
 2
   namespace DSU{
        const int MAXN = 1e5 + 3;
 3
 4
        int F[MAXN];
 5
        int getfa(int u){
            return u = F[u] ? u : F[u] = getfa(F[u]);
 6
        }
 7
 8
 9
   namespace Dual{
        const int MAXN = 1e5 + 3;
10
11
        const int MAXM = 1e5 + 3;
        int A[MAXM], B[MAXM], W[MAXM], I[MAXM], n, m;
12
        int outer;
13
        bool cmp(int a, int b){
14
            return W[a] < W[b];</pre>
15
16
        }
       vector <pair<int, int> > E[MAXN];
17
18
        const int MAXT = 20 + 3;
        int F[MAXN][MAXT], G[MAXN][MAXT], D[MAXN], h = 20;
19
20
        void dfs(int u, int f){
```

```
21
            D[u] = D[f] + 1;
22
            for(int i = 1; i \leq h; ++ i)
                F[u][i] = F[F[u][i - 1]][i - 1],
23
                G[u][i] = max(G[u][i - 1], G[F[u][i - 1]][i - 1]);
24
25
            for(auto \delta[v, w] : E[u]) if(v \neq f){
26
                G[v][0] = w;
                F[v][0] = u;
27
28
                dfs(v, u);
29
            }
30
31
        void build(){
32
            for(int i = 1; i \leq n; ++ i)
33
                DSU :: F[i] = i;
34
            for(int i = 1; i \leq m; ++ i)
                I[i] = i;
35
            sort(I + 1, I + 1 + m, cmp);
36
37
            for(int i = 1; i \leq m; ++ i){
                int a = A[I[i]];
38
                int b = B[I[i]];
39
40
                int w = W[I[i]];
                int fa = DSU :: getfa(a);
41
42
                int fb = DSU :: getfa(b);
                if(fa \neq fb){
43
                     DSU :: F[fa] = fb;
44
                     E[a].push_back({b, w});
45
46
                     E[b].push_back({a, w});
47
            }
48
49
            dfs(1, 0);
50
        int solve(int u, int v){
51
            if(u = outer || v = outer)
52
53
                return -1;
54
            int ans = 0;
            if(D[u] < D[v]) swap(u, v);
55
            for(int i = h; i \ge 0; -- i)
56
                if(D[F[u][i]) \ge D[v]){
57
58
                     ans = max(ans, G[u][i]);
                     u = F[u][i];
59
60
                 }
            if(u = v) return ans;
61
            for(int i = h; i \ge 0; -- i)
62
                if(F[u][i] \neq F[v][i]){
63
64
                     ans = max(ans, G[u][i]);
                     ans = max(ans, G[v][i]);
65
                     u = F[u][i];
66
                     v = F[v][i];
67
68
69
            ans = max(ans, G[u][0]);
70
            ans = \max(ans, G[v][0]);
71
            return ans;
        }
72
73
74
   namespace Planer{
75
        const int MAXN = 1e5 + 3 + 3;
        const int MAXE = 2e5 + 3;
76
        const int MAXG = 1e5 + 3;
77
```

```
const int MAXQ = 2e5 + 3;
 78
 79
         point P[MAXN];
         using edge = tuple<int, int>;
 80
 81
         double gety(int a, int b, double x){
             return P[a].y + (x - P[a].x) / (P[b].x - P[a].x) * (P[b].y - P[a].x)
 82
                a].y);
         }
 83
 84
         double scanx;
 85
         struct Cmp1{
             bool operator ()(const pair<edge, int> l1, const pair<edge, int</pre>
 86
                 > 12) const{
                 const edge &e1 = l1.first;
 87
 88
                 const edge &e2 = l2.first;
                 double h1 = gety(get<0>(e1), get<1>(e1), scanx);
 89
                 double h2 = gety(get<0>(e2), get<1>(e2), scanx);
 90
 91
                 return h1 < h2;</pre>
 92
             };
 93
         };
 94
         struct Cmp2{
             bool operator ()(const pair<edge, int> l1, const pair<edge, int</pre>
 95
                 > l2) const{
 96
                 if(l1.second = l2.second)
 97
                      return false;
                 const edge &e1 = l1.first;
 98
                 const edge &e2 = l2.first;
 99
                 vec v1 = P[get<1>(e1)] - P[get<0>(e1)];
100
                 vec v2 = P[get<1>(e2)] - P[get<0>(e2)];
101
                 if(sign(v1.y) \neq sign(v2.y)){
102
                      return v1.y > 0;
103
                 } else {
104
                      return sign(mulx(v1, v2)) = 1;
105
106
107
             };
         };
108
         vector <pair<edge, int> > E[MAXN];
109
110
         vector <int> G[MAXG];
         int L[MAXE], R[MAXE], W[MAXE], n, m, q, o;
111
112
         double theta;
         int outer;
113
114
         void rotate(){
             srand(time(0));
115
             theta = PI * rand() / RAND_MAX;
116
117
         int add(double x, double y){
118
119
             srand(time(0));
             P[++ n] = rotate(vec(x, y), theta);
120
             return n;
121
122
         int link(int u, int v, int w){
123
124
             ++ m;
             E[u].push_back(\{\{u, v\}, ++ o\});
125
             L[o] = u, R[o] = v, W[o] = w;
126
             E[v].push_back(\{\{v, u\}, ++ o\});
127
128
             L[o] = v, R[o] = u, W[o] = w;
129
             return m;
130
         int I[MAXE];
131
```

```
int polys;
132
         pair<edge, int> findleft(int l, int r){
133
             auto it = lower_bound(E[r].begin(), E[r].end(), make_pair(edge())
134
                 r, l), 0), Cmp2());
             if(it = E[r].begin())
135
136
                 return E[r].back();
             else
137
138
                 return *(it - 1);
         }
139
         void leftmost(){
140
141
             for(int i = 1; i \leq n; ++ i){
                  sort(E[i].begin(), E[i].end(), Cmp2());
142
143
             for(int p = 1; p \leq n; ++ p){
144
                 for(auto &[e1, id1] : E[p]){
145
                      auto \delta[x, y] = e1;
146
147
                      if(!I[id1]){
                          int l = x;
148
149
                          int r = y;
150
                          I[id1] = ++ polys;
                          G[polys].push_back(id1);
151
152
                          while (r \neq p)
                               auto [e2, id2] = findleft(l, r);
153
                              auto [a, b] = e2;
154
                               I[id2] = polys;
155
156
                              G[polys].push_back(id2);
157
                              l = r;
                               r = b;
158
159
                          }
                      }
160
                 }
161
162
             for(int i = 1;i ≤ polys; ++ i){
163
164
                 double area = 0;
                 for(int j = 0; j < G[i].size(); ++ j){</pre>
165
                      area += mulx(P[L[G[i][j]]], P[R[G[i][j]]]);
166
167
168
                 if(area < 0)
                      outer = i;
169
170
             }
171
         void dual(){
172
173
             Dual :: n = polys;
             Dual :: m = 0;
174
             for(int i = 1; i \leq m; ++ i){
175
                 int u = I[2 * i - 1], v = I[2 * i], w = W[2 * i];
176
                 if(u = outer || v = outer)
177
                      w = 1e9L + 1;
178
                 ++ Dual :: m;
179
180
                 Dual :: A[Dual :: m] = u;
                 Dual :: B[Dual :: m] = v;
181
                 Dual :: W[Dual :: m] = w;
182
183
184
             Dual :: build();
185
             Dual :: outer = outer;
186
         set <pair<edge, int>, Cmp1> S;
187
```

```
vector <pair<double, int> > T;
188
         vector <pair<double, int> > Q;
189
190
         double X[MAXQ], Y[MAXQ];
191
         int
                Z[MAXQ];
         int ask(double x, double y){
192
193
             point p = rotate(vec(x, y), theta);
194
195
             X[q] = p.x;
196
             Y[q] = p.y;
197
             return q;
198
         void locate(){
199
200
             T.clear(), Q.clear(), S.clear();
201
             for(int i = 1; i \leq q; ++ i){
202
                 Q.push_back(make_pair(X[i], i));
203
204
             for(int i = 1;i ≤ polys; ++ i){
205
                 for(auto &e : G[i]){
                      int u = L[e];
206
207
                      int v = R[e];
                      if(P[u].x > P[v].x){
208
209
                          T.push_back(make_pair(P[v].x + 1e-5, e));
                          T.push back(make pair(P[u].x - 1e-5, -e));
210
                      }
211
                 }
212
213
             sort(T.begin(), T.end());
214
             sort(Q.begin(), Q.end());
215
216
             int p1 = 0, p2 = 0;
             scanx = -1e9;
217
             Cmp1 CMP;
218
             while(p1 < Q.size() || p2 < T.size()){</pre>
219
                 // for(auto it1 = S.begin(), it2 = next(S.begin()); it2 \neq
220
                     S.end(); ++ it1, ++ it2)
                         assert(CMP(*it1, *it2));
221
                 double x1 = p1 < Q.size() ? Q[p1].first : 1e9;</pre>
222
                 double x2 = p2 < T.size() ? T[p2].first : 1e9;</pre>
223
224
                 scanx = min(x1, x2);
                 if(equal(scanx, x1)){
225
226
                      auto &x = X[Q[p1].second];
227
                      auto &y = Y[Q[p1].second];
228
                      auto &z = Z[Q[p1].second];
229
                      P[n + 1] = point(-1e9, y);
                      P[n + 2] = point(1e9, y);
230
                      auto it = S.lower_bound(\{\{n + 1, n + 2\}, 0\});
231
232
                      if(it = S.end())
233
                          z = outer;
234
                      else
235
                          z = it \rightarrow second;
236
                      ++ p1;
237
                 if(equal(scanx, x2)){
238
                      int g = T[p2].second;
239
240
                      if(g > 0){
                          assert(!S.count({{L[g], R[g]}, I[g]}));
241
242
                          S.insert(\{\{L[g], R[g]\}, I[g]\});
243
                      } else {
```

```
244
                          g = -g;
                          assert( S.count({{L[g], R[g]}, I[g]}));
245
                          S.erase ({{L[g], R[g]}, I[g]});
246
247
248
                      ++ p2;
                 }
249
             }
250
         }
251
252
253
    const int MAXN = 1e5 + 3;
254
    int A[MAXN], B[MAXN];
255
    int main(){
256
    #ifndef ONLINE_JUDGE
         freopen("test.in", "r", stdin);
257
         freopen("test.out", "w", stdout);
258
259
    #endif
         int n, m, q;
260
261
         Planer :: rotate();
262
         cin >> n >> m;
263
         for(int i = 1; i \leq n; ++ i){
             double x, y;
264
265
             cin \gg x \gg y;
266
             Planer :: add(x, y);
267
268
         for(int i = 1; i \leq m; ++ i){
269
             int u, v, w;
270
             cin >> u >> v >> w;
271
             Planer :: link(u, v, w);
272
273
         Planer :: leftmost();
274
         Planer :: dual();
275
         cin >> q;
276
         for(int i = 1; i \leq q; ++ i){
277
             double a1, b1, a2, b2;
278
             cin >> a1 >> b1;
             A[i] = Planer :: ask(a1, b1);
279
280
             cin \gg a2 \gg b2;
281
             B[i] = Planer :: ask(a2, b2);
282
283
         Planer :: locate();
         for(int i = 1; i \leq q; ++ i)
284
             A[i] = Planer :: Z[A[i]],
285
286
             B[i] = Planer :: Z[B[i]];
         for(int i = 1; i \leq q; ++ i){
287
288
             int ans = Dual :: solve(A[i], B[i]);
             cout << ans << endl;</pre>
289
290
291
         return 0;
292
```

### 9.4 二维基础

```
#include <bits/stdc++.h>
using namespace std;
using i64 = long long;
const int INF = 1e9;
const i64 INFL = 1e18;
```

```
6 int gread();
7
   const double EPS = 1e-9;
   const double PI = acos(-1);
9
   bool equal(double a, double b){
       return fabs(a - b) < EPS;</pre>
10
11
   int sign(double a){
12
13
       if(equal(a, 0))
14
           return 0;
15
       return a > 0 ? 1 : -1;
16
   double sqr(double x){
17
18
       return x * x;
19
   struct vec{ // 二维向量
20
21
       double x;
22
       double y;
23
       vec(){}
       vec(double _x, double _y) : x(_x), y(_y){}
24
25
   };
   vec operator +(const vec &a, const vec &b){
26
27
       return vec(a.x + b.x, a.y + b.y);
28
29
   vec operator -(const vec &a, const vec &b){
30
       return vec(a.x - b.x, a.y - b.y);
31
32
   double mulp(const vec &a, const vec &b){
33
       return a.x * b.x + a.y * b.y;
34
35
   double mulx(const vec &a, const vec &b){
36
       return a.x * b.y - a.y * b.x;
37
38
   vec mul(const double &r, const vec &a){
39
       return vec(r * a.x, r * a.y);
40
   bool equal(vec a, vec b){
41
42
       return equal(a.x, b.x) & equal(a.y, b.y);
43
   using point = vec;
44
45
   point rotate(point a, double t){
46
       double c = cos(t);
47
       double s = sin(t);
48
       return point(a.x * c - a.y * s, a.y * c + a.x * s);
49
50
   bool cmpx(point a, point b){
       return sign(a.x - b.x) = -1;
51
52
53
   bool cmpy(point a, point b){
54
       return sign(a.y - b.y) = -1;
55
   struct line{
                    // 有向直线
56
57
       point o;
58
       vec p;
59
       line(point _{o}, vec _{p}) : o(_{o}), p(_{p}){}
60
  };
61
   struct segm{
                   // 有向线段
62
       point a, b;
```

```
segm(point _a, point _b) : a(_a), b(_b){}
63
64
    };
65
    int side(line l, point p){
        return sign(mulx(l.p, p - l.o));
66
67
68
    int side(segm s, point p){
69
        return sign(mulx(s.b - s.a, p - s.a));
70
    bool parallel(line a, line b){
71
72
        return equal(0, mulx(a.p, b.p));
73
    double abs(vec a){
74
75
        return sqrt(a.x * a.x + a.y * a.y);
76
    double dis(point a, point b){
77
        return sqrt(sqr(a.x - b.x) + sqr(a.y - b.y));
78
79
80
    double abs(segm s){
        return dis(s.a, s.b);
81
82
    double dis(line a, point p){
83
84
        return abs(mulx(p - a.o, a.p)) / abs(a.p);
85
86
    point intersection(line a, line b){
87
        return b.o + mul(mulx(b.o - a.o, a.p) / mulx(a.p, b.p), b.p);
88
    bool intersect(double l1, double r1, double l2, double r2){
89
90
        if(l1 > r1) swap(l1, r1);
91
        if(l2 > r2) swap(l2, r2);
        if(equal(r1, l2) || equal(r2, l1))
92
93
             return true;
        return !equal(max(r1, r2) - min(l1, l2), r1 - l1 + r2 - l2);
94
95
    bool intersect(segm s1, segm s2){
96
97
        bool fx = intersect(s1.a.x, s1.b.x, s2.a.x, s2.b.x);
98
        if(!fx) return false;
        bool fy = intersect(s1.a.y, s1.b.y, s2.a.y, s2.b.y);
99
100
        if(!fy) return false;
        bool g1 = side(s1, s2.a) * side(s1, s2.b) = 1;
101
        if(g1) return false;
102
        bool g2 = side(s2, s1.a) * side(s2, s1.b) = 1;
103
        if(g2) return false;
104
105
        return true;
106
    struct circ{ // 二维圆形
107
108
        point o;
109
        double r;
110
    struct poly{ // 二维多边形
111
112
        vector <point> P;
113
    double area(point a, point b, point c){
114
115
        return abs(mulx(b - a, c - a)) / 2;
116
    double area(const poly &P){
117
        double ans = 0;
118
        for(int i = 0;i < P.P.size();++ i){</pre>
119
```

# 10 其他

## 10.1 笛卡尔树

```
1 #include <bits/stdc++.h>
 2 using namespace std;
 3 using i64 = long long;
 4 const int INF = 1e9;
 5 const i64 INFL = 1e18;
 6 | int gread();
   const int MAXN = 1e7 + 3;
   int n, L[MAXN], R[MAXN], A[MAXN];
   void build(){
9
10
        stack <int> S;
       A[n + 1] = -1e9;
11
        for(int i = 1; i \le n + 1; ++ i){
12
13
            int v = 0;
14
            while(!S.empty() & A[S.top()] > A[i]){
15
                auto u = S.top();
                R[u] = v;
16
17
                v = u;
18
                S.pop();
19
20
            L[i] = v;
            S.push(i);
21
        }
22
23
   int main(){
24
25
        n = qread();
26
        for(int i = 1; i \leq n; ++ i)
            A[i] = qread();
27
28
        build();
        long long ans1 = 0, ans2 = 0;
29
        for(int i = 1;i ≤ n;++ i){
30
            // cout << L[i] << " " << R[i] << endl;
31
            ans1 ^{-1} 1ll * i * (L[i] + 1);
32
            ans2 ^{-1}1 | * i * (R[i] + 1);
33
34
        cout << ans1 << " " << ans2 << endl;</pre>
35
36
        return 0;
37
   }
```

## 10.2 CDQ 分治

#### 10.2.1 例题

给定三元组序列  $(a_i,b_i,c_i)$ , 求解  $f(i) = \sum_i [a_i \le a_i \land b_j \le b_i \land c_j \le c_i]$ 。

```
#include<bits/stdc++.h>
 1
   #define up(l, r, i) for(int i = l, END##i = r; i \leq END##i; ++ i)
   #define dn(r, l, i) for(int i = r, END##i = l; i \ge END##i; -- i)
 4 using namespace std;
   typedef long long i64;
 5
   const int INF = 2147483647;
 6
   const int MAXN= 1e5 + 3;
7
8
   const int MAXM= 2e5 + 3;
   struct Node{
9
10
        int id, a, b, c;
   }A[MAXN], B[MAXN];
11
   bool cmp(Node a, Node b){
12
        if(a.a \neq b.a) return a.a < b.a;
13
14
        if(a.b \neq b.b) return a.b < b.b;
        if(a.c ≠ b.c) return a.c < b.c;
15
        return a.id < b.id;</pre>
16
17
   int K[MAXN], H[MAXN];
18
19
   int gread();
   int n, m, D[MAXM];
   namespace BIT{
21
22
        void increase(int x, int w){
23
            while(x \leq m) D[x] += w, x += x & -x;
24
25
        void decrease(int x, int w){
26
            while(x \leq m) D[x] -= w, x += x & -x;
27
28
        void query(int x, int &r){
29
            while(x) r += D[x], x -= x & -x;
30
        }
31
   void cdq(int l, int r){
32
33
        if(l \neq r){
            int t = l + r >> 1; cdq(l, t), cdq(t + 1, r);
34
35
            int p = l, q = t + 1, u = l;
            while (p \le t \& q \le r)
36
                if(A[p].b \leq A[q].b)
37
                    BIT :: increase(A[p].c, 1), B[u ++] = A[p ++];
38
39
                 else
                    BIT :: query(A[q].c, K[A[q].id]), B[u ++] = A[q ++];
40
41
            while(p \le t) BIT :: increase(A[p].c, 1),
42
                                                               B[u \leftrightarrow] = A[p]
               ++];
            while(q \leq r) BIT :: query(A[q].c, K[A[q].id]), B[u \leftrightarrow] = A[q
43
               ++];
44
            up(l, t, i) BIT :: decrease(A[i].c, 1);
45
            up(l, r, i) A[i] = B[i];
        }
46
47
   int main(){
48
        n = qread(), m = qread();
49
        up(1, n, i) A[i].id = i, A[i].a = qread(), A[i].b = qread(), A[i].c
50
            = qread();
51
        sort(A + 1, A + 1 + n, cmp), cdq(1, n);
52
        sort(A + 1, A + 1 + n, cmp);
        dn(n, 1, i){
53
            if(A[i].a = A[i + 1].a \& A[i].b = A[i + 1].b \& A[i].c = A[i]
54
```

```
i + 1].c)
K[A[i].id] = K[A[i + 1].id];
H[K[A[i].id]] ++;

up(0, n - 1, i) printf("%d\n", H[i]);
return 0;
}
```

## 10.3 自适应辛普森

### 10.3.1 例题

计算

$$\int_0^{+\infty} x^{(a/x)-x}$$

```
1 #include<bits/stdc++.h>
2 using namespace std;
3 using i64 = long long;
4 const int INF = 1e9;
5 const i64 INFL = 1e18;
6 double simpson(double (*f)(double), double l, double r){
       double mid = (l + r) / 2;
7
       return (r - l) * (f(l) + 4 * f(mid) + f(r)) / 6.0;
8
9
   double adapt_simpson(double (*f)(double), double l, double r, double
10
      EPS, int step){
       double mid = (l + r) / 2;
11
       double w0 = simpson(f, l, r);
12
       double w1 = simpson(f, l, mid);
13
       double w2 = simpson(f, mid, r);
14
       if(fabs(w0 - w1 - w2) < EPS & step < 0)
15
            return w1 + w2;
16
       else
17
            return adapt_simpson(f, l, mid, EPS, step - 1) +
18
                   adapt_simpson(f, mid, r, EPS, step - 1);
19
20
21
   double a, l, r;
   double fun(double x){
23
       return pow(x, a / x - x);
24
   int main(){
25
       cin >> a;
26
27
       if(a < 0)
            cout << "orz" << endl;</pre>
28
       else {
29
30
           l = 1e-9;
            r = 150:
31
32
            cout << fixed << setprecision(5) << adapt_simpson(fun, l, r, 1e</pre>
               -9, 15);
       }
33
34
```

#### 10.4 模拟退火

#### 10.4.1 例题

给定 n 个物品挂在洞下,第 i 个物品坐标  $(x_i, y_i)$  重量为  $w_i$ 。询问平衡点。

```
#include<bits/stdc++.h>
 2 using namespace std;
 3 const double T0 = 2e3, Tk = 1e-14, delta = 0.993, R = 1e-3;
 4 mt19937 MT(114514);
   double distance(double x, double y, double a, double b){
 5
       return sqrt(pow(a - x, 2) + pow(b - y, 2));
 6
7
8
   const int MAXN = 1e3 + 3;
   double X[MAXN], Y[MAXN], W[MAXN]; int n;
   double calculate(double x, double y){
10
       double gx, gy, a;
11
12
       for(int i = 0; i < n; ++i){
            a = atan2(y - Y[i], x - X[i]);
13
14
            gx += cos(a) * W[i];
15
            gy += sin(a) * W[i];
16
       return pow(gx, 2) + pow(gy, 2);
17
18
19
   double ex, ey, eans = 1e18;
20
   void SA(){
21
       double T = T0, x = 0, y = 0, ans = calculate(x, y);
       double ansx, ansy;
22
       uniform_real_distribution<double> U;
23
24
       while(T > Tk){
            double nx, ny, nans;
25
26
            nx = x + 2 * (U(MT) - .5) * T;
            ny = y + 2 * (U(MT) - .5) * T;
27
28
            if((nans = calculate(nx, ny)) < ans){</pre>
29
                ans = nans;
30
                ansx = x = nx;
31
                ansy = y = ny;
            } else if(exp(-distance(nx, ny, x, y) / T / R) > U(MT)){
32
33
                x = nx, y = ny;
34
35
           T *= delta;
36
        if(ans < eans) eans = ans, ex = ansx, ey = ansy;
37
38
   int main(){
39
40
       cin >> n;
        for(int i = 0; i < n; ++ i)
41
            cin \gg X[i] \gg Y[i] \gg W[i];
42
43
        cout << fixed << setprecision(3);</pre>
44
        if(n = 1){
            cout << X[0] << " " << Y[0] << endl;
45
        } else {
46
47
            SA(), SA(), SA();
            cout << ex << " " << ey << endl;
48
49
50
       return 0;
51
```

# 10.5 伪随机生成

```
1 #include<bits/stdc++.h>
2 using namespace std;
3 using u32 = uint32_t;
4 using u64 = uint64_t;
5 u32 xorshift32(u32 &x){
6
       x ^ x < 33;
7
       x ^ x > 17;
       x ^= x << 5;
8
9
       return x;
10 }
11 u64 xorshift64(u64 &x){
12
       x ^  x << 13;
       x ^  x > 7;
13
       x ^  x < 17;
14
15
       return x;
16 }
```