

Pallas

HPC Trace Analysis at scale

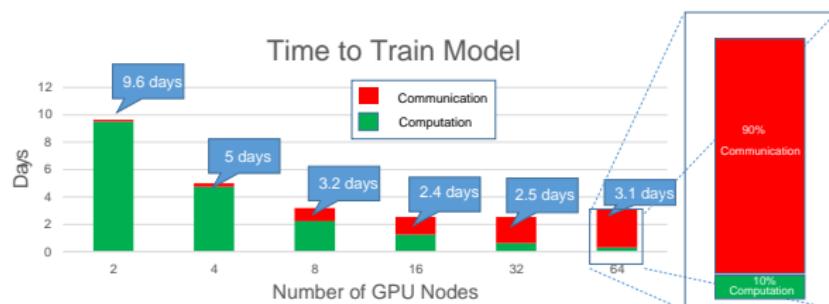
Catherine Guelque
Supervisors: Francois Trahay & Valentin Honoré

Télécom SudParis



Scalability issues in HPC

- Load-balancing
- Concurrent access to resources
- Interactions between threads
- Non-negligable communication times



How to scale, debug and optimize applications on such systems ?

Tracing

Trace → timeline of the execution

Idea: Intercept function calls (MPI, OMP, CUDA) → Event

- **Timestamp**
- **Additional Data:** Arguments, callstack, etc.



Figure: An OTF2 Trace visualised with Vampir.

Trace analysis

Visualisation aren't scalable (too many threads, memory issue)
Statistical analysis need to be done quick and cheap



Figure: A sizable trace in ViTE

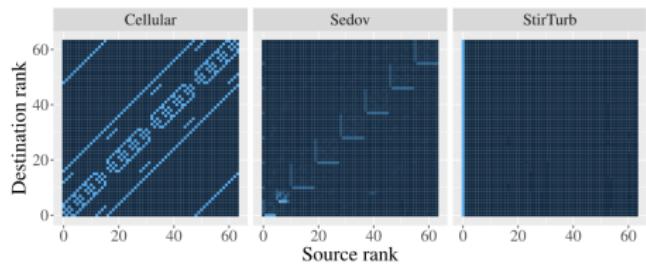


Figure: Communication matrices generated by Pilgrim.

Types of traces

Sequential

Array of events in chronological order

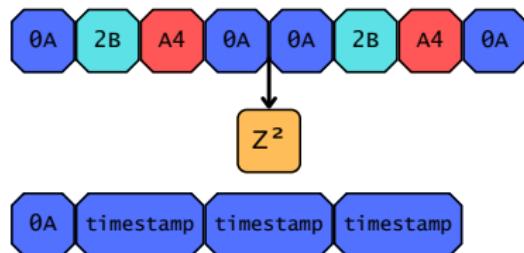
- Straightforward to read & write
- Redundancy → heavy traces



Structural

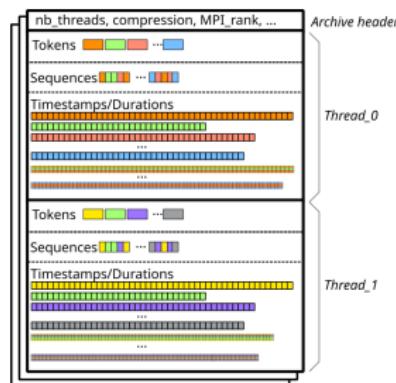
HPC apps are predictable → include the structure of the program

- Better compression
- More information
- Easier analysis



Pallas

- Structural trace format
- Efficient data retrieval
 - Independent data and metadata loading
 - Parallelisable read and write
- Efficient encoding and compression



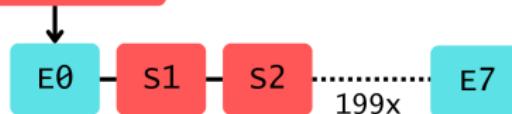
Appendix

General idea

- Events are intercepted and stored as Tokens
- Intercepted function calls create **Blocks** = Hierarchy

```
void foo() {  
    MPI_Send(...);  
    MPI_Recv(...);  
}  
  
int main() {  
    for (int i = 0; i < 200; i++) {  
        foo();  
    }  
}
```

s0 = main()



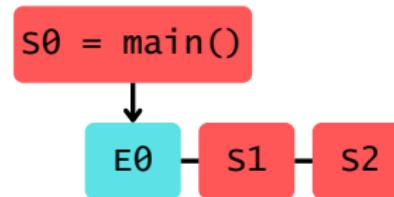
Structure detection

Each Block has its own buffer of tokens

When a token is added → Basic loop detection algorithm

If repetition of tokens is detected:

- Check already existing Sequences with hashing function
- Replace repeating Tokens with new Loop token



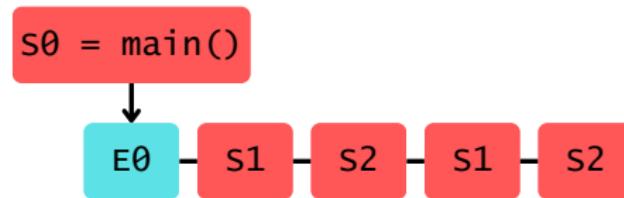
Structure detection

Each Block has its own buffer of tokens

When a token is added → Basic loop detection algorithm

If repetition of tokens is detected:

- Check already existing Sequences with hashing function
- Replace repeating Tokens with new Loop token



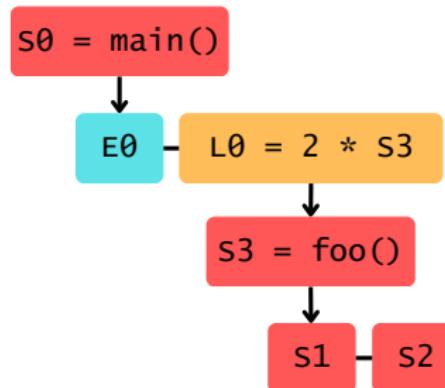
Structure detection

Each Block has its own buffer of tokens

When a token is added → Basic loop detection algorithm

If repetition of tokens is detected:

- Check already existing Sequences with hashing function
- Replace repeating Tokens with new Loop token



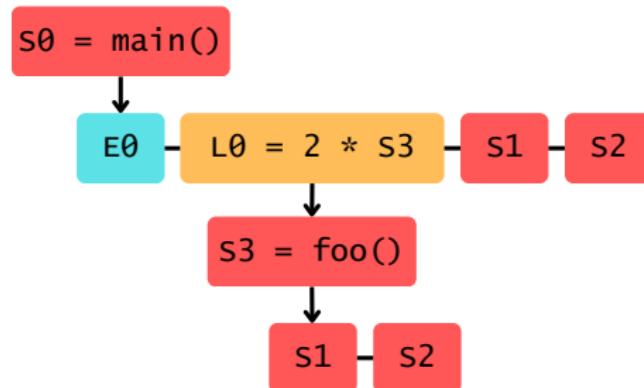
Structure detection

Each Block has its own buffer of tokens

When a token is added → Basic loop detection algorithm

If repetition of tokens is detected:

- Check already existing Sequences with hashing function
- Replace repeating Tokens with new Loop token



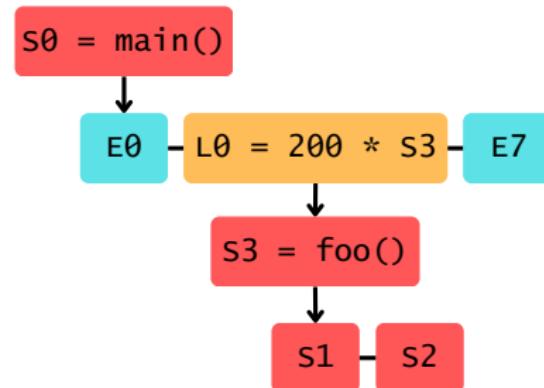
Structure detection

Each Block has its own buffer of tokens

When a token is added → Basic loop detection algorithm

If repetition of tokens is detected:

- Check already existing Sequences with hashing function
- Replace repeating Tokens with new Loop token



Reading the structure

The structure can be read independently from the timestamps → quick and lightweight overview of the program

```
void ping_pong() {
    MPI_Send(...);
    MPI_Recv(...);
}

void main() {
    int i;
    for(i = 0; i < 100; i++) {
        ping_pong();
    }
    for(i = 0; i < 10 000; i++) {
        ping_pong();
    }
    MPI_Barrier();
}
```

```
0 Reading events for thread 0 (P#0T#0):
1 Tag      Event
2 S6      E0_S L1 Eb_S
3   E0_S  THREAD_BEGIN()
4     L1  2 * S5 = L0 S4
5     L0  (100, 10_000) * S3 = S1 S2
6     S1  E1_E E2_S E3_L
7       E1_E  Enter 0 (MPI_Send)
8       E2_S  MPI_SEND(dest=1, comm=0, tag=0, len=16)
9       E3_L  Leave 0 (MPI_Send)
10    S2  E4_E E5_S E6_L
11      E4_E  Enter 1 (MPI_Recv)
12      E5_S  MPI_RECV(src=1, comm=0, tag=0, len=16)
13      E6_L  Leave 1 (MPI_Recv)
14    S4  E7_E E8_S E9_S Ea_L
15      E7_E  Enter 2 (MPI_Barrier)
16      E8_S  MPI_COLLECTIVE_BEGIN()
17      E9_S  MPI_COLLECTIVE_END(op=0, comm=0, root=-1, sent=0, recved=0)
18      Ea_L  Leave 2 (MPI_Barrier)
19 Eb_S  THREAD_END()
```

Reading the timestamps

Durations can be loaded selectively

Timestamps are calculated on the fly using the sequences' durations

```
void ping_pong() {
    MPI_Send(...);
    MPI_Recv(...);
}

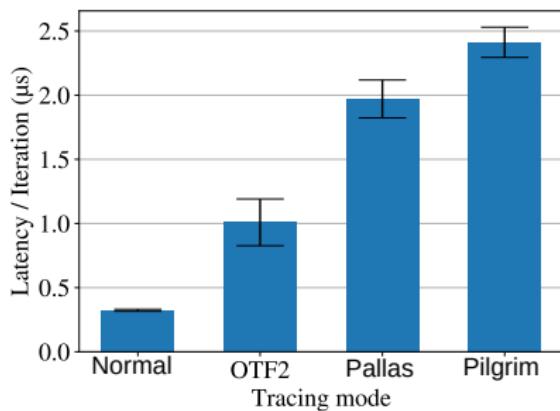
void main() {
    int i;
    for(i = 0; i < 100; i++) {
        ping_pong();
    }
    for(i = 0; i < 10 000; i++) {
        ping_pong();
    }
    MPI_Barrier();
}
```

Tag	Duration	Event
S7	2.00004882	E0_S S6 Ed_S
└─E0_S	0.00003671	THREAD_BEGIN()
└─S6	2.00000076	E1_E L1 Ec_L
└─E1_E	0.00001130	Enter 1 (Working)
└─L1	1.99998946	2 * S5 = L0 S4
└─S5	0.09816532	E8_E E9_S Ea_S Eb_L
└─S5	1.90182414	E8_E E9_S Ea_S Eb_L
└─Ec_L	0.00001135	Leave 1 (Working)
└─Ed_S	THREE_THREAD_END()	

Experimental parameters

- MPI Ping-Pong, NAS Parallel Benchmarks, AMG, MiniFE, Lulesh & Quicksilver
- 5 iterations, plotted the means with standard deviations
- Tested on
 - OTF2 using EZTrace
 - Pallas using EZTrace
 - Pilgrim

Event recording cost

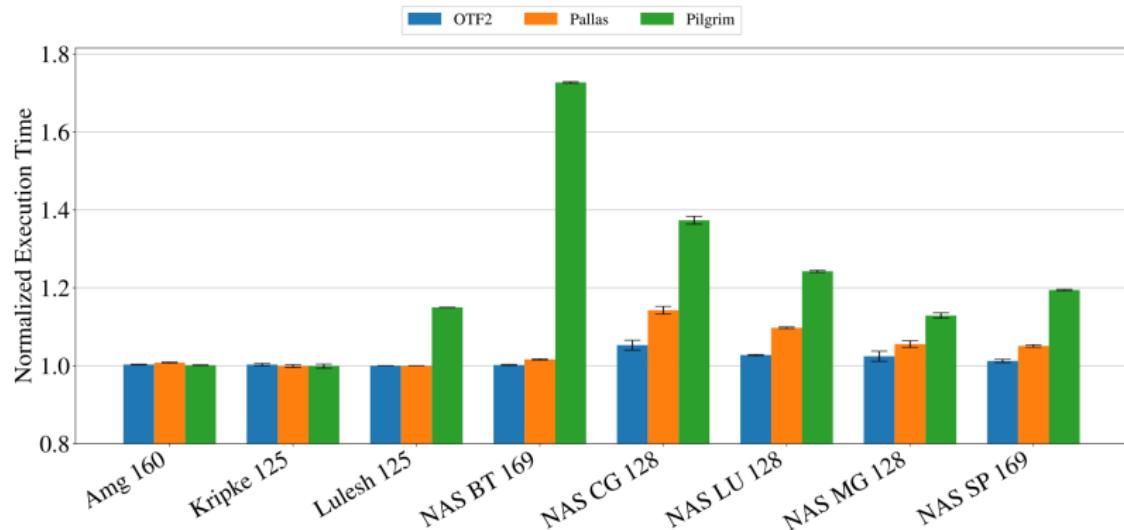


	Added Latency (ns)	Events/Iter	Latency / Event (ns)
OTF2	686	6	114
Pallas	1 647	6	275
Pilgrim	2 089	2	1 044

Figure: **Top:** Latency (in μ s) of the MPI Ping-Pong program, using OTF2, Pilgrim et Pallas. **Bottom:** Table of the latency per event.

Overhead

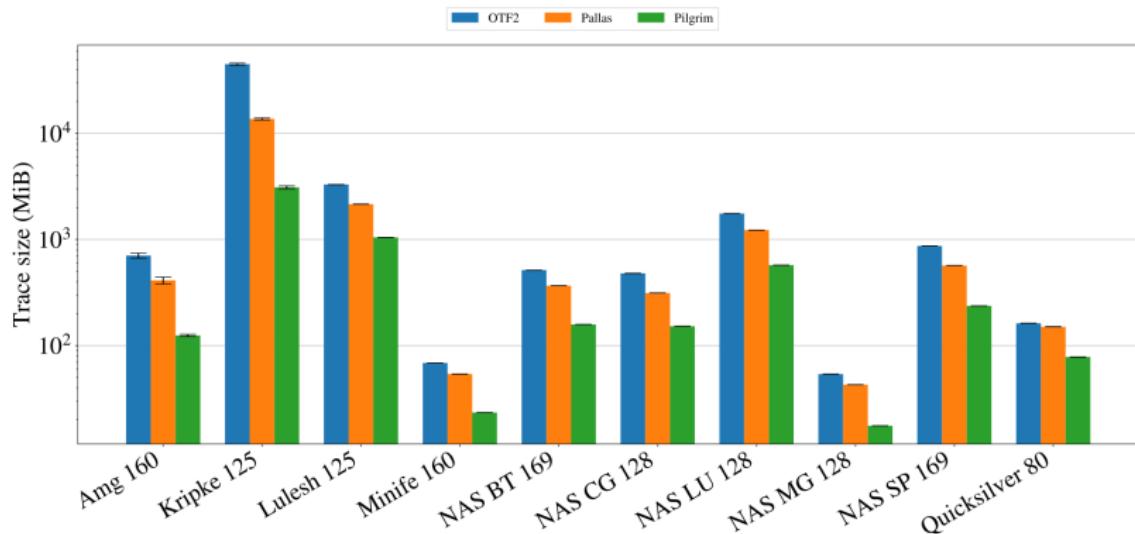
Overhead for different Kernels.



- OTF2, Pallas and Pilgrim almost have the same overhead.

Trace size

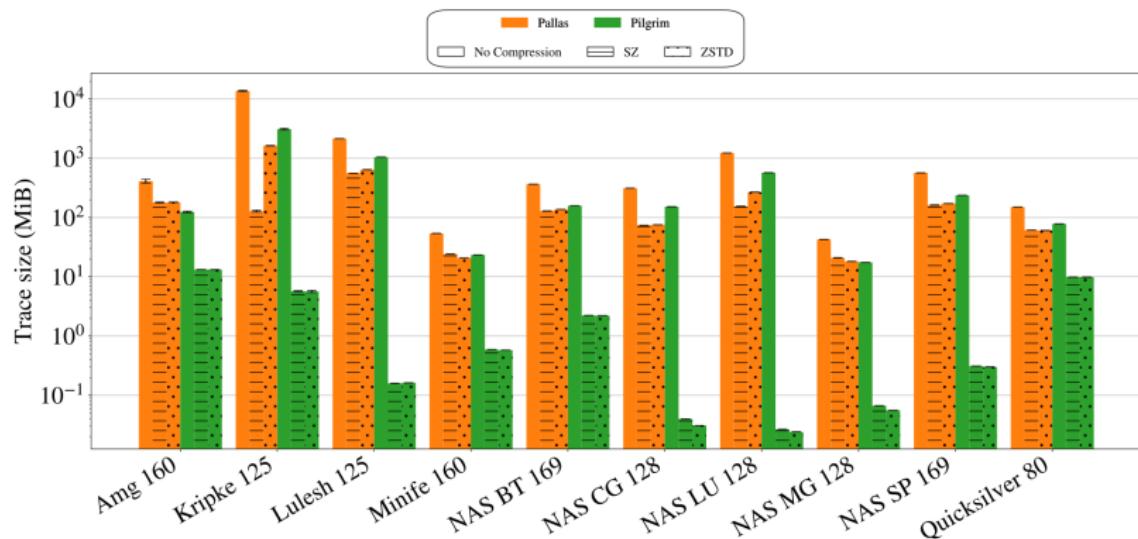
Size of traces for different kernels.



- Without compression, up to 10x difference between OTF2 and Pilgrim
- Possibility: Pilgrim collects less information than EZTrace (Strings, Thread names, etc.)

Trace size

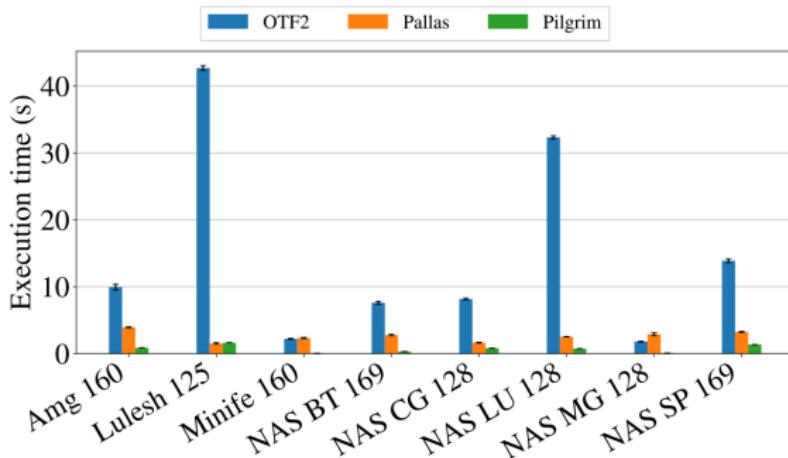
Size of traces for different kernels with different compressions.



- With compression, Pilgrim shows substantially better results **because all the timestamps are compressed together.**

Analysis speed: Communication Matrix

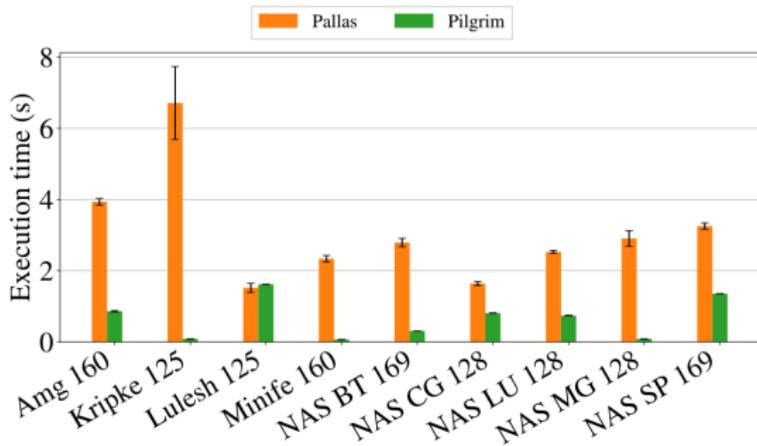
Time to plot a communication matrix from different trace formats.



- No represented: Kripke analysis took 450s for OTF2.
- OTF2 (sequential trace) analysis is slow.

Analysis speed: Communication Matrix

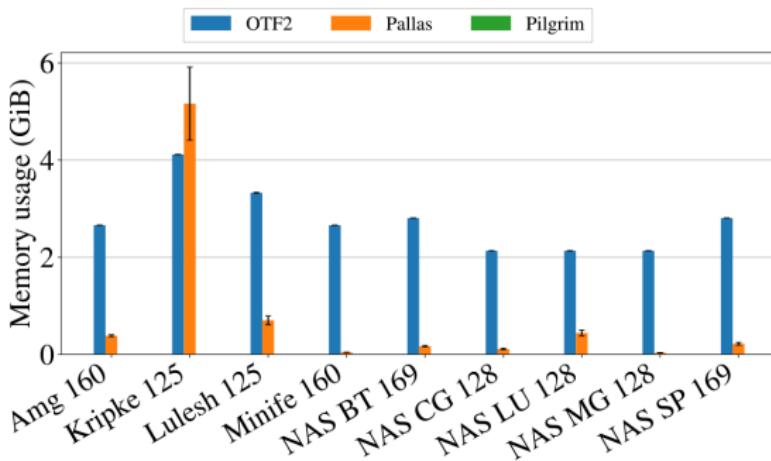
Time to plot a communication matrix from different trace formats.



- Analysis speed **of the structure** isn't quite as fast as Pilgrim yet.
- Analysis speed uncorrelated with actual trace size (MiniFE was 10-50 times lighter than Lulesh)
- We're supposed to be faster... (Work in Progress)

Memory consumption: Contention detection

Memory consumption to detect contention from different traces.



- Very low consumption compared to OTF2... Except for Kripke.
- No data for Pilgrim yet (should consume more memory).
- Some debugging of Pallas is still in order.

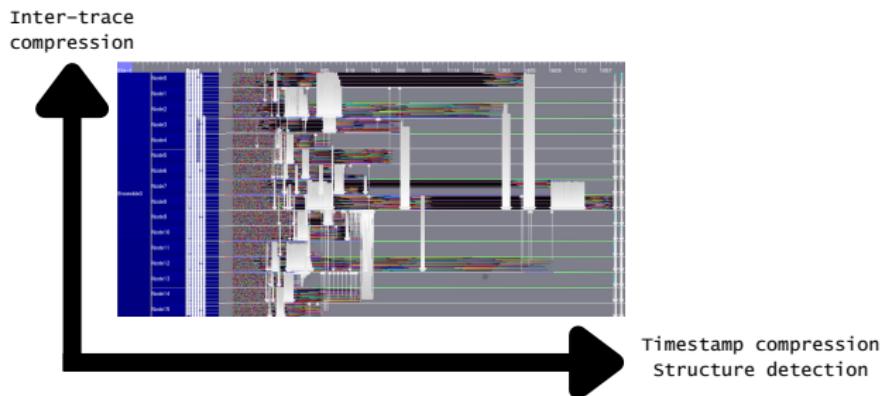
Conclusion

Pallas:

- ✓ Low Overhead
- ✓ Structure detection
- 👉 Compact structures
- ✓ Efficient timestamp storage with compression / encoding
- ✗ Efficient compression
- ✓ Basic scalable analysis
- 👉 Good performance for analysis
- ✓ On demand-trace loading and exploration

Future developments

- Inter-trace compression → "Vertical" scalability
- More efficient compression techniques
- Better trace reading performance
- Tracing non-MPI kernels
- More complex and scalable analysis



Timestamp compression & encoding

Durations are similar → easily compressible

Different storage options:

- No timestamps (Structure only)
 - Encoding:
 - Removed leading 0s
 - Replace leading 0s (as presented before)
 - Compression:
 - ZSTD
 - SZ
 - ZFP
 - Bin-based (similar to QSDG)
 - Histogram-based (same thing but Gaussian distribution)
- 

Lossy compression

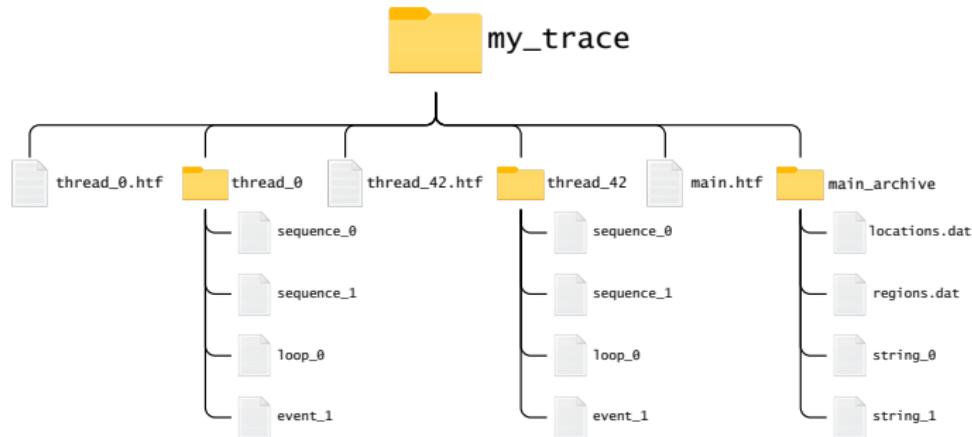
Folder & file structure

Each Event/Loop/Sequence are stored separately:

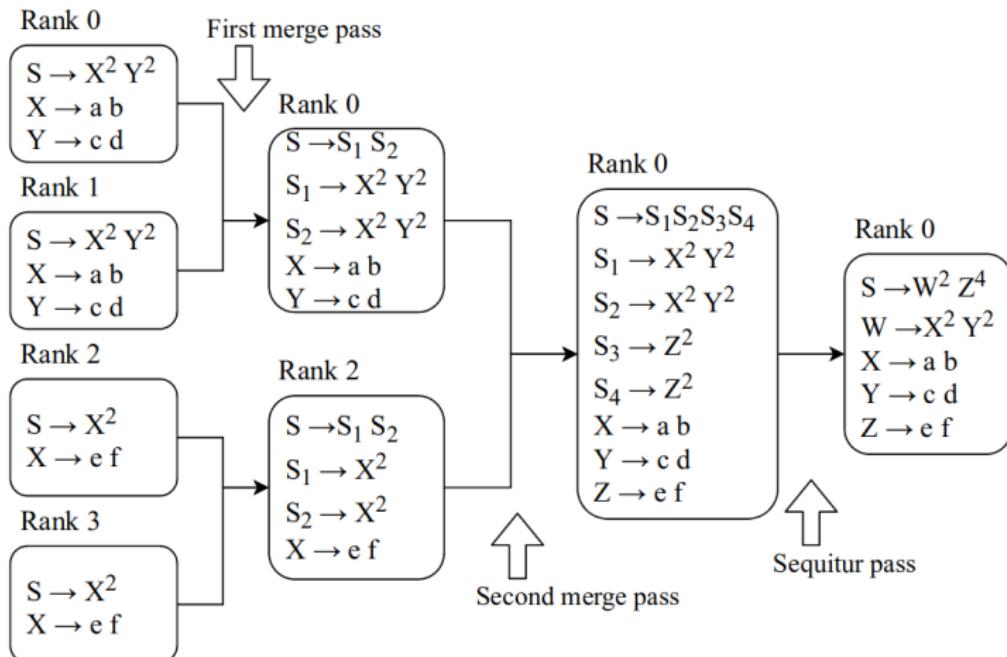
- Loops and runtime variable: stored as-is
- Sequences and Events: stored with array of **durations**

Timestamps are calculated when reading the trace

1 file & folder / thread → (almost) no concurrent writing



(Pilgrim) Inter-trace compression



Pallas_info

```

Archive ffffffe;
  trace_name: ping_pong.htrf
  Strings {.nb_strings: 26} :
    // Définitions de 26 Strings

Regions {.nb_regions: 6} :
  // Définitions de 6 Régions (Sections de code)

Location_groups {.nb_lg: 2}:
  // Définitions des 2 Location Groups (Processus)

Locations {.nb_loc: 2}:
  // Définitions des 2 Locations (Fils d'exécutions)

Threads {.nb_threads: 2}:
  0: {.archive=0, .nb_events=12, .nb_sequences=7, .nb_loops=2}
    // 2nd threads, identique au premier

Archives {.nb_archives: 2}

Thread 0 {.archive: 0}
  Events {.nb_events: 12}
    E0  THREAD_BEGIN()           {.nb_events: 1}
    E1  Enter 0 (MPI_Send)       {.nb_events: 1100}
    E2  MPI_SEND(dest=1, comm=0, tag=0, len=16) {.nb_events: 1100}
    E3  Leave 0 (MPI_Send)       {.nb_events: 1100}
    E4  Enter 1 (MPI_Recv)       {.nb_events: 1100}
    E5  MPI_RECV(src=1, comm=0, tag=0, len=16) {.nb_events: 1100}
    E6  Leave 1 (MPI_Recv)       {.nb_events: 1100}
    E7  Enter 2 (MPI_Barrier)    {.nb_events: 2}
    E8  MPI_COLLECTIVE_BEGIN()   {.nb_events: 2}
    E9  MPI_COLLECTIVE_END(op=0, comm=0, root=-1, sent=0, recv=0) {.nb_events: 2}

    Ea  Leave 2 (MPI_Barrier)    {.nb_events: 2}
    Eb  THREAD_END()            {.nb_events: 1}

Sequences {.nb_sequences: 7}
  S0  {S6}                   S1     {E1, E2, E3}
  S2  {E4, E5, E6}
  S3  {S1, S2}
  S4  {E7, E8, E9, Ea}
  S5  {S3, S4}
  S6  {E0, S5, Eb}
  L0  {.nb_loops: 2, .nb_allocated_loops: 2}
    L0  {.nb_loops: 2, .token: S3, .nb_iterations: [100, 1000]}
    L1  {.nb_loops: 1, .token: S5, .nb_iterations: [2]}

Loops {.nb_loops: 2, .nb_allocated_loops: 2}
  L0  {.nb_loops: 2, .token: S3, .nb_iterations: [100, 1000]}
  L1  {.nb_loops: 1, .token: S5, .nb_iterations: [2]}

```

Using ncurses

```
ncdu 1.11 ~ Use the arrow keys to navigate, press ? for help
-- /etc --
3,6 MiB [#####] /brltty
1,8 MiB [##] /apparmor.d
1,1 MiB [##] /ssl
636,0 KiB [#] /asciidoc
524,0 KiB [#] /xdg
496,0 KiB [#] /X11
484,0 KiB [#] /init
392,0 KiB [#] /fonts
344,0 KiB [ ] /ssh
332,0 KiB [ ] /sane.d
328,0 KiB [ ] /init.d
232,0 KiB [ ] /ImageMagick-6
220,0 KiB [ ] /dbus-1
168,0 KiB [ ] /lynx
152,0 KiB [ ] /java-8-openjdk
144,0 KiB [ ] /console-setup
140,0 KiB [ ] /default
136,0 KiB [ ] /ld.so.cache
120,0 KiB [ ] /pam.d
112,0 KiB [ ] /systemd
100,0 KiB [ ] /ppp
Total disk usage: 14,9 MiB Apparent size: 9,2 MiB Items: 3311
```

```
0 --- Reading bran_quicksilver.htm -----
1 In S6
2   100 µs [ ] E0_S      THREAD_BEGIN()
3   5.53 s [#####] L1      3 * S5 = L0 S4
4   100 µs [ ] Eb_S      THREAD_END()
5
6 =====
7 --- Reading bran_quicksilver.htm -----
8 In S6/L1
9   525 ms [##] S5      L0(1.000) S4
10  4.00 s [#####] S5      L0(8.000) S4
11  1.00 s [##] S5      L0(2.000) S4
```