

# **Utility/Privacy Trade-off through the lens of Optimal Transport**

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# Outline

Utility/Privacy Trade-off through the lens of Optimal Transport

- Online Repeated auction Example
- Privacy Regularized Policy (PRP) Model
- Sinkhorn Loss with PRP
- DC program with PRP
- Experiments

# Example

## Online repeated auctions

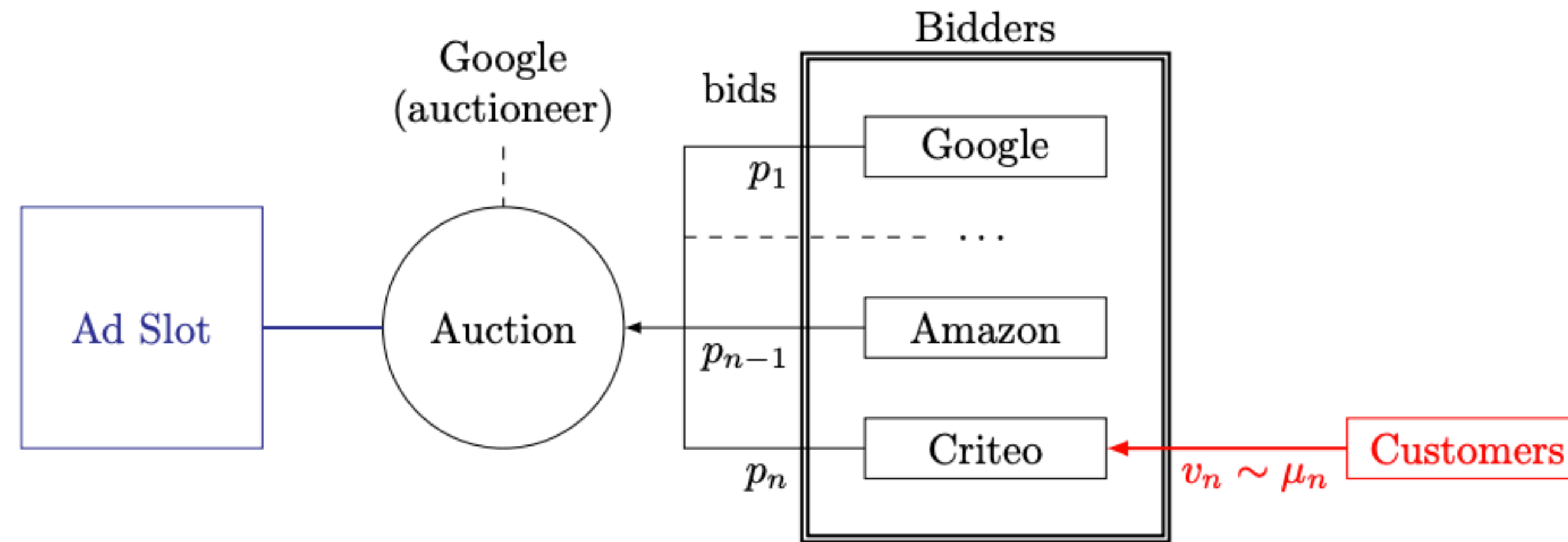


Figure 1: Online advertisement auction system.

### Scene:

auctioneer might bid the slot for its customer(so auctioneer is a direct concurrent bidder)

auctioneer is the only one who knows the true value

competitors observe a distribution of bids

If auctioneer doesn't consider the privacy, competitors can infer the true value.

# Example

## Online repeated auctions

### Suppose:

$x$  is the agent's public action,  $x \in \mathcal{X}$

type  $k$  is the hidden information,  $k \in [K]$

$c_k$  is the loss vector

### Utility Loss — without privacy concern

the **utility loss**

$$x^\top c_k$$

the optimal problem is

$$\min_{x \in \mathcal{X}} x^\top c_k$$

the **optimal solution** is

$$x_k^*$$

for each given  $k$ .

# Example

## Online repeated auctions

### Denote:

$p_0$  — the prior of type  $k$

$p_x$  — posterior distribution of the hidden type  $k$

$\mu_k$  — the agent's strategy's probability distribution

### Privacy Loss:

If the agent plays deterministically  $x_k^*$  when type is  $k$ , then the adversary could infer the true value of  $k$  based on the played action  $x$  (**Bayes rule**).

So agent should hide the true type for the long-term utility, that is to **take a strategy** to control **the amount of information** given to the adversary.

# Example

## Online repeated auctions

**Measure information loss — KL divergence between prior and posterior**

$$\text{KL} (p_x, p_0) = \sum_{k=1}^K \log \left( \frac{p_x(k)}{p_0(k)} \right) p_x(k)$$

where

$$p_x(k) = \frac{p_0(k) \mu_k(x)}{\sum_{l=1}^K p_0(l) \mu_l(x)}$$

classical cost of information in **economics**.

# Toy Example

**Total loss = utility loss + information loss:**

$$x^\top c_k + \lambda \text{KL} (p_x, p_0) \quad (\lambda > 0)$$

**Global objective:**

$$\min_{\mu_1, \dots, \mu_K} \sum_{k=1}^K p_0(k) \mathbb{E}_{x \sim \mu_k} \left[ x^\top c_k + \lambda \text{KL} (p_x, p_0) \right]$$

# Toy Example

## global objective

$$\min_{\mu_1, \dots, \mu_K} \sum_{k=1}^K p_0(k) \mathbb{E}_{x \sim \mu_k} \left[ x^\top c_k + \lambda \text{KL} (p_x, p_0) \right]$$

- if  $\lambda = 0$  **totally revealing strategy**, the best strategy is to deterministically play  $x_k^*$  given each  $k$
- if  $\lambda = \infty$  so called **non-revealing strategy** in game theory, the best strategy is to play

$$\arg \min_x x^\top c [p_0]$$

where

$$c [p_0] = \sum_{k=1}^K p_0(k) c_k$$

- if  $0 < \lambda < \infty$  **partially revealing strategy** the behavior interpolates between two extreme strategies



# General Model

## Denote:

the agent's strategy

$$\mathcal{Y} \rightarrow \mathcal{P}(\mathcal{X}) = X | Y \in \mathcal{P}(\mathcal{X})^{\mathcal{Y}}$$

**utility loss** for playing  $x$  with type  $y$

$$c(x, y)$$

**information cost**

$$c_{priv}(X, Y) = \mathbb{E}_{x \sim X} D(p_x, p_0)$$

where

$y \in \mathcal{Y}$  is the private type

$\mathcal{P}(\mathcal{X})$  is a set of distributions over  $\mathcal{X}$

$(X, Y)$  is the joint distribution of action and type

$D(p_x, p_0)$  is the measurement of information cost

## Objective:

$$\inf_{X|Y \in \mathcal{P}(\mathcal{X})^{\mathcal{Y}}} \mathbb{E}_{(x,y) \sim (X,Y)} [c(x, y)] + \lambda c_{priv}(X, Y) \quad (1)$$

# General Model

**Suppose:**

$\gamma$  is a **joint distribution** in  $\mathcal{P}(\mathcal{X} \times \mathcal{Y})$

$\pi_{1\#}\gamma$  is marginal distribution of  $X$ ,

$$\pi_{1\#}\gamma(A) = \gamma(A \times \mathcal{Y})$$

$\pi_{2\#}\gamma$  is marginal distribution of  $Y$ ,

$$\pi_{2\#}\gamma(B) = \gamma(\mathcal{X} \times B)$$

here

$$\pi_{2\#}\gamma = p_0$$

**Privacy Regularized Policy:**

$$\inf_{\gamma \in \mathcal{P}(\mathcal{X} \times \mathcal{Y}), \pi_{2\#}\gamma = p_0} \int_{\mathcal{X} \times \mathcal{Y}} \left[ c(x, y) + \lambda D(p_x, p_0) \right] d\gamma(x, y) \quad (\text{PRP})$$

# PRP's theoretical properties

## Definition 1.

$D$  is a  **$f$ -divergence** if for all distributions  $P, Q$  such that  $P$  is absolutely continuous w.r.t.  $Q$ ,

$$D(P, Q) = \int_{\mathcal{Y}} f\left(\frac{dP(y)}{dQ(y)}\right) dQ(y)$$

where  $f$  is a **convex function** defined on  $\mathbb{R}_+^*$  with  $f(1) = 0$ .

**common f-divergence:**

KL/ reverse KL /total variation distance

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## Why $f$ -divergence costs?

1. considered in **non-Bayesian** cases
2. good properties of convexity, composition and post-processing invariance

# PRP's theoretical properties

## Theorem 1.

If  $D$  is a  $f$ -divergence, PRP  $\rightarrow$  a **convex minimization problem** in  $\gamma \in \mathcal{P}(\mathcal{X} \times \mathcal{Y})$

Here,

suppose  $D$  is always a  $f$ -divergence in the remaining part

minimum can be found by classical optimization methods such as **gradient descent**

# PRP's theoretical properties

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## Analyze

$\mathcal{P}(\mathcal{X} \times \mathcal{Y})$  has generally an infinite dimension

$\mathcal{P}(\mathcal{X} \times \mathcal{Y})$  is **dimensionally finite**  $\leftarrow$  if both sets  $\mathcal{X}$  and  $\mathcal{Y}$  are discrete

# PRP theoretical properties

## Discrete type

**Suppose**  $p_0$  is a **discrete** prior of size  $K$

$$p_0 = \sum_{k=1}^K p_0^k \delta_{y_k}$$

**Define** — (here  $\mathcal{X}$  is an infinite space)

$$\mu_k(A) = \gamma \left( A \times \{y_k\} \right), \text{ for any } A \subset \mathcal{X}$$

$$\mu = \sum_{k=1}^K \mu_k = \pi_{1\#}\gamma$$

$$p^k(x) = \frac{d\mu_k(x)}{d\mu(x)}$$

**PRP** is equivalent to

$$\inf_{\substack{\mu, (p^k(\cdot))_{1 \leq k \leq K} \\ p^k \geq 0, \sum_{l=1}^K p^l(\cdot) = 1}} \int_{\mathcal{X}} \left[ p^k(x) c(x, y_k) + \lambda p_0^k f\left(\frac{p^k(x)}{p_0^k}\right) \right] d\mu(x)$$

such that  $\forall k \leq K$ ,

$$\int_{\mathcal{X}} p^k(x) d\mu(x) = p_0^k$$

# PRP theoretical properties

Discrete type - optimal solutions existence

## Theorem 2.

**If the prior is discrete of size  $K$** , for all  $\epsilon > 0$ , **(PRP) has an  $\epsilon$ -optimal solution** such that  $\pi \# \gamma = \mu$  has a finite support of at most  $K + 2$  points.

Furthermore, if  $X$  is compact and  $c(\cdot, y_k)$  is lower semicontinuous for every  $k$ , then it also holds for  $\epsilon = 0$ .



# PRP theoretical properties

Discrete type - optimal solutions existence

## Corollary 1.

In the case of a discrete prior, (PRP) is equivalent to:

$$\inf_{(\gamma, x) \in \mathbb{R}_+^{(K+2) \times K} \times \mathcal{X}^{K+2}} \sum_{i,k} \gamma_{i,k} c(x_i, y_k) + \lambda \sum_{i,k} \gamma_{i,k} D(p_{x_i}, p_0)$$

such that  $\forall k \leq K$ ,

$$\sum_i \gamma_{i,k} = p_0^k$$

where

$$\gamma_{i,k} := \gamma \left( \left\{ (x_i, y_k) \right\} \right), \text{ if } \gamma \in \left\{ (x_i, y_k) \mid 1 \leq i \leq K+2, 1 \leq k \leq K \right\}.$$

## PRP in Experiments

It is not jointly convex in  $(\gamma, x)$ .

# Sinkhorn Loss minimization

Discrete type - optimal solutions form

**Sinkhorn loss:**

$$\text{OT}_{c,\lambda}(\mu, \nu) := \min_{\gamma \in \Pi(\mu, \nu)} \int c(x, y) d\gamma(x, y) + \lambda \int \log \left( \frac{d\gamma(x, y)}{d\mu(x) d\nu(y)} \right) d\gamma(x, y) \quad (2)$$

where

$$\Pi(\mu, \nu) = \{ \gamma \in \mathcal{P}(\mathcal{X} \times \mathcal{Y}) \mid \pi_{1\#}\gamma = \mu, \pi_{2\#}\gamma = \nu \}$$

given distributions

$$(\mu, \nu) \in \mathcal{P}(\mathcal{X}) \times \mathcal{P}(\mathcal{Y})$$

here, the last part is **the regularization term** added to speed up computations.

# Sinkhorn Loss minimization

Discrete type - optimal solutions form

## Sinkhorn Algorithm

Sinkhorn algorithm has a **linear convergence rate** to compute  $\text{OT}_{c,\lambda}(\mu, \nu)$  for distributions

$$\mu = \sum_{i=1}^n \alpha_i \delta_{x_i}$$
$$\nu = \sum_{j=1}^m \beta_j \delta_{y_j}$$

the unique matrix  $\gamma$  solution of the Problem (2) has the form  **$\text{diag}(u)K \text{diag}(v)$**  in the discrete case, where

$$K_{i,j} = e^{-\frac{c(x_i, y_j)}{\lambda}}$$

it updates

$$(u, v) \leftarrow (\alpha / Kv, \beta / K^\top u)$$

for  $n$  iterations or until convergence.

# Sinkhorn Loss minimization

Discrete type - optimal solutions form

**PRP:**

$$\inf_{\mu \in \mathcal{P}(\mathcal{X})} \text{OT}_{c,\lambda}(\mu, p_0)$$

posterior probability (Bayes rule)

$$dp_x(y) = \frac{d\gamma(x, y)}{d\mu(x)}$$

where

$$D = \text{KL}$$

$$\nu = p_0$$

**with additional constraint**

$$\pi_{1\#}\gamma = \mu$$

Minimizing without this constraint is thus equivalent to minimizing the Sinkhorn loss over all action distributions  $\mu$ .

It is a new interpretation of Sinkhorn loss:

regularization term  $\rightarrow$  privacy loss

# Sinkhorn Loss minimization

Discrete type - optimal solutions form

**PRP:**

$$\inf_{\mu \in \mathcal{P}(\mathcal{X})} \text{OT}_{c,\lambda}(\mu, p_0)$$

where

$$D = \text{KL}$$

$$\nu = p_0$$

**with additional constraint**

$$\pi_{1\#}\gamma = \mu$$

When  $\mu$  and  $\nu$  are both fixed, the optimal transport plan  $\gamma^*$  remains the same.

But, here  **$\mu$  is varying**, and it is much more complex.

# Sinkhorn Loss minimization

## Discrete type - optimal solutions form

Consider discrete support, we can look for a distribution

$$\mu = \sum_{j=1}^{K+2} \alpha_j \delta_{x_j}$$

**minimization problem over tuple  $(\alpha, x)$**

$$\inf_{(\alpha, x) \in \Delta_{K+2} \times \mathcal{X}^{K+2}} \text{OT}_{c, \lambda} \left( \sum_{i=1}^{K+2} \alpha_i \delta_{x_i}, p_0 \right) \quad (3)$$

## Sinkhorn in Experiments

Given  $\alpha$ , Sinkhorn algorithm can get the unique optimal solution  $\gamma^*$ .

# Sinkhorn Loss minimization

## Discrete type - optimal solutions form

Consider discrete support, we can look for a distribution

$$\mu = \sum_{j=1}^{K+2} \alpha_j \delta_{x_j}$$

**minimization problem over tuple  $(\alpha, x)$**

$$\inf_{(\alpha, x) \in \Delta_{K+2} \times \mathcal{X}^{K+2}} \text{OT}_{c, \lambda} \left( \sum_{i=1}^{K+2} \alpha_i \delta_{x_i}, p_0 \right) \quad (3)$$

Given  $\gamma^*$ , compute  $\nabla \text{OT}_{c, \lambda}$  to get optimal  $\alpha^*$ .

# Sinkhorn Loss minimization

## discrete minimization algorithm

### Gradient computation

Computing  $\nabla \text{OT}_{c,\lambda}$  is a known difficult task!

Solution - **the dual solution of Sinkhorn loss Problem**

It is **fast** as it does not need to store all the Sinkhorn iterations in memory and backpropagate through them afterwards.

**Convergence** of Sinkhorn algorithm is guaranteed to provide **an accurate approximation of the gradient.**

[reference] *G. Peyre and M. Cuturi. Computational optimal transport. Foundations and Trends in Machine Learning, 11(5 – 6) : 355 – 607, 2019.*



# DC program with PRP

linear utility cost over hyperrectangle

**Definition:** standard DC program is of the form

$$\min_{x \in \mathcal{X}} f(x) - g(x)$$

where both  $f$  and  $g$  are convex functions.

Methods:

- **DCA** (a local minimum)

[reference] *P. Tao and L. An Convex analysis approach to DC programming : Theory, algorithms and applications .Acta mathematica vietnamica,22(1) : 289 – 355,1997.*

# DC program with PRP

linear utility cost over hyperrectangle

## Theorem 3.

If  $\mathcal{X} = \prod_{l=1}^d [a_l, b_l]$  and  $c(x, y) = x^\top y$  then (PRP) is equivalent to the following *DC* program:

$$\min_{\gamma \in \mathbb{R}_+^{(K+2) \times K}} \lambda \sum_{i,k} p_0^k h_k(\gamma_i) - \sum_{i=1}^{K+2} \left\| \sum_{k=1}^K \gamma_{i,k} \phi(y_k) \right\|_1$$

such that  $\forall k \leq K$ ,

$$\sum_{i=1}^{K+2} \gamma_{i,k} = p_0^k$$

with

$$\phi(y)^l := (b_l - a_l) y^l / 2$$

$$h_k(\gamma_i) := \left( \sum_{m=1}^K \gamma_{i,m} \right) f \left( \frac{\gamma_{i,k}}{p_0^k \sum_{m=1}^K \gamma_{i,m}} \right)$$

# Experiments

## convergence rates of usual non-convex optimization

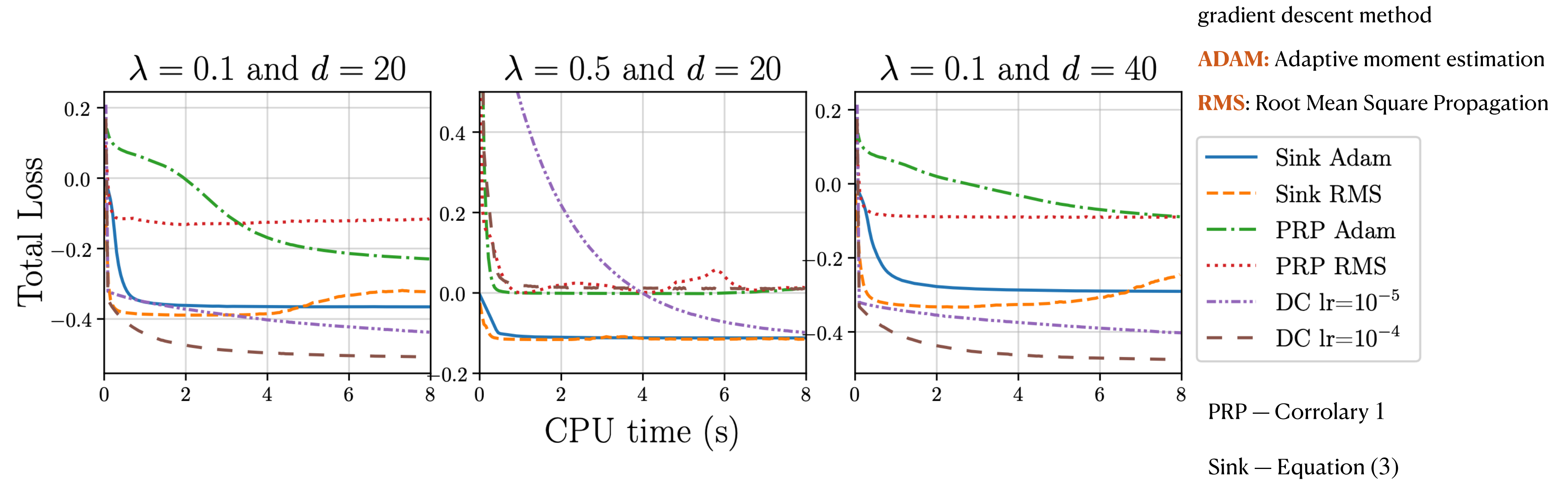
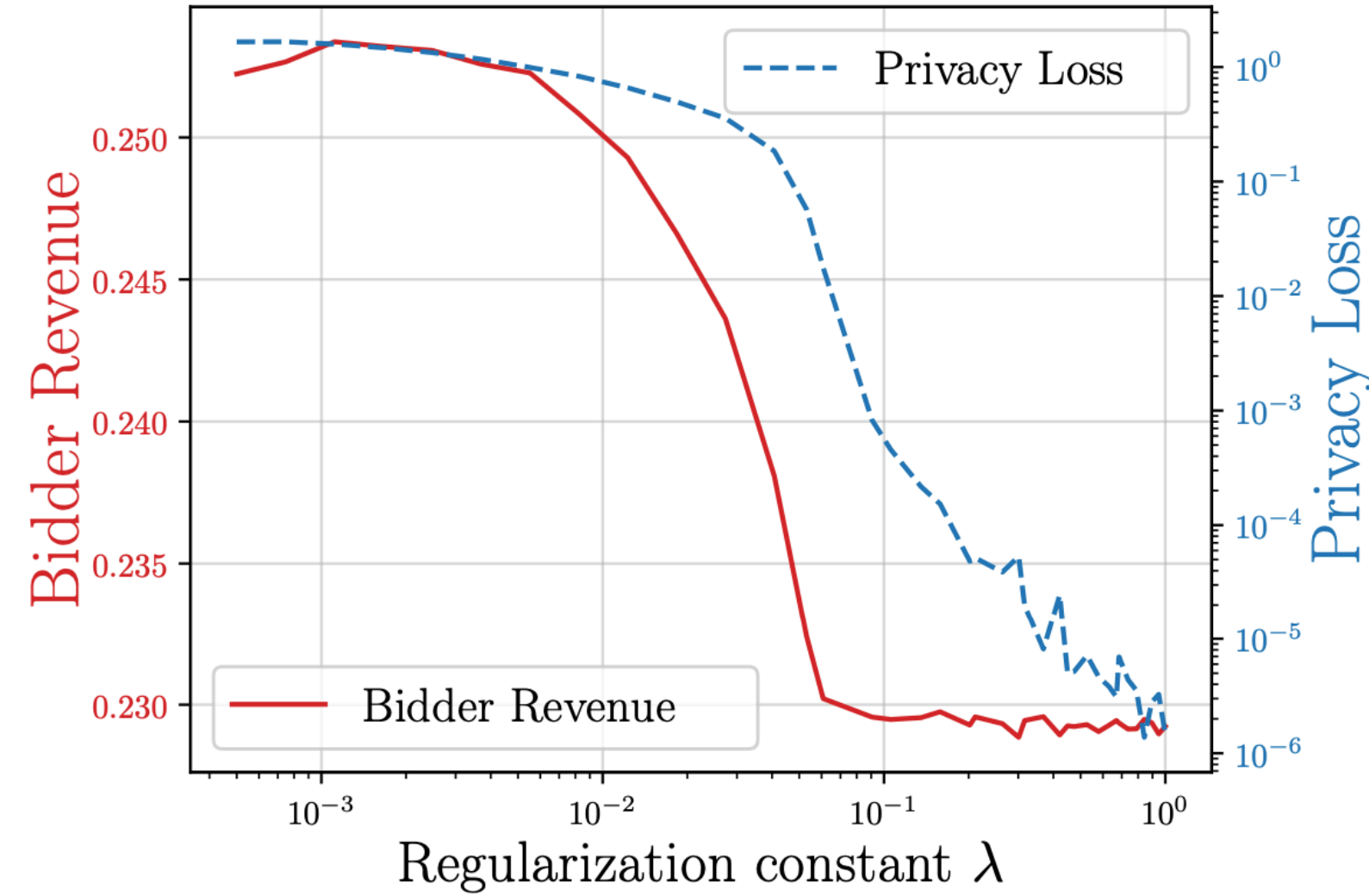


Figure 2: Comparison of optimization schemes.

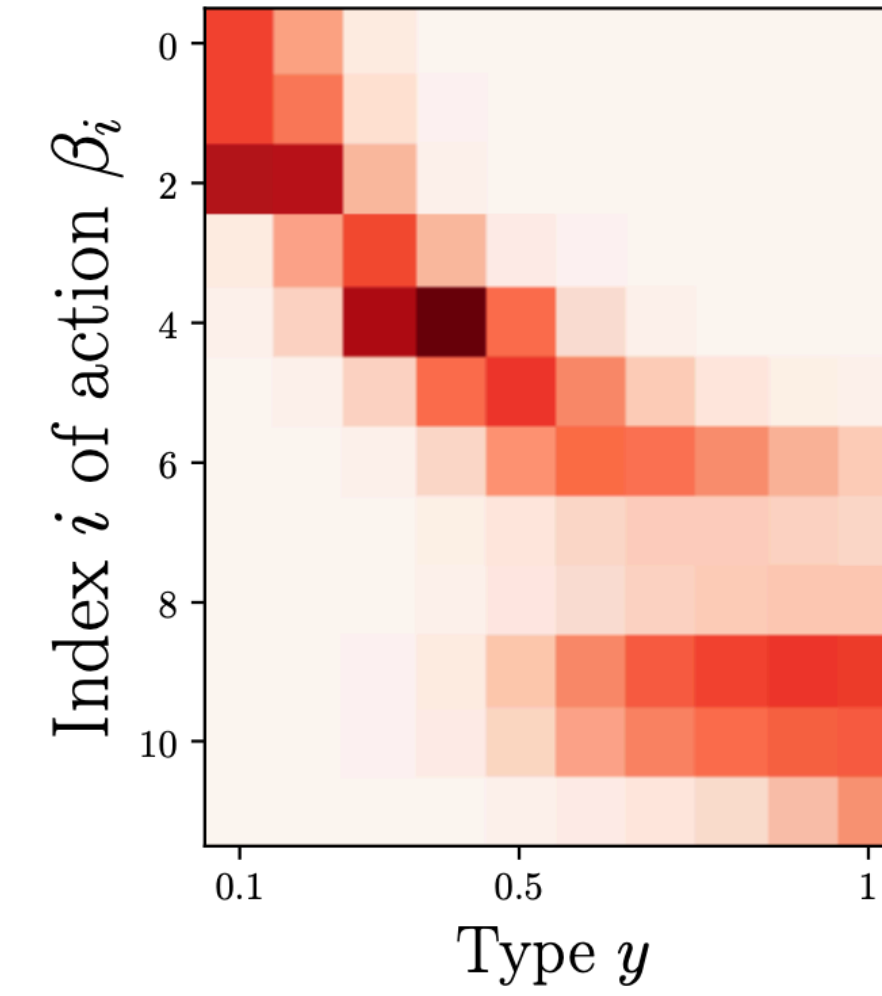
- PRP is **more sensitive to problem parameters**
- for **larger values of  $\lambda$** , Sinkhorn performs well when the privacy cost is predominant
- for **larger values of  $d$** , PRP converges to **worse spurious local minima**
- DC finds better local minima than the other ones

# Experiments

## utility-privacy in repeated auctions



(a) Evolution of privacy-utility with  $\lambda$ .



(b) Joint distribution map for  $\lambda = 0.01$ . The intensity of a point  $(i, j)$  corresponds to the value of  $\gamma(\beta_i, y_j)$ .

Figure 3: Privacy-utility trade-off in online repeated auctions.

- Figure 3a both **decrease with  $\lambda$** , significantly drop at a critical point near 0.05, which can be seen as the cost of information here.
- Figure 3b **partially revealing strategy** that randomizes the type over neighboring types and reveals more information when the revenue is sensible to action

# Experiments

## utility-privacy in repeated auctions

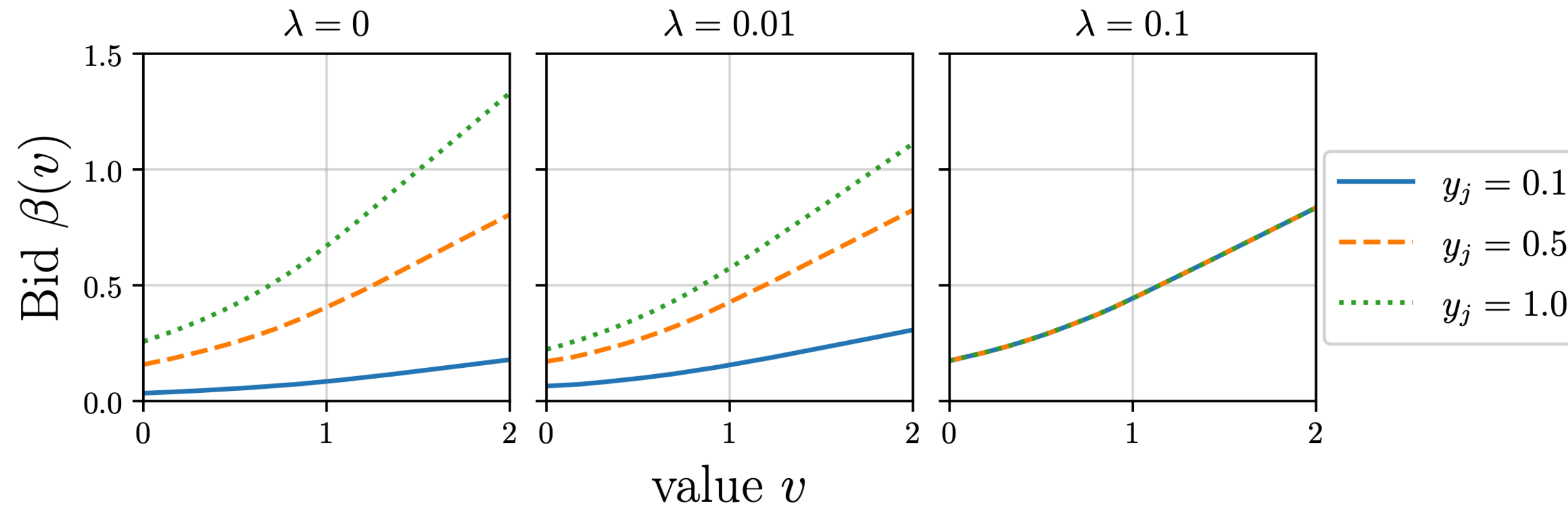


Figure 4: Evolution of the bidding strategy with the type and the regularization constant.

- revealing strategy — action significantly scales with type
- partially revealing strategy — action scales less with type

Thanks!