Annealed Training for Combinatorial Optimization on Graphs

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- ► Introduction
- ► Annealed Training for Combinatorial Optimization
- ► Case Study
- **►** Experiments
- ► Summary

This paper proposed an annealed training framework for Combinatorial Optimization (CO) problems.

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What are the Combinatorial Optimization (CO) problems?



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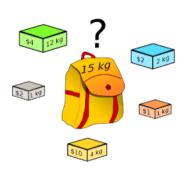
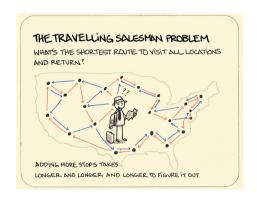


Figure 2: The Knapsack problem

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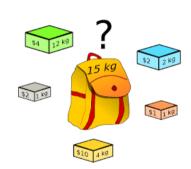


Figure 1: The Travelling salesman problem

Figure 2: The Knapsack problem

Combinatorial Optimization is the process of searching for maxima (or minima) of an objective function F whose domain is a discrete but large configuration space.

We denote the set of combinatorial optimization problems as \mathcal{I} . An instance $I \in \mathcal{I}$ is

$$I = (c(\cdot), \{\psi_i\}_{i=1}^m) := \mathop{\arg\min}_{\boldsymbol{x} \in \{0,1\}^n} c(\boldsymbol{x}) \quad \text{s.t. } \psi_i(\boldsymbol{x}) = 0, \quad i = 1, \dots, m \tag{1}$$

where $c(\cdot)$ is the objective function we want to minimize and $\psi_i \in \{0,1\}$ indicates whether the i-th constraint is satisfied or not.

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We could rewrite the constrained problem into an equivalent unconstrained form via big M method:

$$\underset{\boldsymbol{x} \in \{0,1\}^n}{\arg\min} f^{(I)}(\boldsymbol{x}) := c(\boldsymbol{x}) + \sum_{i=1}^m \beta_i \psi_i(\boldsymbol{x}), \quad \beta_i \ge 0$$
(2)

Using unbiased $f^{(I)}$ to measure the fitness of a solution x, we can define the unbiased energy-based models (EBMs):

$$P_{ au}^{(I)}(\mathbf{x}) \propto e^{-f^{(I)}(\mathbf{x})/ au}$$
 (3)

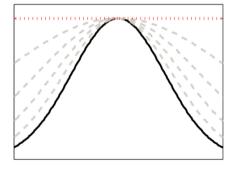
Using unbiased $f^{(I)}$ to measure the fitness of a solution x, we can define the unbiased energy-based models (EBMs):

$$P_{\tau}^{(I)}(\mathbf{x}) \propto e^{-f^{(I)}(\mathbf{x})/ au}$$
 (3)

The temperature τ can be used to control the smoothness of the distribution.

$$P_{ au}^{(I)}(oldsymbol{x}) \propto e^{-f^{(I)}(oldsymbol{x})/ au}$$

- When $\tau \to +$ inf, the EBM P_{τ} converges to a uniform distribution over the whole state space $\{0, 1\}$;
- ▶ When $\tau \to 0$, the EBM P_{τ} converges to a uniform distribution over the optimal solutions.



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Given an instance $I \in \mathcal{I}$, $G_{\theta}(I) = \phi$ generates a vector ϕ that determines a variational distribution Q_{ϕ} to approximate the target distribution $P_{\tau}^{(I)}$.

¹Zhuwen Li, Qifeng Chen, and Vladlen Koltun. "Combinatorial optimization with graph convolutional networks and guided tree search". In: Advances in neural information processing systems 31 (2018).

²Hanjun Dai et al. "A Framework For Differentiable Discovery Of Graph Algorithms". In: (2020).

Given an instance $I \in \mathcal{I}$, $G_{\theta}(I) = \phi$ generates a vector ϕ that determines a variational distribution Q_{ϕ} to approximate the target distribution $P_{\tau}^{(I)}$. In this work, we consider the variational distribution as a product distribution:

$$Q_{\phi}(x) = \prod_{i=1}^{n} (1 - \phi_i)^{1 - x_i} \phi_i^{x_i}$$
(4)

Such a form is a popular choice in learning graphical neural networks for combinatorial optimization¹²

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We want to minimize the KL-divergence:

$$D_{\mathrm{KL}}(Q_{\phi}||P_{\tau}^{(I)}) = \int Q_{\phi}(\boldsymbol{x}) \left(\log Q_{\phi}(\boldsymbol{x}) - \log \frac{e^{-f^{(I)}(\boldsymbol{x})/\tau}}{\sum_{\boldsymbol{x} \in \{0,1\}^n} e^{-f^{(I)}(\boldsymbol{x})/\tau}} \right) d\boldsymbol{x}$$

$$= \frac{1}{\tau} \mathbb{E}_{\boldsymbol{x} \sim Q_{\phi}(\cdot)}[f^{(I)}(\boldsymbol{x})] - H(Q_{\phi}) + \log \sum_{\boldsymbol{x} \in \{0,1\}^n} e^{-f^{(I)}(\boldsymbol{x})/\tau}$$
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We want to minimize the KL-divergence:

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$$= \frac{1}{\tau} \mathbb{E}_{\boldsymbol{x} \sim Q_{\phi}(\cdot)}[f^{(I)}(\boldsymbol{x})] - H(Q_{\phi}) + \log \sum_{\boldsymbol{x} \in \{0,1\}^n} e^{-f^{(I)}(\boldsymbol{x})/\tau}$$
(5)

Remove the terms not involving ϕ and multiply the constant τ , the annealed loss functions for ϕ and τ can be defined as:

$$L_{\tau}(\phi, I) = \mathbb{E}_{\boldsymbol{x} \sim Q_{\phi}(\cdot)}[f^{(I)}(\boldsymbol{x})] - \tau H(Q_{\phi})$$

$$L_{\tau}(\theta) = \mathbb{E}_{I \sim \mathcal{I}}\left[\mathbb{E}_{\boldsymbol{x} \sim Q_{G_{\theta}(\cdot)}}[f^{(I)}(\boldsymbol{x})] - \tau H(Q_{G_{\theta}(I)})\right]$$
(6)

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What is Simulated Annealing?

Simulated Annealing is a stochastic global search optimization algorithm.

Figure 3: The Simulated Annealing process

Annealed Training

To address the non-convexity in training, we employ an annealed training. In particular, we use a large initial temperature τ_0 to smooth the loss function and reduce τ_t gradually to zero during training.

$$\tau_k = \frac{\tau_0}{1 + \alpha k} \tag{7}$$

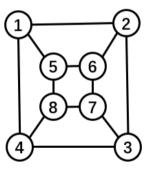
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Maximum Independent Set

An independent set is a subset of the vertices $S \subseteq V$, such that for arbitrary $i, j \in S$, $(i, j) \notin E$.

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An independent set is a subset of the vertices $S \subseteq V$, such that for arbitrary $i, j \in S$, $(i, j) \notin E$.

The maximum independent set problem is to find an independent set S having the largest weight. If we denote $x_i = 1$ to indicate $i \in S$ and $x_i = 0$ to indicate $i \notin S$, the problem can be formulated as:

$$\underset{x \in \{0,1\}^n}{\arg\min} c(x) := -\sum_{i=1}^n w_i x_i, \quad \text{s.t. } x_i x_j = 0, \forall (i,j) \in E$$
(8)

We define the corresponding energy function:

$$f(x) := -\sum_{i=1}^{n} w_i x_i + \sum_{(i,j) \in E} \beta_{ij} x_i x_j$$
(9)

Proof - Maximum Independent Set

$$f(x) := -\sum_{i=1}^{n} w_i x_i + \sum_{(i,j) \in E} \beta_{ij} x_i x_j$$
 (10)

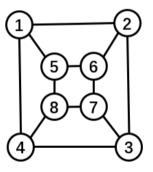
Proposition A.1. If $\beta_{ij} \ge \min\{w_i, w_j\}$ for all $(i,j) \in E$, then for any $x \in \{0, 1\}^n$, there exists a $x' \in \{0, 1\}^n$ that satisfies the constraints in Eqn.8 and has lower energy: $f(x') \le f(x)$.

Maximum Clique

A clique is a subset of the vertices $S \subseteq V$, such that every two distinct $i, j \in S$ are adjacent: $(i, j) \in E$.

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(11)

where $E^c=\{(i,j)\in V\times V:i\neq j,(i,j)\not\in E\}$ is the set of complement edges on graph G.

We define the corresponding energy function:

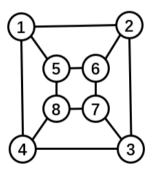
$$f(x) := -\sum_{i=1}^{n} w_i x_i + \sum_{(i,j) \in E^c} \beta_{ij} x_i x_j$$
 (12)

Minimum Dominate Set

A dominate set is a subset of the vertices $S \subseteq V$, where for any $v \in V$, there exists $u \in S$ such that $(u, v) \in E$ or u = v.

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The minimum dominate set problem is to find a dominate set S having the minimum weight. If we denote $x_i = 1$ to indicate $i \in S$ and $x_i = 0$ to indicate $i \notin S$, the problem can be formulated as:

$$\underset{x \in \{0,1\}^n}{\arg\min} c(x) := \sum_{i=1}^n w_i x_i, \quad \text{s.t. } (1 - x_i) \prod_{j \in N(i)} (1 - x_j) = 0, \forall i \in V$$
(13)

We define the corresponding energy function:

$$f(x) := \sum_{i=1}^{n} w_i x_i + \sum_{i=1}^{n} \beta_i (1 - x_i) \prod_{j \in N(i)} (1 - x_j)$$
(14)

Proof - Minimum Dominate Set

$$f(x) := \sum_{i=1}^{n} w_i x_i + \sum_{i=1}^{n} \beta_i (1 - x_i) \prod_{i \in N(i)} (1 - x_j)$$
(15)

Proposition A.3. If $\beta_i \ge \min\{w_k : k \in N(i) \text{ or } k = i\}$, then for any $x \in \{0, 1\}^n$, there exists a $x' \in \{0, 1\}^n$ that satisfies the constraints in Eqn.13 and has lower energy: $f(x') \le f(x)$

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Experiments

Table 1: Evaluation of Maximum Independent Set

Size	small		large		Collab		Twitter	
Method	ratio	time (s)	ratio	time (s)	ratio	time (s)	ratio	time (s)
Erdos Our's	0.805 ± 0.052 0.898 ± 0.030	$0.156 \\ 0.165$	0.781 ± 0.644 0.848 ± 0.529	$2.158 \\ 2.045$	$0.986 \pm 0.056 \\ 0.997 \pm 0.020$	$0.010 \\ 0.010$	0.975 ± 0.033 0.986 ± 0.012	$0.020 \\ 0.020$
Greedy MFA RUNCS G(0.5s) G(1.0s)		0.002 0.042 1.936 0.723 1.063	$\begin{array}{c} 0.720 \pm 0.046 \\ 0.747 \pm 0.056 \\ 0.587 \pm 0.312 \\ 0.632 \pm 0.176 \\ 0.635 \pm 0.176 \end{array}$	0.009 0.637 7.282 1.199 1.686	$\begin{array}{c} 0.996 \pm 0.017 \\ 0.998 \pm 0.007 \\ 0.912 \pm 0.101 \\ 1.000 \pm 0.000 \\ 1.000 \pm 0.000 \end{array}$	0.001 0.002 0.254 0.029 0.029	$\begin{array}{c} 0.957 \pm 0.037 \\ 0.994 \pm 0.010 \\ 0.845 \pm 0.184 \\ 0.950 \pm 0.191 \\ 1.000 \pm 0.000 \end{array}$	0.006 0.003 4.429 0.441 0.462

Table 2: Evaluation of Maximum Clique

Size	small		large		Collab		Twitter	
Method	ratio	time (s)	ratio	time (s)	ratio	time (s)	ratio	time (s)
Erdos Our's	0.813 ± 0.067 0.901 ± 0.055	$0.279 \\ 0.262$	0.735 ± 0.084 0.831 ± 0.078	$0.622 \\ 0.594$	0.960 ± 0.019 0.988 ± 0.011	$0.139 \\ 0.143$	0.822 ± 0.085 0.920 ± 0.083	$0.222 \\ 0.213$
Greedy MFA RUNCS G(0.5s) G(1.0s)		0.002 0.144 2.045 0.599 0.705	$\begin{array}{c} 0.727 \pm 0.038 \\ 0.710 \pm 0.045 \\ 0.574 \pm 0.299 \\ 0.812 \pm 0.087 \\ 0.847 \pm 0.101 \end{array}$	0.014 0.147 7.332 0.617 1.077	$\begin{array}{c} 0.999 \pm 0.002 \\ 1.000 \pm 0.000 \\ 0.887 \pm 0.134 \\ 0.997 \pm 0.035 \\ 0.999 \pm 0.015 \end{array}$	0.001 0.005 0.164 0.061 0.062	$\begin{array}{c} 0.959 \pm 0.034 \\ 0.994 \pm 0.010 \\ 0.832 \pm 0.153 \\ 0.976 \pm 0.065 \\ 0.997 \pm 0.029 \end{array}$	0.001 0.010 4.373 0.382 0.464

Experiments

Table 3: Evaluation of Minimum Dominate Set

Size	small		large		Collab		Twitter	
Method	ratio	time (s)	ratio	time (s)	ratio	time (s)	ratio	time (s)
	0.909 ± 0.037 0.954 ± 0.006	$0.121 \\ 0.120$	0.889 ± 0.017 0.931 ± 0.015	$0.449 \\ 0.453$	$0.982 \pm 0.070 \\ 0.993 \pm 0.062$	$0.007 \\ 0.006$	0.924 ± 0.098 0.952 ± 0.074	$0.015 \\ 0.016$
MFA 0 G(0.5s) 0	0.743 ± 0.053 0.926 ± 0.032 0.993 ± 0.014 0.999 ± 0.005	0.254 0.213 0.381 0.538	0.735 ± 0.026 0.910 ± 0.016 0.994 ± 0.013 0.999 ± 0.005	3.130 3.520 0.384 0.563	0.661 ± 0.406 0.895 ± 0.210 1.000 ± 0.000 1.000 ± 0.000	0.028 0.030 0.042 0.042	$\begin{array}{c} 0.741 \pm 0.142 \\ 0.952 \pm 0.076 \\ 1.000 \pm 0.000 \\ 0.839 \pm 0.000 \end{array}$	0.079 0.099 0.084 0.084

Table 4: Evaluation of Minimum Cut

Size	SF-295		Faceboo	k	Twitter	
Method	ratio	time (s)	ratio	time (s)	ratio	time (s)
Erdos Our's	$egin{array}{l} 0.124 \pm 0.001 \ 0.135 \pm 0.011 \end{array}$	$0.22 \\ 0.23$	0.156 ± 0.026 0.151 ± 0.045	$289.3 \\ 290.5$	0.292 ± 0.009 0.201 ± 0.007	$6.17 \\ 6.16$
L1 GNN L2 GNN Pagerank-Nibble CRD MQI Simple-Local G(10s)	$\begin{array}{c} 0.188 \pm 0.045 \\ 0.149 \pm 0.038 \\ 0.375 \pm 0.001 \\ 0.364 \pm 0.001 \\ 0.659 \pm 0.000 \\ 0.650 \pm 0.024 \\ 0.105 \pm 0.000 \end{array}$	0.02 0.01 1.48 0.03 0.03 0.05 0.16	$\begin{array}{c} 0.571 \pm 0.191 \\ 0.305 \pm 0.082 \\ \text{N/A} \\ 0.301 \pm 0.097 \\ 0.935 \pm 0.024 \\ 0.961 \pm 0.019 \\ 0.961 \pm 0.010 \\ \end{array}$	13.83 13.83 N/A 596.46 408.52 1787.79 1787.79	$\begin{array}{c} 0.318 \pm 0.077 \\ 0.388 \pm 0.074 \\ 0.603 \pm 0.005 \\ 0.502 \pm 0.020 \\ 0.887 \pm 0.007 \\ 0.895 \pm 0.006 \\ 0.535 \pm 0.006 \end{array}$	0.53 0.53 20.62 20.35 0.71 0.84 52.98

Parameter Change Distance

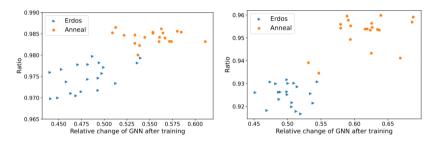


Figure 2: Distance in MIS

Figure 3: Distance in MDS

The relative change is calculated as $\frac{\|u-v\|_2}{\|v\|_2}$, where v and u are vectors flattened from the parameters of GNN before and after training.

Ablation Study

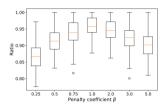
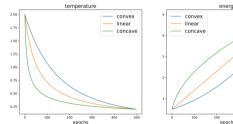


Figure 4: Ablation for β



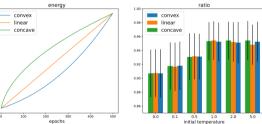


Figure 5: Ablation for annealing schedule

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- ► This paper is poorly written...