# Improved Complexity Bounds in Wasserstein Barycenter Problem

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- 1 Introduction
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- 3 Algorithm 1: Mirror Prox for Wasserstein Barycenter
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## OT: Arithmetic calculations problem

- Simplex method or interior point method:  $\widetilde{O}(n^3)$
- Sinkhorn algorithm:  $\widetilde{O}(n^2 \|C\|_{\infty}^2/\varepsilon^2)$ , with  $\varepsilon$ -precision. C is cost matrix, the regularization parameter before negative entropy is of order  $\varepsilon$ .
- Accelerated Sinkhorn algorithm:  $\widetilde{O}\left(n^{2.5}\|C\|_{\infty}/\varepsilon\right)$ . In practice, it has better dependence on  $\varepsilon$  but not on n.
- All entropy-regularized based approaches are numerically unstable when the regularizer parameter  $\gamma$  before negative entropy is small (this also means that precision  $\varepsilon$  is high as  $\gamma$  must be selected proportional to  $\varepsilon$ .

The recent work (Jambulapati et al. (2019)) privides an optimal method:  $\widetilde{O}\!\left(n^2\|C\|_\infty/\varepsilon\right)$ 

- Based on dual extrapolation and area-convexity.
- Without additional penalization.



## Wasserstein Barycenter problem

- Iterative Bregman projections (IBP) algorithm: The IBP is an extension of the Sinkhorn's algorithm for m measures, and hence, its complexity is m times more than the Sinkhorn complexity:  $\widetilde{O}\left(mn^2\|C\|_{\infty}^2/\varepsilon^2\right)$ .
- Accelerated IBP algorithm:  $\widetilde{O}(mn^{2.5}||C||_{\infty}/\varepsilon)$ .
- Another fast version of IBP, FastIBP:  $\widetilde{O}\left(mn^{\frac{7}{3}}\|C\|_{\infty}^{\frac{4}{3}}/\varepsilon^{\frac{4}{3}}\right)$ .

## Contribution

#### The first contribution:

- Propose an algorithm which does not suffer from a small value of the regularization parameter.
- Convergence rate:  $\widetilde{O}(mn^{2.5}||C||_{\infty}/\varepsilon)$ , not worse than the celebrated accelerated IBP.
- Based on mirror prox with specific prox-function.

#### The second contribution:

- Propose an algorithm that has better complexity than the accelerated IBP.
- Convergence rate:  $\widetilde{O}(mn^2||C||_{\infty}/\varepsilon)$ .
- Based on rewriting the WB problem as a saddle-point problem and further application of the dual extrapolation scheme under the weaker convergence requirements of area-convexity.

In some sense, the first algorithm can be seen as a simplified version of the second algorithm.

## Contribution

Table 1: Algorithms and their rates of convergence for the Wasserstein barycenter problem

Approach	Paper	Complexity
IBP	(Kroshnin et al., 2019)	$\widetilde{O}\left(\frac{mn^2\ C\ _{\infty}^2}{\varepsilon^2}\right)$ $\widetilde{O}\left(\frac{mn^2\sqrt{n}\ C\ _{\infty}}{mn^2\sqrt{n}\ C\ _{\infty}}\right)$
Accelerated IBP	(Guminov et al., 2019)	$\widetilde{O}\left(\frac{mn^2\sqrt{n}\ C\ _{\infty}}{\varepsilon}\right)$
FastIBP	(Lin et al., 2020)	$\widetilde{O}\left(\frac{mn^2\sqrt[3]{n}\ C\ _{\infty}^{4/3}}{\varepsilon\sqrt[3]{\varepsilon}}\right)$
Mirror prox with specific norm	This work	$\widetilde{O}\left(\frac{mn^2\sqrt{n}\ C\ _{\infty}}{\varepsilon}\right)$
Dual extrapolation with area-convexity	This work	$\widetilde{O}\left(\frac{mn^2\ C\ _{\infty}}{\varepsilon}\right)$

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## Problem Statement

• Let  $\Delta_n = \{ p \in \mathbb{R}^n_+ : \sum_{i=1}^n p_i = 1 \}$  be the probability simplex. Given two histograms  $p, q \in \Delta_n$  and ground cost  $C \in \mathbb{R}^{n*n}_+$ , the OT problem is formulated as follows

$$W(p,q) = \min_{X \in \mathcal{U}(p,q)} \langle C, X \rangle \tag{1}$$

- ullet where X is a transport plan.  ${\mathcal U}$  is the transport polytope.
- Let d be vectorized cost matrix of C, x be vectorized transport plan of X,  $b = \begin{pmatrix} p \\ q \end{pmatrix}$ , and  $A = \{0,1\}^{2n \times n^2}$  be an incidence matrix.
- $\bullet$  As  $\sum_{i,j=1}^{n} X_{ij} = 1,$  (Jambulapati et al. (2019)) rewrite (1) as

$$W(p, q) = \min_{x \in \Lambda_n 2} \max_{y \in [-1, 1]^{2n}} \left\{ d^{\mathsf{T}} x + 2||d||_{\infty} \left( y^{\mathsf{T}} A x - b^{\mathsf{T}} y \right) \right\}$$
(2)



#### Problem Statement

• Given histograms  $q_1, q_2, ..., q_m \in \Delta_n$ , a WB of those measures is a solution of the following problem:

$$p^* = \arg\min_{p \in \Delta_n} \frac{1}{m} \sum_{i=1}^m W(p, q_i)$$
 (3)

• Rewrite (3) using the reformulation (2) of OT as follows:

$$\min_{p \in \Delta_n} \frac{1}{m} \sum_{i=1}^m \min_{x_i \in \Delta_n 2} \max_{y_i \in [-1,1]^{2n}} \left\{ d^{\mathsf{T}} x_i + 2 ||d||_{\infty} \left( y_i^{\mathsf{T}} A x_i - b_i^{\mathsf{T}} y_i \right) \right\}$$
(4)

## Problem Statement

• Define spaces  $\mathcal{X} \triangleq \prod^m \Delta_{n^2} \times \Delta_n$  and  $\mathcal{Y} \triangleq [-1, 1]^{2mn}$ , where  $\prod^m \Delta_{n^2} \times \Delta_n = \underbrace{\Delta_{n^2} \times \ldots \times \Delta_{n^2}}_{m} \times \Delta_n$ . Reweite (4) for column vectors  $\mathbf{x} = \left(x_1^{\mathsf{T}}, \ldots, x_m^{\mathsf{T}}, p^{\mathsf{T}}\right)^{\mathsf{T}} \in \mathcal{X}$  and  $\mathbf{y} = \left(y_1^{\mathsf{T}}, \ldots, y_m^{\mathsf{T}}\right)^{\mathsf{T}} \in \mathcal{Y}$  as follows:

$$\min_{\mathbf{x} \in \mathcal{X}} \max_{\mathbf{y} \in \mathcal{Y}} F(\mathbf{x}, \mathbf{y}) \triangleq \frac{1}{m} \left\{ \mathbf{d}^{\mathsf{T}} \mathbf{x} + 2 ||\mathbf{d}||_{\infty} \left( \mathbf{y}^{\mathsf{T}} \mathbf{A} \mathbf{x} - \mathbf{c}^{\mathsf{T}} \mathbf{y} \right) \right\}$$
(5)

- As objective  $F(\mathbf{x}, \mathbf{y})$  in (5) is convex in x and concave in y, problem (5) is a saddle-point problem.
- $\begin{aligned} \bullet & \text{ where } \mathbf{d} = (d^\intercal, \dots, d^\intercal, 0_n^\intercal)^\intercal , \mathbf{c} = \begin{pmatrix} 0_n^\intercal, q_1^\intercal, \dots, 0_n^\intercal, q_m^\intercal \end{pmatrix}^\intercal \text{ and } \\ \mathbf{A} = \begin{pmatrix} \hat{A} & \mathcal{E} \end{pmatrix} \in \{-1, 0, 1\}^{2mn\times(mn^2+n)} \text{ with block-diagonal matrix } \mathbf{A} \text{ of m} \\ & \text{blocks } \hat{A} = \begin{pmatrix} A & 0_{2n\times n^2} & \cdots & 0_{2n\times n^2} \\ 0_{2n\times n^2} & A & \cdots & 0_{2n\times n^2} \\ \vdots & \vdots & \ddots & \vdots \\ 0_{2n\times n^2} & 0_{2n\times n^2} & \cdots & A \end{pmatrix} \text{ and matrix } \\ \mathcal{E}^\intercal = \underbrace{((-I_n & 0_{n\times n})\underbrace{(-I_n & 0_{n\times n})} \cdots \underbrace{(-I_n & 0_{n\times n})}}_{} \cdots \underbrace{(-I_n & 0_{n\times n})}_{} ). \end{aligned}$

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## Setup

- Notation: For a prox-function d(x), define the corresponding Bregman divergence  $B(x, y) = d(x) d(y) \langle \nabla d(y), x y \rangle$ . For example, the Euclidean  $\mathbf{l}_2$ -norm  $\|\mathbf{y}\|_2$ , prox-function  $d_{\mathbf{Y}}(\mathbf{y}) = \frac{1}{2}\|\mathbf{y}\|_2^2$ , and the  $B_{\mathbf{y}}(\mathbf{y}, \check{\mathbf{y}}) = \frac{1}{2}\|\mathbf{y} \check{\mathbf{y}}\|_2^2$
- For space  $X \triangleq \prod^m \Delta_{n^2} \times \Delta_n$ , we choose the following specific norm  $\|\mathbf{x}\|_{\mathcal{X}} = \sqrt{\sum_{i=1}^m \|x_i\|_1^2 + m\|p\|_1^2}$  for  $\mathbf{x} = (x_1, \dots, x_m, p)^T$ . Given X with prox-function  $d_X(\mathbf{x}) = \sum_{i=1}^m \langle x_i, \ln x_i \rangle + m \langle p, \ln p \rangle$  and  $B_X(\mathbf{x}, \check{\mathbf{x}}) = \sum_{i=1}^m \langle x_i, \ln (x_i/\check{x}_i) \rangle \sum_{i=1}^m \mathbf{1}^\top (x_i \check{x}_i) + m \langle p, \ln (p/\check{p}) \rangle m \mathbf{1}^\top (p \check{p})$ .
- Define  $R_{\chi}^2 = \sup_{\mathbf{x} \in \chi} d_{\chi}(\mathbf{x}) \min_{\mathbf{x} \in \chi} d_{\chi}(\mathbf{x})$  and  $R_{\chi}^2 = \sup_{\mathbf{y} \in \mathcal{Y}} d_{\chi}(\mathbf{y}) \min_{\mathbf{y} \in \mathcal{Y}} d_{\chi}(\mathbf{y}).$
- Definition:  $f(\widetilde{\mathbf{x}}, \mathbf{y})$  is  $(L_{xx}, L_{xy}, L_{yx}, L_{yy})$ -smooth if for any  $x, x' \in X$  and  $y, y' \in \mathcal{Y}$   $\|\nabla f(\mathbf{x}, \mathbf{y}) \nabla f(\mathbf{x}', \mathbf{y})\|_{L^{\infty}} \leq L_{x} \|\mathbf{y} \mathbf{y}'\|_{L^{\infty}}$

$$\begin{aligned} &\|\nabla_{\mathbf{x}}f(\mathbf{x},\mathbf{y}) - \nabla_{\mathbf{x}}f(\mathbf{x}',\mathbf{y})\|_{\mathcal{X}^{s}} \leq L_{\mathbf{x}\mathbf{x}} \|\mathbf{x} - \mathbf{x}'\|_{\mathcal{X}}, \\ &\|\nabla_{\mathbf{x}}f(\mathbf{x},\mathbf{y}) - \nabla_{\mathbf{x}}f(\mathbf{x},\mathbf{y}')\|_{\mathcal{X}^{s}} \leq L_{\mathbf{x}\mathbf{y}} \|\mathbf{y} - \mathbf{y}'\|_{\mathcal{Y}}, \\ &\|\nabla_{\mathbf{y}}f(\mathbf{x},\mathbf{y}) - \nabla_{\mathbf{y}}f(\mathbf{x},\mathbf{y}')\|_{\mathcal{Y}^{s}} \leq L_{\mathbf{y}\mathbf{y}} \|\mathbf{y} - \mathbf{y}'\|_{\mathcal{Y}}, \\ &\|\nabla_{\mathbf{y}}f(\mathbf{x},\mathbf{y}) - \nabla_{\mathbf{y}}f(\mathbf{x}',\mathbf{y})\|_{\mathcal{Y}^{s}} \leq L_{\mathbf{y}\mathbf{x}} \|\mathbf{x} - \mathbf{x}'\|_{\mathcal{X}}. \end{aligned}$$

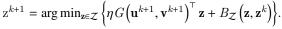


## Implementation

• As problem (5) is a saddle-point problem, we will evaluate the quality of an algorithm that outputs a pair of solutions  $(\widetilde{\mathbf{x}}, \widetilde{\mathbf{y}}) \in (\mathcal{X}, \mathcal{Y})$  through the so-called duality gap

$$\max_{\mathbf{y} \in \mathcal{Y}} F(\widetilde{\mathbf{x}}, \mathbf{y}) - \min_{\mathbf{x} \in \mathcal{X}} F(\mathbf{x}, \widetilde{\mathbf{y}}) \le \varepsilon \tag{6}$$

• The first algorithm is based on mirror prox (MP) algorithm (Nemirovski, 2004) on space  $\mathcal{Z} \triangleq \mathcal{X} \times \mathcal{Y}$  with prox-function  $d_{\mathcal{Z}}(\mathbf{z}) = a_1 d_{\mathcal{X}}(\mathbf{x}) + a_2 d_{\mathcal{Y}}(\mathbf{y})$  and  $B_{\mathcal{Z}}(\mathbf{z}, \check{\mathbf{z}}) = a_1 B_{\mathcal{X}}(\mathbf{x}, \check{\mathbf{x}}) + a_2 B_{\mathcal{Y}}(\mathbf{y}, \check{\mathbf{y}})$ , where  $a_1 = \frac{1}{R_{\mathcal{X}}^2}$ ,  $a_2 = \frac{1}{R_{\mathcal{Y}}^2}$ ,  $\left(\begin{array}{c} \mathbf{u}^{k+1} \\ \mathbf{v}^{k+1} \end{array}\right) = \arg\min_{\mathbf{z} \in \mathcal{Z}} \left\{ \eta G(\mathbf{x}^k, \mathbf{y}^k)^{\mathsf{T}} \mathbf{z} + B_{\mathcal{Z}}(\mathbf{z}, \mathbf{z}^k) \right\}$ ,





## Implementation

• Here  $\eta$  is learning rate,  $\mathbf{z}^1 = \arg\min_{\mathbf{z} \in \mathcal{Z}} d_{\mathcal{Z}}(\mathbf{z})$  and  $G(\mathbf{x}, \mathbf{y})$  is a gradient operator defined as follows

$$G(\mathbf{x}, \mathbf{y}) = \begin{pmatrix} \nabla_{\mathbf{x}} F(\mathbf{x}, \mathbf{y}) \\ -\nabla_{\mathbf{y}} F(\mathbf{x}, \mathbf{y}) \end{pmatrix} = \frac{1}{m} \begin{pmatrix} d + 2||d||_{\infty} \mathbf{A}^{\top} \mathbf{y} \\ 2||d||_{\infty} (\mathbf{c} - \mathbf{A}\mathbf{x}) \end{pmatrix}$$
(7)

• If  $F(\mathbf{x}, \mathbf{y})$  is  $(L_{xx}, L_{xy}, L_{yx}, L_{yy})$ -smooth, then to satify (6) with  $\widetilde{\mathbf{x}} = \frac{1}{N} \sum_{k=1}^{N} \mathbf{u}^{k}, \widetilde{\mathbf{y}} = \frac{1}{N} \sum_{k=1}^{N} \mathbf{v}^{k}$  one needs to perform

$$N = \frac{4}{\varepsilon} \max \left\{ L_{xx} R_{\chi}^2, L_{xy} R_{\chi} R_{y}, L_{yx} R_{y} R_{\chi}, L_{yy} R_{y}^2 \right\}$$
 (8)

• iterations of MP(Bubeck, 2014) with

$$\eta = 1/\left(2\max\left\{L_{xx}R_{\mathcal{X}}^{2}, L_{xy}R_{\mathcal{X}}R_{\mathcal{Y}}, L_{yx}R_{\mathcal{Y}}R_{\mathcal{X}}, L_{yy}R_{\mathcal{Y}}^{2}\right\}\right)$$
(9)

# Complexity Bound

- Lemma: Objective  $F(\mathbf{x}, \mathbf{y})$  in (5) is  $(L_{xx}, L_{xy}, L_{yx}, L_{yy})$ -smooth with  $L_{xx} = L_{yy} = 0$  and  $L_{xy} = L_{yx} = 2\sqrt{2}||d||_{\infty}/m$ .
- Theorem: Assume that  $F(\mathbf{x}, \mathbf{y})$  in (5) is  $(0, 2\sqrt{2}||d||_{\infty}/m, 2\sqrt{2}||d||_{\infty}/m, 0)$ -smooth and  $R_{\mathcal{X}} = \sqrt{3m\ln n}, R_{\mathcal{Y}} = \sqrt{mn}$ . Then after  $N = 8||d||_{\infty}\sqrt{6n\ln n}/\varepsilon$  iterations, Algorithm 1 with  $\eta = \frac{1}{4||d||_{\infty}\sqrt{6n\ln n}}$  outputs a pair  $(\widetilde{\mathbf{u}}, \widetilde{\mathbf{v}}) \in (\mathcal{X}, \mathcal{Y})$  such that  $\max_{\mathbf{v} \in \mathcal{Y}} F(\widetilde{\mathbf{u}}, \mathbf{v}) \min_{\mathbf{x} \in \mathcal{X}} F(\mathbf{x}, \widetilde{\mathbf{v}}) \leq \varepsilon$ .
- The total complexity of Algorithm 1 is  $O(mn^2 \sqrt{n \ln n} ||d||_{\infty} \varepsilon^{-1})$ . The complexity of one iteration of Algorithm 1 is  $O(mn^2)$  as the number of non-zero elements in matrix A is  $2n^2$ , and m is the number of vector-components in  $\mathbf{y}$  and  $\mathbf{x}$ . Multiplying this by the number of iterations N, we get the result.
- As d is the vectorized cost matrix of C, we may reformulate the complexity results of Theorem with respect to C as  $O(mn^2 \sqrt{n \ln n} ||C||_{\infty} \varepsilon^{-1})$ .

## Algorithm 1

#### Algorithm 1 Mirror Prox for Wasserstein Barycenters

Input: measures  $q_1, ..., q_m$ , linearized cost matrix d, incidence matrix A, step  $\eta$ , starting points  $p^1 = \frac{1}{n} \mathbf{1}_n$ ,  $x_1^1 = ... = x_m^1 = \frac{1}{n^2} \mathbf{1}_{n^2}, y_1^1 = ... = y_m^1 = \mathbf{0}_{2n}$ 

- 1:  $\alpha = 2||d||_{\infty} \eta n$ ,  $\beta = 6||d||_{\infty} \eta \ln n$ ,  $\gamma = 3m\eta \ln n$ .
- 2: for  $k = 1, 2, \dots, N 1$  do
  - 3: for  $i = 1, 2, \dots, m \text{ do}$
- 4:  $v_i^{k+1} = y_i^k + \alpha \left(Ax_i^k \begin{pmatrix} p^k \\ q_i \end{pmatrix}\right)$ ,
  - Project  $v_i^{k+1}$  onto  $[-1, 1]^{2n}$

$$u_i^{k+1} = \frac{x_i^k \odot \exp\left\{-\gamma \left(d + 2\|d\|_{\infty} A^{\top} y_i^k\right)\right\}}{\sum\limits_{l=1}^{n^2} [x_i^k]_l \exp\left\{-\gamma \left([d]_l + 2\|d\|_{\infty} [A^{\top} y_i^k]_l\right)\right\}}$$

6: end for

5:

7:

$$s^{k+1} = \frac{p^k \odot \exp\left\{\beta \sum_{i=1}^m [y_i^k]_{1...n}\right\}}{\sum_{l=1}^n [p^k]_l \exp\left\{\beta \sum_{i=1}^m [y_i^k]_l\right\}}$$

- 8: for  $i = 1, 2, \dots, m$  do
- 9:  $y_i^{k+1} = y_i^k + \alpha \left( A u_i^{k+1} {s^{k+1} \choose q_i} \right)$ 
  - Project  $y_i^{k+1}$  onto  $[-1, 1]^{2n}$

$$x_i^{k+1} = \frac{x_i^k \odot \exp\left\{-\gamma \left(d + 2\|d\|_{\infty} A^{\top} v_i^{k+1}\right)\right\}}{\sum\limits_{l = 1}^{n^2} [x_i^k]_l \exp\left\{-\gamma \left([d]_l + 2\|d\|_{\infty} [A^{\top} v_i^{k+1}]_l\right)\right\}}$$

- 11: end for
- 12:

10:

$$p^{k+1} = \frac{p^k \odot \exp \left\{\beta \sum_{i=1}^m [v_i^{k+1}]_{1...n}\right\}}{\sum_{l=1}^n [p^k]_l \exp \left\{\beta \sum_{i=1}^m [v_i^{k+1}]_l\right\}}$$

- 13: end for
- Output:  $\widetilde{\mathbf{u}} = \sum_{k=1}^{N} \begin{pmatrix} u_1^k \\ \vdots \\ u_m^k \\ o_k^k \end{pmatrix}, \widetilde{\mathbf{v}} = \sum_{k=1}^{N} \begin{pmatrix} v_1^k \\ \vdots \\ v_m^k \end{pmatrix}$

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#### MNIST and notMNIST

• The result of IBP with regularizing parameter  $\gamma$  is numberically unstable, as  $\gamma$  must be selected proportional to  $\varepsilon$ .

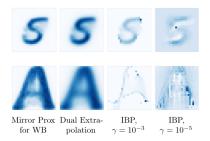
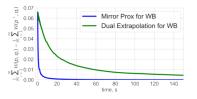


Figure 1: WBs of hand-written digit '5' (first row) and of letters 'A' (second row) computed by Algorithm 1 (Mirror Prox for WB), Algorithm 4 (Dual Extrapolation for WB) and the IBP with small values of the regularizing parameter.

#### Gaussian measures

- To compare the convergence of the proposed algorithms, we randomly generated 10 Gaussian measures with equally spaced support of 100 points in [-10,10], mean from [-5,5] and variance from [0.8,1.8].
- Figure 2 presents the convergence with respect to the function optimality gap  $\frac{1}{m} \sum_{i=1}^{m} \mathcal{W}(p, q_i) \frac{1}{m} \sum_{i=1}^{m} \mathcal{W}(p^*, q_i)$ . Here  $p^*$  is the true bartcenter.



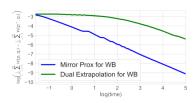
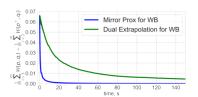


Figure 2: Convergence of Algorithm 1 (Mirror Prox for WB) and Algorithm 4 (Dual Extrapolation for WB) to the true barycenter of Gaussian measures w.r.t the function optimality gap  $\frac{1}{m} \sum_{i=1}^{m} \mathcal{W}(p, q_i) - \frac{1}{m} \sum_{i=1}^{m} \mathcal{W}(p^*, q_i)$ . Here  $p^*$  is the true barycenter.

#### Gaussian measures

- Algorithm 4 has better complexity bound, Algorithm 1 has better convergence in practice. The slope ration -1 for the convergence of Algorithm 1 in log-scale perfectly fits the theoretical dependence of working time (iteration number N) on the desired accuracy  $\varepsilon$  ( $N \sim \varepsilon^{-1}$  from Theorem).
- For Algorithm 4, this slope ratio −1 is achieved only after a number of iterations but this is due to the need of solving practically computationally costly subproblems.



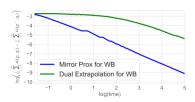


Figure 2: Convergence of Algorithm 1 (Mirror Prox for WB) and Algorithm 4 (Dual Extrapolation for WB) to the true barycenter of Gaussian measures w.r.t the function optimality gap  $\frac{1}{m}\sum_{i=1}^{m} \mathcal{W}(p,q_i) - \frac{1}{m}\sum_{i=1}^{m} \mathcal{W}(p^*,q_i)$ . Here  $p^*$  is the true barycenter.

#### Gaussian measures

• Figure 3 illustrates the convergence of the barycenters obtained by Algorithms 1 and 4 to the true barycenter.

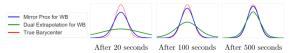


Figure 3: Convergence of the barycenters obtained by Algorithm 1 (Mirror Prox for WB) and Algorithm 4 (Dual Extrapolation for WB) to the true barycenter of Gaussian measures.

• Figure 4 illustrates better approximations of the true Gaussian barycenter by Algorithms 1 and 4 compared to the *gamma*-regularized IBP barycenter.

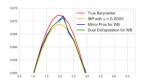


Figure 4: Convergence of the barycenters obtained by Algorithm 1 (Mirror Prox for WB), Algorithm 4 (Dual Extrapolation for WB), and the IBP to the true barycenter of Gaussian measures.