编译原理实验

徐达烽 16340260 数字媒体技术

思路

- 1. 使用规范名称,提高代码可读性。
- 2. 将大于等于,小于等于和不等于号修改成>=,<=,<>,并修改这些符号的处理逻辑。
- 3. 修改源程序中存在的bug,使其成功运行。
- 4. 添加RED, WRT命令, 从而实现read, write函数。

第一部分

1. 编译程序源代码

```
program PLO;
{带有代码生成的PLO编译程序}
const
    kReservedWords = 11; {保留字的个数}
    kIdentsMax = 100; {标识符表长度}
    kNumLengthMax = 14; {数字的最大位数}
    kIdentLengthMax = 10; {标识符的长度}
    kAddrMax = 2047; {最大地址}
    kNestingLayersMax = 3; {程序体嵌套的最大深度}
    kInstructionsMax = 200; {代码数组的大小}
    kDebugMessageOn = 1;
type
    Symbol = (NUL, IDENT, NUMBER, PLUS, MINUS, TIMES, SLASH, ODDSYM,
        EQL, NEQ, LSS, LEQ, GTR, GEQ, LPAREN, RPAREN, COMMA, SEMICOLON,
        PERIOD, BECOMES, BEGINSYM, ENDSYM, IFSYM, THENSYM,
        WHILESYM, DOSYM, CALLSYM, CONSTSYM, VARSYM, PROCSYM);
    Identifier = packed array [1..kIdentLengthMax] of char;
    ObjectType = (kConstant, kVariable, kProcedure);
    SymbolSet = set of Symbol;
    FunctionCode = (LIT, OPR, LOD, STO, CAL, INT, JMP, JPC);
```

```
{functions}
    Instruction = packed record
       func: FunctionCode; {功能码}
       level: 0..kNestingLayersMax; {相对层数}
       adr: 0..kAddrMax; {相对地址}
   end;
   {LIT 0,a: 取常数a
   OPR 0,a : 执行运算a
   LOD 1,a: 取层差为1的层、相对地址为a的变量
   STO 1,a: 存到层差为1的层、相对地址为a的变量
   CAL 1,a: 调用层差为1的过程
   INT O,a: t寄存器增加a
   JMP 0,a:转移到指令地址a处
    JPC 0,a : 条件转移到指令地址a处 }
var
   intermediate: text;
   stack data: text;
    curr char: char; {最近读到的字符}
    curr symbol : Symbol; {最近读到的符号}
   id: Identifier; {最近读到的标识符}
    curr ident: Identifier; {当前标识符的字符串}
   num : integer; {最近读到的数}
    char count : integer; {当前行的字符计数}
    line length : integer; {当前行的长度}
    error count : integer;
   code count : integer; {代码数组的当前下标}
   line : array [1..81] of char; {当前行}
   code: array [0..kInstructionsMax] of Instruction; {中间代码数组}
   words: array [1..kReservedWords] of Identifier; {存放保留字的字符串}
    words symbol : array [1..kReservedWords] of Symbol; {存放保留字的记号}
    ssym : array [char] of Symbol; {存放算符和标点符号的记号}
    mnemonic : array [FunctionCode] of string;
    {中间代码算符的字符串}
    declare symbols, stat begin symbols, factor begin symbols : SymbolS
et;
    table: array [0..kIdentsMax] of {符号表}
           record
              name : Identifier;
              case kind : ObjectType of
              kConstant : (val : integer);
```

```
kVariable, kProcedure : (level, adr : integer)
                 end;
     procedure ExitWithError(message: string);
     begin
         writeln('Fatal Error: ', message);
         halt;
     end;
    procedure error (n : integer);
     begin
         writeln('****', ' ' : char_count - 1, '^', n : 2);
         {当前行已读的字符数}
         error count := error count + 1;
         {错误数err加1}
         //halt;
    end {error};
85. procedure GetSymbol; {Lexical Analyzer}
     var i, j, k : integer;
    procedure GetChar; {取下一字符}
    begin
         if char count = line length then {如果cc指向行末}
         begin
             {如果已到文件尾}
             if eof(input) then ExitWithError('PROGRAM INCOMPLETE');
             {读新的一行}
             line length := 0;
             char count := 0;
             //writeln('char count reset');
             write(code count : 5, ' '); {code count : 5位数}
             while not eoln(input) do {如果不是行末}
             begin
                 line length := line length + 1;
                 read(curr char);
                 write(curr char);
                 line[line_length] := curr_char; {一次读一行入line}
             end;
```

```
writeln;
        line length := line length + 1;
        //writeln('line length: ',line length);
        read(line[line length]); {line[line length]中是行末符}
    char count := char count + 1;
    curr char:= line[char count]; {取line中下一个字符}
    //writeln('Getchar: ', ord(curr_char));
end {GetChar};
begin {GetSymbol}
    while curr_char in [' ', #13, #9, #10] do GetChar; {跳过无用空白}
    if curr char in ['a'..'z'] then
   begin {标识符或保留字}
       k := 0;
        repeat {处理字母开头的字母、数字串}
           if k < kIdentLengthMax then</pre>
           begin
               k := k + 1;
               curr ident[k] := curr char;
               //write(curr char);
           end;
           GetChar;
       until not(curr char in ['a'..'z', '0'..'9']);
       //writeln;
       {id中存放当前标识符或保留字的字符串}
       id := curr ident;
       curr ident := '';
       i := 0;
        j := kReservedWords + 1;
        {用二分查找法在保留字表中找当前的标识符id}
        repeat
           k := (i + j) div 2;
           if words[k] >= id then j := k
           else i := k
       until i + 1 >= j;
        {如果找到, 当前记号sym为保留字, 否则sym为标识符}
       if (j = kReservedWords + 1) or (words[j] \iff id) then
       begin
           curr symbol := IDENT;
           //writeln('find indent: ', id);
        end
```

```
else
    begin
        curr symbol := words symbol[j] ;
        //writeln('find reserved: ', id);
    end
end
else if curr char in ['0'...'9'] then
begin {数字}
   k := 0;
    num := 0;
    curr symbol := NUMBER; {当前记号sym为数字}
   repeat {计算数字串的值}
        num := 10*num + (ord(curr char)-ord('0'));
        k := k + 1;
       GetChar;
    until not(curr char in ['0'..'9']);
    {当前数字串的长度超过上界,则报告错误}
    if k > kNumLengthMax then error(30);
    //writeln('find number: ', num);{debug}
end
else if curr char= ':' then {处理赋值号}
begin
    GetChar;
    if curr char= '=' then
    begin
       curr symbol := BECOMES;
       GetChar
    end
    else
       curr symbol := NUL;
end
else if curr char = '<' then</pre>
begin
    GetChar;
    if curr char = '>' then {处理不等号}
    begin
       curr symbol := NEQ;
       GetChar;
    end
    else if curr char = '=' then {处理小于等于号}
    begin
       curr symbol := LEQ;
```

```
GetChar;
             end
             else curr symbol := LSS;
          end
          else if curr char = '>' then
          begin
             Getchar;
             if curr char = '=' then
             begin
                 curr symbol := GEQ;
                 GetChar;
             end
              else curr symbol := GTR;
         end
         else {处理其它算符或标点符号}
         begin
              //writeln('curr_symbol curr_char: ', ord(curr_char));
              curr symbol := ssym[curr char];
             GetChar;
          end;
221. end {GetSymbol};
     procedure GenerateCode(next func : FunctionCode; next level,
      next addr : integer);
     begin
          {如果当前指令序号>代码的最大长度}
          if code count > kInstructionsMax then ExitWithError('PROGRAM TOO
     LONG');
          with code[code count] do {生成一条新代码}
         begin
             func := next func; {功能码}
             level := next level; {层号}
             adr := next addr {地址}
          code count := code count + 1 {指令序号加1}
     end {GenerateCode};
      procedure Test(s1, s2 : SymbolSet; n : integer);
          {如果当前记号不属于集合S1,则报告错误n,跳过一些记号,直到当前记号属于S1US2}
      begin
```

```
if not (curr symbol in s1) then
    begin
        error(n);
        s1 := s1 + s2;
        while not (curr symbol in s1) do GetSymbol
end {Test};
 procedure Block(lev, table top: integer; symbol set: SymbolSet); {程
 序体}
var
    data top: integer; {本过程数据空间分配下标} {栈顶指针}
    symbol start : integer; {本过程标识表起始下标}
    code start : integer; {本过程代码起始下标}
procedure Enter(k : ObjectType);
begin {把obj填入符号表中}
    table top := table top + 1; {符号表指针加1}
    with table[table top] do{在符号表中增加新的一个条目}
    begin
        name := id; {当前标识符的名字}
        kind := k; {当前标识符的种类}
        case k of
           kConstant:
               begin {当前标识符是常数名}
                   if num > kAddrMax then {当前常数值大于上界,则出错}
                      error(30);
                      num := 0
                   end;
                   val := num
               end;
           kVariable :
               begin {当前标识符是变量名}
                   level := lev; {定义该变量的过程的嵌套层数}
                   adr := data top; {变量地址为当前过程数据空间栈顶}
                   data top := data top +1; {栈顶指针加1}
               end;
            kProcedure :
               level := lev {本过程的嵌套层数}
```

```
end
    end
 end {Enter};
function position(id : Identifier) : integer; {返回id在符号表的入口}
var
    i : integer;
begin
    {在标识符表中查标识符id}
    table[0].name := id; {在符号表栈的最下方预填标识符id}
    i := table top; {符号表栈顶指针}
   while table[i].name <> id do
        i := i - 1;
    {从符号表栈顶往下查标识符id}
    position := i {若查到,i为id的入口,否则i=0 }
end {position};
 procedure ConstDeclaration;
begin
     if curr symbol = IDENT then {当前记号是常数名}
    begin
        GetSymbol;
        if curr symbol in [EQL, BECOMES] then {当前记号是等号或赋值号}
        begin
            if curr symbol = BECOMES then error(1);
            {如果当前记号是赋值号,则出错}
            GetSymbol;
            if curr symbol = NUMBER then {等号后面是常数}
            begin
               Enter(kConstant); {将常数名加入符号表}
               GetSymbol
            end
            else error(2) {等号后面不是常数出错}
        else error(3) {标识符后不是等号或赋值号出错}
    end
     else error(4) {常数说明中没有常数名标识符}
end {ConstDeclaration};
 procedure VarDeclaration;
```

```
331. begin
         if curr symbol = IDENT then {如果当前记号是标识符}
             Enter(kVariable); {将该变量名加入符号表的下一条目}
             GetSymbol
         else error(4) {如果变量说明未出现标识符,则出错}
    end {VarDeclaration};
341. procedure ListCode;
     {列出本程序体生成的代码}
    var i : integer;
    begin
         {code_start: 本过程第一个代码的序号,cx-1: 本过程最后一个代码的序号}
         for i := code start to code count - 1 do
             with code[i] do {打印第i条代码}
                writeln(intermediate, i:3, mnemonic[func]:5, level : 3, adr
      : 5)//
        {i: 代码序号;
          mnemonic[f]: 功能码的字符串;
          1: 相对层号(层差);
          a: 相对地址或运算号码}
    end {ListCode};
    procedure Statement(symbol set : SymbolSet);
     var i, next node, next node 2 : integer;
    procedure Expression(symbol set : SymbolSet);
     var addop : Symbol;
    procedure Term(symbol set : SymbolSet);
     var mulop : Symbol;
     procedure Factor(symbol set : SymbolSet);
     var i : integer;
     begin
         Test (factor begin symbols, symbol set, 24);
         {测试当前的记号是否因子的开始符号, 否则出错, 跳过一些记号}
         while curr symbol in factor begin symbols do
            {如果当前的记号是否因子的开始符号}
         begin
             if curr symbol = IDENT then {当前记号是标识符}
             begin
```

```
i := position(id); {查符号表,返回id的入口}
           if i = 0 then
               error(11)
               {若在符号表中查不到id,则出错,否则,做以下工作}
           else
               with table[i] do
               case kind of
                   kConstant : GenerateCode(LIT, 0, val);
                      {若id是常数,生成指令,将常数val取到栈顶}
                  kVariable : GenerateCode(LOD, lev-level, adr);
                      {若id是变量,生成指令,将该变量取到栈顶;
                          lev: 当前语句所在过程的层号;
                         level: 定义该变量的过程层号;
                         adr: 变量在其过程的数据空间的相对地址}
                  kProcedure : error(21)
                      {若id是过程名,则出错}
               end;
           GetSymbol {取下一记号}
        end
        else if curr symbol = NUMBER then {当前记号是数字}
        begin
           if num > kAddrMax then {若数值越界,则出错}
           begin
               error(30);
               num := 0
           end;
           GenerateCode (LIT, 0, num); {生成一条指令, 将常数num取到栈顶}
           GetSymbol {取下一记号}
        end
        else if curr symbol = LPAREN then {如果当前记号是左括号}
        begin
           GetSymbol; {取下一记号}
           Expression([RPAREN]+symbol set); {处理表达式}
           if curr symbol = RPAREN then GetSymbol
           {如果当前记号是右括号,则取下一记号,否则出错}
           else error(22)
        end;
        Test (symbol set, [LPAREN], 23)
        {测试当前记号是否同步, 否则出错, 跳过一些记号}
    end {while}
end {Factor};
```

```
begin {Term}
     Factor(symbol set+[TIMES, SLASH]); {处理项中第一个因子}
    while curr symbol in [TIMES, SLASH] do
        {当前记号是"乘"或"除"号}
    begin
        mulop := curr symbol; {运算符存入mulop}
        GetSymbol; {取下一记号}
        Factor(symbol set+[TIMES, SLASH]); {处理一个因子}
        if mulop = TIMES then GenerateCode (OPR, 0, 4)
        {若mulop是"乘"号,生成一条乘法指令}
                        else GenerateCode (OPR, 0, 5)
        {否则, mulop是除号, 生成一条除法指令}
    end
end {Term};
begin {Expression}
    if curr symbol in [PLUS, MINUS] then {若第一个记号是加号或减号}
    begin
        addop := curr symbol; {"+"或"-"存入addop}
        GetSymbol;
        Term(symbol set+[PLUS, MINUS]); {处理一个项}
        if addop = MINUS then GenerateCode (OPR, 0, 1)
        {若第一个项前是负号,生成一条"负运算"指令}
    else Term(symbol set+[PLUS, MINUS]);
        {第一个记号不是加号或减号,则处理一个项}
    while curr symbol in [PLUS, MINUS] do {若当前记号是加号或减号}
    begin
        addop := curr symbol; {当前算符存入addop}
        GetSymbol; {取下一记号}
        Term(symbol set+[PLUS, MINUS]); {处理一个项}
        if addop = PLUS then GenerateCode(OPR, 0, 2)
        {若addop是加号,生成一条加法指令}
                   else GenerateCode (OPR, 0, 3)
        {否则, addop是减号, 生成一条减法指令}
     end
 end {Expression};
procedure Condition(symbol set : SymbolSet);
var relop : Symbol;
begin {Condition}
    if curr symbol = ODDSYM then {如果当前记号是"odd"}
```

```
begin
       GetSymbol; {取下一记号}
       Expression(symbol set); {处理算术表达式}
       GenerateCode (OPR, 0, 6) {生成指令,判定表达式的值是否为奇数,
       是,则取"真";不是,则取"假"}
   else {如果当前记号不是"odd"}
   begin
       Expression([EQL, NEQ, LSS, GTR, LEQ, GEQ] + symbol set);
       {外理算术表达式}
       if not (curr symbol in [EQL, NEQ, LSS, LEQ, GTR, GEQ]) then
       {如果当前记号不是关系符,则出错;否则,做以下工作}
          error(20)
       else
       begin
          relop := curr symbol; {关系符存入relop}
          GetSymbol; {取下一记号}
          Expression(symbol set); {处理关系符右边的算术表达式}
          case relop of
              EQL : GenerateCode (OPR, 0, 8);
                 {生成指令,判定两个表达式的值是否相等}
              NEQ : GenerateCode(OPR, 0, 9);
                 {生成指令, 判定两个表达式的值是否不等}
              LSS : GenerateCode (OPR, 0, 10);
                 {生成指令,判定前一表达式是否小干后一表达式}
              GEQ : GenerateCode(OPR, 0, 11);
                 {生成指令,判定前一表达式是否大干等干后一表达式}
              GTR : GenerateCode (OPR, 0, 12);
                 {生成指令,判定前一表达式是否大于后一表达式}
              LEQ: GenerateCode (OPR, 0, 13);
                 {生成指令,判定前一表达式是否小于等于后一表达式}
          end
       end
   end
end {Condition};
begin {Statement}
   if curr symbol = IDENT then {处理赋值语句}
   begin
       i := position(id); {在符号表中查id,返回id在符号表中的入口}
       if i = 0 then error(11) {若在符号表中查不到id,则出错}
       else if table[i].kind <> kVariable then {对非变量赋值,则出错}
       begin
          error (12);
```

```
i := 0;
       end;
       GetSymbol; {取下一记号}
       if curr symbol = BECOMES then GetSymbol else error(13);
       {若当前是赋值号,取下一记号,否则出错}
       Expression(symbol set); {处理表达式}
       if i <> 0 then {若赋值号左边的变量id有定义}
           with table[i] do GenerateCode(STO, lev-level, adr)
   end
   else if curr symbol = CALLSYM then {处理过程调用语句}
   begin
       GetSymbol; {取下一记号}
       if curr symbol <> IDENT then error(14) {下一记号不是标识符(过程名),
出错}
       else
       begin
           i := position(id); {查符号表,返回id在表中的位置}
           if i = 0 then error(11) {在符号表中查不到,出错}
           else
              with table[i] do
                  if kind = kProcedure then GenerateCode(CAL, lev-lev
el, adr)
                  {如果在符号表中id是过程名}
                  else error(15); {若id不是过程名,则出错}
          GetSymbol {取下一记号}
       end
   end
   else if curr symbol = IFSYM then {处理条件语句}
   begin
       GetSymbol; {取下一记号}
       Condition([THENSYM, DOSYM]+symbol set); {处理条件表达式}
       if curr symbol = THENSYM then GetSymbol else error(16);
       {如果当前记号是"then",则取下一记号;否则出错}
       next node := code count; {next node记录下一代码的地址}
       GenerateCode (JPC, 0, 0); {生成指令,表达式为"假"转到某地址(待填),
       否则顺序执行}
       Statement (symbol set); {处理一个语句}
       code[next node].adr := code count
       {将下一个指令的地址回填到上面的 ipc指令地址栏}
   end
   else if curr symbol = BEGINSYM then {处理语句序列}
   begin
```

```
GetSymbol;
       Statement([SEMICOLON, ENDSYM]+symbol set);
          {取下一记号,处理第一个语句}
       while curr symbol in [SEMICOLON]+stat begin symbols do
          {如果当前记号是分号或语句的开始符号,则做以下工作}
       begin
          if curr symbol = SEMICOLON then GetSymbol else error(10);
              {如果当前记号是分号,则取下一记号,否则出错}
          Statement([SEMICOLON, ENDSYM]+symbol set) {处理下一个语句}
       end;
       if curr symbol = ENDSYM then GetSymbol else error(17)
          {如果当前记号是"end",则取下一记号,否则出错}
   end
   else if curr symbol = WHILESYM then {处理循环语句}
   begin
       next node := code count; {next node记录下一指令地址,即条件表达式的
       第一条代码的地址}
       GetSymbol; {取下一记号}
       Condition([DOSYM]+symbol set); {处理条件表达式}
       next node 2 := code count; {记录下一指令的地址}
       GenerateCode(JPC, 0, 0); {生成一条指令,表达式为"假"转到某地
       址(待回填),否则顺序执行}
       if curr symbol = DOSYM then GetSymbol else error(18);
       {如果当前记号是"do",则取下一记号, 否则出错}
       Statement (symbol set); {处理"do"后面的语句}
       GenerateCode(JMP, 0, next node); {生成无条件转移指令, 转移到"while"
后的
       条件表达式的代码的第一条指令处}
       code[next node 2].adr := code count
       {把下一指令地址回填到前面生成的jpc指令的地址栏}
   end;
   Test(symbol set, [], 19)
       {测试下一记号是否正常, 否则出错, 跳过一些记号}
end {Statement};
begin {Block}
   data top := 3; {本过程数据空间栈顶指针}
   symbol start := table top; {标识符表的长度(当前指针)}
   table[table top].adr := code count; {本过程名的地址,即下一条指令的序号}
   GenerateCode (JMP, 0, 0); {生成一条转移指令}
   if lev > kNestingLayersMax then error(32);
       {如果当前过程层号>最大层数,则出错}
   repeat
```

```
if curr symbol = CONSTSYM then {处理常数说明语句}
begin
   GetSymbol;
   repeat
       ConstDeclaration; {处理一个常数说明}
       while curr symbol = COMMA do {如果当前记号是逗号}
       begin
          GetSymbol;
          ConstDeclaration
       end; {处理下一个常数说明}
       if curr symbol = SEMICOLON then GetSymbol else error(5)
       {如果当前记号是分号,则常数说明已处理完, 否则出错}
   until curr symbol <> IDENT
   {跳过一些记号,直到当前记号不是标识符(出错时才用到)}
end:
if curr symbol = VARSYM then {当前记号是变量说明语句开始符号}
begin
   GetSymbol;
   repeat
       VarDeclaration; {处理一个变量说明}
       while curr symbol = COMMA do {如果当前记号是逗号}
       begin
          GetSymbol;
          VarDeclaration
       end;
          {处理下一个变量说明}
       if curr symbol = SEMICOLON then GetSymbol else error(5)
          {如果当前记号是分号,则变量说明已处理完,否则出错}
   until curr symbol <> IDENT;
       {跳过一些记号,直到当前记号不是标识符(出错时才用到)}
end;
while curr symbol = PROCSYM do {处理过程说明}
begin
   GetSymbol;
   if curr symbol = IDENT then {如果当前记号是过程名}
   begin
       Enter(kProcedure);
       GetSymbol
   end {把过程名填入符号表}
   else error(4); {否则,缺少过程名出错}
   if curr symbol = SEMICOLON then GetSymbol else error(5);
       {当前记号是分号,则取下一记号,否则,过程名后漏掉分号出错}
```

```
Block(lev+1, table top, [SEMICOLON]+symbol set); {处理过程体}
              {lev+1: 过程嵌套层数加1; table top: 符号表当前栈顶指针,也是新
过程符号表起始位置; [SEMICOLON]+symbol set: 过程体开始和末尾符号集}
          if curr symbol = SEMICOLON then {如果当前记号是分号}
          begin
              GetSymbol; {取下一记号}
              Test(stat begin symbols+[IDENT, PROCSYM], symbol set, 6
                 {测试当前记号是否语句开始符号或过程说明开始符号,
                 否则报告错误6,并跳过一些记号}
          end
          else error(5) {如果当前记号不是分号,则出错}
       end:
       //writeln('Ha??');
       Test (stat begin symbols+[IDENT], declare symbols, 7)
          {检测当前记号是否语句开始符号, 否则出错, 并跳过一些记号}
   until not (curr symbol in declare symbols);
   {回到说明语句的处理(出错时才用),直到当前记号不是说明语句
   的开始符号
   code[table[symbol_start].adr].adr := code_count; {table[symbol_sta
rt].addr是本过程名的第1条
       代码 (JMP, 0, 0) 的地址,本语句即是将下一代码 (本过程语句的第
       1条代码)的地址回填到该jmp指令中,得(JMP, 0, code count)}
   with table[symbol start] do {本过程名的第1条代码的地址改为下一指令地址cx}
       adr := code count; {代码开始地址}
   end;
   code start := code count; {code start记录起始代码地址}
   GenerateCode(INT, 0, data top); {生成一条指令, 在栈顶为本过程留出数据空间
}
   Statement([SEMICOLON, ENDSYM]+symbol set); {处理一个语句}
   GenerateCode (OPR, 0, 0); {生成返回指令}
   Test(symbol set, [], 8); {测试过程体语句后的符号是否正常,否则出错}
   ListCode; {打印本过程的中间代码序列}
end {Block};
procedure Interpret;
const kStackSize = 500; {运行时数据空间(栈)的上界}
var pc, base, top: integer; {程序地址寄存器,基地址寄存器,栈顶地址寄存器}
```

```
683. i : Instruction; {指令寄存器}
          stack: array [1..kStackSize] of integer; {数据存储栈}
     function BaseOf(lev : integer) : integer;
    var b1 : integer;
    begin {BaseOf}
         b1 := base; {顺静态链求层差为lev的外层的基地址}
         while lev > 0 do
        begin
           b1 := stack[b1];
            lev := lev - 1
         end;
         BaseOf := b1
    end; {BaseOf}
    begin {Interpret}
        writeln('START PL/0');
         top := 0; {栈顶地址寄存器}
         base := 1; {基地址寄存器}
         pc := 0; {程序地址寄存器}
         stack[1] := 0;
        stack[2] := 0;
         stack[3] := 0;
            {最外层主程序数据空间栈最下面预留三个单元}
            {每个过程运行时的数据空间的前三个单元是:SL, DL, RA;
            SL: 指向本过程静态直接外层过程的SL单元;
            DL: 指向调用本过程的过程的最新数据空间的第一个单元;
            RA: 返回地址 }
        repeat
            i := code[pc]; {i取程序地址寄存器p指示的当前指令}
            pc := pc+1; {程序地址寄存器p加1,指向下一条指令}
            with i do
               case func of
                   LIT :
                       begin {当前指令是取常数指令(LIT, 0, a)}
                          top := top+1;
                          stack[top] := adr
                       end; {栈顶指针加1, 把常数a取到栈顶}
                   OPR :
                       case adr of {当前指令是运算指令(OPR, 0, a)}
                          0 : begin {a=0时,是返回调用过程指令}
                                 top := base-1; {恢复调用过程栈顶}
                                 pc := stack[top+3]; {程序地址寄存器p取返回地
     址}
```

```
base := stack[top+2];
                                   {基地址寄存器b指向调用过程的基地址}
                           end;
                       1 : stack[top] := -stack[top]; {一元负运算,栈顶元
素的值反号
                       2 : begin {加法}
                               top := top-1;
                               stack[top] := stack[top] + stack[top+1]
                           end;
                       3 : begin {减法}
                               top := top-1;
                               stack[top] := stack[top]-stack[top+1]
                           end;
                       4: begin {乘法}
                               top := top-1;
                               stack[top] := stack[top] * stack[top+1]
                           end;
                       5 : begin {整数除法}
                               top := top-1;
                               stack[top] := stack[top] div stack[top+1
1
                           end;
                       6 : stack[top] := ord(odd(stack[top])); {算s[top]
]是否奇数, 是则s[top]=1, 否则s[top]=0}
                       8 : begin
                               top := top-1;
                               stack[top] := ord(stack[top] = stack[top
+1])
                           end; {判两个表达式的值是否相等,
                               是则s[top]=1, 否则s[top]=0}
                       9: begin
                               top := top-1;
                               stack[top] := ord(stack[top] <> stack[to
p+1])
                           end; {判两个表达式的值是否不等,
                               是则s[top]=1, 否则s[top]=0}
                       10: begin
                               top := top-1;
                               stack[top] := ord(stack[top] < stack[top</pre>
+11)
                           end; {判前一表达式是否小于后一表达式,
                               是则s[top]=1, 否则s[top]=0}
```

```
11: begin
                            top := top-1;
                            stack[top] := ord(stack[top] >= stack[to
p+1])
                         end; {判前一表达式是否大于或等于后一表达式,
                            是则s[top]=1, 否则s[top]=0}
                     12: begin
                            top := top-1;
                            stack[top] := ord(stack[top] > stack[top
+11)
                        end; {判前一表达式是否大于后一表达式,
                            是则s[top]=1, 否则s[top]=0}
                     13: begin
                            top := top-1;
                            stack[top] := ord(stack[top] <= stack[to</pre>
p+1])
                        end; {判前一表达式是否小于或等于后一表达式,
                            是则s[top]=1, 否则s[top]=0}
                  end;
              LOD :
                 begin {当前指令是取变量指令(LOD, 1, a)}
                     top := top + 1;
                     stack[top] := stack[BaseOf(level) + adr]
                     {栈顶指针加1,根据静态链SL,将层差为1,相对地址
                     为a的变量值取到栈顶}
                 end;
              STO:
                 begin {当前指令是保存变量值(STO, 1, a)指令}
                     stack[BaseOf(level) + adr] := stack[top];
                     writeln(stack data, stack[top]);
                     {根据静态链SL,将栈顶的值存入层差为1,相对地址
                     为a的变量中 }
                     top := top-1 {栈顶指针减1}
                 end;
              CAL :
                 begin {当前指令是(CAL, 1, a)}
                         {为被调用过程数据空间建立连接数据}
                     stack[top+1] := BaseOf(level);
                         {根据层差1找到本过程的静态直接外层过程的数据空间
的SL单元,将其地址存入本过程新的数据空间的
                        SL单元}
                     stack[top+2] := base;
                     {调用过程的数据空间的起始地址存入本过程DL单元}
```

```
stack[top+3] := pc;
                       {调用过程cal指令的下一条的地址存入本过程RA单元}
                       base := top+1; {b指向被调用过程新的数据空间起始地址}
                       pc := adr {指令地址寄存储器指向被调用过程的地址a}
                   end;
                INT : top := top + adr;
                   {若当前指令是(INT, 0, a),则数据空间栈顶留出a大小的空间}
                JMP : pc := adr;
                   {若当前指令是(JMP, 0, a),则程序转到地址a执行}
                JPC :
                   begin {当前指令是(JPC, 0, a)}
                       if stack[top] = 0 then pc := adr;
                       {如果当前运算结果为"假"(0),程序转到地址a
                       执行,否则顺序执行}
                       top := top-1 {数据栈顶指针减1}
                   end
            end {with, case}
    until pc = 0;
        {程序一直执行到p取最外层主程序的返回地址0时为止}
    writeln('END PL/0');
end; {Interpret}
begin {主程序}
     assign(input, 'pl0_src.pas');
    reset (input);
    assign(intermediate, 'intermediate_code.txt');
    rewrite(intermediate);
    assign(stack data, 'stack data.txt');
    rewrite(stack data);
    for curr char:= 'a' to ';' do ssym[curr char] := NUL;
    {ASCII码的顺序}
    words[1] := 'begin';
    words[2] := 'call';
    words[3] := 'const';
    words[4] := 'do';
    words[5] := 'end';
    words[6] := 'if';
    words[7] := 'odd';
    words[8] := 'procedure';
    words[9] := 'then';
    words[10] := 'var';
    words[11] := 'while';
```

```
words symbol[1] := BEGINSYM; words symbol[2] := CALLSYM;
   words symbol[3] := CONSTSYM; words symbol[4] := DOSYM;
   words symbol[5] := ENDSYM;
                               words symbol[6] := IFSYM;
   words symbol[7] := ODDSYM;
                               words symbol[8] := PROCSYM;
   words symbol[9] := THENSYM;
                              words symbol[10] := VARSYM;
   words symbol[11] := WHILESYM;
   ssym['+'] := PLUS;
                          ssym['-'] := MINUS;
   ssym['*'] := TIMES;
                          ssym['/'] := SLASH;
   ssym['('] := LPAREN;
                          ssym[')'] := RPAREN;
   ssym['='] := EQL;
                          ssym[','] := COMMA;
   ssym['.'] := PERIOD;
   ssym['<'] := LSS;
                         ssym['>'] := GTR;
   ssym[';'] := SEMICOLON;
   {算符和标点符号的记号}
   mnemonic[LIT] := 'LIT';
                            mnemonic[OPR] := 'OPR';
   mnemonic[LOD] := 'LOD';
                           mnemonic[STO] := 'STO';
   mnemonic[CAL] := 'CAL';
                           mnemonic[INT] := 'INT';
   mnemonic[JMP] := 'JMP';
                           mnemonic[JPC] := 'JPC';
   {中间代码指令的字符串}
   declare symbols := [CONSTSYM, VARSYM, PROCSYM];
   {说明语句的开始符号}
   stat begin symbols := [BEGINSYM, CALLSYM, IFSYM, WHILESYM];
   {语句的开始符号}
   factor begin symbols := [IDENT, NUMBER, LPAREN];
   {因子的开始符号}
   error count := 0; {发现错误的个数}
   char count := 0; {当前行中输入字符的指针}
   code count := 0; {代码数组的当前指针}
   line length := 0; {输入当前行的长度}
   curr char:= ' '; {当前输入的字符}
   GetSymbol; {取下一个记号}
   Block(0, 0, [PERIOD] + declare symbols + stat begin symbols); {处理
程序体}
   if curr symbol <> PERIOD then error(9);
   {如果当前记号不是句号,则出错}
   if error count = 0 then Interpret
   {如果编译无错误,则解释执行中间代码}
   else writeln(error count, ' ERROR(S) IN PL/O PROGRAM');
```

```
896.
897. close(intermediate);
898. close(stack_data);
899. end.
```

2. PLO源程序代码

```
1. const m = 7, n = 85;
2. var x, y, z, q, r;
4. procedure multiply;
5. var a, b;
6. begin
       a := x;
        b := y;
       z := 0;
       while b > 0 do
       begin
           if odd b then z := z + a;
           a := 2*a ;
          b := b/2 ;
       end
16. end;
18. procedure divide;
19. var w;
20. begin
      r := x;
       q := 0;
       w := y;
       while w \le r do w := 2*w;
        while w > y do
      begin
          q := 2*q;
           w := w/2;
           if w <= r then
           begin
              r := r-w;
            q := q+1
           end
       end
35. end;
```

```
37. procedure gcd;
    var f, g;
39. begin
      f := x;
        g := y;
        while f <> g do
        begin
           if f < g then g := g-f;
           if g < f then f := f-g;
        end;
        z := f
48. end;
50. begin
      x := m;
        y := n;
        call multiply;
        x := 25;
        y := 3;
        call divide;
        x := 84;
        y := 36;
        call gcd;
60. end.
```

3. 中间代码

```
2 INT 0 5
3 LOD 1 3
```

4 STO 0 3

5 LOD 1 4

6 STO 0 4

7 LIT 0 0

8 STO 15

9 LOD 0 4

10 LIT 0 0

11 OPR 0 12

12 JPC 0 29

13 LOD 0 4

- 14 OPR 0 6
- 15 JPC 0 20
- 16 LOD 15
- 17 LOD 0 3
- 18 OPR 0 2
- 19 STO 15
- 20 LIT 0 2
- 21 LOD 0 3
- 22 OPR 04
- 23 STO 0 3
- 24 LOD 0 4
- 25 LIT 0 2
- 26 OPR 0 5
- 27 STO 0 4
- 28 JMP 0 9
- 29 OPR 0 0
- 31 INT 0 4
- 32 LOD 1 3
- 33 STO 1 7
- 34 LIT 0 0
- 35 STO 16
- 36 LOD 14
- 37 STO 0 3
- 38 LOD 0 3
- 39 LOD 1 7
- 40 OPR 0 13
- 41 JPC 0 47
- 42 LIT 0 2
- 43 LOD 0 3
- 44 OPR 0 4
- 45 STO 0 3
- 46 JMP 0 38
- 47 LOD 0 3

- 48 LOD 1 4
- 49 OPR 0 12
- 50 JPC 0 72
- 51 LIT 0 2
- 52 LOD 16
- 53 OPR 04
- 54 STO 16
- 55 LOD 0 3
- 56 LIT 0 2
- 57 OPR 0 5
- 58 STO 0 3
- 59 LOD 0 3
- 60 LOD 17
- 61 OPR 0 13
- 62 JPC 0 71
- 63 LOD 17
- 64 LOD 0 3
- 65 OPR 0 3
- 66 STO 17
- 67 LOD 16
- 68 LIT 0 1
- 69 OPR 0 2
- 70 STO 16
- 71 JMP 0 47
- 72 OPR 0 0
- 74 INT 0 5
- 75 LOD 1 3
- 76 STO 0 3
- 77 LOD 14
- 78 STO 0 4
- 79 LOD 0 3
- 80 LOD 0 4
- 81 OPR 09

- 82 JPC 0 100
- 83 LOD 0 3
- 84 LOD 0 4
- 85 OPR 0 10
- 86 JPC 0 91
- 87 LOD 0 4
- 88 LOD 0 3
- 89 OPR 0 3
- 90 STO 0 4
- 91 LOD 0 4
- 92 LOD 0 3
- 93 OPR 0 10
- 94 JPC 0 99
- 95 LOD 0 3
- 96 LOD 0 4
- 97 OPR 03
- 98 STO 0 3
- 99 JMP 0 79
- 100 LOD 0 3
- 101 STO 1 5
- 102 OPR 0 0
- 103 INT 0 8
- 104 LIT 0 7
- 105 STO 0 3
- 106 LIT 0 85
- 107 STO 04
- 108 CAL 0 2
- 109 LIT 0 25
- 110 STO 03
- 111 LIT 0 3
- 112 STO 04
- 113 CAL 0 31
- 114 LIT 0 84

115 STO 0 3

116 LIT 0 36

117 STO 0 4

118 CAL 0 74

119 OPR 0 0

4. 栈中的数据

第二部分

1. 编译程序源代码

```
1. program PL0;
2. {<mark>支持</mark>read, write<mark>函数的</mark>PL0编译程序}
```

```
3. const
         kReservedWords = 13; {保留字的个数}
         kIdentsMax = 100; {标识符表长度}
         kNumLengthMax = 14; {数字的最大位数}
         kIdentLengthMax = 10; {标识符的长度}
         kAddrMax = 2047; {最大地址}
         kNestingLayersMax = 3; {程序体嵌套的最大深度}
         kInstructionsMax = 200; {代码数组的大小}
         kDebugMessageOn = 1;
     type
         Symbol = (NUL, IDENT, NUMBER, PLUS, MINUS, TIMES, SLASH, ODDSYM,
             EQL, NEQ, LSS, LEQ, GTR, GEQ, LPAREN, RPAREN, COMMA, SEMICOLON,
             PERIOD, BECOMES, BEGINSYM, ENDSYM, IFSYM, THENSYM,
             WHILESYM, DOSYM, CALLSYM, CONSTSYM, VARSYM, PROCSYM, READSYM, W
     RITESYM);
         Identifier = packed array [1..kIdentLengthMax] of char;
         ObjectType = (kConstant, kVariable, kProcedure);
         SymbolSet = set of Symbol;
         FunctionCode = (LIT, OPR, LOD, STO, CAL, INT, JMP, JPC, RED, WRT);
     {functions}
         Instruction = packed record
             func: FunctionCode; {功能码}
             level: 0..kNestingLayersMax; {相对层数}
             adr : 0..kAddrMax; {相对地址}
         end;
         {LIT 0,a : 取常数a
         OPR 0,a : 执行运算a
         LOD 1,a: 取层差为1的层、相对地址为a的变量
         STO 1,a: 存到层差为1的层、相对地址为a的变量
         CAL 1,a: 调用层差为1的过程
         INT 0,a: t寄存器增加a
         JMP 0,a:转移到指令地址a处
         JPC 0,a:条件转移到指令地址a处 }
     var
         intermediate: text;
         stack data: text;
         pl0 input: text;
         curr char: char; {最近读到的字符}
         curr symbol : Symbol; {最近读到的符号}
         id: Identifier; {最近读到的标识符}
         curr ident : Identifier; {当前标识符的字符串}
         num : integer; {最近读到的数}
         char count : integer; {当前行的字符计数}
         line length : integer; {当前行的长度}
```

```
error count : integer;
     code count : integer; {代码数组的当前下标}
     line : array [1..81] of char; {当前行}
     code: array [0..kInstructionsMax] of Instruction; {中间代码数组}
    words: array [1..kReservedWords] of Identifier; {存放保留字的字符串}
     word symbol : array [1..kReservedWords] of Symbol; {存放保留字的记号}
     ssym : array [char] of Symbol; {存放算符和标点符号的记号}
     mnemonic : array [FunctionCode] of string;
     {中间代码算符的字符串}
     declare symbols, stat begin symbols, factor begin symbols : SymbolS
 et;
     table: array [0..kIdentsMax] of {符号表}
            record
                name : Identifier;
                case kind : ObjectType of
                kConstant : (val : integer);
                kVariable, kProcedure : (level, adr : integer)
            end;
 procedure ExitWithError(message: string);
 begin
    writeln('Fatal Error: ', message);
    halt;
 end;
 procedure error (n : integer);
 begin
    writeln('****', ' ' : char count - 1, '^', n : 2);
    {当前行已读的字符数}
    error count := error count + 1;
     {错误数err加1}
    //halt;
end {error};
 procedure GetSymbol; {Lexical Analyzer}
var i, j, k : integer;
 procedure GetChar; {取下一字符}
 begin
```

```
if char count = line length then {如果cc指向行末}
     begin
         {如果已到文件尾}
         if eof(input) then ExitWithError('PROGRAM INCOMPLETE');
         {读新的一行}
         line length := 0;
         char count := 0;
         //writeln('char count reset');
         write(code count : 5, ' '); {code count : 5位数}
         while not eoln(input) do {如果不是行末}
         begin
            line length := line length + 1;
            read(curr char);
             write(curr char);
             line[line length] := curr char; {一次读一行入line}
         end;
         writeln;
         line length := line length + 1;
         //writeln('line length: ',line length);
         read(line[line length]); {line[line length]中是行末符}
    end;
     char count := char count + 1;
     curr char:= line[char count]; {取line中下一个字符}
     //writeln('Getchar: ', ord(curr char));
end {GetChar};
begin {GetSymbol}
     while curr_char in [' ', #13, #9, #10] do GetChar; {跳过无用空白}
     if curr char in ['a'..'z'] then
     begin {标识符或保留字}
         k := 0;
         repeat {处理字母开头的字母、数字串}
             if k < kIdentLengthMax then</pre>
            begin
                k := k + 1;
                curr ident[k] := curr char;
             end;
             GetChar;
         until not(curr char in ['a'...'z', '0'...'9']);
         {id中存放当前标识符或保留字的字符串}
         id := curr ident;
```

```
curr ident := '';
    i := 0;
    j := kReservedWords + 1;
    {用二分查找法在保留字表中找当前的标识符id}
    repeat
       k := (i + j) div 2;
       if words[k] >= id then j := k
       else i := k
    until i + 1 >= j;
    {如果找到, 当前记号sym为保留字, 否则sym为标识符}
    if (j = kReservedWords + 1) or (words[j] <> id) then
   begin
       curr symbol := IDENT;
       //writeln('find indent: ', id);
    end
   else
   begin
       curr_symbol := word_symbol[j] ;
       //writeln('find reserved: ', id);
    end
end
else if curr char in ['0'..'9'] then
begin {数字}
   k := 0;
    num := 0;
    curr symbol := NUMBER; {当前记号sym为数字}
   repeat {计算数字串的值}
       num := 10*num + (ord(curr char)-ord('0'));
       k := k + 1;
       GetChar;
   until not(curr char in ['0'...'9']);
   {当前数字串的长度超过上界,则报告错误}
   if k > kNumLengthMax then error(30);
   //writeln('find number: ', num);{debug}
end
else if curr char= ':' then {处理赋值号}
begin
   GetChar;
    if curr char= '=' then
```

```
begin
                  curr symbol := BECOMES;
                  GetChar
              end
              else
               curr_symbol := NUL;
          end
          else if curr char = '<' then</pre>
          begin
              GetChar;
              if curr_char = '>' then {处理不等号}
              begin
                  curr symbol := NEQ;
                  GetChar;
              end
              else if curr char = '=' then {处理小于等于号}
              begin
                 curr symbol := LEQ;
                  GetChar;
              end
              else curr symbol := LSS;
          end
          else if curr char = '>' then
          begin
              Getchar;
              if curr char = '=' then
              begin
                  curr symbol := GEQ;
                  GetChar;
              end
              else curr symbol := GTR;
          end
         else {处理其它算符或标点符号}
          begin
              //writeln('curr symbol curr char: ', ord(curr char));
              curr symbol := ssym[curr char];
              GetChar;
          end;
220. end {GetSymbol};
223. procedure GenerateCode (next func : FunctionCode; next level,
      next addr : integer);
```

```
224. begin
          {如果当前指令序号>代码的最大长度}
          if code count > kInstructionsMax then ExitWithError('PROGRAM TOO
      LONG');
          with code[code count] do {生成一条新代码}
          begin
             func := next func; {功能码}
             level := next level; {层号}
             adr := next addr {地址}
         end;
         code count := code count + 1 {指令序号加1}
    end {GenerateCode};
      procedure Test(s1, s2 : SymbolSet; n : integer);
          {如果当前记号不属于集合S1,则报告错误n,跳过一些记号,直到当前记号属于S1US2}
     begin
          if not (curr symbol in s1) then
          begin
             error(n);
             s1 := s1 + s2;
             while not (curr symbol in s1) do GetSymbol
          end
     end {Test};
      procedure Block(lev, table top: integer; symbol set: SymbolSet); {程
      序体 }
    var
          data top: integer; {本过程数据空间分配下标} {栈顶指针}
          symbol start: integer; {本过程标识表起始下标}
          code start : integer; {本过程代码起始下标}
      procedure Enter(k : ObjectType);
     begin {把obj填入符号表中}
          table top := table top + 1; {符号表指针加1}
          with table[table top] do{在符号表中增加新的一个条目}
         begin
             name := id; {当前标识符的名字}
             kind := k; {当前标识符的种类}
             case k of
                 kConstant :
                    begin {当前标识符是常数名}
```

```
if num > kAddrMax then {当前常数值大于上界,则出错}
                        begin
                           error(30);
                           num := 0
                        end;
                       val := num
                    end;
                kVariable :
                    begin {当前标识符是变量名}
                        level := lev; {定义该变量的过程的嵌套层数}
                        adr := data top; {变量地址为当前过程数据空间栈顶}
                        data top := data top +1; {栈顶指针加1}
                    end:
                kProcedure :
                    level := lev {本过程的嵌套层数}
             end
         end
    end {Enter};
    function position(id : Identifier) : integer; {返回id在符号表的入口}
     var
         i : integer;
293. begin
         {在标识符表中查标识符id}
         table[0].name := id; {在符号表栈的最下方预填标识符id}
         i := table top; {符号表栈顶指针}
        while table[i].name <> id do
             i := i - 1;
         {从符号表栈顶往下查标识符id}
         position := i {若查到,i为id的入口,否则i=0 }
    end {position};
    procedure ConstDeclaration;
    begin
         if curr symbol = IDENT then {当前记号是常数名}
         begin
             GetSymbol;
             if curr symbol in [EQL, BECOMES] then {当前记号是等号或赋值号}
             begin
```

```
if curr symbol = BECOMES then error(1);
                {如果当前记号是赋值号,则出错}
                GetSymbol;
                if curr symbol = NUMBER then {等号后面是常数}
                begin
                   Enter(kConstant); {将常数名加入符号表}
                   GetSymbol
                end
                else error(2) {等号后面不是常数出错}
             end
             else error(3) {标识符后不是等号或赋值号出错}
         end
         else error(4) {常数说明中没有常数名标识符}
    end {ConstDeclaration};
    procedure VarDeclaration;
    begin
         if curr symbol = IDENT then {如果当前记号是标识符}
         begin
             Enter(kVariable); {将该变量名加入符号表的下一条目}
             GetSymbol
        end
        else error(4) {如果变量说明未出现标识符,则出错}
337. end {VarDeclaration};
    procedure ListCode;
341. {列出本程序体生成的代码}
     var i : integer;
343. begin
         {code start: 本过程第一个代码的序号,cx-1: 本过程最后一个代码的序号}
         for i := code start to code count - 1 do
            with code[i] do {打印第i条代码}
                writeln(intermediate, i:3, mnemonic[func]:5, level : 3, adr
     : 5)//
        {i: 代码序号;
          mnemonic[f]: 功能码的字符串;
         1: 相对层号(层差);
          a: 相对地址或运算号码}
352. end {ListCode};
    procedure Statement(symbol set : SymbolSet);
```

```
var i, next node, next node 2 : integer;
 procedure Expression(symbol set : SymbolSet);
 var addop : Symbol;
 procedure Term(symbol set : SymbolSet);
 var mulop : Symbol;
 procedure Factor(symbol set : SymbolSet);
var i : integer;
begin
     Test (factor begin symbols, symbol set, 24);
     {测试当前的记号是否因子的开始符号, 否则出错, 跳过一些记号}
    while curr symbol in factor begin symbols do
       {如果当前的记号是否因子的开始符号}
    begin
        if curr symbol = IDENT then {当前记号是标识符}
        begin
            i := position(id); {查符号表,返回id的入口}
            if i = 0 then
               error(11)
               {若在符号表中查不到id,则出错,否则,做以下工作}
            else
               with table[i] do
               case kind of
                   kConstant : GenerateCode(LIT, 0, val);
                       {若id是常数,生成指令,将常数val取到栈顶}
                   kVariable : GenerateCode(LOD, lev-level, adr);
                       {若id是变量,生成指令,将该变量取到栈顶;
                          lev: 当前语句所在过程的层号;
                          level: 定义该变量的过程层号;
                          adr: 变量在其过程的数据空间的相对地址}
                   kProcedure : error(21)
                       {若id是过程名,则出错}
               end;
            GetSymbol {取下一记号}
        end
        else if curr symbol = NUMBER then {当前记号是数字}
        begin
            if num > kAddrMax then {若数值越界,则出错}
            begin
               error(30);
               num := 0
            end;
```

```
GenerateCode (LIT, 0, num); {生成一条指令, 将常数num取到栈顶}
            GetSymbol {取下一记号}
        end
        else if curr symbol = LPAREN then {如果当前记号是左括号}
        begin
            GetSymbol; {取下一记号}
            Expression([RPAREN]+symbol set); {处理表达式}
            if curr symbol = RPAREN then GetSymbol
            {如果当前记号是右括号,则取下一记号,否则出错}
            else error(22)
        end;
        Test (symbol set, [LPAREN], 23)
        {测试当前记号是否同步, 否则出错, 跳过一些记号}
    end {while}
end {Factor};
begin {Term}
    Factor(symbol set+[TIMES, SLASH]); {处理项中第一个因子}
    while curr symbol in [TIMES, SLASH] do
        {当前记号是"乘"或"除"号}
    begin
        mulop := curr symbol; {运算符存入mulop}
        GetSymbol; {取下一记号}
        Factor(symbol set+[TIMES, SLASH]); {处理一个因子}
        if mulop = TIMES then GenerateCode (OPR, 0, 4)
        {若mulop是"乘"号,生成一条乘法指令}
                        else GenerateCode (OPR, 0, 5)
        {否则, mulop是除号, 生成一条除法指令}
    end
end {Term};
 begin {Expression}
    if curr symbol in [PLUS, MINUS] then {若第一个记号是加号或减号}
    begin
        addop := curr symbol; {"+"或"-"存入addop}
        GetSymbol;
        Term(symbol set+[PLUS, MINUS]); {处理一个项}
        if addop = MINUS then GenerateCode (OPR, 0, 1)
        {若第一个项前是负号,生成一条"负运算"指令}
    end
    else Term(symbol set+[PLUS, MINUS]);
        {第一个记号不是加号或减号,则处理一个项}
```

```
while curr_symbol in [PLUS, MINUS] do {若当前记号是加号或减号}
    begin
        addop := curr symbol; {当前算符存入addop}
        GetSymbol; {取下一记号}
        Term(symbol set+[PLUS, MINUS]); {处理一个项}
        if addop = PLUS then GenerateCode(OPR, 0, 2)
        {若addop是加号,生成一条加法指令}
                  else GenerateCode (OPR, 0, 3)
        {否则, addop是减号, 生成一条减法指令}
    end
end {Expression};
procedure Condition(symbol set : SymbolSet);
var relop : Symbol;
begin {Condition}
    if curr symbol = ODDSYM then {如果当前记号是"odd"}
    begin
        GetSymbol; {取下一记号}
        Expression(symbol set); {处理算术表达式}
        GenerateCode (OPR, 0, 6) {生成指令,判定表达式的值是否为奇数,
        是,则取"真";不是,则取"假"}
    end
    else {如果当前记号不是"odd"}
    begin
        Expression([EQL, NEQ, LSS, GTR, LEQ, GEQ] + symbol set);
        {处理算术表达式}
        if not (curr symbol in [EQL, NEQ, LSS, LEQ, GTR, GEQ]) then
        {如果当前记号不是关系符,则出错;否则,做以下工作}
           error(20)
        else
        begin
           relop := curr symbol; {关系符存入relop}
           GetSymbol; {取下一记号}
           Expression(symbol set); {处理关系符右边的算术表达式}
           case relop of
               EQL : GenerateCode (OPR, 0, 8);
                  {生成指令,判定两个表达式的值是否相等}
               NEQ : GenerateCode(OPR, 0, 9);
                   {生成指令,判定两个表达式的值是否不等}
               LSS : GenerateCode (OPR, 0, 10);
                   {生成指令,判定前一表达式是否小于后一表达式}
               GEQ : GenerateCode (OPR, 0, 11);
                   {生成指令,判定前一表达式是否大于等于后一表达式}
```

```
GTR : GenerateCode (OPR, 0, 12);
                   {生成指令,判定前一表达式是否大于后一表达式}
               LEQ : GenerateCode (OPR, 0, 13);
                   {生成指令,判定前一表达式是否小于等于后一表达式}
           end
        end
    end
end {Condition};
begin {Statement}
    if curr symbol = IDENT then {处理赋值语句}
    begin
        i := position(id); {在符号表中查id,返回id在符号表中的入口}
        if i = 0 then error(11) {若在符号表中查不到id,则出错}
        else if table[i].kind <> kVariable then {对非变量赋值,则出错}
        begin
           error (12);
           i := 0;
        end;
        GetSymbol; {取下一记号}
        if curr symbol = BECOMES then GetSymbol else error(13);
        {若当前是赋值号,取下一记号,否则出错}
        Expression(symbol set); {处理表达式}
        if i <> 0 then {若赋值号左边的变量id有定义}
           with table[i] do GenerateCode(STO, lev-level, adr)
    end
    else if curr symbol = CALLSYM then {处理过程调用语句}
    begin
        GetSymbol; {取下一记号}
        if curr symbol <> IDENT then error(14) {下一记号不是标识符(过程名),
 出错}
        else
        begin
           i := position(id); {查符号表,返回id在表中的位置}
           if i = 0 then error(11) {在符号表中查不到,出错}
           else
               with table[i] do
                  if kind = kProcedure then GenerateCode(CAL, lev-lev
el, adr)
                  {如果在符号表中id是过程名}
                   else error(15); {若id不是过程名,则出错}
```

```
GetSymbol {取下一记号}
       end
   end
   else if curr symbol = IFSYM then {处理条件语句}
   begin
       GetSymbol; {取下一记号}
       Condition([THENSYM, DOSYM]+symbol set); {处理条件表达式}
       if curr symbol = THENSYM then GetSymbol else error(16);
       {如果当前记号是"then",则取下一记号; 否则出错}
       next node := code count; {next node记录下一代码的地址}
       GenerateCode(JPC, 0, 0); {生成指令,表达式为"假"转到某地址(待填),
       否则顺序执行}
       Statement (symbol set); {处理一个语句}
       code[next node].adr := code count
       {将下一个指令的地址回填到上面的 jpc指令地址栏}
   end
   else if curr symbol = BEGINSYM then {处理语句序列}
   begin
       GetSymbol;
       Statement([SEMICOLON, ENDSYM]+symbol set);
          {取下一记号,处理第一个语句}
       while curr symbol in [SEMICOLON]+stat begin symbols do
          {如果当前记号是分号或语句的开始符号,则做以下工作}
       begin
          if curr symbol = SEMICOLON then GetSymbol else error(10);
              {如果当前记号是分号,则取下一记号,否则出错}
          Statement([SEMICOLON, ENDSYM]+symbol set) {处理下一个语句}
       end;
       if curr symbol = ENDSYM then GetSymbol else error(17)
           {如果当前记号是"end",则取下一记号,否则出错}
   else if curr symbol = WHILESYM then {处理循环语句}
   begin
       next node := code count; {next node记录下一指令地址,即条件表达式的
       第一条代码的地址}
       GetSymbol; {取下一记号}
       Condition([DOSYM]+symbol set); {处理条件表达式}
       next node 2 := code count; {记录下一指令的地址}
       GenerateCode(JPC, 0, 0); {生成一条指令,表达式为"假"转到某地
       址(待回填),否则顺序执行}
       if curr symbol = DOSYM then GetSymbol else error(18);
       {如果当前记号是"do",则取下一记号, 否则出错}
       Statement (symbol set); {处理"do"后面的语句}
       GenerateCode(JMP, 0, next node); {生成无条件转移指令, 转移到"while"
后的
```

```
条件表达式的代码的第一条指令处}
        code[next node 2].adr := code count
        {把下一指令地址回填到前面生成的jpc指令的地址栏}
    end
    else if curr symbol = READSYM then {处理读入语句}
    begin
       GetSymbol;
        if curr symbol <> LPAREN then error(10)
       else
       begin
           repeat
               GetSymbol;
               if curr symbol <> IDENT then error(4)
               else
               begin
                   i := position(id);
                                      {在符号表中查id,返回id在符号表中的
入口}
                   if i = 0 then error(11) {若在符号表中查不到id,则出错}
                   else if table[i].kind <> kVariable then error(12) {
对非变量赋值,则出错}
                   else
                       with table[i] do GenerateCode(RED, lev-level, ad
r);
               end;
               GetSymbol;
           until curr symbol <> COMMA;
            if curr symbol <> RPAREN then error(22);
           GetSymbol;
        end
    end
    else if curr_symbol = WRITESYM then {处理输出语句}
    begin
        GetSymbol;
        if curr symbol <> LPAREN then error(10)
       else
       begin
           repeat
               GetSymbol;
               Expression([RPAREN, COMMA] + symbol set);
               GenerateCode (WRT, 0, 0);
           until curr symbol <> COMMA;
           if curr symbol <> RPAREN then error(22);
           GetSymbol;
        end
    end;
```

```
Test (symbol set, [], 19)
       {测试下一记号是否正常, 否则出错, 跳过一些记号}
end {Statement};
begin {Block}
   data top := 3; {本过程数据空间栈顶指针}
   symbol start := table top; {标识符表的长度(当前指针)}
   table [table top].adr := code count; {本过程名的地址,即下一条指令的序号}
   GenerateCode (JMP, 0, 0); {生成一条转移指令}
   if lev > kNestingLayersMax then error(32);
       {如果当前过程层号>最大层数,则出错}
   repeat
       if curr symbol = CONSTSYM then {处理常数说明语句}
       begin
          GetSymbol;
          repeat
              ConstDeclaration; {处理一个常数说明}
              while curr symbol = COMMA do {如果当前记号是逗号}
              begin
                 GetSymbol;
                 ConstDeclaration
              end; {处理下一个常数说明}
              if curr symbol = SEMICOLON then GetSymbol else error(5)
              {如果当前记号是分号,则常数说明已处理完, 否则出错}
          until curr symbol <> IDENT
          {跳过一些记号,直到当前记号不是标识符(出错时才用到)}
       end;
       if curr symbol = VARSYM then {当前记号是变量说明语句开始符号}
       begin
          GetSymbol;
          repeat
              VarDeclaration; {处理一个变量说明}
              while curr symbol = COMMA do {如果当前记号是逗号}
              begin
                 GetSymbol;
                 VarDeclaration
              end;
                 {处理下一个变量说明}
              if curr symbol = SEMICOLON then GetSymbol else error(5)
                 {如果当前记号是分号,则变量说明已处理完,否则出错}
          until curr symbol <> IDENT;
              {跳过一些记号,直到当前记号不是标识符(出错时才用到)}
```

```
end;
       while curr symbol = PROCSYM do {处理过程说明}
      begin
          GetSymbol;
          if curr symbol = IDENT then {如果当前记号是过程名}
          begin
             Enter(kProcedure);
             GetSymbol
          end {把过程名填入符号表}
          else error(4); {否则,缺少过程名出错}
          if curr symbol = SEMICOLON then GetSymbol else error(5);
              {当前记号是分号,则取下一记号,否则,过程名后漏掉分号出错}
          Block(lev+1, table top, [SEMICOLON]+symbol set); {处理过程体}
             {lev+1: 过程嵌套层数加1; table top: 符号表当前栈顶指针,也是新
过程符号表起始位置; [SEMICOLON]+symbol set: 过程体开始和末尾符号集}
          if curr symbol = SEMICOLON then {如果当前记号是分号}
          begin
             GetSymbol; {取下一记号}
             Test(stat begin symbols+[IDENT, PROCSYM], symbol set, 6
                 {测试当前记号是否语句开始符号或过程说明开始符号,
                 否则报告错误6,并跳过一些记号}
          end
          else error(5) {如果当前记号不是分号,则出错}
      end;
      //writeln('Ha??');
      Test(stat begin symbols+[IDENT], declare symbols, 7)
          {检测当前记号是否语句开始符号, 否则出错, 并跳过一些记号}
   until not (curr symbol in declare symbols);
   {回到说明语句的处理(出错时才用),直到当前记号不是说明语句
   的开始符号
   code[table[symbol start].adr].adr := code count; {table[symbol sta
rt].addr是本过程名的第1条
      代码(JMP, 0, 0)的地址,本语句即是将下一代码(本过程语句的第
      1条代码)的地址回填到该jmp指令中,得(JMP, 0, code count)}
   with table[symbol start] do {本过程名的第1条代码的地址改为下一指令地址cx}
   begin
       adr := code count; {代码开始地址}
   end;
```

```
code start := code count; {code start记录起始代码地址}
     GenerateCode(INT, 0, data top); {生成一条指令, 在栈顶为本过程留出数据空间
     Statement([SEMICOLON, ENDSYM]+symbol set); {处理一个语句}
     GenerateCode (OPR, 0, 0); {生成返回指令}
     Test(symbol set, [], 8); {测试过程体语句后的符号是否正常,否则出错}
     ListCode; {打印本过程的中间代码序列}
end {Block};
 procedure Interpret;
const kStackSize = 500; {运行时数据空间(栈)的上界}
var pc, base, top : integer; {程序地址寄存器,基地址寄存器,栈顶地址寄存器}
      i: Instruction; {指令寄存器}
      stack: array [1..kStackSize] of integer; {数据存储栈}
function BaseOf(lev : integer) : integer;
var b1 : integer;
begin {BaseOf}
     b1 := base; {顺静态链求层差为lev的外层的基地址}
     while lev > 0 do
     begin
        b1 := stack[b1];
        lev := lev - 1
     end;
     BaseOf := b1
end; {BaseOf}
begin {Interpret}
     writeln('START PL/0');
     top := 0; {栈顶地址寄存器}
    base := 1; {基地址寄存器}
    pc := 0; {程序地址寄存器}
     stack[1] := 0;
    stack[2] := 0;
    stack[3] := 0;
        {最外层主程序数据空间栈最下面预留三个单元}
        {每个过程运行时的数据空间的前三个单元是:SL, DL, RA;
        SL: 指向本过程静态直接外层过程的SL单元;
        DL: 指向调用本过程的过程的最新数据空间的第一个单元;
        RA: 返回地址 }
    repeat
        i := code[pc]; {i取程序地址寄存器p指示的当前指令}
        pc := pc+1; {程序地址寄存器p加1,指向下一条指令}
```

```
with i do
           case func of
               LIT :
                   begin {当前指令是取常数指令(LIT, 0, a)}
                       top := top+1;
                      stack[top] := adr
                   end; {栈顶指针加1,把常数a取到栈顶}
               OPR :
                   case adr of {当前指令是运算指令(OPR, 0, a)}
                      0 : begin {a=0时,是返回调用过程指令}
                              top := base-1; {恢复调用过程栈顶}
                              pc := stack[top+3]; {程序地址寄存器p取返回地
址}
                              base := stack[top+2];
                                  {基地址寄存器b指向调用过程的基地址}
                          end;
                      1 : stack[top] := -stack[top]; {一元负运算, 栈顶元
素的值反号
                      2: begin {加法}
                              top := top-1;
                              stack[top] := stack[top] + stack[top+1]
                          end;
                      3 : begin {减法}
                              top := top-1;
                              stack[top] := stack[top]-stack[top+1]
                          end;
                      4 : begin {乘法}
                              top := top-1;
                              stack[top] := stack[top] * stack[top+1]
                          end;
                      5 : begin {整数除法}
                              top := top-1;
                              stack[top] := stack[top] div stack[top+1
                          end;
                       6 : stack[top] := ord(odd(stack[top])); {算s[top
] 是否奇数, 是则s[top]=1, 否则s[top]=0}
                      8 : begin
                              top := top-1;
                              stack[top] := ord(stack[top] = stack[top
+1])
                          end; {判两个表达式的值是否相等,
                              是则s[top]=1, 否则s[top]=0}
```

```
9: begin
                              top := top-1;
                              stack[top] := ord(stack[top] <> stack[to
p+1])
                          end; {判两个表达式的值是否不等,
                              是则s[top]=1, 否则s[top]=0}
                      10: begin
                              top := top-1;
                              stack[top] := ord(stack[top] < stack[top</pre>
+1])
                          end; {判前一表达式是否小于后一表达式,
                              是则s[top]=1, 否则s[top]=0}
                      11: begin
                              top := top-1;
                              stack[top] := ord(stack[top] >= stack[to
p+1])
                          end; {判前一表达式是否大于或等于后一表达式,
                              是则s[top]=1, 否则s[top]=0}
                      12: begin
                              top := top-1;
                              stack[top] := ord(stack[top] > stack[top
+11)
                          end; {判前一表达式是否大于后一表达式,
                              是则s[top]=1, 否则s[top]=0}
                      13: begin
                              top := top-1;
                              stack[top] := ord(stack[top] <= stack[to</pre>
p+1])
                          end; {判前一表达式是否小于或等于后一表达式,
                              是则s[top]=1, 否则s[top]=0}
                   end;
               LOD :
                  begin {当前指令是取变量指令(LOD, 1, a)}
                      top := top + 1;
                      stack[top] := stack[BaseOf(level) + adr]
                      {栈顶指针加1,根据静态链SL,将层差为1,相对地址
                      为a的变量值取到栈顶
                  end;
               STO:
                   begin {当前指令是保存变量值(STO, 1, a)指令}
                      stack[BaseOf(level) + adr] := stack[top];
```

```
writeln(stack data, stack[top]);
                    {根据静态链SL,将栈顶的值存入层差为1,相对地址
                    为a的变量中 }
                    top := top - 1 {栈顶指针减1}
                 end;
              CAL :
                 begin {当前指令是(CAL, l, a)}
                        {为被调用过程数据空间建立连接数据}
                    stack[top+1] := BaseOf(level);
                        {根据层差1找到本过程的静态直接外层过程的数据空间
的SL单元,将其地址存入本过程新的数据空间的
                        SL单元}
                    stack[top+2] := base;
                    {调用过程的数据空间的起始地址存入本过程DL单元}
                    stack[top+3] := pc;
                    {调用过程cal指令的下一条的地址存入本过程RA单元}
                    base := top+1; {b指向被调用过程新的数据空间起始地址}
                    pc := adr {指令地址寄存储器指向被调用过程的地址a}
                 end;
             INT : top := top + adr;
                 {若当前指令是(INT, 0, a),则数据空间栈顶留出a大小的空间}
              JMP : pc := adr;
                 {若当前指令是(JMP, 0, a),则程序转到地址a执行}
              JPC :
                 begin {当前指令是(JPC, 0, a)}
                    if stack[top] = 0 then pc := adr;
                    {如果当前运算结果为"假"(0), 程序转到地址a
                    执行, 否则顺序执行}
                    top := top-1 {数据栈顶指针减1}
                 end;
             RED :
                 begin
                    read(pl0 input, stack[BaseOf(level)+adr]);
                 end;
             WRT:
                 begin
                    writeln(stack[top]);
                    top := top - 1
                 end
          end {with, case}
   until pc = 0;
       {程序一直执行到p取最外层主程序的返回地址0时为止}
   writeln('END PL/0');
end; {Interpret}
```

```
begin {主程序}
    assign(input, 'pl0 src.pas');
    reset(input);
    assign(intermediate, 'intermediate code2.txt');
    rewrite (intermediate);
    assign(stack data, 'stack_data2.txt');
    rewrite(stack data);
    assign(pl0 input, 'input.txt');
    reset (pl0 input);
    for curr char:= 'a' to ';' do ssym[curr char] := NUL;
    {ASCII码的顺序}
    words[1] := 'begin';
    words[2] := 'call';
    words[3] := 'const';
    words[4] := 'do';
    words[5] := 'end';
    words[6] := 'if';
    words[7] := 'odd';
    words[8] := 'procedure';
    words[9] := 'read';
    words[10] := 'then';
    words[11] := 'var';
    words[12] := 'while';
    words[13] := 'write';
    word symbol[1] := BEGINSYM;
    word symbol[2] := CALLSYM;
    word symbol[3] := CONSTSYM;
    word symbol[4] := DOSYM;
    word symbol[5] := ENDSYM;
    word symbol[6] := IFSYM;
    word symbol[7] := ODDSYM;
    word symbol[8] := PROCSYM;
    word symbol[9] := READSYM;
    word symbol[10] := THENSYM;
    word symbol[11] := VARSYM;
    word symbol[12] := WHILESYM;
    word symbol[13] := WRITESYM;
    ssym['+'] := PLUS;
                            ssym['-'] := MINUS;
    ssym['*'] := TIMES;
                            ssym['/'] := SLASH;
    ssym['('] := LPAREN;
                            ssym[')'] := RPAREN;
    ssym['='] := EQL;
                            ssym[','] := COMMA;
```

```
ssym['.'] := PERIOD;
    ssym['<'] := LSS;
                        ssym['>'] := GTR;
    ssym[';'] := SEMICOLON;
    {算符和标点符号的记号}
    mnemonic[LIT] := 'LIT';
                            mnemonic[OPR] := 'OPR';
    mnemonic[CAL] := 'CAL';
                           mnemonic[INT] := 'INT';
    mnemonic[JMP] := 'JMP';
                           mnemonic[JPC] := 'JPC';
    {中间代码指令的字符串}
    declare symbols := [CONSTSYM, VARSYM, PROCSYM];
    {说明语句的开始符号}
    stat begin symbols := [BEGINSYM, CALLSYM, IFSYM, WHILESYM, READSYM,
 WRITESYM];
    {语句的开始符号}
    factor begin symbols := [IDENT, NUMBER, LPAREN];
    {因子的开始符号}
    error count := 0; {发现错误的个数}
    char count := 0; { 当前行中输入字符的指针 }
    code count := 0; {代码数组的当前指针}
    line length := 0; {输入当前行的长度}
    curr char:= ' '; {当前输入的字符}
    GetSymbol; {取下一个记号}
    Block(0, 0, [PERIOD] + declare symbols + stat begin symbols); {处理
 程序体 }
    if curr symbol <> PERIOD then error(9);
    {如果当前记号不是句号,则出错}
    if error count = 0 then Interpret
    {如果编译无错误,则解释执行中间代码}
    else writeln(error count, ' ERROR(S) IN PL/0 PROGRAM');
    close(intermediate);
    close(stack data);
end.
```

2. PLO源程序代码

这是一个输入x和y,输出x除以y的商和余数的程序。

```
1. var x, y, q, r;
2. procedure divide;
   var w;
    begin
       r := x;
        q := 0;
        w := y;
       while w \le r do w := 2 * w;
        while w > y do
        begin
         q := 2 * q;
           w := w / 2;
           if w <= r then
           begin
              r := r - w;
               q := q + 1
           end
       end
19. end;
20. begin
     read(x); read(y);
     call divide;
     write(q);
     write(r);
25. end.
```

3. 中间代码

```
2 INT 0 4
3 LOD 1 3
4 STO 1 6
5 LIT 0 0
6 STO 1 5
7 LOD 1 4
8 STO 0 3
9 LOD 0 3
10 LOD 1 6
```

- 11 OPR 0 13
- 12 JPC 0 18
- 13 LIT 0 2
- 14 LOD 0 3
- 15 OPR 04
- 16 STO 0 3
- 17 JMP 0 9
- 18 LOD 0 3
- 19 LOD 1 4
- 20 OPR 0 12
- 21 JPC 0 43
- 22 LIT 0 2
- 23 LOD 1 5
- 24 OPR 0 4
- 25 STO 1 5
- 26 LOD 0 3
- 27 LIT 0 2
- 28 OPR 0 5
- 29 STO 0 3
- 30 LOD 0 3
- 31 LOD 16
- 32 OPR 0 13
- 33 JPC 0 42
- 34 LOD 1 6
- 35 LOD 0 3
- 36 OPR 0 3
- 37 STO 16
- 38 LOD 15
- 39 LIT 0 1
- 40 OPR 0 2
- 41 STO 15
- 42 JMP 0 18
- 43 OPR 0 0

44 INT 0 7

45 0 3

46 0 4

47 CAL 0 2

48 LOD 0 5

49 0 0

50 LOD 0 6

5100

52 OPR 0 0

4. 输入输出的数据

输入:298

输出:35