

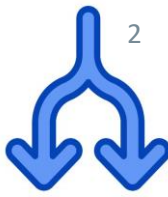


Practical Concurrent and Parallel Programming I

Intro to Concurrency and Mutual Exclusion

Raúl Pardo, Jørgen Staunstrup

Agenda



- Course General Info
- Introduction to Concurrency
- Java Threads
- Mutual Exclusion
- Java Locks

- **Course manager: Raúl Pardo**

- PhD from Chalmers University 2017
- Postdoc at Inria 2017 & ITU 2019
- Research interest: Privacy & Security, Formal Methods, Probabilistic Reasoning, Concurrency



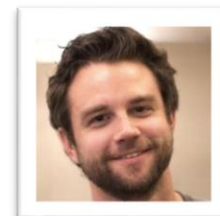
- **Co-teacher: Jørgen Staunstrup**

- He will introduce himself shortly

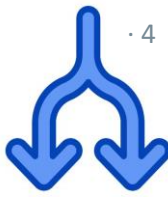


- **Teaching Assistants**

- Holger Stadel Borum
- Amund Ranheim Lome
- You will meet them after the lecture



Standard weekly plan



- Lectures
 - Mondays 10-12
 - We publish the readings for the lecture a week before the lecture (approx.)
- Exercise sessions
 - Mondays 12-14
 - Every week we publish a set of exercises covering the material of the lecture



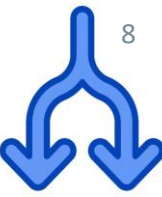
- The learnIT website (<https://learnit.itu.dk/course/view.php?id=3020335>) contains information about the course and submission links for the assignments
- The material of the course is hosted in our github repository (<https://github.itu.dk/jst/PCPP2021-public>)
 - Lecture slides
 - Readings
 - Example code
 - Exercises
 - Accompanying code for the exercises



- You are expected to work on **groups of 2/3 people**
 - Today in the exercise session we will start by forming groups with the help of TAs
- **Solutions are submitted bi-weekly**
 - Every two weeks you submit the solution to the two exercise sets of the two previous weeks
- Assignments submissions **on Mondays before the lecture**
 - We made sure to minimize clashing with other courses in the CS programme
 - An assignment is hand-in including solutions to two exercise sets
- **Oral Feedback and Assessment**
 - Feedback and assessment of assignments if provided in oral sessions
 - You must **book an oral feedback slot with a TA/teacher using the scheduler in learnIT**
 - Today in the exercise session you can already book a slot (<https://learnit.itu.dk/mod/scheduler/view.php?id=143279>)
- **In total there are 12 exercise sets, you must pass 10 to be entitled to take the exam**
- Exercises are divided into mandatory and challenging
 - Doing the challenging exercises will increase your chances of getting high marks in the exam
 - Challenging exercises are optional

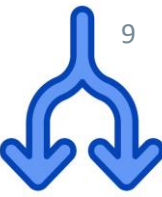


- The main channel of communication is the Questions and Answers Forum in LearnIT
 - <https://learnit.itu.dk/mod/hsuforum/view.php?id=143401>
 - **We strongly encourage you to use the forum!**
 - We will check the forum regularly
 - But, we also encourage you to answers questions yourself!
 - The forum is meant to be a discussion platform to boost learning and share knowledge
- The only constraint: *do not directly post solutions to exercises*



- Oral exam
 - We will ask questions covering all the content in the readings and lecture slides

About Jørgen



Datalog/Computer scientist Aarhus University: 1975

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Taught first course on Concurrency: 1978 (USC, Los Angeles)

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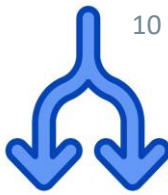
It was an on-line course !!

Joined ITU in 2001, retired in 2014

Teaching MSc course Mobile App Development at ITU since 2016

Concurrency

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It takes one person 2 hours to dig 2 meters of ditch.
How long will it take 2 people to dig 4 meters of ditch?

What if they only have one shovel?

What if they have to dig a hole (and not a ditch)?

How fast can a 100 persons dig 1 meter of ditch?

...

The first computers were sequential



In my first years at university, we used the Danish Gier (1965)

The first computers were sequential



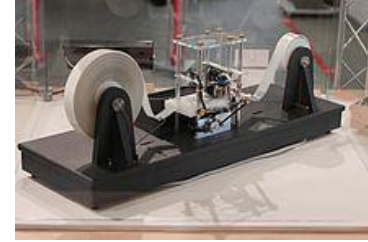
The Turing machine (1936) - a mathematical model

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The Turing machine (1936) - a mathematical model



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11



The Turing machine (1936) - a mathematical model

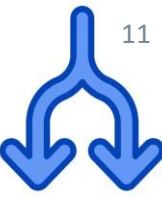


Eniac (1945)

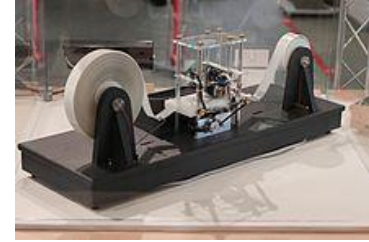


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Transistormaskinen GIER, 1961

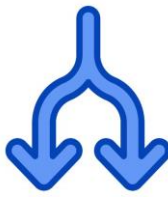
Timesharing - my first programs (1969)



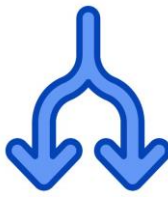
Via a number of terminals several users shared the computer

My first (and best) question:

What happens if two users print simultaneously?



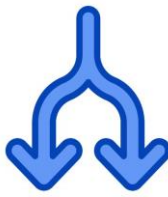
Concurrency



```
public class RemoveDuplicateInArrayExample{
    public static int removeDuplicateElements(int arr[], int n){
        if (n==0 || n==1){ return n; }
        int[] temp = new int[n];
        int j = 0;
        for (int i=0; i<n-1; i++){
            if (arr[i] != arr[i+1]){
                temp[j++] = arr[i];
            }
        }
        temp[j++] = arr[n-1];
        // Changing original array
        for (int i=0; i<j; i++){
            arr[i] = temp[i];
        }
        return j;
    }

    public static void main (String[] args) {
        int arr[] = {10,20,20,30,30,40,50,50};
        int length = arr.length;
        length = removeDuplicateElements(arr, length);
        //printing array elements
        for (int i=0; i<length; i++) System.out.print(arr[i]+" ");
    }
}
```

Concurrency



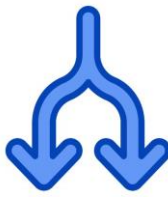
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}
```

Single stream

```
-----
----
-----
-----
---
```

Concurrency



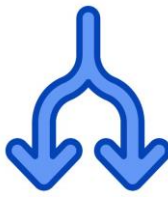
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```



Single stream

Concurrency



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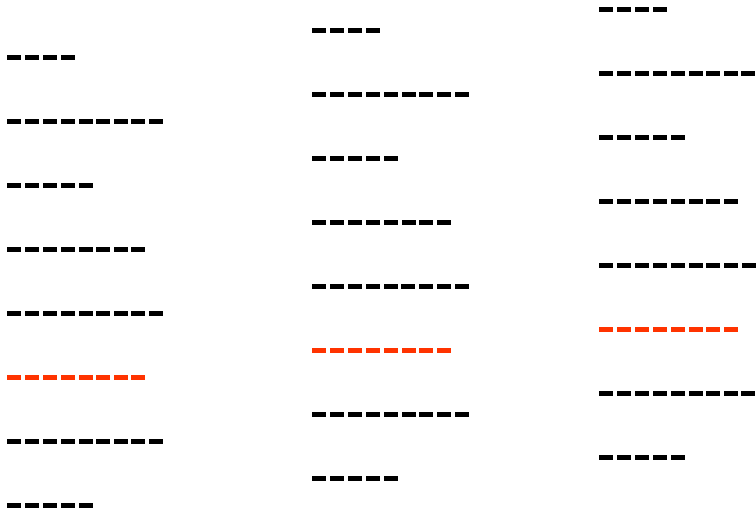
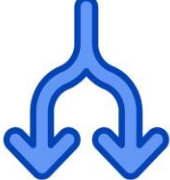
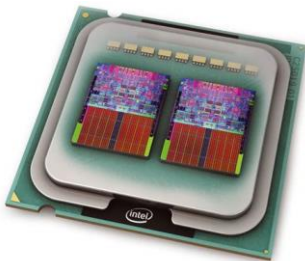
Single stream



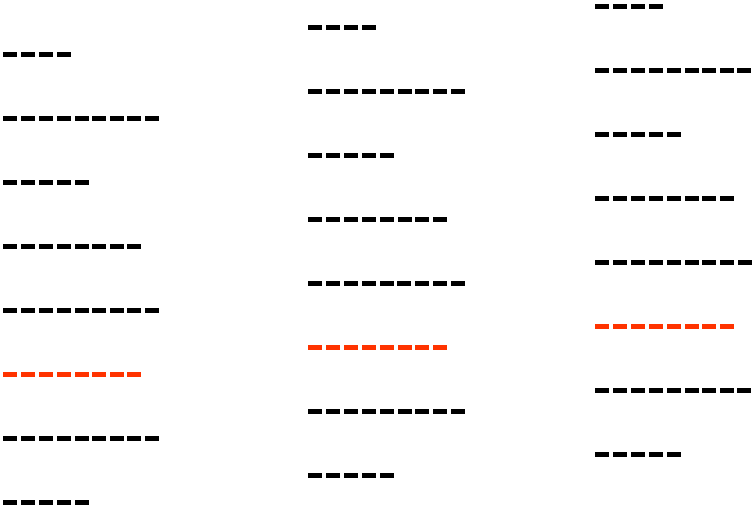
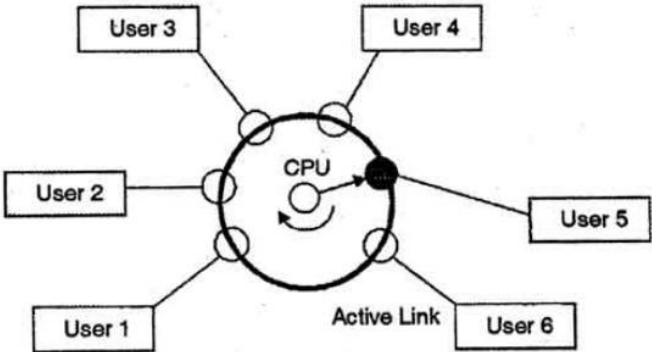
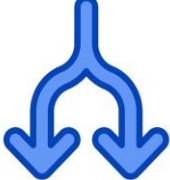
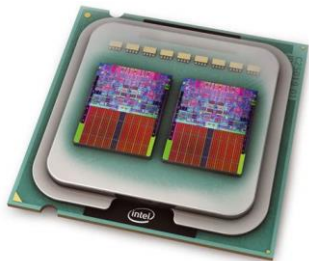
Concurrent (multiple streams)

-----	-----	-----
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Multiple streams

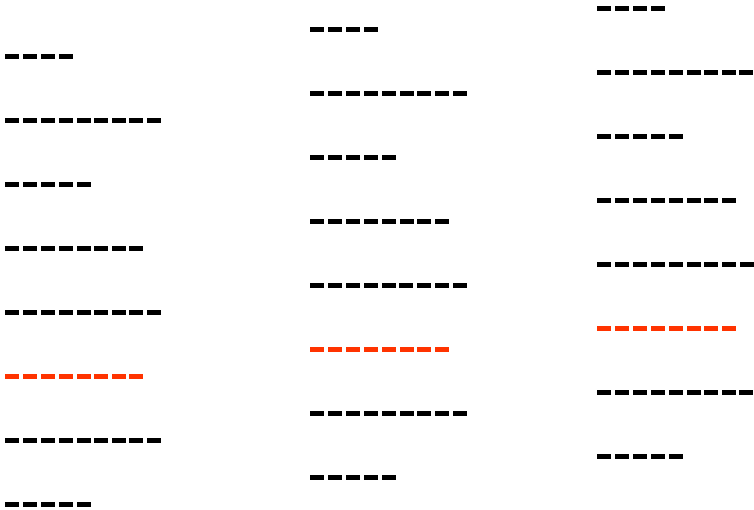
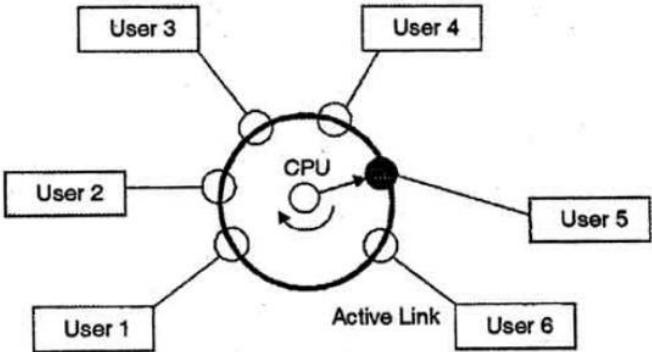
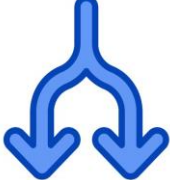
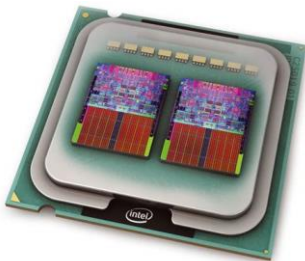


Multiple streams



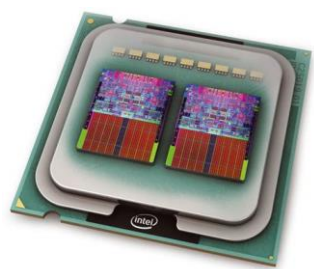
Multiple streams

Exploitation

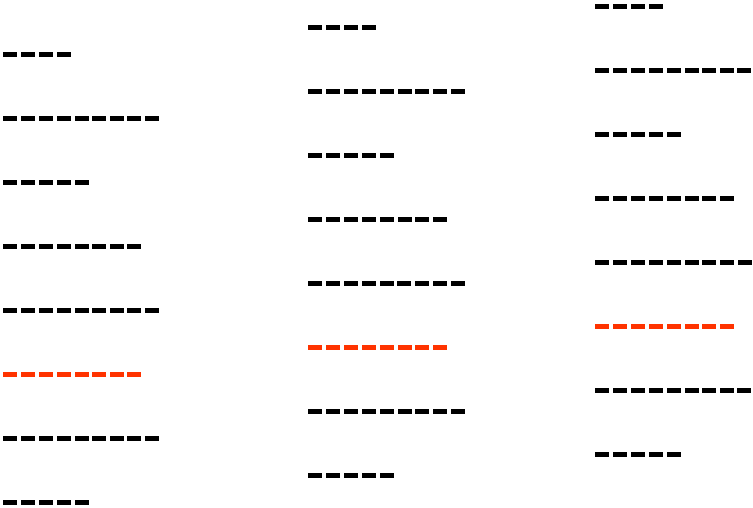
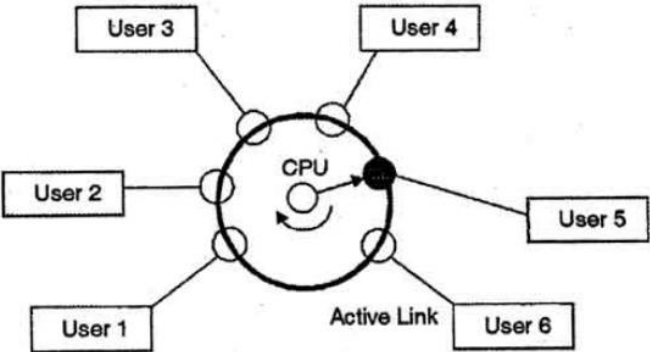
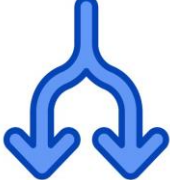


Multiple streams

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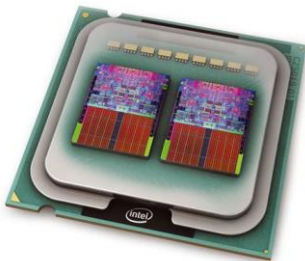


Inherent

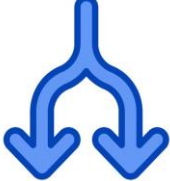


Multiple streams

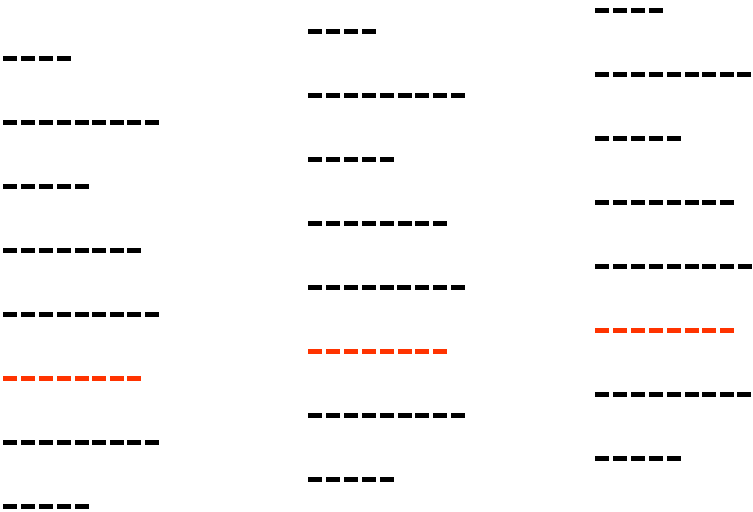
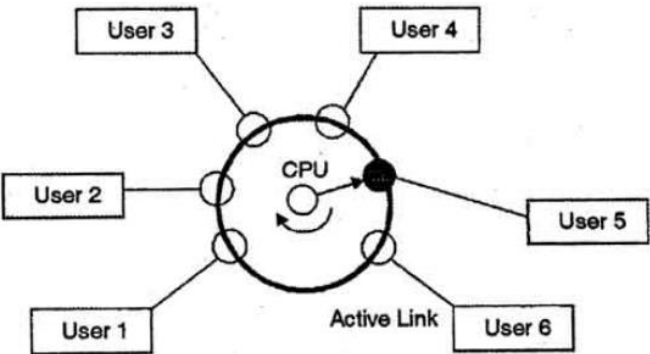
Exploitation



Inherent

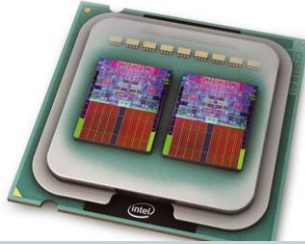


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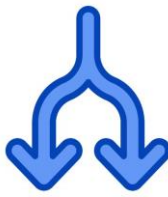


Multiple streams

Exploitation

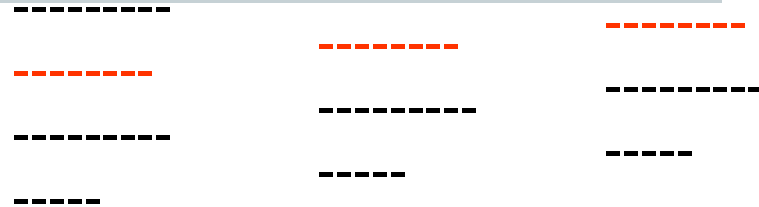


Inherent



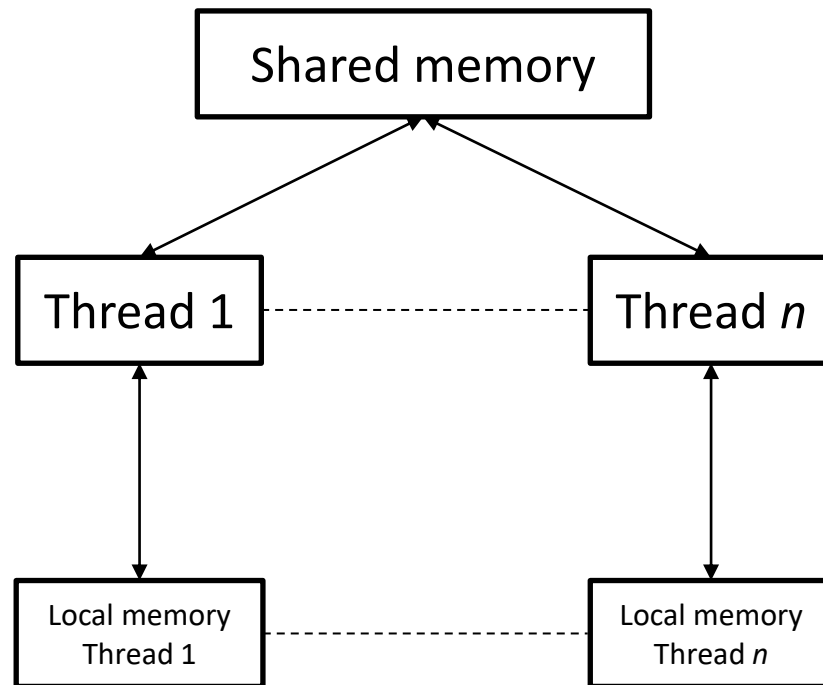
Concurrency is an abstraction for all of these

(and more)





- A *thread* is a stream of program statements executed sequentially
- Several threads can be executed at the same time, i.e., *concurrently*
- Each threads works at its own speed
- Each thread has its own local memory
- Threads can communicate via shared memory



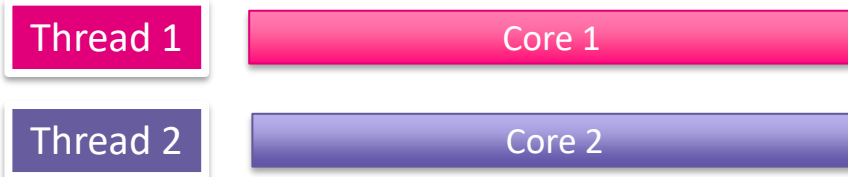
Sequential vs Parallel vs Concurrent



- **Sequential** execution (in 1 processing core)



- **Parallel** execution (in 2 processing cores)



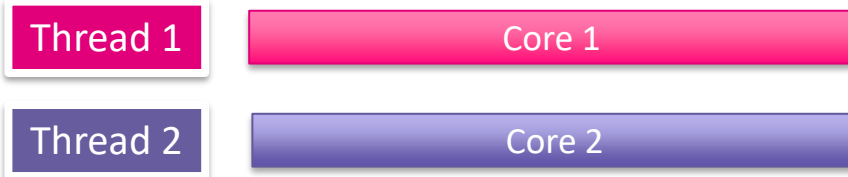
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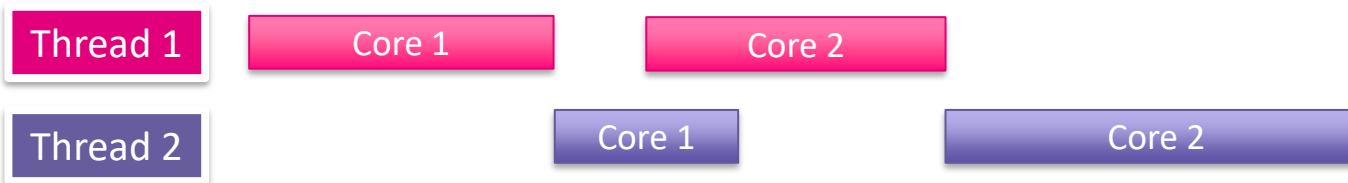
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- **Concurrent** execution (in 2 processing cores)



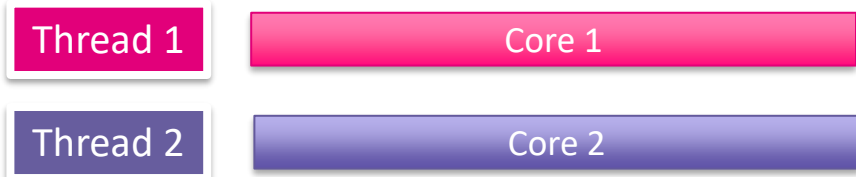
Sequential vs Parallel vs Concurrent



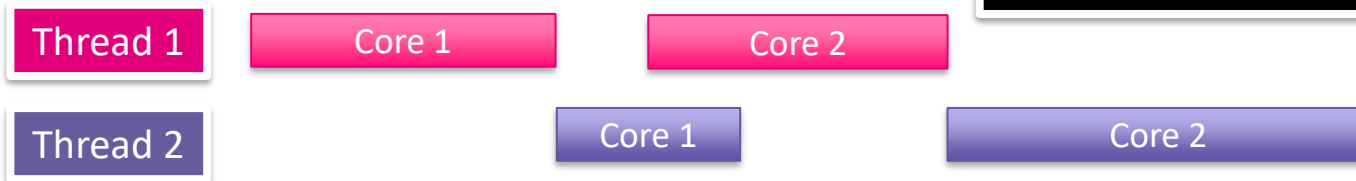
- **Sequential** execution (in 1 processing core)



- **Parallel** execution (in 2 processing cores)



- **Concurrent** execution (in 2 processing cores)



- Computation is split in arbitrary blocks (as opposed to sequential)
- Concurrent execution can happen with one core (as opposed to parallel)
- When there are more than two cores different threads may run in parallel

Threads in Java – Example

Tivoli entrance turnstile



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Threads in Java - Example

19



- Java threads can be created either by implementing **Runnable** or extending from **Thread**
- The behaviour of the thread is in the **run()** method (override)

```
final long PEOPLE = 10_000;  
long counter = 0;
```

...

```
public class Turnstile extends Thread {  
    public void run() {  
        for (int i = 0; i < PEOPLE; i++) {  
            counter++;  
        }  
    }  
}
```

This thread simulates 10000
people entering to Tivoli

Threads in Java - Example

19



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- The behaviour of the thread is in the **run ()** method (override)

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final long PEOPLE = 10_000;  
long counter = 0;  
...
```

Shared memory

This thread simulates 10000 people entering to Tivoli

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Threads in Java - Example

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```
final long PEOPLE = 10_000;  
long counter = 0;  
...
```

Shared memory

This thread simulates 10000
people entering to Tivoli

```
public class TurnstileThread {  
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        for (int i = 0; i < PEOPLE; i++) {  
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        }  
    }  
}
```

Local memory

Threads in Java - Example

19



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- The behaviour of the thread is in the **run ()** method (override)

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Shared memory

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public class TurnstileThread {  
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        for (int i = 0; i < PEOPLE; i++) {  
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        }  
    }  
}
```

Local memory

Behaviour of
the thread

Threads in Java - Example



- The function **start()** starts the execution of the thread
- The function **join()** waits for the thread to terminate

```
Turnstile turnstile = new Turnstile();
```

Create an instance
of the thread

```
turnstile.start();
```

Start execution of
the thread

```
turnstile.join();
```

Wait until the
thread terminates

Print the value of
counter

```
System.out.println(counter+" people entered");
```



- Altogether (not executable, see `CounterThreads.java` for the executable program)

```
class CounterThreads {
    long counter = 0;
    final long PEOPLE = 10_000;

    // main thread behaviour
    Turnstile turnstile = new Turnstile();
    turnstile.start();
    turnstile.join();
    System.out.println(counter+" people entered");

    // inner class for accessing shared variables
    public class Turnstile extends Thread {
        public void run() {
            for (int i = 0; i < PEOPLE; i++) {
                counter++;
            }
        }
    }
}
```



- Altogether (not executable, see `CounterThreads.java` for the executable program)

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            for (int i = 0; i < PEOPLE; i++) {  
                counter++;  
            }  
        }  
    }  
}
```

Shared memory

Create, start, wait
till termination
and print results

Definition of the
thread's behaviour

Threads in Java - Example

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What value of **counter** will this program print?
Executable, see **CounterThreads.java** for (program)

```
class CounterThreads {  
    long counter = 0;  
    final long PEOPLE = 10_000;  
  
    // main thread behaviour  
    Turnstile turnstile = new Turnstile();  
    turnstile.start();  
    turnstile.join();  
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        }  
    }  
}
```

Shared memory

Create, start, wait
till termination
and print results

Definition of the
thread's behaviour



Other ways to define threads

Runnable object in the thread constructor

```
long counter = 0;
final long PEOPLE = 10_000;

Thread t = new Thread(new Runnable() {
    public void run() {
        for (int i=0; i<PEOPLE; i++){
            counter++;
        }
    }
});
t.start();
t.join();
System.out.println(counter+" people entered");
```

Using Java lambda expressions

```
long counter = 0;
final long PEOPLE = 10_000;

Thread t = new Thread(() -> {
    for (int i=0; i<PEOPLE; i++){
        counter++;
    }
});
t.start();
t.join();
System.out.println(counter+" people entered");
```



Other ways to define threads

Runnable object in the thread constructor

```
long counter = 0;
final long PEOPLE = 10_000;

Thread t = new Thread(new Runnable() {
    public void run() {
        for (int i=0; i<PEOPLE; i++){
            counter++;
        }
    }
});
t.start();
t.join();
System.out.println(counter+" people entered");
```

Using Java lambda expressions

```
long counter = 0;
final long PEOPLE = 10_000;

Thread t = new Thread(() -> {
    for (int i=0; i<PEOPLE; i++){
        counter++;
    }
});
t.start();
t.join();
System.out.println(counter+" people entered");
```

I would only recommend these when the thread's code is small, e.g., without several methods. And when the local state of the thread is minimal. WARNING: Possibly bias opinion!

Threads in Java – Example

Tivoli entrance turnstile



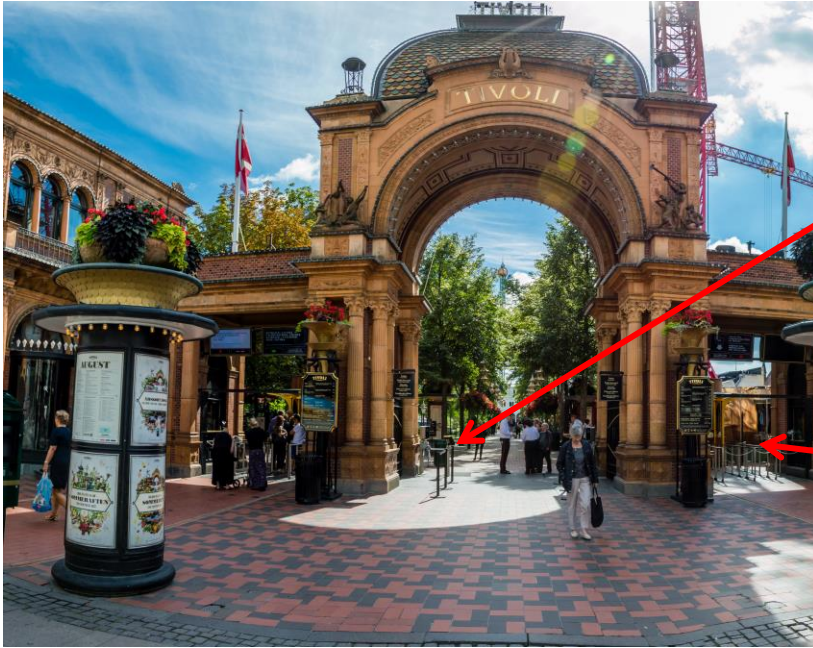
23



Threads in Java – Example II

Tivoli entrance turnstile

24





- Altogether (not executable, see `CounterThreads2.java` for the executable program)

```
long counter = 0;
final long PEOPLE = 10_000;

Turnstile turnstile1 = new Turnstile();
Turnstile turnstile2 = new Turnstile();
turnstile1.start();turnstile2.start();
turnstile2.join();turnstile2.join();
System.out.println(counter+" people entered");

public class Turnstile extends Thread {
    public void run() {
        for (int i = 0; i < PEOPLE; i++) {
            counter++;
        }
    }
}
```

We simply add another
Turnstile thread

Threads in Java – Example II

25



What value of **counter** will this program print? (cutable, see **CounterThreads2.java** for program)

```
long counter = 0;
final long PEOPLE = 10_000;

Turnstile turnstile1 = new Turnstile();
Turnstile turnstile2 = new Turnstile();
turnstile1.start();turnstile2.start();
turnstile2.join();turnstile2.join();
System.out.println(counter+" people entered");

public class Turnstile extends Thread {
    public void run() {
        for (int i = 0; i < PEOPLE; i++) {
            counter++;
        }
    }
}
```

We simply add another
Turnstile thread

Threads in Java – Example II

25



What value of **counter** will this program print? (cutable, see **CounterThreads2.java** for program)



```
long counter = 0;
final long PEOPLE = 10_000;

Turnstile turnstile1 = new Turnstile();
Turnstile turnstile2 = new Turnstile();
turnstile1.start();turnstile2.start();
turnstile2.join();turnstile2.join();
System.out.println(counter+" people entered");

public class Turnstile extends Thread {
    public void run() {
        for (int i = 0; i < PEOPLE; i++) {
            counter++;
        }
    }
}
```

We simply add another Turnstile thread

Threads in Java – Example II

25



What value of **counter** will this program print? (cutable, see **CounterThreads2.java** for program)



```
long counter = 0;
final long PEOPLE = 10_000;

Turnstile turnstile1 = new Turnstile();
Turnstile turnstile2 = new Turnstile();
turnstile1.start();turnstile2.start();
turnstile2.join();turnstile2.join();
System.out.println(counter+" people entered");
```

We simply add another Turnstile thread

```
public class Turnstile extends Thread {
    public void run() {
        i++) {
    }
}
```

VERY HARD: What is the minimum value of **counter** that this program can print?

- What was is the problem in the previous program?
- To answer this question we need to understand
 - Atomicity
 - States of a thread
 - Non-determinism
 - Interleavings



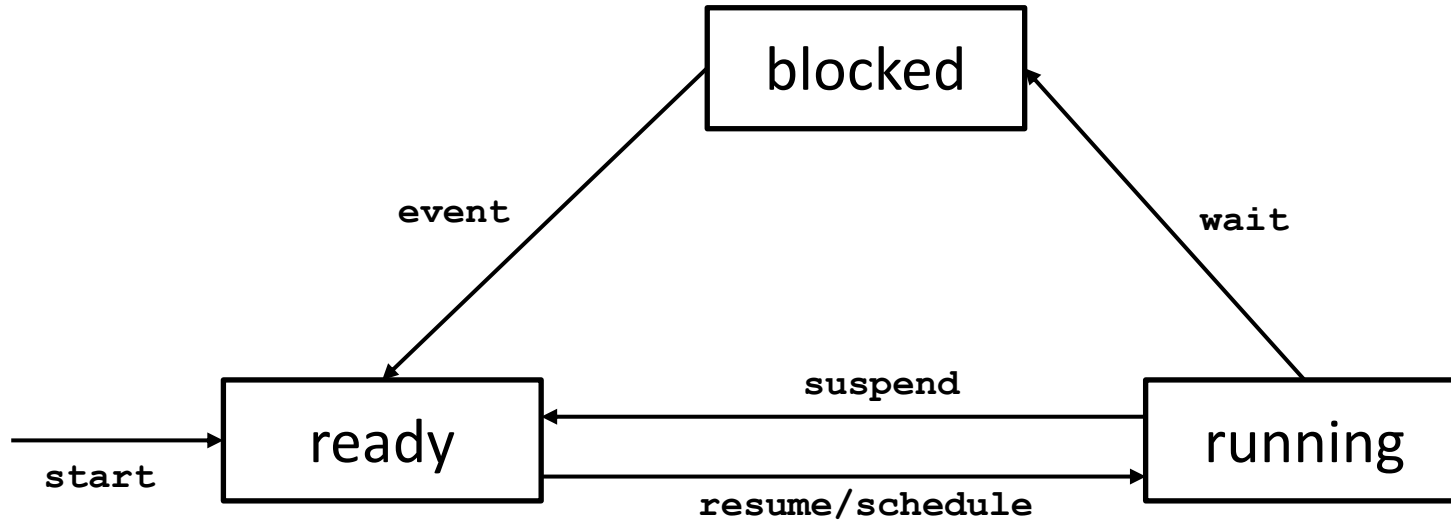
- The program statement **counter++** is not *atomic*
- Atomic statements are executed as a single (indivisible) operation

```
public class Turnstile extends Thread {  
    public void run() {  
        for (int i = 0; i < PEOPLE; i++) {  
            counter++;  
        }  
    }  
}
```

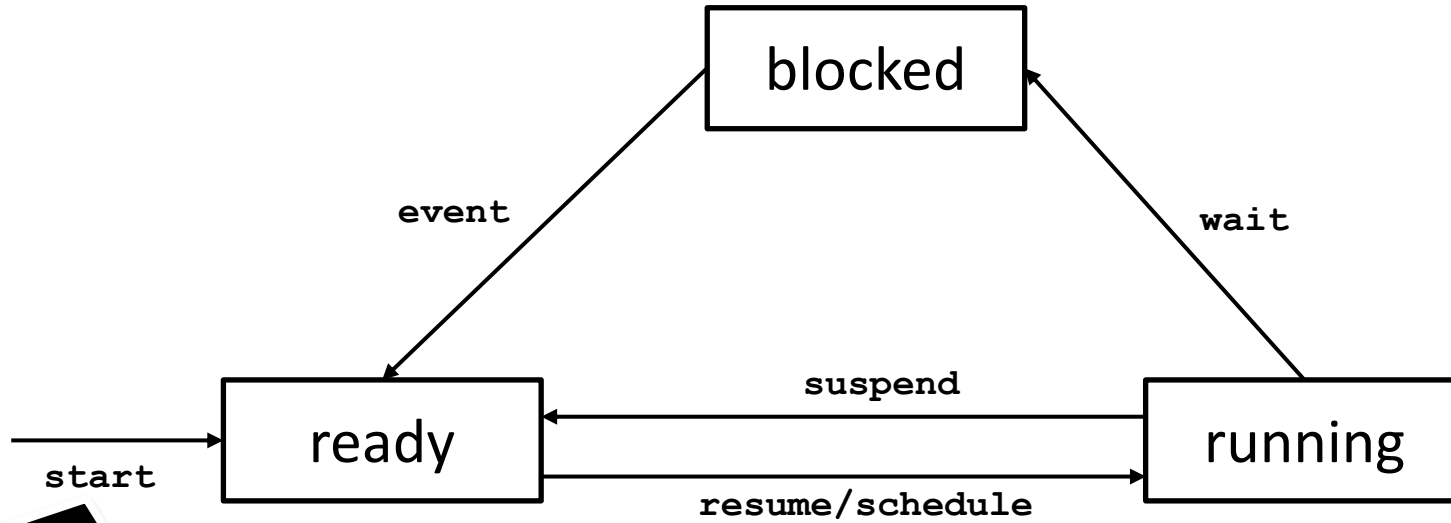
`int temp = counter;
counter = temp + 1;`

Watchout: Just because a program statement is a one-liner, it doesn't mean that it is atomic

States of a thread (simplified)

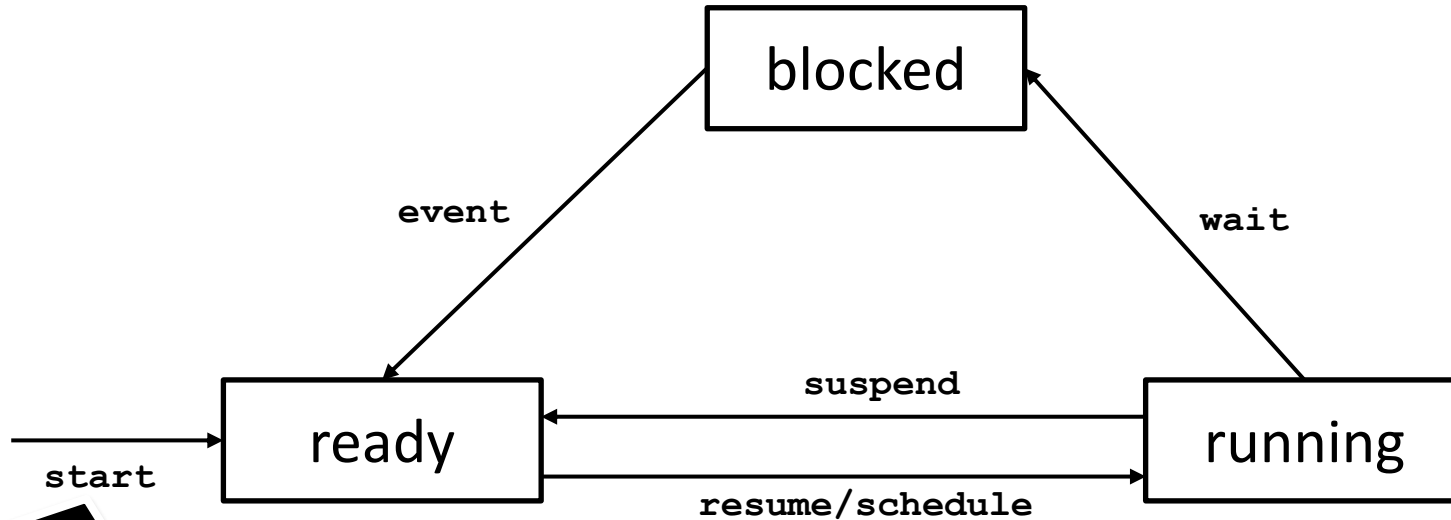


States of a thread (simplified)



This event corresponds to after executing `start()`

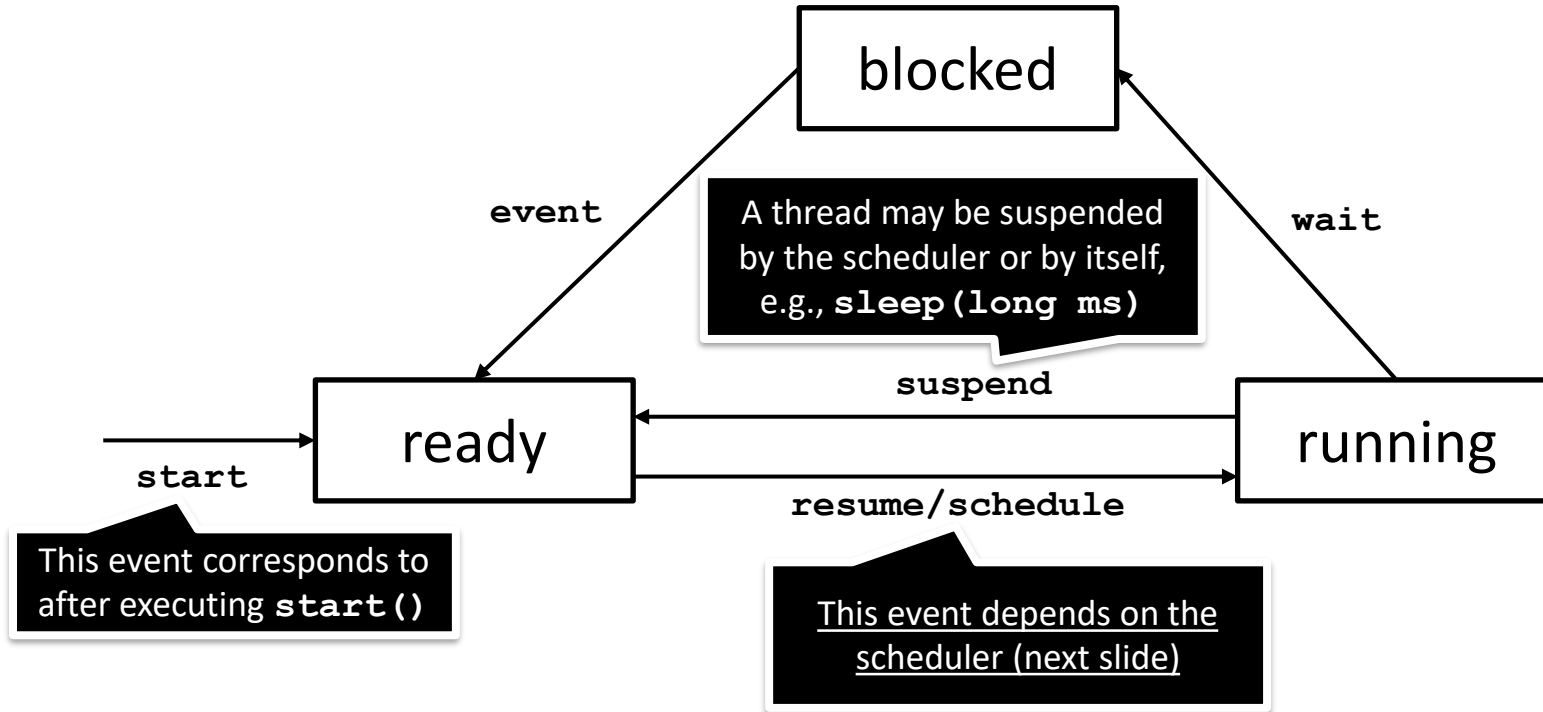
States of a thread (simplified)



This event corresponds to after executing `start()`

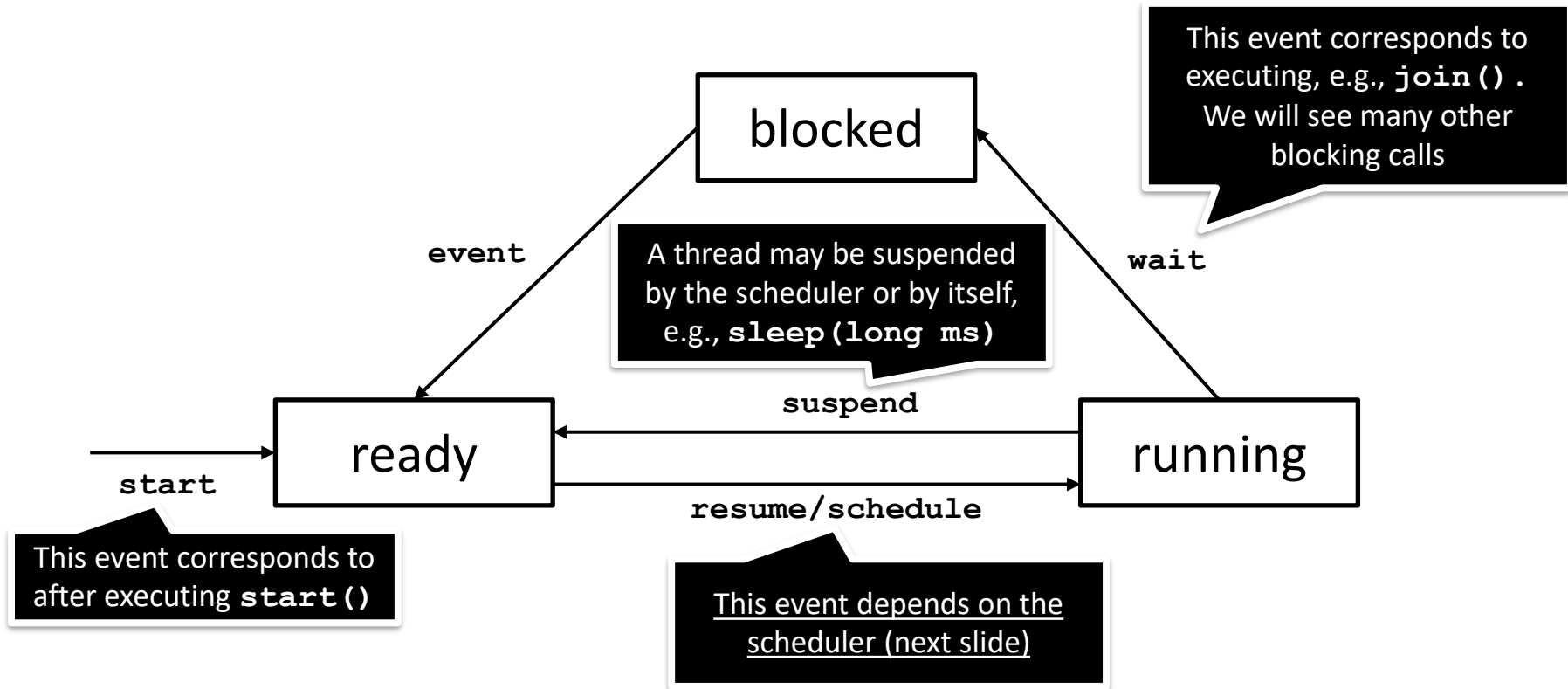
This event depends on the scheduler (next slide)

States of a thread (simplified)



States of a thread (simplified)

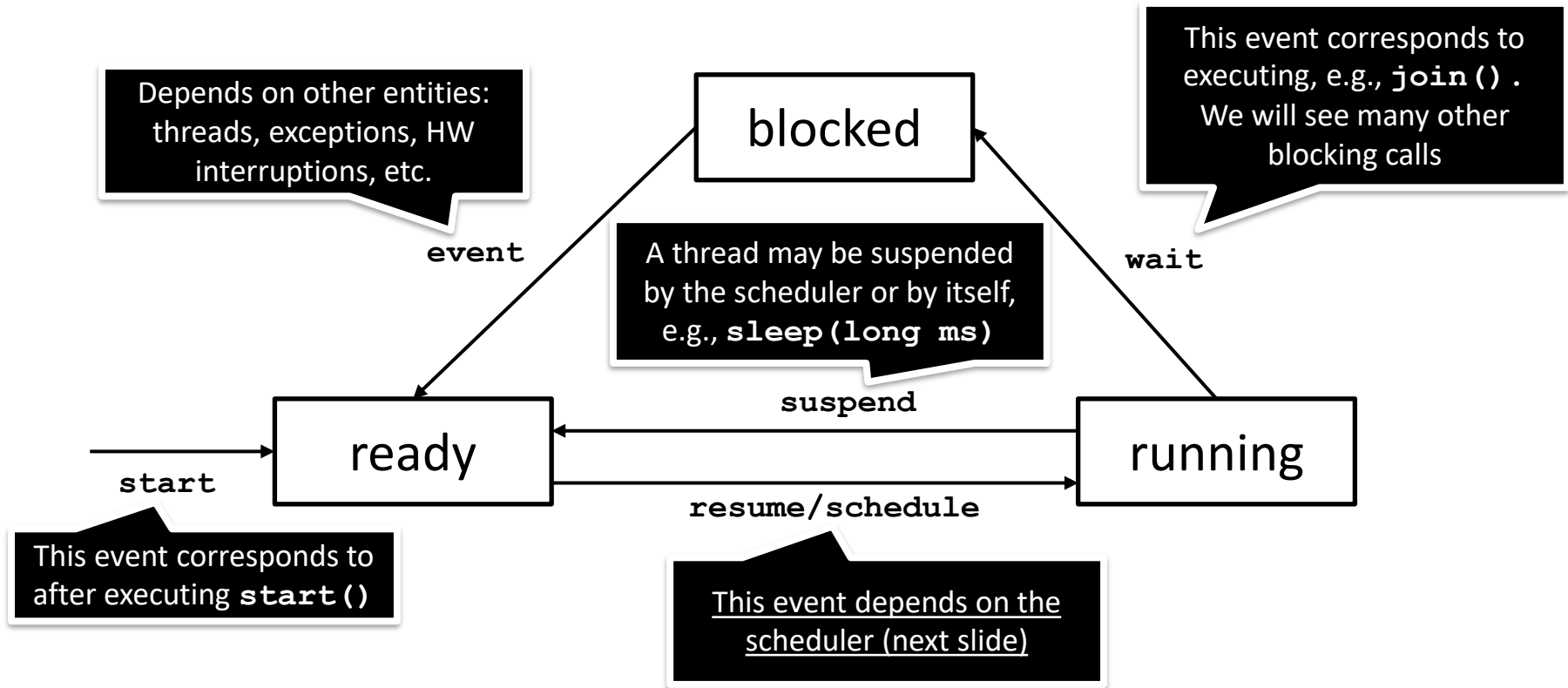
· 28



States of a thread (simplified)



· 28





- In all operating systems/executing environments a *scheduler* selects the processes/threads under execution
 - Threads are selected *non-deterministically*, i.e., no assumptions can be made about what thread will be executed next
- Consider two threads $t1$ and $t2$ in the ready state; $t1(ready)$ and $t2(ready)$
 1. $t1(running) \rightarrow t1(ready) \rightarrow t1(running) \rightarrow t1(ready) \rightarrow \dots$
 2. $t2(running) \rightarrow t2(ready) \rightarrow t2(running) \rightarrow t2(ready) \rightarrow \dots$
 3. $t1(running) \rightarrow t1(ready) \rightarrow t2(running) \rightarrow t2(ready) \rightarrow \dots$
 4. Infinitely many different executions!

- The statements in a thread are executed when the thread is in its “running” state
- An *interleaving* is a possible sequence of operations for a concurrent program
 - Note this: a sequence of operations for a concurrent program, not for a thread. Concurrent programs are composed by 2 or more threads.

Interleaving – Example I

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main

```
long counter = 0;
final long PEOPLE = 10_000;

Turnstile turnstile1 = new Turnstile();
Turnstile turnstile2 = new Turnstile();
turnstile1.start(); turnstile1.start();
```

time

turnstile1

```
// first iteration loop
int i = 0
if (i < PEOPLE)
int temp = counter // counter == 0
counter = temp + 1 // counter == 1
```

time

turnstile2

```
// first iteration loop
int i = 0
if (i < PEOPLE)
int temp = counter // counter == 1
counter = temp + 1 // counter == 2
...
```

time

Interleaving – Example II

· 32



main

```
long counter = 0;  
final long PEOPLE = 10_000;
```

```
Turnstile turnstile1 = new Turnstile();  
Turnstile turnstile2 = new Turnstile();  
turnstile1.start(); turnstile1.start();
```

time

turnstile1

```
// first iteration loop  
int i = 0  
if (i < PEOPLE)  
int temp = counter // counter == 0
```

```
counter = temp + 1 // counter == 1
```

time



turnstile2

```
// first iteration loop  
int i = 0  
if (i < PEOPLE)  
int temp = counter // counter==0
```

```
counter = temp + 1 // counter == 1  
i++  
...
```

time

Interleaving – Example II

· 32



main

```
long counter = 0;  
final long PEOPLE = 10_000;
```

```
Turnstile turnstile1 = new Turnstile();  
Turnstile turnstile2 = new Turnstile();  
turnstile1.start(); turnstile1.start();
```

Are there any other interleavings?

time

turnstile1

```
// first iteration loop  
int i = 0  
if (i < PEOPLE)  
int temp = counter // counter == 0
```

```
counter = temp + 1 // counter == 1
```

time



turnstile2

```
// first iteration loop  
int i = 0  
if (i < PEOPLE)  
int temp = counter // counter==0
```

```
counter = temp + 1 // counter == 1  
i++  
...
```

time

Interleaving – Example II

· 32



main

```
long counter = 0;
final long PEOPLE = 10_000;

Turnstile turnstile1 = new Turnstile();
Turnstile turnstile2 = new Turnstile();
turnstile1.start(); turnstile1.start();
```

Are there any other interleavings?

There are as many interleavings as possible ways to order the operations in the program

turnstile1

```
// first iteration loop
int i = 0
if (i < PEOPLE)
int temp = counter // counter == 0
```

```
counter = temp + 1 // counter == 1
```

time ↓



turnstile2

```
// first iteration loop
int i = 0
if (i < PEOPLE)
int temp = counter // counter==0
```

```
counter = temp + 1 // counter == 1
i++
...
```

time ↓

Concurrency Humour...

- *Knock, knock*
- *Race condition*
- *Who's there?*



- *A **race condition** occurs when the result of the computation depends on the interleavings of the operations*



- *A **data race** occurs when two concurrent threads:*
 - *Access a shared memory location*
 - *At least one access is a write*

Race Conditions vs Data Races



Not all race conditions are data races

- Threads may not access shared memory
- Threads may not write on shared memory

Not all data races result in race conditions

- The result of the program may not change based on the writes of threads

- What was is the problem in the previous program?
 - The statement **counter++** is not atomic
 - Some interleavings result in threads reading stale (outdated) data
 - Consequently, the program has race conditions
- In what follows, we will see how to tackle this type of problems



- A *critical section* is a part of the program that only one thread can execute at the same time
 - Useful to avoid race conditions in concurrent programs

```
public class Turnstile extends Thread {  
    public void run() {  
        for (int i = 0; i < PEOPLE; i++) {  
            // start critical section  
            int temp = counter;  
            counter = temp + 1;  
            // end critical section  
        }  
    }  
}
```




- A *critical section* is a part of the program that only one thread can execute at the same time
 - Useful to avoid race conditions in concurrent programs

```
public class Turnstile extends Thread {  
    public void run() {  
        for (int i = 0; i < PEOPLE; i++) {  
            // start critical section  
            int temp = counter;  
            counter = temp + 1;  
            // end critical section  
        }  
    }  
}
```

Critical sections should cover the parts of the code handling shared memory



- A *critical section* is a part of the program that only one thread can execute at the same time

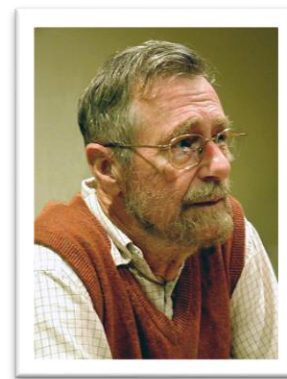
... can cause race conditions in concurrent programs

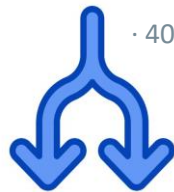
Shouldn't the critical section start before the **for**?

```
Countstile extends Thread {  
    public void run() {  
        for (int i = 0; i < PEOPLE; i++) {  
            // start critical section  
            int temp = counter;  
            counter = temp + 1;  
            // end critical section  
        }  
    }  
}
```

Critical sections should cover the parts of the code handling shared memory

- The *mutual exclusion* property states that
 - *Two or more threads cannot be executing their critical section at the same time*
- Mutual exclusion was first formulated by EW Dijkstra (see optional readings)
 - He devised a protocol to ensure mutual exclusion (*solving the mutual exclusion problem*)
 - He laid down the properties for a satisfactory solution to ensuring mutual exclusion





- An ideal solution to the mutual exclusion problem must ensure:
 - Mutual exclusion: at most one thread executing the critical section at the same time
 - Absence of *deadlock*: threads eventually exit the critical section allowing other threads to enter
 - Absence of *starvation*: if a thread is ready to enter the critical section, it must eventually do so



- An ideal solution to the mutual exclusion problem must ensure:
 - Mutual exclusion: at most one thread executing the critical section at the same time
 - Absence of *deadlock*: threads eventually exit the critical section allowing other threads to enter
 - Absence of *starvation*: if a thread is ready to enter the critical section, it must eventually do so

In practice, we will see that it is not always possible to achieve absence of starvation



- In Java, mutual exclusion can be achieved using the **Lock** interface in the `java.util.concurrent.locks` package
 - **lock()**
 - Acquires the lock if available, otherwise it blocks
 - It is blocking
 - **unlock()**
 - Releases the lock, if there are other threads waiting for the lock it signals one of them
 - It is not blocking



- Simple protocol: call `lock()` before entering the critical section, and `unlock()` after exiting
- Each critical section must have a lock associated to it, but many critical sections may use the same lock.
- Simplified, see `CounterThreadsLock.java`

```
Lock l = new Lock();

public class Turnstile extends Thread {
    public void run() {
        for (int i = 0; i < PEOPLE; i++) {
            l.lock()           // start critical section
            int temp = counter;
            counter = temp + 1;
            l.unlock()         // end critical section
        }
    }
}
```

Java Locks | Critical sections

· 42



- Simple protocol: call `lock()` before entering the critical section, and `unlock()` after exiting
- Each critical section must have a lock associated to it. However, critical sections may use the same lock.
- Simplified, see `CounterThreadsLock.java`

```
Lock l = new Lock();

public class Turnstile extends Thread {
    public void run() {
        for (int i = 0; i < PEOPLE; i++) {
            l.lock()           // start critical section
            int temp = counter;
            counter = temp + 1;
            l.unlock()         // end critical section
        }
    }
}
```

- We do not focus on the implementation of `lock()/unlock()`, but in their use to solve concurrency problems.
- See Sections 2.3, 2.5, 2.7 and Section 7.3 onwards in Herlihy for implementation details of `lock()/unlock()`

Java Locks | Interleavings

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main

```
long counter = 0;
final long PEOPLE = 10_000;
Lock l = new Lock()

Turnstile turnstile1 = new Turnstile();
Turnstile turnstile2 = new Turnstile();
turnstile1.start(); turnstile1.start();
```

turnstile1

```
// first iteration loop
int i = 0
```

```
if (i < PEOPLE)
```

```
l.lock() // begin critical section
int temp = counter // counter == 0
counter = temp + 1 // counter == 1
l.unlock() // end critical section
```

Operations in the critical section
are always executed sequentially
by the same thread

turnstile2

```
// first iteration loop
int i = 0
```

```
if (i < PEOPLE)
```

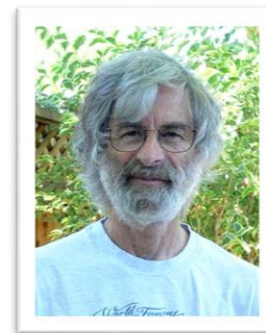
```
l.lock() // begin critical section
int temp = counter // counter == 1
counter = temp + 1 // counter == 2
l.unlock() // end critical section
```



- In fact, we can now characterize an order of execution between some of the operations of a program
- We say that an operation a *happens-before* than operation b , denoted as $a \rightarrow b$, iff
 - a and b belong to the same thread and a appears before b in the thread definition
 - a is an **unlock()** and b is a **lock()** on the same lock
- In the absence of *happens-before* relation between operations, the JVM is free to choose any execution order
 - In that case we say that operations are executed *concurrently*
 - Sometimes denoted as $a \parallel b$



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 - Sometimes denoted as $a \parallel b$
- *Happens-before* is a *pre-order* over operations of concurrent programs
 - Transitive
 - Anti-symmetric
- “Happened-before” was first introduced by Leslie Lamport for distributed systems
 - See optional readings





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 - Anti-symmetric
- “Happened-before”
 - See optional r



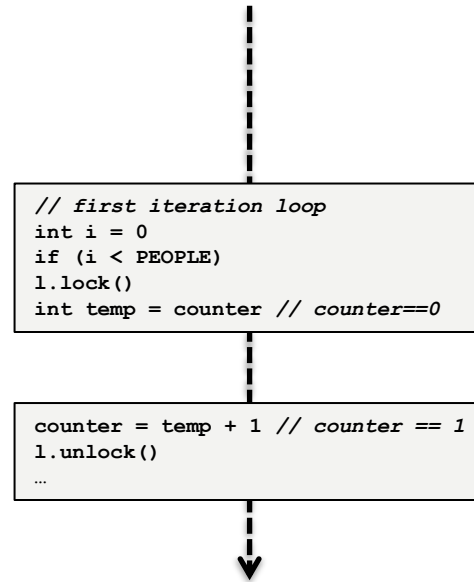
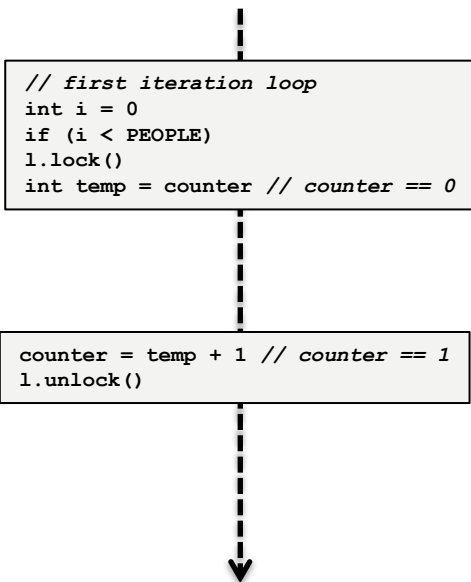
“Don’t be brainwashed by programming languages. Free your mind with mathematics.” Time for one question before we go straight to the next talk.

#HLF18



- We use locks to remove undesired interleavings, and happens-before can help us reasoning about correctness

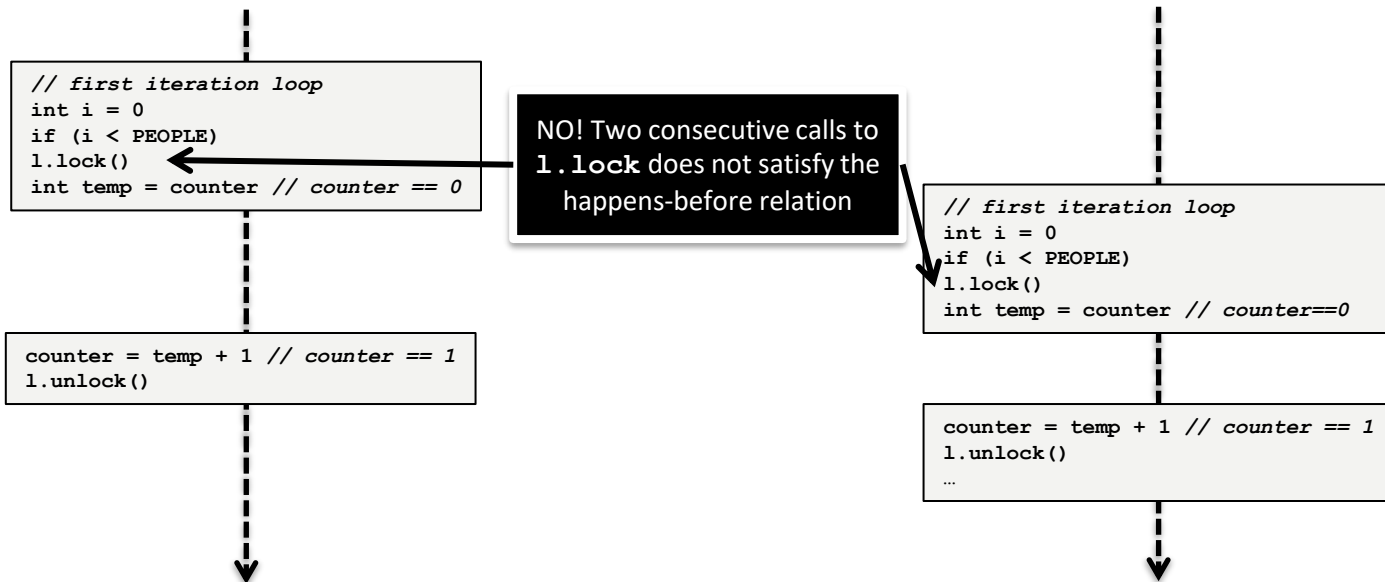
Is this a valid interleaving?





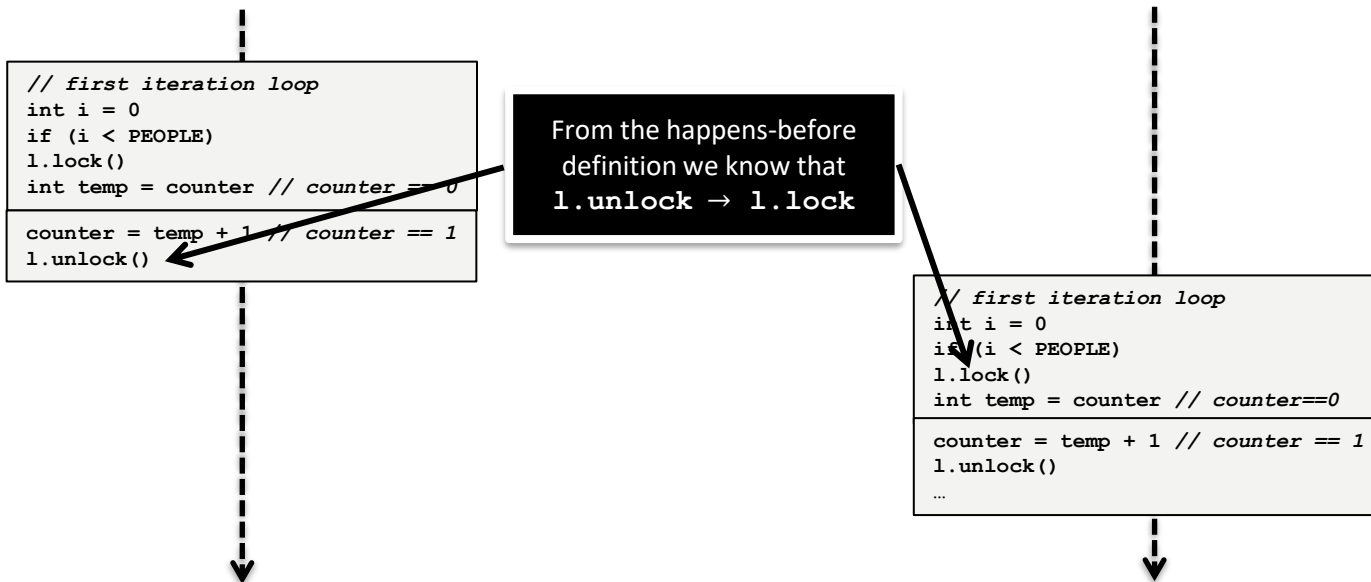
- We use locks to remove undesired interleavings, and happens-before can help us reasoning about correctness

Is this a valid interleaving?



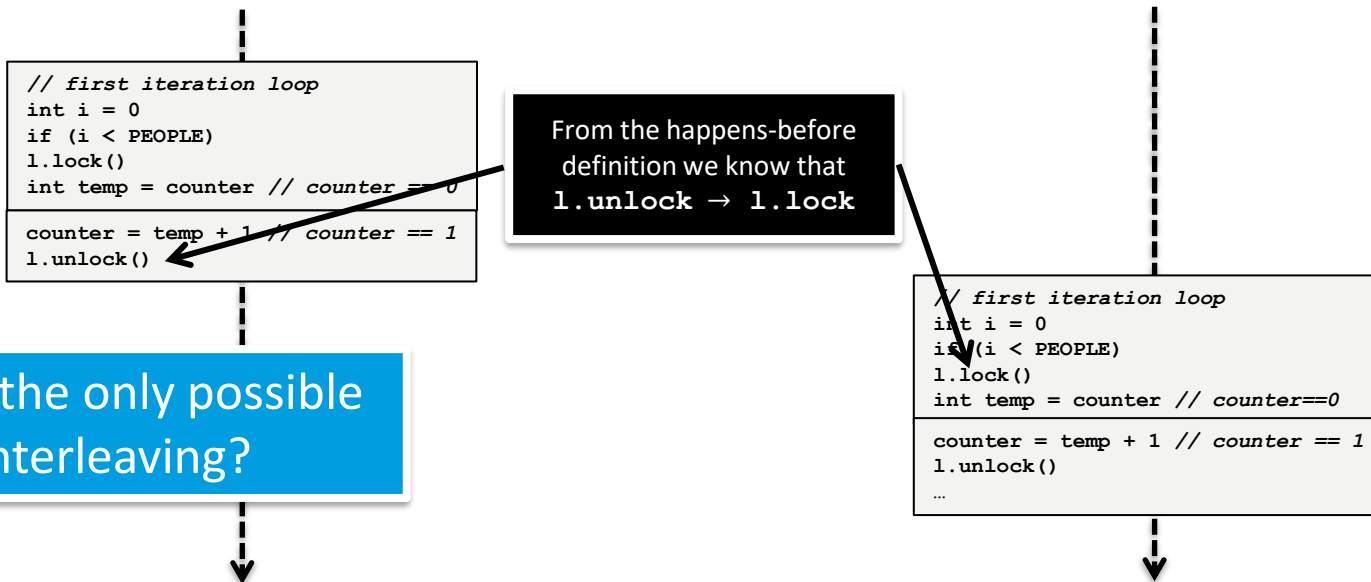


- We use locks to remove undesired interleavings, and happens-before can help us reasoning about correctness





- We use locks to remove undesired interleavings, and happens-before can help us reasoning about correctness





- The solution to the turnstile problem ensures mutual exclusion **but does it ensure absence of deadlock?**

```
Lock l = new Lock();

public class Turnstile extends Thread {
    public void run() {
        for (int i = 0; i < PEOPLE; i++) {
            l.lock()           // start critical section
            int temp = counter;
            counter = temp + 1;
            l.unlock()         // end critical section
        }
    }
}
```

Absence of *deadlock*: threads eventually exit the critical section allowing other threads to enter

(Naïve) examples of deadlocks



- Locking twice within the thread

```
for (int i = 0; i < PEOPLE; i++) {  
    l.lock()  
    int temp = counter;  
    counter = temp + 1;  
    l.lock() // blocks forever  
            // or forgetting to unlock()  
}
```

Thread 1

This is not unrealistic.
For instance, see these bug
reposts in the Linux kernel

[1] <https://github.com/torvalds/linux/commit/e1db4ce>
[2] <https://github.com/torvalds/linux/commit/2904207>

- Exception during execution

```
for (int i = 0; i < PEOPLE; i++) {  
    l.lock()  
    l2.lock()  
    int temp = counter;  
    throw new Exception();  
    // If exception handling doesn't  
    // unlock, blocks forever  
}
```

Thread 1

```
for (int i = 0; i < PEOPLE; i++) {  
    l2.lock() // blocks forever  
    ...  
}
```

Thread 2

(Naïve) examples of deadlocks

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- Locking twice within the thread

```
for (int i = 0; i < PEOPLE; i++) {  
    l.lock()  
    int temp = counter;  
    counter = temp + 1;  
    l.lock() // blocks forever  
            // or forgetting to unlock()  
}
```

Thread 1

This is not unrealistic.
For instance, see these bug
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[1] <https://github.com/torvalds/linux/commit/e1db4ce>
[2] <https://github.com/torvalds/linux/commit/2904207>

- Exception

Why did I write 1 and 12
in this example?

```
for (int i = 0; i < PEOPLE; i++) {  
    l.lock()  
    l2.lock()  
    int temp = counter;  
    throw new Exception();  
    // If exception handling doesn't  
    // unlock, blocks forever  
}
```

Thread 1

```
for (int i = 0; i < PEOPLE; i++) {  
    l2.lock() // blocks forever  
    ...  
}
```

Thread 2



- When using Java locks, it is recommended to use the idiom on the right
- This prevents deadlocks problems if the thread finish unexpectedly due to exception
 - Solutions not following this idiom cannot be deemed as free from deadlocks (note for assignments)
- We use this idiom in **CounterThreadLocks.java**

```
Lock l = new Lock();  
  
l.lock()  
try {  
    // critical section code  
} finally {  
    l.unlock()  
}
```

- Java **Lock** is an interface, so it cannot be used as we showed in the examples today
- We use an implementation of the **Lock** interface, namely **ReentrantLock**
 - Reentrant locks act like a regular lock, except that they allow locks to be locked more than once by the same thread

```
Lock l = new ReentrantLock();

for (int i = 0; i < PEOPLE; i++) {
    l.lock()
    int temp = counter;
    counter = temp + 1;
    l.lock();    // it doesn't block
    l.unlock();  // still holds the lock
    l.unlock();  // now the lock is free
}
```