

SYSTEM PROGRAMMING

CHAPTER ONE: INTRODUCTION

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CONTENTS

- Basic terminology
- System software and machine architecture
- The Simplified Instructional Computer (SIC)
 - SIC and SIC/XE Machine architecture
 - Data and instruction formats
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 - I/O
 - Programming Examples

BASIC TERMINOLOGY

- System: Group of related things that works towards common goal.
 - An organized collection of parts(sub systems)are integrated together to accomplish goal
 - E.g. computer(made from integrated components like I/O, cpu , storage etc.)
- Program/Software: executable sw runs on a machine like computer
 - Consists a compiled code that can run directly from the computer's operating system.
- Types of software: there are two types
 - Application SW
 - System SW

CON...

Application Software

- A set of programs written in a specific programming language to solve a particular problem.
 - It is independent of the machine on which it is operated.
 - Concerned mainly with the solution of the problem and makes use of the computer software as a tool.
 - Purpose of supporting or improving the user's work with a software (it gives services for user)
 - The focus is on the application
- Example:
- Word processing :MS word, WordPad
 - Database :Oracle, MS Access
 - Spreadsheet :Excel, Apple Numbers
 - Multimedia :Real Player, Media Player
 - Presentation Graphics :MS PowerPoint

CON...

System Software

- It is software designed to provide a platform for other software.
 - A set of programs that creates the environment to facilitate working of an application software.
 - Able effective execution of application Software
 - Acts as an intermediary between computer hardware and application programs.
 - Controls the computer system and enhances its performance.
 - Focus is on the system.
 - Machine dependent(Instruction Set, Instruction Format, Addressing Mode)
 - However , general design and logic of an assembler and compiler, code optimization techniques are machine independent
- System programming means development of computer programs that allow the entire system to function as a single unit.

EXAMPLES OF SYSTEM SOFTWARE

- Operating **systems**: like macOS, Linux, Android and Microsoft Windows, computational science **software**, game engines, industrial automation,
- Text Editor
 - Allows user to create/modify a file having only plain text.
 - Edit contents to get an immediate visual feedback
 - Notepad (Microsoft)
- Compiler
 - Translates high level language to machine language.
 - Source code into data, that computer can understand.
 - Source code to object code
 - ADA, BASIC, Fortran, Pascal
 - C, C++, C#
 - Open Source like java
- Linker
 - Gathers 1 or more objects codes generated by the compiler & combines into a single EXE

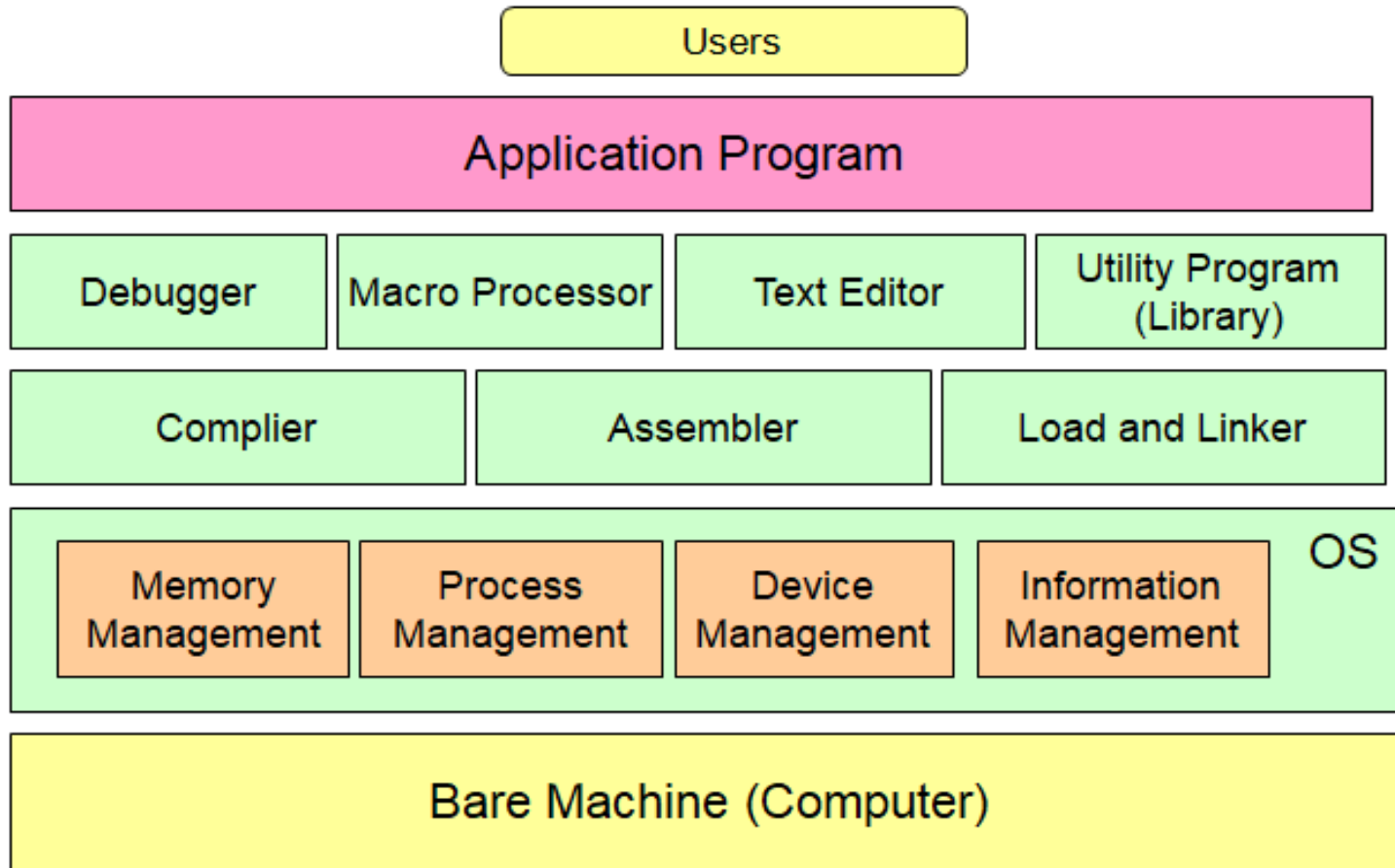
CON...

- Loader
 - Loads object code into memory for execution
- Debugger
 - Tests & debugs errors
 - Code run on Instruction Set Simulator (ISS)
 - DEBUG – Microsoft DOS built-in debugger
 - Nemiver – Graphical C/C++ debugger
 - Ups – Fortran, C
 - VB Watch – Visual Basic
 - Jswat – Open source Java
 - Xdebug – PHP
- Assembler
 - Assembly language to machine language
 - Mnemonics to machine code , Produces executable machine code
- Interpreter
 - High level language to machine language

NEED FOR SYSTEM SOFTWARE

- Controls some of the aspects of operation of computers.
 - File operation
 - I/O operation
 - Memory management
 - Etc.

SYSTEM SOFTWARE CONCEPT



SYSTEM SOFTWARE AND MACHINE ARCHITECTURE

- One characteristic in which most system software differ from application software is **machine dependency**.
 - e.g. assembler translate mnemonic instructions into machine code
 - e.g. compilers must generate machine language code
 - e.g. operating systems are directly concerned with the management of nearly all of the resources of a computing system
- Some aspects of system software that do not directly depend upon the type of computing system
 - e.g. general design and logic of an assembler and compiler
 - e.g. code optimization techniques

THE SIMPLIFIED INSTRUCTIONAL COMPUTER (SIC)

- **SIC** is a hypothetical computer system introduced in *System Software*.
 - It includes the hardware features most often found on real machines while avoiding unusual or irrelevant complexities
 - Since it is difficult to learn system programming on real computer system SIC introduced.
- **Two versions of SIC**
 - SIC - standard model
 - SIC/XE (extra equipments or extra expensive)
 - SIC program can be executed on SIC/XE (upward compatible)
- The two versions has been designed to be upward compatible.
 - Object program for the standard SIC machine will execute properly on SIC/XE system.

SIC Machine Architecture

○ Memory

- Memory consists of 8-bit , 3 consecutive bytes form a word (24 bits), SIC is designed as a 24-bit machine. Total of 32768 bytes (32 KB) in the computer memory.
- A word is addressed by its lowest numbered byte (i.e., addressing starts at byte 0).

○ Register: SIC machines have 5 registers

Register	Number	Known as	Use
A	0	Accumulator	Arithmetic operations
X	1	Index register	Addressing
L	2	Linkage register	jumping to specific memory addresses and storing return addresses
PC	8	Program counter	Contains the address of the next instruction to execute
SW	9	Status word	Contains a variety of information, such as carry or overflow flags

- Register : is a temporary storage area built into a CPU, **registers** are used to pass data from memory to the processor
- Flags : is a value that acts as a signal for a function or process

SIC MACHINE ARCHITECTURE ...

○ Data Formats

- Integers are stored as 24-bit binary numbers
- 2's complement for negative values
- 8-bit ASCII for characters
- No floating point hardware in standard SIC

○ Instruction Formats

- 24-bit format
- x indicates indexed addressing mode

8

1

15

opcode

X

address

SIC MACHINE ARCHITECTURE ...

◦ Addressing modes

- 2 addressing modes available
 - Indicated by x bit in the instruction

Mode	Indication	Target address calculation
Direct	$x = 0$	$TA = \text{address}$
Indexed	$x = 1$	$TA = \text{address} + (x)$

- (x) represents the contents of register x

SIC MACHINE ARCHITECTURE ...

◦ Instruction Set

- **load and store:** To move or store data from accumulator to memory or vice-versa example : LDA, LDX, STA, STX ...
- **Arithmetic operations:** ADD, SUB, MUL, DIV, etc.
 - Used to perform operations on accumulator and memory and store result in accumulator (register A)
- **comparison:** COMP
 - compares the value in register A with a word in memory, this instruction sets a condition code CC to indicate the result
- **conditional jump instructions:** JLT, JEQ, JGT
 - these instructions test the setting of condition code CC and jump accordingly
- **subroutine linkage:** JSUB, RSUB
 - JSUB jumps to the subroutine, placing the return address in register L
 - RSUB returns by jumping to the address contained in register L
 - a **subroutine** is a sequence of program instructions that performs a specific task

SIC MACHINE ARCHITECTURE ...

◉ Input and Output

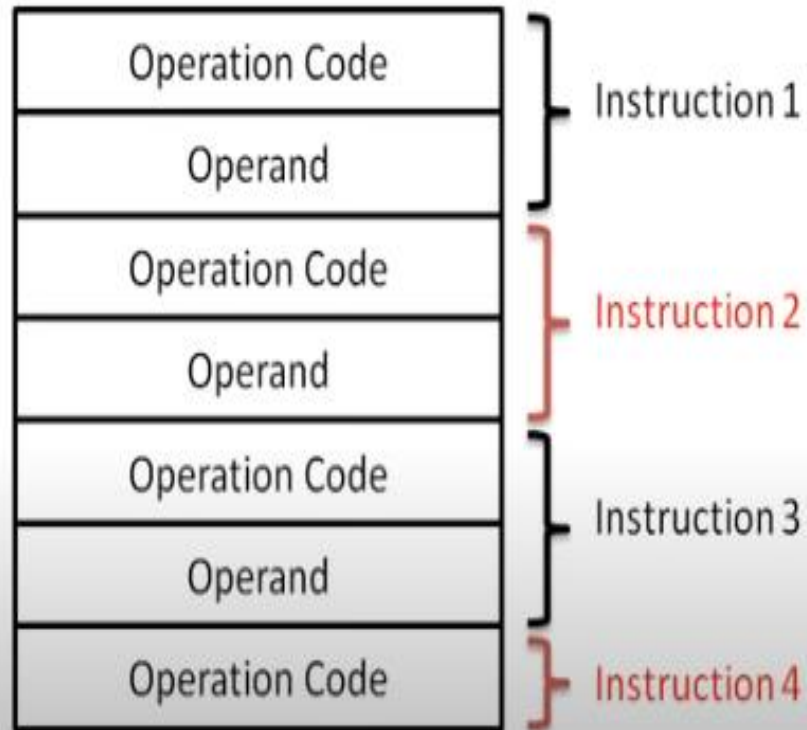
- Input and output are performed by transferring 1 byte at a time to or from the rightmost 8 bits of register A
- **The Test Device (TD):** instruction tests whether the addressed device is ready to send or receive a byte of data
 - *“Less Than” if device is ready; “Equal” if device is busy.*
- **Read Data (RD)**
 - ◉ read a byte from the device to register A
- **Write Data (WD)**
 - ◉ write a byte from register A to the device

MACHINE CODE INSTRUCTION

- **Machine code** is a set of **instructions** executed directly by a **computer's** central processing unit
- Each **instruction** performs a very specific task, such as a load, a jump, or an ALU operation on a unit of data in a CPU register or memory.
- Machine code instruction usually consists operation code(op code) &operands
- Some instructions consists only operation code
- Op code is a mnemonic used to refer to a microprocessor instruction in assembly language
- Operand is A quantity to which an operator is applied (in $3 - x$, the operands of the subtraction operator are 3 and x).
- All operation code and operands are binary code

CON...

A machine code program consists of a sequence of **Op Codes** and **Operands** stored in the computers memory (e.g. RAM)



SIC/XE MACHINE ARCHITECTURE

- Memory
 - Larger Memory
 - 2^{20} bytes (1 megabyte) in the computer memory
 - Leads to change in instruction formats & addressing modes
- Registers
 - Additional registers provided

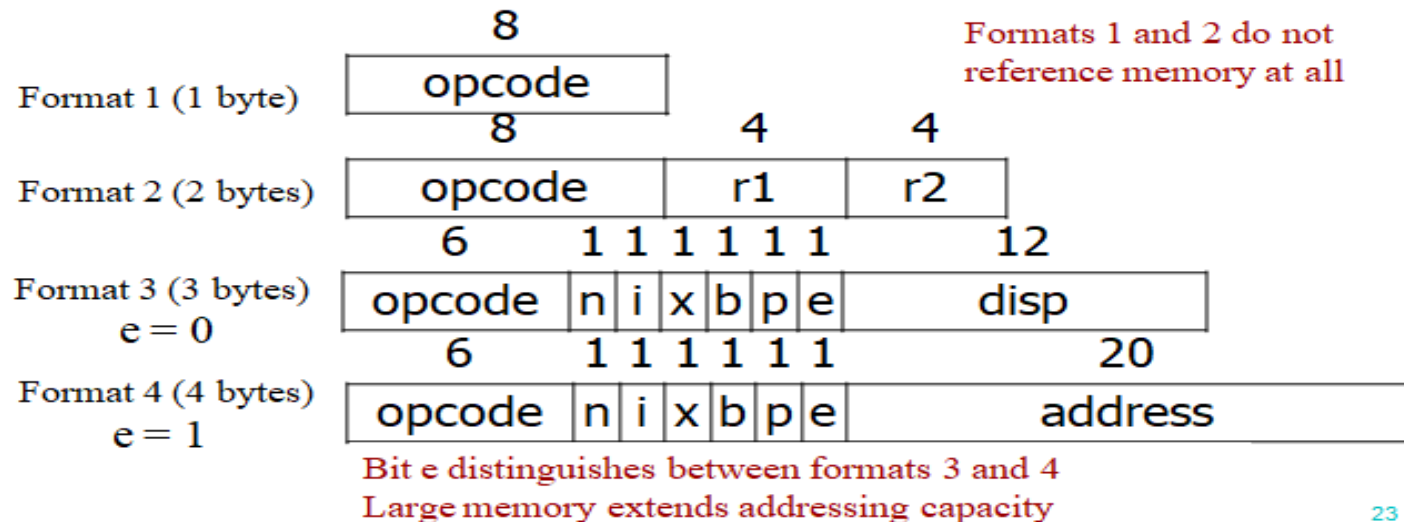
Mnemonic	Number	Use
B	3	Base register – used for addressing
S	4	General purpose register – no special use
T	5	General purpose register – no special use
F	6	Floating point accumulator (48 bits instead 24 bits)

SIC/XE MACHINE ARCHITECTURE...

○ Data format

- 24-bit binary number for integer, 2's complement for negative values
- 48-bit floating-point data type
- The exponent is between 0 and 2047
- value represented = $f \cdot 2^{(\text{exponent}-1024)}$, f is fraction b/n 0 & 1

○ Instruction format



SIC/XE MACHINE ARCHITECTURE...

- **Format 1:** Consists of 8 bits of allocated memory to store instructions. Don't refer memory location
- **Format 2:** Consists of 16 bits of allocated memory to store 8 bits of instructions and two 4-bits segments to store operands(source and destination register)
- **Format 3:** Consists of 6 bits to store an instruction, 6 bits of flag values, and 12 bits of displacement.
- **Format 4:** Only valid on SIC/XE machines, consists of the same elements as format 3, but instead of a 12-bit displacement, stores a 20-bit address.
- Both format 3 and format 4 have six-bit flag values in them, consisting of the following flag bits:
 - **n:** Indirect addressing flag
 - **i:** Immediate addressing flag
 - **x:** Indexed addressing flag
 - **b:** Base address-relative flag
 - **p:** Program counter-relative flag
 - **e:** Extended for Format 4 instruction flag

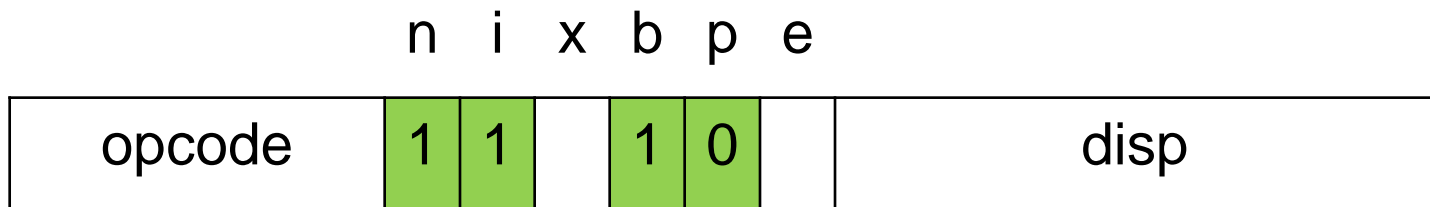
SIC/XE MACHINE ARCHITECTURE...

■ Addressing modes

- **Base relative** ($n=1, i=1, b=1, p=0$)
- **Program-counter relative** ($n=1, i=1, b=0, p=1$)
- **Direct** ($n=1, i=1, b=0, p=0$)
- **Immediate** ($n=0, i=1, x=0$)
- **Indirect** ($n=1, i=0, x=0$)
- **Indexing** (both n & $i = 0$ (SIC) or 1 (SIC/EX), $x=1$)
- **Extended** ($e=1$) for format 4 and $e=0$ for format 3
 - Note: Indexing cannot be used with immediate or indirect addressing.

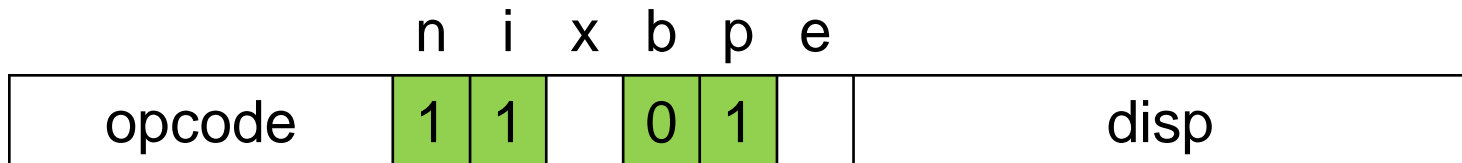
SIC/XE MACHINE...

■ Base Relative Addressing Mode



$n=1, i=1, b=1, p=0$ $TA=(B)+disp$ $(0 \leq disp \leq 4095)$

■ Program-Counter Relative Addressing Mode

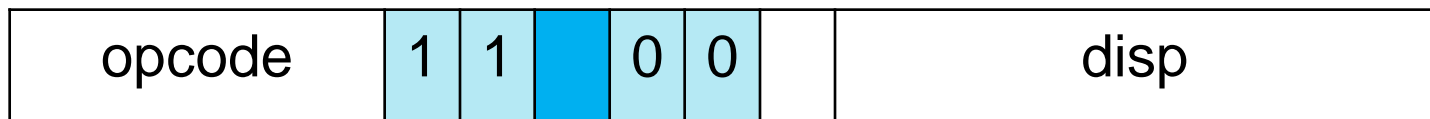


$n=1, i=1, b=0, p=1$ $TA=(PC)+disp$ $(-2048 \leq disp \leq 2047)$

SIC/XE MACHINE...

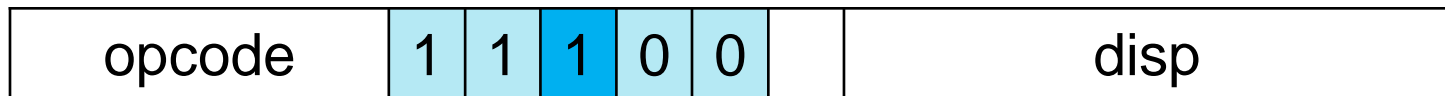
■ Direct Addressing Mode

n i x b p e



$n=1, i=1, b=0, p=0$ $TA=disp$ $(0 \leq disp \leq 4095)$

n i x b p e



$n=1, i=1, b=0, p=0$ $TA=(X)+disp$
(with index addressing mode)

SIC/XE MACHINE...

■ Immediate Addressing Mode

n i x b p e



$n=0, i=1, x=0$

TA = disp

■ Indirect Addressing Mode

n i x b p e



$n=1, i=0, x=0$

TA=(disp)

CON..

- PC-relative simple addressing: **(PC) + disp**
- Base-relative indexed simple addressing: **(B) + disp + (X)**
- PC-relative indirect addressing: **(PC) + disp**
- Immediate addressing: **disp**
- Indirect addressing :**disp**

SIC/XE MACHINE ARCHITECTURE...

○ Instruction Set

- new registers: LDB, STB, etc.
- floating-point arithmetic: ADDF, SUBF, MULF, DIVF
- register move: RMO
- register-register arithmetic: ADDR, SUBR, MULR, DIVR
- supervisor call: SVC
 - generates an interrupt for OS

○ Input/output

- SIO, TIO, HIO: start, test, halt the operation of I/O device

SIC ASSEMBLY SYNTAX

- SIC uses a special assembly language with its own operation codes that hold the hex values needed to assemble and execute programs.
- A sample program is provided below to get an idea of what a SIC program might look like.
- The **first column** represents a **forwarded symbol** that will **store its location in memory**.
- The **second column** denotes either a **SIC instruction** (opcode) or a **constant** value (BYTE or WORD).
- The **third column** takes the symbol value obtained by going through the first column and uses it to run the operation specified in the second column.
- This process creates an object code, and all the object codes are put into an object file to be run by the SIC machine.

SIC AND SIC/XE PROGRAMMING EXAMPLES

- Data movement operations
- Arithmetic operations
- Looping and indexing operations
- Data input and output operations



PROGRAMMING EXAMPLES

WORD – Integer

constant

BYTE – Character

constant

```

LDA      FIVE
STA      ALPHA
LDCH     CHARZ
STCH     C1
.
.
.
ALPHA    RESW      1
FIVE     WORD      5
CHARZ    BYTE     C'Z'
C1       RESB      1
    
```

```

LOAD CONSTANT 5 INTO REGISTER A
STORE IN ALPHA
LOAD CHARACTER 'Z' INTO REGISTER A
STORE IN CHARACTER VARIABLE C1
    
```

```

ONE-WORD VARIABLE
ONE-WORD CONSTANT
ONE-BYTE CONSTANT
ONE-BYTE VARIABLE
    
```

(a)

```

LDA      #5
STA      ALPHA
LDA      #90
STCH     C1
.
.
.
ALPHA    RESW      1
C1       RESB      1
    
```

```

LOAD VALUE 5 INTO REGISTER A
STORE IN ALPHA
LOAD ASCII CODE FOR 'Z' INTO REG A
STORE IN CHARACTER VARIABLE C1
    
```

```

ONE-WORD VARIABLE
ONE-BYTE VARIABLE
    
```

(b)

**RESW - Reserve indicated
number of words for a data
area**

**RESB - Reserve indicated
number of bytes for a data
area**

Sample data movement operations for (a) SIC and

PROGRAMMING EXAMPLES

	LDA	ALPHA	LOAD ALPHA INTO REGISTER A
	ADD	INCR	ADD THE VALUE OF INCR
	SUB	ONE	SUBTRACT 1
	STA	BETA	STORE IN BETA
	LDA	GAMMA	LOAD GAMMA INTO REGISTER A
	ADD	INCR	ADD THE VALUE OF INCR
	SUB	ONE	SUBTRACT 1
	STA	DELTA	STORE IN DELTA
	.		
	.		
	.		
ONE	WORD	1	ONE-WORD CONSTANT
.			ONE-WORD VARIABLES
ALPHA	RESW	1	
BETA	RESW	1	
GAMMA	RESW	1	
DELTA	RESW	1	
INCR	RESW	1	

$$\text{BETA} \leftarrow (\text{ALPHA} + \text{INCR} - 1)$$

$$\text{DELTA} \leftarrow (\text{GAMMA} + \text{INCR} - 1)$$

Fig: Arithmetic operations
SIC

SIC PROGRAMMING EXAMPLE LOOP

	LDA	ZERO	initialize index value to 0
	STA	INDEX	
ADDLP	LDX	INDEX	load index value to reg X
	LDA	ALPHA,X	load word from ALPHA into reg A
	ADD	BETA,X	
	STA	GAMMA,X	store the result in a word in GAMMA
	LDA	INDEX	
	ADD	THREE	add 3 to index value
	STA	INDEX	
	COMP	K300	compare new index value to 300
	JLT	ADDLP	loop if less than 300
	...		
	...		
INDEX	RESW	1	
ALPHA	RESW	100	array variables—100 words each
BETA	RESW	100	
GAMMA	RESW	100	
ZERO	WORD	0	one-word constants
THREE	WORD	3	
K300	WORD	300	

$$\text{Gamma[]} = \text{Alpha[]} + \text{Beta[]}$$

SIC PROGRAMMING EXAMPLE I/O

○ Input and output

INLOOP	TD	INDEV	test input device
	JEQ	INLOOP	loop until device is ready
	RD	INDEV	read one byte into register A
	STCH	DATA	
	.		
OUTLP	.		
	TD	OUTDEV	test output device
	JEQ	OUTLP	loop until device is ready
	LDCH	DATA	
	WD	OUTDEV	write one byte to output device
INDEV	.		
	.		
	BYTE	X' F1'	input device number
	BYTE	X' 05'	output device number
DATA	RESB	1	

SIC/EX PROGRAMMING EXAMPLES

Arithmetic operation.

$BETA \leftarrow (ALPHA + INCR - 1)$

$DELTA \leftarrow (GAMMA + INCR - 1)$

```
LDS      INCR
LDA      ALPHA
ADDR     S,A
SUB      #1
STA      BETA
LDA      GAMMA
ADDR     S,A
SUB      #1
STA      DELTA
```

```
LOAD VALUE OF INCR INTO REGISTER S
LOAD ALPHA INTO REGISTER A
ADD THE VALUE OF INCR
SUBTRACT 1
STORE IN BETA
LOAD GAMMA INTO REGISTER A
ADD THE VALUE OF INCR
SUBTRACT 1
STORE IN DELTA
```

ONE WORD VARIABLES

```
ALPHA    RESW    1
BETA     RESW    1
GAMMA    RESW    1
DELTA    RESW    1
INCR     RESW    1
```

This program will execute faster because it need not load INCR from memory each time when INCR is needed.

SIC/EX Programming Examples

Looping and indexing

	LDS	#3	INITIALIZE REGISTER S TO 3
	LDT	#300	INITIALIZE REGISTER T TO 300
	LDX	#0	INITIALIZE INDEX REGISTER TO 0
ADDLPS	LDA	ALPHA,X	LOAD WORD FROM ALPHA INTO REGISTER A
	ADD	BETA,X	ADD WORD FROM BETA
	STA	GAMMA,X	STORE THE RESULT IN A WORD IN GAMMA
	ADDR	S,X	ADD 3 TO INDEX VALUE
	COMPR	X,T	COMPARE NEW INDEX VALUE TO 300
	JLT	ADDLPS	LOOP IF INDEX VALUE IS LESS THAN 300
	.		
	.		
	.		
.			ARRAY VARIABLES--100 WORDS EACH
ALPHA	RESW	100	
BETA	RESW	100	
GAMMA	RESW	100	

This program will execute faster because it uses register-to-register

This program will execute faster because it uses register-to-register add to reduce memory accesses.

Fig: Looping & indexing
SIC/XE

SIC/EX Programming Examples

Input /output

	JSUB	READ	CALL READ SUBROUTINE
	.		
	.		
	.		
	.		
			SUBROUTINE TO READ 100-BYTE RECORD
READ	LDX	#0	INITIALIZE INDEX REGISTER TO 0
	LDT	#100	INITIALIZE REGISTER T TO 100
RLOOP	TD	INDEV	TEST INPUT DEVICE
	JEQ	RLOOP	LOOP IF DEVICE IS BUSY
	RD	INDEV	READ ONE BYTE INTO REGISTER A
	STCH	RECORD,X	STORE DATA BYTE INTO RECORD
	TIXR	T	ADD 1 TO INDEX AND COMPARE TO 100
	JLT	RLOOP	LOOP IF INDEX IS LESS THAN 100
	RSUB		EXIT FROM SUBROUTINE
	.		
	.		
	.		
INDEV	BYTE	X'F1'	INPUT DEVICE NUMBER
RECORD	RESB	100	100-BYTE BUFFER FOR INPUT RECORD

Fig: Subroutine call & record input operations

THANK YOU!

Question?