

Implementation of Deque using circular array

Deque or Double Ended Queue is a generalized version of **Queue data structure** that allows insert and delete at both ends. In previous post we had discussed introduction of deque. Now in this post we see how we implement deque Using circular array.

Operations on Deque:

Mainly the following four basic operations are performed on queue:

insetFront(): Adds an item at the front of Deque.

insertRear(): Adds an item at the rear of Deque.

deleteFront(): Deletes an item from front of Deque.

deleteRear(): Deletes an item from rear of Deque.

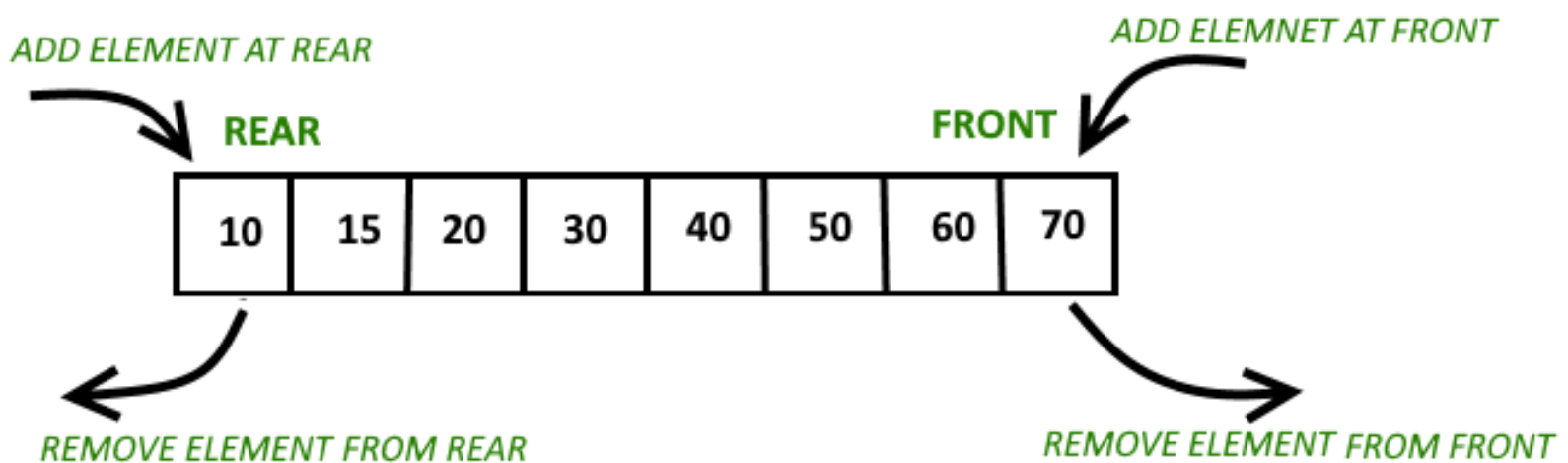
In addition to above operations, following operations are also supported

getFront(): Gets the front item from queue.

getRear(): Gets the last item from queue.

isEmpty(): Checks whether Deque is empty or not.

isFull(): Checks whether Deque is full or not.



Recommended: Please try your approach on **{IDE}** first, before moving on to the solution.

Circular array implementation deque

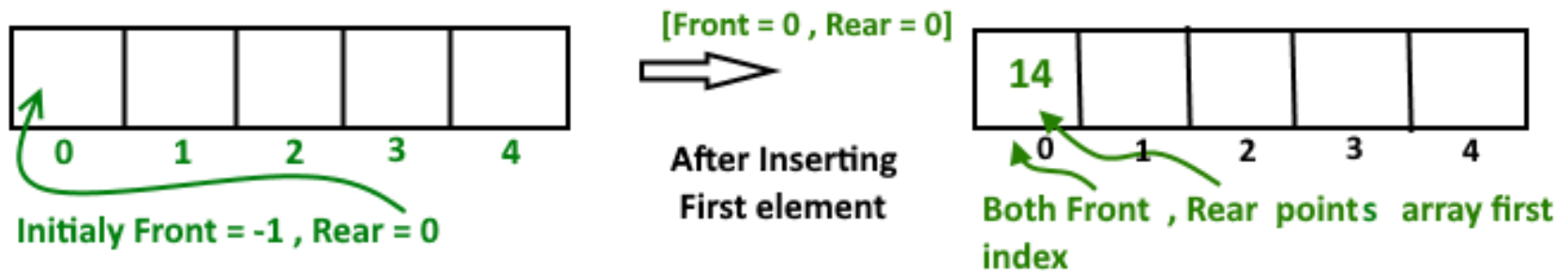
For implementing deque, we need to keep track of two indices, front and rear. We enqueue(push) an item at the rear or the front end of queue and dequeue(pop) an item from both rear and front end.

Working

1. Create an empty array 'arr' of size 'n'

initialize **front = -1** , **rear = 0**

Inserting First element in deque either front end or rear end they both lead to the same result.



After insert **Front** Points = 0 and **Rear** points = 0

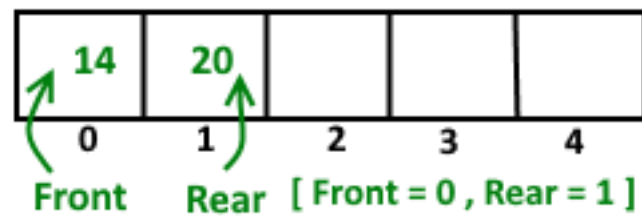
Insert Elements at Rear end

- a). First we check deque if Full or Not
- b). IF `Rear == Size-1`
 then reinitialize `Rear = 0` ;
 Else increment `Rear` by '1'
 and push current key into `Arr[rear] = key`
 Front remain same.

Insert Elements at Front end

- a). First we check deque if Full or Not
- b). IF `Front == 0` || initial position, move Front
 to points last index of array
 `front = size - 1`
 Else decremented front by '1' and push
 current key into `Arr[Front] = key`
 Rear remain same.

Insert element at Rear



Insert element at Front end Now Front points last index

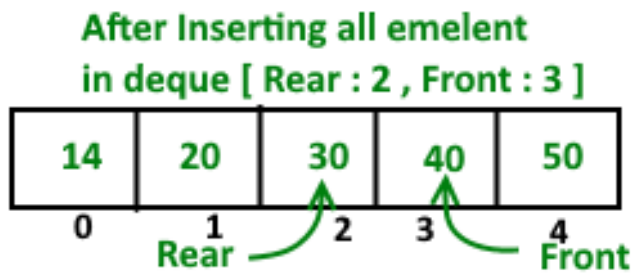


Delete Element From Rear end

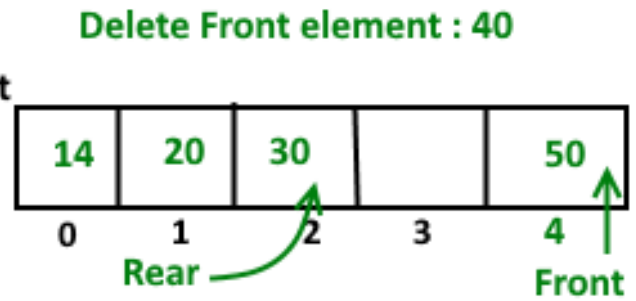
- a). first Check deque is Empty or Not
- b). If deque has only one element
 `front = -1 ; rear = -1 ;`
 Else IF Rear points to the first index of array
 it's means we have to move rear to points
 last index [now first inserted element at
 front end become rear end]
 `rear = size-1 ;`
 Else || decrease rear by '1'
 `rear = rear-1;`

Delete Element From Front end

- a). first Check deque is Empty or Not
- b). If deque has only one element
 `front = -1 ; rear = -1 ;`
 Else IF front points to the last index of the array
 it's means we have no more elements in array so
 we move front to points first index of array
 `front = 0 ;`
 Else || increment Front by '1'
 `front = front+1;`



Delete Element from
Front end , New front



Below c++ implementation of above idea.

```
// C++ implementation of De-queue using circular
```

```
// array
```

```
#include<iostream>
```

```
using namespace std;
```

```
// Maximum size of array or Dequeue
```

```
#define MAX 100
```

```
// A structure to represent a Deque
```

```
class Deque
```

```
{
```

```
    int arr[MAX];
```

```
    int front;
```

```
    int rear;
```

```
    int size;
```

```
public :
```

```
    Deque(int size)
```

```
{
```

```
        front = -1;
```

```
        rear = 0;
```

```
        this->size = size;
```

```
}
```

```
// Operations on Deque:
```

```
void insertfront(int key);
```

```
void insertrear(int key);
```

```
void deletefront();
```

```
void deleterear();
```

```
bool isFull();
```

```
bool isEmpty();
```

```
int getFront();
```

```
int getRear();
```

```
};
```

```
// Checks whether Deque is full or not.
```

```
bool Deque::isFull()
```

```
{
```

```
    return ((front == 0 && rear == size-1)||
            front == rear+1);
```

```
}
```

```
// Checks whether Deque is empty or not.
```

```
bool Deque::isEmpty ()
```

```
{
```

```
    return (front == -1);
```

```
}
```

```

// Inserts an element at front
void Deque::insertfront(int key)
{
    // check whether Deque if full or not
    if (isFull())
    {
        cout << "Overflow\n" << endl;
        return;
    }

    // If queue is initially empty
    if (front == -1)
    {
        front = 0;
        rear = 0;
    }

    // front is at first position of queue
    else if (front == 0)
        front = size - 1 ;

    else // decrement front end by '1'
        front = front-1;

    // insert current element into Deque
    arr[front] = key ;
}

// function to inset element at rear end
// of Deque.
void Deque ::insertrear(int key)
{
    if (isFull())
    {
        cout << " Overflow\n " << endl;
        return;
    }

    // If queue is initially empty
    if (front == -1)
    {
        front = 0;
        rear = 0;
    }

    // rear is at last position of queue
    else if (rear == size-1)
        rear = 0;

    // increment rear end by '1'
    else
        rear = rear+1;

    // insert current element into Deque
    arr[rear] = key ;
}

// Deletes element at front end of Deque
void Deque ::deletefront()
{
    // check whether Deque is empty or not
    if (isEmpty())
    {
        cout << "Queue Underflow\n" << endl;
    }
}

```

```

        return ;
    }

    // Deque has only one element
    if (front == rear)
    {
        front = -1;
        rear = -1;
    }
    else
        // back to initial position
        if (front == size -1)
            front = 0;

        else // increment front by '1' to remove current
            // front value from Deque
            front = front+1;
}

// Delete element at rear end of Deque
void Deque::deleterear()
{
    if (isEmpty())
    {
        cout << " Underflow\n" << endl ;
        return ;
    }

    // Deque has only one element
    if (front == rear)
    {
        front = -1;
        rear = -1;
    }
    else if (rear == 0)
        rear = size-1;
    else
        rear = rear-1;
}

// Returns front element of Deque
int Deque::getFront()
{
    // check whether Deque is empty or not
    if (isEmpty())
    {
        cout << " Underflow\n" << endl;
        return -1 ;
    }
    return arr[front];
}

// function return rear element of Deque
int Deque::getRear()
{
    // check whether Deque is empty or not
    if (isEmpty() || rear < 0)
    {
        cout << " Underflow\n" << endl;
        return -1 ;
    }
    return arr[rear];
}

// Driver program to test above function

```

```

int main()
{
    Deque dq(5);
    cout << "Insert element at rear end : 5 \n";
    dq.insertrear(5);

    cout << "insert element at rear end : 10 \n";
    dq.insertrear(10);

    cout << "get rear element " << " "
         << dq.getRear() << endl;

    dq.deleterear();
    cout << "After delete rear element new rear"
         << " become " << dq.getRear() << endl;

    cout << "insert element at front end \n";
    dq.insertfront(15);

    cout << "get front element " << " "
         << dq.getFront() << endl;

    dq.deletefront();

    cout << "After delete front element new "
         << "front become " << dq.getFront() << endl;
    return 0;
}

```

[Run on IDE](#)

Output:

```

insert element at rear end : 5
insert element at rear end : 10
get rear element : 10
After delete rear element new rear become : 5
insert element at front end : 15
get front element : 15
After delete front element new front become : 5

```

Time Complexity: Time complexity of all operations like insertfront(), insertlast(), deletefront(), delete-last() is $O(1)$.

In next post we will discuss deque implementation using Doubly linked list.

This article is contributed by **Nishant Singh**. If you like GeeksforGeeks and would like to contribute, you can also write an article using contribute.geeksforgeeks.org or mail your article to contribute@geeksforgeeks.org. See your article appearing on the GeeksforGeeks main page and help other Geeks.

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