## <u>Abstract:</u>

analyzing the behavior of numerical methods: Gaussian-elimination, LU decomposition, Gaussian-Jordan and Gauss-Seidel.

# Linear equations system solver

Matlab implementation

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## II. Description and Pseudocode:

## 1. Gaussian-elimination:

Gaussian elimination is an algorithm for solving a system of linear equations. To perform row reduction on a matrix, one uses a sequence of elementary row operations to modify the matrix until the lower left-hand corner of the matrix is filled with zeros. And the by using backward substitutions the solution will be found.

```
gaussElimination(coeff, values, tolerance)
   n = size(coeff)
                           //number of variables
   scale = [ ]
    //array containing scaling factor of each row
   //These 2 for loops are used to get the scaling value of each row.
   for i = 1 : n
       //finding the scaling value of the current row.
       scale(i) = abs(coeff(i, 1));
       for j = 2 : n
           if (abs(coeff(i, j)) > scale(i))
                scale(i) = abs(coeff(i, j))
           end if
   end loop
   [coeff, values, err] = eliminate(coeff, values, scale, n, tolerance);
  //This if statement checks if something wrong has happened as division b zero
  if (err == -1)
       return
   end if
   x = substitute(a, n, b)
end
```

#### 2. Gauss-Seidel:

It is an iterative method used to solve a linear system of equations. It can be applied to any matrix with non-zero elements on the diagonals, convergence is only guaranteed if the matrix is either diagonally dominant.

If we have a system of equations like:  $a_{11}X_1 + a_{12}X_2 + a_{13}X_3 = b_1$ 

$$a_{21}X_1 + a_{22}X_2 + a_{23}X_3 = b_2$$

$$a_{31}X_1 + a_{32}X_2 + a_{33}X_3 = b_3$$

Then the solution of the variables for the new iterations will be:

$$X_1^{new} = \frac{b_1 - a_{12} X_2^{old} - a_{13} X_3^{old}}{a_{11}}, X_2^{new} = \frac{b_2 - a_{21} X_1^{new} - a_{23} X_3^{old}}{a_{22}}, X_3^{new} = \frac{b_3 - a_{31} X_1^{new} - a_{32} X_2^{new}}{a_{33}}$$

Iterations are repeated until the convergence criterion is satisfied:

$$\left| \varepsilon_{a,i} \right| = \left| \frac{X_i^{j} - X_i^{j-1}}{X_i^{j}} \right| \leq \varepsilon_s$$

If the coefficient matrix A Diagonally-Dominant then Gauss-Seidel is guaranteed to converge.

$$|a_{ii}| \geq \sum_{j=1, j\neq i}^{n} |a_{i,j}|$$

```
seidel(a, b, x, iterations, eps)
   //reordering the equations such that the matrix becomes diagonally
   //a is coefficient matrix, b is values if equations array, x is the initial guesses
   [a, b, flag] = reorder(a, b);
   if flag == 1
       "Can't be diagonally-dominant"
   end if
   n = size(a) //n is the number of equations
   iterations matrix = zeros(0, n * 2);
         //2 columns for each variable one for the value at
         //this iteration and the other is error
   for i = 1 : n
       iterations matrix(1, 2 * i - 1) = x(i)
       iterations_matrix(1, 2 * i) = 0
   end loop
   for iter = 2 : iterations
```

```
stop = 1
    //calculates each X at this iteration and the error
    for var_x = 1 : n
        sum = b(var x);
        for i = 1 : sz
            if i ~= var_x
                index = iter;
                if (var_x < i)</pre>
                     index = iter - 1
                end if
                sum = sum - a(var_x, i) * iterations_matrix(index, 2 * i - 1)
            end if
        end loop
        if a(var_x, var_x) ~= 0
            iterations_matrix(iter, 2 * var_x - 1) = sum / a(var_x, var_x)
             //calculate
                             X_(var_x)
            iterations_matrix(iter, 2 * var_x) =
            abs((iterations_matrix(iter, 2 * var_x - 1) -
                iterations_matrix(iter - 1, 2 * var_x 1)) /
             iterations_matrix(iter, 2 * var_x - 1)) //calculate the |\varepsilon_a(a, iter)|
        else
            return "Division by Zero"
        if iterations_matrix(iter, 2 * var_x) > eps //cannot stop if |\epsilon_a(a,i)| > \epsilon_s
            stop = ∅;
        end if
    end loop
   if stop == 1
        break;
  end if
end loop
row_size = size(iterations_matrix, 2);
last_row_index = size(iterations_matrix, 1);
final_ans = zeros(1, row_size / 2);
//copy last row of values only in the final_ans array.
```

```
for i = 1: row_size
    if (mod(i, 2) == 1)
        final_ans(ceil(i / 2)) = iterations_matrix(last_row_index, i);
    end if
end loop
end
```

#### 3. Gaussian-Jordan:

This method allows the isolation of the coefficients of a system of linear equations. In Gauss-Jordan method, given matrix can be fetched to row echelon form and made simpler. We can say that the transformation of augmented matrix of the given system into reduced row-echelon(scaling) form is done by the use of row operations.

In Gauss-Jordan method, elementary row operations are utilized in order to solve given system of linear equations.

#### Elementary row operations:

- i) Switching of Rows
- ii) Row multiplication
- iii) Addition of Rows

#### Row echelon form:

- 1) The first nonzero element of a row must be 1 which is known as leading 1.
- 2) The leading 1 of a row must be positioned at the right side of the leading 1 of its previous row.
- 3) In case of a column containing a leading 1, all the remaining elements in that column must be 0.
- **4)** At the bottom of the matrix, there is a rows having only zeros as the elements. If any of the above rules are violated, the matrix is not said to be in reduced row echelon form.

```
exception = MException('jordan:division_by_zero', 'Division by zero');
            throw(exception);
        end
        for j = 1:widthAug
            augmented(i, j) = augmented(i, j) / pivot;
        end
        for j = 1:length
            if i ~= j
                augmented(j, :) = augmented(j, :) - augmented(i, :) .* augmented(j, i)
                ./ augmented(i, i);
            end
        end
    end
   for i = 1:length
        x(i) = augmented(i, widthAug);
   end
end
```

## 4. LU Decomposition:

Let *A* be a square matrix. An **LU factorization** refers to the factorization of A, with proper row and/or column orderings or permutations, into two factors, a lower triangular matrix *L* and an upper triangular matrix *U*,

#### A=L U

In the lower triangular matrix all elements above the diagonal are zero, in the upper triangular matrix, all the elements below the diagonal are zero. For example, for a 3-by-3 matrix A, its LU decomposition looks like this:

$$\begin{bmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \end{bmatrix} = \begin{bmatrix} l_{11} & 0 & 0 \\ l_{21} & l_{22} & 0 \\ l_{31} & l_{32} & l_{33} \end{bmatrix} \begin{bmatrix} u_{11} & u_{12} & u_{13} \\ 0 & u_{22} & u_{23} \\ 0 & 0 & u_{33} \end{bmatrix}.$$

Without a proper ordering or permutations in the matrix, the factorization may fail to materialize. For example, it is easy to verify (by expanding the matrix multiplication)

that . If , then at least one of and has to be zero, which implies either *L* or *U* is singular. This is impossible if *A* is nonsingular. This is a procedural problem. It can be removed by simply reordering the rows of *A* so that the first element of the permuted matrix is nonzero. The same problem in subsequent factorization steps can be removed the same way; see the basic procedure below.

```
function [method_name, x, err] = luDecomposition(a, b, tol)
    method_name = 'LU Decomposition';
    err = 0;
    n = size(a, 1);
    [a, ind, err] = decompose(a, tol);
    if err == - 1
        return;
    end
    x = luSubstitute(a, b, ind, n);
end
```

#### 5. Intermediate Methods:

#### -Eliminate Method:

This method is used to modify the matrix of coefficients until the lower left-hand corner of the matrix is filled with zeros. This operation can be done by scaling each row first and then subtracting a multiple of that row from the rows below it such that the lower left corner becomes a zero triangle.

```
eliminate(coeff, values, scale, n, tolerance)
    err = 0 //there is no error at the beginning such error turns to -1 when there is a

value less than tolerance
    //These 2 for loops are used to convert the coefficients matrix into a triangular matrix
    for i = 1 : n - 1
        [coeff, values, scale] = pivot(coeff, values, scale, n, i)
        if (abs(a(i, i) / scale(i)) < tolerance)
        err = -1
        return
    end if
    for j = i + 1 : n
        factor = coeff(j, i) / coeff(i, i)
        for k = i : n
            coeff(j, k) = coeff(j, k) - factor * coeff(i, k)
        end if</pre>
```

```
values(j) = values(j) - factor * values(i)
end loop
end loop
if abs(coeff(n, n)) < tolerance
err = - 1;
end if</pre>
```

#### -Pivoting Method:

This method is used such that the diagonal(pivots) after scaling them as large as possible. That is done to avoid dividing by zero as possible but there is still a probability of division by zero. This method swaps 2 rows such that the pivot after scaling it is as large as possible.

```
pivot(coeff, values, scale, n, row)
   //setting the pivot as the normal current pivot
   piv = row
   largest = abs(coeff(row, row) / scale(k))
   //This for loop finds the largest scaled pivot
   for i = row + 1 : n
       temp = abs(coeff(i, row) / scale(i))
       if temp > largest
            piv = i
            largest = temp
        end if
   end loop
   //This for loop swaps the row with the largest scaled pivot with the current row
   for i = row : n
       temp = coeff(k, i)
        coeff(k, i) = coeff(piv, i);
        coeff(piv, i) = temp;
   end loop
   // swap the values of the current row and the current row
   temp = values(row);
   values(row) = values(piv);
   values(piv) = temp;
```

```
//swap the scaling values if the current row and the pivot row
temp = scale(row);
scale(k) = scale(piv);
scale(piv) = temp;
end
```

#### -Substitution Method:

This method uses backward substitution to calculate the values of different unknowns. Such that at the last row there is only one unknown which will be obtained. The row before the last has 2 variables one is unknown and the other is known so we can get its value and so on.

#### -Reorder Method:

This method is used to reorder the equations in way that makes the matric Diagonally-Dominant such that it will be guaranteed that it will converge. Its algorithm runs in way such that in each row the max coefficient of the row is determined and then get the number of the column(j) of that coefficient and then we swap  $row_i$  with  $row_i$ .

```
reOrder(unOrdered, b)
   n = size(unOrdered) //unknowns of the system of equations
   orderdB = [] //ordered values of each equation
```

```
taken = [] //check if the row containing this already having an equation or is
empty
   orderd = zeros(0, n * 2) //the ordered matrix
   err = 0
   for i = 1:n
        taken(i) = 0 //initialize all the rows to be empty
   end loop
    //reoders the matrix
   for i = 1:n
        maxIndex = 1
        maxElement = abs(unOrdered(i, 1))
        //select the row to be filled in the ordered matrix
        for j = 1:n
            if abs(unOrdered(i, j)) >= maxElement
                maxIndex = j;
                maxElement = abs(unOrdered(i, j))
            end if
        end loop
        if taken(maxIndex) == 1
            err = 1
            ordered = unordered
            orderedB = b
            break
        end if
        taken(maxIndex) = 1
        //fill the selected row of the ordered matrix with <code>[row] _i</code>
        for j = 1:n
            ordered(maxIndex, j) = unOrdered(i, j)
        end loop
        orderdB(maxIndex) = b(i)
    end
end
```

III. Analysis for the behavior of different examples using the analysis template, conclusion about the behavior of each method.

## **Example 1:**

## Number of equations: 7

```
147*x1 + 3*x2 - 1*x3 + 7*x4 + 6*x5 - 6*x6 - 7*x7 == 97

7*x1 + 147*x2 - 5*x3 - 6*x4 - 0*x5 + 9*x6 - 2*x7 == 146

8*x1 - 5*x2 + 196*x3 + 4*x4 - 1*x5 + 1*x6 - 9*x7 == 12

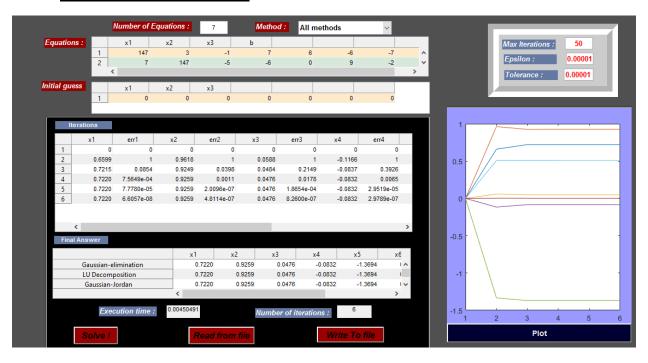
5*x1 + 0*x2 + 0*x3 - 392*x4 - 6*x5 + 9*x6 - 5*x7 == 49

6*x1 - 6*x2 + 2*x3 - 0*x4 - 98*x5 - 8*x6 - 3*x7 == 129

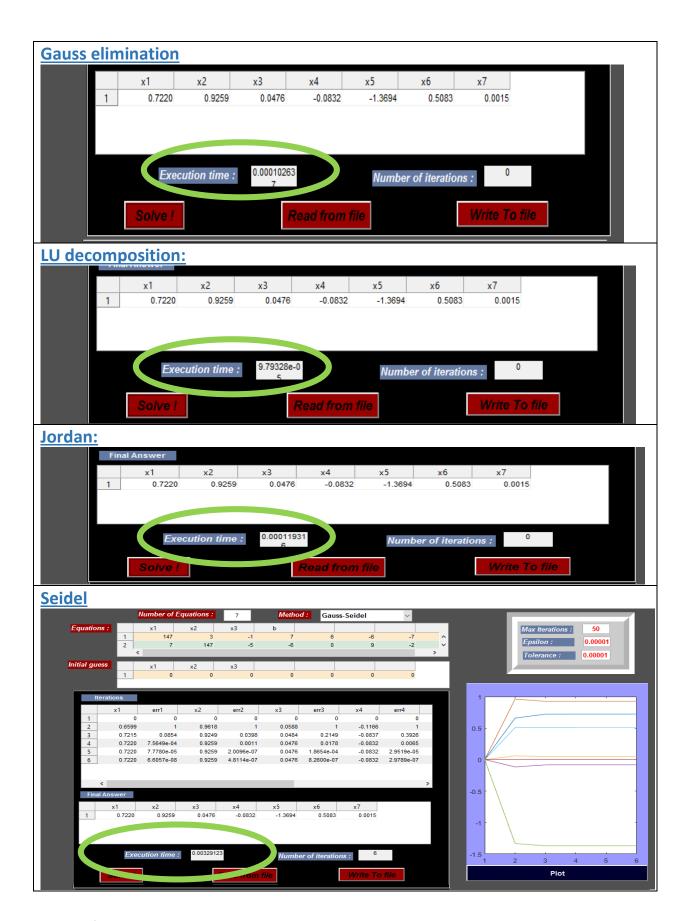
3*x1 + 5*x2 - 8*x3 + 9*x4 + 6*x5 + 245*x6 + 5*x7 == 122

2*x1 + 0*x2 + 2*x3 - 9*x4 - 5*x5 + 5*x6 - 441*x7 == 11
```

## 1. Solve using All methods:



In this example: LU decomposition is faster



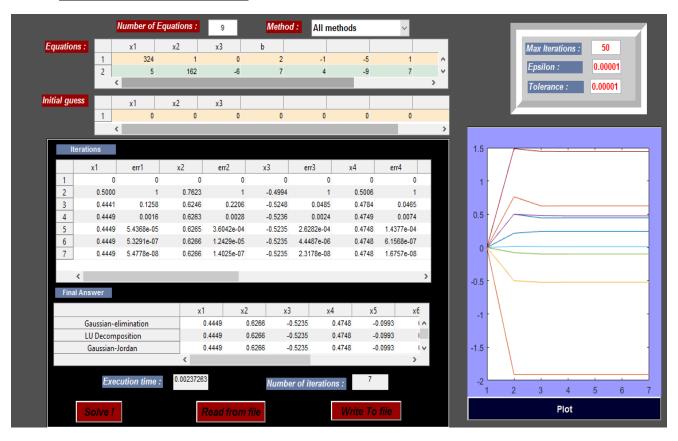
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## **Example 2:**

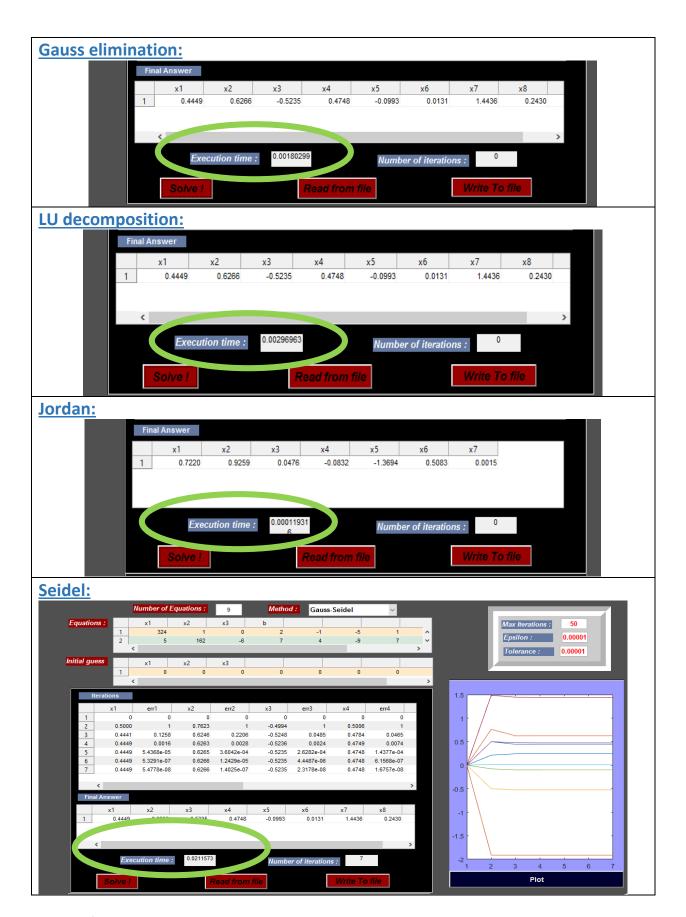
### Number of equations:9

```
324*x1 + 1*x2 + 0*x3 + 2*x4 - 1*x5 - 5*x6 + 1*x7 - 2*x8 - 8*x9 == 162
5*x1 + 162*x2 - 6*x3 + 7*x4 + 4*x5 - 9*x6 + 7*x7 + 2*x8 - 3*x9 == 126
4*x1 - 9*x2 - 218*x3 - 8*x4 + 3*x5 - 0*x6 - 2*x7 + 3*x8 - 0*x9 == 104
9*x1 - 8*x2 - 3*x3 + 162*x4 - 1*x5 - 7*x6 - 9*x7 + 5*x8 - 8*x9 == 81
1*x1 - 0*x2 - 8*x3 - 2*x4 - 729*x5 + 6*x6 - 9*x7 - 5*x8 + 1*x9 == 60
6*x1 + 2*x2 - 2*x3 + 4*x4 + 5*x5 + 729*x6 - 4*x7 + 9*x8 - 4*x9 == 20
6*x1 + 6*x2 - 0*x3 + 6*x4 - 0*x5 + 0*x6 + 81*x7 + 4*x8 - 2*x9 == 131
9*x1 - 0*x2 - 8*x3 + 4*x4 - 1*x5 - 2*x6 + 2*x7 + 405*x8 + 6*x9 == 100
6*x1 + 3*x2 - 8*x3 - 9*x4 - 3*x5 - 9*x6 - 1*x7 + 4*x8 - 81*x9 == 159
```

## 1. Solve using All methods:



# Gauss elimination is faster in this example



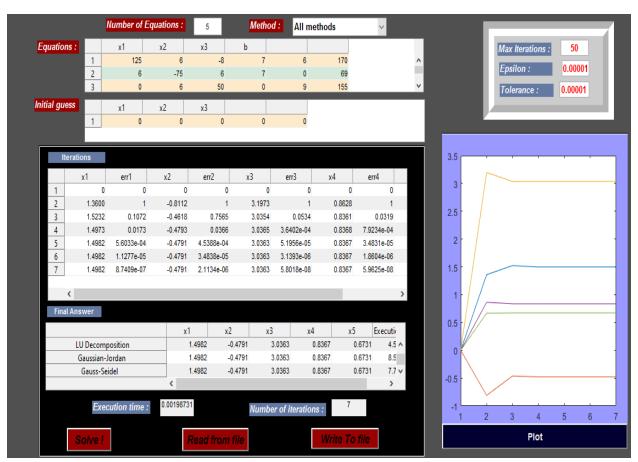
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## **Example 3:**

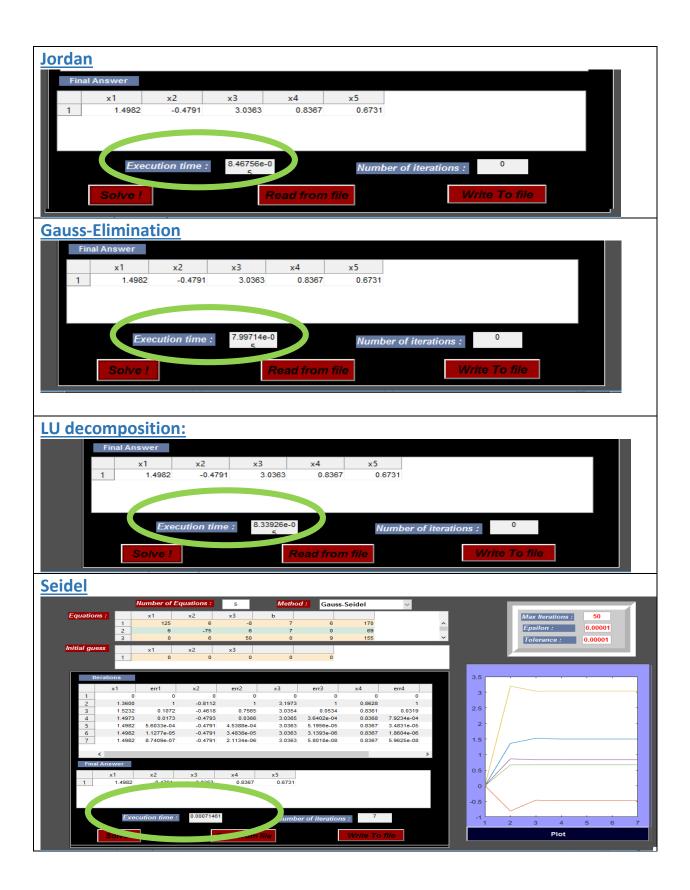
## Number of equations:5

```
125*x1 + 6*x2 - 8*x3 + 7*x4 + 6*x5 == 170
6*x1 - 75*x2 + 6*x3 + 7*x4 - 0*x5 == 69
0*x1 + 6*x2 + 50*x3 - 0*x4 + 9*x5 == 155
1*x1 + 3*x2 - 9*x3 + 125*x4 + 1*x5 == 78
6*x1 - 5*x2 - 0*x3 - 8*x4 + 100*x5 == 72
```

## 1. Solve using All methods:



In this example: Jordan is faster



By comparing the execution time for each method, we can analyze the behavior of each method in each example

IV. Problematic functions, the reason for their misbehavior and suggestions (if exists).

#### **Pitfalls of Gauss-elimination:**

- 1-<u>Division by zero:</u> It is possible that during both elimination and backward substitution phases a division by zero can occur. It can be enhanced by using pivoting and scaling but still there is probability of division by zero.
- 2-**Round-off errors**: computers carry only a limited number of significant digits, round-off errors will occur and they will *propagate* from one iteration to the next. This problem is clearly obvious when there are a great number of equations are to be solved. It is advised to use double precision to decrease this problem but it leads to slow computation. It preferred to substitute the results back into the original equations and check whether a substantial error has occurred.
- 3-<u>III-conditioned:</u> small changes in the coefficients results in large changes in the solution. It is advised to scale the coefficients (to multiply all the coefficients of some equation with some number) so that the small changes in the coefficients can appear. Time complexity of gauss-elimination is  $O(N^3)$ .

#### **Improvement strategies:**

The division by zero case is handled using pivoting and scaling to reduce round off errors Round off errors won't appear due to high accuracy.

#### **Pitfalls of Gauss-Seidel:**

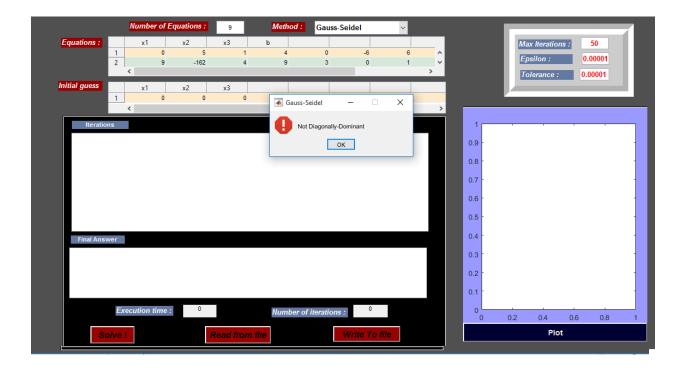
Not all system of equations will converge. As a suggestion to fix that is to reorder the equations in a way such that it will be **Diagonally-Dominant**. A Diagonally-Dominant system of equations is a system of equations that will always converges. A system of equations is called Diagonally-Dominant if  $|a_{ii}| \geq \sum_{j=1, j \neq i}^{n} |a_{i,j}|$ .

The time complexity of Gauss-Seidel is  $O(N^2)$ .

But this time complexity is multiplied by the number of iterations which might be infinite because it may diverge.

## **Example:**

```
Number of equations :9 0*x1 + 5*x2 + 1*x3 + 4*x4 + 0*x5 - 6*x6 + 6*x7 + 8*x8 - 1*x9 == 129
9*x1 - 162*x2 + 4*x3 + 9*x4 + 3*x5 + 0*x6 + 1*x7 - 3*x8 - 1*x9 == 170
3*x1 + 3*x2 + 0*x3 - 7*x4 - 3*x5 - 7*x6 + 4*x7 - 1*x8 + 7*x9 == 2
1*x1 + 7*x2 - 4*x3 + 81*x4 + 8*x5 - 7*x6 - 8*x7 + 9*x8 + 4*x9 == 5
7*x1 - 9*x2 + 1*x3 + 6*x4 + 648*x5 - 7*x6 + 3*x7 + 1*x8 + 5*x9 == 154
7*x1 + 3*x2 + 0*x3 - 6*x4 + 9*x5 + 162*x6 + 3*x7 + 6*x8 - 1*x9 == 57
8*x1 - 9*x2 + 9*x3 - 0*x4 - 4*x5 - 4*x6 - 405*x7 - 9*x8 - 9*x9 == 50
0*x1 + 3*x2 - 6*x3 - 4*x4 - 0*x5 - 8*x6 - 5*x7 + 405*x8 - 3*x9 == 170
8*x1 - 9*x2 + 9*x3 + 0*x4 - 2*x5 - 4*x6 - 8*x7 + 8*x8 + 405*x9 == 35
```



# **Pitfalls of LU decomposition:**

- 1. Matrices don't always have an LU decomposition. Sometimes it is impossible to write a matrix in the form "lower triangular" × "upper triangular".
- 2. in LU decomposition equation matrix is turned to upper and lower matrices L and U A = L \* U but as a result of "round-off error" when we multiply L \* U we don't get the equation matrix

Multiplication of lower and upper matrices won't give the original matrix

3. Not all matrices have an inverse (in case of getting inverse using LU decomposition)

4. Effect of pivoting: we have to remember swapped rows in the original matrix after decomposing into lower and upper

#### **Pitfalls of Gauss Jordan:**

- 1.Division by zero
- 2. round off error
- 3. ill conditioned systems

Same as those explained in Gaussian elimination

#### Improvement strategies:

Same as those used in the Gauss elimination

-Use pivoting and scaling to avoid division-by-zero and to reduce round-off error.

v. Sample runs and snapshots from GUI

## Number of equations: 7

```
147*x1 + 3*x2 - 1*x3 + 7*x4 + 6*x5 - 6*x6 - 7*x7 == 97
7*x1 + 147*x2 - 5*x3 - 6*x4 - 0*x5 + 9*x6 - 2*x7 == 146
8*x1 - 5*x2 + 196*x3 + 4*x4 - 1*x5 + 1*x6 - 9*x7 == 12
5*x1 + 0*x2 + 0*x3 - 392*x4 - 6*x5 + 9*x6 - 5*x7 == 49
6*x1 - 6*x2 + 2*x3 - 0*x4 - 98*x5 - 8*x6 - 3*x7 == 129
3*x1 + 5*x2 - 8*x3 + 9*x4 + 6*x5 + 245*x6 + 5*x7 == 122
2*x1 + 0*x2 + 2*x3 - 9*x4 - 5*x5 + 5*x6 - 441*x7 == 11
```

