

Triangulated Lecture 9

0. Recap

Left / right localization systems in categories.

Left system:

1) $\text{id}_x \in S \quad \forall x \in \mathcal{C}, S \text{ closed under comp}$

2)
$$\begin{array}{ccc} X & \xrightarrow{f} & Y \\ s \downarrow & \downarrow t & s, t \in S \\ Z & \dashrightarrow w \\ g & & \end{array}$$

3)
$$X \xrightarrow{g} Y \xrightarrow{f} Z \dashrightarrow w$$

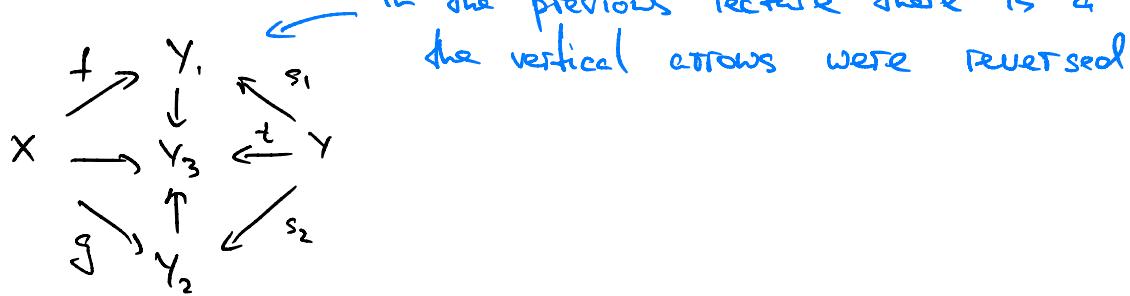
Dually - right localization system.

Define $S^{-1}\mathcal{C}$ (also denoted $\mathcal{C}[S^{-1}]$):

$$\text{Ob} = \text{Ob } \mathcal{C}$$

$$\text{Hom}(X, Y) = \left\{ \begin{array}{c} f: X \rightarrow Y \\ s: Y \rightarrow X \end{array} \mid s \in S \right\} / \sim$$

Equivalence relation:



Need to show that it's an equivalence relation,
define composition (show that it's well-def).

In order to show transitivity of \sim , need to show
that



Two steps : 1) use $\gamma \rightarrow \gamma_1$
 \downarrow |
 $\gamma_2 \rightsquigarrow \gamma_4$

2) need to coequalize $\gamma \rightarrow \gamma_3 \rightrightarrows \gamma_4 \rightarrow \gamma_5$.

Application: construct A/B , B - Serre subcategory
 (full, closed under extensions).

Defined B -iso's as $f: A \rightarrow B$ s.t. $\text{ker } f, \text{coker } f \in B$.

One checks that B -iso - both left & right localization system. $A/B \neq A\{B\text{-iso}'\}$ satisfy the same UP.

Q: Why is $A\{B\text{-iso}'\}$ abelian? Easiest way - check that A/B is abelian (you have explicit descriptions of B -iso (B -mono/ B -epi)).

Finally, we defined $\mathcal{D}(\mathcal{A}) = \mathcal{C}(\mathcal{A}) [Q^{-1}]$

category of class of
complexes quasi-isom's

1. Complexes in abelian categories

Def $\mathcal{C}(\mathcal{A})$ consists of objects

$$\dots \xrightarrow{d^{i-1}} A^i \xrightarrow{d^i} A^{i+1} \xrightarrow{d^{i+1}} \dots \quad d^{i+1} \circ d^i = 0 \quad \forall i$$

morphisms

$$\begin{array}{ccccccc} \dots & \xrightarrow{\quad} & A^i & \xrightarrow{d} & A^{i+1} & \xrightarrow{\quad} & \dots \\ & & \downarrow f^i & \lhd & \downarrow f^{i+1} & & \\ \dots & \xrightarrow{\quad} & B^i & \xrightarrow{\quad} & B^{i+1} & \xrightarrow{\quad} & \dots \end{array}$$

composition is obvious.

Ex $\mathcal{C}(\mathcal{A})$ is abelian if \mathcal{A} is abelian.

Sketch \oplus are defined term-wise: $A^i, B^i \in C(\mathcal{A})$, then

$$A^i \oplus B^i : \rightarrow A^i \oplus B^i \rightarrow A^{i+1} \oplus B^{i+1} \rightarrow$$

$\ker f$ & $\text{Coker } f$ are also term-wise!

$$\begin{array}{ccccc} & \overset{\circ}{\downarrow} & & \overset{\circ}{\downarrow} & \\ \rightarrow & k^i & \rightarrow & k^{i+1} & \text{ker } f \\ \downarrow & & \downarrow & & \\ \rightarrow & A^i & \rightarrow & A^{i+1} & \\ \downarrow f^i & & \downarrow f^{i+1} & & \\ \rightarrow & B^i & \rightarrow & B^{i+1} & \\ \downarrow & & \downarrow & & \\ \rightarrow & C^i & \rightarrow & C^{i+1} & \text{Coker } f \\ \downarrow & & \downarrow & & \\ 0 & & 0 & & \end{array}$$

Please, check all the details.

$$A^\cdot \in C(\mathcal{A}) \rightsquigarrow H^i(A^\cdot) = \ker d^i / \text{Im } d^{i-1}$$

$\ker d^i = Z^i(A^\cdot)$ called cycles

$\text{Im } d^{i-1} = B^i(A^\cdot)$ called boundaries

Get a collection of functors $H^i : C(\mathcal{A}) \rightarrow \mathcal{A}$.

Prop Let $0 \rightarrow X^\cdot \rightarrow Y^\cdot \rightarrow Z^\cdot \rightarrow 0$ be a short exact sequence in $C(\mathcal{A})$. There exist morphisms $s : H^i(Z^\cdot) \rightarrow H^{i+1}(Z^\cdot)$ s.t. the following is exact

$$\dots \xrightarrow{s} H^i(X^\cdot) \rightarrow H^i(Y^\cdot) \rightarrow H^i(Z^\cdot) \xrightarrow{s} H^{i+1}(X^\cdot) \rightarrow \dots$$

and given $0 \rightarrow X_1^\cdot \rightarrow Y_1^\cdot \rightarrow Z_1^\cdot \rightarrow 0$
 $0 \rightarrow X_2^\cdot \rightarrow Y_2^\cdot \rightarrow Z_2^\cdot \rightarrow 0$

there is a morphism of LES's \Leftrightarrow

all the squares

$$h^i(z_i) \xrightarrow{s} h^{i+1}(x_i)$$

$$\downarrow$$

$$h^i(z_i) \xrightarrow{s} h^{i+1}(x_i)$$

$$\downarrow$$

commute.

Pf $0 \rightarrow x \rightarrow y \rightarrow z \rightarrow 0$ is a SES, thus

$$0 \rightarrow z^i(x) \rightarrow z^i(y) \rightarrow z^i(z)$$

$$\downarrow$$

$$\downarrow$$

$$\downarrow$$

$$0 \rightarrow x^i \rightarrow y^i \rightarrow z^i \rightarrow 0$$

$$\downarrow$$

$$\downarrow$$

$$\downarrow$$

$$0 \rightarrow x^{i+1} \rightarrow y^{i+1} \rightarrow z^{i+1} \rightarrow 0$$

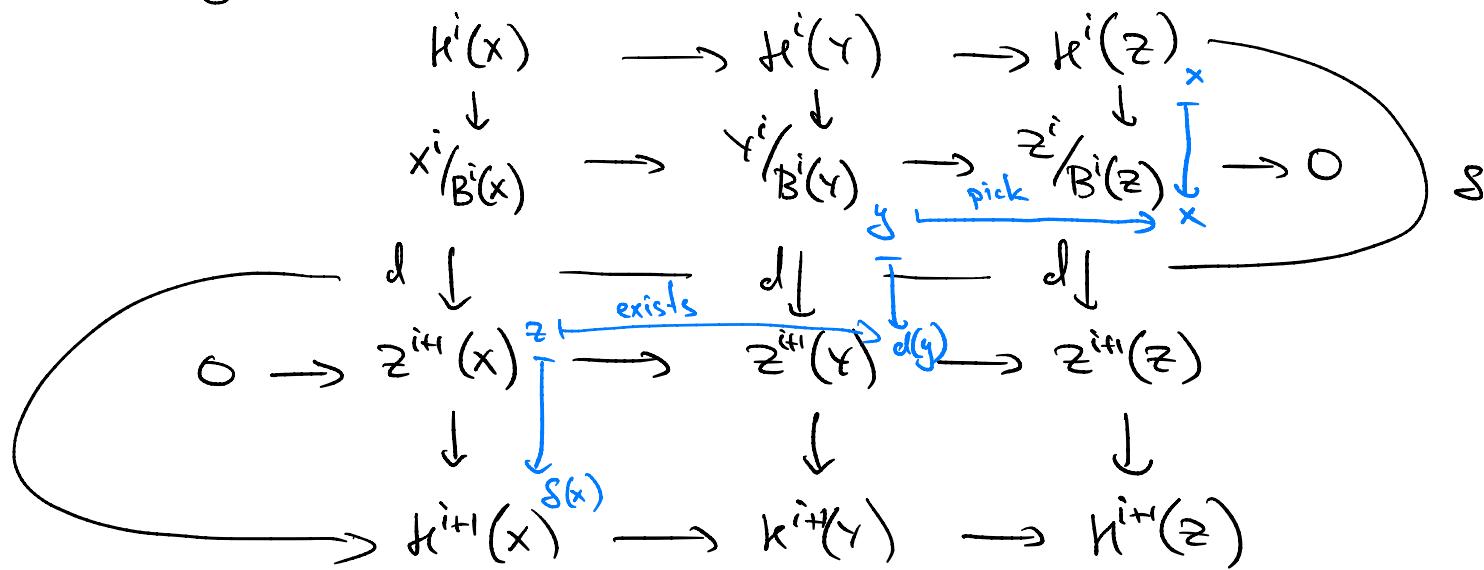
$$\downarrow$$

$$\downarrow$$

$$\downarrow$$

$$x^{i+1}/B^{i+1}(x) \rightarrow y^{i+1}/B^{i+1}(y) \rightarrow z^{i+1}/B^{i+1}(z) \rightarrow 0$$

Stick these sequences into another s -lemma type diagram:



How is S constructed?

Conclusion:

$$\begin{array}{ccccccc}
 0 & \xrightarrow{x^i} & y^i & \xrightarrow{z^i} & 0 \\
 \downarrow & & \downarrow & & \downarrow \\
 0 & \xrightarrow{x^{i+1}} & y^{i+1} & \xrightarrow{z^{i+1}} & 0
 \end{array}$$

$z \in f^i(z) \rightsquigarrow$

- 1) pick a representative $\tilde{z} \in z^i$
- 2) pull it back to y^i
- 3) apply d_y^i
- 4) pull it back to x^{i+1} .

□

Recall that long time ago we defined covers of morphisms.

Shift functors $\Sigma^n: C(\mathcal{A}) \rightarrow C(\mathcal{A})$

$$(x^{\cdot} \Sigma^n)^i = x^{i+n}, \quad d_{x \Sigma^n} = (-)^n d_x.$$

Gives an action of \mathbb{Z} on $C(\mathcal{A})$: $\mathbb{Z} \rightarrow \text{Aut}(C(\mathcal{A}))$.

Cone $f: X^i \rightarrow Y^i$ in $C(\mathcal{A}) \rightsquigarrow C(f) \in C(\mathcal{A})$

$$C(f) = X^{i+1} \oplus Y^i, \quad d = \begin{pmatrix} -d_X^{i+1} & 0 \\ f^i & d_Y^i \end{pmatrix}$$

Ex Check that $d^2 = 0$

More importantly, there is a SES of complexes:

$$0 \rightarrow Y^* \rightarrow C(f) \rightarrow X[\Sigma\beta] \rightarrow 0$$

The associated LES of cohomology is

$$\dots \rightarrow H^i(Y^*) \rightarrow H^i(C(f)) \rightarrow H^i(X[\Sigma\beta]) \xrightarrow{\delta} H^{i+1}(Y) \rightarrow \dots$$

\Downarrow

$$H^{i+k}(X)$$

Exc Using the explicit description of δ , check
that for $0 \rightarrow Y \rightarrow C(f) \rightarrow X[\Sigma\beta] \rightarrow 0$
the connecting homomorphism is $\delta = f_*$,
the morphism $H^i(X) \rightarrow H^i(Y)$!

Recall that $f: X \rightarrow Y$ is a quasi-iso if
 $H^i(f): H^i(X) \rightarrow H^i(Y)$ - iso $\forall i \in \mathbb{Z}$.

We conclude that $f: X \rightarrow Y$ - quasi-iso \Leftrightarrow
 $\Leftrightarrow H^i(C(f)) = 0 \quad \forall i.$

$$H^i(f)$$

$$\tilde{\rightarrow} \dots \rightarrow \underbrace{H^i(C(f))}_{\text{must be } 0!} \rightarrow H^{i+1}(X) \xrightarrow{\sim} H^{i+1}(Y) \rightarrow H^{i+1}(C(f)) \rightarrow \dots$$

Run the "analogy" with quotients by Serre subcategories.
 We want to invert quasi-isom's, looks like it is enough
 to kill acyclic complexes.

$$f: X \rightarrow Y \rightsquigarrow 0 \rightarrow Y \rightarrow C(f) \rightarrow X\Sigma I \rightarrow 0$$

↑ if qis ⇒ must go to 0?

Also remark that $x^\circ \in C(\mathcal{A})$ is acyclic (all $H^i(x^\circ) = 0$)
 iff $0 \rightarrow x^\circ$ is a quasi-isomorphism.

Inverting qis's should kill all acyclics.

Problem In $C(A)$ qis's do not form a localization system (at least, not obvious).

2. Homotopy category

Given $x^\cdot, y^\cdot \in C(\mathcal{A})$, one can produce a complex

$$\text{Hom}^\cdot(x^\cdot, y^\cdot) \in C(\mathcal{Ab}) :$$

$\text{Hom}^n(x^\cdot, y^\cdot) = \prod \text{Hom}(x^i, y^{i+n}) \leftarrow$ collection of $x^i \rightarrow y^{i+n}$,
no relations / commutation with d's.

Given $\varphi \in \text{Hom}^n(x^\cdot, y^\cdot)$, put $d\varphi$ by the rule

$$\begin{array}{ccccccc} \dots & \rightarrow & x^i & \rightarrow & x^{i+1} & \rightarrow & \dots \\ & & \swarrow \varphi^i & & \swarrow \varphi^{i+1} & & \\ \dots & \rightarrow & y^{i+n} & \rightarrow & y^{i+n+1} & \rightarrow & \dots \end{array}$$

$$(d\varphi)^i = d\varphi^i - (-1)^i \varphi^{i+1} d$$

Exc Check that $\text{Hom}^*(X, Y)$ is indeed a complex.

Look at its lower terms.

$$Z^0(\text{Hom}^*(X, Y)) = \left\{ \varphi^i : x^i \rightarrow y^i \mid d\varphi^i - \varphi^{i+1}d = 0 \right\}$$

\uparrow
 $\text{Hom}_{C(d)}(X, Y)$

$$B^0(\text{Hom}^*(X, Y)) = \left\{ dh^i + h^{i+1}d : x^i \rightarrow y^i \mid h^i : x^i \rightarrow y^{i+1} \right\}$$

\nearrow

morphisms homotopic to zero!

We discussed that there is an equivalence relation on $\text{Hom}(X, Y)$ $f \sim g \Leftrightarrow f - g$ is homotopic to zero.

Morphisms homotopic to 0 form an ideal.

$$f \sim g \Rightarrow hf \sim hg, fw \sim gw$$

Observation $f \sim 0 \Rightarrow h^i(f) = 0$ for all $i \in \mathbb{Z}$.

Thus, $f \sim g$, then $h^i(f) = h^i(g)$ for all $i \in \mathbb{Z}$.

Define $K(\mathcal{A})$ as the category whose objects are complexes, morphisms - morphisms \sim .

$$\text{Hom}_{K(\mathcal{A})}(X, Y) = H^0(\text{Hom}^*(X, Y)).$$

Problem 5 Show that in $K(\mathcal{A})$ quasi-isomorphisms form a localization system.

$$\begin{array}{ccc}
 X & \xrightarrow{f} & Y \\
 \text{qis } s\downarrow & & \downarrow t \\
 Z & \dashrightarrow & W \\
 & & g
 \end{array}
 \quad tf \sim \overset{\text{homotopic}}{\underset{\text{not equal}}{\rightsquigarrow}} gs$$

Exc Check that $K(\mathcal{A})[Q_{\text{is}}^{-1}] = D(\mathcal{A}) = C(\mathcal{A})\Sigma Q_{\text{is}}^{-1}$!

Hint $k(\lambda)$ also satisfies some UP.

3. Resolutions

Recall that we have a functor $\mathcal{A} \rightarrow \mathcal{C}(\mathcal{A})$, exact, which sends $X \mapsto \dots \rightarrow 0 \rightarrow X \rightarrow 0 \rightarrow \dots$

Def A resolution of X is any complex $R \in C(\alpha)$ qis to X .

Def A projective resolution of X - resolution of the form

$$\rightarrow P^3 \rightarrow P^2 \rightarrow P^1 \rightarrow P^0 \rightarrow 0 \rightarrow \dots , \quad P^i - \text{proj.}$$

an injective resolution of X - resolution
of the form

$\dots \rightarrow 0 \rightarrow I^0 \rightarrow I^1 \rightarrow \dots$ I^i -injective.

Recall \mathcal{A} has enough projectives/injectives if
 $\forall A \in \mathcal{A} \exists P \rightarrowtail A / A \rightarrowtail I$ with
P-projective/I-injective.

If \mathcal{A} has enough projectives/injectives \Rightarrow
 \Rightarrow every object has a projective/injective
resolution.

Prop $\text{Mod-}A$ & $A\text{-Mod}$ have enough projectives.

Pf Free modules are projective, every module
is a quotient of a free module.

Then $A\text{-Mod}$ & $\text{Mod-}A$ have enough injectives.

Side remark: projectives are easier to deal with
(morally), but in many important situations
sheaves of \mathbb{A}^1 s e.g., there is not enough of
them.

Prop (Bayer's criterion)

$I \in A\text{-Mod}$ is injective \Leftrightarrow satisfies the UP
for all $I \hookrightarrow A$, I -ideal.

Pf \Rightarrow immediate

$$\begin{array}{ccc} \hookleftarrow & & \\ 0 \rightarrow X \rightarrow Y & & \\ f \downarrow & & \\ I & & \end{array}$$

consider the maximal submodule in Y containing X
for which the ext exists (Zorn).

WLOG we may assume that f does not extend
to any larger submodule.

If $Y \neq X \Rightarrow \exists g \in Y \setminus X$.

Consider $0 \rightarrow I \rightarrow X \oplus A \xrightarrow{X'} \circ$

$$\begin{array}{ccc} I & \xrightarrow{f} & X \\ \downarrow & \exists \quad \dashrightarrow & \downarrow \\ A & \xrightarrow{\quad g \quad} & I \end{array}$$

$X \oplus (g)$

This square is both cartesian and cocartesian. \square

Rmk A very similar argument (tough on ordinals) shows that in any Grothendieck category injective envelopes exist!

Application to $A\text{-Mod}$:

- 1) Check that \mathbb{Q}/\mathbb{Z} is injective in Ab !
- 2) $F: \mathcal{A} \rightarrow \mathcal{B}$, $G: \mathcal{B} \rightarrow \mathcal{A}$ are adjoint &
 F is exact, then G sends injectives to injectives.

3) Consider functors

$$F: A\text{-Mod} \rightarrow (\text{Mod-}A)^{\circ} \quad G: (\text{Mod-}A)^{\circ} \rightarrow A\text{-Mod}$$
$$M \mapsto \text{Hom}_{A\text{-Ab}}(M, \mathbb{Q}/\mathbb{Z})$$

Check that they are adjoint & exact.

Use that $(\text{Mod-}A)^{\circ}$ has enough injectives:
same as projectives in $\text{Mod-}A$.