

Triangulated and Derived Categories in Algebra and Geometry

Lecture 3

0. Adjoint Functors Revised

Recall (F, G) , $F \dashv G$ $F: A \leftrightarrows G: B$ if

$$\text{Hom}_B(F(-), -) \xrightarrow{\sim} \text{Hom}_A(-, G(-))$$

An isom of bifunctors $A^{\text{op}} \times B \rightarrow \text{Sets}$.

Examples • $\text{Ab} \subset \text{Grp} \hookleftarrow$ full embedding

Left adjoint:

$$\text{Hom}_{\text{Ab}}(?, A) \simeq \text{Hom}_{\text{Grp}}(G, A)$$

quotient by the commutator

• Grothendieck group construction

Consider $\text{AbMon} \supset \text{Ab}$, the full subcategory of abelian groups inside abelian monoids.

Left adjoint functor $M \rightsquigarrow K(A)$, where $K(A)$ is the quotient of

- 1) consider formal expressions / pairs $(m, n) \in M \times M$, think of (m, n) as a formal difference " $m - n$ ".

It inherits an operation : $(m, n) + (m', n') = (m+m', n+n')$

- 2) Want " $m - n$ " = " $(m+k) - (n+k)$ "

Consider the equivalence relation

$$(m, n) \sim (m', n') \iff \exists k \in M \text{ s.t. } m + n' + k = m' + n + k.$$

$$"m - n" = "m' - n" \iff m + n' = m' + n$$

since cancellation might not hold, need a less strong condition

$$\underline{\text{Ex}} \quad K(\mathbb{N}) \simeq \mathbb{Z}$$

Exc Check that K is a functor $\text{AbMon} \rightarrow \text{Ab}$, it's left adjoint to $\text{Ab} \hookrightarrow \text{AbMon}$.

Unit & Counit

$\text{Hom}_{\mathcal{B}}(F(X), Y) \simeq \text{Hom}_A(X, F(Y))$, plug in $Y = F(X)$

(consider the functor $A^{\text{op}} \times A \xrightarrow{\text{Id} \times F} A^{\text{op}} \times \mathcal{B}$).

$$\text{Hom}_{\mathcal{B}}(F(X), F(X)) \xrightarrow{\sim} \text{Hom}_A(X, (F \circ F)(X)) \quad \text{Id}_A \rightarrow F \circ F$$

Similarly, you obtain a nat transformation

$$\eta: F \circ G \rightarrow \text{Id}_{\mathcal{B}}$$

$$\begin{array}{ccccc} \text{Lm} & F & \xrightarrow{F \circ \epsilon} & F \circ G \circ F & \xrightarrow{\eta \circ F} F \\ & G & \xrightarrow{\epsilon \circ G} & G \circ F \circ G & \xrightarrow{G \circ \eta} G \end{array} \leftarrow \begin{array}{l} \text{this composition is the identity} \\ \text{morphism} \end{array}$$

Exc Prove it!

Hint There is a diagram (functoriality!)

$$\phi: F(x) \rightarrow Y$$

$$\text{Hom}_{\mathcal{F}}(F(x), F(x))$$

$$\begin{matrix} \downarrow \\ \phi \end{matrix}$$

$$\text{Hom}_{\mathcal{F}}(F(x), Y)$$

adjunction

$$\xrightarrow{\sim} \text{Hom}_A(x, (G \circ F)(x))$$

$$\downarrow G(\phi)$$

$$\xrightarrow{\sim} \text{Hom}_A(x, G(Y))$$

1. Monoidal categories

Def A monoidal category \mathcal{M} is a category \mathcal{M} , a bifunctor $\otimes: \mathcal{M} \times \mathcal{M} \rightarrow \mathcal{M}$, an object $I \in \mathcal{M}$ called the unit object, and three natural isomorphisms:

$$1) - \otimes (- \otimes -) \xrightarrow{\sim} (- \otimes -) \otimes -,$$

- 2) $I \otimes - \xrightarrow{\sim} Id$,
 3) $- \otimes I \xrightarrow{\sim} Id$.

Subject to some properties:

a) $A \otimes (B \otimes (C \otimes D)) \xrightarrow{\sim} (A \otimes B) \otimes (C \otimes D) \xrightarrow{\sim} ((A \otimes B) \otimes C) \otimes D$

$$\begin{array}{ccc} & \downarrow & \\ & & \hookdownarrow \\ A \otimes ((B \otimes C) \otimes D) & \xrightarrow{\sim} & (A \otimes (B \otimes C)) \otimes D \\ & \uparrow & \\ & & \uparrow s \end{array}$$

b) $A \otimes (I \otimes B) \xrightarrow{\sim} (A \otimes I) \otimes B$

$$\begin{array}{ccc} & \swarrow & \searrow \\ & A \otimes B & \end{array}$$

- Main examples
- Sets ; \otimes is the Cartesian product , \mathbb{Z}^{op} .
 - Cat ; \otimes is the product of cat's ; \mathcal{O}^{op} .

- $A\text{-Mod}$, \otimes - tensor product
In particular, $Ab \simeq \mathbb{Z}\text{-Mod}$ and $\text{Vect-}k$.

2. Enriched categories

Def Given a monoidal category M , an enriched category \mathcal{C} (enriched in M) consists of a set of objects $\text{Ob } \mathcal{C}$, $\forall X, Y \in \text{Ob } \mathcal{C}$ an object $\text{Hom}_{\mathcal{C}}(X, Y) \in M$,
 $\forall X$ a morphism $\text{id}_X : I \rightarrow \text{Hom}_{\mathcal{C}}(X, X)$, $\forall X, Y, Z \in \mathcal{C}$

$\circ : \text{Hom}_{\mathcal{C}}(Y, Z) \otimes \text{Hom}_{\mathcal{C}}(X, Y) \rightarrow \text{Hom}_{\mathcal{C}}(X, Z)$, such that...

Rmk Associativity of \circ is embedded in associativity of \otimes .

Rmk If you want some set of morphisms $X \rightarrow Y$,
 Look at arrows $I \rightarrow \text{Hom}_{\mathcal{C}}(X, Y)$ in M .

- Ex
- Categories are enriched categories over $(\text{Sets}, \times, \text{id})$.
 - 2-categories are enriched over $(\text{Cat}, \times, \circ)$.
 - Preadditive categories are enriched categories over (Ab, \otimes) .

3. Preadditive categories

Def A preadditive category is a category \mathcal{A} s.t. every $\text{Hom}_{\mathcal{A}}(X, Y)$ is an abelian group, the composition is bilinear

$$\text{Hom}_{\mathcal{A}}(Y, Z) \times \text{Hom}_{\mathcal{A}}(X, Y) \longrightarrow \text{Hom}_{\mathcal{A}}(X, Z)$$

$$\begin{array}{ccc} & \searrow & \nearrow \\ \text{Hom}_{\mathcal{A}}(Y, Z) & \otimes_{\mathbb{Z}} & \text{Hom}_{\mathcal{A}}(X, Y) \end{array} .$$

Lm Let \mathcal{A} be preadditive. The following are equivalent for an object $x \in \mathcal{A}$.

- (1) x is initial,
- (2) x is final,
- (3) $\text{id}_x = 0$ in $\mathcal{A}(x, x)$.

Def An object which is both initial and final is called a zero object, denoted by 0 .

Pf (of the lemma)

(1), (2) \Rightarrow (3) is trivial since in both cases $\text{Hom}(x, x)$ consists of a single element.

Assume $\text{id}_x = 0$ in $\text{Hom}(x, x)$

$\forall Y$ consider

$$\text{Hom}(x, x) \times \text{Hom}(Y, x) \rightarrow \text{Hom}(Y, x)$$

then $\forall f \in \text{Hom}(Y, x)$ bilinearity

$$f = \text{id}_x \circ f = 0 \circ f = 0$$

(χ is a bilinear form, then $\chi(0, v) = \chi(0+0, v) = 2\chi(0, v) \Rightarrow$
 $\Rightarrow \chi(0, v) = 0.$)

Thus, $f=0 \Rightarrow \text{Hom}(Y, X) = \{0\} \Rightarrow X$ is final.

Similarly, (3) \Rightarrow (1). □

Cor If a preadditive category has an initial / final object,
 then it has a zero object.

Lm Let \mathcal{A} be preadditive, $X, Y \in \mathcal{A}.$ If the product $X \times Y$ exists, so does the coproduct. Moreover, $X \sqcup Y \cong X \times Y.$

Proof Let $Z = X \times Y.$ Let's show that Z is a coproduct.

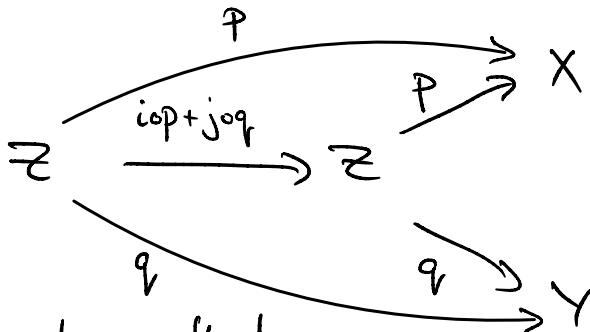
$$\begin{array}{ccc} X & \xrightarrow{i} & Z \\ & \searrow & \downarrow p \\ & Y & \xrightarrow{j} Y \end{array}$$

Let's define i and j by
 the the UP of Z

$$\begin{array}{ccccc} X & \xrightarrow{\text{id}} & Z & \xrightarrow{\quad} & X \\ & \xrightarrow{3!i} & Z & \xrightarrow{\quad} & Y \\ & & 0 & \xrightarrow{\quad} & Y \end{array}$$

These maps satisfy: $p \circ i = \text{id}_x$, $q \circ i = 0$,
 $q \circ j = \text{id}_y$, $p \circ j = 0$.

Claim $i \circ p + j \circ q = \text{id}_z$



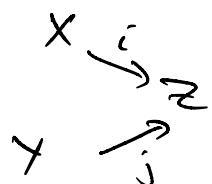
Enough to show that
the compositions with $p \neq q$ are $p \neq q$ again.

$$p \circ (i \circ p + j \circ q) = (p \circ i) \circ p + (p \circ j) \circ q = \text{id}_x \circ p + 0 \circ q = p.$$

id_x 0

Same for q .

Check that



satisfies the UP of $X \sqcup Y$.

□

Rmk 1) We could have done the same proof for coproduct \Rightarrow product!

2) What we actually proved is TFAE

(1) $X \times Y$ exists

(2) $X \sqcup Y$ exists

(3)

$$\begin{array}{ccc} X & \xleftarrow{i} & X \\ & \downarrow p & \downarrow q \\ Z & \xrightarrow{j} & Y \end{array}$$

$p \circ i = 1, p \circ j = 0$
 $q \circ i = 0, q \circ j = 1$
 $i \circ p + j \circ q = 1$

Def Let \mathcal{A}, \mathcal{B} be preadditive. A functor $F: \mathcal{A} \rightarrow \mathcal{B}$ is additive if the induced maps on hom's are homomorphisms of abelian groups: $\forall f, g: X \rightarrow Y$
 $F(f+g) = F(f) + F(g).$

Cor An additive functor preserves products, coproducts and zero objects.

Pf . \mathbb{Z} is zero $\Leftrightarrow \text{id}_{\mathbb{Z}} = 0$. If $\text{id}_{\mathbb{Z}} = 0$, then
 $F(\text{id}_{\mathbb{Z}}) = \text{id}_{F(\mathbb{Z})} = F(0) = 0$,

- for products use this diagram:

$$\begin{array}{ccc} X & \xrightarrow{\quad} & Z & \xrightarrow{\quad} & X \\ Y & \xrightarrow{\quad} & Z & \xrightarrow{\quad} & Y \end{array}$$

□

Def An additive category is a preadditive category in which all finite products exist (including the empty = final obj).

Notation In an additive category $X \times Y$ is denoted by $X \oplus Y$ and is called the direct sum.

Main examples

- Ab - abelian groups
- A - ring, $A\text{-Mod}$ and $\text{Mod-}A$
 - $A\text{-Mod}$ ↑
left modules
 - $\text{Mod-}A$ ↑
right modules
- look at free or f.g. free A -modules

- Let \mathcal{A} be an additive category, then \mathcal{A}^{op} is additive!
- Given a category \mathcal{I} & an additive category \mathcal{A} ,
the category of functors $\text{Fun}(\mathcal{I}, \mathcal{A})$ is additive!
Abelian groups of natural transformations (pointwise addition).
Direct sums are object-wise: $F, G : \mathcal{I} \rightarrow \mathcal{A}$,
 $(F \oplus G)(x) = F(x) \oplus G(x)$.

Exc Work out the details.

- X - topological space $\rightsquigarrow \text{Op}(X)$
objects are $U \subset X$ - open, $U \subseteq V \Leftrightarrow U \subseteq V$
- $\text{PSh}_{\mathcal{A}}(X) = \text{Fun}(\text{Op}(X)^{\text{op}}, \mathcal{A})$
presheaves on X with values in \mathcal{A} .

Def Let \mathcal{A} be additive, $X \in \mathcal{A}$. The diagonal
 $\Delta : X \rightarrow X \oplus X$ is induced by $\begin{array}{ccc} X & \xrightarrow{s} & X \\ & \downarrow & \\ & \xrightarrow{t} & X \end{array}$.

$\nabla: X \oplus X \rightarrow X$ is induced by $\begin{array}{ccc} X & \xrightarrow{f} & X \\ \downarrow g & & \downarrow g \\ X & \xrightarrow{g} & X \end{array}$.

codiagonal

Exc Show that the following commutes $\forall f, g: X \rightarrow Y$

$$\begin{array}{ccccccc} X & \xrightarrow{\Delta_X} & X \oplus X & \xrightarrow{f+g} & Y \oplus Y & \xrightarrow{\nabla_Y} & Y \\ & \searrow & & \nearrow & & & \\ & & & f+g & & & \end{array}$$

Exc Using show that a functor $F: \mathcal{A} \rightarrow \mathcal{B}$ is additive \Leftrightarrow it preserves products:

$$F(X \oplus Y) \cong F(X) \oplus F(Y).$$

Exc Redefine additive categories as categories with zero objects and direct sums ($\mathbb{N} = \mathbb{U}$). Reconstruct the abelian group structure on hom spaces, show bilinearity.

E.g.

$$\begin{array}{ccc} X & \xrightarrow{0} & Y \\ \downarrow 0 & \nearrow & \end{array}$$

Def Let \mathcal{A} be additive, $f: X \rightarrow Y$.

The kernel (if exists) of f is the equalizer $E_f(0, f)$.

The cokernel (if exists) of f is the coequalizer $C_{eq}(0, f)$.

The coimage is the cokernel of the kernel.

The image is the kernel of the cokernel.

Exe Show that $\text{ker}(f)$ is what you are more used to:

$$\begin{array}{ccc} \text{Ker}(f) & \xrightarrow{\epsilon} & X \xrightarrow{f} Y \\ \exists ! h \uparrow & \nearrow g & \\ Z & & \end{array} \quad \begin{array}{l} f \circ \epsilon = 0 \\ \forall g \text{ s.t. } f \circ g = 0 \\ \exists ! h \text{ s.t. } \epsilon \circ h = g. \end{array}$$

Next time Abelian categories.

Before Try and construct $\text{Coim}(f) \rightarrow \text{Im}(f)$
if everything exists.