



# SQL介绍

## Introduction to SQL

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- IBM **Sequel language** is developed as part of **System R project** at the IBM San Jose Research Laboratory in the early 1970s
- Renamed **Structured Query Language (SQL)**
- ANSI (美国国家标准学会) and ISO (国际标准化组织) standard SQL
  - SQL-86
  - SQL-89
  - SQL-92
  - SQL-99
  - SQL: 2003, 2006, 2008, 2011, 2016
- Commercial systems offer most SQL-92 features, plus varying feature sets from later standards and special proprietary features
  - Not all examples here may work on particular DBMS systems

- **数据定义语言 (Data definition language, DDL)**
  - Relation schemas
  - Integrity constraints
  - View
  - Authorization
- **数据操纵语言 (Data manipulation language, DML)**
  - Queries
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- **定义关系及其相关信息，包括：**
  - 关系的模式
  - 属性的域(domain)
  - 完整性约束
  - 索引结构
  - 安全性和权限信息
  - 物理存储结构

- **char( $n$ )**
  - fixed length character string with user-specified length  $n$
- **varchar( $n$ )**
  - variable length character strings with user-specified maximum length  $n$
- **int**
  - Integer, a finite subset of the integers that is machine-dependent
- **smallint**
  - small integer, a machine-dependent subset of the integer domain type
- **numeric( $p$ ,  $d$ )**
  - fixed point number (定点数), with user-specified precision of  $p$  digits, and with  $d$  digits to the right of decimal point
  - E.g., numeric(3,1) allows 44.5 to be stored exactly, but neither 444.5 nor 0.32 can be stored exactly



- **real, double precision**
  - floating point and double-precision floating point numbers, with machine-dependent precision
- **float(*n* )**
  - floating point number with user-specified precision of at least *n* digits
- **null value**
  - allowed in all domain types. Declaring an attribute to be not null prohibits null values for that attribute
- In SQL-92, we can create user-defined domain types
  - **create domain** person\_name char(20) **not null**

- Relation is defined using the *create table* command:

*create table*  $r(A_1 D_1, A_2 D_2, \dots, A_n D_n,$   
 $(integrity\_constraint_1), \dots, (integrity\_constraint_k))$

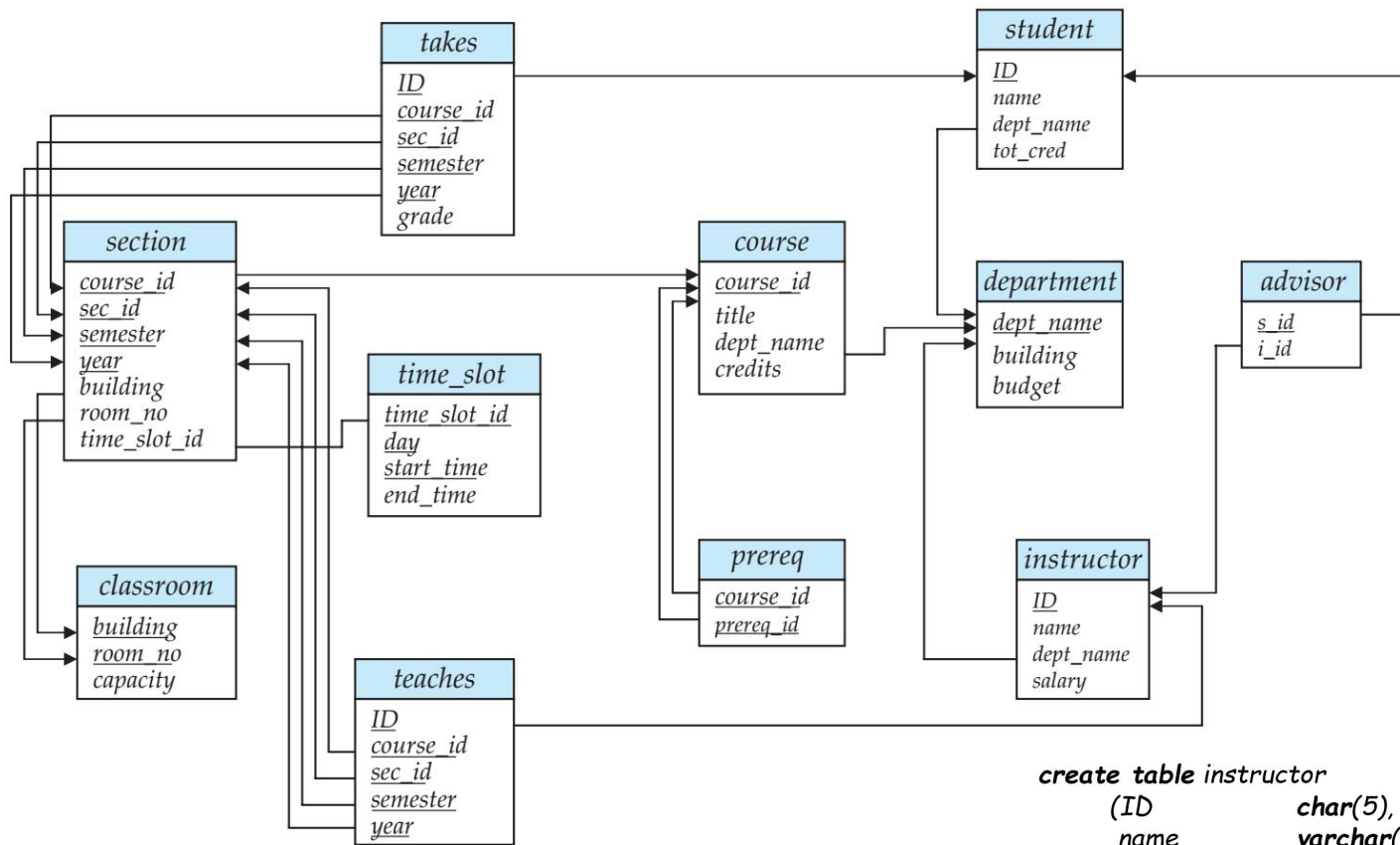
- $r$ : the relation name
  - $A_i$ : an attribute name in the schema of relation  $r$
  - $D_i$ : the data type of attribute  $A_i$
- 例:

```
create table department(  
    dept_name    varchar(20),  
    building     varchar(15),  
    budget       numeric(12, 2),  
    primary key (dept_name));
```

- not null
- primary key ( $A_1, \dots, A_n$ )
- foreign key ( $A_{k1}, A_{k2}, \dots, A_{kn}$ ) references  $s$
- check ( $P$ ), where  $P$  is a predicate

```
create table instructor(  
    ID          char(5),  
    name        varchar(20) not null,  
    dept_name    varchar(20),  
    salary       numeric(8, 2),  
    primary key (ID),  
    foreign key (dept_name) references department,  
    check (salary >= 0));
```

**注意:** **Primary key** declaration on an attribute automatically ensures **not null** and **unique** in SQL-92 onwards, needs to be explicitly stated in SQL-89



思考：用SQL定义student和advisor?

```

create table instructor
(ID          char(5),
 name       varchar(20) not null,
 dept_name  varchar(20),
 salary     numeric(8, 2),
 primary key (ID),
 foreign key (dept_name) references department,
 check (salary >= 0));
  
```

- **drop table**: deletes all information about the dropped relation from the database
- **alter table**: add attributes to a relation or drop attributes from a relation
  - alter table r add A D*
    - All tuples in the relation are assigned null as the value for the new attribute
  - alter table r drop A*
    - Dropping of attributes is not supported by some databases

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- 典型的SQL查询结构:

```
select  $A_1, A_2, \dots, A_n$   
from  $r_1, r_2, \dots, r_m$   
where  $P$ 
```

- The result of an SQL query is a relation and each query is equivalent to the relational algebra expression:

$$\Pi_{A_1, A_2, \dots, A_n}(\sigma_P(r_1 \times r_2 \times \dots \times r_m))$$

```
for each tuple  $t_1$  in relation  $r_1$   
  for each tuple  $t_2$  in relation  $r_2$   
    ...  
    for each tuple  $t_m$  in relation  $r_m$   
      Concatenate  $t_1, t_2, \dots, t_m$  into a single tuple  $t$   
      Add  $t$  into the result relation
```

- The select clause lists the attributes desired in the result of a query
  - correspond to the **projection (投影)** operation of the RA (relational algebra)
  - E.g., find the names of all departments in the instructor relation

```
select dept_name  
from instructor
```

- 对应的关系代数表达式:

$$\Pi_{dept\_name}(instructor)$$

- **注意:** SQL is case **insensitive**



- SQL allows **duplicates** in relations. To eliminate the duplicates, insert the keyword **distinct** after select
  - Find the names of all departments in the instructor relation, and remove duplicates

```
select distinct dept_name  
from instructor
```

- The keyword **all** specifies that duplicates should not be removed

```
select all dept_name  
from instructor
```

- An asterisk in the select clause denotes “all attributes”  
***select*** \*  
***from*** *instructor*
- The select clause can contain **arithmetic expressions** involving the operation +, −, \* , and /, and operating on constants or attributes of tuples
- E.g.,  
***select*** *ID, dept\_name, salary \* 1.1*  
***from*** *instructor*

- The **from clause** lists the relations involved in the query
  - corresponds to the **Cartesian product** operation of the RA

- 例: find the Cartesian product  $borrower \times loan$

```
select *  
from borrower, loan
```

- 例: find the name, loan number and loan amount of all the customers that have a loan at the Jiading branch

```
select customer_name, borrower.loan_number, amount  
from borrower, loan  
where borrower.loan_number=loan.loan_number and  
branch_name='Jiading'
```

- The where clause specifies conditions that the result must satisfy
  - correspond to the selection predicate of the RA
  - 例: *find the names of those instructors who are with the Department of Computer Science and have the salary larger than 70000*

```
select name  
from instructor  
where dept_name = 'Computer Science' and salary > 70000
```

Comparison results can be combined using the logical connectives **and**, **or**, and **not**

- SQL includes a **between** comparison operator
  - 例: *find the loan numbers of those loans with loan amount between \$90,000 and \$100,000*

```
select loan_number  
from loan  
where amount between 90000 and 100000
```

## ► 自然连接(Natural Join)



```
select  $A_1, A_2, \dots, A_n$   
from  $r_1$  natural join  $r_2$  natural join ... natural join  $r_m$   
where  $P$ ;
```

```
select name, course_id  
from instructor natural join teaches;
```

```
select name, title  
from instructor natural join teaches, course  
where teaches.course_id=course.course_id;
```

## ► join ... using (...)



- natural join of instructor and teaches
  - (ID, name, dept\_name, salary, course\_id, sec\_id)
- course
  - (course\_id, title, dept\_name, credits)

结果不等 {

```
select name, title
from instructor natural join teaches, course
where teaches.course_id= course.course_id;
```

结果相等

```
select name, title
from instructor natural join teaches natural join course;

select name, title
from (instructor natural join teaches) join course using (course_id);
```

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- The SQL allows renaming relations and attributes using the **as** clause:  
*old\_name as new\_name*
- Find the *name*, *loan\_number* and *loan\_amount* of all customers, and rename the column name *loan\_number* as *loan\_id*:

```
select customer_name, borrower.loan_number as loan_id, amount  
from borrower, loan  
where borrower.loan_number = loan.loan_number
```

- 找出每位老师的姓名和所授所有课程的ID

```
select T.name, S.course_id  
from instructor as T, teaches as S  
where T.ID = S.ID
```

- 找出满足下列条件的所有教师的姓名，他们的工资至少比Biology系的某一位教师的工资高

```
select distinct T.name  
from instructor as T, instructor as S  
where T.salary > S.salary and S.dept_name = 'Biology'
```

- SQL includes a string-matching operator for comparisons on character strings
  - percent (%): matches any substring
  - underscore (\_): matches any character
- **like/not like**: 找出所在建筑名称中包含子串“Watson”的所有系名

```
select dept_name  
from department  
where building like '%Watson%'
```
- 要匹配的字符中有百分号的情况，需要转义

```
like 'Watson\%'
```
- “\*” denote “all attributes” : *select instructor.\* from instructor*
- SQL supports a variety of string operations such as
  - concatenation (串联) (using “||”)
  - converting from upper case to lower case (and vice versa)
  - finding string length, extracting substrings, etc.

- List in alphabetic order the names of all customers having a loan in Jiading branch  

```
select distinct customer_name  
from borrower, loan  
where borrower.loan_number = loan.loan_number and  
        branch_name = 'Jiading'  
order by customer_name
```
- Specify **desc** for descending order or **asc** (default) for ascending order, for each attribute  

```
select *  
from loan  
order by amount desc, loan-number asc
```

- SQL includes a **between/not between** comparison operator
  - Example: find the names of all instructors with salary between \$90,000 and \$100,000

```
select name  
from instructor  
where salary between 90000 and 100000
```

- Tuple comparison

等价

```
select name, course_id  
from instructor, teaches  
where (instructor.ID, dept_name) = (teaches.ID, 'Biology');
```

```
select name, course_id  
from instructor, teaches  
where instructor.ID=teaches.ID and dept_name='Biology'
```

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- The set operations **union**, **intersect**, and **except** operate on relations and correspond to the relational algebra operations  $\cup$ ,  $\cap$ ,  $-$
- Each set operation **automatically eliminates duplicates**. To retain all duplicates use the corresponding multiset versions **union all**, **intersect all** and **except all**
  - Suppose a tuple occurs  $m$  times in  $r$  and  $n$  times in  $s$ , then, it occurs:
    - $m + n$  times in  $r$  union all  $s$
    - $\min(m, n)$  times in  $r$  intersect all  $s$
    - $\max(0, m - n)$  times in  $r$  except all  $s$

- Find all customers who have a loan, an account, or both:  
    (select customer\_name from depositor)  
    **union [all]**  
    (select customer\_name from borrower)
- Find all customers who have both a loan and an account.  
    (select customer\_name from depositor)  
    **intersect [all]**  
    (select customer\_name from borrower)
- Find all customers who have an account but no loan.  
    (select customer\_name from depositor)  
    **except [all]**  
    (select customer\_name from borrower)



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- It is possible for tuples to have a **null** value, signifying an **unknown** value or a value that does not exist
- The predicate **is null** can be used to check for null values

```
select loan_number  
from loan  
where amount is null
```

- The result of any arithmetic expression involving null is null
  - E.g.,  $5 + \text{null}$  returns null



- Calculate the sum of all loan amounts  
***select sum (amount)***  
***from loan***
  - Above statement ignores null amounts
  - Result is null if there is no non-null amount
- All aggregate operations **except count(\*)** ignore tuples with null values on the aggregated attributes

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## ► 聚集函数(Aggregate Functions)



- These functions receive as input a set of values, and return a value
  - **avg**: average value
  - **min**: minimum value
  - **max**: maximum value
  - **sum**: sum of values
  - **count**: number of values

- 找出Computer Science系教师的平均工资  
***select avg (salary)***  
***from instructor***  
***where dept\_name= 'Computer Science'***
- 找出在2018年春季学期授课的教师总数  
***select count (distinct ID)***  
***from teaches***  
***where semester = 'Spring' and year=2018***

- Find the number of depositors for each branch  
***select branch\_name, count (distinct customer\_name)***  
***from depositor, account***  
***where depositor.account\_number = account.account\_number***  
***group by branch\_name***
- 注意: Attributes in select clause outside of aggregate functions must appear in group by list  
*/\*erroneous query\*/*  
***select dept\_name, ID, avg(salary)***  
***from instructor***  
***group by dept\_name***

## ► 聚集函数：Having子句



- At times, it is useful to state a condition that applies to groups rather than to tuples
- 例：find the names of all branches where the average account balance is more than \$1,200

```
select branch_name, avg (balance)
from account
group by branch_name
having avg (balance) > 1200
```

- 注意：
  - predicates in the having clause are applied **after** the information of groups
  - predicates in the where clause are applied **before** forming groups

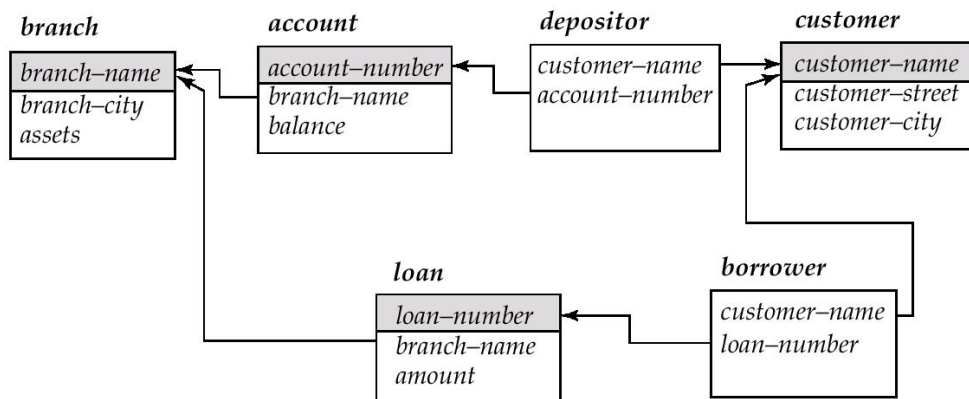


## ► 聚集函数：Having子句(续)



- 例：find the average balance for each customer who lives in Shanghai and has at least three accounts

```
select depositor.customer_name, avg (balance)
from depositor, account, customer
where depositor.account_number=account.account_number
      and depositor.customer_name=customer.customer_name
      and customer_city='Shanghai'
group by depositor.customer_name
having count(distinct depositor.account_number) >=3
```



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## ► 嵌套子查询(Nested Subqueries)



- A subquery is a **select-from-where** expression that is nested within another query in the where clause or from clause
- A common use of subqueries is to perform
  - tests for **set membership** (测试集合成员资格)
  - make **set comparisons** (集合比较)
  - determine **set cardinality** (确定集合基数)

- Find all customers who have **both** an account and a loan at the bank  
***select distinct customer\_name***  
***from borrower***  
***where customer\_name in (select customer\_name***  
***from depositor)***
  - Find all customers who have a loan but do not have an account at the bank  
***select distinct customer\_name***  
***from borrower***  
***where customer\_name not in (select customer\_name***  
***from depositor)***
- 
- select distinct name***  
***from instructor***  
***where name not in ('Mozart', 'Einstein');***

- Find all customers who have both an account and a loan at the Jiading branch

```
select distinct customer_name
from borrower, loan
where borrower.loan_number=loan.loan_number and
      branch_name="Jiading" and
      (branch_name, customer_name) in
      (select branch_name, customer_name
       from depositor, account
       where depositor.account_number =account.account_number)
```

- Find all branches that have greater assets than some branch located in Brooklyn

```
select distinct T.branch_name  
from branch as T, branch as S  
where T.assets > S.assets and S.branch_city = 'Brooklyn'
```

- Same query using **>some** clause

```
select branch_name  
from branch  
where assets > some  
    (select assets  
     from branch  
     where branch_city = 'Brooklyn')
```

## ► Some子句



- $E < \text{comp} > \text{some } r \Leftrightarrow \exists t \in r \text{ such that } (E < \text{comp} > t)$ , where  $\text{comp}$  can be:  $<, \leq, >, \geq, =, \neq$

$(5 < \text{some } \begin{array}{|c|} \hline 0 \\ \hline 5 \\ \hline 6 \\ \hline \end{array}) = \text{true}$  (read: 5 < some tuple in the relation)

$(5 < \text{some } \begin{array}{|c|} \hline 0 \\ \hline 5 \\ \hline \end{array}) = \text{false}$

$(5 = \text{some } \begin{array}{|c|} \hline 0 \\ \hline 5 \\ \hline \end{array}) = \text{true}$

$(5 \neq \text{some } \begin{array}{|c|} \hline 0 \\ \hline 5 \\ \hline \end{array}) = \text{true}$  (since  $0 \neq 5$ )

(= some)  $\equiv$  in  
However,  
( $\neq$  some)  $\not\equiv$  not in, why?

- $E < \text{comp} > \text{all } r \Leftrightarrow \forall t \in r (E < \text{comp} > t)$

$$(5 < \text{all } \begin{array}{|c|} \hline 0 \\ \hline 5 \\ \hline 6 \\ \hline \end{array}) = \text{false}$$

$$(5 < \text{all } \begin{array}{|c|} \hline 6 \\ \hline 10 \\ \hline \end{array}) = \text{true}$$

$$(5 = \text{all } \begin{array}{|c|} \hline 4 \\ \hline 5 \\ \hline \end{array}) = \text{false}$$

$$(5 \neq \text{all } \begin{array}{|c|} \hline 4 \\ \hline 6 \\ \hline \end{array}) = \text{true (since } 5 \neq 4 \text{ and } 5 \neq 6)$$

( $\neq$  all)  $\equiv$  not in  
However,  
( $=$  all)  $\neq$  in, why?



- Find the names of all branches that have greater assets than all the branches located in Brooklyn.

```
select branch_name  
from branch  
where assets > all  
    (select assets  
     from branch  
     where branch_city = 'Brooklyn')
```

```
select dept_name  
from instructor  
group by dept_name  
having avg (salary) >= all (select avg (salary)  
                             from instructor  
                             group by dept_name);
```

思考：左边SQL语句的查询内容是什么？

- The **exists** construct returns the value true if the argument subquery is nonempty
  - $\text{exists } r \Leftrightarrow r \neq \emptyset$
  - $\text{not exists } r \Leftrightarrow r = \emptyset$
- E.g., find all the customers who have both account(s) and loan(s) at the bank

```
select distinct customer_name  
from borrower  
where exists (  
    select *  
    from depositor  
    where depositor.customer_name=borrower.customer_name)
```



- We can write “relation A contains relation B” as “not exists (B except A).”
- E.g., find those students who have taken all the courses offered by the Biology Department

```
select distinct S.ID, S.name  
from student as S  
where not exists ((select course_id  
                   from course  
                   where dept_name = 'Biology')  
except  
(select T.course_id  
 from takes as T  
 where S.ID = T.ID));
```



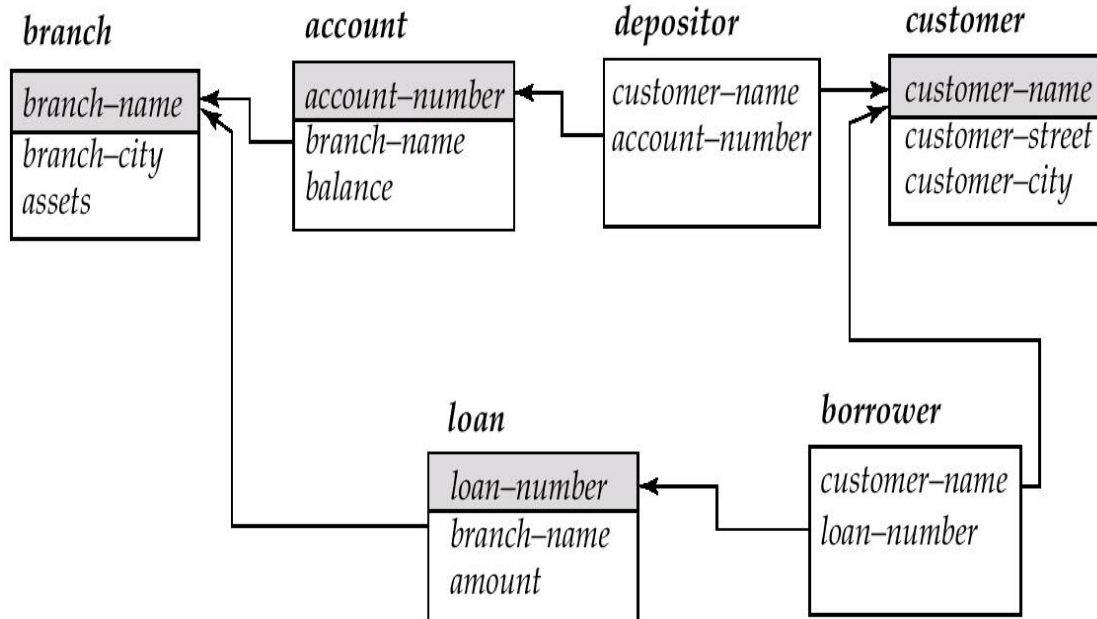
- Find all the customers who have accounts at all branches located in Brooklyn

```
select distinct S.customer_name
from depositor as S
where not exists (
    (select branch_name    /* all branches in Brooklyn */
     from branch
     where branch_city = 'Brooklyn')
    except
    (select R.branch_name /* finds all the branches at which customer
S.customer_name has an account */
     from depositor as T, account as R
     where T.account_number = R.account_number and
           T.customer_name = S.customer_name))
```

- Note:**  $\text{not exists } (X - Y) \Leftrightarrow X - Y = \emptyset \Leftrightarrow X \subseteq Y$

- The **unique** construct tests whether a subquery has any duplicate tuples in its result
- E.g., find all the customers who have at most one account at the Jiading branch

```
select T.customer_name
from depositor as T    /* Better to use from customer. why? */
where unique (
    select R.customer_name
    from account, depositor as R
    where T.customer_name = R.customer_name and
        R.account_number = account.account_number and
        account.branch_name = 'Jiading')
```



- 找出所有在2017年最多开设一次的课程

```
select T.course_id  
from course as T  
where unique (select R.course_id  
              from section as R  
              where T.course_id= R.course_id and  
                    R.year = 2017);
```

等价

```
select T.course_id  
from course as T  
where 1 >= (select count(R.course_id)  
            from section as R  
            where T.course_id= R.course_id and  
                  R.year = 2017);
```

- Find all customers who have at least two accounts at the Jiading branch.

```
select distinct T.customer_name
from depositor T
where not unique(
    select R.customer_name
    from account, depositor as R
    where T.customer_name = R.customer_name and
        R.account-number = account.account_number and
        account.branch_name = 'Jiading')
```

```
select T.course_id
from course as T
where not unique (select R.course_id
    from section as R
    where T.course_id= R.course_id and
        R.year = 2017);
```

找出所有在2017年至少开设  
两次的课程

**思考：找出所有在2017年至  
少开设三次的课程？**



- In some cases, it is not desirable for all users to see the entire logical model, i.e., all the actual relations stored in the database
- Consider a person who needs to know a customer's name, loan number and branch name, but has no need to see the loan amount

```
select customer_name, borrower.loan_number, branch_name  
from borrower, loan  
where borrower.loan_number = loan.loan_number
```

- A **view** provides a mechanism to hide certain data from the view of certain users. It's a **virtual relation**

- A view is defined using the create view statement

***create view v as < query expression >***

where  $v$  is the view name, and  $\langle \text{query expression} \rangle$  is any legal SQL expression

- Once a view is defined, the view name can be used to refer to the virtual relation, and **the query expression is stored in the database**. The expression is substituted into queries when the view is used.

## ► 视图：举例



- A view consisting of branches and their customers

**create view** *all\_customer* as

**(select** *branch\_name, customer\_name*

**from** *depositor, account*

**where** *depositor.account\_number = account.account\_number)*

**union**

**(select** *branch\_name, customer\_name*

**from** *borrower, loan*

**where** *borrower.loan\_number = loan.loan\_number)*

- Find all customers of the Jiading branch

**select** *customer\_name*

**from** *all\_customer*

**where** *branch\_name = 'Jiading'*

- **Derived Relations**

- Find the average account balance of those branches which have the average account balance greater than \$1200

```
select branch_name, avg (balance)  
from account  
group by branch_name  
having avg (balance) > 1200
```

```
select branch_name, avg_balance  
from (select branch_name, avg (balance)  
from account  
group by branch_name)  
as result (branch_name, avg_balance)  
where avg_balance > 1200
```

- **Note:** we do not need to use the having clause, since we compute the temporary (view) relation result in the from clause, and the attributes of result can be used directly in the where clause

- Find the maximum total balance across all branches

```
select max(tot_balance)
from (select branch_name, sum (balance)
       from account
       group by branch_name)
as branch_total (branch_name, tot_balance)
```

- With clause allows views to be defined **locally** to a query, rather than globally. Analogous to procedures in a programming language
- Find all accounts with the maximum balance

```
with max_balance(value) as  
    select max(balance)  
    from account
```

```
select account_number  
from account, max_balance  
where account.balance = max_balance.value
```



## 使用with子句的复杂查询



- Find all branches where the total account deposit is greater than the average of the total account deposits at all branches

```
with branch_total (branch_name, value) as  
select branch_name, sum (balance)  
from account  
group by branch_name
```

```
with branch_total_avg (value) as  
select avg (value)  
from branch_total
```

```
select branch_name  
from branch_total, branch_total_avg  
where branch_total.value >= branch_total_avg.value
```

## ► 标量子查询(Scalar Subquery)



- Scalar subquery (标量子查询) is used where **a single value** is expected
- List all departments along with the number of instructors in each department

```
select dept_name,  
      (select count(*)  
       from instructor  
       where department.dept_name = instructor.dept_name)  
      as num_instructors  
from department;
```

- Runtime error if subquery returns more than one tuple



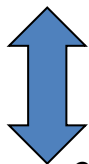
- SQL概览
- SQL数据定义
- SQL查询的基本结构
- 附加基本运算
- 集合运算
- 空值
- 聚集函数
- 嵌套子查询
- **数据库的修改**

## ► 数据库的修改: Deletion



- Delete all accounts at every branch located in Needham city

```
delete from account  
where branch_name in (select branch_name  
                        from branch  
                        where branch_city = 'Needham')
```



```
delete from depositor  
where account_number in  
      (select account_number  
       from branch, account  
       where branch_city = 'Needham'  
       and branch.branch_name = account.branch_name)
```

**思考：删除操作是否存在问题？**

## ► 数据库的修改: Deletion(续)



- Delete the records of all accounts with balances below the average at the bank  
*delete from account*  
*where balance < (select avg(balance)*  
*from account)*
- **Note:** as we delete tuples from account, the average balance changes
- Solution used in SQL:
  - First, compute avg balance and find all tuples to delete
  - Next, delete all tuples found above (without recomputing avg or retesting the tuples)

- Add a new tuple to account

***insert into account values*** ('A-9732', 'Perryridge', 1200)

or equivalently

***insert into account (branch\_name, balance, account\_number)***  
***values*** ('Perryridge', 1200, 'A-9732')

- Add a new tuple to account with balance set to null

***insert into account***  
***values*** ('A-777', 'Perryridge', null)

## ► 数据库的修改: Insertion(续)



- Provide as a gift for all loan customers of the Jiading branch, i.e., a \$200 saving account. Let the loan number serve as the account number for the new saving account  
*insert into account*  
    *select loan\_number, branch\_name, 200*  
    *from loan*  
    *where branch\_name = 'Jiading'*  
*insert into depositor*  
    *select customer\_name, loan\_number*  
    *from loan, borrower*  
    *where loan.loan\_number = borrower.loan\_number*  
        *and branch\_name = 'Jiading'*
- The select from where statement is **fully evaluated** before its results are inserted into the relation. Otherwise, queries like *insert into table1 select \* from table2* would cause problems

## ► 数据库的修改: Updates



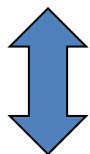
- Increase all accounts with balances over \$10,000 by 6%, and all other accounts receive an increase of 5%.

- Write two update statements:

***update account***

***set balance = balance \* 1.05***

***where balance ≤ 10000***



***update account***

***set balance = balance \* 1.06***

***where balance > 10000***

**思考: 插入操作是否存在问题?**

- The **order** is important
- Can be done better using the case statement (next slide)

- Same query: increase all accounts with balances over \$10,000 by 6%, and all other accounts receive 5%.

```
update account  
  set balance =  
  case  
    when balance <= 10000 then balance *1.05  
    else balance * 1.06  
  end
```

```
case  
  when  $pred_1$  then  $result_1$   
  when  $pred_2$  then  $result_2$   
  ...  
  when  $pred_n$  then  $result_n$   
  else  $result_0$   
end
```

- **SQL**
  - DDL + DML
- **DDL**
  - 数据库模式、完整性约束等
- **SQL查询**
  - select子句、from子句、where子句
  - natural join, join ... using (...)
- **SQL附加运算**
  - 更名运算、字符串运算、排序order by
- **集合运算**
  - union、intersect、except
- **空值**
- **聚集函数**
  - avg、min、max、sum、count
  - group by、having
- **嵌套子查询**
  - 集合包含：in、not in
  - 集合比较：some子句、all子句、exists、unique
  - 视图、导出关系、with子句、标量子查询
- **数据库的修改**
  - Deletion、Insertion、Updates



- **Exercises**
  - 3.8, 3.9 (选择其中一个)
  - 3.15, 3.16, 3.17, 3.21 (选择其中两个)
- **Submission**
  - Canvas上提交, 上传单个PDF文件
  - Deadline: 待定