Read-only rootfs

Theory and practice

Chris Simmonds

Embedded Linux Conference Europe 2016



License



These slides are available under a Creative Commons Attribution-ShareAlike 3.0 license. You can read the full text of the license here

http://creativecommons.org/licenses/by-sa/3.0/legalcode You are free to

- copy, distribute, display, and perform the work
- make derivative works
- · make commercial use of the work

Under the following conditions

- Attribution: you must give the original author credit
- Share Alike: if you alter, transform, or build upon this work, you may distribute the resulting work only
 under a license identical to this one (i.e. include this page exactly as it is)
- For any reuse or distribution, you must make clear to others the license terms of this work

About Chris Simmonds



- Consultant and trainer
- Author of Mastering Embedded Linux Programming
- Working with embedded Linux since 1999
- Android since 2009
- Speaker at many conferences and workshops

"Looking after the Inner Penguin" blog at http://2net.co.uk/



https://uk.linkedin.com/in/chrisdsimmonds/



https://google.com/+chrissimmonds

Overview

- Why you need a read-only rootfs
- Where it all goes wrong
- Putting it right

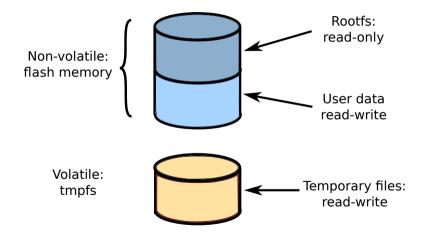
Why you need a read-only rootfs

- Reduce wear on flash memory
- Eliminate system file corruption
- Avoid accidents
- Enable rootfs image to be updated (manually or OTA)
- Make reset to factory defaults much easier

Where it all goes wrong

- You can't just mount the rootfs ro
- Some files have to be created or updated at runtime
- Examples:
 - Passwords
 - Random seed
 - SSH keys
 - Network configuration
 - wpa_supplicant parameters

Categories of storage



Becoming stateless

- The state (stuff that changes) is either
 - Non-volatile: state that needs to be preserved
 - Volatile: state that is only needed for the session
- To create a stateless rootfs:
 - Move non-volatile state into an area of permanent storage reserved for that purpose
 - Move Volatile state to temporary storage (tmpfs RAM disk)

The ideal of a stateless rootfs

- No state stored in rootfs
- Each component can revert to sensible default configuration if there is no local configuration

Factory reset

Just wipe the non-volatile state

System update

- Simply write a new rootfs image
- Two common mechanisms:
 - A/B rootfs partitions: A is live while B is being updated, then swap
 - Or, normal/recovery rootfs: boot into the recovery rootfs to update the normal rootfs
- This is the subject of several other presentations this week
 - One of which is mine: Software Update for IoT: The Current State of Play, Wednesday 14:00

Identifying state

- Which files are being written to?
- And by whom?
- Tools
 - diskstats
 - block_dump

diskstats

- /proc/diskstats contains information about disk reads and writes
- Useful stuff is field 1: reads completed and field 5: writes completed
 - Format is documented in Documentation/iostats.txt
 - vmstat -d prints the same information but better
- Example

Number of reads = 2640: Number of writes = 127

block_dump

- But, what is the cause of those writes? Which files?
- Turn on block system kernel logging using block_dump
 - See Documentation/laptops/laptop-mode.txt in kernel source
- Example (rootfs is on /dev/vda):

```
# echo 1 > /proc/sys/vm/block_dump
# echo hello > world.txt
# dmesg
sh(234): dirtied inode 701 (.ash_history) on vda
sh(234): dirtied inode 701 (.ash_history) on vda
sh(234): dirtied inode 701 (.ash_history) on vda
sh(234): dirtied inode 694 (world.txt) on vda
sh(234): dirtied inode 694 (world.txt) on vda
sh(234): dirtied inode 694 (world.txt) on vda
jbd2/vda-8(78): WRITE block 802 on vda (2 sectors)
jbd2/vda-8(78): WRITE block 804 on vda (2 sectors)
```

Logging block I/O from bootup

- Set block_dump in early boot script (may need to bump the size of the kernel log buffer to capture everything)
- · Then filter out the junk:

```
# dmesg | grep dirtied | grep "on vda" | sort
chown(150): dirtied inode 548 (passwd) on vda
dd(298): dirtied inode 716 (random-seed) on vda
dd(298): dirtied inode 716 (random-seed) on vda
dropbearkey(325): dirtied inode 717 (dropbear_rsa_host_key) on vda
dropbearkey(325): dirtied inode 717 (dropbear_rsa_host_key) on vda
gzip(197): dirtied inode 714 (udev-cache.tar.gz) on vda
gzip(197): dirtied inode 714 (udev-cache.tar.gz) on vda
gzip(197): dirtied inode 714 (udev-cache.tar.gz) on vda
login(347): dirtied inode 543 (motd) on vda
[...]
```

Only the file name is displayed, not full path name, but it's enough

Common problems

- A lot of stuff happens on first boot
 - e.g. Dropbear writes SSH keys to /etc/dropbear, udev snapshot saved to /etc/udev-cache.tar.gz
 - Some of these things can be done at build time
 - Others need be changed to put stuff in a non-volatile state partition
- System config, including network parameters
- Saving random-seed
- Log files

Putting it right

- In reality, packages store state in various places
- · Pragmatic solutions include:
 - Adding symlinks from rootfs to non volatile storage: e.g. /etc to /data/etc
 - · Using overlayfs or similar to overlay rootfs with non volatile storage

Putting it right: first pass

- Create a new partition for non-volatile state
 - For example /data (an idea borrowed from Android)
- Remount rootfs readonly
- Use tmpfs for volatile state (/run and /var/volatile)

/etc/fstab

```
/dev/root
                          auto
                                      defaults.ro
/dev/vdb
           /data
                          ext4
                                      defaults
                                      mode=0755.nodev.nosuid.strictatime 0 0
tmpfs
           /run
                          tmpfs
           /var/volatile
                          tmpfs
                                      defaults
tmpfs
                                                               0
[...]
```

log files

- Many daemons write log files to /var/log
 - Including syslogd
- Usually, such log files are unnecessary for embedded
- Solution: make log files part of the volatile state by mounting a tmpfs on /var/log

Poky core-image-minimal does this already:

```
# ls -ld /var/log
lrwxrwxrwx 1 root root 12 Oct 6 14:11 /var/log -> volatile/log
# mount
tmpfs on /var/volatile type tmpfs (rw,relatime)
```

random-seed

- Saving random-seed at shutdown is necessary when using /dev/urandom
- Solution: make it part of non-volatile state by symlinking /var/lib/urandom to /data/var/lib/urandom
- Or, modify the shell script that creates and restores random-seed

ssh keys

- The Dropbear ssh daemon stores keys in /etc/dropbear or /var/lib/dropbear
- Keys should be unique per device
- Solution: either move to non-volatile state by symlinking /var/lib/dropbear to /data
- · Or, generate and pre load keys when device is provisioned

Doing less on first boot

- · Build packages with sensible defaults so they work without config files
- Or, pre-load config files into the /data partition
- Generate per-device configuration at build-time rather than at first-boot
- Populate configuration files that are common for many packages (e.g. passwd, group with the composite requirements of all

What about Android/Brillo?

Android is a good example

- rootfs moved into /system, which is read only and stateless
- Non-volatile state in /data
- Has factory reset
- Has OTA update

What about Yocto Project?

Current versions of Yocto Project have a partial solution

Add to your config:

```
IMAGE_FEATURES = "read-only-rootfs"
```

- Sets ROOTFS_READ_ONLY=yes in /etc/default/rcS
- Mounts rootfs ro and /var/lib as tmpfs
- · But, nowhere to store non volatile state

Conclusion

- Read-only rootfs makes for better embedded systems
- Make rootfs stateless by moving state into non-volatile and volatile filesystems
 - · diskstats and block_dump will help identify state

Questions?

Slides on Slide Share

```
http://www.slideshare.net/chrissimmonds/readonly-rootfs-theory-and-practice
```

Containers

- Containerised operating systems (e.g CoreOS) are a popular way to implement Internet/cloud services
- They have a similar problem:
 - Removing state from containers makes deployment easier
- Consequently, they are leading the way on statelessness
- Containers may be a useful way to deploy embedded
- swupd from Clear Linux
- resoin.io

Embedded != cloud service

- Use cases differ
- Cloud
- can assume high speed network to supply config (via LDAP, NFS, DHCP, etc)
- Emphasis on run-time integration of different app images into a base os
- Embedded
- Static app image (in a monolithic rootfs)
- Have to function stand-alone