FsLab实验报告

姓名: 梅祎航 学号: 2022201459

本次实验旨在完成一个VSFS,我通过仔细的思考,设计了一套VSFS的体系。能保证完成所有的功能要求,并通过了所有的16个测试点。同时为一些函数加上了互斥锁,以防止多线程的错误操作。

整体结构

如上图,整体的结构设计如上,接下来我会一一讲解每个部分以及相关的辅助函数

输入输出

由于disk_read和disk_write没有报错功能,所以我对这两个函数在外面包装了一层报错,定义为safe_read和safe_write,之后的读入和写入都用我包装后的函数(虽然直到写完了这个报错也没派上用场)

```
int safe_read(int block_id, void* buffer) {
    if (disk_read(block_id, buffer) != 0) {
        fprintf(stderr, "Error reading block %d\n", block_id);
        return -1;
    }
    return 0;
}
//safe_write差不多,我就不放上来了
```

锁

添加了三种锁,分别锁住对bitmap, inode和目录的操作

```
pthread_mutex_t bitmap_mutex = PTHREAD_MUTEX_INITIALIZER;
pthread_mutex_t inode_mutex = PTHREAD_MUTEX_INITIALIZER;
```

```
pthread_mutex_t dir_mutex = PTHREAD_MUTEX_INITIALIZER;

void init_locks() {
    pthread_mutex_init(&bitmap_mutex, NULL);
    pthread_mutex_init(&inode_mutex, NULL);
    pthread_mutex_init(&dir_mutex, NULL);
}

void destroy_locks() {
    pthread_mutex_destroy(&bitmap_mutex);
    pthread_mutex_destroy(&inode_mutex);
    pthread_mutex_destroy(&dir_mutex);
}
```

Super Block

Super Block的设计

对于Super Block, 我的想法是存放根目录的编号以及函数fs_statfs所需要的信息, 因此有以下的结构体:

```
struct superblock {
    unsigned long f_bsize; // 块大小
    uint32_t f_blocks; // 块数量
    uint32_t f_bfree; // 空闲块数量
    uint32_t f_bavail; // 可用块数量
    uint32_t f_files; // 文件节点数
    uint32_t f_ffree; // 空闲节点数
    uint32_t f_favail; // 可用节点数
    unsigned long f_namemax; // 文件名长度上限
    uint32_t root_inode; // 根目录的 i 节点编号
};
```

辅助函数

有初始化函数

```
void init_superblock() {
   struct superblock sb;
   sb.f_bsize = BLOCK_SIZE;
   sb.f_blocks = BLOCK_NUM;
   sb.f_bfree = BLOCK_NUM - (1 + 2 + 1 + INODE_BLOCKS); // 初始化时空闲块

数
   sb.f_bavail = sb.f_bfree;
   sb.f_files = MAX_INODE; // 总文件节点数
   sb.f_ffree = sb.f_files; // 初始化时空闲节点数
   sb.f_favail = sb.f_ffree;
   sb.f_namemax = NAME_MAX;
   sb.root_inode = ROOT_INODE; // 根目录的 i 节点编号
   safe_write(0,&sb);
}
```

最初我还定义了四个辅助函数data_plus(),data_minus(),inode_plus(),inode_minus()用来实现f_bavail, f_bfree, f_ffree, f_favail的增减,但后来发现还不如在fs_statfs函数被调用时直接统计Bitmap的0的个数,故而放弃了上面四个函数。。。

BitMap

BitMap的设计

由于有65536个block, 所以Data Block Bitmap有65536个Bit, 也就是65536/8/4096 = 2Block

由于需要容纳32768个inode, 所以Inode Bitmap有32768个Bit, 也就是32768/8/4096=1Block

故而将Block1和Block2作为Data Block Bitmap,将Block3作为Inode Bitmap。

辅助函数

对于Bitmap, 主要有三种辅助函数:

• 第一种,对Bitmap直接操作的,以下三个函数分别是对指定的Bit进行赋1,赋0,检验是否为1:

```
void set_bit(uint8_t *bitmap, int index) {
    bitmap[index / 8] |= (1 << (index % 8));
}

void clear_bit(uint8_t *bitmap, int index) {
    bitmap[index / 8] &= ~(1 << (index % 8));
}

int test_bit(uint8_t *bitmap, int index) {
    return bitmap[index / 8] & (1 << (index % 8));
}</pre>
```

• 第二种,在Bitmap上分配或释放一个Bit

```
// BITMAP上分配1个新的数据块
int allocate block in bitmap() {
    pthread mutex lock(&bitmap mutex);
    uint8 t data block bitmap[BLOCK SIZE * 2];
    safe read(1, data block bitmap);
    safe read(2, data block bitmap + BLOCK SIZE);
    for (int i = 0; i < BLOCK NUM; ++i) {
        if (!test bit(data block bitmap, i)) {
            set bit(data block bitmap, i);
            safe write(1,data block bitmap);
            safe write(2, data block bitmap + BLOCK SIZE);
            pthread mutex unlock(&bitmap mutex);
            return i;
        }
    pthread mutex unlock(&bitmap mutex);
    return -1;
// BITMAP上释放1个数据块
void free block in bitmap(int block id) {
    pthread mutex lock(&bitmap mutex);
    uint8 t data block bitmap[BLOCK SIZE * 2];
    safe read(1, data block bitmap);
    safe read(2, data block bitmap + BLOCK SIZE);
    clear bit(data block bitmap, block id);
    safe write(1,data block bitmap);
    safe write(2, data block bitmap + BLOCK SIZE);
   pthread mutex unlock(&bitmap mutex);
// BITMAP上分配新的1个inode
int get_new_inode_bit_in_bitmap() {
    pthread_mutex_lock(&bitmap_mutex);
    uint8 t inode bitmap[BLOCK SIZE];
    safe_read(3, inode_bitmap);
    for (int i = 0; i < MAX_INODE; ++i) {</pre>
        if (!test_bit(inode_bitmap, i)) {
            set_bit(inode_bitmap, i);
            safe_write(3,inode_bitmap);
            pthread_mutex_unlock(&bitmap_mutex);
            return i;
    pthread_mutex_unlock(&bitmap_mutex);
    return -1; // 没有空闲 inode
// BITMAP上释放1个inode
void free_inode_in_bitmap(int inode_id) {
```

```
pthread_mutex_lock(&bitmap_mutex);
uint8_t inode_bitmap[BLOCK_SIZE];
safe_read(3, inode_bitmap);
clear_bit(inode_bitmap, inode_id);
safe_write(3,inode_bitmap);
pthread_mutex_unlock(&bitmap_mutex);
}
```

• 第三种,专门供函数fs_statfs用的,用来统计多少Bit是空闲的

```
//bitmap上数有多少1
int count_free_bits(uint8_t *bitmap, int size) {
    int count = 0;
    for (int i = 0; i < size * 8; ++i) {
        if (!test_bit(bitmap, i)) {
            count++;
          }
    }
    return count;
}
```

Inode

inode的设计

Inode的设计是比较重要的,特别是指针的设计。

注意到每个Block都可以用0~65535的数字表示,所以在本次实验中,可以用一个2Byte的数字做"指针"。

因此,一个indirect_pointer指向的Block可以存放4096/2 = 2048个指针,而这2048个Block可以存放2048*4096B = 8MB,正好符合最多8MB的空间要求,所以Inode中只需要一个indirect pointer即可,不需要pointer或double pointer。有以下示意图:

```
Inode
|---indirect_pointer
|---Blocki
|--Blocki1
|--Blocki2
|...
|--Blocki2048
```

由上述论述设计的Inode如下,按照内存从小到大声明,使得对齐浪费的字节只有1Byte,总共32Byte:

```
struct inode {
   uint8_t is_file; // 是不是普通文件
                                                         // 1
byte
   uint16 t indirect pointer; // 一级间接块编号
                                                         // 2
bytes
                  // 文件大小(字节)
   uint32 t size;
                                                        // 4
   time_t st_Atime;
                           // 被访问的时间
                                                         // 8
bytes
   time t st Mtime;
                           // 被修改的时间
                                                         // 8
bytes
                           // 状态改变时间
                                                         // 8
   time t st Ctime;
bytes
};//共32 byte
```

所以一个Block能放下4096/32 = 128个Inode,要放下32768需要32768/128=256个Block。因此将Block4~Block259作为Inode区。

目录的设计

对于目录文件,其indirect_pointer与普通文件不同。

首先我设计了如下结构体作为目录的条目。这样设计既能巧妙地对齐32字节,又使得name后面的字节一定为0,杜绝了name字符串末尾不为0的问题(这个问题不这样解决的话十分麻烦)。:

这样的话,一个Block能存放4096/32=128个entry,根据以下的示意图,一个目录文件能记录2048*128 = 262144个文件,其总量超过了Inode的总数32768,能记录下所有的Inode。

辅助函数

这部分的辅助函数, 注释解释地很清楚了, 我就不赘述了:

```
//创建一个新的inode并写入inode id指定的inode区
void create_inode(int inode_id, uint8_t is_file) {
   struct inode a;
    a.is file = is file;
    a.size = 0;
    a.indirect_pointer = allocate_block_in_bitmap();
    clear block(a.indirect pointer);
    a.st Atime = time(NULL);
    a.st Mtime = a.st Atime;
   a.st Ctime = a.st Atime;
   write_inode(inode_id, &a);
//把栈上的inode写入inode index指定的inode区
int write_inode(int inode_index, struct inode *a) {
    pthread mutex lock(&inode mutex);
    uint8 t inode block[BLOCK SIZE];
    int block_num = 4 + (inode_index / INODE_NUM);
    int offset = (inode index % INODE NUM) * INODE SIZE;
    if (safe read(block num, inode block) != 0) {
       pthread mutex unlock(&inode mutex);
       return -1;
   memcpy(inode block + offset, a, sizeof(struct inode));
    if (safe_write(block_num, inode_block) != 0) {
       pthread_mutex_unlock(&inode_mutex);
       return -1;
    pthread_mutex_unlock(&inode_mutex);
    return 0;
//把inode_index指定的inode写入栈上的inode
int read_inode(int inode_index, struct inode *inode) {
    pthread_mutex_lock(&inode_mutex);
    uint8 t buffer[BLOCK_SIZE];
    int block_num = 4 + (inode_index / INODE_NUM);
    int offset = (inode_index % INODE_NUM) * INODE_SIZE;
    if (safe_read(block_num, buffer) != 0) {
       pthread_mutex_unlock(&inode_mutex);
```

```
return -1;
    }
    memcpy(inode, buffer + offset, sizeof(struct inode));
    pthread mutex unlock(&inode mutex);
   return 0;
}
//通过path找inode,成功则返回inode的index
int get inode by path(const char *path, struct inode *inode) {
    if (strncmp(path, "/", 1) == 0 && strlen(path) == 1) {
       return read inode(ROOT INODE, inode);
    int inode index = ROOT INODE;
    if (read inode(inode index, inode) != 0) {
       return -ENOENT;
    char *path copy = strdup(path);
    if (!path copy) {
       return -ENOMEM;
    char *token = strtok(path copy, "/");
    while (token != NULL) {
       if(inode->is file)
            return -ENOENT;
        uint16 t pointers[BLOCK SIZE / sizeof(uint16 t)];
        if (safe read(inode->indirect pointer, pointers) != 0) {
           free(path copy);
           return -ENOENT;
        }
        int found = 0;
        for (int i = 0; i < BLOCK_SIZE / sizeof(uint16_t); ++i) {</pre>
            if (pointers[i] == 0) {
               continue;;
            }
            uint8_t buffer[BLOCK_SIZE];
            if (safe read(pointers[i], buffer) != 0) {
               free(path_copy);
               return -ENOENT;
            }
            struct dir_entry *entries = (struct dir_entry *)buffer;
            for (int j = 0; j < DIR ENTRY PER BLOCK; ++j) {</pre>
```

```
if(!entries[j].is used)
                    continue;
                if (strcmp(entries[j].name, token) == 0) {
                    inode index = entries[j].inode_index;
                    if (read inode(inode index, inode) != 0) {
                        free(path copy);
                       return -ENOENT;
                    found = 1;
                   break;
               }
            }
            if (found) {
               break;
            }
        }
        if (!found) {
           free(path copy);
           return -ENOENT;
        }
        token = strtok(NULL, "/");
    free(path copy);
   return inode index;
//在parent inode这个父文件夹中删除文件name的相关信息并更新时间
int delete_directory_entry(struct inode *parent_inode, int
parent inode index, const char *name) {
    pthread mutex lock(&dir mutex);
    uint8 t pointers[BLOCK_SIZE];
    if (safe read(parent inode->indirect pointer, pointers) != 0) {
       pthread_mutex_unlock(&dir_mutex);
       return -EIO;
    }
    uint16 t *dir block pointers = (uint16 t *)pointers;
    for (int i = 0; i < BLOCK_SIZE / sizeof(uint16_t); ++i) {</pre>
        if (dir block pointers[i] == 0) {
           continue;
       uint8 t block[BLOCK SIZE];
        if (safe_read(dir_block_pointers[i], block) != 0) {
            pthread mutex unlock(&dir mutex);
```

```
return -EIO;
        }
        struct dir entry *entries = (struct dir entry *)block;
        for (int j = 0; j < DIR ENTRY PER BLOCK; ++j) {</pre>
            if (entries[j].is used && strcmp(entries[j].name, name) ==
0) {
                int inode index = entries[j].inode index;
                entries[j].is used = 0;
                entries[j].inode index = 0;
                if (safe write(dir block pointers[i], block) != 0) {
                    pthread mutex unlock(&dir mutex);
                    return -EIO;
                parent inode->st Mtime = parent inode->st Ctime =
time(NULL);
                if (write inode(parent inode index, parent inode) != 0)
                    pthread mutex unlock(&dir mutex);
                    return -EIO;
                pthread_mutex_unlock(&dir mutex);
                return inode index;
        }
    pthread mutex unlock(&dir mutex);
    return -ENOENT;
//往parent inode这个父文件夹中写入文件name的相关信息并更新时间
int add_directory_entry(struct inode *parent_inode, int
parent inode index, const char *name, int inode index) {
    pthread_mutex_lock(&dir_mutex);
    uint8_t pointers[BLOCK SIZE];
    if (safe_read(parent_inode->indirect_pointer, pointers) != 0) {
       pthread mutex unlock(&dir mutex);
       return -EIO;
    uint16_t *dir_block_pointers = (uint16_t *)pointers;
    for (int i = 0; i < BLOCK SIZE / sizeof(uint16 t); ++i) {</pre>
        if (dir_block_pointers[i] == 0) {
            dir_block_pointers[i] = allocate_block_in_bitmap();
            if (dir block pointers[i] < 0) {</pre>
                pthread mutex unlock(&dir mutex);
                return -ENOSPC;
            }
```

```
clear_block(dir_block_pointers[i]);
            safe write(parent inode->indirect pointer, pointers);
        uint8 t block[BLOCK SIZE];
        if (safe read(dir block pointers[i], block) != 0) {
            pthread mutex unlock(&dir mutex);
            return -EIO;
        struct dir entry *entries = (struct dir entry *)block;
        for (int j = 0; j < DIR_ENTRY PER BLOCK; ++j) {</pre>
            if (!entries[j].is used) {
                strncpy(entries[j].name, name, NAME MAX);
                entries[j].inode index = inode index;
                entries[j].is used = 1;
                if (safe write(dir block pointers[i], block) != 0) {
                    pthread mutex unlock(&dir mutex);
                    return -EIO;
                parent inode->st Mtime = parent inode->st Ctime =
time(NULL);
                if (write inode(parent inode index, parent inode) != 0)
                    pthread mutex unlock(&dir mutex);
                    return -EIO;
                pthread mutex unlock(&dir mutex);
                return 0;
    pthread_mutex_unlock(&dir_mutex);
    return -ENOSPC;
//修改inode的大小为new size
int resize file(struct inode *inode, size t new size) {
    size t current blocks = (inode->size + BLOCK SIZE - 1) / BLOCK SIZE;
    size t new blocks = (new size + BLOCK SIZE - 1) / BLOCK SIZE;
    uint16_t pointers[BLOCK_SIZE / sizeof(uint16_t)];
    if (safe read(inode->indirect pointer, pointers) != 0) {
       return -EIO;
    if (new size > inode->size) { // 扩展文件
        for (size_t i = current_blocks; i < new_blocks; ++i) {</pre>
            int new block = allocate block in bitmap();
```

```
if (new_block < 0) {
    return -ENOSPC;
}
clear_block(new_block);
pointers[i] = new_block;
}
} else if (new_size < inode->size) { // 缩小文件
    for (size_t i = new_blocks; i < current_blocks; ++i) {
        clear_block(pointers[i]);
        free_block_in_bitmap(pointers[i]);
        pointers[i] = 0;
}

if (safe_write(inode->indirect_pointer, pointers) != 0) {
    return -EIO;
}

inode->size = new_size;
return 0;
}
```

Data Block

辅助函数

这一部分主要就有两个辅助函数:

- 一个是清空一个Block
- 一个是清空一个inode,包括其indirect_pointer指向的空间,indirect_pointer指向空间内存放的指针指向的空间,并在Bitmap上把相应的inode和data block置0

```
free_block_in_bitmap(pointers[i]);
}

clear_block(inode->indirect_pointer);
free_block_in_bitmap(inode->indirect_pointer);
free_inode_in_bitmap(inode_index);
}
```

主体函数

mkfs

做了五件事,初始化锁,Super Block,Bitmap,Inode,然后创建根目录的Inode。

```
int mkfs()
    init locks();
   //初始化super blcok
    struct superblock sb;
    init superblock(&sb);
    // 初始化data Bitmap
    uint8 t data block bitmap[BLOCK SIZE * 2] = {0};
    for (int i = 0; i \le 259; i++) {
        set bit(data block bitmap, i);
    if (safe write(1, data block bitmap) != 0 || safe write(2,
data_block_bitmap + BLOCK_SIZE) != 0) {
       return -1;
    }
    // 初始化inode Bitmap
    uint8_t inode_bitmap[BLOCK_SIZE] = {0};
    set bit(inode bitmap, ROOT INODE);
    if (safe_write(3, inode_bitmap) != 0) {
       return -1;
    // 初始化inode
    uint8_t inode_blocks[BLOCK_SIZE * INODE_BLOCKS] = {0};
    for (int i = 0; i < INODE BLOCKS; i++) {</pre>
        if (safe_write(4 + i, inode_blocks + i * BLOCK_SIZE) != 0) {
           return -1;
```

```
//写入根目录inode
create_inode(0,0);
return 0;
}
```

fs_getattr

这个函数很简单,只需要通过path找到对应的inode,然后根据inode填写attr即可。

```
int fs_getattr (const char *path, struct stat *attr)
{
    printf("fs_getattr is called:%s\n",path);
    struct inode inode;
    if (get_inode_by_path(path, &inode) < 0) {
        return -ENOENT;
    }

    attr->st_mode = inode.is_file ? REGMODE : DIRMODE;
    attr->st_nlink = 1;
    attr->st_uid = getuid();
    attr->st_gid = getgid();
    attr->st_size = inode.size;
    attr->st_atime = inode.st_Atime;
    attr->st_mtime = inode.st_Mtime;
    attr->st_ctime = inode.st_Ctime;

    return 0;
}
```

fs_readdir

主要做了以下几件事:

- 根据path找到inode, 并检验其是目录
- 读取indirect_pointer指向的Block并存入pointers, 遍历pointers中的2048个指针, 若指针
 - 。 是0,则跳过
 - 。 非0, 读取指针指向的Block到entries, 遍历128个entries
 - 如果is_used==1, 则filler(buffer, entries[j].name, NULL, 0)
- 最后更新inode的Atime, 并调用write inode写入磁盘

```
int fs readdir(const char *path, void *buffer, fuse_fill_dir_t filler,
off t offset, struct fuse file info *fi) {
    printf("fs_readdir is called:%s\n",path);
    struct inode dir inode;
    int inode index = get inode by path(path, &dir inode);
    if (inode index < 0) {</pre>
       return inode index;
    }
    if (dir inode.is file) {
       return -ENOTDIR;
    }
    uint16 t pointers[BLOCK SIZE / sizeof(uint16 t)];
    if (safe read(dir inode.indirect pointer, pointers) != 0) {
       return -EIO;
    }
    for (int i = 0; i < BLOCK SIZE / sizeof(uint16 t); ++i) {</pre>
        if (pointers[i] == 0) {
           continue;
        }
        uint8 t block[BLOCK SIZE];
        if (safe read(pointers[i], block) != 0) {
           return -EIO;
        struct dir entry *entries = (struct dir entry *)block;
        for (int j = 0; j < DIR ENTRY PER BLOCK; ++j) {</pre>
            if(!entries[j].is used)
                continue;
            if (filler(buffer, entries[j].name, NULL, 0) != 0) {
               return -ENOMEM;
    }
    dir inode.st Atime = time(NULL);
    if (write inode(inode index, &dir inode) != 0) {
       return -EIO;
    }
   return 0;
```

fs_read

这个函数和上个函数差不多,主要干了以下的事情:

- 根据path找inode
- 把indirect_block指向的Block存入data_block_pointers, 然后计算初始读取的位置 block_offset, 从data_block_pointers[block_offset]指向的Block开始读,一直读完size个字节。把读入的内容全存入buffer
- 最后更新Atime

```
int fs read(const char *path, char *buffer, size t size, off t offset,
struct fuse file info *fi) {
    printf("fs read is called:%s\n",path);
    struct inode file inode;
    int inode index = get inode by path(path, &file inode);
    if (inode index < 0) {</pre>
       return inode index;
    }
    if (!file inode.is file) {
       return -EISDIR;
    if (offset >= file inode.size) {
       return 0;
    size t bytes to read = min(size, file inode.size - offset);
    uint8 t indirect block[BLOCK SIZE];
    if (safe read(file inode.indirect pointer, indirect block) != 0) {
       return -EIO;
    uint16 t *data block pointers = (uint16 t *)indirect block;
    size_t bytes_read = 0;
    size t block offset = offset / BLOCK SIZE;
    size_t block_start_offset = offset % BLOCK_SIZE;
    while (bytes_read < bytes_to_read) {</pre>
        if (data_block_pointers[block_offset] == 0) {
           return -EIO;
        }
        uint8_t data_block[BLOCK_SIZE];
        if (safe_read(data_block_pointers[block_offset], data_block) !=
0) {
           return -EIO;
        size t bytes from block = min(BLOCK SIZE - block start offset,
```

fs_mknod

这个函数是创建一个常规文件的函数, 主要干了以下的事情:

- 使用dirname和basename把path分成父文件夹路径dir_path和文件名字base_name, 然后根据dir_path找到父文件夹的inode为dir_inode
- 调用get_new_inode_bit_in_bitmap为新文件在Bitmap上找一个空位 new_inode_index
- 调用create_inode在new_inode_index创建一个常规文件
- 调用add_directory_entry在父文件夹中添加这个新文件的信息

```
int fs mknod(const char *path, mode t mode, dev t dev) {
    printf("fs mknod is called:%s\n",path);
    char *path copy1 = strdup(path);
    char *path copy2 = strdup(path);
    if (!path copy1 || !path copy2) {
        free(path copy1);
        free (path copy2);
        return -ENOMEM;
    char *dir path = dirname(path copy1);
    char *base name = basename(path copy2);
    struct inode dir inode;
    int parent inode index = get inode by path(dir path, &dir inode);
    if (parent_inode_index < 0) {</pre>
        free(path copy1);
        free (path copy2);
        return parent_inode_index;
    if (dir inode.is file) {
        free (path copy1);
        free(path copy2);
        return -ENOTDIR;
    int new inode index = get new inode bit in bitmap();
    if (new inode index < 0) {</pre>
        free (path_copy1);
        free(path copy2);
        return -ENOSPC;
    create_inode(new_inode_index, 1);
```

```
int ret = add_directory_entry(&dir_inode, parent_inode_index,
base_name, new_inode_index);
  free(path_copy1);
  free(path_copy2);
  return ret;
}
```

fs_mkdir

这个函数是创建一个目录文件,和上一个函数基本一样的,除了一点:

• create_inode的第二个参数为0,代表此时创建的是目录文件

```
int fs mkdir(const char *path, mode t mode) {
    printf("fs mkdir is called:%s\n",path);
    char *path copy1 = strdup(path);
    char *path copy2 = strdup(path);
    if (!path copy1 || !path copy2) {
        free(path copy1);
        free(path copy2);
        return -ENOMEM;
    }
    char *dir path = dirname(path copy1);
    char *base name = basename(path copy2);
    struct inode dir inode;
    int parent inode index = get inode by path(dir path, &dir inode);
    if (parent inode index < 0) {</pre>
        free (path_copy1);
        free(path copy2);
       return parent_inode_index;
    }
    if (dir inode.is file) {
       free (path_copy1);
       free(path_copy2);
       return -ENOTDIR;
    }
    int new_inode_index = get_new_inode_bit_in_bitmap();
    if (new inode index < 0) {</pre>
       free (path_copy1);
        free(path copy2);
        return -ENOSPC;
```

```
create_inode(new_inode_index, 0);

int ret = add_directory_entry(&dir_inode, parent_inode_index,
base_name, new_inode_index);
  free(path_copy1);
  free(path_copy2);
  return ret;
}
```

fs_rmdir

这个函数是删除一个目录文件, 主要流程是:

- 使用dirname和basename把path分成父文件夹路径dir_path和文件名字base_name,然后根据dir_path找到父文件夹的inode为dir_inode,根据base_name找到目标文件的inode为target_inode
- 调用delete_directory_entry, 在dir_inode中删除目标文件的信息
- 调用free_inode_and_blocks清空目标文件

```
int fs rmdir(const char *path) {
    printf("fs rmdir is called:%s\n",path);
    char *path copy1 = strdup(path);
    char *path copy2 = strdup(path);
    if (!path copy1 || !path copy2) {
        free(path copy1);
       free(path copy2);
       return -ENOMEM;
    char *dir_path = dirname(path_copy1);
    char *base name = basename(path copy2);
    struct inode dir inode;
    int parent_inode_index = get_inode_by_path(dir_path, &dir_inode);
    if (parent inode index < 0) {</pre>
        free (path_copy1);
       free (path_copy2);
       return parent inode index;
    }
    if (dir_inode.is_file) {
       free(path copy1);
       free(path_copy2);
       return -ENOTDIR;
```

```
struct inode target inode;
    int target inode index = get inode by path(path, &target inode);
    if (target_inode_index < 0) {</pre>
        free(path copy1);
        free(path copy2);
        return target inode index;
    }
    if (target inode.is file) {
        free(path copy1);
        free(path copy2);
       return -ENOTDIR;
    }
    int ret = delete directory entry(&dir inode, parent inode index,
base name);
    if (ret < 0) {
        free(path copy1);
        free(path copy2);
       return ret;
    }
    free inode and blocks (&target inode, target inode index);
    free(path copy1);
    free(path copy2);
    return 0;
}
```

fs_unlink

这个函数就是删一个常规文件, 和上一个函数一模一样

```
int fs_unlink(const char *path) {
    printf("fs_unlink is called:%s\n",path);
    char *path_copy1 = strdup(path);
    char *path_copy2 = strdup(path);
    if (!path_copy1 || !path_copy2) {
        free(path_copy1);
        free(path_copy2);
        return -ENOMEM;
    }
    char *dir_path = dirname(path_copy1);
    char *base_name = basename(path_copy2);
    struct inode dir_inode;
```

```
int parent_inode_index = get_inode_by_path(dir_path, &dir_inode);
    if (parent inode index < 0) {</pre>
        free (path_copy1);
        free(path copy2);
        return parent inode index;
    }
    if (dir inode.is file) {
        free(path copy1);
       free(path copy2);
        return -ENOTDIR;
    }
    struct inode target inode;
    int target inode index = get inode by path(path, &target inode);
    if (target inode index < 0) {</pre>
        free(path copy1);
       free(path copy2);
       return target inode index;
    }
    if (!target inode.is file) {
       free(path copy1);
       free(path copy2);
       return -EISDIR;
    }
    int ret = delete directory entry(&dir inode, parent inode index,
base name);
    if (ret < 0) {
       free(path copy1);
       free(path_copy2);
       return ret;
    }
    free inode and blocks (&target inode, target inode index);
    free(path copy1);
    free (path_copy2);
    return 0;
```

fs_rename

这个函数是把oldpath的文件转移到newpath, 主要流程:

- 通过oldpath和newpath分别找到四个inode,即两个路径的父文件夹和目标文件,记为old_dir_inode(老父文件夹),target_inode(老文件),new_dir_inode(新父文件夹),check_inode(新文件)
- 调用add_directory_entry, 在new_dir_inode中添加新文件的名字和inode编号
- 调用delete_directory_entry, 在old_dir_inode中删除target_inode

```
int fs rename(const char *oldpath, const char *newpath) {
    printf("fs rename is called:%s\n",oldpath);
    char *oldpath copy1 = strdup(oldpath);
    char *oldpath copy2 = strdup(oldpath);
    if (!oldpath copy1 || !oldpath copy2) {
        free(oldpath copy1);
        free(oldpath copy2);
       return -ENOMEM;
    }
    char *old dir path = dirname(oldpath copy1);
    char *old base name = basename(oldpath copy2);
    struct inode old dir inode;
    int old parent inode index = get inode by path(old dir path,
&old dir inode);
    if (old parent inode index < 0) {</pre>
        free(oldpath copy1);
        free(oldpath copy2);
       return old parent inode index;
    }
    struct inode target inode;
    int target inode index = get inode by path(oldpath, &target inode);
    if (target inode index < 0) {</pre>
        free(oldpath copy1);
        free(oldpath_copy2);
        return target_inode_index;
    }
    char *newpath copy1 = strdup(newpath);
    char *newpath_copy2 = strdup(newpath);
    if (!newpath_copy1 || !newpath_copy2) {
        free(oldpath_copy1);
        free(oldpath_copy2);
        free(newpath copy1);
        free(newpath_copy2);
        return -ENOMEM;
```

```
char *new dir path = dirname(newpath copy1);
    char *new base name = basename(newpath copy2);
    struct inode new dir inode;
    int new parent inode index = get inode by path(new dir path,
&new dir inode);
    if (new parent inode index < 0) {</pre>
        free(oldpath copy1);
        free(oldpath copy2);
        free(newpath copy1);
       free(newpath copy2);
       return new parent inode index;
    struct inode check inode;
    int check result = get inode by path(newpath, &check inode);
    if (check result > 0) {
        free(oldpath copy1);
        free(oldpath copy2);
        free(newpath copy1);
       free(newpath copy2);
       return -EEXIST;
    int ret = add directory entry(&new dir inode,
new parent inode index, new base name, target inode index);
    if (ret < 0) {
        free(oldpath copy1);
        free(oldpath copy2);
        free(newpath copy1);
       free(newpath copy2);
       return ret;
    }
    ret = delete_directory_entry(&old_dir_inode, old_parent_inode_index,
old base name);
    if (ret < 0) {
       free(oldpath copy1);
        free(oldpath copy2);
        free(newpath copy1);
       free(newpath copy2);
       return ret;
    }
    free(oldpath copy1);
    free(oldpath copy2);
    free(newpath copy1);
    free (newpath copy2);
    return 0;
```

}

fs write

这个和fs_read差不多,不过是操作反过来了而已,不再赘述

```
int fs write(const char *path, const char *buffer, size t size, off t
offset, struct fuse file info *fi) {
    printf("fs write is called:%s\n",path);
    struct inode file inode;
    int inode index = get inode by path(path, &file inode);
    if (inode index < 0) {</pre>
        return 0;
    if (!file inode.is file) {
       return 0;
    if (fi->flags & O APPEND) {
        offset = file inode.size;
    }
    size t end offset = offset + size;
    if(end offset>file inode.size){
        if (resize file(&file inode, end offset) < 0) {
            return 0;
    size t bytes written = 0;
    uint16 t pointers[BLOCK SIZE / sizeof(uint16 t)];
    if (safe_read(file_inode.indirect_pointer, pointers) != 0) {
        return 0;
    while (bytes_written < size) {</pre>
        size_t block_index = (offset + bytes_written) / BLOCK_SIZE;
        size_t block_offset = (offset + bytes_written) % BLOCK_SIZE;
        size t bytes to write = min(BLOCK SIZE - block offset, size -
bytes_written);
        uint8 t block[BLOCK SIZE];
        if (safe_read(pointers[block_index], block) != 0) {
            return 0;
        }
        memcpy(block + block_offset, buffer + bytes_written,
bytes_to_write);
```

```
if (safe_write(pointers[block_index], block) != 0) {
    return 0;
}

bytes_written += bytes_to_write;
}

file_inode.st_Mtime = file_inode.st_Ctime = time(NULL);
if (write_inode(inode_index, &file_inode) != 0) {
    return 0;
}

return bytes_written;
}
```

fs_truncate

这个函数是修改文件大小的:

- 先通过path找到inode
- 再调用resize_file
- 最后更新下Ctime

```
int fs truncate(const char *path, off t size) {
    printf("fs truncate is called:%s\n",path);
    struct inode file inode;
    int inode index = get inode by path(path, &file inode);
    if (inode index < 0) {</pre>
       return inode index;
    }
    if (!file_inode.is_file) {
       return -EISDIR;
    }
    int ret = resize_file(&file_inode, size);
    if (ret < 0) {
       return ret;
    }
    file_inode.st_Ctime = time(NULL);
    if (write_inode(inode_index, &file_inode) != 0) {
       return -EIO;
    return 0;
```

fs_utime

这个函数是修改文件时间的:

- 先通过path找到inode
- 修改时间

```
int fs_utime(const char *path, struct utimbuf *buffer) {
    printf("fs_utime is called:%s\n",path);
    struct inode file_inode;
    int inode_index = get_inode_by_path(path, &file_inode);
    if (inode_index < 0) {
        return inode_index;
    }

    file_inode.st_Atime = buffer->actime;
    file_inode.st_Mtime = buffer->modtime;

    file_inode.st_Ctime = time(NULL);

    if (write_inode(inode_index, &file_inode) != 0) {
        return -EIO;
    }

    return 0;
}
```

fs_statfs

这个函数就是调用count_free_bits两次去数一数inode和data block的bitmap有多少0,并给stat赋值。

但是我既然设计了Super Block也不能完全不管它,所以我在这个函数中顺便更新了Super Block

```
int fs_statfs(const char *path, struct statvfs *stat) {
   printf("fs_statfs is called:%s\n",path);
   struct superblock sb;
   if (safe_read(0, &sb) != 0) {
      return -EIO;
   }
  memset(stat, 0, sizeof(struct statvfs));
```

```
stat->f_bsize = sb.f_bsize;
    stat->f blocks = sb.f blocks;
    stat->f_files = sb.f_files;
    stat->f namemax = sb.f namemax;
    uint8 t block bitmap[BLOCK SIZE * 2];
    if (safe read(1, block bitmap) != 0 || safe read(2, block bitmap +
BLOCK SIZE) != 0) {
       return -EIO;
    uint8 t inode bitmap[BLOCK SIZE];
    if (safe read(3, inode bitmap) != 0) {
       return -EIO;
    int free blocks = count free bits(block bitmap, BLOCK SIZE * 2);
    int free inodes = count free bits(inode bitmap, BLOCK SIZE);
    sb.f bfree=sb.f bavail=stat->f bfree = stat->f bavail = free blocks;
    sb.f ffree=sb.f favail=stat->f ffree = stat->f favail = free inodes;
    if (safe write(0, &sb) != 0) {
       return -EIO;
   return 0;
```

fs_open

这个函数没啥说的,就是根据fi->flags和O_APPEND来设置inode的size

```
int fs_open(const char *path, struct fuse_file_info *fi) {
   printf("fs_open is called:%s\n",path);
   struct inode file_inode;
   int inode_index = get_inode_by_path(path, &file_inode);
   if (inode_index < 0) {
      return inode_index;
   }

   if (!file_inode.is_file) {
      return -EISDIR;
   }

   if (fi->flags & O_APPEND) {
      fi->fh = file_inode.size;
   } else {
```

```
fi->fh = 0;
}
return 0;
}
```

遇到的问题

本次实验进行的比较顺利,没遇到太多问题。也就以下三个问题:

- dir_entry起初是没有设计is_used这个参数的。导致fs_readdir函数中,把一个Block的128个dir_entry全都调用了filler函数指针,而其中大部分的dir_entry是空的,这让 ls指令出现了问题
- dir_entry—开始设计的时候, name后面是is_used。而对活跃的dir_entry, is_used=1。所以当文件名字正好是24字节时(2.sh), strcmp(name,...)会出现严重的错误。
- 在add_directory_entry等函数中,有时需要修改inode的indirect_pointer指向的那个Block,为其增添一个指针。但是我修改后没有写入磁盘中,导致了一系列错误

附加思考

1.多个线程可能会同时访问和修改共享资源,如全局变量、文件系统结构等。这可能导致以下问题:

• 竞争: 多个线程同时读取和修改共享数据,导致数据不一致。

• 死锁: 两个或多个线程相互等待对方释放资源,导致程序停止运行。

特别对于文件系统来说,有一个经典的读一写问题,如果不妥善处理会使读入和写入磁盘的操作有巨大的问题

2.在文件系统中,以下几个地方需要互斥访问:

位图操作:分配和释放数据块和 inode 时,需要互斥访问位图。

inode 读写:读写 inode 信息时需要互斥访问,以防止数据不一致。

目录操作:添加和删除目录时需要互斥访问。

可以对相应的辅助函数加锁,例如:

```
pthread_mutex_t bitmap_mutex;
pthread_mutex_t inode_mutex;
pthread_mutex_t dir_mutex;

void init_locks() {
    pthread_mutex_init(&bitmap_mutex, NULL);
    pthread_mutex_init(&inode_mutex, NULL);
```

```
pthread_mutex_init(&dir_mutex, NULL);
void destroy locks() {
    pthread mutex destroy(&bitmap mutex);
    pthread mutex destroy(&inode mutex);
    pthread mutex destroy(&dir mutex);
}
int allocate block in bitmap() {
    pthread mutex lock(&bitmap mutex);
    // .....
    pthread mutex unlock(&bitmap mutex);
void free block in bitmap(int block id) {
    pthread mutex lock(&bitmap mutex);
    // .....
    pthread mutex unlock(&bitmap mutex);
int write inode(int inode index, struct inode *a) {
    pthread mutex lock(&inode mutex);
    pthread mutex unlock(&inode mutex);
int read inode(int inode index, struct inode *inode) {
    pthread mutex lock(&inode mutex);
    // .....
    pthread mutex unlock(&inode mutex);
int add directory entry(struct inode *parent inode, int
parent_inode_index, const char *name, int inode_index) {
    pthread mutex lock(&dir mutex);
    // .....
    pthread mutex unlock(&dir mutex);
}
int delete_directory_entry(struct inode *parent_inode, int
parent_inode_index, const char *name) {
    pthread mutex lock(&dir mutex);
    pthread mutex unlock(&dir mutex);
```

样例测试结果

本次实验我的代码能通过所有的16个样例点,具体输入如下:

• 0: 没有输出

• 1

```
dir1 dir2
dir3
```

• 2

```
file1 file2
abcdefghijklmnopqrstuvwx file1 file2
```

• 3

```
0
32
10000
24
```

• 4:这里时区不同,所以样例是+0005,而本地输出是+0001

```
2012-12-21 00:00:36.000000000 +00002012-12-21 00:00:36.000000000 +0000 2012-12-21 00:00:36.000000000 +0000 +00002012-12-21 00:00:36.000000000 +0000
```

• 5

```
dir1 dir3
dir3
dir1 dir3
```

• 6

```
file1 file2 picture
file1 file2
file1
```

```
dir2 file1
dir2 file2
dir3 file2
dir
dir3 file2
```

• 8

```
dir1-1 dir1-2
dir1 dir2
dir2
dir2-1
dir1-1 dir1-2
dir1-1 dir2-2
```

• 9

```
fuse
```

• 10

```
12345
2333333333
2333333333
abcdefg
```

- 11: 没有任何输出,因为fs.c和fs-test.c相同, diff fs.c fs-test.c不会有输出
- 12: 这里的输出和样例有些许不同,File没有带引号,uid和gid也不同,但这是因为系统不同导致的

```
File: file1
Size: 0 Blocks: 0 IO Block: 4096 regular
empty file
Access: (0644/-rw-r--r--) Uid: (1087/u2022201459) Gid: (1002/
2022ICS)
File: file1
Size: 6 Blocks: 0 IO Block: 4096 regular file
Access: (0644/-rw-r--r--) Uid: (1087/u2022201459) Gid: (1002/
2022ICS)
```

- 13: 和第11题类似,没有任何输出,因为test.in和test.in.out相同,diff test.in test.in.out不会有输出
- 14

```
dir0 dir145 dir213 dir282 dir350 dir419 dir488 dir556 dir624
dir693 dir761 dir83 dir899 dir967 file112 file181 file25
file318 file387 file455 file523 file592 file660 file729 file798
file866 file934
      dir146 dir214 dir283 dir351 dir42 dir489 dir557 dir625
dir1
dir694 dir762 dir830 dir9 dir968 file113 file182 file250
file319 file388 file456 file524 file593 file661 file73 file799
file867 file935
dir10 dir147 dir215 dir284 dir352 dir420 dir49 dir558 dir626
dir695 dir763 dir831 dir90 dir969 file114 file183 file251
file32 file389 file457 file525 file594 file662 file730 file8
file868 file936
dir100 dir148 dir216 dir285 dir353 dir421 dir490 dir559 dir627
dir696 dir764 dir832 dir900 dir97 file115 file184 file252
file320 file39 file458 file526 file595 file663 file731 file80
file869 file937
dir1000 dir149 dir217 dir286 dir354 dir422 dir491 dir56
dir697 dir765 dir833 dir901 dir970 file116 file185 file253
file321 file390 file459 file527 file596 file664 file732 file800
file87 file938
dir1001 dir15 dir218 dir287 dir355 dir423 dir492 dir560 dir629
dir698 dir766 dir834 dir902 dir971 file117 file186 file254
file322 file391 file46 file528 file597 file665 file733 file801
file870 file939
dir1002 dir150 dir219 dir288 dir356 dir424 dir493 dir561 dir63
dir699 dir767 dir835 dir903 dir972 file118 file187 file255
file323 file392 file460 file529 file598 file666 file734 file802
file871 file94
dir1003 dir151 dir22 dir289 dir357 dir425 dir494 dir562 dir630
dir7 dir768 dir836 dir904 dir973 file119 file188 file256
file324 file393 file461 file53 file599 file667 file735 file803
file872 file940
dir1004 dir152 dir220 dir29 dir358 dir426 dir495 dir563 dir631
dir70 dir769 dir837 dir905 dir974 file12 file189 file257
file325 file394 file462 file530 file6 file668 file736 file804
file873 file941
dir1005 dir153 dir221 dir290 dir359 dir427 dir496 dir564 dir632
dir700 dir77 dir838 dir906 dir975 file120 file19 file258
file326 file395 file463 file531 file60 file669 file737 file805
file874 file942
dir1006 dir154 dir222 dir291 dir36 dir428 dir497 dir565 dir633
dir701 dir770 dir839 dir907 dir976 file121 file190 file259
file327 file396 file464 file532 file600 file67 file738 file806
file875 file943
dir1007 dir155 dir223 dir292 dir360 dir429 dir498 dir566 dir634
dir702 dir771 dir84 dir908 dir977 file122 file191 file26
file328 file397 file465 file533 file601 file670 file739 file807
file876 file944
```

```
dir1008 dir156 dir224 dir293 dir361 dir43 dir499 dir567 dir635
dir703 dir772 dir840 dir909 dir978 file123 file192 file260
file329 file398 file466 file534 file602 file671 file74 file808
file877 file945
dir1009 dir157 dir225 dir294 dir362 dir430 dir5 dir568 dir636
dir704 dir773 dir841 dir91 dir979 file124 file193 file261
file33 file399 file467 file535 file603 file672 file740 file809
file878 file946
dir101 dir158 dir226 dir295 dir363 dir431 dir50 dir569 dir637
dir705 dir774 dir842 dir910 dir98 file125 file194 file262
file330 file4 file468 file536 file604 file673 file741 file81
file879 file947
dir1010 dir159 dir227 dir296 dir364 dir432 dir500 dir57 dir638
dir706 dir775 dir843 dir911 dir980 file126 file195 file263
file331 file40 file469 file537 file605 file674 file742 file810
file88 file948
dir1011 dir16 dir228 dir297 dir365 dir433 dir501 dir570 dir639
dir707 dir776 dir844 dir912 dir981 file127 file196 file264
file332 file400 file47 file538 file606 file675 file743 file811
file880 file949
dir1012 dir160 dir229 dir298 dir366 dir434 dir502 dir571 dir64
dir708 dir777 dir845 dir913 dir982 file128 file197 file265
file333 file401 file470 file539 file607 file676 file744 file812
file881 file95
dir1013 dir161 dir23 dir299 dir367 dir435 dir503 dir572 dir640
dir709 dir778 dir846 dir914 dir983 file129 file198 file266
file334 file402 file471 file54 file608 file677 file745 file813
file882 file950
dir1014 dir162 dir230 dir3 dir368 dir436 dir504 dir573 dir641
dir71 dir779 dir847 dir915 dir984 file13 file199 file267
file335 file403 file472 file540 file609 file678 file746 file814
file883 file951
dir1015 dir163 dir231 dir30 dir369 dir437 dir505 dir574 dir642
dir710 dir78 dir848 dir916 dir985 file130 file2 file268
file336 file404 file473 file541 file61 file679 file747 file815
file884 file952
dir1016 dir164 dir232 dir300 dir37 dir438 dir506 dir575 dir643
dir711 dir780 dir849 dir917 dir986 file131 file20 file269
file337 file405 file474 file542 file610 file68 file748 file816
file885 file953
dir1017 dir165 dir233 dir301 dir370 dir439 dir507 dir576 dir644
dir712 dir781 dir85 dir918 dir987 file132 file200 file27
file338 file406 file475 file543 file611 file680 file749 file817
file886 file954
dir1018 dir166 dir234 dir302 dir371 dir44 dir508 dir577 dir645
dir713 dir782 dir850 dir919 dir988 file133 file201 file270
file339 file407 file476 file544 file612 file681 file75 file818
file887 file955
```

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dir1019 dir167 dir235 dir303 dir372 dir440 dir509 dir578 dir646
dir714 dir783 dir851 dir92 dir989 file134 file202 file271
file34 file408 file477 file545 file613 file682 file750 file819
file888 file956
dir102 dir168 dir236 dir304 dir373 dir441 dir51 dir579 dir647
dir715 dir784 dir852 dir920 dir99 file135 file203 file272
file340 file409 file478 file546 file614 file683 file751 file82
file889 file957
dir1020 dir169 dir237 dir305 dir374 dir442 dir510 dir58 dir648
dir716 dir785 dir853 dir921 dir990 file136 file204 file273
file341 file41 file479 file547 file615 file684 file752 file820
file89 file958
dir1021 dir17 dir238 dir306 dir375 dir443 dir511 dir580 dir649
dir717 dir786 dir854 dir922 dir991 file137 file205 file274
file342 file410 file48 file548 file616 file685 file753 file821
file890 file959
dir1022 dir170 dir239 dir307 dir376 dir444 dir512 dir581 dir65
dir718 dir787 dir855 dir923 dir992 file138 file206 file275
file343 file411 file480 file549 file617 file686 file754 file822
file891 file96
dir1023 dir171 dir24 dir308 dir377 dir445 dir513 dir582 dir650
dir719 dir788 dir856 dir924 dir993 file139 file207 file276
file344 file412 file481 file55 file618 file687 file755 file823
file892 file960
dir103 dir172 dir240 dir309 dir378 dir446 dir514 dir583 dir651
dir72 dir789 dir857 dir925 dir994 file14 file208 file277
file345 file413 file482 file550 file619 file688 file756 file824
file893 file961
dir104 dir173 dir241 dir31 dir379 dir447 dir515 dir584 dir652
dir720 dir79 dir858 dir926 dir995 file140 file209 file278
file346 file414 file483 file551 file62 file689 file757 file825
file894 file962
dir105 dir174 dir242 dir310 dir38 dir448 dir516 dir585 dir653
dir721 dir790 dir859 dir927 dir996 file141 file21 file279
file347 file415 file484 file552 file620 file69 file758 file826
file895 file963
dir106 dir175 dir243 dir311 dir380 dir449 dir517 dir586 dir654
dir722 dir791 dir86 dir928 dir997 file142 file210 file28
file348 file416 file485 file553 file621 file690 file759 file827
file896 file964
dir107 dir176 dir244 dir312 dir381 dir45 dir518 dir587 dir655
dir723 dir792 dir860 dir929 dir998 file143 file211 file280
file349 file417 file486 file554 file622 file691 file76 file828
file897 file965
dir108 dir177 dir245 dir313 dir382 dir450 dir519 dir588 dir656
dir724 dir793 dir861 dir93 dir999 file144 file212 file281
file35 file418 file487 file555 file623 file692 file760 file829
```

file898 file966

- dir109 dir178 dir246 dir314 dir383 dir451 dir52 dir589 dir657
 dir725 dir794 dir862 dir930 file0 file145 file213 file282
 file350 file419 file488 file556 file624 file693 file761 file83
 file899 file967
- dir11 dir179 dir247 dir315 dir384 dir452 dir520 dir59 dir658
 dir726 dir795 dir863 dir931 file1 file146 file214 file283
 file351 file42 file489 file557 file625 file694 file762 file830
 file9 file968
- dir110 dir18 dir248 dir316 dir385 dir453 dir521 dir590 dir659
 dir727 dir796 dir864 dir932 file10 file147 file215 file284
 file352 file420 file49 file558 file626 file695 file763 file831
 file90 file969
- dir111 dir180 dir249 dir317 dir386 dir454 dir522 dir591 dir66
 dir728 dir797 dir865 dir933 file100 file148 file216 file285
 file353 file421 file490 file559 file627 file696 file764 file832
 file900 file97
- dir112 dir181 dir25 dir318 dir387 dir455 dir523 dir592 dir660
 dir729 dir798 dir866 dir934 file1000 file149 file217 file286
 file354 file422 file491 file56 file628 file697 file765 file833
 file901 file970
- dir113 dir182 dir250 dir319 dir388 dir456 dir524 dir593 dir661
 dir73 dir799 dir867 dir935 file1001 file15 file218 file287
 file355 file423 file492 file560 file629 file698 file766 file834
 file902 file971
- dir114 dir183 dir251 dir32 dir389 dir457 dir525 dir594 dir662
 dir730 dir8 dir868 dir936 file1002 file150 file219 file288
 file356 file424 file493 file561 file63 file699 file767 file835
 file903 file972
- dir115 dir184 dir252 dir320 dir39 dir458 dir526 dir595 dir663
 dir731 dir80 dir869 dir937 file1003 file151 file22 file289
 file357 file425 file494 file562 file630 file7 file768 file836
 file904 file973
- dir116 dir185 dir253 dir321 dir390 dir459 dir527 dir596 dir664
 dir732 dir800 dir87 dir938 file1004 file152 file220 file29
- file358 file426 file495 file563 file631 file70 file769 file837 file905 file974
- dir117 dir186 dir254 dir322 dir391 dir46 dir528 dir597 dir665
 dir733 dir801 dir870 dir939 file1005 file153 file221 file290
 file359 file427 file496 file564 file632 file700 file77 file838
 file906 file975
- dir118 dir187 dir255 dir323 dir392 dir460 dir529 dir598 dir666
 dir734 dir802 dir871 dir94 file1006 file154 file222 file291
 file36 file428 file497 file565 file633 file701 file770 file839
 file907 file976
- dir119 dir188 dir256 dir324 dir393 dir461 dir53 dir599 dir667
 dir735 dir803 dir872 dir940 file1007 file155 file223 file292
 file360 file429 file498 file566 file634 file702 file771 file84
 file908 file977

- dir12 dir189 dir257 dir325 dir394 dir462 dir530 dir6 dir668
 dir736 dir804 dir873 dir941 file1008 file156 file224 file293
 file361 file43 file499 file567 file635 file703 file772 file840
 file909 file978
- dir120 dir19 dir258 dir326 dir395 dir463 dir531 dir60 dir669
 dir737 dir805 dir874 dir942 file1009 file157 file225 file294
 file362 file430 file5 file568 file636 file704 file773 file841
 file91 file979
- dir121 dir190 dir259 dir327 dir396 dir464 dir532 dir600 dir67
 dir738 dir806 dir875 dir943 file101 file158 file226 file295
 file363 file431 file50 file569 file637 file705 file774 file842
 file910 file98
- dir122 dir191 dir26 dir328 dir397 dir465 dir533 dir601 dir670
 dir739 dir807 dir876 dir944 file1010 file159 file227 file296
 file364 file432 file500 file57 file638 file706 file775 file843
 file911 file980
- dir123 dir192 dir260 dir329 dir398 dir466 dir534 dir602 dir671
 dir74 dir808 dir877 dir945 file1011 file16 file228 file297
 file365 file433 file501 file570 file639 file707 file776 file844
 file912 file981
- dir124 dir193 dir261 dir33 dir399 dir467 dir535 dir603 dir672
 dir740 dir809 dir878 dir946 file1012 file160 file229 file298
 file366 file434 file502 file571 file64 file708 file777 file845
 file913 file982
- dir125 dir194 dir262 dir330 dir4 dir468 dir536 dir604 dir673
 dir741 dir81 dir879 dir947 file1013 file161 file23 file299
 file367 file435 file503 file572 file640 file709 file778 file846
 file914 file983
- dir126 dir195 dir263 dir331 dir40 dir469 dir537 dir605 dir674
 dir742 dir810 dir88 dir948 file1014 file162 file230 file3
 file368 file436 file504 file573 file641 file71 file779 file847
 file915 file984
- dir127 dir196 dir264 dir332 dir400 dir47 dir538 dir606 dir675
 dir743 dir811 dir880 dir949 file1015 file163 file231 file30
 file369 file437 file505 file574 file642 file710 file78 file848
- file916 file985 dir128 dir197 dir265 dir333 dir401 dir470 dir539 dir607 dir676 dir744 dir812 dir881 dir95 file1016 file164 file232 file300
- file37 file438 file506 file575 file643 file711 file780 file849 file917 file986
- dir129 dir198 dir266 dir334 dir402 dir471 dir54 dir608 dir677
 dir745 dir813 dir882 dir950 file1017 file165 file233 file301
 file370 file439 file507 file576 file644 file712 file781 file85
 file918 file987
- dir13 dir199 dir267 dir335 dir403 dir472 dir540 dir609 dir678
 dir746 dir814 dir883 dir951 file1018 file166 file234 file302
 file371 file44 file508 file577 file645 file713 file782 file850
 file919 file988

- dir130 dir2 dir268 dir336 dir404 dir473 dir541 dir61 dir679
 dir747 dir815 dir884 dir952 file1019 file167 file235 file303
 file372 file440 file509 file578 file646 file714 file783 file851
 file92 file989
- dir131 dir20 dir269 dir337 dir405 dir474 dir542 dir610 dir68
 dir748 dir816 dir885 dir953 file102 file168 file236 file304
 file373 file441 file51 file579 file647 file715 file784 file852
 file920 file99
- dir132 dir200 dir27 dir338 dir406 dir475 dir543 dir611 dir680
 dir749 dir817 dir886 dir954 file1020 file169 file237 file305
 file374 file442 file510 file58 file648 file716 file785 file853
 file921 file990
- dir133 dir201 dir270 dir339 dir407 dir476 dir544 dir612 dir681
 dir75 dir818 dir887 dir955 file1021 file17 file238 file306
 file375 file443 file511 file580 file649 file717 file786 file854
 file922 file991
- dir134 dir202 dir271 dir34 dir408 dir477 dir545 dir613 dir682
 dir750 dir819 dir888 dir956 file1022 file170 file239 file307
 file376 file444 file512 file581 file65 file718 file787 file855
 file923 file992
- dir135 dir203 dir272 dir340 dir409 dir478 dir546 dir614 dir683
 dir751 dir82 dir889 dir957 file1023 file171 file24 file308
 file377 file445 file513 file582 file650 file719 file788 file856
 file924 file993
- dir136 dir204 dir273 dir341 dir41 dir479 dir547 dir615 dir684
 dir752 dir820 dir89 dir958 file103 file172 file240 file309
 file378 file446 file514 file583 file651 file72 file789 file857
 file925 file994
- dir137 dir205 dir274 dir342 dir410 dir48 dir548 dir616 dir685
 dir753 dir821 dir890 dir959 file104 file173 file241 file31
 file379 file447 file515 file584 file652 file720 file79 file858
 file926 file995
- dir138 dir206 dir275 dir343 dir411 dir480 dir549 dir617 dir686
 dir754 dir822 dir891 dir96 file105 file174 file242 file310
 file38 file448 file516 file585 file653 file721 file790 file859
 file927 file996
- dir139 dir207 dir276 dir344 dir412 dir481 dir55 dir618 dir687
 dir755 dir823 dir892 dir960 file106 file175 file243 file311
 file380 file449 file517 file586 file654 file722 file791 file86
 file928 file997
- dir14 dir208 dir277 dir345 dir413 dir482 dir550 dir619 dir688
 dir756 dir824 dir893 dir961 file107 file176 file244 file312
 file381 file45 file518 file587 file655 file723 file792 file860
 file929 file998
- dir140 dir209 dir278 dir346 dir414 dir483 dir551 dir62 dir689
 dir757 dir825 dir894 dir962 file108 file177 file245 file313
 file382 file450 file519 file588 file656 file724 file793 file861
 file93 file999

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dir141 dir21 dir279 dir347 dir415 dir484 dir552 dir620 dir69
dir758 dir826 dir895 dir963 file109 file178 file246 file314
file383 file451 file52 file589 file657 file725 file794 file862
file930
dir142 dir210 dir28 dir348 dir416 dir485 dir553 dir621 dir690
dir759 dir827 dir896 dir964 file11 file179 file247 file315
file384 file452 file520 file59 file658 file726 file795 file863
file931
dir143 dir211 dir280 dir349 dir417 dir486 dir554 dir622 dir691
dir76 dir828 dir897 dir965 file110 file18 file248 file316
file385 file453 file521 file590 file659 file727 file796 file864
file932
dir144 dir212 dir281 dir35 dir418 dir487 dir555 dir623 dir692
dir760 dir829 dir898 dir966 file111 file180 file249 file317
file386 file454 file522 file591 file66 file728 file797 file865
file933
```

• 15: 这一题没有输出,但我修改了一下样例使之有输出:

```
cd mnt
for ((i=0;i<128;++i)); do
    mkdir dir
    ls
    cd dir
    done</pre>
```

输出见下, 共128个dir:

```
dir
```

dir dir

dir dir dir dir dir dir dir dir dir dir dir dir dir dir dir dir dir dir dir dir dir dir dir dir dir dir dir dir dir dir dir dir dir dir dir dir dir dir dir dir dir dir dir dir dir dir dir dir dir

dir			
dir			