Chapter 26

Presuppositions

(Joint work with Koji Mineshima)

A sentence generally has a foreground and background contents, the former of which is called its *assertion* while the latter of which is called its *presupposition*. For example, the sentence (682) has *John was the one who did it* as its assertion and *someone broke my iPhone* as its presupposition.

(682) It is John who broke my iPhone.

Assertion: John was the one who did it.

- The foreground content: the main point of an utterance
- The main point/foreground of an utterence

Presupposition: Someone broke my iPhone.

- the background content: its truth is usually (normally) assumed to be taken for granted by the participants in the conversation
- something that is already agreed on by the participants in the conversation

The background contents, presupposision, are known to exhibit two puzzles.

26.1 The first puzzule: Presupposition Projection

We can naturally infer (683) from (683a–e), as well as (682) not only from (682) but also from (683a–e). (682) presupposes (683)

- (682) It was John who broke my iPhone.
- (683) Someone broke my iPhone. ((682) presupposes (683))
 - (683) projects out of all the embedded contexts in (684a-e).
- (684) a. It wasn't John who broke my iPhone.

negation

b. Maybe it was John who broke my iPhone.

modal

the

- c. If it was John who broke my iPhone, then he has to fix it.

 antecedent of a conditional
 - ~...ation
- d. Was it John who broke my iPhone?e. Suppose that it was John who broke my iPhone.

question

1 41 4 1

hypothetical assumption

26.1.1 The Case of Entailment

These examples show a striking contrast with the case of entailment as exemplified in the following examples.

- (685) John is an American pianist.
- (686) John is American. ((685) entails (686))
 - (686) does not survive in the contexts (687a-e).
- (687) a. John is not an American pianist.

negation

b. Maybe John is an American pianist.

modal

c. If John is an American pianist, he is skillful.

the antecedent of a conditional

d. Is John an American pianist?

question

e. Suppose that John is an American pianist.

hypothetical assumption

The sentence (??a) entails that John is American, but this entailment does not survive in the environments (??b,c), in contrast to the case of presupposition.

- (685) John is an American pianist. american(john) ∧ pianist(john)
- (687) a. John is not an American pianist. $\neg(\mathbf{american}(\mathbf{j}) \land \mathbf{pianist}(\mathbf{j}))$
 - b. Maybe John is an American pianist.
 - $\Diamond(\operatorname{american}(\mathbf{j}) \land \operatorname{pianist}(\mathbf{j}))$
 - c. If John is an American pianist, he is skillful. $\mathbf{american}(\mathbf{j}) \land \mathbf{pianist}(\mathbf{j}) \rightarrow \mathbf{skillful}(\mathbf{j})$

Standard semantics correctly predicts these patterns:

- $(685) \vdash american(john)$
- (687a) ⊬ american(john)
- (687b) ⊬ american(john)
- (687c) ⊬ american(john)

26.1.2 The Case of Presupposition

- (682) It was John who broke my iPhone. SR_1
- (684) a. It wasn't John who broke my iPhone. $\neg SR_1$
 - b. Maybe it was John who broke my iPhone. $\Diamond SR_1$
 - c. If it was John who broke my iPhone, he has to fix it. $SR_1 \rightarrow \cdots$

What SR accounts for the following inference patterns?

• $SR_1 \Vdash \exists x(\mathbf{broke}(x, \mathbf{my_iphone}))$

- $\neg SR_1 \Vdash \exists x(\mathbf{broke}(x, \mathbf{my_iphone}))$
- $\Diamond SR_1 \Vdash \exists x(\mathbf{broke}(x, \mathbf{my_iphone}))$
- $SR_1 \rightarrow A \Vdash \exists x (\mathbf{broke}(x, \mathbf{my_iphone}))$
- Q: Can "\=" be defined as a standard consequence relation "\="?
- A: No. If that were the case, then $\exists x(\mathbf{broke}(x, \mathbf{my_iphone}))$ was a tautology (under the classical setting).

26.2 The second puzzule: Presupposition Filteration

(688) presupposes that someone broke the window, but the conditional in (689) does not inherit this presupposition. such as (688a) and (689a). the second sentence in the conjunction (??a) or in the conditional (??b).

- (688) It was John who broke the window.
 - \Rightarrow Someone broke the window
- (689) If the window was broken, it was John who broke it.
 - ⇒ Someone broke the window

Similarly for (690) and (691).

- (690) The king of France is wise.
 - \Rightarrow France has a king.
- (691) If France has a king, the king of France is wise.
 - \Rightarrow France has a king.

A presupposition is **filtered** when it occurs in certain contexts.

The problem posed by these examples is to account for the fact that while a simple sentence *The king of France is wise* presupposes that France has a king, neither (??a) nor (??b) inherits this presupposition.

In general, if S' entails the presuppositions of S, constructions like S' and S and If S' then S do not inherit the presuppositions of S.

26.3 Presupposition triggers

- (692) a. The elevator in this building is clean. Description
 - b. There is an elevator in this building.
- (693) a. John's sister is happy. Possessive
 - b. John has a sister.
- (694) a. Bill **regrets** that he lied to Mary. **Factive**
 - b. Bill lied to Mary.
- (695) a. John has **stopped** beating his wife. **Aspectual**
 - b. John has beaten his wife.
- (696) a. Harry managed to find the book. Implicative
 - b. Finding the book required some effort.

(697) a. Sam broke the window **again** today.

Iterative

- b. Sam broke the window before.
- (698) a. It was Sam who broke the window.

Cleft

- b. Someone broke the window.
- (699) a. What John broke was his typewriter.

Pseudo-cleft

- b. John broke something.
- (700) a. $[Pat]_F$ is leaving, **too**. (Focus on Pat)

Additive

b. Someone other than Pat is leaving.

For classical examples of presupposition triggers, see Levinson (1983), Soames (1989), Geurts (1999), and Beaver (2001), among others.

26.4 "Presupposition Is Anaphora" hypothesis

(van der Sandt 1992; Geurts 1996; Kripke 2009): There are striking parallels between anaphoric expressions and presupposition triggers. (van der Sandt, 1992; Geurts, 1999)

Presupposition filtering

- (701) a. John has children and **John's children** are wise.
 - b. If John has children, **John's children** are wise.
- (702) a. The window was broken and it was John who broke it.
 - b. If the window was broken, it was John who broke it.

Why are the presuppositions are filtered out in such contexts as (701) and (702)? Compare (701) and (702) with the paradigm examples of anaphora resolution. which motivated dynamic semantics and discourse representation theory:

Anaphora resolution

- (703) a. John owns a donkey and he beats it.
 - b. If John owns a donkey, he beats it.

How a pronoun it can be an aphorically linked to a quantificational expression a donkey?

26.5 Analysis by DTS

26.5.1 On Filtering

The present account can explain the filtering of presupposition without further stipulation.

(704) If France has a king, the king of France is wise.

The relevant derivation for (??a) goes in the same way as the case of anaphora resolution for the mini-discourse (??). Here we will take a brief look at the case of a conditional sentence in (??b). To begin with, the SR of the sentence (??b) is compositionally obtained via the following derivation tree.

(705)
$$\left(u: \left[\begin{array}{c} x: \mathbf{e} \\ \mathbf{kingOf}(x, fr) \end{array}\right]\right) \rightarrow \left[\begin{array}{c} @\left[\begin{array}{c} x: \mathbf{e} \\ \mathbf{kingOf}(x, fr) \end{array}\right] \right]$$

(706) Type check diagram of (704)

$$\frac{u : \begin{bmatrix} x : e \\ \mathbf{kingOf}(x, fr) \end{bmatrix}^{1}}{u : \begin{bmatrix} x : e \\ \mathbf{kingOf}(x, fr) \end{bmatrix}} : \text{type} \quad \frac{\begin{bmatrix} x : e \\ \mathbf{kingOf}(x, fr) \end{bmatrix} : \text{type}}{\begin{bmatrix} x : e \\ \mathbf{kingOf}(x, fr) \end{bmatrix}} : \text{type} \quad \frac{\begin{bmatrix} x : e \\ \mathbf{kingOf}(x, fr) \end{bmatrix}}{\begin{bmatrix} wise(\pi_{1}v)[?/v] \\ wise(\pi_{1}v) \end{bmatrix}}$$

$$\frac{\left(u : \begin{bmatrix} x : e \\ \mathbf{kingOf}(x, fr) \end{bmatrix}\right)}{\left(u : \begin{bmatrix} x : e \\ \mathbf{kingOf}(x, fr) \end{bmatrix}\right)} : \text{type}$$

$$\frac{\left(u : \begin{bmatrix} x : e \\ \mathbf{kingOf}(x, fr) \end{bmatrix}\right)}{\mathbf{wise}(\pi_{1}v)} : \text{type}$$

(707) @-elimination of (704)

$$\frac{\vdots}{\left[\begin{array}{c} x : \mathbf{e} \\ \mathbf{kingOf}(x, fr) \end{array} \right] : \mathbf{type}} \frac{\mathbf{wise} : \mathbf{e} \to \mathbf{type}}{\mathbf{wise}(\pi_{1}u) : \mathbf{type}} \stackrel{(CON)}{\underbrace{\left[\begin{array}{c} x : \mathbf{e} \\ \mathbf{kingOf}(x, fr) \end{array} \right]^{1}}}{\mathbf{m}_{1}u : \mathbf{e}} \stackrel{(DE)}{\underbrace{\left[\begin{array}{c} x : \mathbf{e} \\ \mathbf{kingOf}(x, fr) \end{array} \right]}} \\ \underbrace{\left[\begin{array}{c} x : \mathbf{e} \\ \mathbf{kingOf}(x, fr) \end{array} \right] : \mathbf{type}}_{(DE)}$$

Type checking algorithm returns a fully-specified semantic representation. Presupposition filtering is performed via exactly the same process as anaphora resolution.

Proposition 464 (A takeaway message from DTS on Presupposition Filtering). A presupposition is filtered in the same way as an anaphoric expression is resolved.

26.5.2 On Projection

The projection of presupposition is naturally accounted for using DTS. Consider how to derive the presupposition projected out of negation. Recall that negation is defined to be an implication of the form $\neg A \equiv A \to \bot$ where \bot is the absurdity type, i.e., the type that has no inhabitants. Given the formation rule for the bottom type shown on the left below, the formation rule for negation can be derived as on the right:

$$\frac{A : \text{type}}{\bot : \text{type}} \ (\bot F) \qquad \frac{A : \text{type}}{\neg A : \text{type}} \ (\neg F)$$

We analyze the definite article the as having the CCG category NP^*/N and the semantic representation in (708).

(708)

The underspecified term $@_1$ is a function that takes a local context c as argument. A term of the form $@_i\Lambda$ is called *underspecified term* and specifies that the term $@_i$

has type Λ . In the case of (708), the term $@_i$ is annotated with a Σ -type $(x:e) \times nx$. This means that the underspecified term $@_i$ is a term having the Σ -type. In this case, such a term is a pair of an entity x and a proof that x satisfies the condition n. Then its first projection, i.e., an entity x, is applied to a given predicate p. The annotated Σ -type here means that the underspecified term $@_i$ introduced by the definite article the requires a pair of an entity x and a proof that x satisfies the condition n provided by the restrict of the category N, given a local context c. the type $(x:\mathbf{E}) \times \mathbf{door}(x)$ from a given local context c. Such a pair consists of some entity x and a proof that x is a door, Then its first projection, i.e., an entity x, is applied to a given predicate. This means that for the entire sentence containing an expression of the form the N to be well-typed, one needs to construct a proof of the existence of such an N. Intuitively, this captures the existence presupposition triggered by a definite description.

According to the formation rule (ΠF) , the proposition A and its negation $\neg A$ have the same presupposition.

(709) It is not the case that the king of France is bald.

(710)
$$\neg \left[\begin{array}{c} @ \left[\begin{array}{c} x : e \\ \mathbf{king}(x, fr) \end{array} \right] \\ \mathbf{bold}(\pi_1 u) \end{array} \right]$$

The anaphora/presupposition resolution for the semantic representation (SR) A is triggered by the judgement $\vdash A$: type. This means that the presupposition resolution is amount to proving that the SR in question is well-typed given the signature.

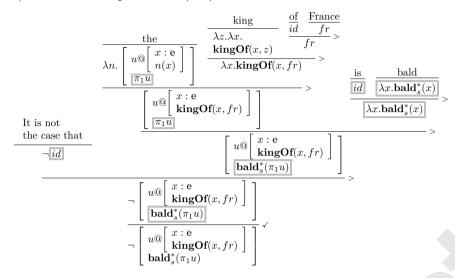
Assuming that wise: $e \to type$ is in the signature, the proof that the SR yielded by the sentence *The king of France is not wise* is well-typed runs as follows.*1

(711) Syntactic structure of (709)

$$\frac{\text{the}}{NP^{\star}/N} \stackrel{\text{king}}{=} \frac{\frac{\text{of}}{PP_{of}/NP} \stackrel{\text{France}}{NP}}{\frac{NP}{NP}} > \frac{\text{is}}{(S \backslash NP)(S \backslash NP))^{\star}} \stackrel{\text{bald}}{(S \backslash NP)^{\star}} > \frac{\text{the case that}}{\frac{S^{\star}/S}{S}} > \frac{\frac{S^{\star}}{S} \checkmark}{\frac{S^{\star}}{S}} > \frac{\frac{S^{\star}}{S} \checkmark}{\frac{S^{\star}}{S}} > \frac{\frac{S^{\star}}{S} \checkmark}{\frac{S^{\star}}{S} \checkmark}$$

^{*1} To be explicit, we the sequent-style of natural deduction, where the context on which each step in a derivation depends is indicated on the left-hand side of ⊢.

(712) Semantic composition of (709)



(713) Type check diagram of (709)

$$\frac{\left[\begin{array}{c} x : \mathbf{e} \\ \mathbf{kingOf}(x,fr) \end{array}\right] : \mathbf{type} \quad ? : \left[\begin{array}{c} x : \mathbf{e} \\ \mathbf{kingOf}(x,fr) \end{array}\right] \quad \mathbf{bald}_{s}^{*}(\pi_{1}u) ? / u] : \mathbf{type}}{\left[\begin{array}{c} u^{@} \left[\begin{array}{c} x : \mathbf{e} \\ \mathbf{kingOf}(x,fr) \end{array}\right] \right] : \mathbf{type}} \\ \frac{\left[\begin{array}{c} u^{@} \left[\begin{array}{c} x : \mathbf{e} \\ \mathbf{kingOf}(x,fr) \end{array}\right] \right] : \mathbf{type}}{\neg \left[\begin{array}{c} u^{@} \left[\begin{array}{c} x : \mathbf{e} \\ \mathbf{kingOf}(x,fr) \end{array}\right] \right] : \mathbf{type}} \\ \neg \left[\begin{array}{c} u^{@} \left[\begin{array}{c} x : \mathbf{e} \\ \mathbf{kingOf}(x,fr) \end{array}\right] \right] : \mathbf{type} \\ \neg \left[\begin{array}{c} u^{@} \left[\begin{array}{c} x : \mathbf{e} \\ \mathbf{kingOf}(x,fr) \end{array}\right] \right] : \mathbf{type} \\ \neg \left[\begin{array}{c} u^{@} \left[\begin{array}{c} x : \mathbf{e} \\ \mathbf{kingOf}(x,fr) \end{array}\right] \right] : \mathbf{type} \\ \neg \left[\begin{array}{c} u^{@} \left[\begin{array}{c} x : \mathbf{e} \\ \mathbf{kingOf}(x,fr) \end{array}\right] \right] : \mathbf{type} \\ \neg \left[\begin{array}{c} u^{@} \left[\begin{array}{c} x : \mathbf{e} \\ \mathbf{kingOf}(x,fr) \end{array}\right] \right] : \mathbf{type} \\ \neg \left[\begin{array}{c} u^{@} \left[\begin{array}{c} x : \mathbf{e} \\ \mathbf{kingOf}(x,fr) \end{array}\right] \right] : \mathbf{type} \\ \neg \left[\begin{array}{c} u^{@} \left[\begin{array}{c} x : \mathbf{e} \\ \mathbf{kingOf}(x,fr) \end{array}\right] \right] : \mathbf{type} \\ \neg \left[\begin{array}{c} u^{@} \left[\begin{array}{c} x : \mathbf{e} \\ \mathbf{kingOf}(x,fr) \end{array}\right] \right] : \mathbf{type} \\ \neg \left[\begin{array}{c} u^{@} \left[\begin{array}{c} x : \mathbf{e} \\ \mathbf{kingOf}(x,fr) \end{array}\right] \right] : \mathbf{type} \\ \neg \left[\begin{array}{c} u^{@} \left[\begin{array}{c} x : \mathbf{e} \\ \mathbf{kingOf}(x,fr) \end{array}\right] \right] : \mathbf{type} \\ \neg \left[\begin{array}{c} u^{@} \left[\begin{array}{c} x : \mathbf{e} \\ \mathbf{kingOf}(x,fr) \end{array}\right] \right] : \mathbf{type} \\ \neg \left[\begin{array}{c} u^{@} \left[\begin{array}{c} x : \mathbf{e} \\ \mathbf{kingOf}(x,fr) \end{array}\right] \right] : \mathbf{type} \\ \neg \left[\begin{array}{c} u^{@} \left[\begin{array}{c} x : \mathbf{e} \\ \mathbf{kingOf}(x,fr) \end{array}\right] \right] : \mathbf{type} \\ \neg \left[\begin{array}{c} u^{@} \left[\begin{array}{c} x : \mathbf{e} \\ \mathbf{kingOf}(x,fr) \end{array}\right] \right] : \mathbf{type} \\ \neg \left[\begin{array}{c} u^{@} \left[\begin{array}{c} x : \mathbf{e} \\ \mathbf{kingOf}(x,fr) \end{array}\right] \right] : \mathbf{type} \\ \neg \left[\begin{array}{c} u^{@} \left[\begin{array}{c} x : \mathbf{e} \\ \mathbf{kingOf}(x,fr) \end{array}\right] \right] : \mathbf{type} \\ \neg \left[\begin{array}{c} u^{@} \left[\begin{array}{c} x : \mathbf{e} \\ \mathbf{kingOf}(x,fr) \end{array}\right] \right] : \mathbf{type} \\ \neg \left[\begin{array}{c} u^{@} \left[\begin{array}{c} x : \mathbf{e} \\ \mathbf{kingOf}(x,fr) \end{array}\right] \right] : \mathbf{type} \\ \neg \left[\begin{array}{c} u^{@} \left[\begin{array}{c} x : \mathbf{e} \\ \mathbf{kingOf}(x,fr) \end{array}\right] \right] : \mathbf{type} \\ \neg \left[\begin{array}{c} u^{@} \left[\begin{array}{c} x : \mathbf{e} \\ \mathbf{kingOf}(x,fr) \end{array}\right] \right] : \mathbf{type} \\ \neg \left[\begin{array}{c} u^{@} \left[\begin{array}{c} u^{@} \left[\begin{array}{c} x : \mathbf{e} \\ \mathbf{kingOf}(x,fr) \end{array}\right] \right] : \mathbf{type} \\ \neg \left[\begin{array}{c} u^{@} \left[\begin{array}{c} u^{@} \left[\begin{array}{c} u : \mathbf{e} \\ \mathbf{kingOf}(x,fr) \end{array}\right] \right] : \mathbf{type} \\ \neg \left[\begin{array}{c} u^{@} \left[\begin{array}{c} u : \mathbf{e} \\ \mathbf{kingOf}(x,fr) \end{array}\right] \right] : \mathbf{type} \\ \neg \left[\begin{array}{c} u^{@} \left[\begin{array}{c} u : \mathbf{e} \\ \mathbf{kingOf}(x,fr) \end{array}\right] \right] : \mathbf{type} \\ \neg \left[\begin{array}{c} u^{@} \left[\begin{array}{c} u : \mathbf{e} \\ \mathbf{kingOf}(x,$$

In order for the sentence "The king of France is bald" to be well-formed, the context Γ must be such that the following type inhabits a proof (namely, there exists a king of France).

(714)
$$\Gamma \vdash ? : \left[\begin{array}{c} x : \mathbf{e} \\ \mathbf{kingOf}(x, fr) \end{array} \right]$$

The same inference is triggered for the antecedent of a conditional sentence like (715).

- (715) If the king of France is wise, people will be happy.
- (716) Type check diagram of (715)

$$\frac{\begin{bmatrix} u @ \begin{bmatrix} x : e \\ \mathbf{kingOf}(x, fr) \end{bmatrix}}{\mathbf{wise}(\pi_{1}u)} : \text{type } \mathbf{happy}(people) : \text{type} \\ \frac{u @ \begin{bmatrix} x : e \\ \mathbf{kingOf}(x, fr) \end{bmatrix}}{\mathbf{wise}(\pi_{1}u)} \rightarrow \mathbf{happy}(people) : \text{type} \end{bmatrix}$$

What is presupposed by the original sentence in (??b) can be read off from the open branch ending with the judgment having the underspecified term $@_1$. The context on which this judgement depends can be made explicit in the following way. For the given SR to be well-typed, (which means that the utterance of (??b) is felicitous), one has to find a term by which (717) holds.

(717)
$$\left[\begin{array}{c} x : \mathbf{e} \\ \mathbf{kof}(x) \end{array} \right] true$$

That is to say, one has to find a proof term for the proposition that there is a king of France. In this way, we can derive the existence presupposition for the negated sentence (??b), as well as for the positive counterpart (??a). As is easily seen by the definition $\neg A \equiv A \to \bot$, the same inference is triggered for the antecedent of a conditional sentence in (??c). Thus we can also account for presupposition projection for conditionals as exemplified in (??c).

Proposition 465 (A takeaway message from DTS on Presupposition Projection). A presupposition projects because it's truth is a requirement for a sentence containing it to be *semantically well-formed*, not to be *true*.

Corollary 466. Existence of an antecedent is a requirement for a sentence containing anaphora to be semantically well-formed, not to be true.

26.6 Accommodation

Consider (718), taken from Heim (1983).*2

(718) No nation cherishes its king.

The sentence (718) is interpreted in two distinct ways regarding its presupposition, as described in Heim (1983, p. 403) as follows:

What about quantifiers other than universal? Concerning "no," we find conflicting factual claims in the literature. According to Cooper (1983), (24) [=our (718)] should presuppose that every nation (in the relevant domain of discourse) has a king; for ?, it presuppose that some nation does.

One reading of (718) that implies every nation has a king is obtained via the *global accommodation*, whereas another reading of (718) implying Some nation has a king is obtained via the *local accommodation*.

We assume that the lexical item of its is defined as (719) in almost the same way as his in (??) except for its presuppositional content:

(719)
$$[its \vdash NP^*/N] \stackrel{def}{=} \lambda n. \begin{bmatrix} u@ \begin{bmatrix} x : e \\ \neg \mathbf{human}(x) \end{bmatrix} \\ v@ \begin{bmatrix} y : e \\ n(y) \\ \mathbf{of}(y, \pi_1 u) \end{bmatrix} \end{bmatrix}$$

^{*2} Sentence (24) on p. 403.

Since the syntactic structure of (718) is almost the same as (??), we obtain its semantic representation via the homomorphic semantic composition, which is (720) corresponding to the BVA reading for (718).

$$(720) \quad (w:(x:\mathbf{e})\times\mathbf{nation}(x)) \to \neg \left[\begin{array}{l} u@(x:\mathbf{e})\times\neg\mathbf{human}(x) \\ v@(y:\mathbf{e})\times\mathbf{king}(y)\times\mathbf{of}(y,\pi_1u) \\ \mathbf{cherish}(\pi_1v,\pi_1w) \end{array} \right]$$

The type checking for (720) factors through (721).

$$(721) \quad \Gamma, \ w:(x:\mathbf{e}) \times \mathbf{nation}(x) \vdash \left[\begin{array}{l} u@(x:\mathbf{e}) \times \neg \mathbf{human}(x) \\ v@(y:\mathbf{e}) \times \mathbf{king}(y) \times \mathbf{of}(y,\pi_1 u) \\ \mathbf{cherish}(\pi_1 v, \pi_1 w) \end{array} \right] : \mathbf{type}$$

which will invoke the @-rule (??) and call for the type checking (722a), which obviously succeeds, and then the proof search (722b).

(722) a.
$$\Gamma$$
, $w:(x:e) \times \mathbf{nation}(x) \vdash (x:e) \times \neg \mathbf{human}(x) : \text{type}$
b. Γ , $w:(x:e) \times \mathbf{nation}(x) \vdash ?:(x:e) \times \neg \mathbf{human}(x)$

If we assume world knowledge that every nation is non-human (i.e., the constant term **nnh**: $(x : e) \to \mathbf{nation}(x) \to \neg \mathbf{human}(x)$), the proof search for (722b) will find the term $(\pi_1 w, \mathbf{nnh}(\pi_1 w)(\pi_2 w))$ that corresponds to the BVA reading of (718). The next step is the type checking (723), which breaks down to the type checking (724a) and the proof search (724b).

(723)
$$\Gamma$$
, $w:(x:e) \times \mathbf{nation}(x) \vdash \begin{bmatrix} v@(y:e) \times \mathbf{king}(y) \times \mathbf{of}(y, \pi_1 w) \\ \mathbf{cherish}(\pi_1 v, \pi_1 w) \end{bmatrix}$: type

(724) a.
$$\Gamma$$
, $w:(x:e) \times \mathbf{nation}(x) \vdash (y:e) \times \mathbf{king}(y) \times \mathbf{of}(y, \pi_1 w)$: type
b. Γ , $w:(x:e) \times \mathbf{nation}(x) \vdash ?:(y:e) \times \mathbf{king}(y) \times \mathbf{of}(y, \pi_1 w)$

Again, (724a) is a routine, and the proof search for (724b) is a search for the king of y. Remember that, in the case of (??) and (??), we assumed that there exists a world knowledge **fatherOf**, that is, every boy has a father. The question here is what happens if we do not have a knowledge that every nation has a king, and thus the proof searching for (724b) fails.

The first strategy for our type checker is global accommodation, defined in Bekki (2014), which can be restated as the following instruction.

Definition 467 (Global accommodation). When the proof search

$$\llbracket \Gamma \vdash ? : A \rrbracket$$

fails for the type A such that $fv(A) = x_1, \ldots, x_n$, the type checker may add the constant term of type $(x_1 : A_1) \to \cdots \to (x_n : A_n) \to A$ (where $x_1 : A_1, \ldots, x_n : A_n \subseteq \Gamma$) to the signature and re-run the type checking.

Applying Definition 467 to the case of (724b), the constant term added to the signature is $(w : (x : e) \times \mathbf{nation}(x)) \to (y : e) \times \mathbf{king}(y) \times \mathbf{of}(y, \pi_1 w)$, which is exactly the knowledge that every nation has a king.

The second strategy for our type checker is local accommodation, which, under the definition of underspecified types, is given the following definition.

Definition 468 (Local accommodation). When the proof search

$$\llbracket \Gamma \vdash ? : A' \rrbracket$$

required for the type checking $\llbracket \Gamma \vdash (x@A) \times B : \mathbf{type} \rrbracket$ fails, the type checker may replace the result of $\llbracket \Gamma \vdash (x@A) \times B : \mathbf{type} \rrbracket$ with the result of type checking $\llbracket \Gamma \vdash (x : A) \times B : \mathbf{type} \rrbracket$.*3

One benefit of employing the underspecified types is that it allows us to define the operation of local accommodation in a straightforward way. This is due to their definition that takes a closer form to that of the Σ -formation rule: when the proof search $\llbracket\Gamma \vdash ?:A\rrbracket$ fails in the type checking $\llbracket\Gamma \vdash (x@A) \times B: \text{type}\rrbracket$, the semantic system may optionally replace $(x@A) \times B$ with $(x:A) \times B$ and re-run the type checking. Note that this transformation is safe owing to the similarity between the verification conditions of the @-rule and Σ -formation rule.

In other words, the local accommodation in DTS is understood as accommodating the presupposed content of an underspecified type by existentially quantifying it. Applying Definition 468 to (724b), replacing the type checking for (723) with the following:

(725)
$$\Gamma$$
, $w:(x:e) \times \mathbf{nation}(x) \vdash \begin{bmatrix} v:(y:e) \times \mathbf{king}(y) \times \mathbf{of}(y, \pi_1 w) \\ \mathbf{cherish}(\pi_1 v, \pi_1 w) \end{bmatrix}$: type

thus the whole semantic representation for (718) turns into (726).

(726)
$$\Gamma \vdash (w : (x : e) \times \mathbf{nation}(x)) \rightarrow \neg \begin{bmatrix} v : (y : e) \times \mathbf{king}(y) \times \mathbf{of}(y, \pi_1 w) \\ \mathbf{cherish}(\pi_1 v, \pi_1 w) \end{bmatrix} :$$
type

In words, (726) claims that No nation has a king and cherishes it, which exactly corresponds to the locally accommodated reading of (718).

26.7 Cancellation of Presuppositions

26.8 Discussion

History and Further Readings

^{*3} This operation should not be applied to all types of anaphora and presupposition: for example, it is known that pronouns in general do not undergo local accommodation when their antecedents are missing. For this purpose, we may want to add a binary feature to underspecified types that tells us whether local accommodation is applicable to them, an approach pursued by Yana et al. (2024). Alternatively, we may argue that the prohibition of local accommodation for pronouns is based rather on pragmatic factors. This is a controversial issue and I would like to leave it open.