



ECE 458

Assignment 2



JULY 5, 2021

Question 1

For question 1, our implementation goes through every secret key and returns the public key.

```
def question1(g, p, sks):  
    for sk in sks:  
        #goes through all secret keys and generates public key  
        pki = generate_public_key(g, sk, p)  
        pk.append(pki)  
        #print(sk, pki)  
    #print(pk)
```

In our generate_public_key method we use the pow function which allows us to replicate $g^{sk} \bmod p$ in python. We used this function as it is computationally less intense than actually taking powers and mods in python.

```
def generate_public_key(g, sk, p):  
    #print("GENRATING PUBLIC KEY")  
    #pow does gives us  $g^{sk} \% p$   
    output = pow(g, sk, p)  
    return output
```

We got the following results:

pairs of (ski, pki) for i = 1, 2, 3:

Sk1: 432398415306986194693973996870836079581453988813,

pk1:

49336018324808093534733548840411752485726058527829630668967480568854756416567496216
29491905191014868618662270686970232166446509470324736864650682101529030248099045013
02806169292269172462551470632923017242976806834012586361821855991241311700775484507
54294083728885075516985144944984920010138492897272069257160

Sk2: 165849943586922055423650237226339279137759546603,

pk2:

15805870909801658055470664266537756807095291079892225065940408602107787041560306349
91442689557657923705109503128714241617779465259782688524304338594046010559983239609

11988506849369081405990373344476931374270404283587717926798415978556358354906886367
573116683374249583660883550584998781456103573121625443428385

Sk3: 627658512551971075308886219669315148725310346887,

Pk3:

10748734721713476921614873611718904745721634734864852672984965178578400403592975468
63771191907252113917147569678898094618007043330822196002059738924118429587340918403
87580674995858195843168769210392825540392494396460556417290804549088996408888054754
496901159723813244712885638670114254415946917214251683052844

Question 2

In question 2, we are converting the input into an int, and then to a hex value. Next, the leading zeros are stripped out and turned to bytes. The hashobj is created by using the *sha3_224()* method from the *hashlib* library. Finally, the hashed value is obtained by using the hashobj and returned in binary.

SHA3-224 for number 10:

10011110100001101111111011010010101010101111001010100101011110100000001011111000
01000000100100110100101101000010110110011100011100011101010000001100110110011000010
0111101110100101000010111101001000001010110110101010001011

SHA3-224 for number 20:

1011111110000110011101110010001001100011101010110101010000111010100100101111001001
11101101010110100100110000111111011111001100001001100111110110111100110111100101000
1111111011010010110011000000000011000011010110100010111101

SHA3-224 for number 30:

11101011111101101000100000100001110101101010000100010110010101111010101001100000000
10111101000100001110000111101001000001011101111000010011010000111010000101110110110
01010110000000001111001110000010100011001111011000001111101

Question 3

For question 3 we generate the m1 and m2 using the needed primary keys for each message and the amnt value. We pass these values in from our array into the question3 function.

```
elif selected == '3':  
    #gets messages for both m1 and m2. Question 1 populates primary key array  
    question1(g,p,sks)  
    print("M1")  
    question3(pk[0], pk[1], 2)
```

```
print("M2")
question3(pk[1], pk[2], 3)
```

The question3 function calls the generate message function.

```
def question3(pk1, pk2, amnt):
    #calls the generate message function to generate the concatenated message
    print(generate_message(pk1,pk2,amnt))
```

The generate_message function gets the binary representation of all inputs, concatenates together and then returns an integer.

```
def generate_message(pk1, pk2, amnt):
    #take the primary keys, convert to binary keeping significant digits of 399, and then add the amnt to it. We then
    #convert back into an integer
    pk1 = bin(pk1)[2:401]
    pk2 = bin(pk2)[2:401]
    amnt = bin(amnt)[2:] if amnt != 1 else "01" #special case for 1.

    m1 = int(pk1 + pk2 + amnt, 2)
    return m1
```

Thus we get our messages. Below is our outputs:

M1:

```
36599492524161102438426956351931861753323140623812778402756972807061815843875480980
13985066122238476554562154647781507975854150384723431517763123752624375814447084742
521086657783968078531722961438847114809276456296127939875396313500323209774
```

M2:

```
58627233595607699023061607981987080781902518631747317209768530288285875986900794737
40178204613802974429824697161497463490301926377746107083836495644583074647011520782
397841059567377684346054784824548991914382152504386907016111662295782911303
```

Question 4

For question 4 we must both generate a digital signature and then verify it.

We get the signature pair using the method given for digital signatures. So we first find a random k in between 0 and p . We then take the message and hash it. From there we determine r and s using the respective inputs needed to determine their value.

```

def generate_signature(message, p, q, g, sk):
    #get random k in between 0 and p
    k = randint(0, p-1)

    #convert message into hex and then bytes
    integer = hex(message)
    integer = integer[2:]
    integer = bytes(integer, encoding='utf-8')

    #hash the message
    hashobj.update(integer)
    hashed = str(hashobj.hexdigest())
    #convert from hex string to int
    hashed_num = int(hashed, 16) #this is our h(m)

    y = pow(g, sk, p) # public key
    #print('public key ', y)

    r = pow(g, k, p) #first integer of signature pair

    # hashed_num = (sk1 * r + k * s) % q
    s = ((hashed_num - sk * r) * mod_inverse(k, q)) % q #second half of signature pair

    #print((hashed_num - sk * r) * mod_inverse(k, q))
    print('r ', r)
    print('s ', s)
    #send signature pair and public key abck
    return hashed_num, r, s, y

```

This is how we can get Sig1 and Sig2.

We verify by calling the `verify_signature` method. We first check to make sure that the parameters fall in our expected ranges. If it is we go onto verifying it by back calculating an `r` value(`w`) and comparing it to the `r` value we found in our initial signature generation. That way we know it is working. We do this by using the method discussed in the lecture notes.

```
def verify_signature(hashd_message, r, s, g, p, q, pk):  
    #check if values fall in expected rangers  
    if r < 0 or r > p or s < 0 or s > p:  
        print("verification rejected ")  
        return False  
  
    s1 = mod_inverse(s, q)  
    u = hashd_message * mod_inverse(s, q) % q  # u = hashd_num * s inverse mod q,  
    v = (-r * s1) % q  # and v = -r * s inverse mod q  
  
    arg1 = pow(g, u, p) #g ^ u mod p  
    arg2 = pow(pk, v, p) #pk ^ v mod p  
  
    w = arg1 * arg2  
    w = pow(w, 1, p) # g^u y^v mod p  
    #this value is our generated r  
    print("w", w)  
  
    #if the r is the same we can verify that we generated a valid signature  
    if r == w:  
        print("verification succesful ")  
        return True
```

This is our output for the question.

Output

```

DaivikMac:Assignment 2 daivikgoel$ python3 assignment2.py
WHICH QUESTION DO YOU WANT: 4
r 810589125034580997099852498362216118417379205273175690438538889284414086982589954
986254662771404106957874315758119966017050449656555561538146520728206922553162171656
097046443968765626237342479373744818062848943408402330146411598857450760811471223162
96371162777820868479779412405529957693476205858345707057298
s 82312537251777474857406676638257961685160321020
w 8105891250345809970998524983622161184173792052731756904385388892844140869825899549
862546627714041069578743157581199660170504496565555615381465207282069225531621716560
970464439687656262373424793737448180628489434084023301464115988574507608114712231629
6371162777820868479779412405529957693476205858345707057298
verification succesful
r 115981462476884199700876396428684859792112075633525002670666855814338887454096389
906776852598820252800120147217242565372530136539194514517329317404599330240464474409
939216097532918199328733951954168650826976898839799716232874206186807558435369756590
845271784511671833379279608277153701556839005025112976455530
s 668174398700767433498639110966695275529316861058

```

Question 5

For this question we want to find a nonce value that gives us 24 leading zeroes. First we need to create a message and hashed message of the previous for our Ti. Hence first we go through every message and generate a hashed equivalent (Note: Since T0 would use the hash of amnt0 our hashed messages array has that value first. This creates an offset of 1.

```

def find_nonce():
    found = 0

    NONCE = 0

    question1(g,p,sks) #generate the public keys

    messages = [] #save m1, m2, m3 etc
    hashed_messages = [question2(amnt[0])] #Offset by 1 as h(m0) is h(amnt0)

    for i in range(len(pk)-1):
        #create a message using primary keys
        messages.append(generate_message(pk[i], pk[i+1], amnt[i+1]))
        #create hashed message using these messages
        hashed_messages.append(question2(generate_message(pk[i], pk[i+1], amnt[i+1])))
    nonces = []

```

From there we start a clock to determine the time. Our function will get both nonce 1 and 2 and keep going till we find a value with 24 leading zeroes. We add a nonce value, our mi and h(m(i-1)) and hash it.

If this value has the 24 leading zeroes we stop else we keep iterating the nonce value by 1. We figure out if it has 24 leading zeroes if the hash values's length is 200. This is because python will not show leading zeroes so we know should the length be 200, we had 24 leading zeroes.

```
start_time = datetime.datetime.now()

for i in range(2):
    #find nonce1 and nonce2
    found = 0
    NONCE = 0
    print(i)
    #for y in [2,4,6,8,10,12,14,16,18,20,22]: #this is for testing purposes on effort required
        #found = 0
    while found == 0:
        #because the array is offset by 1 in hashedmessages it represents the equation  $\text{NONCE} + m_i + h(m_{i-1})$ 
        z = NONCE + messages[i] + int(hashed_messages[i])
        #find hash of the concatenation.

        new_hash = question2(z) #new_hash is T_i
        # since the first 24 bits aren't shown in binary representation, when the length of the new hash is exactly 24
        #bits short then we know we had 24 leading zeroes
        if 224 - len(new_hash) == 24:
            #if we have 24 leading zeroes we have found a nonce
            found = 1
            end_time = datetime.datetime.now()

            print("elapsed time for finding nonce: ", (end_time - start_time))
            nonces.append(NONCE)
            NONCE += 1

    #show both nonces
    print(nonces)
```

With this method we got the following nonce values:


```
DaivikMac:Assignment 2 daivikgoel$ python3 assignment2.py
WHICH QUESTION DO YOU WANT: 5
0
elapsed time for finding nonce: 0:04:03.778100
1
elapsed time for finding nonce: 0:09:34.984521
[46488426, 63390458]
```

Question 7

Number of exhaustive searches required for varying nonce in order to get hash value that satisfies $k=24$ leading zeros.

Since the hash function is collision resistant, we have no way of knowing in advance what the hashed output would be prior to trying out all the input values. Thus, we would need to brute force at most 2^k times to ensure that we land on a hashed value with k leading zeros.

For $k = 24$, the number of iterations taken to find the nonce 1 was 46488426 while the number of iterations taken to find nonce 2 was 63390458. Their respective times taken being 0:04:03 and 0:09:34.

effort needed if we wanted to find each nonce with 32 leading zeros:

The effort to find each nonce with 32 leading zeros would be 2^{32} . ($2^{32} / 2^{24} = 2^8 = 256$). Thus, it takes roughly 256 times more effort to generate nonce with 32 leading zeros compared to generating the nonce with 24 leading zeros. To further prove this, the figure below was made by plotting the values we received for increasing values of k , starting from $k = 2$ we can see that for increasing values of k , nonce increases exponentially.

