

Iterative Methods

Splitting

Let $A = M - N$ $M \approx A$

Then $Ax = b$

$$(M - N)x = b$$

$$Mx = Nx + b$$

Define an iterative method by:

$$Mx^{k+1} = Nx^k + b$$

Then

$$\begin{aligned} x^{k+1} &= M^{-1}Nx^k + M^{-1}b \\ &= M^{-1}(M - A)x^k + M^{-1}b \\ &= x^k + M^{-1}(b - Ax^k) \end{aligned}$$

Note: If $M = A$, then $x^{k+1} = x^k + A^{-1}(b - Ax^k)$

$$= x^k + x - x^k = x$$

→ one step convergence

But one needs to compute $A^{-1}(b - Ax^k)$

Goals:

(1) $M \approx A$

(2) M^{-1} is easy to compute

Richardson iteration

- $M = 1/\theta I$ (θ is appropriately chosen)

Then $M^{-1} = \theta I$

Thus $x^{k+1} = x^k + \theta I (b - A x^k)$

Consider the i -th entry of x^{k+1} :

$$x_i^{k+1} = x_i^k + \theta(b_i - \sum_{j=1}^n a_{i,j} x_j^k)$$

Algorithm

x^0 = initial guess

for $k = 0, 1, 2, \dots$

for $i = 1, 2, \dots, n$

$$x_i^{k+1} = x_i^k + \theta(b_i - \sum_{j=1}^n a_{i,j} x_j^k)$$

end

end

Note

Need 2 separate vectors x^k, x^{k+1} .

Jacobi iteration

$$\bullet M = D = \text{diagonal of } A = \begin{bmatrix} a_{1,1} & & \\ & \ddots & \\ & & a_{n,n} \end{bmatrix}$$

$$\begin{aligned} \text{Thus } x^{k+1} &= x^k + D^{-1}(b - Ax^k) \\ x_i^{k+1} &= x_i^k + \frac{1}{a_{i,i}}(b_i - \sum_{j=1}^n a_{i,j}x_j^k) \\ &= x_i^k + \frac{1}{a_{i,i}}(b_i - \sum_{j \neq i} a_{i,j}x_j^k - a_{i,i}x_i^k) \\ &= \frac{1}{a_{i,i}}(b_i - \sum_{j \neq i} a_{i,j}x_j^k) \end{aligned}$$

Interpretation

Let $r^k = b - Ax^k$ (residual vector of x^k)

Then $x^k = x \Leftrightarrow r^k = 0$

Thus $\|r^k\|_2 \approx 0 \Rightarrow x^k \approx x$

$$\text{Consider } r_i^k = b_i - \sum_{j=1}^n a_{i,j}x_j^k = b_i - \sum_{j \neq i} a_{i,j}x_j^k - a_{i,i}x_i^k$$

In general, $r_i^k \neq 0$. Now modify x_i^k so that $r_i^k = 0$.

i.e.

$$b_i - \sum_{j \neq i} a_{i,j}x_j^k - a_{i,i}x_i^{k+1} = 0$$

$$x_i^{k+1} = \frac{1}{a_{i,i}}(b_i - \sum_{j \neq i} a_{i,j}x_j^k)$$

Gauss-Seidel iteration

Let $A = D - L - U$, where D = diagonal of A , L = strictly lower Δ part, U = strictly upper Δ part

i.e.

$$A = \begin{bmatrix} \ddots & & -U \\ & D & \\ -L & & \ddots \end{bmatrix}$$

- $M = D - L$ = lower Δ part of A .

i.e. $x^{k+1} = x^k + (D-L)^{-1} (b - A x^k)$

Interpretation

Modify x_i^k so that $r_i^k = 0$. Use the new x_j^{k+1} , $j < i$, from the previous updates.

i.e.
$$r_i^k = b_i - \sum_{j < i} a_{i,j} x_j^{k+1} - a_{i,i} x_i^{k+1} - \sum_{j > i} a_{i,j} x_j^k = 0$$

Thus
$$x_i^{k+1} = \frac{1}{a_{i,i}} (b_i - \sum_{j < i} a_{i,j} x_j^{k+1} - \sum_{j > i} a_{i,j} x_j^k)$$

Backward GS

- $M = D - U$: $x^{k+1} = x^k + (D-U)^{-1} (b - A x^k)$

Symmetric GS

A forward sweep followed by a backward sweep:

SOR iteration

- Successive over-relaxation
- Weighted average of x_i^k and $x_i^{k+1} = \text{GS}(x_i^k)$:

$$x_i^{k+1} = (1 - \omega)x_i^k + \frac{\omega}{a_{i,i}}(b_i - \sum_{j < i} a_{i,j}x_j^{k+1} - \sum_{j > i} a_{i,j}x_j^k)$$

$$\text{SOR} = \begin{cases} \text{GS} & \omega = 1 \\ \text{Under-relaxation} & \omega < 1 \\ \text{Over-relaxation} & \omega > 1 \end{cases}$$

- $M = 1/\omega D - L$
- For a suitably chosen $\omega (>1)$, SOR can be much better than GS.

Convergence Analysis

Q1: Under what condition does the iteration converge?

Q2: If the iteration converges, how fast is it?

Def: λ is an eigenvalue and v an eigenvector of A if

$$A v = \lambda v \quad v \neq 0$$

Def: The spectral radius of A , $\rho(A) = \{|\lambda| : \lambda \text{ eig of } A\}$, is the largest absolute value of the eigenvalues of A .

Theorem: The iterative method:

$$x^{k+1} = x^k + M^{-1} (b - A x^k)$$

is convergent for any x^0 and b if and only if

$$\rho(I - M^{-1} A) < 1$$

- $I - M^{-1} A$ is called the iteration matrix
- $\rho(I - M^{-1} A)$ is called the rate of convergence

Minimization formulation

Assume A is SPD. Consider the functional:

$$F(x) \equiv \frac{1}{2} x^T A x - b^T x \quad x \in \mathbb{R}^n$$

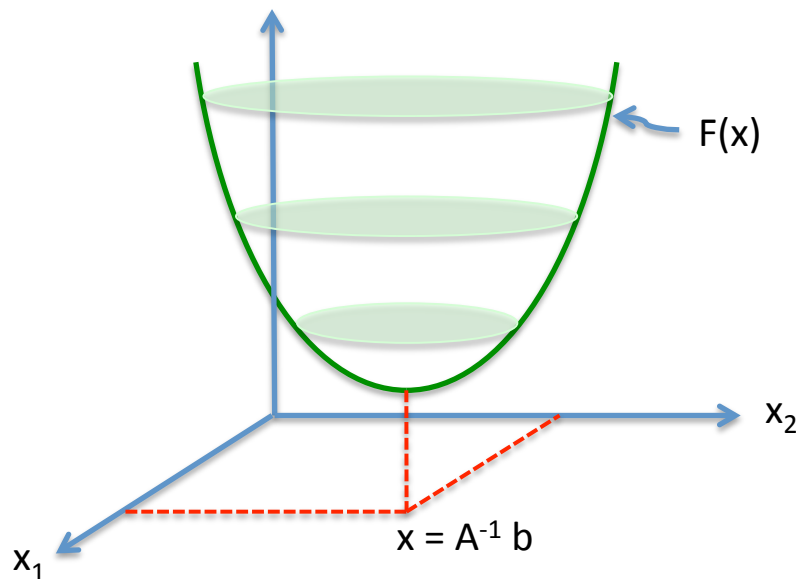
Theorem: The solution of $A x = b$ is equivalent to the solution of the minimization problem:

$$\min_x F(x)$$

Pf: The minimizer satisfies $F'(x) = 0$ i.e. $\frac{\partial F}{\partial x_k} = 0$ and $F'(x) = Ax - b$

Note: $F(x)$ is convex. So a local min = global min.

e.g. $n = 2$



Search directions

Idea: $\min F(x)$ along a direction p ($p \neq 0$)

Let x^k = current approx. Define $x^{k+1} = x^k + \alpha p$

Determine α by $\min F(x^{k+1})$ along p

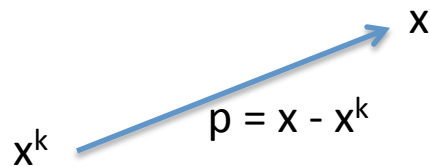
$$\begin{aligned}\text{Let } f(\alpha) &\equiv F(x^k + \alpha p) \\ &= \frac{1}{2}(x^k + \alpha p)^T A(x^k + \alpha p) - (x^k + \alpha p)^T b \\ &= \frac{1}{2}(x^k)^T A x^k + \frac{\alpha}{2} p^T A x^k + \frac{\alpha}{2} (x^k)^T A p + \frac{\alpha^2}{2} p^T A p \\ &\quad - (x^k)^T b - \alpha p^T b \\ &= \frac{1}{2}(x^k)^T A x^k - (x^k)^T b + \alpha p^T A x^k - \alpha p^T b + \frac{\alpha^2}{2} p^T A p \\ &= F(x^k) - \alpha p^T (b - A x^k) + \frac{\alpha^2}{2} p^T A p\end{aligned}$$

$$0 = f'(\alpha) = -p^T (b - A x^k) + \alpha p^T A p$$

$$\Rightarrow \alpha = \frac{p^T (b - A x^k)}{p^T A p} = \frac{p^T r^k}{p^T A p}$$

Notes

- 1) A is SPD $\Rightarrow p^T A p > 0$
- 2) What is the optimal search direction p ?



However x is unknown.

Steepest descent method

- Local optimal direction

Consider $f'(\alpha) = p^T (A x^k - b) + \alpha p^T A p$

Then $f'(0) = p^T F'(x^k) \quad (F'(x) = A x - b)$
= changes in F at x^k in the direction of p

Idea: make $f'(0)$ as negative as possible by varying p

Assume $\|p\| = 1$. Then

$f'(0)$ max if $p = F'(x^k) / \|F'(x^k)\|$ = steepest ascent

$f'(0)$ min if $p = -F'(x^k) / \|F'(x^k)\|$ = steepest descent
 $= r^k / \|r^k\| \quad (F'(x) = -r)$

Steepest descent method:

$$x^{k+1} = x^k + \alpha^k r^k \quad \alpha^k = \text{step length}$$

The optimal $\alpha^k = (r^k)^T (r^k) / (r^k)^T A r^k \quad (\alpha = p^T r^k / p^T A p)$

Also

$$\begin{aligned}
 r^{k+1} &= b - A x^{k+1} \\
 &= b - A (x^k + \alpha_k r^k) \\
 &= b - A x^k - \alpha_k A r^k \\
 &= r^k - \alpha_k A r^k
 \end{aligned}$$

Algorithm

Given x^0 , compute $r^0 = b - A x^0$

for $k = 0, 1, 2, \dots$

$$\alpha_k = (r^k)^T (r^k) / (r^k)^T A r^k$$

$$x^{k+1} = x^k + \alpha_k r^k$$

$$r^{k+1} = r^k - \alpha_k A r^k$$

end

Notes

- 1) Only 1 matrix-vector product ($A r^k$) per iteration
- 2) “Nonlinear” iterative method:

$$x^{k+1} = x^k + \alpha_k (b - A x^k)$$

i.e. $M = M^k = 1/\alpha_k I$