

# THEORY OF INFERENCE

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## 1 Overview

## 2 Rules of Inference

Let  $A$  and  $B$  be two statement formula. We say that "  $B$  logically follows from  $A$ " or " $B$  is a valid conclusion of the premise  $A$ " iff  $A \rightarrow B$  is a Tautology i.e.,  $A \implies B$

To demonstrate that a particular formula is valid consequence of a given set of premises, we use the follow rules of inference.

*Rule P : A premise may be introduced at any point in the derivation*

*Rule T : A formula  $S$  may be introduced in a derivation if  $S$  is tautologically implied by any one or more of the preceding formulas in the derivation*

**Question :** Demonstrate that  $r$  is a valid inference from the premises

$p \rightarrow q$ ,  $q \rightarrow r$  and  $p$

**Solution :**

$p \rightarrow q$  (Rule P)

$p$  (Rule P)

$q$  (Rule T,  $p \wedge (p \rightarrow q) \implies q$ )

$q \rightarrow r$  (Rule P)

$r$  (Rule T,  $q \wedge (q \rightarrow r) \implies r$ )

**Question :** RVS follows logically from the premises  $C \wedge D$ ,  $C \vee D \rightarrow \neg H$ ,  $\neg H \rightarrow (A \wedge \neg B)$ ,  $(A \wedge \neg B) \rightarrow (R \vee S)$ .

**Solution :**

$$C \vee D \rightarrow \neg H \quad (\text{Rule } P)$$

$$\neg H \rightarrow A \wedge \neg B \quad (\text{Rule } P)$$

$$C \vee D \rightarrow A \wedge \neg B \quad (\text{Rule } T)$$

$$A \wedge \neg B \rightarrow R \vee S \quad (\text{Rule } P)$$

$$C \vee D \rightarrow R \vee S \quad (\text{Rule } T)$$

$$C \vee D \quad (\text{Rule } P)$$

$$R \vee S \quad (\text{Rule } T)$$

**Question :** Show that  $S \vee r$  is tautologically implied by  $(p \vee q) \wedge (p \rightarrow r) \wedge (q \rightarrow s)$

**Solution :**

$$\begin{array}{ll}
 p \vee q & (\text{Rule } P) \\
 \neg p \rightarrow q & (\text{Rule } T \text{ i.e., } p \rightarrow q \Leftrightarrow \neg p \vee q) \\
 q \rightarrow S & (\text{Rule } P) \\
 \neg p \rightarrow S & (\text{Rule } T) \\
 \neg S \rightarrow p & (\text{Rule } T, p \rightarrow q \Leftrightarrow \neg q \rightarrow \neg p) \\
 p \rightarrow r & (\text{Rule } P) \\
 \neg S \rightarrow r & (\text{Rule } T) \\
 S \vee r & (\text{Rule } T, p \rightarrow q \Leftrightarrow \neg(\neg p \vee q))
 \end{array}$$

**Exercise Q1** :  $R \wedge (p \vee q)$  is a valid conclusion from the premises  $p \vee q$ ,  $q \rightarrow r$ ,  $p \rightarrow M$  and  $\neg M$ .



**Exercise Q2** : If A works hard, then either B or C will enjoy themselves. If B enjoys himself, then A will not work hard. If D enjoys himself, then C will not, therefore prove that if A works hard, D will not enjoy himself.

**Solution hint** :

*A : A works hard*

*B : B will enjoy himself*

*C : C will enjoy himself*

*D : D will enjoy himself*

To prove  $A \rightarrow \neg D$  follows from  $A \rightarrow B \vee C$ ,  $B \rightarrow \neg A$  and  $D \rightarrow \neg C$