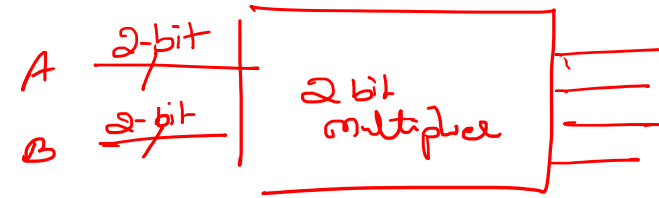


# Multipliers and Magnitude comparators

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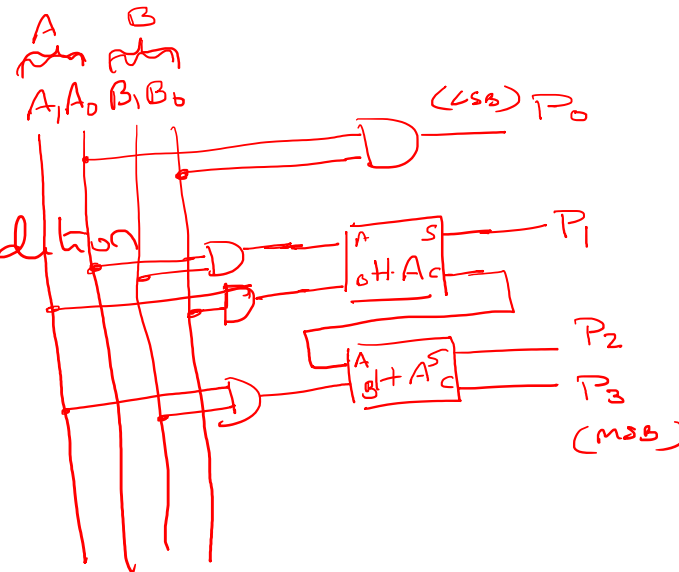
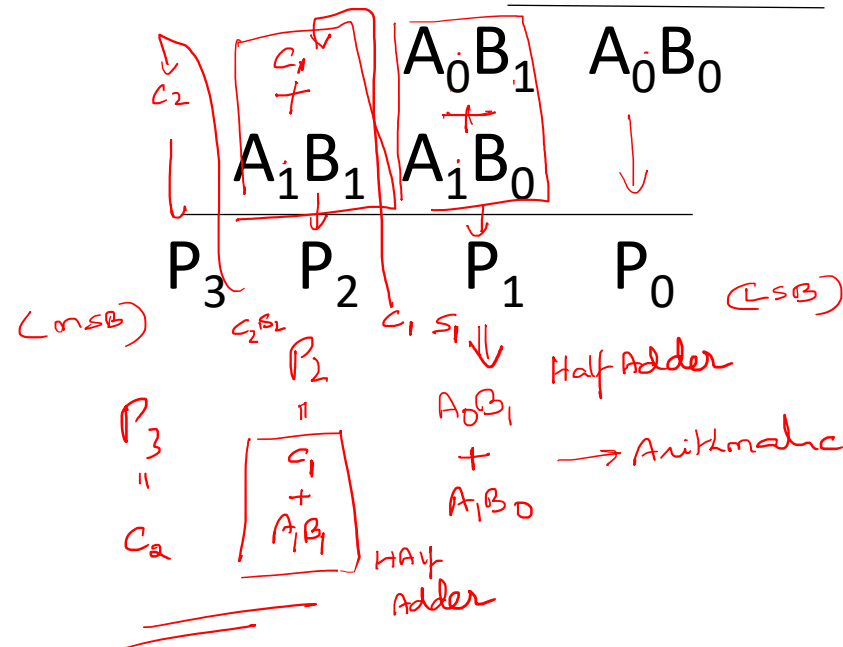
# Binary Multiplier

- 2 bit  $\times$  2 bit binary multiplier using adders and external gates.



$$\begin{matrix} & B & & A \\ & B_1 B_0 & \times & A_1 A_0 \end{matrix}$$

Ex:  $\begin{matrix} b_1 b_0 & a_1 a_0 \\ 10 & \times 11 \\ \hline b_1 a_1 & b_0 a_1 & b_1 a_0 & b_0 a_0 \\ 1 \times 1 & 0 \times 1 & 1 \times 1 & 0 \times 1 \end{matrix}$



Bit by bit multiplication

logical	$a_0$	$b_0$	Product
	0	0	0
	0	1	0
	1	0	0
	1	1	1

Arithmetic bit by bit multiplication } resembles AND operation

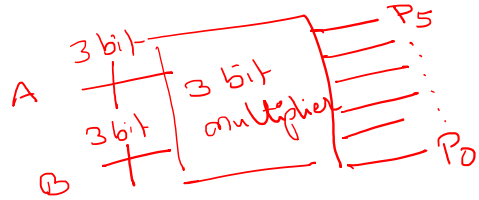
logic

$a_0$   $b_0$

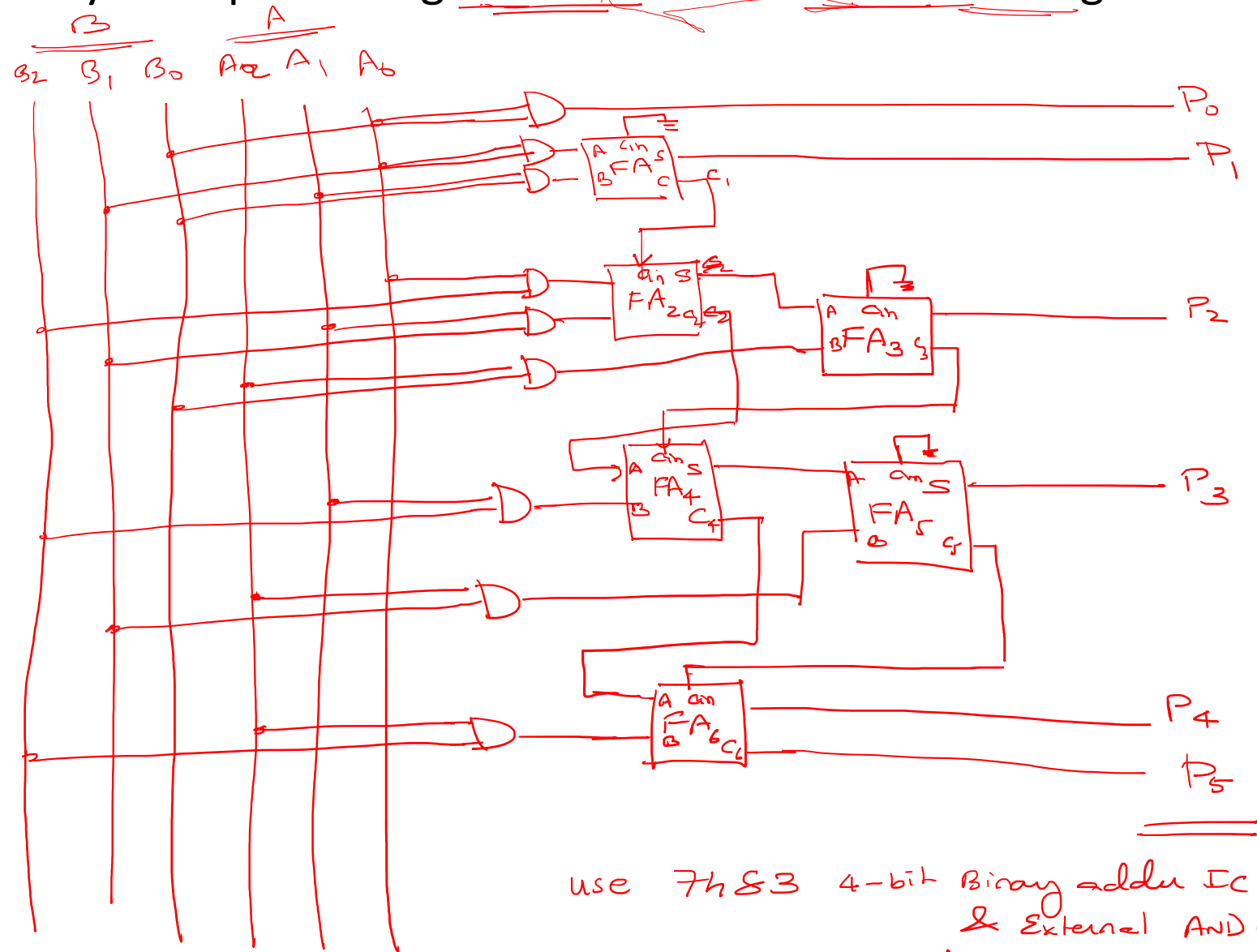
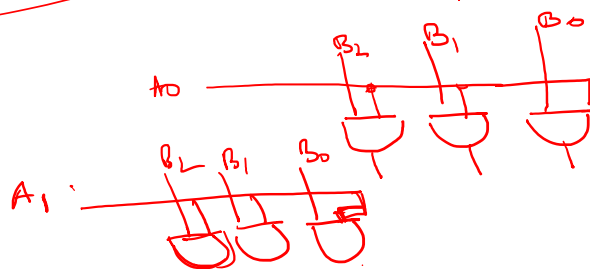
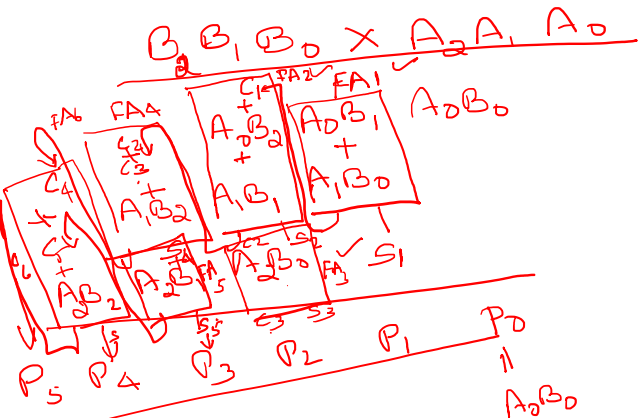
Product



- Design a 3 bit  $\times$  3 bit binary multiplier using Full adders and external AND gates.



$$\begin{array}{r} A \text{ or } B \\ 111 \\ 7 \\ \hline 7 \times 7 = 49 \\ 4 \rightarrow 15 \quad 1111 \\ 5 \rightarrow 31 \quad 11111 \\ \hline 6 \rightarrow 63 \end{array}$$



use 7483 4-bit Binary adder IC  
 & External AND gates  
 to achieve 3x3 bit  
 multiplication.



$$15 \times 7 = 105$$

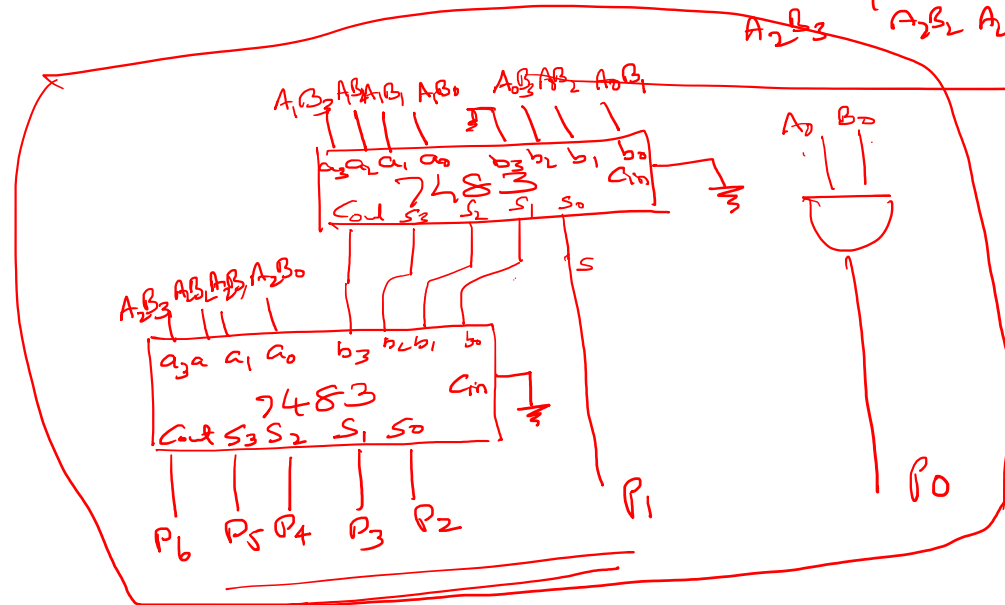
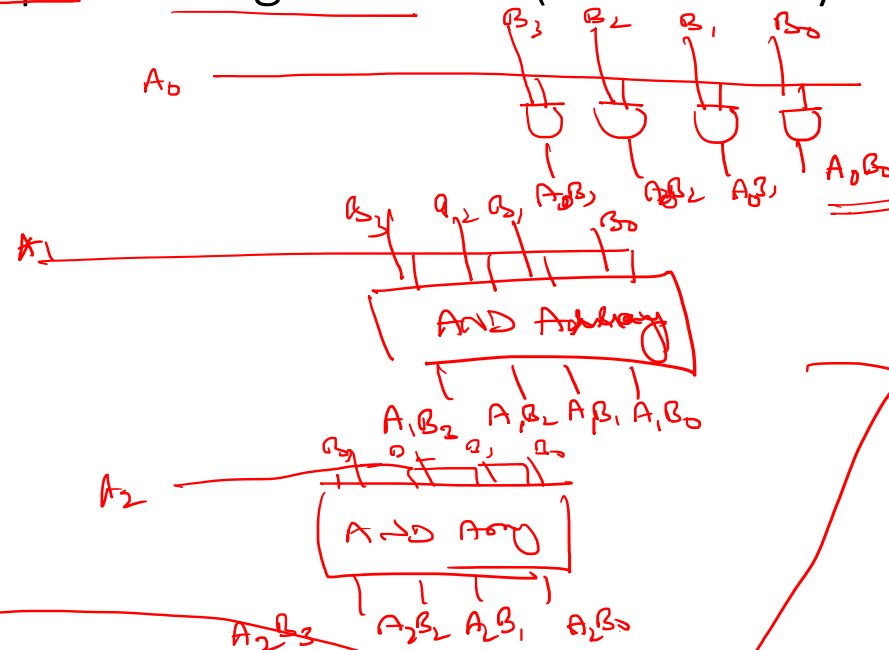
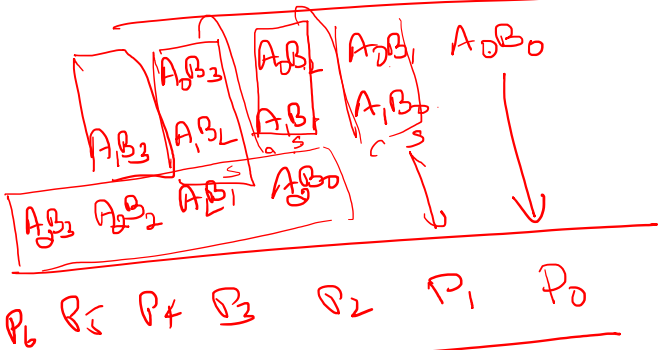
1111 x 111

minimum 7 bits

Design a 4 bit  $\times$  3 bit binary multiplier using 7483 ICs (4 bit binary adders) and external AND gates.

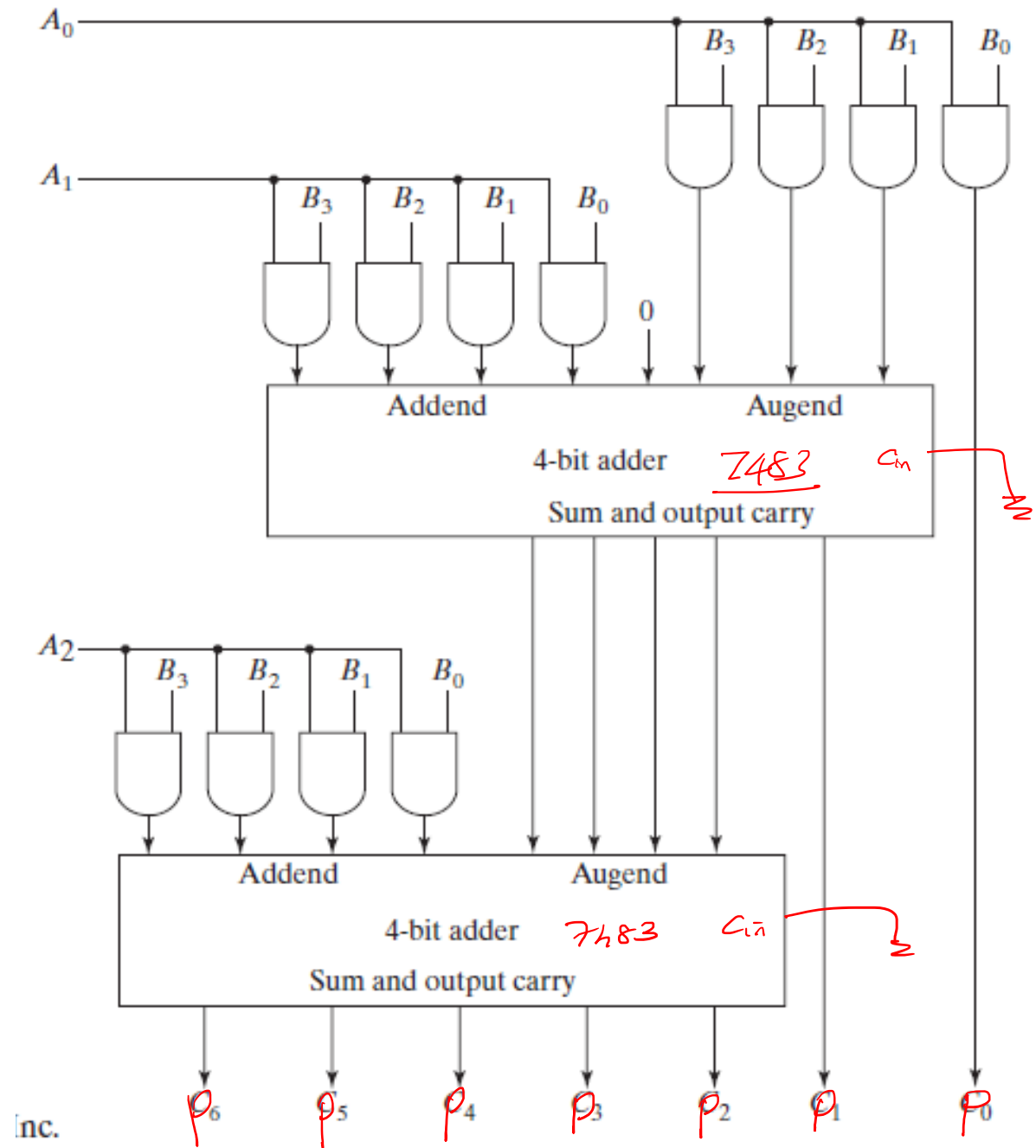
$$B (4bit) \times A (3bit)$$

$$B_3 B_2 B_1 B_0 \times A_2 A_1 A_0$$





## 4-bit by 3-bit binary multiplier





# Magnitude Comparator

- 1 bit Magnitude comparator

Truth table

Input		Output		
A	B	A<B	A=B	A>B
0	0	0	1	0
0	1	1	0	0
1	0	0	0	1
1	1	0	1	0

As Example

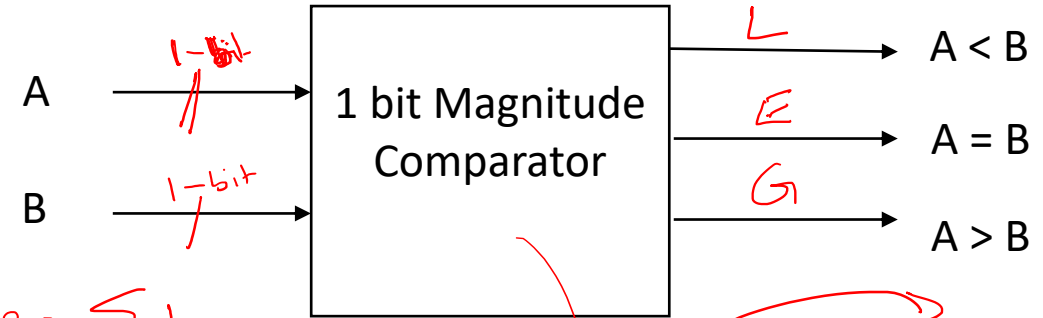
if to compare A & B  $\Rightarrow$  whether  $A > B$   
 $A = B$

① Draw i/p & o/p relation using combinational logic  $A < B$

② subtraction  $A - B$

$$\begin{array}{r} A \\ - B \\ \hline \text{Ans} \end{array}$$
 $\neq 0 \& +ve \quad A > B$   
 $\neq 0 \& -ve \quad A < B$   
 $= 0 \quad A = B$

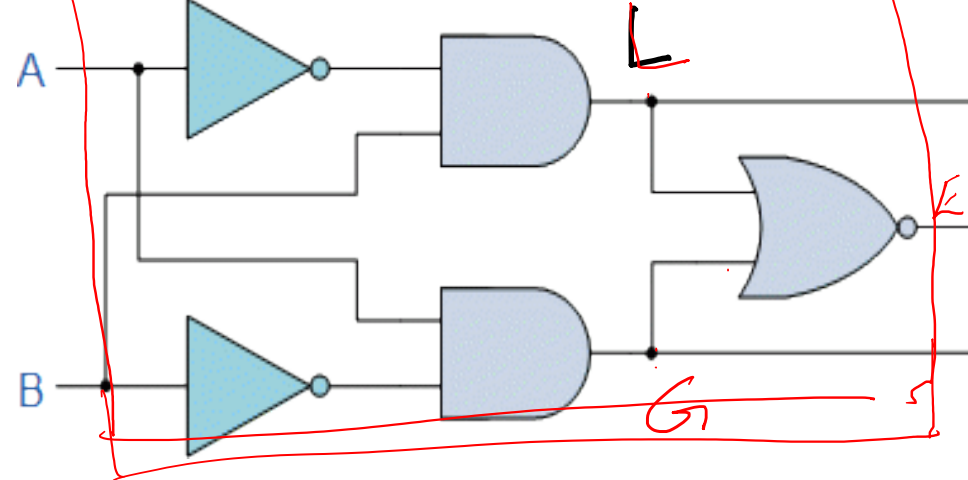
1-bit



$$A < B = \sum_{m_1} 1 = \bar{A}B \checkmark$$

$$A = B = \sum_{m_0, m_3} 1 = \bar{A}\bar{B} + AB = \overline{A \oplus B} = \bar{A}\bar{B} + AB = \bar{A}B + A\bar{B}$$

$$A > B = \sum_{m_2} 1 = A\bar{B}$$



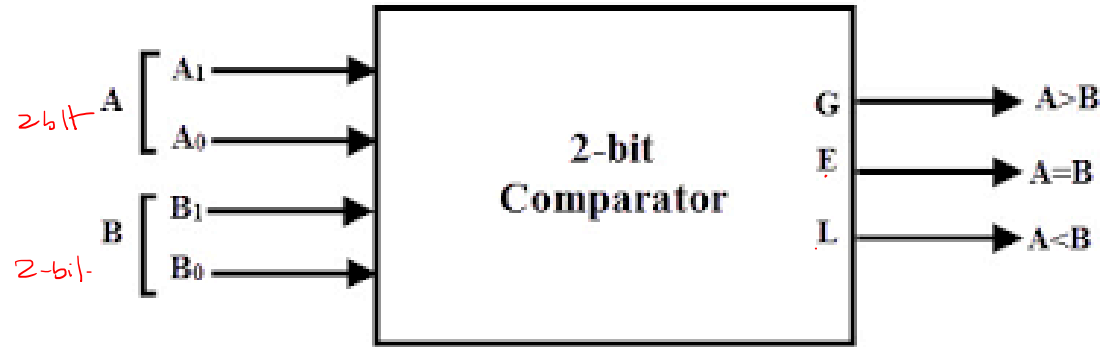
$$= \bar{A}B \checkmark \Rightarrow \underline{\underline{A < B}}$$

$$= \bar{A}\bar{B} + AB \Rightarrow \underline{\underline{A = B}} = \overline{A \oplus B}$$

$$= A\bar{B} \checkmark \Rightarrow \underline{\underline{A > B}}$$

# • 2 bit magnitude comparator

	Inputs				Outputs		
	A <sub>1</sub>	A <sub>0</sub>	B <sub>1</sub>	B <sub>0</sub>	A>B	A=B	A<B
m <sub>0</sub>	0	0	0	0	0	1	0
m <sub>1</sub>	0	0	0	1	0	0	1
m <sub>2</sub>	0	0	1	0	0	0	1
m <sub>3</sub>	0	0	1	1	0	0	1
m <sub>4</sub>	0	1	0	0	1	0	0
m <sub>5</sub>	0	1	0	1	0	1	0
m <sub>6</sub>	0	1	1	0	0	0	1
m <sub>7</sub>	0	1	1	1	0	0	1
m <sub>8</sub>	1	0	0	0	1	0	0
m <sub>9</sub>	1	0	0	1	1	0	0
m <sub>10</sub>	1	0	1	0	0	1	0
m <sub>11</sub>	1	0	1	1	0	0	1
m <sub>12</sub>	1	1	0	0	1	0	0
m <sub>13</sub>	1	1	0	1	1	0	0
m <sub>14</sub>	1	1	1	0	1	0	0
m <sub>15</sub>	1	1	1	1	0	1	0



$$G = A > B = \sum m(4, 8, 9, 12, 13, 14)$$

$$E = A = B = \sum m(0, 5, 10, 15)$$

$$L = A < B = \sum m(1, 2, 3, 6, 7, 11)$$

A <sub>1</sub> A <sub>0</sub> \ B <sub>1</sub> B <sub>0</sub>	00	01	11	10
00	0	0	0	0
01	0	0	0	0
11	0	0	0	0
10	0	0	0	0

$$G = \sum m(4, 8, 9, 12, 13, 14)$$

$$G = A_1 \bar{B}_1 + A_0 \bar{B}_1 \bar{B}_0 + A_1 A_0 \bar{B}_0$$

$$A > B = (A > B) + (A = B) (A > B)$$

$$= A_1 \bar{B}_1 + (\bar{A} \oplus B) (A_0 \bar{B}_0)$$

$$G = A_1 \bar{B}_1 + A_0 \bar{B}_0 [A_1 + \bar{B}_1]$$

$$= A_1 \bar{B}_1 + A_0 \bar{B}_0 (\bar{A}_1 \bar{B}_1 + A_1 B_1)$$

$$= A_1 \bar{B}_1 + A_0 \bar{B}_0 (A_1 \oplus B_1)$$

# 2-bit magnitude comparator

$$A=B \Rightarrow E = \underbrace{(A_1=B_1)}_{E_1 \cdot E_0} \cdot (A_0=B_0) = (\overline{A_1 \oplus B_1}) \cdot (\overline{A_0 \oplus B_0})$$

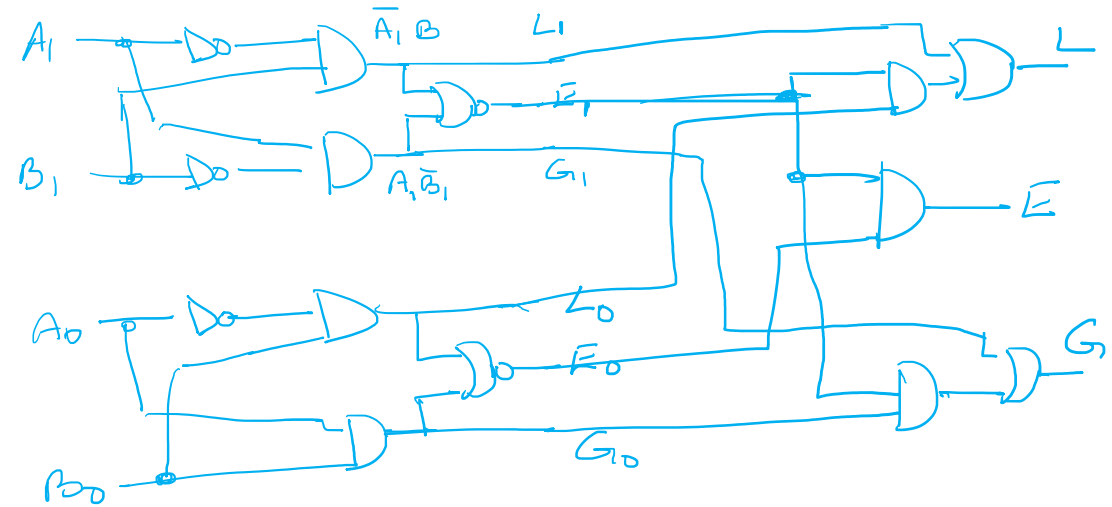
$$A < B \Rightarrow L = (A_1 < B_1) + (A_1 = B_1) \cdot (A_0 < B_0)$$

$$L = \underbrace{(A_1 B_1)}_{L_1} + \underbrace{(\overline{A_1 \oplus B_1}) \cdot (\overline{A_0 B_0})}_{E_1 \cdot L_0}$$

$$A > B = G = (A_1 > B_1) + (A_1 = B_1) \cdot (A_0 > B_0)$$

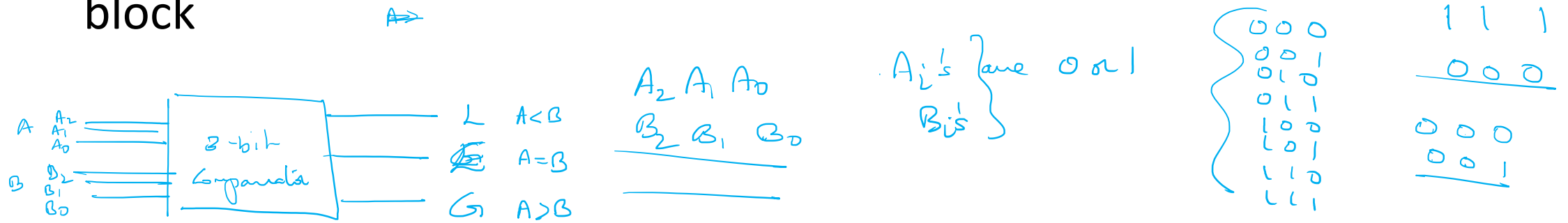
$$G = \underbrace{A_1 \overline{B_1}}_{G_1} + \underbrace{(\overline{A_1 \oplus B_1}) \cdot (A_0 \overline{B_0})}_{E_1 \cdot G_0}$$

$$\overline{A_1 \oplus B_1} = \overline{A_1 B_1 + A_1 \overline{B_1}}$$



Two-bit magnitude comparator  
using 1-bit magnitude  
comparators

- 3 bit magnitude comparator using 1-bit magnitude comparator block



$$\underline{G \text{ or } A > B} = G_2 + E_2 G_1 + E_2 E_1 G_0 = A_2 \bar{B}_2 + (A_2 \oplus B_2) \cdot A_1 \bar{B}_1 + (A_2 \oplus B_2) (A_1 \oplus B_1) A_0 \bar{B}_0$$

$$(A_2 > B_2) + (A_2 = B_2) (A_1 > B_1) + (A_2 = B_2) (A_1 = B_1) (A_0 > B_0)$$

$$\underline{E \text{ or } A = B} = E_2 \cdot E_1 \cdot E_0 = (\overline{A_2 \oplus B_2}) \cdot (\overline{A_1 \oplus B_1}) \cdot (\overline{A_0 \oplus B_0})$$

$$(A_2 = B_2) \cdot (A_1 = B_1) \cdot (A_0 = B_0)$$

$$\underline{L \text{ or } A < B} = L_2 + E_2 L_1 + E_2 E_1 L_0 = \bar{A}_2 B_2 + (\overline{A_2 \oplus B_2}) \cdot \bar{A}_1 B_1 + (\overline{A_2 \oplus B_2}) \cdot (\overline{A_1 \oplus B_1}) \cdot \bar{A}_0 B_0$$

$$(A_2 < B_2) + (A_2 = B_2) (A_1 < B_1) + (A_2 = B_2) (A_1 = B_1) (A_0 < B_0)$$

- 4 bit magnitude comparator



$A_3 A_2 A_1 A_0$   
 $B_3 B_2 B_1 B_0$

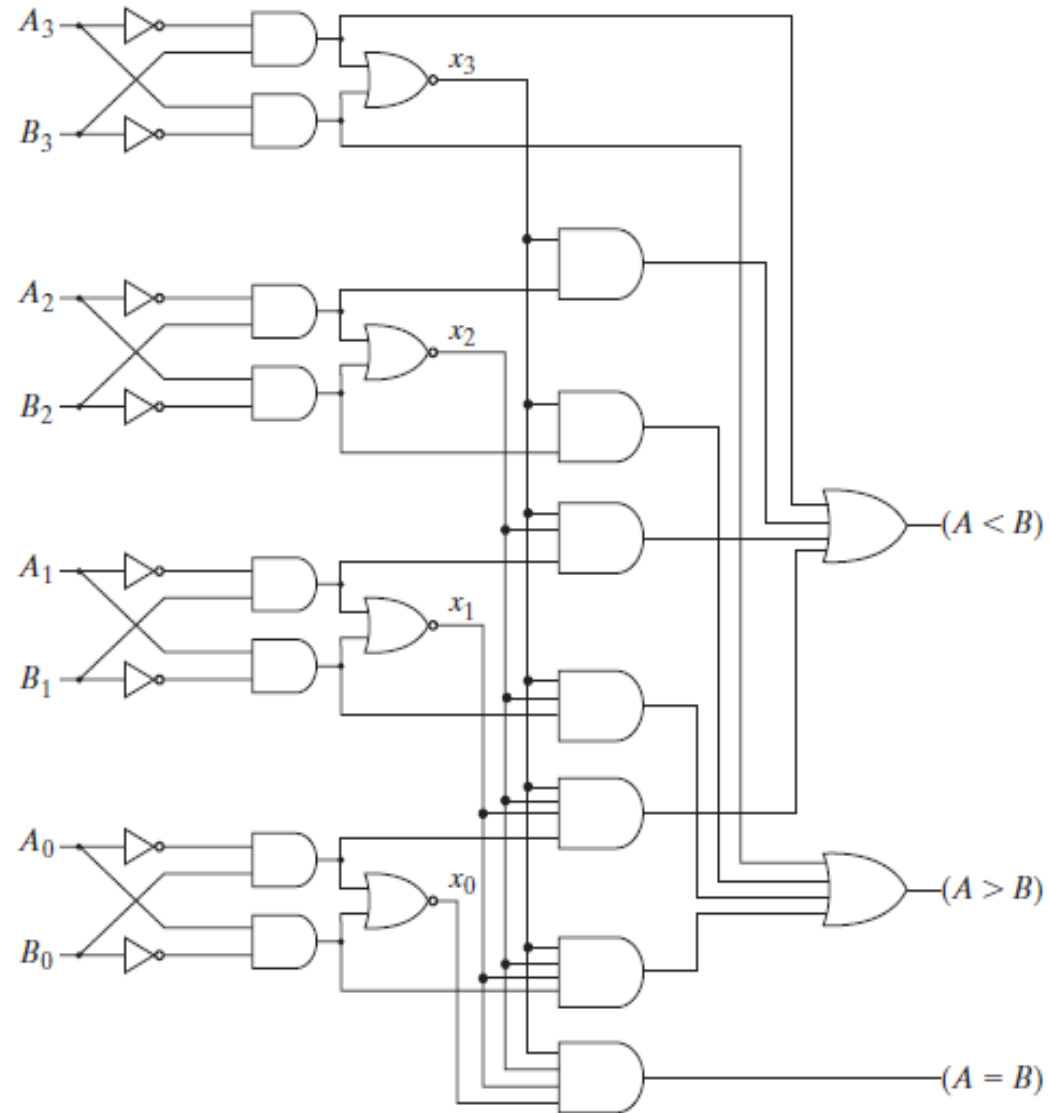
$$G = A > B = G_3 + \bar{E}_3 G_2 + \bar{E}_3 \bar{E}_2 G_1 + \bar{E}_3 \bar{E}_2 \bar{E}_1 G_0$$

$$A_3 \bar{B}_3 + (\overline{A_3 \oplus B_3}) A_2 \bar{B}_2 + \dots$$

$$E = A = B \Rightarrow \bar{E}_3 \cdot \bar{E}_2 \cdot \bar{E}_1 \cdot \bar{E}_0$$

$$L = A < B \Rightarrow L_3 + \bar{E}_3 L_2 + \bar{E}_3 \bar{E}_2 L_1 + \bar{E}_3 \bar{E}_2 \bar{E}_1 L_0$$

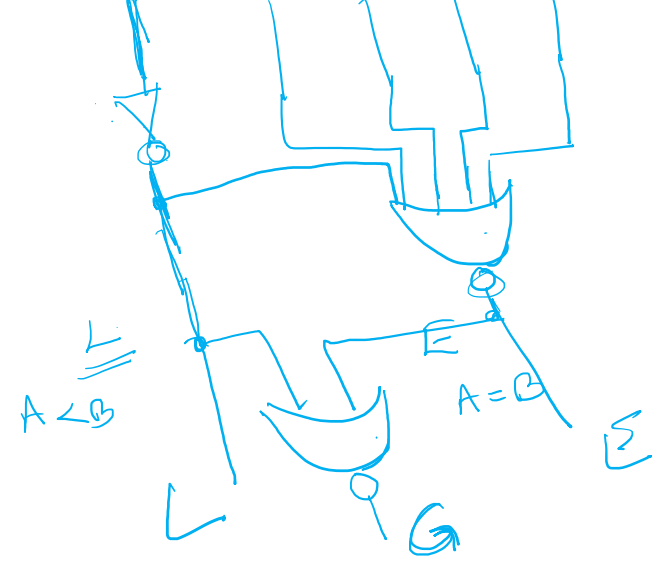
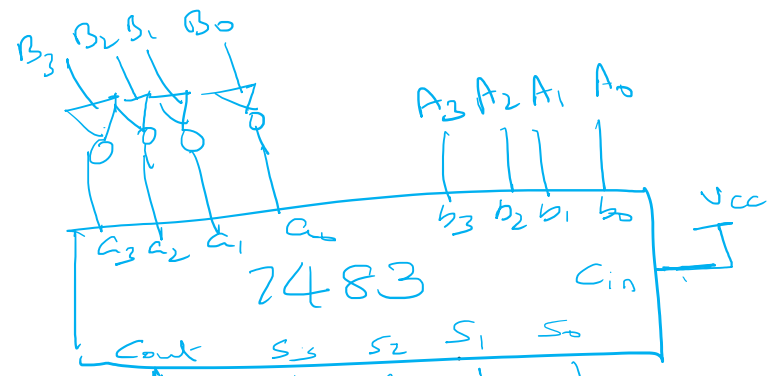
# 4-bit magnitude comparator



only NoR

# Design a 4-bit magnitude comparator using 7483 IC and external gates

What is 7483 IC? → it is 4-bit Binary Adder



$A \rightarrow 0110$   
 $-B \rightarrow 1011$   
 $+2's \text{ } \overline{B}$   
 $-0100 \rightarrow 0010$  (2's complement)

EQUAL?

$4 \rightarrow 0100$   
 $-4 \rightarrow 1011$  (2's complement)  
 $0000$  (Ans)

G=?

$$G = \overline{L} \cdot \overline{E}$$

$4 \rightarrow 0100$   
 $-6 \rightarrow 0110$  (Carry → +ve)  
 No carry → -ve

$0100$   
 $1001$   
 $1110$  (Ans - 2's complement)

$$L + E = \overline{L} \cdot \overline{E}$$