

\* Poset :  $(A, R)$  s.t  $R$  is

- i) reflexive
- ii) antisym
- iii) transitive

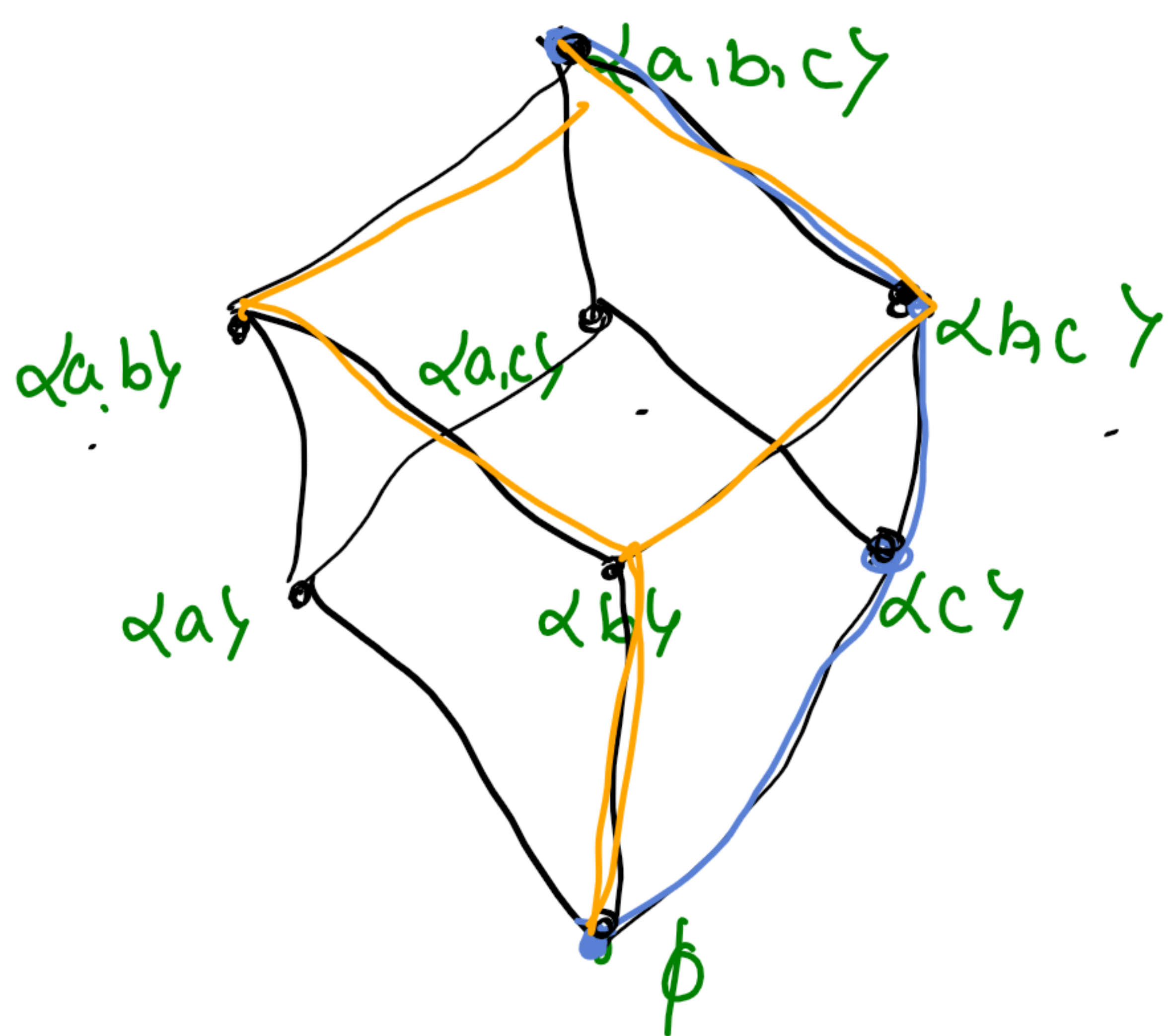
\* Totally ordered set / chain :- For any two elts  $a$  &  $b$ , either  $aRb$  &  $bRa$

\* Antichain :- For any 2 elts  $a$  &  $b$ , neither  $aRb$  nor  $bRa$

\* maximal elts  
minimal elts

\* Hasse Diagram :

\*  $(P(S), \subseteq)$  is a poset, where  $S = \{a, b, c\}$   
 $P(S)$  = powerset



$\phi, \{c\}, \{b, c\}, \{a, b, c\}$

Length of Chain :- No of elts  
in the chain

Length of the longest chain = 4

$\{ \phi \}$ ,  $\{a\}, \{b\}, \{c\}$ ,  $\{a, b\}, \{b, c\}, \{a, c\}$ ,  $\{a, b, c\}$



**Theorem :** Let  $(P, \leq)$  be a poset. Suppose length of the longest chain in  $P$  is ' $n$ '. Then elements in  $P$  can be partitioned into  $n$  disjoint antichains

**Proof :-** Induction on the length of the longest chain ' $n$ '.  
When  $n = 1$  : No 2 elts are related, clearly they form an antichain

Assume that the result is true when the length of the longest chain in the poset is  $(n-1)$

Let  $P$  be a poset in which length of the longest chain is ' $n$ '.

Let  $M$  be the set of all maximal elts. Clearly  $M$  is a nonempty & an antichain (no 2 elts of  $M$  are related)

consider  $(P-M, \leq)$

Since there is no chain of length ' $n$ ' in  $(P-M)$ , the length of the longest chain would be at most  $(n-1)$ .

Further, since no 2 elts of  $M$  are related, obviously the length of the longest chain in  $(P-M)$  is exactly  $(n-1)$

By induction hypothesis,  $(P-M)$  can be partitioned into  $(n-1)$  disjoint antichains. Thus,  $P$  can be partitioned into  $n$  disjoint antichains

$$* a \leq a \vee b \quad \& \quad a \wedge b \leq a$$

$$* \text{ If } a \leq b \text{ \& } c \leq d, \text{ then } a \vee c \leq b \vee d \\ a \wedge c \leq b \wedge d$$



Lattice :- unique lub & glb

Basic prop<sup>s</sup> of algebraic sys defined by lattices

Let  $(A, \leq, \vee, \wedge)$  be the algebraic sys defined by the lattice  $(A, \leq)$

i) Commutative law :- For any 2 elts  $a, b \in A$

i)  $a \vee b = b \vee a$

ii)  $a \wedge b = b \wedge a$

ii) associative law :- for any 3 elts  $a, b, c \in A$

i)  $a \vee (b \vee c) = (a \vee b) \vee c$

ii)  $a \wedge (b \wedge c) = (a \wedge b) \wedge c$

Proof :-

Let  $\underline{a \vee (b \vee c) = g} \quad (a \vee b) \vee c = h$

To show that  $g = h$ , let me show  $g \leq h$  &  $h \leq g$   
 $\Rightarrow g = h$  (antisym)

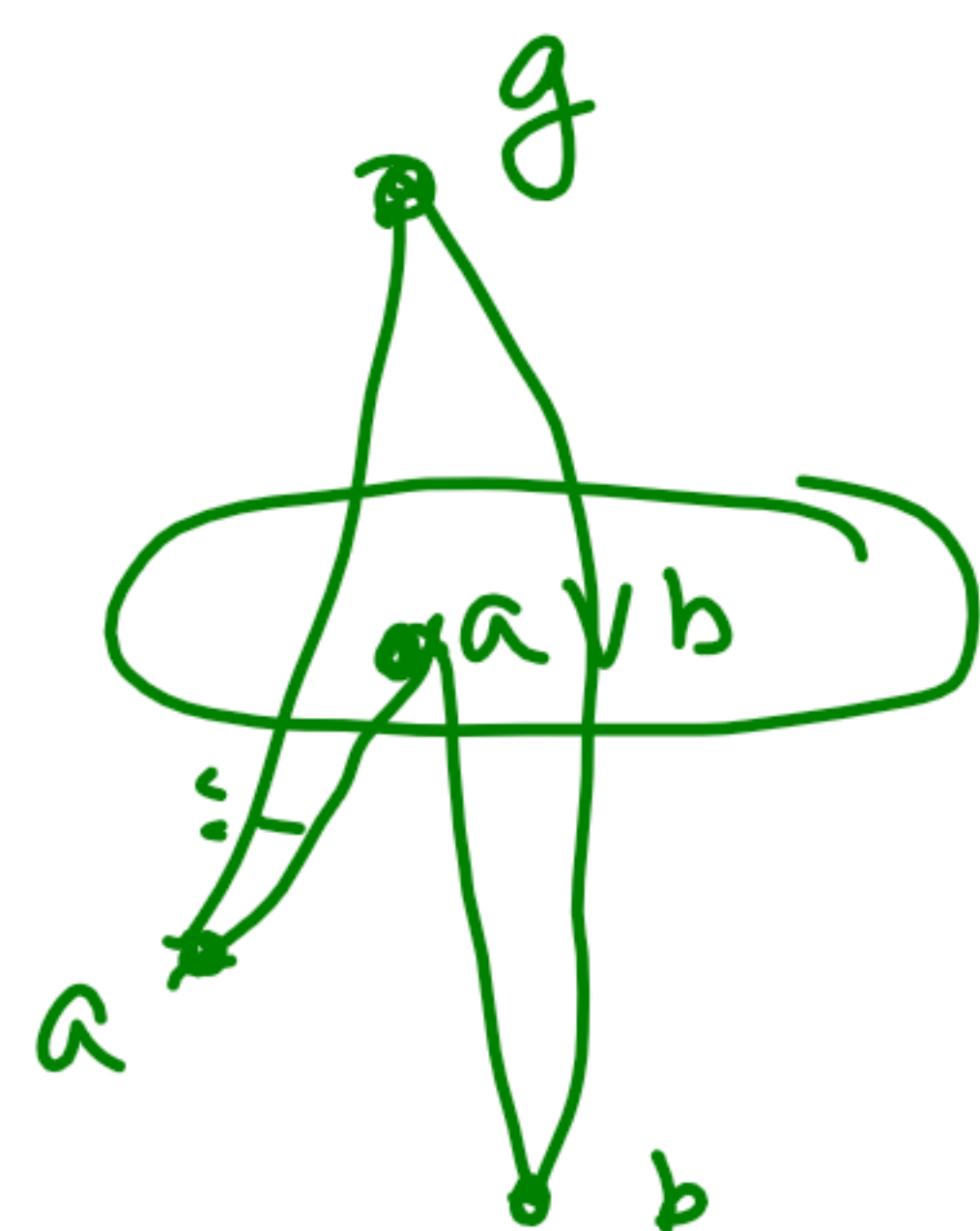
$a \leq a \vee (b \vee c)$  |  $a \leq a \vee b$

$a \leq g$  ——— ①

$(b \vee c) \leq g$

$b \leq b \vee c \leq g \Rightarrow b \leq g$  ——— ②

$c \leq b \vee c \leq g \Rightarrow c \leq g$  ——— ③



Since  $(a \vee b)$  is the lub of  $a$  &  $b$ ,  $g$  is some ub of both  $a$  &  $b \Rightarrow a \vee b \leq g$



$$a \vee b \leq g \quad \text{and} \quad c \leq g$$

$$(a \vee b) \vee c \leq g$$

$(a \vee b) \vee c$  is the least upper bound of  $(a \vee b)$  &  $c$ .  
 But  $g$  is some upper bound.

$$h \leq g$$

Similarly one show  $g \leq h$  (Try)

$$\therefore \underline{\underline{g = h}}$$

$$a \vee (b \vee c) = (a \vee b) \vee c$$

On taking dual;

$$\underline{\underline{a \wedge (b \wedge c) = (a \wedge b) \wedge c}}$$

③ Idempotent law:  $a \vee a = a$   $\forall$  all  $a \in A$   
 $a \wedge a = a$

④ Absorption Law:

$$\begin{aligned} \text{i) } & \underline{a \vee (a \wedge b) = a} \\ \text{ii) } & a \wedge (a \vee b) = a \end{aligned} \quad \forall \text{ all } a, b \in A$$

Proof:-

$$a \leq a \quad (\text{reflexive}) \quad \text{--- ①}$$

$$a \wedge b \leq a \quad (\text{known}) \quad \text{--- ②}$$

$$a \vee (a \wedge b) \leq a \vee a$$

$$a \vee (a \wedge b) \leq a \quad (\because \text{Idempotent})$$

$\because$   
 $a \leq b$  &  $c \leq d$   
 then  
 $a \vee c \leq b \vee d$



## Distributive Lattice:

A lattice is said to be a distributive lattice if the meet operation is distributive over the join operation & the join operation is distributive over the meet operation

For any elements  $a, b, c \in A$

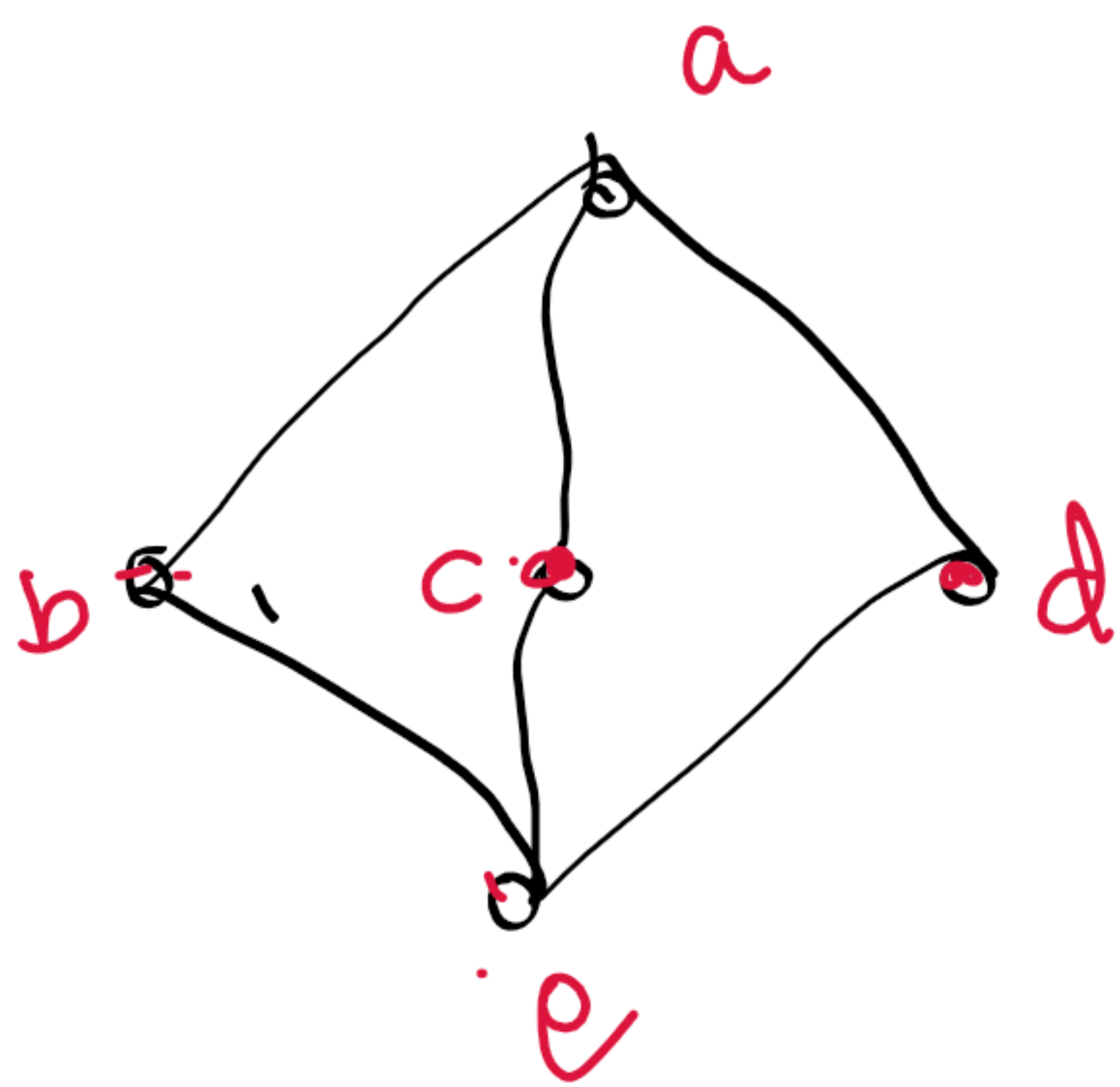
$$i) a \wedge (b \vee c) = (a \wedge b) \vee (a \wedge c)$$

$$ii) a \vee (b \wedge c) = (a \vee b) \wedge (a \vee c)$$

Distributive Laws

Ex:- ①  $(PCS, \subseteq)$ :

②



$$b \wedge (c \vee d) = b \wedge a = b$$

$$(b \wedge c) \vee (b \wedge d) = e \wedge e = e$$

Not distributive

$$LHS \neq RHS$$

Theorem :-

If the meet operation is distributive over the join operation in a lattice, then join is distributive over the meet operation. Similarly, if join is distributive over meet, then meet is distributive over join.

Proof

Consider the lattice  $(A, \leq)$

Let  $a \wedge (b \vee c) = (a \wedge b) \vee (a \wedge c)$ , I've to show that



$$a \vee (b \wedge c) = (a \vee b) \wedge (a \vee c)$$

consider <sup>RHS =</sup>  $\underbrace{(a \vee b)} \wedge \underbrace{(a \vee c)}$   
 $= \underbrace{[(a \vee b) \wedge a]} \vee [(a \vee b) \wedge c] \quad (\because \text{meet is distributive})$

$$= a \vee [(a \vee b) \wedge c] \quad (\because \text{absorption})$$

$$= a \vee [c \wedge (a \vee b)] \quad (\because \text{comm prop})$$

$$= \underbrace{a \vee}_{\text{meet is dist over join}} [(c \wedge a) \vee (c \wedge b)] \quad (\because \text{meet is dist over join})$$

$$= \underbrace{[a \vee (c \wedge a)]} \vee (c \wedge b) \quad (\because \text{associative})$$

$$= a \vee (c \wedge b) \quad (\because \text{absorption})$$

$$= a \vee (b \wedge c) \quad (\because \text{commutative})$$

$$= \text{LHS}$$

$$\begin{aligned} A \vee (B \vee C) \\ = (A \vee B) \vee C \end{aligned}$$

Another statement follows from the duality principle

\* ① Let  $a$  &  $b$  be two elements in the lattice  $(A, \leq)$   
 s.t.  $a \wedge b = b$  iff  $a \vee b = a$

proof:-

Let  $a \wedge b = b$ , prove that  $a \vee b = a$

consider  $a \vee (a \wedge b) = a \quad (\because \text{absorption})$

$$a \vee b = a \quad (\because a \vee b = b)$$



conversely, let  $a \vee b = a$ , I've prove  $a \wedge b = b$

consider  $b \wedge (a \vee b) = b$  ( $\because$  absorption)

$$b \wedge a = b$$

$$\underline{\underline{a \wedge b = b}}$$

② Let  $a, b, c$  be elements of lattice  $(A, \leq)$ . Show that  
if  $a \leq b$ , then  $a \vee (b \wedge c) \leq b \wedge (a \vee c)$

Proof:-

$$\text{Let } a \leq b \quad \text{--- ①}$$

$$a \leq a \vee c \quad \text{--- ②}$$

$$\text{①} \wedge \text{②} \Rightarrow a \leq b \wedge (a \vee c) \quad \text{--- ③}$$

$$\text{w.k.t } b \leq b \text{ and } c \leq a \vee c$$

$$b \wedge c \leq b \wedge (a \vee c) \quad \text{--- ④}$$

From ③ & ④

$$\underline{\underline{a \vee (b \wedge c) \leq b \wedge (a \vee c)}}$$

Prob

① Let  $a, b, c$  be elements in the lattice  $(A, \leq)$ . Show that

$$\text{i) } a \vee (b \wedge c) \leq (a \vee b) \wedge (a \vee c)$$

$$\text{ii) } (a \wedge b) \vee (a \wedge c) \leq a \wedge (b \vee c)$$



② Let  $(A, \vee, \wedge)$  be an algebraic sys. where  $\vee$  &  $\wedge$  satisfy commutative, associative & absorption laws.

a) Define a binary relation  $\leq$  as follows:

For all  $a, b \in A$ ,  $a \leq b$  iff  $a \wedge b = a$

S.t.  $\leq$  is a partial ordering relation

b) S.t.  $(a \vee b)$  is the lub of  $a$  &  $b$  in  $(A, \leq)$   
 $(a \wedge b)$  is the glb of  $a$  &  $b$  in  $(A, \leq)$