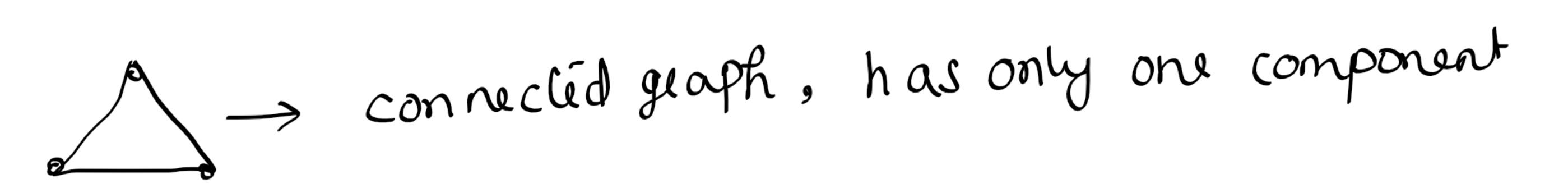
conne led geaphs:

A geaph is said to be connected if there is a path blush every pair of vertices. Else its called a disconnected graph connected component? of a graph is a maximal connected sub geaph



Girth of a graph: denoth of the shortest cycle present in the geaph (Lipany). Denoted by g(G)

ensumference: Length of the longest cycle in G. De noted by

$$g(G) = 3$$
 $c(G) = 4$

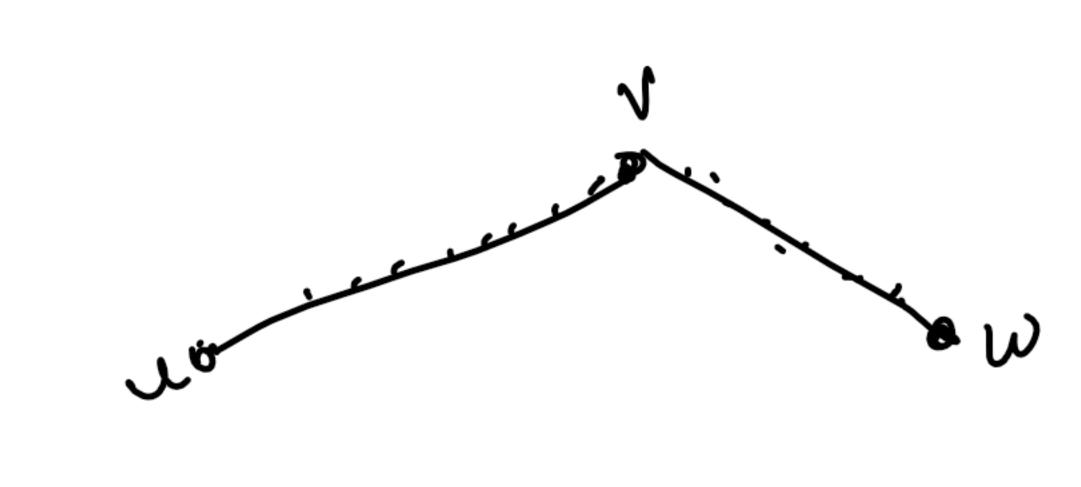
Distance d(u,v) bliss the verties u & v:

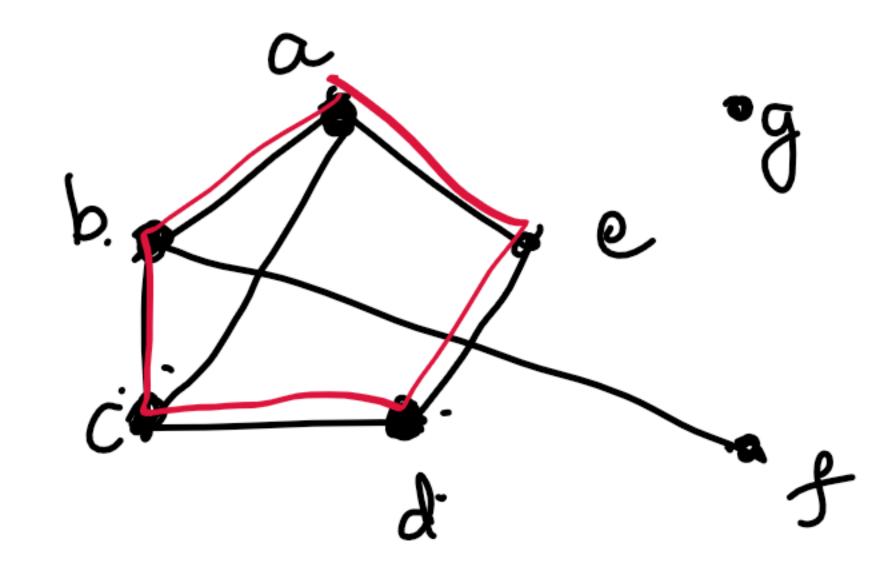
The distance d(u,v) b/wn uf i is the length of the shoutest path joring them if exists. o the wise d(u,v)=00

on a connected geaph, d (u, v) satisfies

3) d(u,v)=0 with d(u,v)=0 2/1 u=V

- II) d(u,v) = d(v,u)
- (iii) d(u,v) + d(v,w) > d(u,w)





$$\Delta(G) = 3$$
ginth $g(G) = 3$

The paths blun a & d -) abcd? acd=

iecumy C(n) = 5

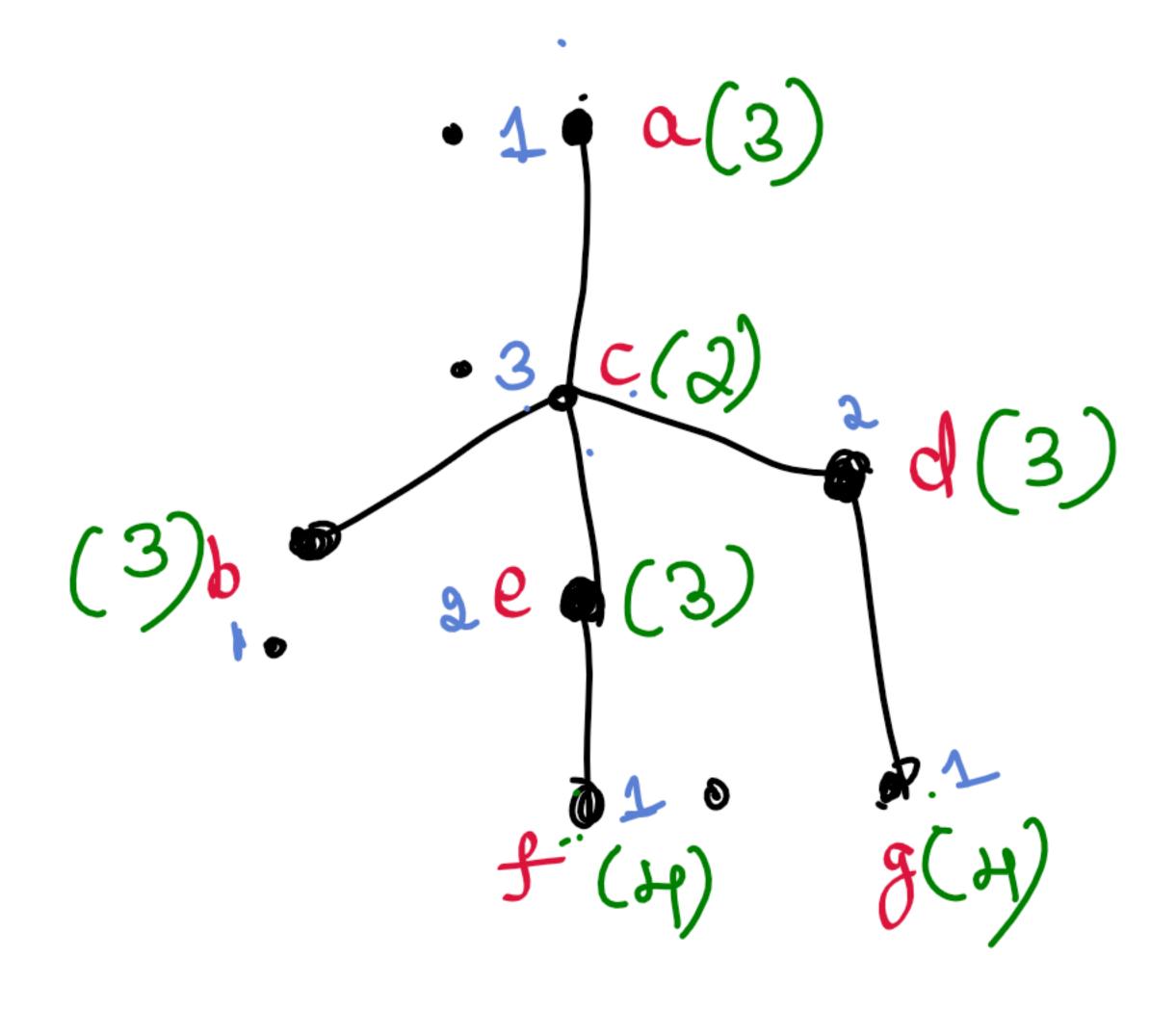
geodesic: shøtest path

Eccentaicity of avertex (e(v)): Distance of v from the farthest vertex

Radius of a geaph: nivinum eccentricity

D'anette of a geaph; maximum e coentricity
091

Cen9th of the Congest geodesic ie maxq d c u,v)}



$$91adius(G) = 2$$
 $diametu(G) = 4$
 $s(G) = 1$

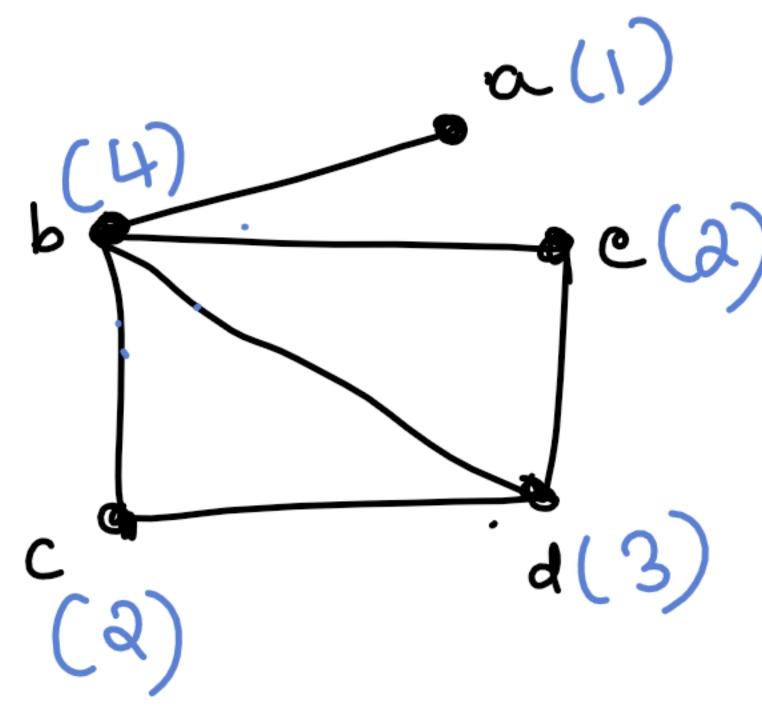
 $\Delta(G) = 3$

minimum degree (SCh)
maximum degree (DCh).

2 Reven 1 productions Then Stages (v) = 29

ver(a) (sumofder of twice the no all vertices — of edges)

Every edge is incident of 2 verties of every edge contributes degree 2 to the Sideg(v)



$$P=5$$
 $q=6$
 $\sum deg(v) = deg(a) + deg(b) + deg(c) + deg(d)$
 $= 1 + 4 + 2 + 3 + 2 = 12$

= 2(6)

Corollaay.

on any geaph, the mo of vertices of degree is even (No of odd deg vertices in any glaph is even)

PROOF:

Se > S degrees of all even deg vertices

So > S of degrees of all odd deg vertices

From prev thme So Se = & 9 (: prev thm)

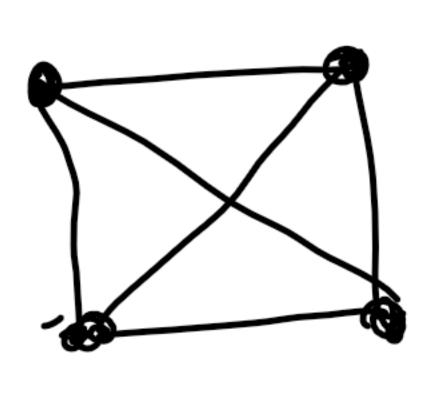
even even

So = & 9 - S = even - even

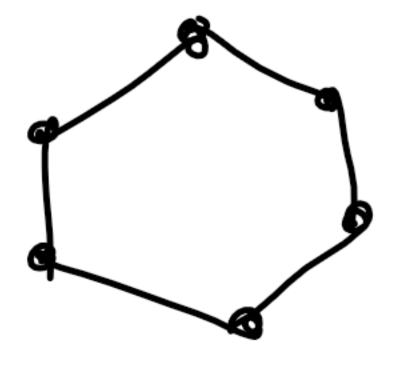
= even

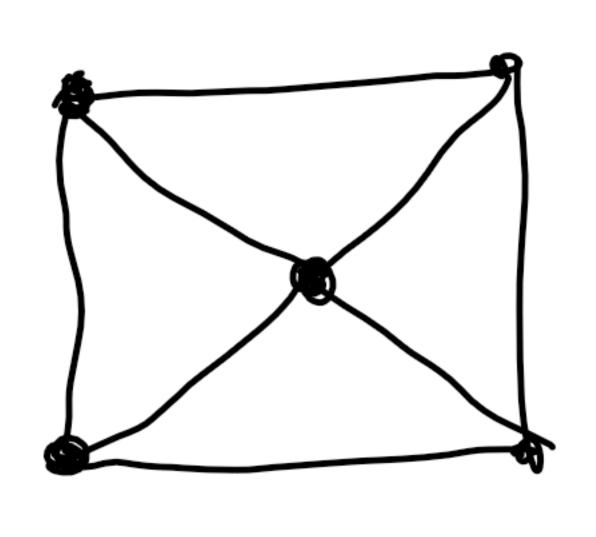
since sois even à each termin sois odd. =) there are even of the s

Peaged geoph: if no two vertices in the geaph are of same degre (De aw ageaph on 5 verties sit each vertex has distinct degree). *. TODaw aperfect georgh on 5 verhões) () deg(') 0 < deg(v) < 4 (a) (3), 4 (e) (DS) (3) (4) Pigeon hole principle Show that no glaph is perfect. (Tsy) Types of geaphs 1) Regular glaph? - Every verlex is of same degree of every rentex has degree a, then we call it 37-regular geaph 2-regular graph



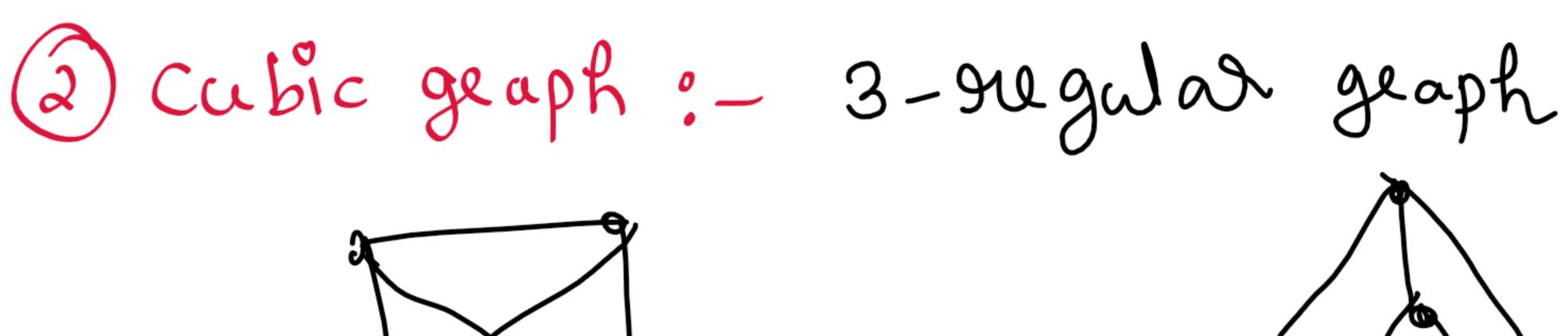
3-regular glaph

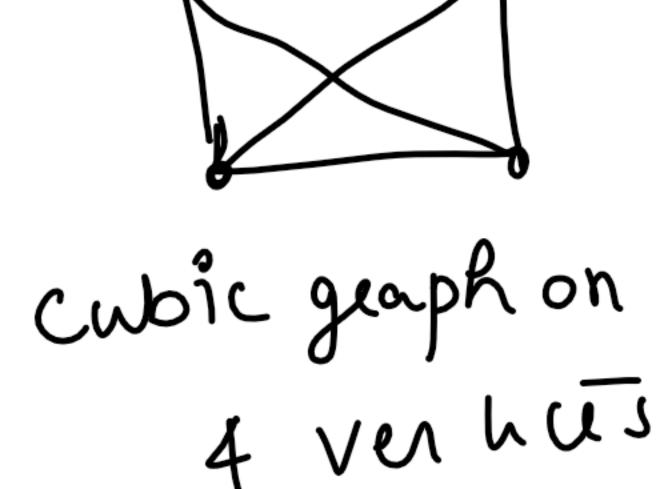


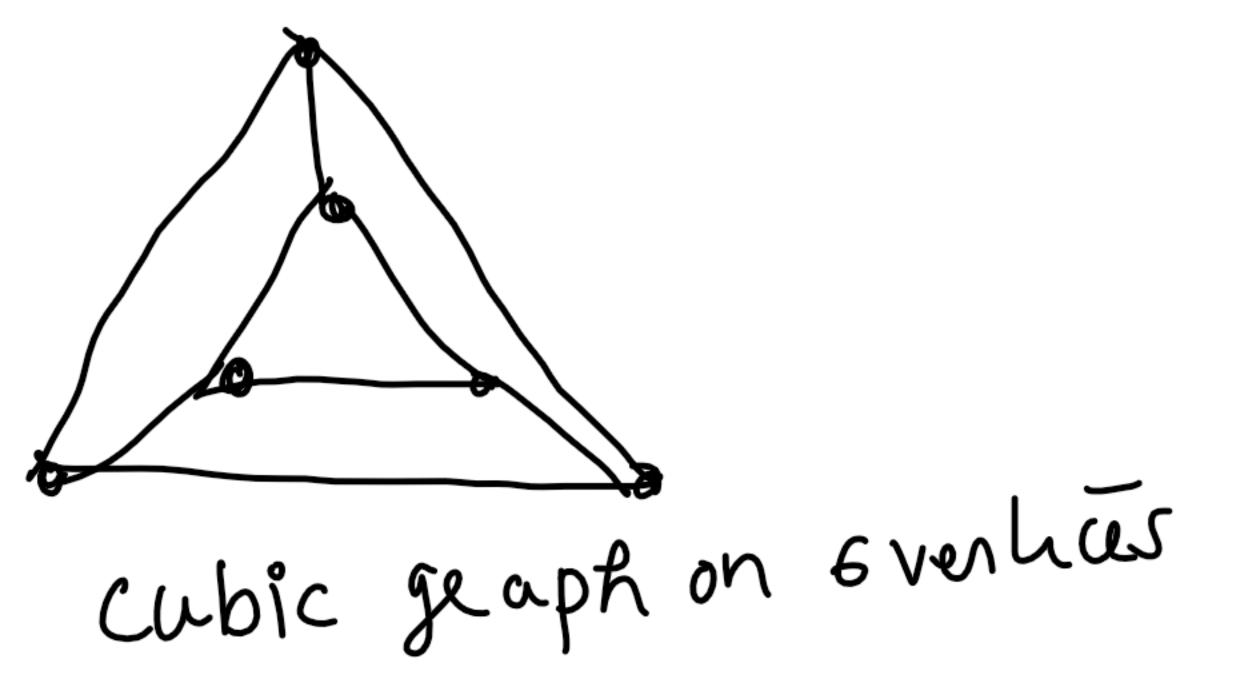


Not a regular geaph

All cycles all 2-regular

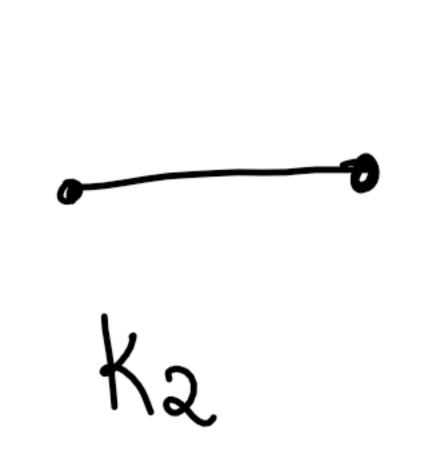


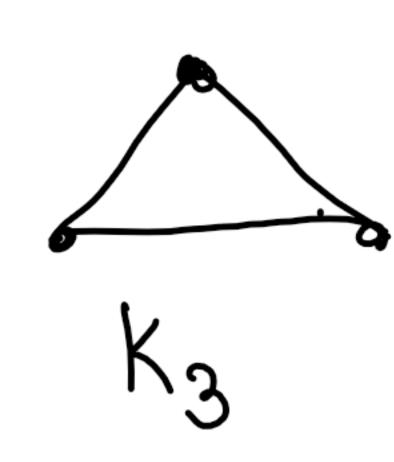


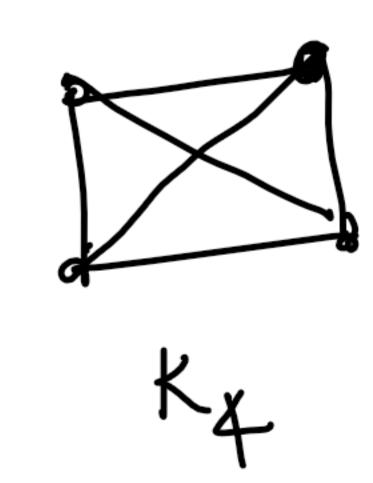


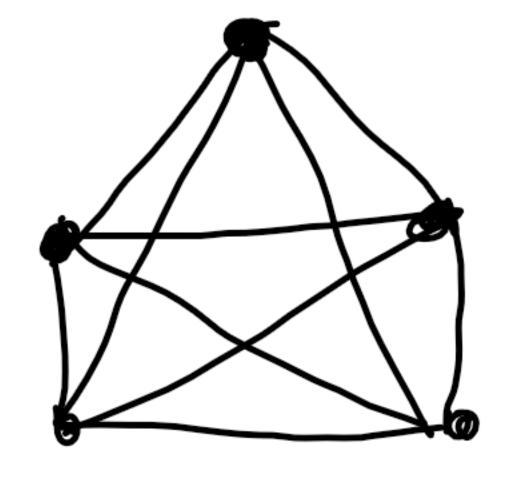
* (Dlaw a cubic graph on 5 vertices) -> Not possible *(No of vertices in any cubic graph must be even Box no of odd deg vertices must be even

3 Complete geaph; if every pair of vertiles are adjacent A complete graph on predicts —> Kp





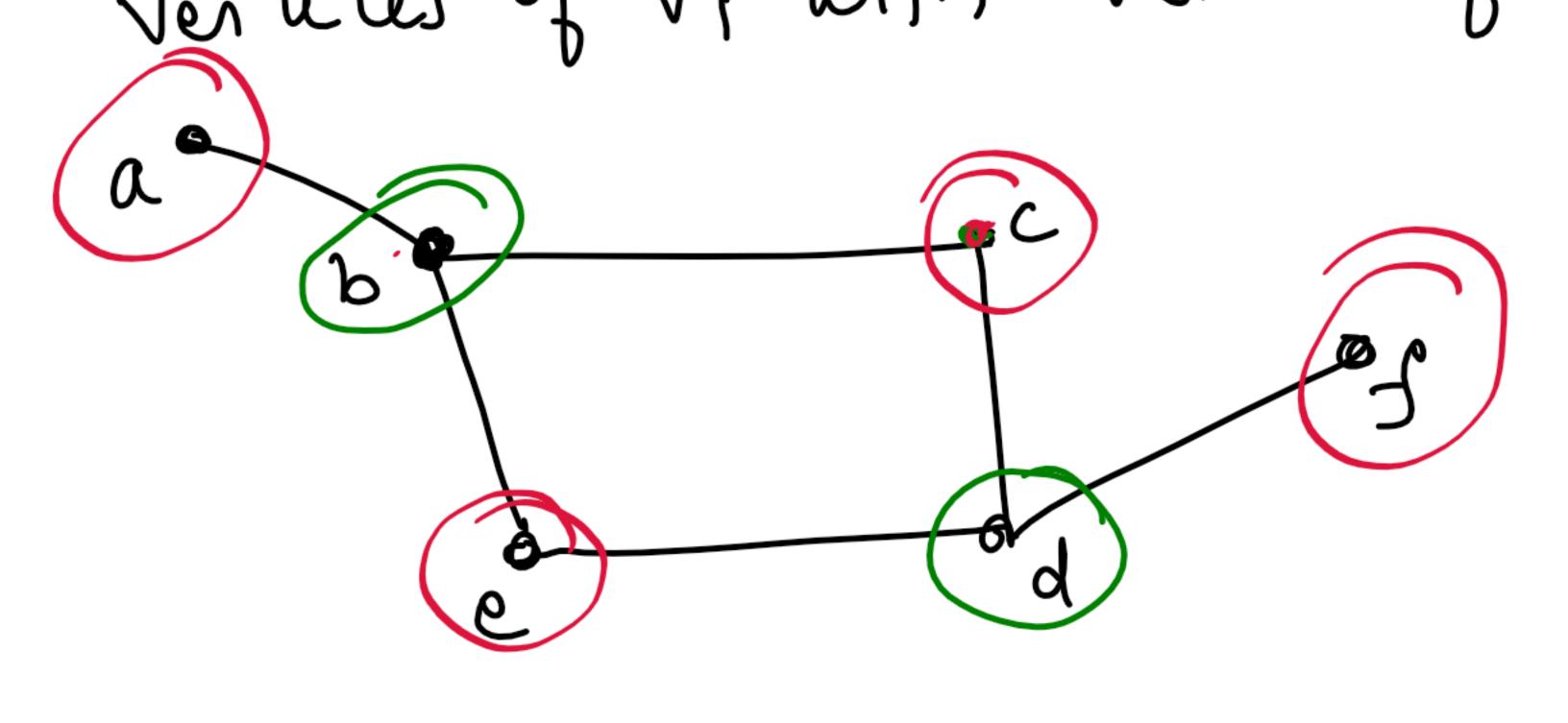


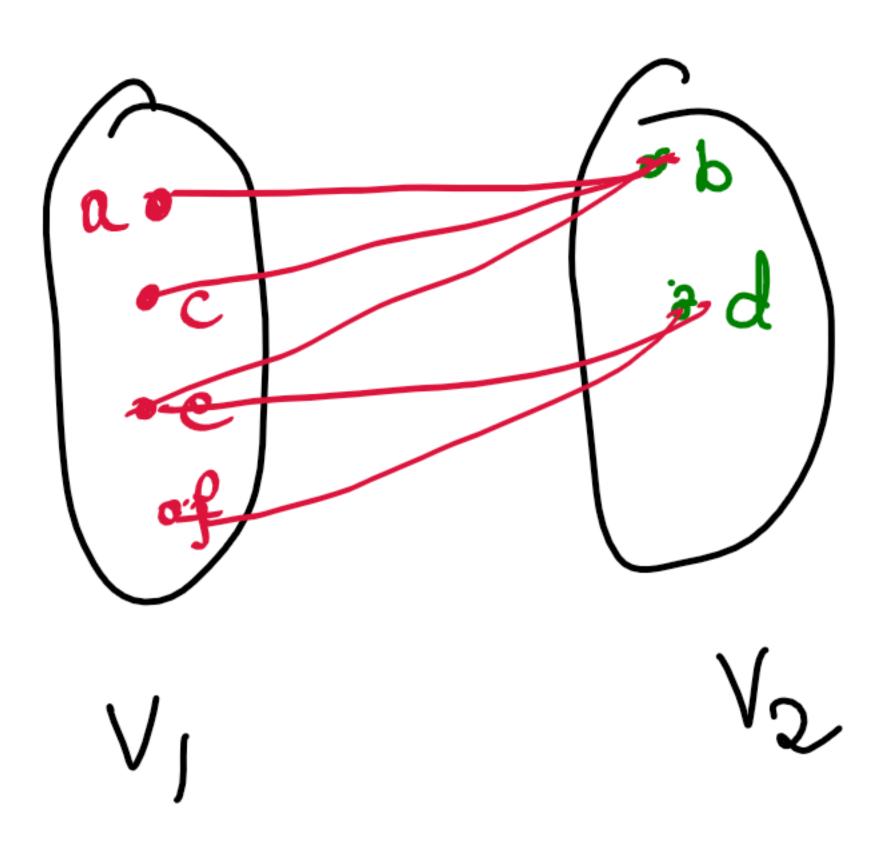


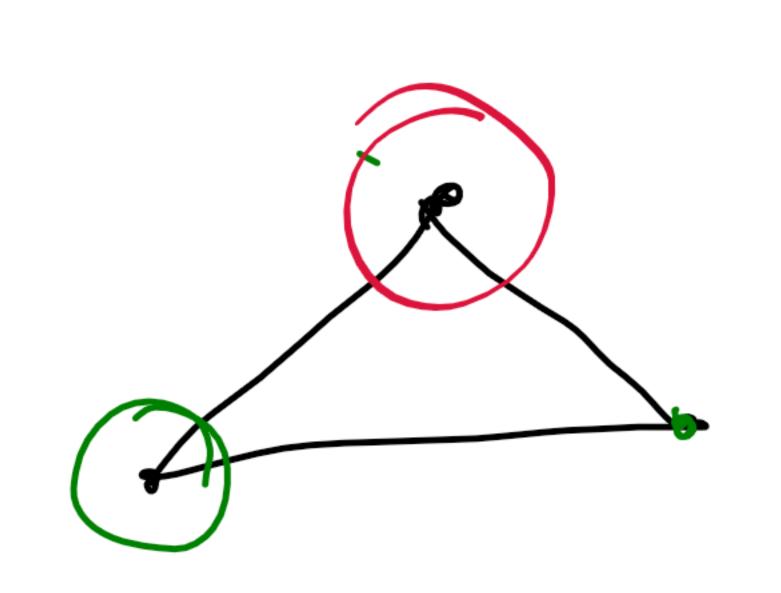
* Kp is $(P-1) - \Re egalax$ geaph * No of edger in Kp = ? 2q = Sum of deg $= (P-1) + (P-1) + \cdots + (P-1)$

(A) Bipartite glaph:

A bigeaph & a bipartite geaph'is the one whose verlex set V(6) can be partitioned into two disjoint subsets VI from such that every edge of the geaph joins the verlex of V, with verlies of V2



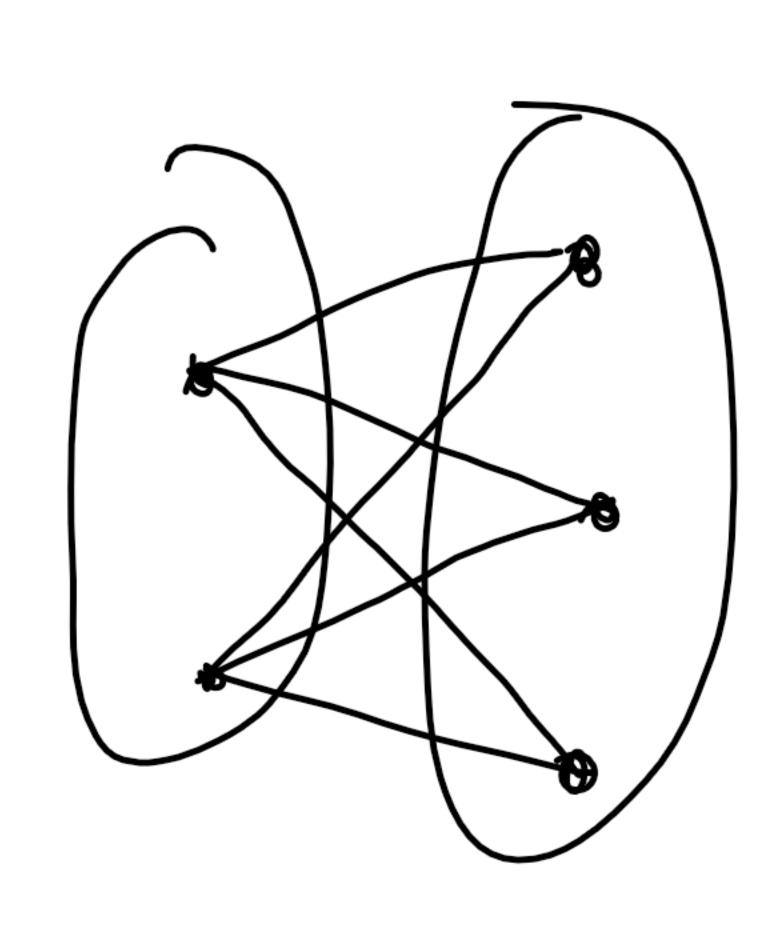




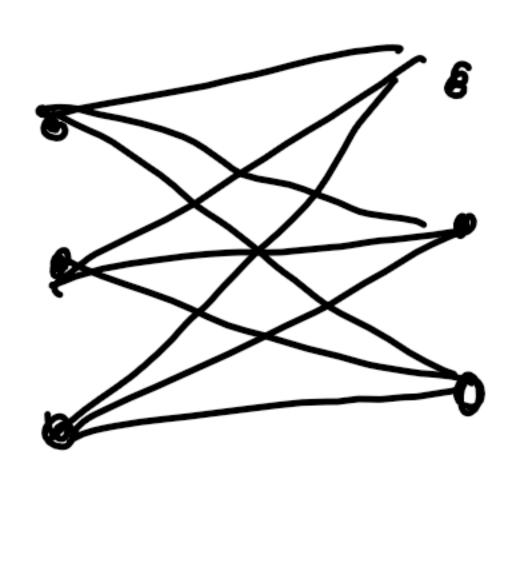
Not abipastite glaph

complete bipartite geaph: - 9f the graph contains every possible edges soining vertices of V_1 with vertices of V_2 $|V_1| = p$, $|V_2| = q$, $\longrightarrow K_{p,q}$

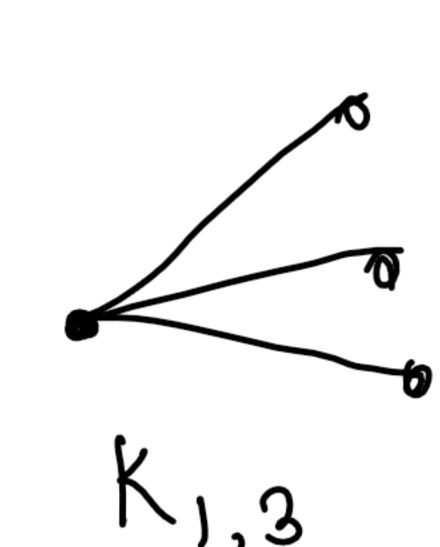
K 2,3

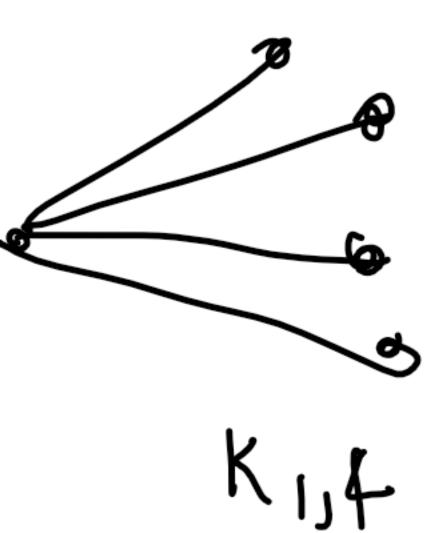


K3,2

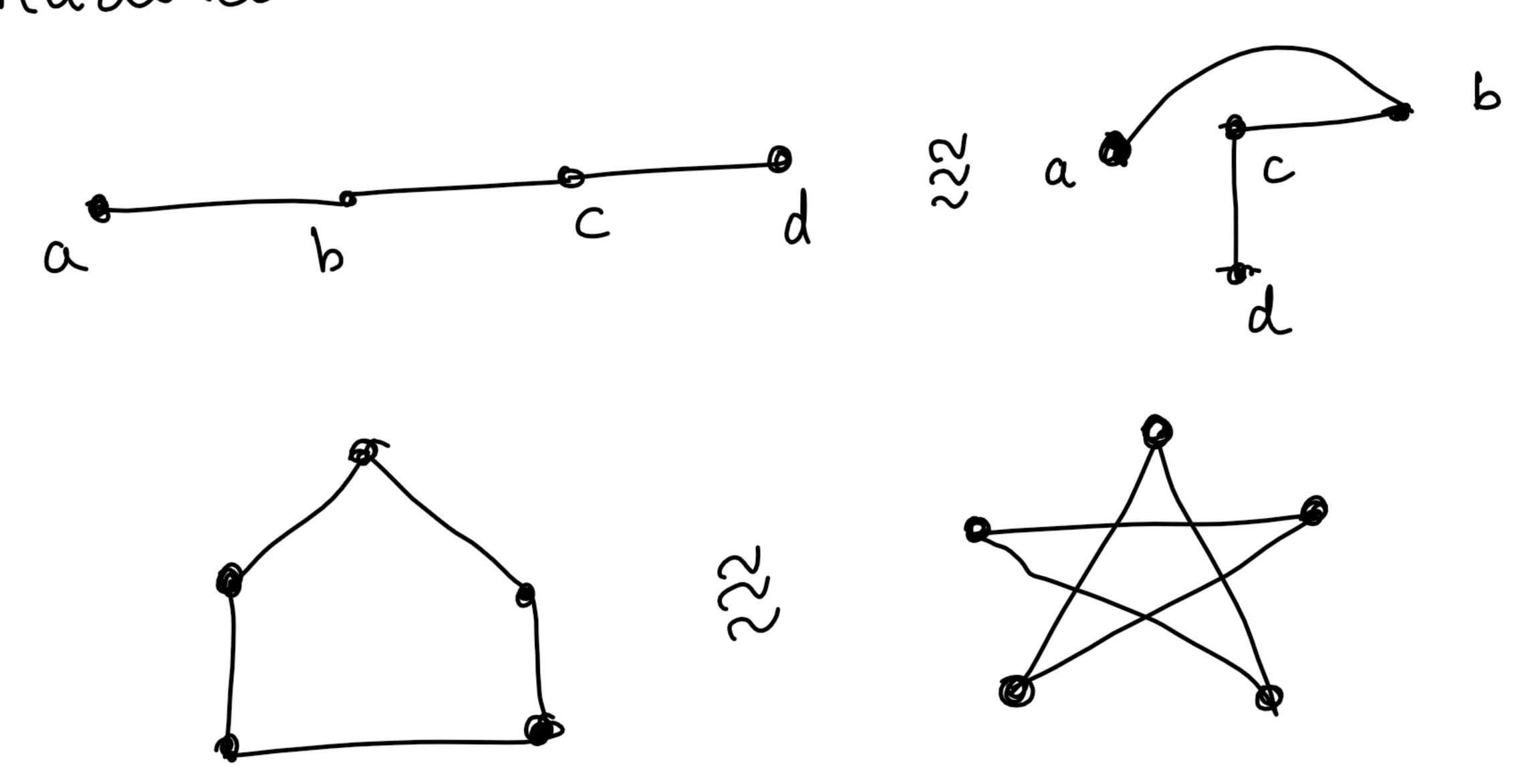


stal geaph's K_{1,p}



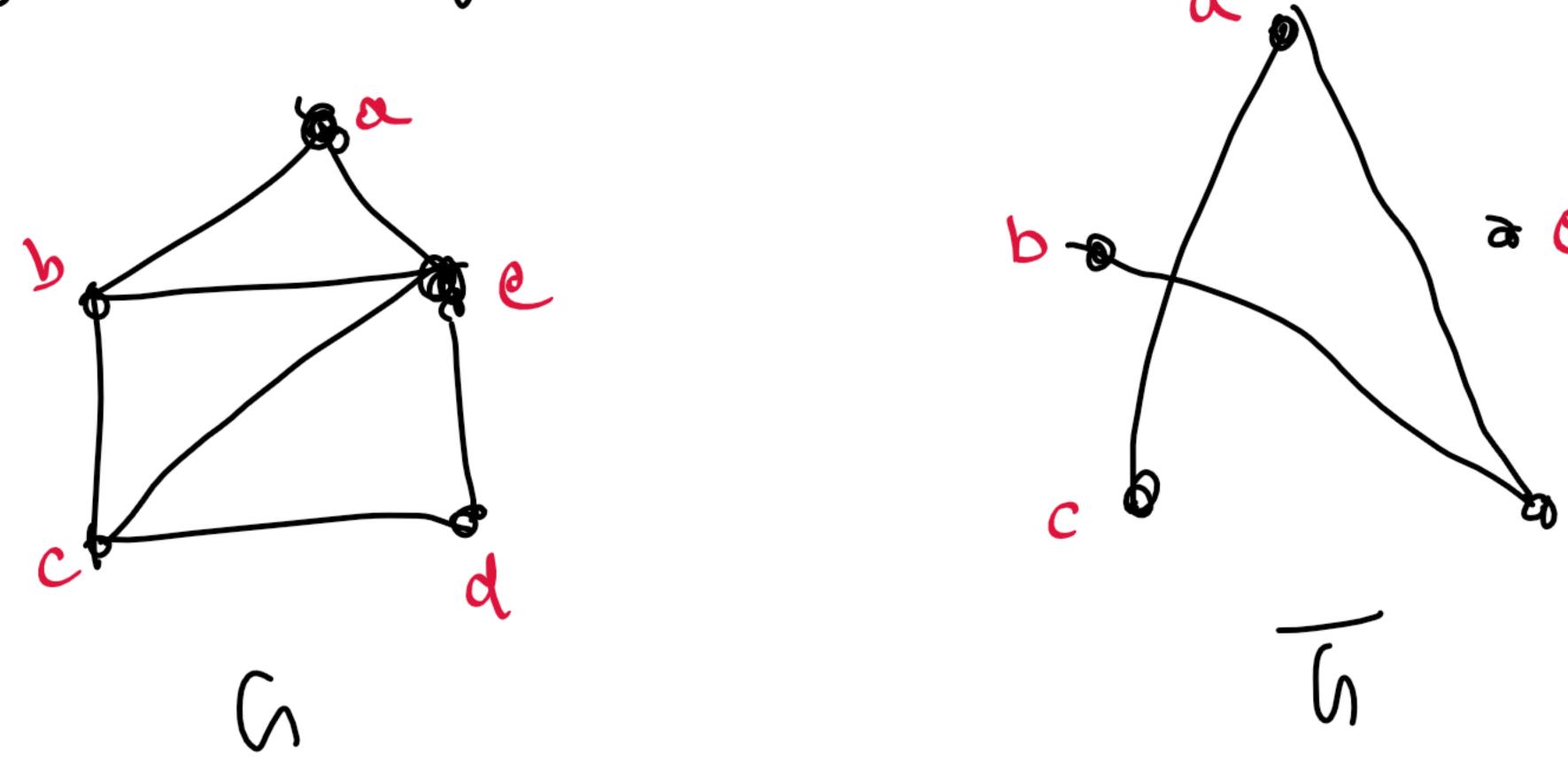


Two graphs are isomorphic: Two graphs hid ha are isomorphic if there is one-to-one correspondence blum the vertices of blum the edges preserving the incidence.

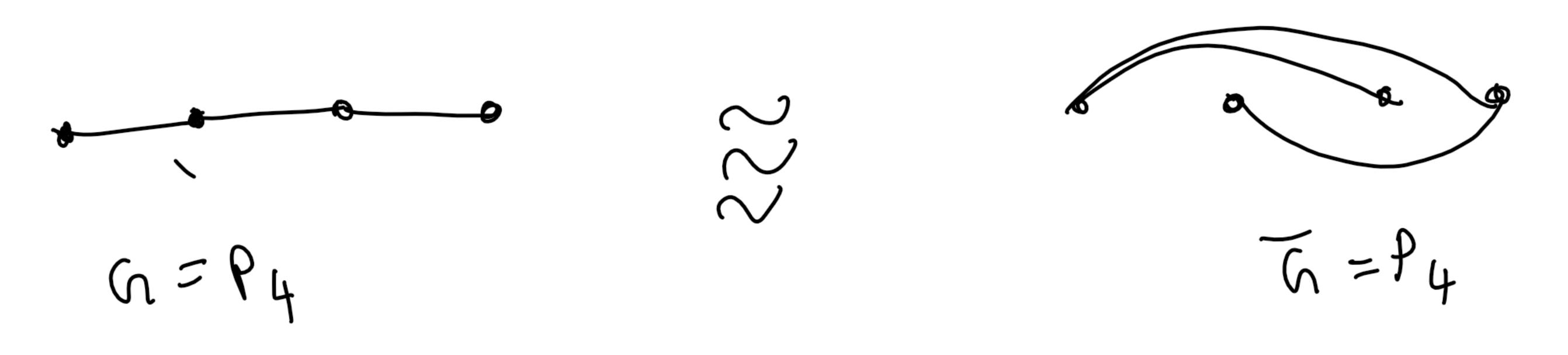


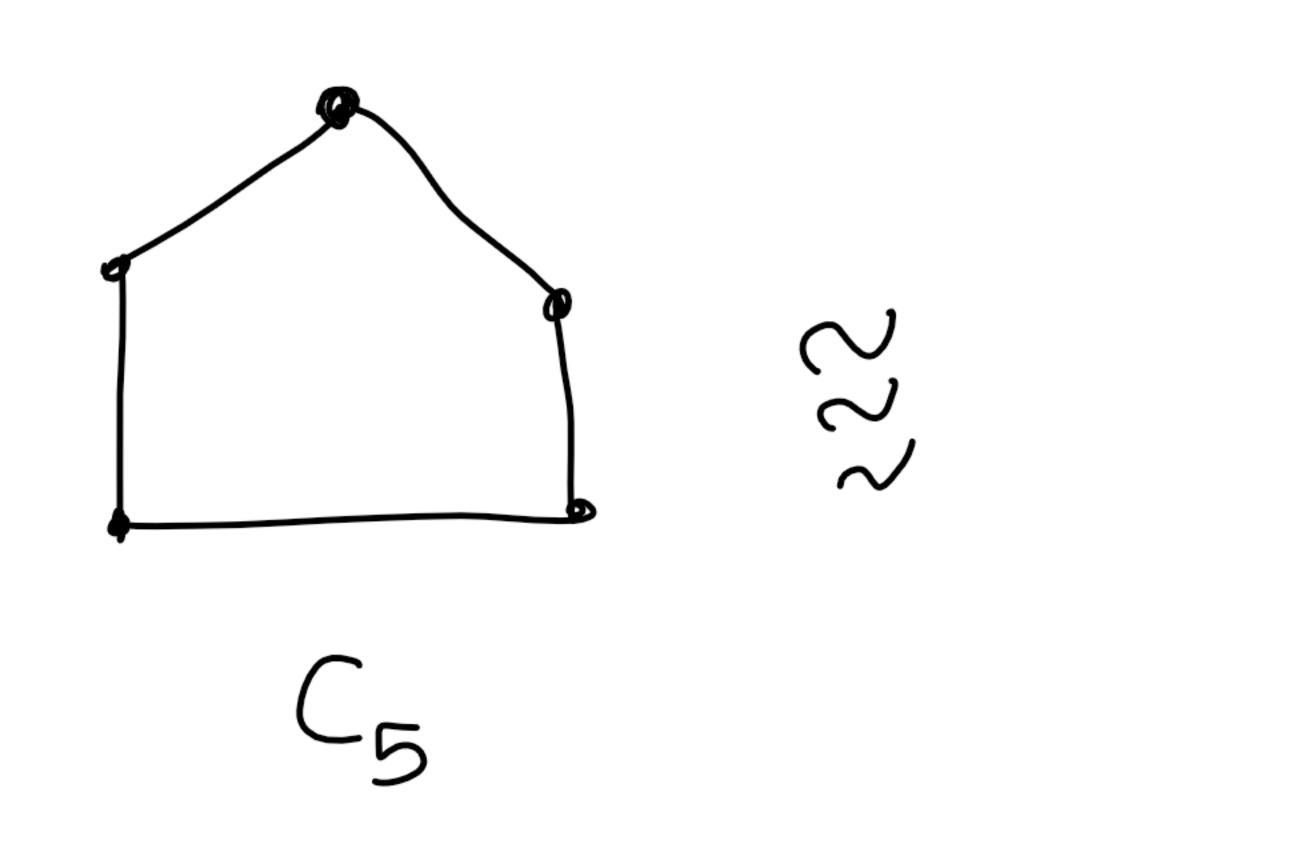
Complement of a geaph

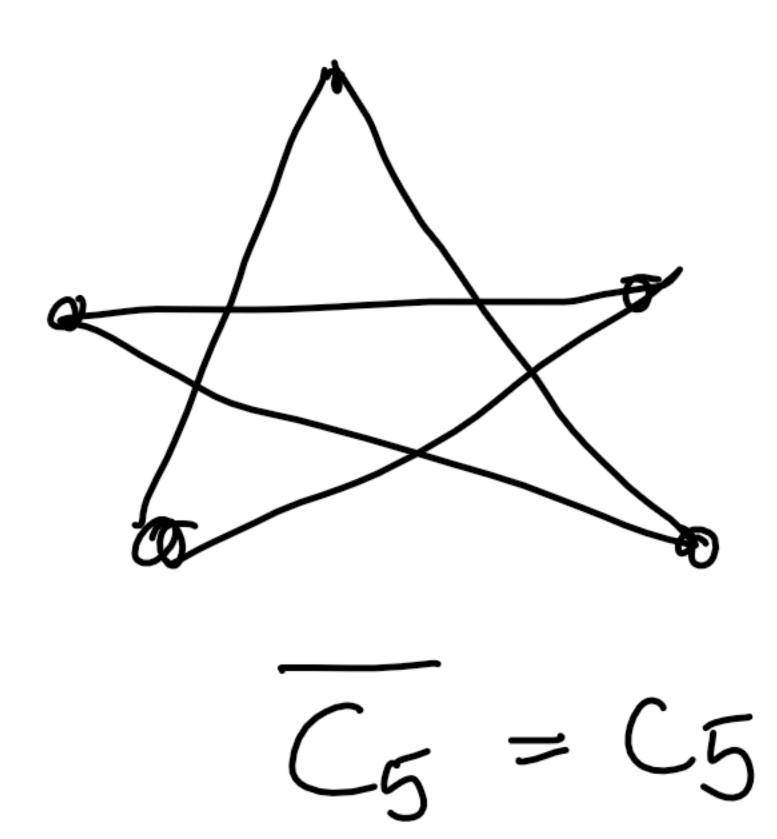
The complement of G, denoted by To is a glaph with V(G) = V(G) and any 2 vertices in G are adjacent if they are not a djacent in G



Self complementaly graph: - if 5 is isomophic to a







Py -> smallest self comp geoph