



Solution Review: Find Two Pairs Such That $a+b = c+d$

This review provides a detailed analysis of the solution to the Find Two Pairs Such That $a+b = c+d$ Challenge.

We'll cover the following

- Solution: Sums Stored as Hash Keys
- Time Complexity

Solution: Sums Stored as Hash Keys

```
1 def find_pair(my_list):
2     result = []
3     # Create a dictionary my_dict with the key being the sum
4     # and the value being a pair, i.e key = 3 , value = {1,2}
5     # Traverse all possible pairs in my_list and store sums in my_dict
6     # If sum already exists then print out the two pairs.
7     my_dict = dict()
8     for i in range(len(my_list)):
9         for j in range(i+1, len(my_list)):
10            added = my_list[i] + my_list[j] # calculate sum
11            # the 'in' operator on dict() item has a complexity of O(1)
12            # This is due to hashing
13            # On a list, the 'in' operator would have the complexity of O(n)
14            if added not in my_dict:
15                # If added is not present in dict then insert it with pair
16                my_dict[added] = [my_list[i], my_list[j]]
17            else:
18                # added already present in the dictionary
19                prev_pair = my_dict.get(added)
20                # Since list elements are distinct, we don't
21                # need to check if any element is common across pairs
```

```

21         # need to check if any element is common among pairs
22         second_pair = [my_list[i], my_list[j]]
23         result.append(prev_pair)
24         result.append(second_pair)
25         return result
26     return result
27
28

```



Each element in `my_list` is summed with all other elements one by one and the pair is stored. The sum becomes the key in the `my_dict` dictionary. At every key, we store the integer pair whose sum generated that key.

Whenever a sum is found such that its key in the dictionary already has an integer pair stored in it, we can conclude that this sum can be made by two different pairs in the list. These two pairs are then returned in the `result` list.

Time Complexity

The algorithm contains a nested loop. However, the inner loop always starts one step ahead of the outer loop:

```

for i in range(len(my_list)):
    for j in range(i+1, len(my_list)):

```

As the outer loop grows, the inner loop gets smaller. First, the inner loop runs $n - 1$ times, then $n - 2$, and so on...

This is an arithmetic series.

After evaluating the series for these values, the time complexity of this algorithm is $O(n^2)$.



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