Virtualising the d-invariant

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Contents

Introduction

Introduction...

very draft In order to keep this text relatively self-contained, we begin in Chapter 1 with the study of knots, for without the bread and butter of knot theory, knot invariants. This provides some motivation for the rest of this chapter in which we slowly build our way up to a specific invariant of alternating knots, the lattice of integer flows of the Tait graph, closely following the work of Greene [Gre11]. We then look at the formulation given by Greene, the d-invariant and how that relates to Heegard-Floer homology and another formulation of the d-invariant. In Chapter 2, we examine virtual knots, a generalisation of knots with a several equivalent formulations that allows knots to have diagrams on an orientable surface of any genus. The aim of this is to extend the lattice of integer flows to the virtual setting, and examine what properties it is able maintain in this new environment. We suspect the virtual d-invariant is less powerful than another invariant known as the Gordon-Litherland Linking Form, and we examine their relationship in Chapter 3, and in Chapter 4we discuss how Gauss codes provide a way for computers to deal with knots. In Chapter 5 we examine an algorithm to compute these invariants, and present a proof that indeed, the virtual d-invariant is not as strong as the Gordon-Litherland Linking Form, and therefore not a complete mutation invariant of alternating knots.

Acknowledgements

Thanks to ...

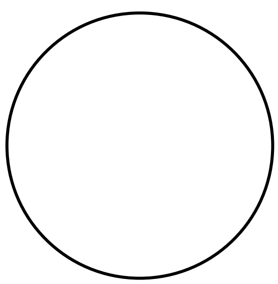
We would like to thank Hans Boden for an insight into the slick proof of the proposition in Chapter 5.

Chapter 1

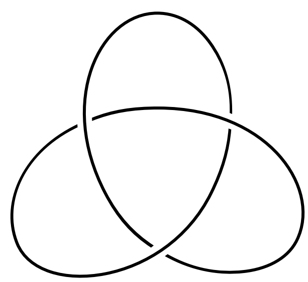
Knots and the Lattice of Integer Flows

1.1 Knots and Knot Invariants

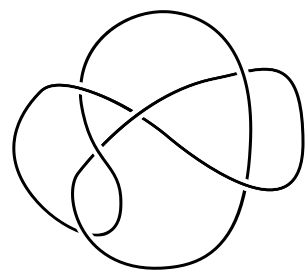
We begin with a swift introduction to the rich and marvellous study of tangled-up pieces of string: the theory of knots. Despite being a complex and intricate field, any child can intuitively grasp the concept of a knot as a closed loop of string sitting in space. To formalise this and remove any pathological examples that are inconsistent with the intuition of our inner child, we define a *knot* to be an injective embedding of the circle into 3-space, $K : \mathbb{S}^1 \hookrightarrow \mathbb{R}^3$. The requirement that the embedding be injective ensures that the string does not intersect itself.



(a) Unknot or Trivial Knot



(b) Trefoil Knot



(c) The knot 6_3

Figure 1: Some examples of knots, presented through knot diagrams.

The example in Fig. 2 shows a knot that can be 'untwisted' to look like the example in Fig. 1a. In general, we do not distinguish between knots if there is some way to deform one, without breaking the circle or passing it through itself, into the other. Hence we say two knots K_1 and K_2 are *equivalent* or equal if there is some homeomorphism of the ambient space $\mathbb{R}^3$ that restricts to a homeomorphism of the knots. We call this notion of equivalence *ambient isotopy*.

Though a central objective of knot theory it to classify knots up to ambient isotopy, as one might imagine, it can be very hard to write down explicit ambient isotopies directly. In practice we find ourselves using other tools to get the job done. Though we claimed that the objects in Fig. 1 were knots, they are really projections of knots onto the plane with markings we call *crossings* to make clear that some *strand* of the knot passes over another strand. We refer to these objects as *knot diagrams* which gives us the terminology we need when we have multiple knot diagrams that represent the same knot such as in Fig. 2.

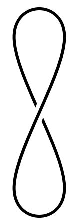
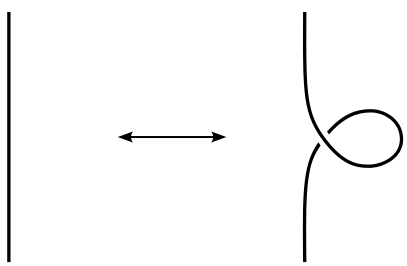


Figure 2: Another diagram for the unknot.

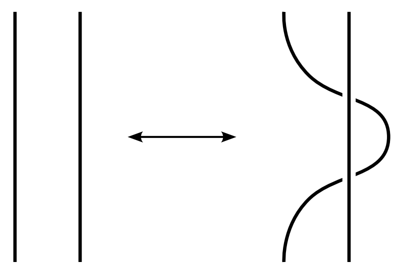
We consider diagrams equivalent under *planar isotopy*: **define**. **It is not obvious when diagrams are equivalent; there are many not-easy-to-tell-are-simple unknot diagrams.** The unknot is not special here – every knot has many diagrams, in fact infinitely many. For brevity we may say 'knot' when we should more strictly have said 'knot diagram'. When we wish to emphasise that we are talking about the knot rather than any specific diagram we refer to the 'knot type'.

Diagrams are much easier to work with than embeddings, and we have the following foundational theorem due to Alexander-Briggs [AB26] and independently Kurt Reidemeister [Rei27].

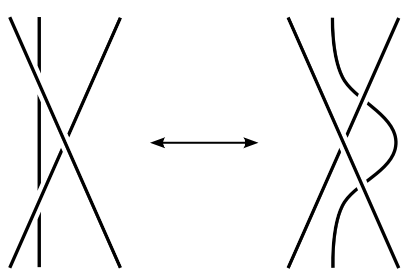
Theorem. *(The Reidemeister Theorem) Two knot diagrams D_1 and D_2 represent equivalent knots if and only if there is some sequence of finitely many moves of any of the types given in Fig. 3 that transform D_1 into D_2.*



(a) RI



(b) RII



(c) RIII

Figure 3: The Reidemeister moves.

The early study of knot theory by pioneers such as P.G. Tait, C.N. Little and T. Kirkman involved trying to find, by hand, some way to show two knots were equivalent. But Tait himself noted that is was impossible by these means alone to ever prove that two knots were distinct. The modern way we do this is by using knot invariants. If we define some map from knot diagrams to some other class of objects, perhaps a truth value, an polynomial, or a group, and can show that none of the Reidemeister moves,the value of this map, then we have a well-defined function on knots – a *knot invariant*. Finally we have a way to prove that two knots are different, for if they take on different values under some invariant they must be distinct. However, invariants are one-sided in nature – taking different values can tell us that two knots are different, but two knots taking on the same value of some invariant doesn't necessitate that they be equivalent knots. An invariant that is an injection from the class of knots is called a *complete* invariant, and while we now know of a zoo of different invariants, we have yet to find what is perhaps the Holy Grail of knot theory: a complete invariant that is also easy to compute.

1.2 Alternating Knots, Knot Mutation and the Tait Graph

We call a knot diagram *alternating* if, traversing the diagram, the crossings alternate under and over. Clearly, for every diagram that is alternating, it is possible to construct a non-alternating diagram of an equivalent knot, one simply needs to apply a type RI Reidemeister move appropriately to any strand of the knot **add figure**. Hence we define an *alternating knot* as a knot that *has* as alternating diagram. Alternating knots are a particularly interesting class of knots. For low crossing-number **define above**, many of the knots are alternating, but this trend quickly reverses as crossing number is increased. Alternating knots have a special connection with a series of conjectures made by Tait in his early attempts to tabulate knots. The most important of these conjectures, in the sense that is implies the others, is known as the flyping conjecture and relates alternating diagrams by moves known as flype moves.

A *tangle* in a knot diagram is a region of the plane that is homeomorphic to a disc, such that the knot crosses the boundary of the disc exactly four times, as in Fig. 5a and Fig. 5b. A *flype* move is a diagrammatic move that flips a tangle but does not alter the knot type, as in Fig. 4.

Tait's Flyping conjecture was proven by Menasco and Thistlethwaite in [MT93] as the following theorem.

Theorem (Tait Flyping Conjecture)**.** *Any two reduced alternating diagrams which present the same knot are related by a sequence of finitely many flypes, as seen in Fig. 4.*

This is like a Reidemeister theorem for alternating knots in that is relates equivalence of diagrams to the existence of a sequence of moves between them.

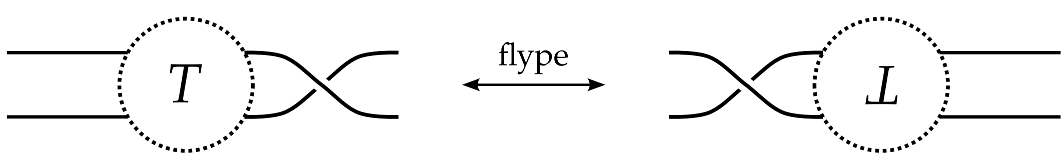
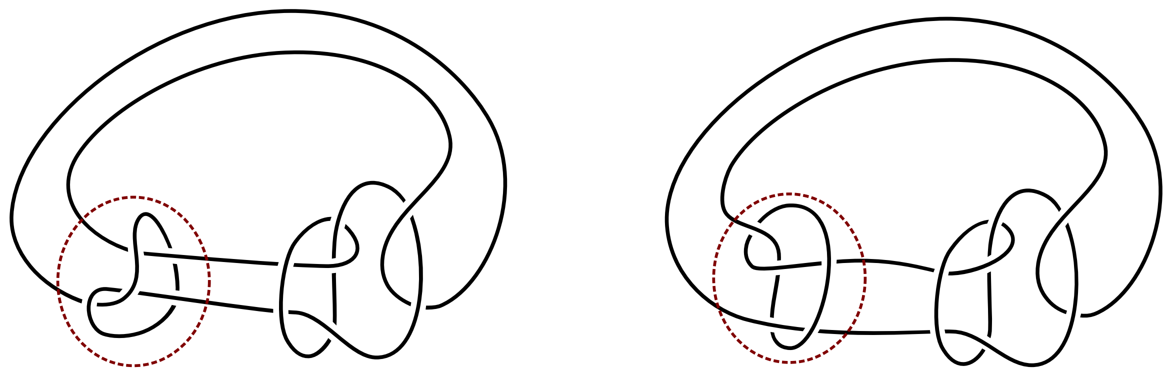


Figure 4: A general flype move.

We also have ways of constructing new knots from existing ones. If we take a diagram, choose a tangle, and then perform some reflection (up to planar isotopy) of that tangle, either reflecting it left-right or up-down, or across one of the diagonals, the corresponding operation on the knot is known as *mutation*, and the two knots known as *mutants*. Mutants are some of the hardest knots to distinguish, as many of their key invariants are the same. An example of this is the Conway Knot and the Kinoshita-Terasaka knot (Fig. 5).



(a) The Kinoshita-Terasaka Knot

(b) The Conway Knot

Figure 5: A famous pair of mutant knots known as the Kinoshita-Terasaka mutants. These knots are the two knots of lowest crossing number that have trivial Alexander polynomial, other than the unknot. The disk of mutation is marked. The projections used for these diagrams were taken from [Ada94].

To every knot diagram, we can associate a graph known as the *Tait graph* as follows. We interpret the knot diagram as a tetravalent planar graph that divides the plane into regions. An application of the Jordan curve theorem that it is possible to colour these regions two colours, black and white, such that adjacent regions are never the same colour. Such a colouring is called a *checkerboard colouring*. Note that in a checkerboard colouring, regions that are diagonal to each other at crossings are necessarily the same colour. To construct the black Tait graph, we place a vertex in every black region of the plane. Each crossing will connect a single pair of these vertices, and for each crossing we draw an edge between said vertices to obtain an undirected, planar graph, potentially with multiple edges.

The white Tait graph is constructed similarly from the vertices corresponding to white regions. Either of these graphs retains enough information to construct the other, as they are planar duals. That is, letting G_1 and G_2 be these graphs, if every face in a planar diagram for G_1 is replaced by a vertex, and edges between vertices in G_1 replaced by edges between the faces they separate, then G_2 has been constructed from G_1 and vice-versa. Hence, we sometimes refer the *the* Tait graph of a knot, as only one is necessary. Later we will examine a more general class of knotted objects for which this duality breaks down. **Add figures.**

If we have the Tait graph of an alternating knot, we can reconstruct the knot from it by **explain the procedure and why it only works for alternating knots**. **This is up to mirror image right?**

1.3 The Lattice of Integer Flows

For the rest of this chapter, we largely follow [Gre11] to introduce the lattice of integer flows and show that it is a complete mutation invariant of alternating knots. This means that it is both an invariant of knots, and among knots, it is both an invariant of mutation, and distinguishes knots that are not mutants. We then talk about the equivalent formulation given in [Gre11] as the d-invariant.

A *lattice* is a finitely generated abelian group L, equipped with an inner product $\langle \cdot, \cdot \rangle : L \times L \longrightarrow \mathbb{R}$. We are primarily interested in *integral lattices*, for which the inner product's image is contained within $\mathbb{Z}$, and for the rest of this text we assume that all lattices are integral. An *isomorphism of lattices* is a bijection $\psi : L_1 \longrightarrow L_2$ that preserves the inner product, that is, $\langle x, y \rangle = \langle \psi(x), \psi(y) \rangle$ for all $x, y \in L$.

Throughout, we let $G = (E, V)$ be a finite, directed, connected graph (in which loops and multiple edges are allowed) with vertex set V and edge set E. In particular, G is a 1-dimensional CW-complex and the boundary map $\partial : C_1(G) \longrightarrow C_0(G)$ is defined by the $|V| \times |E|$ incidence matrix $D : \mathbb{Z}^E \longrightarrow \mathbb{Z}^V$ with entries given by

$$D_{ij} = \begin{cases} +1 & \text{if } e_i \text{ is oriented into } v_j \\ -1 & \text{if } e_i \text{ is oriented out of } v_j \\ 0 & \text{otherwise.} \end{cases}$$

The *lattice of integer flows* of G is the group $\Lambda(G) = \ker D$, along with the inner product induced by the Euclidean inner product on $\mathbb{Z}^E$. Equivalently, $\Lambda(G)$ is the first homology group of G, with inner products taken in $C_1(G)$. While the lattice $\Lambda(G)$ may depend on the orientation of the edges in G, its isomorphism class does not, as the isomorphism class of the homology group is independent of orientation, and the Euclidean inner product is preserved by sending an edge to its negation: $\langle e_i, e_i \rangle = \langle -e_i, -e_i \rangle = 1$, and $\langle e_i, e_j \rangle = \langle -e_i, e_j \rangle = 0$ for $i \neq j$.

TODO: Put an example of a lattice of integer flows?

A 2-*isomorphism* between two graphs $G = (E, V)$ and $G' = (E', V')$ is a bijection $\psi : E \longrightarrow E'$ that preserves cycles, i.e. $\partial(e_i + \cdots + e_j) = 0$ if and only if $\partial\left(\psi(e_i) + \cdots + \psi(e_j)\right) = 0$. We call and edge e of a graph G a *bridge* if the removal of G from e disconnects G, and we say a graph G is 2-*edge-connected* if G has no bridges.

It is well established that a 2-isomorphism implies isomorphic lattices of integer flows; that is $\Lambda(G)$ is a 2-isomorphism invariant of 2-edge-connected graphs [BLN97]. More interestingly, and more recently, due to Su-Wagner [SW10, Theorem 1] and Caporaso-Viviani [CV10, Theorem 3.1.1], for 2-edge-connected graphs, the converse is also true.

Theorem. *For two* 2*-edge-connected graphs G and G', $\Lambda(G) \cong \Lambda(G')$ if and only if G and G' are* 2*-isomorphic. That is, $\Lambda(G)$ is a complete* 2*-isomorphism invariant of* 2*-edge-connected graphs.*

Two graphs G and G' are related by a *Whitney flip* if it is possible to find two disjoint graphs Γ_1, with distinguished vertices u_1 and v_1 and Γ_2 with distinguished vertices u_2 and v_2, such that the identifications $u_1 = u_2 = u$ and $v_1 = v_2 = v$ form G, and the identifications $u_1 = v_2 = u'$ and $v_1 = u_2 = v'$ form G'. An example of graphs related by a Whitney flip is given in Fig. 6. **Is this only for planar graphs? Don't think so but check.**

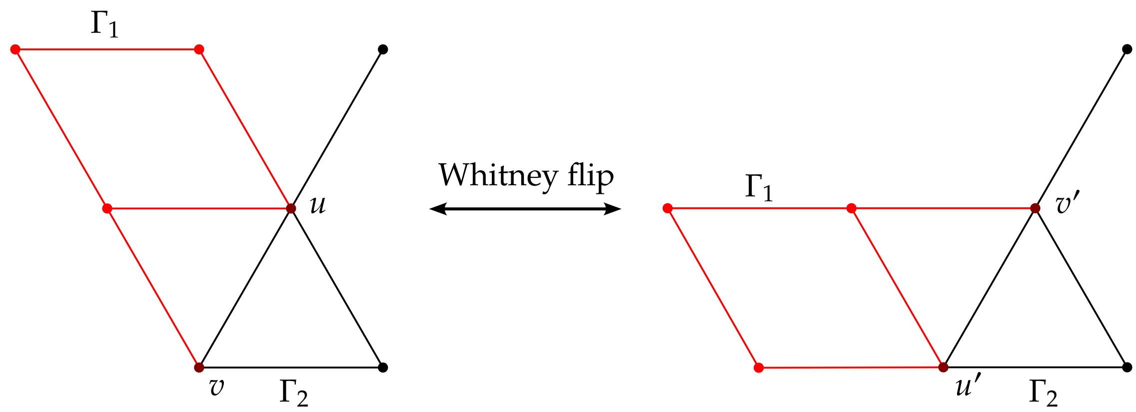


Figure 6: An example of a Whitney flip.

It is clear that sequences of Whitney flips only ever transform graphs within their 2-isomorphism class, as cycles map to cycles. From [Whi33], we have the important converse; a Reidemeister-like theorem for 2-isomorphic graphs.

Theorem (Whitney's Theorem)**.** *Two graphs G and G' are 2-isomorphic if-and-only-if there is a sequence of Whitney flips relating G to G'.*

It can be easily shown that flype moves correspond to Whitney flips of the Tait graphs of the knot. Since all Tait graphs of reduced **define earlier** alternating knots are 2-edge-connected, it is an immediate consequence of Whitney's theorem is that the pair $\Lambda(G_1), \Lambda(G_2)$, where G_1 and G_2 are the Tait graphs of the knot is an invariant of alternating knots.

Greene proves that a mutation of alternating knot diagrams induces at worst a Whitney flip on the the Tait graph, and and a Whitney flip on the Tait graph induces a mutation on the diagram. This completes the proof of the following theorem (Proposition 4.4 in [Gre11]).

Theorem (Greene)**.** *For a knot K, the pair $\Lambda(G_1), \Lambda(G_2)$ of lattices of integer flows of the Tait graphs, which we denote as $\Lambda(K)$ is a complete mutation invariant of alternating knots. That is,*

(1) $\Lambda(K)$ is an invariant of alternating knots, and

(2) Alternating knots A_1 and A_2 are mutants if and only if $\Lambda(A_1) \cong \Lambda(A_2)$.

Give some notes about what it means to have an 'isomorphism of the pair' $\Lambda(G_1), \Lambda(G_2)$ for G_1, G_2 Tait graphs of D.

1.4 The d-invariant and Heegard-Floer Homology

Show that can be compressed into the d-invariant, and nothing is lost.

Chapter 2

Virtual Knots

We now introduce the exciting and relatively new theory of virtual knots. Virtual knots are a generalisation of knots, and there are many different equivalent formulations of them. We start with the most geometric of the formulations, but we also present a combinatorial and computational definition later.

2.1 Knots in Thickened Surfaces

Classical knots, a term which refers specifically to the kind of knots we have introduced prior to this chapter, have diagrams in the plane, $\mathbb{R}^2$, but really they have an extra dimension of 'thickness', encoded in the diagram by the under- and over- crossings. Hence we think of knots as embeddings in $\mathbb{R}^3$. However we didn't really need a whole $\mathbb{R}$'s worth of extra space. We could easily think of classical knots as living in a thickened plane, $\mathbb{R}^2 \times I$ where I is the unit interval $[0,1]$. Thinking of knots as embeddings in $\mathbb{R}^2 \times I$, it becomes natural to ask: what if we replace the plane by another surface; can we have diagrams on other surfaces Σ and therefore knots in *thickened surfaces* $\Sigma \times I$? The answer to these questions is yes, and virtual knots are one such generalisation.

In the context of virtual knots, all surfaces of relevance are closed and orientable. **Except the plane which we identify with the sphere? How do we justify this?** (Note that as we are thinking of surfaces as 2-manifolds embedded in $\mathbb{R}^3$, closed is equivalent to compact.) The classification theorem for compact, orientable surfaces is the following.

Theorem (Classification of compact, orientable surfaces)**.** *Each connected component of a compact, orientable surface is homeomorphic to:*

- *the sphere, or*
- *a connected sum of g tori, for $g \geq 1$.*

Hence there is a bijection between connected components of compact, orientable surfaces given by the *genus*, g of the surface, the number of handles.

We now follow the work of Kuperberg [Kup03] and Carter-Kamada-Saito [CKS00] and give the geometric definition of virtual knots. A *surface knot diagram* on Σ is the analogue of a classical knot diagram, but drawn on a general closed, oriented, connected surface, Σ no longer necessarily the plane. To represent diagrams in this section we use the letter P, for 'projection', another common word for 'diagram', as D will be reserved for disks.

The equivalence relation on surface knot diagrams is given by the Reidemeister moves and surface isotopy (the analogue of planar isotopy) on Σ.

We can define knots on thickened surfaces so that they relate to surface knot diagrams, in the same way knots and knot diagrams are related in the classical context. A *knot in a thickened surface* $\Sigma \times I$ is an embedding $K : S^1 \hookrightarrow \Sigma \times I$ up to ambient isotopy in $\Sigma \times I$.

To define virtual knots in thickened surfaces of any genus under another equivalence relation: stable equivalence, defined based on the following two operations. We give the definitions on the level of surface knot diagrams, but note that they generalise to knots in thickened surfaces accounting for the extra factor of I. The operation of *stabilisation* consists of finding two disks D_1 and D_2 in Σ that do not intersect P. We then remove D_1 and D_2 from Σ and glue a handle whose boundary is $D_1 \cup D_2$. Intuitively, stabilisation is 'adding a handle' to Σ, and any newly added handle does not interact with P. The reverse operation *destabilisation* removes a handle. Performing this operation, we take a cylinder Y that does not intersect P, and such that the circle that Y deformation-retracts to is not null-homologous, removes it, and cap both resulting boundary circles. The point of stable equivalence is to identify two knot diagrams P_1 and P_2 that would otherwise be identical, except for living on surfaces Σ_1 and Σ_2 of different genus g_1 and g_2; for one surface must have a greater genus than the other $g_1 < g_2$. But since the diagrams are identical, there must be at least $g_2 - g_1$ handles in the diagram on Σ_2 that are unnecessary and don't interact in any way with the diagram P_2.

Hence we define a *virtual knot* as an equivalence class of knots on thickened surfaces under stable equivalence. Each virtual knot has a minimum genus surface, in which all of the handles interact with the the knot, and this is known as its virtual genus.

The virtual knots with virtual genus $g_v = 0$ correspond to the classical knots, and the virtual knots with virtual genus $g_v < 0$ are *strictly* virtual. An example of a strictly virtual knot is given in Fig. 7.

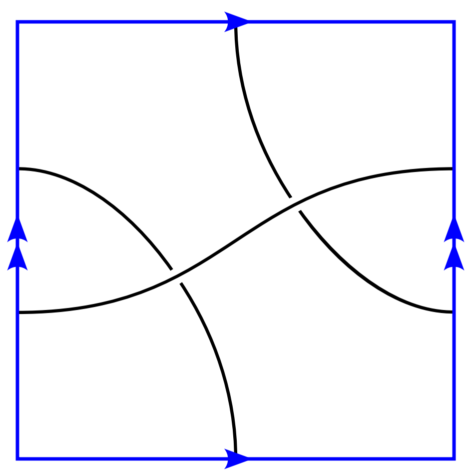


Figure 7: Virtual knot 2_1, a strictly virtual knot that is not checkerboard colourable. This knot is drawn on the gluing diagram of the torus, $\mathbb{T}^2$, its thickened surface of minimal genus.

Instead of drawing virtual knots directly on compact orientable surfaces, for genus $g_v \geq 1$, we often draw them on the gluing diagrams of those surfaces. From a textbook theorem of algebraic topology [Hat00], the compact orientable surface of genus g is obtained from the $4g$-gon by gluing around the polygon with the pattern $aba^{-1}b^{-1}$, $cdc^{-1}d^{-1}$, $\cdots$ continuing on for g iterations, or until all edges have been glued. Figures 7 and 10 are examples of this, with the surfaces being the torus, $\mathbb{T}^2$ and the compact orientable surface of genus 2, $2\mathbb{T}^2$, respectively.

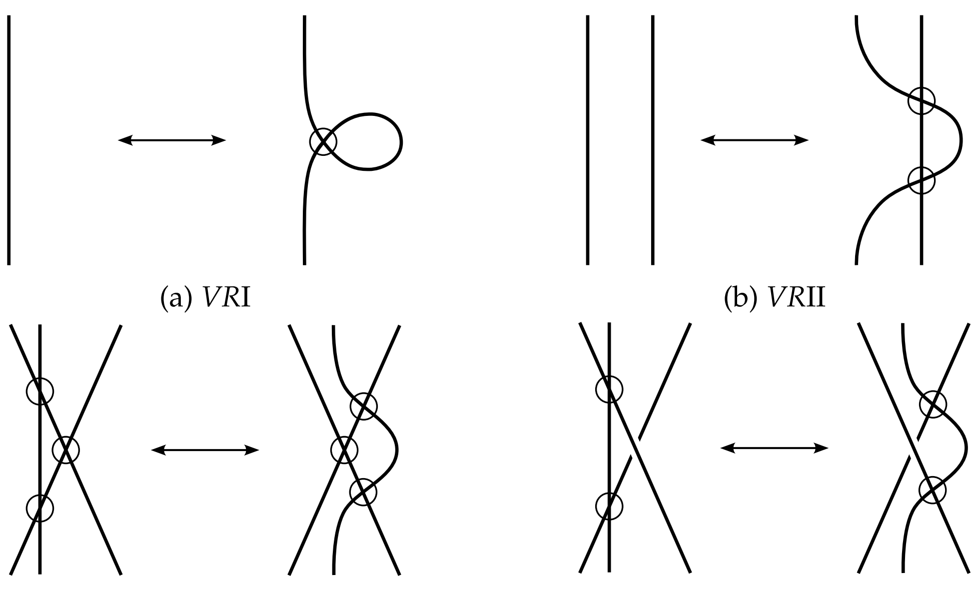
There is no fast algorithm to compute the virtual genus of a virtual knot, the following theorem from [Man13] is useful:

Theorem. *If a diagram of a knot K in a thickened surface $\Sigma \times I$ contains the minimal number of crossings across all diagrams in thickened surfaces, then Σ is the minimal genus surface that can support K. That is, the genus of Σ is the knot's virtual genus.*

Applying this theorem to the classification of virtual knot diagrams up to 6 crossings by Jeremy Green [Gre04] we can determine the virtual genus of virtual knots up to 6 crossings. **More about this computation will be explained in Chapters 4 and 5.**

2.2 Knots with Virtual Crossings

The original formulation of virtual knots (and their discovery) is due to Kauffman in 1996 [Kau99]. This formulation can be related to the formulation of equivalence classes of diagrams on Σ by projecting the surface Σ the onto the plane. Doing this creates two types of crossings. Those that did actually come from a crossing on Σ, we call *classical crossings*, and they have the usual over- and under- strands as determined by the projection. Those that did not exist on Σ but rather are an artefact of the projection we call *virtual crossings*. For strictly virtual knots these virtual crossings will be necessary, as the tetravalent graph that the knot represents in not planar.



(a) *VR*I (b) *VR*II (c) *VR*III (d) *VR*IV

Figure 8: The four additional virtual Reidemeister moves.

The relevant equivalence relation on diagrams with virtual crossings are not hard to deduce. We have the three Reidemeister moves which still hold between classical crossings, three corresponding moves similar to the Reidemeister moves but with all classical crossings replaced by virtual crossings, and finally a 'triangle move' that moves a 'virtual strand' through a crossing.

There is yet another interpretation of virtual knots that we explore in this paper, a computational definition that is integral to computing invariants of virtual knots. We will explore this later a later chapter.

Scaffold

- **The main interpretation we are concerned with is knots in $\Sigma \times I$.**
- **The corresponding definition of mutation is disk-mutation of a diagram on Σ. Defining mutation in this way makes sense and translated back up into $\Sigma \times I$. What is the other definition of mutation? May be worthwhile to mention why surface mutation is not relevant.**
- **Tait graphs are different for virtual knots - only exist if the knot is checkerboard colourable or 'checkerboard'. This has something to do with being Alexander mod numerable and/or $\mathbb{Z}/2$ homologous?**
- **But all alternating knots are certainly checkerboard colourable. Nice little homology proof?**
- **Proof using Kindred Virtual Flyping '22 that lattice of integer flows is still a disc mutation invariant of alternating virtual knots.**
- **But is it still complete? We suspect that an invariant we know to be stronger is complete, which would mean no. We talk about that invariant in the next chapter.**

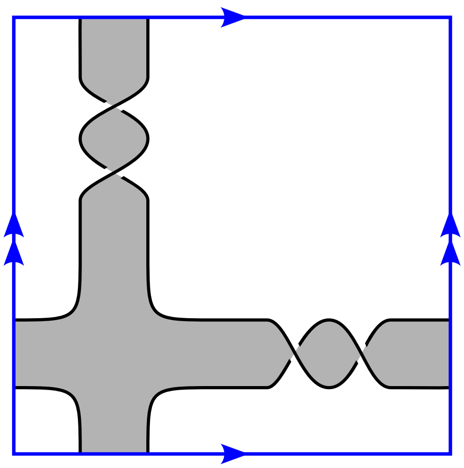
2.3 Mutation and Tait Graph(s) of Virtual Knots

Having introduced virtual knots, the rest of this chapter will be focussed on constructing a generalisation of the d-invariant of [Gre11], or rather, the equivalent lattices of integer flows of the Tait graphs of the knot. First we must introduce the tools we used to define this invariant in Chapter 1 into this new virtual setting. We focus on the definition of virtual knots as knots on thickened surfaces under stable equivalence.

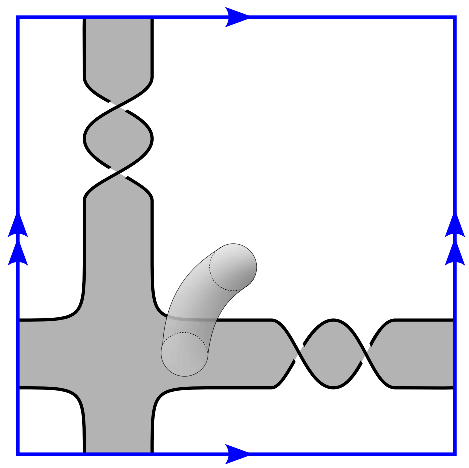
There are two types of mutation of virtual knots: disk mutation and surface mutation. *Disk mutation* is directly analogous to mutation of classical knots. We take a disk $D \subseteq \Sigma$ that contains a tangle and flip or rotate it, and the resulting knot is a disk mutant of the original. This is in contrast to *surface mutation*, in which the chosen subset need only have circular boundary, but could contain handles. Surface mutation is a more invasive operation and we do not consider it in the present text. However, future work may lie in investigating the equivalence classes generated by surface mutation. For the rest of this text, mutation, in the virtual context, refers to disk mutation.

In the classical case, all knots were checkerboard colourable, so the Tait graph could always be produced. Furthermore, only one Tait graph was necessary to encode an alternating knot, as Tait the two graphs were always planar duals to each other. In the virtual case, neither of these facts hold.

If we take a knot in a thickened surface Σ that is checkerboard colourable, such as that in Fig. 9a, it is always possible to put it on a new surface Σ' on which it will not be checkerboard colourable. We do this in Fig. 9b by adding a handle between a black region and a white one, unifying them. The resulting region now needs to be both black and white, so the knot on Σ' is not checkerboard colourable. Hence when trying to checkerboard colour virtual knots, the surface needs to be taken into account.



(a) Virtual knot 4_{105}



(b) A handle is added

Figure 9: Adding a handle can disrupt checkerboard colourability.

We say a knot in a thickened surface is *checkerboard colourable* if there is a checkerboard colourable surface knot diagram that represents it. We have the following classification of checkerboard colourable knots in thickened surfaces [BK19].

Theorem. *Given a knot in a thickened surface, $K \subset \Sigma \times I$, the following are equivalent:*

(i) K is checkerboard colourable,

(ii) K is the boundary of an unoriented spanning surface $F \subset \Sigma \times I$,

(iii) $[K] = 0$ in the homology group $H_1(\Sigma \times I; \mathbb{Z}_2)$.

Likewise, we say a virtual knot is *checkerboard colourable* if there is a representative that is checkerboard colourable (as a knot in a thickened surface). Note that while Fig. 9b is indeed checkerboard colourable as a virtual knot, it is not checkerboard colourable as a knot in a thickened surface.

For the purposes of generalising the d-invariant, we need only concern ourselves with alternating virtual knots – those virtual knots that can be represented by alternating virtual knot diagrams. Fortunately the following theorem tells us that we will not run into any trouble [Kam02].

Theorem (Kamada's Lemma)**.** *Every alternating virtual knot is checkerboard colourable.* ***Check with Hans the correct statement of this lemma.***

Do we want to put the virtual flyping theorem here? Because without it how do we know we don't have flype moves so we don't have the concept of invariants of alternating virtual knots.

Hence for all alternating surface knot diagrams we can define the Tait graph.

2.4 The Virtual d-invariant

Chapter 3

Gordon-Litherland Linking Form

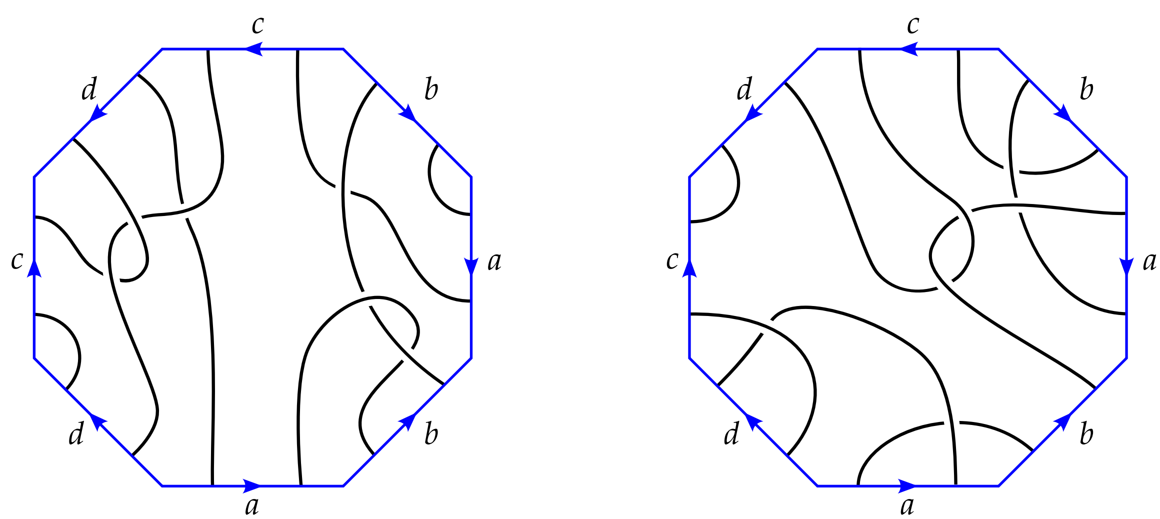
Chapter 4

Gauss Codes and Knot Algorithms

Work in pages already written about the algorithm to find the Tait graph.

Chapter 5

Computing Mock Seifert Matrices



(a) Virtual knot 6_{90101} (b) Virtual knot 6_{90124}

Figure 10: write

From the tables produced, it was noticed that the Kobayashi invariant [SK19], defined as

$$\operatorname{kob} A = \operatorname{tr}(A^\top A^{-1})$$

always satisfied the a relation involving the Alexander polynomial $\Delta_A(t) := \det(tA - A^\top)$ [BK23].

Proposition. *The Kobayashi invariant satisfies*

$$\operatorname{kob} A = -\frac{a_1}{a_0}$$

where a_i is the ith coefficient of the Alexander polynomial $\Delta_A(t)$.

Proof. The matrix A, being a mock Seifert matrix, has odd determinant [BK23] and is therefore invertible. Hence by the multiplicity of determinants,

$$\det(tA - A^\top) = \det(A)\det(tI - A^\top A^{-1})\,.$$

The last factor here is the characteristic polynomial of $A^\top A^{-1}$. It is a well known fact that the 0th and 1st coefficients of the characteristic polynomial of M, are $\det(M)$ and $-\operatorname{tr}(M)$ respectively. Hence,

$$-\frac{a_1}{a_0} = \frac{\operatorname{tr}(A^\top A^{-1})}{\det(A^\top A^{-1})},$$

but as $\det(A^\top A^{-1}) = \det(A^\top)\det(A^{-1}) = \det(A)\det(A^{-1}) = 1$, we have

$$-\frac{a_1}{a_0} = \operatorname{tr}(A^\top A^{-1})$$

as required. $\square$

This means that the Kobayashi invariant is weaker than the Alexander polynomial.

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Appendix A Algorithm