### Exercise 2

Due 24/11/2021, before 23:59

#### 1 Introduction

We continue our journey of building a compiler for the invented object oriented language  ${\bf L}$  . In order to make this document self contained, all the information needed to complete the second exercise is brought here.

### 2 Programming Assignment

The second exercise implements a CUP based parser on top of your JFlex scanner from the exercise 1. The input for the parser is a single text file containing a  $\bf L$  program, and the output is a (single) text file indicating whether the input program is syntactically valid or not. In addition to that, whenever the input program has correct syntax, the parser should internally create the abstract syntax tree (AST). Currently, the course repository contains a simple skeleton parser, that indicates whether the input program has correct syntax, and internally builds an AST for a small subset of  $\bf L$ . As always, you are encouraged to work your way up from there, but feel free to write the whole exercise from scratch if you want to. Note also, that the AST will not be checked in exercise 2. It is needed for later phases (semantic analyzer and code generation) but the best time to design and implement the AST is exercise 2.

# 3 The L Syntax

Table 1 specifies the context free grammar of  ${\bf L}$ . You will need to feed this grammar to CUP, and make sure there are no shift-reduce conflicts. The operator precedence is listed in Table 2.

# 4 Input

The input for this exercise is a single text file, the input L program.

### 5 Output

The output is a *single* text file that contains a *single* word. Either OK when the input program has correct syntax, or otherwise ERROR(*location*), where *location* is the line number of the *first* error that was encountered.

### 6 Submission Guidelines

The skeleton for this exercise can be found here. In your project repository, you need to add the relevant derivation rules and AST constructors. The makefile must be located in the following path:

• ex2/Makefile

This makefile will build the parser (a runnable jar file) in the following path:

• ex2/PARSER

Feel free to reuse the makefile supplied in the skeleton, or write a new one if you want to.

#### 6.1 Command-line usage

PARSER receives 2 parameters (file paths):

- *input* (input file path)
- output (output file path containing the expected output)

#### 6.2 Skeleton

You are encouraged to use the makefile provided by the skeleton. Some files of intereset in the provided skeleton:

- *jflex/LEX\_FILE.lex* (LEX configuration file)
- cup/CUP\_FILE.lex (CUP configuration file)
- src/Main.java
- src/AST/\*.java

To run the skeleton, you might need to install the following packages: \$ sudo apt-get install graphviz eog

To use the skeleton, run the following command (in the src/ex2 directory): \$ make

This performs the following steps:

• Generates the relevant files using jflex/cup

- ullet Compiles the modules into PARSER
- $\bullet$  Runs PARSER on input/Input.txt
- Generates an image of the resulting syntax tree (for debugging only)

## 7 Notes

- The terminals in the grammar are defined according to the token dfinitions of the first exercise. Note the additional token **TYPE\_VOID**.
- The skeleton creates an image file of the AST only for debugging, this should not be included in the final submission.

```
dec^+
Program
dec
                      varDec | funcDec | classDec | arrayTypedef
                ::=
                      TYPE_INT | TYPE_STRING | TYPE_VOID | ID
type
                ::=
                      ARRAY ID = type '['\ ']'
arrayTypedef
               ::=
varDec
                     type ID [ ASSIGN exp ] ';'
                ::=
                     type ID ASSIGN newExp ';'
                     type ID '(' [ type ID [ ',' type ID ]* ] ')' '{' stmt [ stmt ]* '}'
{\rm funcDec}
                ::=
classDec
                      CLASS ID [ EXTENDS ID ] '{' cField [ cField ]* '}'
                ::=
                      var
exp
                ::=
                     '(' exp ')'
                ::=
                    exp BINOP exp
                ::=
                      [ var '.' ] ID '(' [ \exp [ ', ' \exp ]^* ] ')'
                     ['-'] INT | NIL | STRING
                     ID
var
                ::=
                     var '.' ID
                ::=
                      var '[' exp ']'
                ::=
                     varDec
\operatorname{stmt}
                ::=
                     var ASSIGN exp';'
                ::=
                ::= var ASSIGN newExp ';'
                     RETURN [ exp ] ';'
                ::=
                     IF '(' exp ')' '{' stmt [ stmt ]* '}'
                ::=
                     WHILE '(' exp ')' '{' stmt [ stmt ]* '}'
                ::=
                     [var '.' ] ID '(' [exp [', 'exp ]* ]')' ';'
                ::=
                      NEW ID | NEW ID '[' exp ']'
newExp
                ::=
                     varDec | funcDec
cField
BINOP
                     + | - | * | / | < | > | =
                ::=
                      [1-9][0-9]^* \mid 0
INT
                ::=
```

Table 1: Context free grammar for the L programming language.

Precedence	Operator	Description	Associativity
1	:=	assign	
2	=	equals	left
3	<,>		left
4	+,-		left
5	*,/		left
6	[	array indexing	
7	(	function call	
8		field access	left

Table 2: Binary operators of  ${\bf L}$  along with their associativity and precedence. 1 stands for the lowest precedence, and 8 for the highest.