# Dan Ruley CS 3810 Study Assignment #9 December, 2019

## Proof-of-Study:

For this problem I coded my cache simulations in Java, corroborating and checking some of my logic by hand to make sure things work out well. I wrote three separate methods and a few helper functions to simulate Direct Mapped, Fully Associative, and Set Associative caches. I coded the basic ideas and then debugged a bit until I got the same results as Professor Jensen for the set of addresses he used for the practice problems. Before I get to that, here is a breakdown of how some example caches look for the three types:

### Fully Associative with 5 entries and 16 byte data blocks:

Total cache size: 5\*(128(data) + 12(tag) + 3(LRU) + 1(valid)) = 720 bits  $\Rightarrow$  120 bits left over. This is the cache architecture:

Valid(1bit)	Tag(12 bits)	Data(16 bytes)	LRU(3 bits)

### Addresses break down like this:

Tag bits	Offset bits
b15, b14, b13, b12, b11, b10, b9, b8, b7, b6, b5, b4	b3, b2, b1, b0

Memory access time:

1 cycle for a hit

10 + 1 + 16 = 27 cycles for a miss

# **Direct Mapped with 8 rows and 8 byte blocks:**

Total cache size: 8\*(64(data) + 10(tag) + 1(valid)) = 600 bits  $\Rightarrow$  240 bits left over.

This is the cache architecture:

Valid(1bit)	Tag(10 bits)	Data(8 bytes)

Addresses break down like this:

Tag bits	Row bits	Offset bits
b15, b14, b13, b12, b11, b10, b9, b8, b7,		
b6	b5, b4, b3	b2, b1, b0

Memory access time:

1 cycle for a hit

10 + 1 + 8 = 19 cycles for a miss

# Set Associative with 4 rows, 8 byte blocks, and 2 ways:

 $4*(2*(64(data) + 11(tag) + 1(valid) + 1(LRU)) = 616 \text{ bits} \Rightarrow 224 \text{ bits left over.}$ 

This is the cache architecture:

1	Valid(1bit)	Tag(11 bits)	Data(8 bytes)	LRU(1 bit)
ROW 0				
1				
1	Valid(1bit)	Tag(11 bits)	Data(8 bytes)	LRU(1 bit)
ROW 1				
\				
1	Valid(1bit)	Tag(11 bits)	Data(8 bytes)	LRU(1 bit)
ROW 2				
\				
1	Valid(1bit)	Tag(11 bits)	Data(8 bytes)	LRU(1 bit)
ROW 3				
1				

Addresses break down like this:

Tag bits	Row bits	Offset bits
b15, b14, b13, b12, b11, b10, b9, b8, b7, b6,		
b5	b4, b3	b2, b1, b0

And here is the state of the cache after the second pass through the addresses:

1	Valid(1bit)	Tag(11 bits)	Data(8 bytes)	LRU(1 bit)
ROW 0	1	3	d d d d d d d	0
1	1	4	x x x x x x x	1
1	Valid(1bit)	Tag(11 bits)	Data(8 bytes)	LRU(1 bit)
ROW 1	1	3	ylylylylylyly	0
\	1	4	z z z z z z z	1
1	Valid(1bit)	Tag(11 bits)	Data(8 bytes)	LRU(1 bit)
ROW 2	1	3	a a a a a a a	0
1	1	4	p p p p p p p	1
1	Valid(1bit)	Tag(11 bits)	Data(8 bytes)	LRU(1 bit)

ROW 3	1	2	c c c c c c c	0
1	1	3	alalalalalala	1

Memory access time:

1 cycle for a hit

10 + 1 + 8 = 19 cycles for a miss

For the specific results of how each address is read and cached, please see the output of my Java program. The program also outputs the average CPI value for each cache type I simulated, but since there are quite a few of these, I decided to collate the results into three tables that correspond to each cache type.

**Table 1: Average CPI for Direct Mapped Caches** 

Rows	Block Size	Hits	Misses	Avg CPI
16	4	16	16	8
8	8	20	12	7.75
4	16	23	9	8.313
2	32	25	7	10.19
1	64	27	5	12.57

**Table 2: Average CPI for Fully Associative Caches** 

Entries	Block Size	Hits	Misses	Avg CPI
15	4	10	22	10.63
10	8	18	14	8.88
5	16	23	9	8.313
3	32	27	5	7.563
1	64	27	5	12.563

**Table 3: Average CPI for Set Associative Caches** 

Rows	Block Size	Num Ways	Hits	Misses	Avg CPI
16	4	1	16	16	8
8	4	2	15	17	8.44
8	8	1	20	12	7.75
4	4	4	16	16	8
4	8	2	19	13	8.313
4	16	1	23	9	8.313

2	4	8	13	19	9.313
2	8	5	20	12	7.75
2	16	2	22	10	9.13
2	32	1	25	7	10.19
1	4	16	13	19	9.313
1	8	10	20	12	7.75
1	16	6	23	9	8.313
1	32	3	27	5	7.563
1	64	1	27	5	12.563

It is clear that the best performing cache types are a Fully Associative Cache with 3 entries and 32 byte blocks and a Set Associative with 1 row, 3 way association, and 32 byte blocks. They both result in an average CPI of 7.563. It is unsurprising that these two caches yield the same results; a three way set associative cache with one row is identical in performance to a three-entry fully associative cache.

My conclusion from these data is that the spatial locality is primary factor for performance, at least given this specific set of addresses. It is interesting that some other cache types come close to 7.5 CPI but with far more misses. For example, the 8 row, 8 byte block Direct Mapped cache has an average CPI of 7.75 but actually has 12 misses instead of 5. The 32 byte, 3 way Set Associative/Fully Associative seems to hit a "sweet spot" where there are few misses and the large block size takes advantage of the spatial locality. However, increasing the block size any further than this results in huge CPI penalties for misses. In fact, with 32 byte blocks, the penalty is already quite severe but since the number of misses is so low it wins out over other cache architectures.

Please see below for the source code of my Java program, and please excuse any poor software engineering on my part - it had been several months since I wrote any Java before I began this assignment!

```
import java.util.ArrayList;
/*
 * Assumptions: 1 cycle for cache hits
 *
 * 1 + 10 + [1 cycle per byte] for cache misses
 * Total cache size must not exceed 840 bits
 */
```

```
public class StudyAssignment9 {
       public static void main(String[] args) {
               int[] addresses = new int[] { 4, 8, 12, 16, 20, 32, 36, 40, 44, 20, 32, 36, 40, 44, 64,
68, 4, 8, 12, 92, 96,
                              100, 104, 108, 112, 100, 112, 116, 120, 128, 140, 144 };
               directMapped(addresses, 16, 4);
               directMapped(addresses, 8, 8);
               directMapped(addresses, 4, 16);
               directMapped(addresses, 2, 32);
               directMapped(addresses, 1, 64);
               fullyAssociative(addresses, 15,4);
               fullyAssociative(addresses, 10,8);
               fullyAssociative(addresses, 5,16);
               fullyAssociative(addresses, 3,32);
               fullyAssociative(addresses, 1,64);
               //16 row SA
               setAssociative(addresses, 16,4,1);
               //8 row SA's
               setAssociative(addresses, 8,4,2);
               setAssociative(addresses, 8,8,1);
               //4 row SA's
               setAssociative(addresses, 4,4,4);
               setAssociative(addresses, 4,8,2);
               setAssociative(addresses, 4,16,1);
               //2 row SA's
               setAssociative(addresses, 2,4,8);
               setAssociative(addresses, 2,8,5);
               setAssociative(addresses, 2,16,2);
               setAssociative(addresses, 2,32,1);
               //1 row SA's
               setAssociative(addresses, 1,4,16);
               setAssociative(addresses, 1,8,10);
               setAssociative(addresses, 1,16,5);
               setAssociative(addresses, 1,32,3);
```

setAssociative(addresses, 1,64,1);

```
}
/**
* Simulates a Direct Mapped cache
* @param addrs
                      - array of addresses to read/cache
* @param rows
                      - # rows in the cache
* @param block size - # of bytes / block
public static void directMapped(int[] addrs, int rows, int block size) {
       double hits total = 0;
       double misses total = 0;
       int offset bits = logbase2(block size);
       int row bits = logbase2(rows);
       int[] dm cache = new int[rows];
       //First pass to populate the cache
       for (int i = 0; i < addrs.length; i++) {
               int tag = addrs[i] / (int) Math.ceil(Math.pow(2, offset bits + row bits));
               int row = (addrs[i] / (int) Math.ceil(Math.pow(2, offset bits))) % rows;
               dm cache[row] = tag;
       }
       //Second pass for the "real" analysis
       for (int i = 0; i < addrs.length; i++) {
               int tag = addrs[i] / (int) Math.ceil(Math.pow(2, offset bits + row bits));
               int row = (addrs[i] / (int) Math.ceil(Math.pow(2, offset bits))) % rows;
               String result;
               if (dm \ cache[row] == tag) {
                      hits total++;
                      result = "hit from row" + row;
               } else {
                      misses total++;
                      result = "miss - cached to row: " + row;
               }
```

```
System.out.println("address: " + addrs[i] + ", tag: " + tag + ", result: " +
result);
                       dm cache[row] = tag;
               }
               System.out.println("Direct Mapped Cache - rows: " + rows + "
                                                                                  blocksize: "+
block size + "bytes" +
                              "\n" + "Total misses: " + misses total + ", Total hits: "
                              + hits total + "\nCPI for this set of addresses: "
                              + ((hits total + misses total * (11 + block size)) / addrs.length) +
"\n");
       }
        * Simulates a Fully Associative cache.
        * @param addrs
                              - array of addresses to read/cache
        * @param entries
                              - # entries in the cache
        * @param block size - # of bytes / block
        */
       public static void fullyAssociative(int[] addrs, int entries, int block size) {
               double hits total = 0;
               double misses total = 0;
               int offset bits = logbase2(block size);
               int[] fa cache = new int[entries];
               for (int i = 0; i < fa cache.length; i++) {
                       fa cache[i] = -1;
               }
               int j;
               //First pass to get things warmed up
               for (int i = 0; i < addrs.length; i++) {
                      boolean hit = false;
```

```
int tag = addrs[i] / (int) Math.ceil(Math.pow(2, offset_bits));
                       for (j = 0; j < fa cache.length; j++) {
                               if (fa_cache[j] == tag) {
                                       swap(fa_cache, tag, j);
                                       break;
                               }
                               if (!hit) {
                                       swap(fa_cache, tag, -1);
                               }
                       }
               }
               //Do it for real now!
               for (int i = 0; i < addrs.length; i++) {
                       int tag = addrs[i] / (int) Math.ceil(Math.pow(2, offset bits));
                       boolean hit = false;
                       String result = "";
                       for (j = 0; j < fa_cache.length; j++) {
                               if (fa_cache[j] == tag) {
                                       swap(fa_cache, tag, j);
                                       result = "hit from entry" + j;
                                       hits total++;
                                       hit = true;
                                       break;
                               }
                       if (!hit) {
                               swap(fa_cache, tag, -1);
                               misses_total++;
                               result = "miss - cached to entry: " + j;
                       }
                       System.out.println("address:"+addrs[i]+", tag:"+tag+", result:"+
result);
               }
```

```
System.out.println("Fully Associative Cache - entries: " + entries + " blocksize:
" + block size + " bytes" +
                              "\n" + "Total misses: " + misses total + ", Total hits: "
                              + hits total + "\nCPI for this set of addresses: "
                              + ((hits total + misses total * (11 + block size)) / addrs.length) +
"\n");
       }
       /**
        * Simulates a Set Associative cache.
        * @param addrs
                              - array of addresses to read/cache
        * @param rows
                              - # rows in the cache
        * @param block size - # of bytes / block
        * @param setsize
                              - how many entries there are in a set that corresponds to
                       one row (e.g. for two-way, setsize = 2)
        */
       public static void setAssociative(int[] addrs, int rows, int block size, int set size) {
               double hits total = 0;
               double misses total = 0;
               int offset bits = logbase2(block size);
               int row bits = logbase2(rows);
//
               int tagbits = (16 - offset bits - row bits);
               int LRUbits = logbase2(set size);
//
//
               int SIZE = rows * (set size * (block size * 8 + tagbits + LRUbits + 1));
//
               System.out.println(SIZE);
               if (SIZE > 840) {
//
//
                       System.out.println("TOO BIG!!!!" + " rows: " + rows + " blocksize: " +
block_size + " ways: " + set_size);
//
                       return;
//
               }
               ArrayList < int[] > sa cache = new ArrayList <>();
               //build up the array with a set for each row
               for (int i = 0; i < rows; i++) {
```

```
sa cache.add(new int[set_size]);
}
// first pass to populate the cache before analysis
for (int i = 0; i < addrs.length; i++) {
        int tag = addrs[i] / (int) Math.ceil(Math.pow(2, offset bits + row bits));
        int row = (addrs[i] / (int) Math.ceil(Math.pow(2, offset bits))) % rows;
        int[] rowcache = sa cache.get(row);
        boolean hit = false;
        for (int j = 0; j < \text{rowcache.length}; j++) {
               if (rowcache[j] == tag) {
                       swap(rowcache, tag, j);
                       hit = true;
                       break;
               }
        }
        if (!hit) {
               swap(rowcache, tag, -1);
        }
}
//second pass through the address array to perform the analysis.
for (int i = 0; i < addrs.length; i++) {
        int tag = addrs[i] / (int) Math.ceil(Math.pow(2, offset bits + row bits));
        int row = (addrs[i] / (int) Math.ceil(Math.pow(2, offset bits))) % rows;
        int[] rowcache = sa cache.get(row);
        String result = "";
        boolean hit = false;
        for (int j = 0; j < \text{rowcache.length}; j++) {
               if (rowcache[i] == tag) 
                       swap(rowcache, tag, j);
                       hit = true;
                       hits total++;
                       result = "hit from row: " + row;
```

```
break;
                               }
                       }
                               if (!hit) {
                                       swap(rowcache, tag, -1);
                                       result = "miss - cached to row: " + row;
                                       misses total++;
                               }
                               System.out.println("Address: " + addrs[i] + ", tag: " + tag + ",
result: " + result);
                       }
               System.out.println("Set Associative Cache - rows: " + rows + " blocksize: " +
block size + "bytes ways: " + set size + "\n"
               + "Total misses: " + misses_total + ", Total hits: "
                               + hits total + "\nCPI for this set of addresses: "
                               + ((hits total + misses total * (11 + block size)) / addrs.length) +
"\n");
                }
        * Places the tag at the end of the array and shifts everything else down to
        * represent it as being used less recently.
        */
        private static void swap(int[] associative cache, int tag, int index) {
               // If tag was already the most recently used and it's in the set, do nothing
               if (index == associative cache.length - 1)
                       return;
               // If tag was not in the set, shift everything left and replace the oldest entry
               // with tag
               else if (index == -1) {
                       for (int i = 0; i < associative cache.length - 1; <math>i++) {
                               int tmp = associative cache[i];
                               associative cache[i] = associative cache[i + 1];
```

```
associative_cache[i + 1] = tmp;
               }
               associative_cache[associative_cache.length - 1] = tag;
               return;
       }
       // Tag is in the set but it's not the most recent - adjust accordingly.
       else {
               for (int i = index; i < associative_cache.length - 1; i++) {
                       int temp = associative_cache[i + 1];
                       associative\_cache[i + 1] = tag;
                       associative_cache[i] = temp;
}
* Return the log base 2 of input as an integer (seriously, the Java math
* library doesn't have this?!)
private static int logbase2(int x) {
       return (int) (Math.log(x) / Math.log(2) + 1e-10);
}
```

}