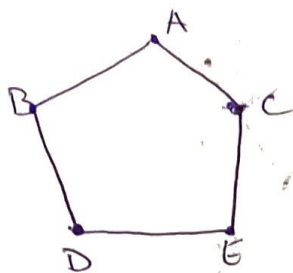
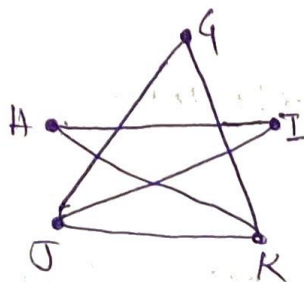


1) Let us consider a  $C_5$  graph below.



Now, we will consider the complement of  $C_5$  graph below is  $\overline{C_5}$



Now consider

$$F: C_5 \rightarrow \overline{C_5}$$

$$F(A) = G$$

$$F(B) = H$$

$$F(C) = I$$

$$F(D) = J$$

$$F(E) = K$$

Now Total number of vertices for  $C_5 = 5$  & Total number of edges for  $C_5 = 5$

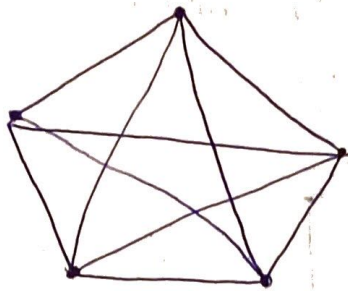
Now total number of vertices for  $\overline{C_5} = 5$  & Total number of edges for  $\overline{C_5} = 5$

where  $F$  is bijection b/w vertex set and edge set, which preserves adjacency.

$\therefore F$  is isomorphic

From above statements, we can say that  $C_5$  graph and its complement  $\overline{C_5}$  is isomorphic. ( $\therefore$  Hence proved).

2) Let us consider the graph  $K_5$



Degree of each vertex = 4

Calculating the all edges with 5 vertices in the graph  $K_5$

$$\text{Total Edges} = 4 + 4 + 4 + 4 + 4 \\ = 20$$

$$e = \text{number of edges} = 10$$

$$v = \text{number of vertices} = 5$$

[ $\therefore 2n$   $K_5$  graph]

Let us assume  $K_5$  is Planar

$$v - e + r = 2$$

$$5 - 10 + r = 2$$

$$r = 2 + 5$$

$$\boxed{r = 7}$$

If a graph is a Planar, then

$$3r \leq 2e$$

$$3(7) \leq 2(10)$$

$$21 \leq 20 \text{ --- (1)}$$

consider (1) which is not true, we can say that our assumption is wrong, which clearly indicates graph  $K_5$  is not Planar.

$\therefore K_5$  is not Planar.

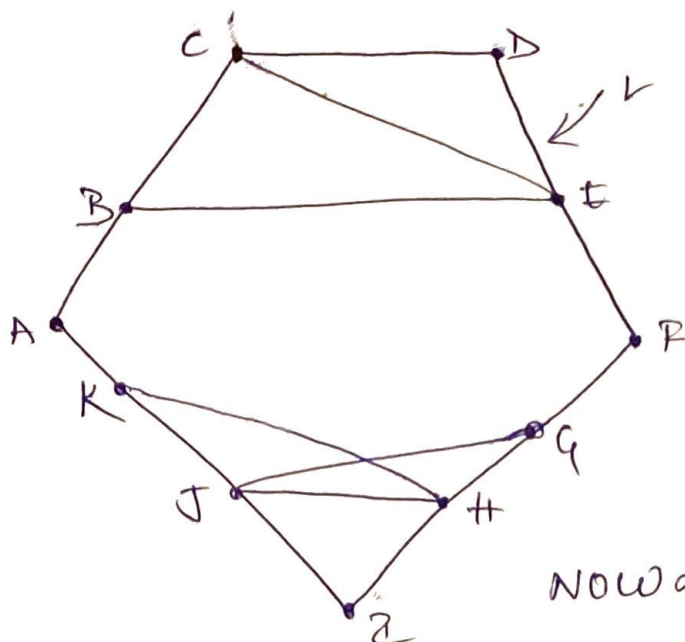
3) Let us explain this by induction. The only tree on 2 vertices is  $K_2$ , which can be clearly a bipartite.

Now, let us assume that every tree on ' $n$ ' vertices is a bipartite graph, where its vertex set can be decomposed into two sets. Let  $T$  be a tree on ' $n+1$ ' vertices, we know that the vertex set of ' $T$ ' contains a leaf ' $v$ '. Furthermore  $T' = T - v$  is a tree on ' $n$ ' vertices by the induction hypothesis the vertex set of ' $T$ ' can be decomposed into disjoint sets  $x$  &  $y$ .

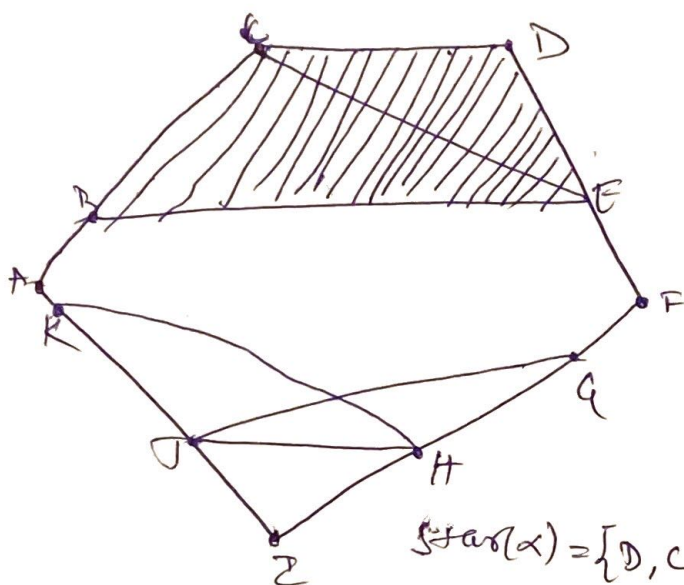
Let ' $w$ ' be the neighbour of  $v$  in  $T$ . Then  $w$  is also a vertex in the graph  $T'$  without loss of generality (wlog), we can assume that  $w \in x \subset V(T')$ .

Then  $x' = x$  &  $y' = y \cup \{v\}$  is a bipartition of  $V(T)$  which makes  $T$  as a bipartite graph.

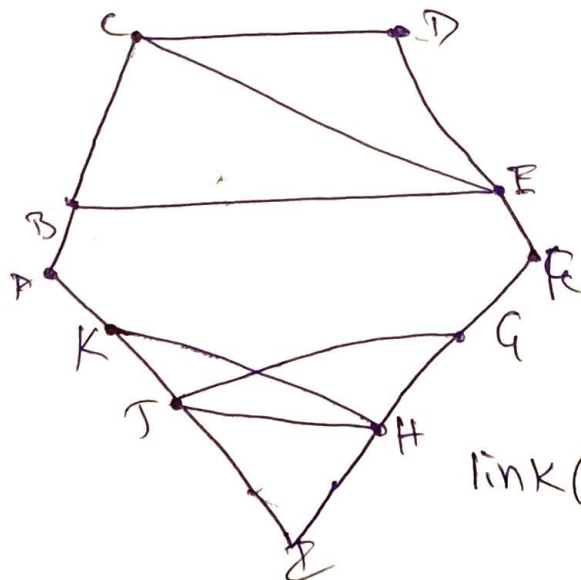
4) consider a graph below



$$\text{NOW } \alpha = \{D, E, \langle D, E \rangle\}$$



$$\text{Star}(\alpha) = \{D, C, E, B, \langle D, C \rangle, \langle D, E \rangle, \langle E, B \rangle\}$$



$$\text{link}(\alpha) = \{B, C, \langle B, C \rangle\}$$