Privacy in NLP

Deep Learning for NLP: Lecture 11

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Importance of Privacy

Importance of Privacy

Question: Why is privacy important?

Two ways to answer:

- 1. Societal perspective
- 2. Research perspective

Today's world: "Data is the new oil"

- Ethical concerns over data collection
- Legal concerns for businesses due to laws and privacy guidelines

Why privacy is important: Research Perspective

Very difficult to convince data holders to provide data (e.g. hospital medical records)

Can 'pip install' MNIST dataset and train a classifier in minutes

Cannot 'pip install cancer-dataset', need a lot of work/resources to get hold of such data

Overview of Lecture

- Why is privacy important and consequences of non-privacy
- Gold standard of privacy: Differential privacy
 - Randomized response
 - Pure differential privacy and the Laplace mechanism
 - Properties of differential privacy
 - Approximate differential privacy and the Gaussian mechanism
- Applying differential privacy for ML: DP-SGD
- Other methods in privacy
 - Secure multiparty computation
 - Federated learning
 - Homomorphic encryption

Attacks on Non-Privatized Data

and Models

Data Anonymization

What if we just anonymize data?

Example

"Robert Smith" \rightarrow "id38729848"

Linkage Attack

Re-identifying anonymized individuals by combining data with background information



Linkage Attack: Adversary acquires private information by correlating multiple datasets

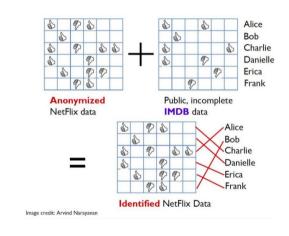
https://www.researchgate.net/figure/
Different-privacy-attack-and-threat-models_fig5_346302647

Consequences of Non-Privacy: Netflix Prize

Netflix: Online streaming service Netflix prize: Challenge between 2006 and 2009, prize of \$1,000,000

- Goal: Create a model for best recommendations of their service
- Data: Anonymized user IDs, movie IDs, ratings and dates

Privacy breach: Match Netflix data (anonymized) with IMDb data (public)

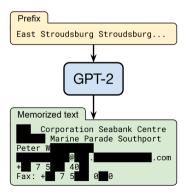


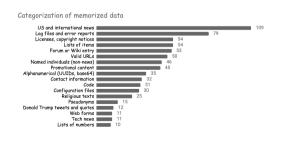
Consequences of Non-Privacy: Memorization in Neural Networks

Let's look at a black-box attack on extracting data from NNs

Neural networks can memorize their training data [Carlini et al., 2021]

We can extract this in multiple ways, one of which is prompting



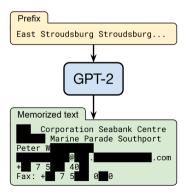


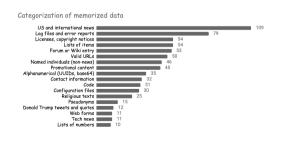
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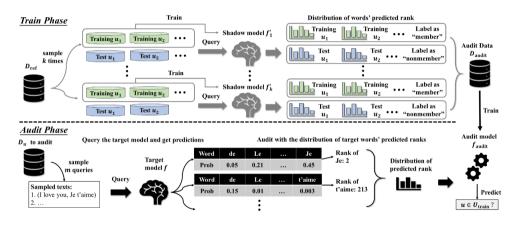
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Consequences of Non-Privacy: Membership Inference



Song and Shmatikov [2019]

Consequences of Non-Privacy: Model inversion for NNs

Model inversion [Fredrikson et al., 2015], an example of a white-box attack Basic idea: Follow the gradient used to adjust the weights of a model, obtain a reverse-engineered example for all represented classes in the model





Figure 1: An image recovered using a new model inversion attack (left) and a training set image of the victim (right). The attacker is given only the person's name and access to a facial recognition system that returns a class confidence score.

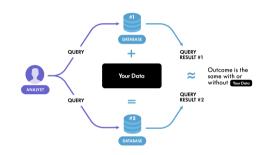
The *gold standard* in leading research with privacy guarantees

Intuitively

Data perturbation, output of algorithm cannot change beyond a very specific amount, when one data point is added/removed

Won the Test of Time Award back in 2016

Used by big companies like Microsoft, Apple, Google and Facebook



https://www.winton.com/research/ using-differential-privacy-to-protect-personal-data

Randomized Response

Differential Privacy: Randomized Response

Oldest DP algorithm (Warner, 1965)

Technique used to collect sensitive information from individuals, while maintaining their confidentiality

Provides plausible deniability

More formally:

n students

 $X_i \in \{0,1\}$: Did individual i cheat on test?

 Y_i : Value that depends on X_i with added randomness

Goal of analyst: Estimate $p = \frac{1}{n} \sum X_i$ (fraction of individuals that cheated)

Method 1: Perfect accuracy, no privacy

$$Y_i = \begin{cases} X_i & \text{w.p. 1} \\ 1 - X_i & \text{w.p. 0} \end{cases}$$

Method 2: Perfect privacy, no accuracy

$$Y_i = \begin{cases} X_i & \text{w.p. } \frac{1}{2} \\ 1 - X_i & \text{w.p. } \frac{1}{2} \end{cases}$$

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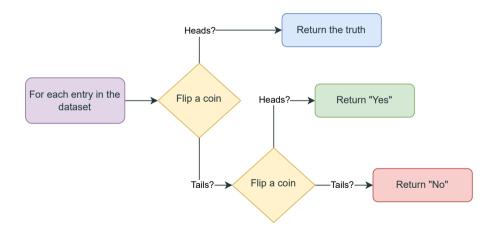
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Method 3



More formally

New parameter $\gamma \in (0,\frac{1}{2})$

$$Y_i = \begin{cases} X_i & \text{w.p. } \frac{1}{2} + \gamma \\ 1 - X_i & \text{w.p. } \frac{1}{2} - \gamma \end{cases}$$

If $\gamma=\frac{1}{2}$, this is Method 1 (no privacy, perfect accuracy)
If $\gamma=0$, this is Method 2 (no accuracy, perfect privacy)

Compromise

Set $\gamma=\frac{1}{4}$ — provides plausible deniability

 $\gamma \rightarrow 0$, maximum deniability

 $\gamma \to \frac{1}{2}$, no deniability (no privacy)

Pure Differential Privacy

Pure Differential Privacy

Differential Privacy (DP)

Data perturbation, where the output of an algorithm cannot change by more than a **specific amount**, when adding/removing/altering one data point in a dataset.

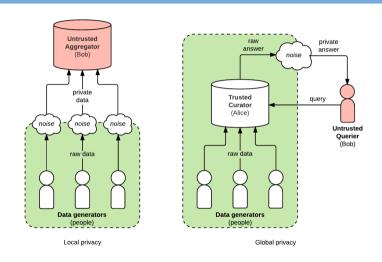
A property of an algorithm, information-theoretic guarantee.

Originally proposed by Dwork et al. [2006], extensively outlined in Dwork and Roth [2013]

Privacy Budget (ε)

The total amount of privacy leakage that is allowed to occur ('amount' of privacy).

Pure Differential Privacy



Pure Differential Privacy

Additional relevant terms:

Query

The 'question' the analyst is asking about the data.

E.g. Mean, sum (simple); gradient of loss function (more complex).

Trusted curator

The aggregator of the data, adds noise to achieve DP guarantee.

Sensitivity

The maximum difference in an algorithm's outputs, when one data point is changed.

Pure DP: Concrete Example

Sum query
Dataset of binary values
Contains n individuals, each associated with 0 or 1
Query: How many people in our
dataset smoke?

Name	Smokes?
Alice	0
Bob	0
Clair	1
÷	:
Zane	1

Pure DP: Concrete Example

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Sensitivity?

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Sensitivity? — 1

Dataset D, consisting of n individuals, x_1, \ldots, x_n

We run mechanism M over this dataset D to get an output M(D)

For $\varepsilon \geq 0$, mechanism M is ε -differentially private if, for all $S \subseteq \mathsf{Range}(M)$, and **neighboring datasets** D and D':

Differential Privacy

$$Pr[M(D) \in S] \le \exp(\varepsilon) Pr[M(D') \in S]$$

Neighboring datasets

A dataset is **neighboring** to another dataset if it differs from it in one row. I.e. $||D - D'||_1 \le 1$

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arepsilon o 0: we approach perfect privacy, but less utility of our algorithm (less difference in output distributions)

 $arepsilon o \infty$: we approach the original non-DP setting (no constraint on output distributions)

How big do we make ε ?

Should be 'fairly small', generally $0.1 \le \varepsilon \le 5$

Randomized response mechanism with $\gamma = \frac{1}{4}$ (throwing a fair coin):

$$\varepsilon = \ln 3 \approx 1.1$$

As arepsilon increases, privacy guarantee gets exponentially worse

Achieving Pure DP: The Laplace Mechanism

How do we achieve this ε -DP guarantee? — Laplace Mechanism

Sensitivity revisited

$$f:D^n\to\mathbb{R}^k$$

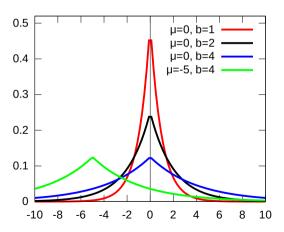
$$l_1$$
-sensitivity of the function f is: $\Delta^{(f)} = \max_{D,D'} ||f(D) - f(D')||_1$

Example Calculation: Sum query

If f sums up a set of bits, $\sum X_i$, where $X_i \in \{0,1\}$, then: $\Delta^{(f)} = 1$

Achieving Pure DP: The Laplace Mechanism

Laplace Distribution: $p(x) = \frac{1}{2b} \exp\left(-\frac{|x|}{b}\right)$



Achieving Pure DP: The Laplace Mechanism

Laplace Mechanism

$$f:D^n o\mathbb{R}^k$$
 $M(D)=f(D)+(Y_1,\ldots,Y_k)$ $Y_i \underset{i.i.d}{\sim} \mathsf{Lap}(rac{\Delta^{(f)}}{arepsilon})$

Add noise to each coordinate, proportional to the L1 sensitivity

Example: Sum query

$$f = \sum X_i, \ \Delta^{(f)} = 1, \ k = 1$$

$$\tilde{p} = f(x) + \mathsf{Lap}(\frac{1}{\varepsilon})$$

Name	Smokes?
Alice	0
Bob	0
Clair	1
÷	÷
Zane	1

Privacy Loss Random Variable

For mechanism $M:D^n\to Y$ and neighboring datasets D,D', we can define the **privacy loss random variable** $\mathcal{L}_{M(D)||M(D')}$

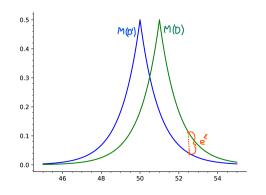
Privacy Loss Random Variable

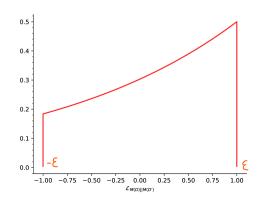
$$\mathcal{L}_{M(D)||M(D')} = \ln\left(\frac{M(D)=\xi}{M(D')=\xi}\right)$$
, distributed by drawing $\xi \sim M(D)$

Privacy Loss Random Variable

We can translate between this privacy loss random variable and our ε -DP:

 $|\mathcal{L}_{M(D)||M(D')}| \leq \varepsilon$ w.p. 1, for all D and D' neighboring datasets





Properties of DP

Properties of DP

- 1. Closed under post-processing
 - If $M:D^n \to Y$ is ε -DP, and $F:Y \to Z$ is another randomized mapping, then $F \circ M$ is ε -DP
- 2. Group privacy
 - If $M:D^n \to Y$ is ε -DP, and D, D' differ in k positions, then for all $S \in \mathsf{Range}(M)$: $Pr[M(D) \in S] \le \exp(k\varepsilon)Pr[M(D') \in S]$
- 3. Basic composition
 - If we run k ε -DP algorithms sequentially through our data, the full process will be $k\varepsilon$ -DP

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Differential Privacy

Approximate Differential Privacy

Approximate Differential Privacy

We can 'loosen' our privacy guarantees a little

Increase utility, not give up that much privacy

Idea

We take our original $\varepsilon\text{-DP}$ privacy guarantee, and add a 'cryptographically small' probability that it will not work

It turns out, this is enough to significantly improve utility of our DP mechanism!

Approximate Differential Privacy

For $\varepsilon \geq 0$, mechanism M is ε, δ -differentially private if, for all $S \subseteq \mathsf{Range}(M)$, and **neighboring datasets** D and D':

Approximate Differential Privacy

$$Pr[M(D) \in S] \le \exp(\varepsilon) Pr[M(D') \in S] + \delta$$

If $\delta = 0$, we go back to our original 'pure' DP definition

How about the privacy loss random variable?

$$|\mathcal{L}_{M(D)||M(D')}| \leq \varepsilon$$
 w.p. $1 - \delta$, for all D and D' neighboring datasets

Approximate Differential Privacy

What should we set δ to?

Good rule of thumb: $\delta \ll \frac{1}{n}$, where n is the size of the dataset D

Example $0, \delta$ -DP Mechanism

M: For each $x \in D$, output x w.p. δ , and do nothing w.p. $1 - \delta$

Probability we do not release anyone's data point: $(1 - \delta)^n$

Probability we **do** release someone's data point: $1 - (1 - \delta)^n$

If δ is around $\frac{1}{n}$, then this is approximately 1 (as δ increases, this approaches 1)

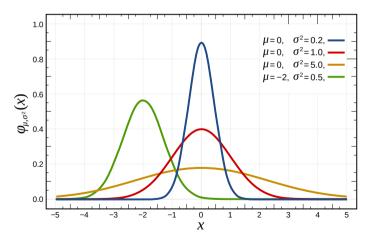
How do we achieve the $\varepsilon, \delta\text{-DP}$ guarantee? — Gaussian Mechanism

l_2 Sensitivity

$$f:D^n\to\mathbb{R}^k$$

$$l_2\text{-sensitivity of function }f$$
 is: $\Delta_2^{(f)} = \max_{D,D'} \lvert\lvert f(D) - f(D') \rvert\rvert_2$

Gaussian Distribution:
$$p(x) = \frac{1}{\sqrt{2\pi\sigma^2}} \exp\left(-\frac{(x-\mu)^2}{2\sigma^2}\right)$$



Gaussian Mechanism
$f:D^n \to \mathbb{R}^k$
$M(D) = f(D) + (Y_1, \dots, Y_k)$
$Y_i \sim_{i.i.d} \mathcal{N}(0, 2\ln\left(\frac{1.25}{\delta}\right)\frac{\Delta_2^2}{\varepsilon^2})$

Add noise to each coordinate, proportional to the L2 sensitivity

Name	Attr. 1	Attr. 2	
Alice	0	1	
Bob	0	1	
Clair	1	0	
÷	:	:	
Zane	1	0	

Example: Sum query with more attributes

Mechanism: $f = \sum X_i$

 $D \in \{0,1\}^{n \times d}$, where n is the number of individuals and d the number of attributes

Worst case for neighboring datasets D, D^{\prime} : D^{\prime} s row has all 1s, $D^{\prime\prime}$ s row has all 0s

$$l_1$$
-sensitivity: $||\mathbf{1} - \mathbf{0}||_1 = ||\mathbf{1}||_1 = d$

$$l_2$$
-sensitivity: $||\mathbf{1}||_2 = \sqrt{d}$

Laplace:
$$\tilde{p} = f(x) + \operatorname{Lap}(\frac{d}{\varepsilon})$$

Gaussian:
$$\tilde{p} \approx f(x) + \mathcal{N}(0, (\frac{\sqrt{d}}{\varepsilon})^2)$$

Benefits of Approximate DP

- 1. Scale of added noise can be significantly less (with higher dimensions)
- 2. Can improve upon basic composition (more advanced composition techniques)
 - If we run $k \in \delta$ -DP algorithms sequentially through our data, the full process will be *less than* $k \in k \delta$ -DP (depending on the composition technique)

Differential Privacy for Machine

Learning

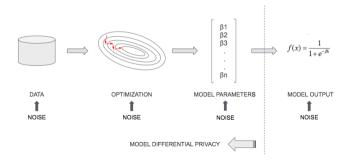
Applying DP for ML

	stars	sentiment	text
0	5	positive	After getting food poisoning at the Palms hote
1	4	positive	"A feast worthy of Gods"\n\nBaccarnal Buffet i
2	4	positive	The crab legs are better than the ones at \mbox{Wick}
3	1	negative	Not worth it! Too salty food and expensive! Th
4	5	positive	I would give this infinite stars if I could. M
10412	5	positive	Best buffet ever! Irma was great, served us be
10413	4	positive	HollIlllyyyy moleyyyy! \n\nThis buffet was one
10414	5	positive	The selection is amazing and all the food is e $% \label{eq:condition}%$
10415	4	positive	One of the best buffets I've had in Vegas. My \dots
10416	4	positive	I got a chance to go to the Bacchanal Buffett

10374 rows × 3 columns

Applying DP for ML

DP and ML: Instead of privatizing an algorithm, we're privatizing a model Need to choose a 'query' as before (associated with the model, dependent on the dataset)



Most common and widely adopted algorithm in differentially private ML: Differentially Private Stochastic Gradient Descent (DP-SGD)

Can apply directly to model training

The 'query' of our DP mechanism: Gradient of the loss function

What's the sensitivity of the gradient?

Most common and widely adopted algorithm in differentially private ML:

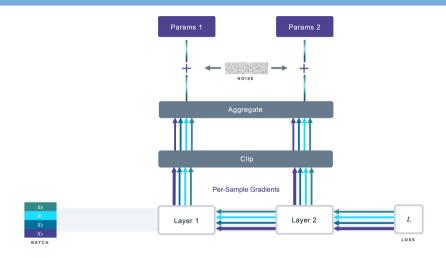
Differentially Private Stochastic Gradient Descent (DP-SGD)

Proposed by Abadi et al. [2016]

Can apply directly to model training

The 'query' of our DP mechanism: Gradient of the loss function

What's the sensitivity of the gradient? — ∞ ...



Algorithm for DP-SGD:

- 1. Select a 'lot' of points
 - Lot: A set of data points, where each point is selected with probability L/n, where L is the 'lot size' and n is the size of the dataset
- 2. For each point in the 'lot', compute the gradient $g_i = \nabla l(\theta_t, x_i, y_i)$, $\forall i \in \text{lot}$
- 3. Clip g_i to l_2 ball of radius C, then average
- 4. Add noise
- 5. Step in negative direction of gradient

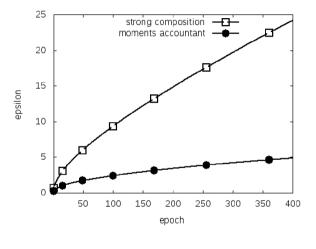
```
Algorithm 1 Differentially private SGD (Outline)
Input: Examples \{x_1,\ldots,x_N\}, loss function \mathcal{L}(\theta) =
   \frac{1}{N}\sum_{i}\mathcal{L}(\theta,x_{i}). Parameters: learning rate \eta_{t}, noise scale
   \sigma, group size L, gradient norm bound C.
   Initialize \theta_0 randomly
   for t \in [T] do
       Take a random sample L_t with sampling probability
       L/N
      Compute gradient
       For each i \in L_t, compute \mathbf{g}_t(x_i) \leftarrow \nabla_{\theta_t} \mathcal{L}(\theta_t, x_i)
      Clip gradient
      \bar{\mathbf{g}}_t(x_i) \leftarrow \mathbf{g}_t(x_i) / \max\left(1, \frac{\|\mathbf{g}_t(x_i)\|_2}{C}\right)
      Add noise
      \tilde{\mathbf{g}}_t \leftarrow \frac{1}{L} \left( \sum_i \bar{\mathbf{g}}_t(x_i) + \mathcal{N}(0, \sigma^2 C^2 \mathbf{I}) \right)
       Descent
      \theta_{t+1} \leftarrow \theta_t - \eta_t \tilde{\mathbf{g}}_t
   Output \theta_T and compute the overall privacy cost (\varepsilon, \delta)
   using a privacy accounting method.
```

Important points about DP-SGD:

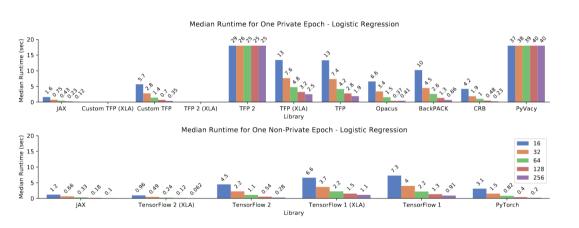
- 1. Use **Poisson sampling** to select lots (**not** simply iterate over batches)
- 2. Moments accountant: Much better bounds on privacy budget
- 3. Clipping can slow down computations

Important points about DP-SGD:

- 1. Use **Poisson sampling** to select lots (**not** simply iterate over batches)
- 2. Moments accountant: Much better bounds on privacy budget
- 3. Clipping can slow down computations



Abadi et al. [2016]



Subramani et al. [2021]

Opacus

```
model = Net()
optimizer = torch.optim.SGD(model.parameters(), lr=0.05)
privacy engine = PrivacyEngine(
  model
  batch size=32.
   sample size=len(train loader.dataset).
   alphas=range(2,32),
   noise multiplier=1.3,
  max_grad_norm=1.0,
privacy engine.attach(optimizer)
```

Other Methods in Privacy Research

Other Methods in Privacy: Secure Multiparty Computation

Secure Multiparty Computation

Combining private inputs from multiple people, in order to compute a function, without revealing anyone's input to the rest

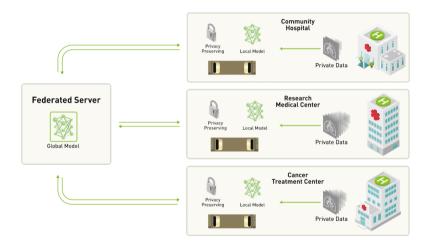
Encryption: These numbers are encrypted, nobody knows their own share

Shared Governance: The numbers can only be decrypted if everyone agrees

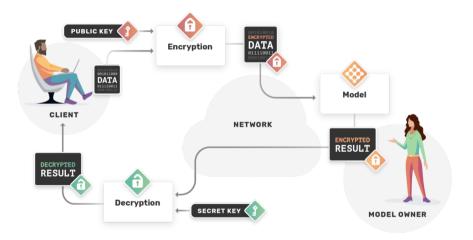
Main aspects:

- 1. Data remains on a remote machine
- 2. Model can be encrypted during training
- 3. Multiple data owners privately combining their data

Other Methods in Privacy: Federated Learning



Other Methods in Privacy: Homomorphic Encryption



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Martin Abadi, Andy Chu, Ian Goodfellow, H. Brendan McMahan, Ilya Mironov, Kunal Talwar, and Li Zhang. Deep Learning with Differential Privacy. In *Proceedings of the 2016 ACM SIGSAC Conference on Computer and Communications Security*, pages 308–318, Vienna, Austria, 2016. ACM. doi: 10.1145/2976749.2978318. URL https://dl.acm.org/doi/10.1145/2976749.2978318.

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