編譯器設計

黄元欣

shin@csie.ntust.edu.tw

Syllabus

Reference Books



J.R. Levine, T. Mason, and D. Brown
O'Reilly 1995
ISBN 1-56592-000-7

Tim Lindbolm - Frank Volin
The Javar Virtual
Machine Specification
Second Edition
Tolera Sures

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The Java™ Virtual Machine Specification, 2nd Ed.
 Tim Lindholm and Framk Yellin
 Addison-Wesley 1999
 ISBN 0-201-43294-3
 http://java.sun.com/docs/books/vmspec

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Text Books



Compilers: Principles, Techniques, and Tools 2e
 A.V. Aho, M.S. Lam, R. Sethi, and J.D. Ullman
 Addison-Wesley 2007
 ISBN 0-321-48681-1



Introduction to the Theory of Computation
 Michael Sipser
 PWS Publishing 1997
 ISBN 0-534-94728-X

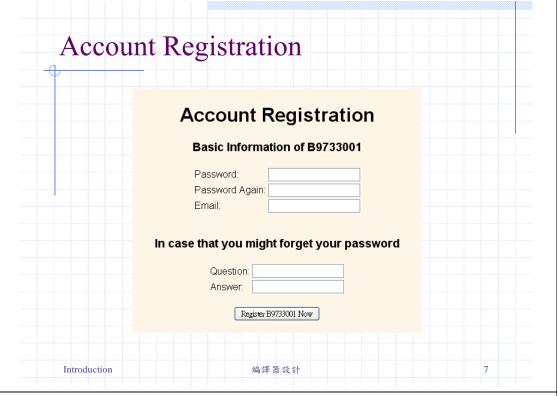
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Course Outline

- Introduction
- Lexical Analysis (Chap. 3)
- Syntax Analysis (Chap. 4)
- Syntax-Directed Translation (Chap. 5)
- Run-time Organization (Chap.7)
- Intermediate Code Generation (Chap.6)
- Code Generation (Chap. 8)

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Syllabus Grading Programming Assignments 30% Midterm 30% Final 40% Office Hours M8, T8 (T4-512, Tel: 6746)



Account Registration

Introduction

You need to get an account at the class home page in order to submit programming assignments and download lecture slides

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http://faculty.csie.ntust.edu.tw/~shin/compilers.html



Grades

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You can get the scores of all your examinations and programming assignments at the class home page:

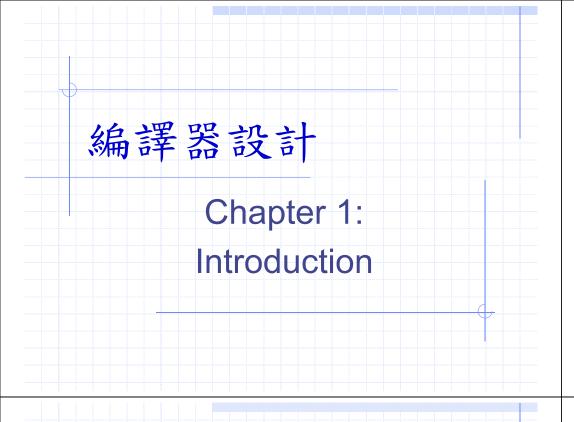
http://faculty.csie.ntust.edu.tw/~shin/compilers.html

• Get Grades

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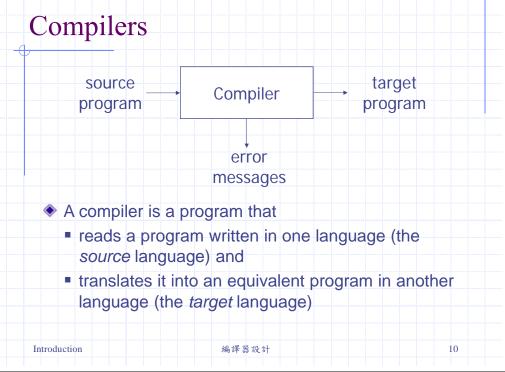
WAAA

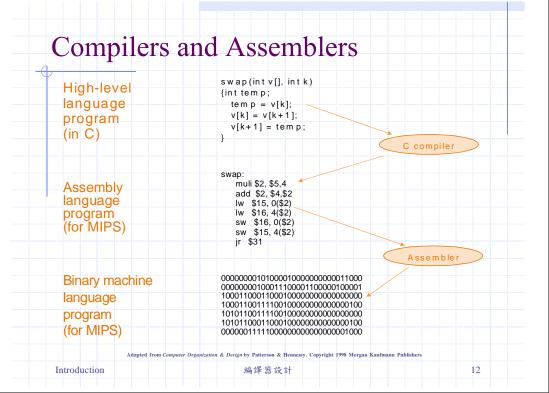


Compilers

- Examples
 - Compilers of C/C++, Fortran, Java, etc
 - Text formatters, e.g. TeX, LaTeX
 - Silicon compilers
 - Query interpreters, e.g SQL compilers
 - Preprocessors
 - Assemblers
 - Browsers
 - Parallelizing compilers

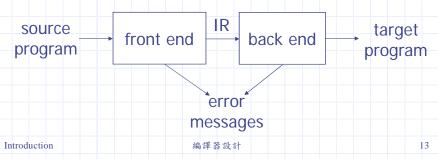
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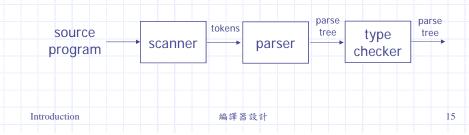
Analysis-Synthesis Model

- There are two parts to compilation
 - Analysis (front end)
 - Breaks up the source program into constituent pieces
 - Creates an intermediate representation (IR)
 - Synthesis (back end)
 - Constructs the desired target program from the IR
 - (Optionally) performs optimizations



Analysis of the Source Program

- Analysis (front end) consists of 3 phases:
 - Linear Analysis (Lexical Analysis)
 - scan characters and group them into tokens
 - Hierarchical Analysis (Syntax Analysis)
 - group tokens into grammatical phrases
 - Semantic Analysis
 - identify semantic errors and gather type information



Analysis-Synthesis Model

- Some tools that perform analysis
 - Structure editors
 - Takes a sequence of commands as input to build a source program, e.g. with the user types while, the editor supplies the matching do
 - Pretty printers
 - Analyzes a program and prints it in such a way that the structure of the program becomes clearly visible
 - Static checkers
 - Reads and analyzes a program, and attempts to discover potential bugs without running the program
 - Interpreters
 - Instead of producing a target program as a translation, an interpreter performs the operations implied by the program

Lexical Analysis

Introduction

- Mapping characters into tokens
 - Tokens: the basic unit of syntax
 - position = initial + rate * 60
 becomes
 (id_position) = (id_initial) + (id_initi

 $\langle id, position \rangle = \langle id, initial \rangle + \langle id, rate \rangle * 60$

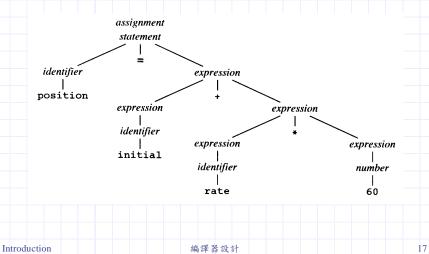
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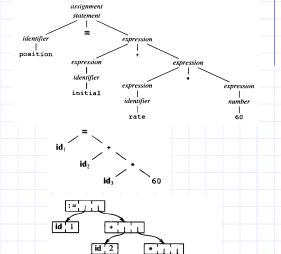
Syntax Analysis

 Tokens are grouped into grammatical phrases that are used to synthesize output



Parse Tree vs. Syntax Tree

- Parse tree
 - describes the syntactic structure of the source program
- Syntax tree
 - A more common internal representation of this syntactic structure
- A compressed representation of the parse tree



num 60

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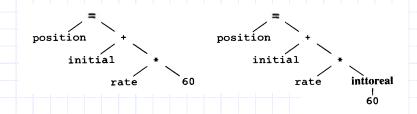
Syntax Analysis

- The hierarchical structure of a program is usually expressed by recursive rules, e.g.
 - Any identifier is an expression
 - Any number is an expression
 - If expression₁ and expression₂ are expressions, so are expression₁ op expression₂ (expression₁)

Semantic Analysis

- Checks for semantic errors
- Gathers type information for the subsequent code-generation phase

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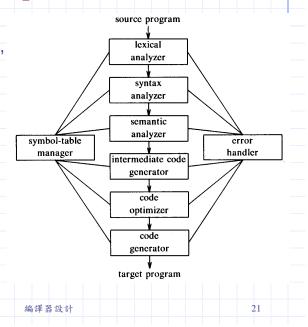
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Phases of a Compiler

A compiler
 operates in phases,
 each of which
 transforms the
 source program
 from one
 representation to
 another



Error Detection and Reporting

- Each phase can encounter errors
 - After detecting an error, a phase must deal with the error, so that compilation can proceed
 - allowing further errors to be detected
 - Lexical phase can detect errors where characters remaining in the input do not form any token
 - Syntax analysis phase detects errors where the token stream violates the syntax of the language
 - Semantic analysis phase tries to detect constructs that have the right syntactic structures but no meaning to the operation involved
 - e.g. a = b + c; where b is an array and c an integer

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Symbol-Table Management

- Essential function of a compiler
 - To record the identifiers used in the source program and collect information about various attributes of each identifier
 - e.g. allocated storage, type, scope, etc.
- Symbol table

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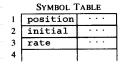
- A data structure containing a record for each identifier, with fields for the attributes
- When an identifier is detected by the lexical analysis, it is entered into the symbol table
- The attributes are determined during syntax analysis and semantic analysis
- e.g. float position, initial, rate;

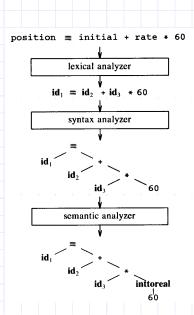
Analysis Phases

Lexical Analysis

Syntax Analysis

Semantic Analysis





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Intermediate Code Generation

- Two properties
 - Easy to produce
 - Easy to translate into the target program
- Examples

Introduction

- Graph representations
- Postfix notation
- Three-address code

id, * intoreal 60

intermediate code generator

temp1 := inttoreal(60)
temp2 := id3 * temp1
temp3 := id2 + temp2
id1 := temp3

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Code Generation

Generates target code

 Consisting of relocatable machine code or assembly code temp1 = id3 * 60.0 id1 = id2 + temp1 v

code generator

MOVF id3, R2

MULF #60.0, R2

MOVF id2, R1

ADDF R2, R1

MOVF R1, id1

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Code Optimization

- Attempts to improve the intermediate code
 - So that faster-running machine code will result

Cousins of the Compiler

Preprocessors

- Preprocessors produce input to compilers
- Macro processing
 - Allows users to define macros
- File inclusion
 - Includes header files into the program text, e.g.
 #include <stdio.h>
- "Rational" preprocessors
 - Augment older languages with more modern flow-of-control and data-structuring facilities
- Language extensions
 - Add capabilities to languages by what amounts to built-in macro

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■ e.g. HPF

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Cousins of the Compiler

Assemblers

mov a, R1 0001 01 00 00000000 add #2, R1 0011 01 10 00000010 mov R1, b 0010 01 00 00000100

Loaders and Link-Editors

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Compiler-Construction Tools

- Some general tools have been created for the automatic design of specific compiler components
 - Parser generators
 - Producing syntax analyzers, normally from input that is based on a context-free grammar
 - Scanner generators
 - Automatically generating lexical analyzers, normally from a specification based on regular expressions
 - Automatic code generators
 - Taking a collection of rules that define the translation of each operation of the intermediate language into the machine language

