

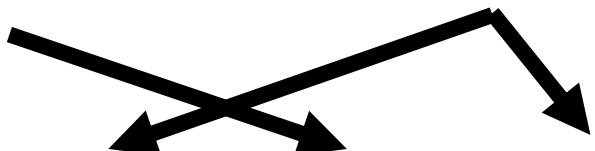
العودية Recursion

هو حل المشكلة بتقسيمها لجزئين

(1) جزء نفس المشكلة لحد أصغر (2) جزء معلوم الحل


مشكلة حساب الفيبوناتشي

(1) فيبوناتشي حد اصغر (2) جمع


$$f(n) = f(n-1) + f(n-2)$$

مشكلة حساب العاملي


(1) عاملي حد اصغر (2) الضرب


$$n! = n \times (n-1)!$$

العودية Recursion

مشكلة حساب الفيبوناتشي

(1) فيبوناتشي حد اصغر 2 جمع


$$f(n) = f(n-1) + f(n-2)$$


+ شرط توقف !!

$$f(0) = 1$$

$$f(1) = 1$$

مشكلة حساب العاملي

(1) عاملي حد اصغر 2 الضرب


$$n! = n \times (n-1)!$$

+ شرط توقف !!

$$0! = 1$$

$$1! = 1$$

العودية Recursion

مشكلة حساب الفيبوناتشي

```
11 f(11 n)
{
    if(n == 0 || n == 1)
        return 1;
    return f(n-1)+f(n-2);
}
```

$\sim O(2^n)$ but why?

مشكلة حساب العاملي

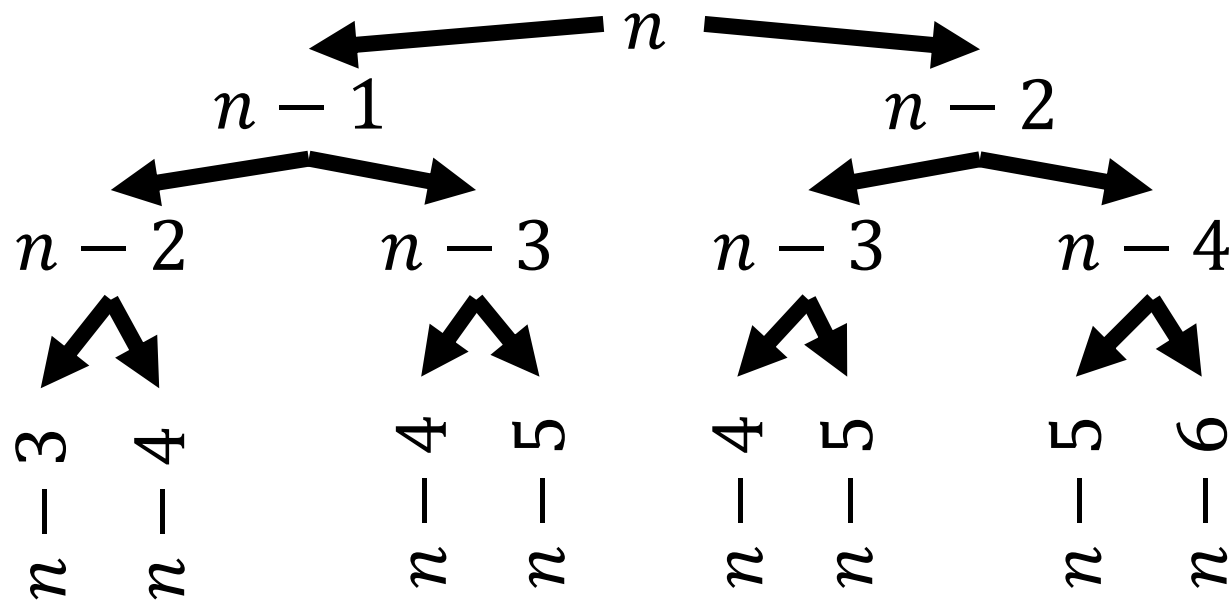
```
11 f(11 n)
{
    if(n == 0 || n == 1)
        return 1;
    return n*f(n-1);
}
```

$O(n)$ but why?

Recursion العودية

مشكلة حساب الفيبوناتشي

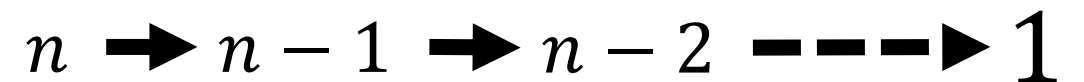
$\sim O(2^n)$ but why?



2^n times

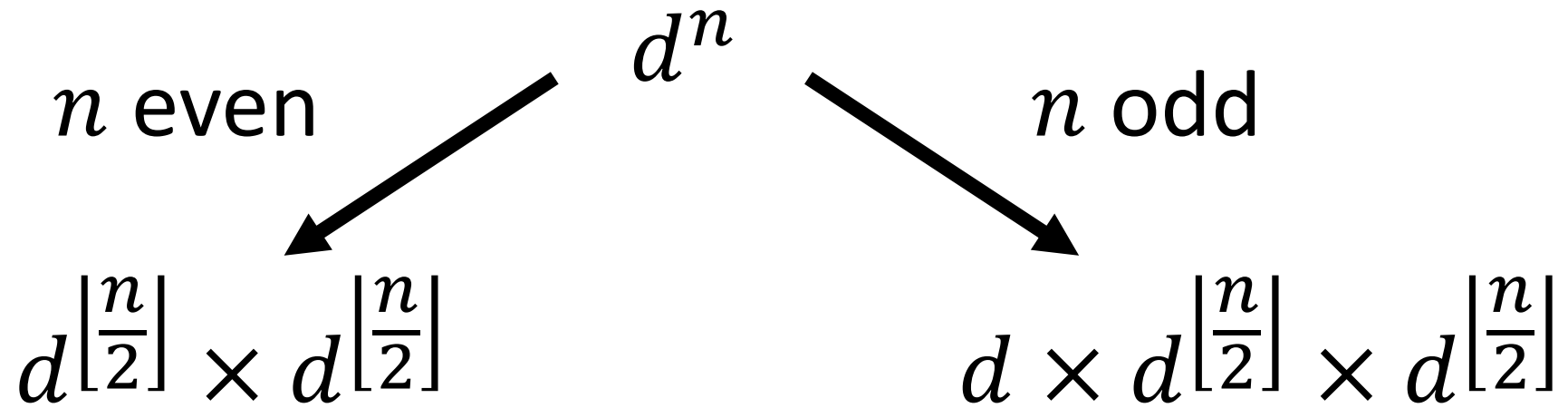
مشكلة حساب العامل

$O(n)$ but why?



n times

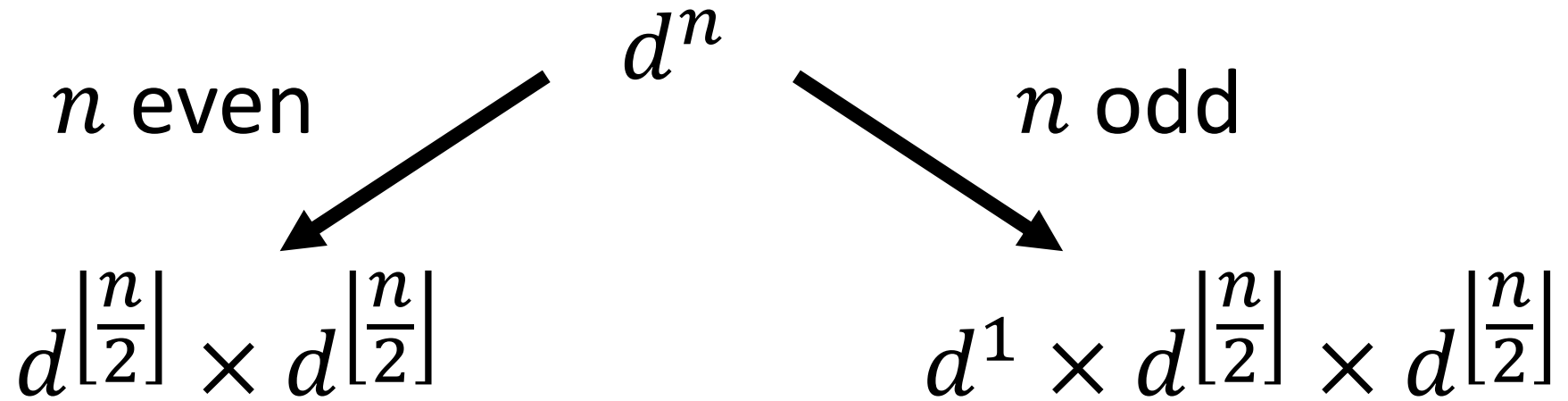
Super Power



Ex: $d^4 = d^2 \times d^2$

$d^5 = d \times d^2 \times d^2$

Super Power



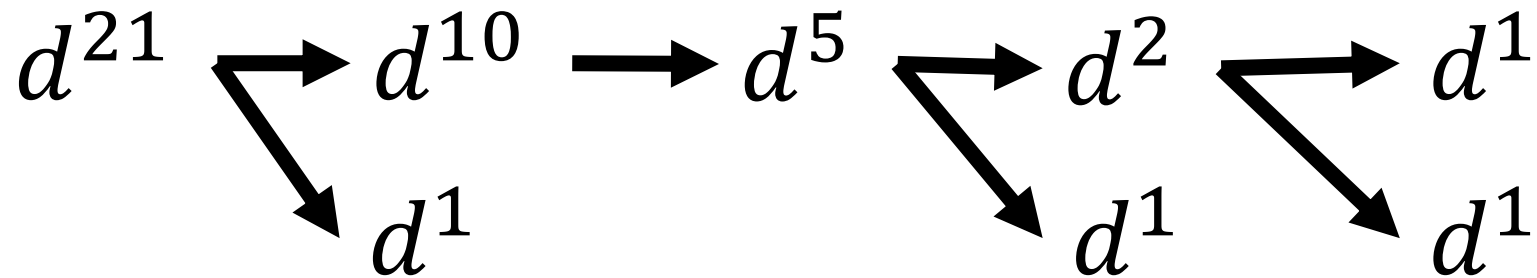
```
11 power(11 d, 11 n) {  
    if (n==0)  
        return 1;  
    11 x = power(d, n/2);  
    if (n%2==1)  
        return x*x*d;  
    return x*x;  
}
```

$O(\log_2 n)$ but why?

Super Power

$O(\log_2(n))$ but why?

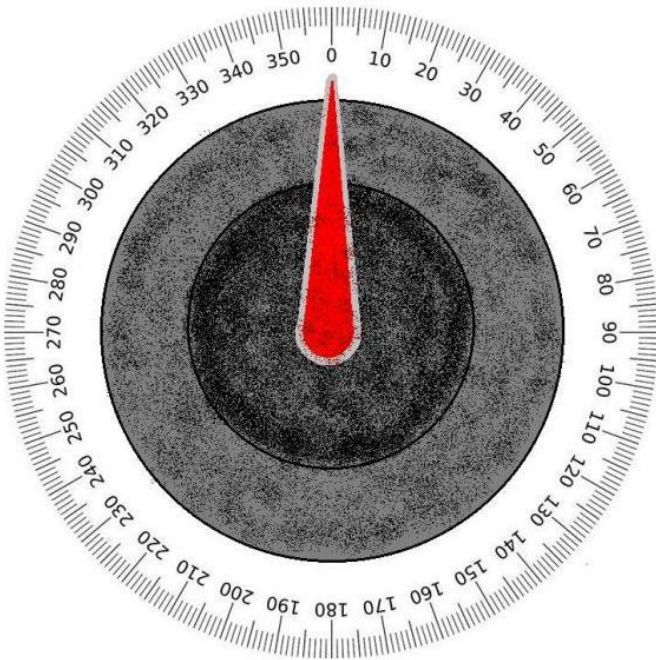
$$n = 10$$



$$n \times \frac{1}{2} \times \frac{1}{2} \times \frac{1}{2} \times \dots = 1 \Rightarrow \frac{n}{2^m} = 1 \Rightarrow \boxed{m = \log_2(n)}$$

B. Petr and a Combination Lock

<https://codeforces.com/contest/1097/problem/B>



Max number of rotations is **15**

$$\begin{array}{r} 10 \\ 20 \\ 30 \\ \hline 0 \end{array}$$

→ YES

$$\begin{array}{r} 10 \\ 10 \\ 10 \\ \hline \neq 0 \end{array}$$

→ NO

B. Petr and a Combination Lock

<https://codeforces.com/contest/1097/problem/B>

Can sum equal X ??

$$\left\{ \begin{array}{c} \pm a_i \\ a_{i+1} \\ a_{i+2} \\ \vdots \\ a_n \end{array} \right\}$$

Can sum equal $X + (-a_i)$ **OR** $X + a_i$??

B. Petr and a Combination Lock

<https://codeforces.com/contest/1097/problem/B>

`possible(0, 0)`

```
bool possible(int i, int X) {  
    if (i == N)  
        return (X % 360) == 0;  
    return possible(i + 1, X + a[i]) || possible(i + 1, X - a[i]);  
}
```

B. Tavas and SaDDas

<https://codeforces.com/problemset/problem/535/B>

The number n is a **Lucky Number** if their decimal representation doesn't contain digits other than **4s** and **7s**.

Ex: 74, 774, 47, 44, 777....

If we **sort** all lucky numbers in **increasing** order, what's the **1-based index** of n ?

Ex: 4 \rightarrow 1

 7 \rightarrow 2

 444 \rightarrow 7

Where n is **Lucky Number**

B. Tavas and SaDDas

<https://codeforces.com/problemset/problem/535/B>

```
set <ll> numbers;  
void f(ll n, ll i) {  
    if (i==10) {  
        return;  
    } else {  
        numbers.insert((n*10)+4);  
        numbers.insert((n*10)+7);  
        f((n*10)+4, i+1);  
        f((n*10)+7, i+1);  
    }  
}
```

f(0, 0);

B. Lucky Numbers (easy)

<https://codeforces.com/problemset/problem/96/B>

The number n is a **Super Lucky Number** if it is a **Lucky Number** and contains **equal** amount of digits 4 and 7.

Ex: 74, 47, 4747, 774474....

B. Lucky Numbers (easy)

<https://codeforces.com/problemset/problem/96/B>

Q: you will get number n .. Find **minimum** number x so that x is a **Super Lucky Number** and $x \geq n$:

Input: 4500

Output: 4747

$$1 \leq n \leq 10^9$$

Input: 70

Output: 74

B. Lucky Numbers (easy)

<https://codeforces.com/problemset/problem/96/B>

```
ll ans = 1e18;
void f(ll sum, ll n4, ll n7)
{
    if(sum > 1e10)
        return;
    if(sum >= n && n4 == n7)
        ans = min(sum, ans);
    f(sum*10+4, n4+1, n7);
    f(sum*10+7, n4, n7+1);
}
```

`f(0, 0, 0);`

B. Lucky Numbers (easy)

<https://codeforces.com/problemset/problem/96/B>

```
long long f(long long sum, long long n4, long long n7)
{
    if(sum > 1e10)
        return 1e10;
    if(sum >= n && n4 == n7)
        return sum;
    return min(f(sum*10+4, n4+1, n7), f(sum*10+7, n4, n7+1));
}
```

`f(0, 0, 0);`

B. Preparing Olympiad

<https://codeforces.com/problemset/problem/550/B>

```
#include<bits/stdc++.h>
using namespace std;
int n, L, R, x;
int a[1000];
vector < vector<int> > subsets;
int solve(int index, int sum, int mn, int mx, vector<int>& cur) {
    // base-case
    if (index == n) {
        if (sum >= L && sum <= R && mx - mn >= x) {
            subsets.push_back(cur);
            return 1;
        }
        return 0;
    }
    // calls
    // pick
    cur.push_back(a[index]);
    int pick = solve(index + 1, sum + a[index], min(mn, a[index]), max(mx, a[index]), cur);
    cur.pop_back();
```

B. Preparing Olympiad

<https://codeforces.com/problemset/problem/550/B>

```
// leave
    int leave = solve(index + 1, sum, mn, mx, cur);
    return pick + leave;
}

int main(){
    ios::sync_with_stdio(0);
    cin >> n >> L >> R >> x;
    for (int i = 0; i < n; ++i)
        cin >> a[i];
    vector<int> cur;
    cout << solve(0, 0, 1e9, -1e9, cur) << '\n';
    cout << "subsets: \n";
    for (vector<int> vec: subsets) {
        for (int x : vec) cout << x << ' ';
        cout << '\n';
    }
    return 0;
}
```

D. a-Good String

<https://codeforces.com/problemset/problem/1385/D>

String S is a c -Good String if:

- The First half if it is a $(c + 1)$ -Good String and the second half is " $cccc\dots$ "
- The Second half if it is a $(c + 1)$ -Good String and the first half is " $cccc\dots$ "
- String S is one string character c

Ex: String $S = \text{"*eeehgff*"}$ is e -Good String

D. a-Good String

<https://codeforces.com/problemset/problem/1385/D>

$S = "e e e e h g f f"$

The diagram illustrates the decomposition of the string $S = "e e e e h g f f"$ into three substrings:

- g-Good:** The substring $h g$ is identified as g-Good.
- f-Good:** The substring $h g f f$ is identified as f-Good.
- e-Good String:** The entire string S is identified as an e-Good String.

D. a-Good String

<https://codeforces.com/problemset/problem/1385/D>

$S = "b\ b\ c\ d\ a\ a\ a\ a"$

The diagram illustrates the nested structure of the string $S = "b\ b\ c\ d\ a\ a\ a\ a"$. It uses three levels of brackets to show how the string is composed of 'a-Good' and 'b-Good' substrings. The innermost bracket, labeled 'c-Good', spans the characters 'c' and 'd'. The middle bracket, labeled 'b-Good', spans the characters 'b', 'b', 'c', and 'd'. The outermost bracket, labeled 'a-Good String', spans the entire string from the first 'b' to the last 'a'.

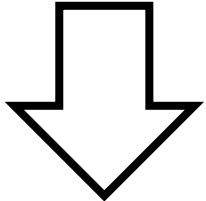
D. a-Good String

<https://codeforces.com/problemset/problem/1385/D>

Q: what is the minimum character we need to change to make S a **a**-Good String

Ex: $S = "b \textcolor{red}{o} c d a a \textcolor{red}{z} a"$

Size of the S
 $\leq 131\,072$

 2 Changes

$S = "b \textcolor{red}{b} c d a a \textcolor{red}{a} a"$

D. a-Good String

<https://codeforces.com/problemset/problem/1385/D>

```
int solve(int l,int r,char ch) {  
    if(l==r) return ch!=s[l];  
    int mid=(l+r)/2;  
    int cntl=0,cntr=0;  
    for(int i=l;i<=r;i++) {  
        if(i<=mid) cntl+=(ch!=s[i]);  
        else cntr+=(ch!=s[i]);  
    }  
    return min(cntl+solve(mid+1,r,ch+1),cntr+solve(1,mid,ch+1));  
}
```

```
cout<<solve(0,n-1,'a')<<endl;
```

next_permutation

```
#include<bits/stdc++.h>
using namespace std;
int main()
{
    ios::sync_with_stdio(0);
    int n; cin >> n;
    vector <int> prem;
    for (int i = 1; i <= n; ++i)
        prem.push_back(i);
    do {
        for (int x: prem)
            cout << x << ' ';
        cout << '\n';
    } while (next_permutation(prem.begin(), prem.end()));

    return 0;
}
```