

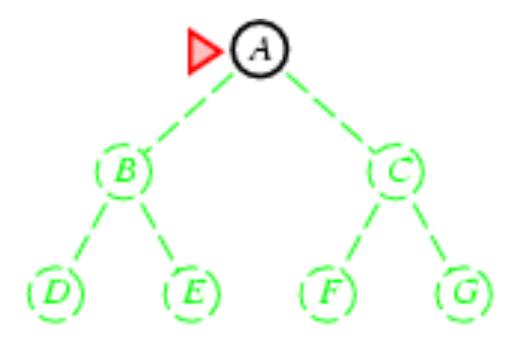
Uninformed Search



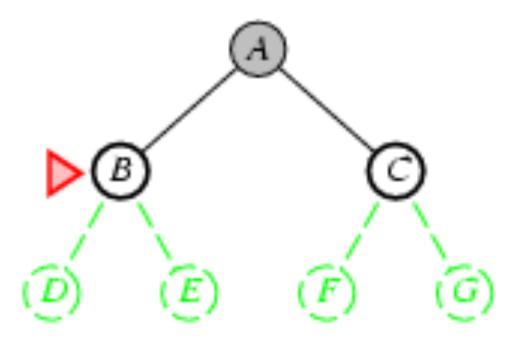
Uninformed search strategies

- ☐ Uninformed search strategies use only the information available in the problem definition
 - Breadth-first search/ Búsqueda en anchura
 - Uniform-cost search/ Búsqueda de coste uniforme
 - Depth-first search/ Búsqueda en profundidad
 - Depth-limited search/ Búsqueda en profundidad limitada
 - Iterative deepening search/Búsqueda de profundización iterativa

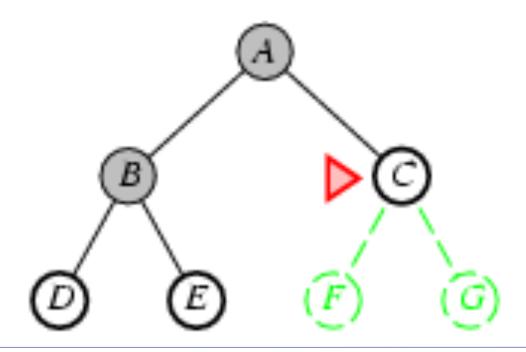
- Expand shallowest unexpanded node
- □ Implementation:
 - fringe is a FIFO queue, i.e., new successors go at end



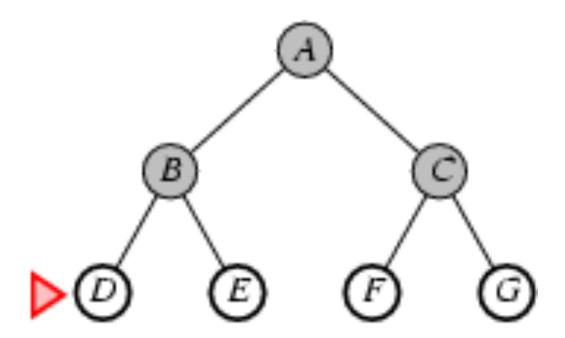
- Expand shallowest unexpanded node
- ☐ Implementation:
 - fringe is a FIFO queue, i.e., new successors go at end



- Expand shallowest unexpanded node
- ☐ Implementation:
 - fringe is a FIFO queue, i.e., new successors go at end



- Expand shallowest unexpanded node
- □ Implementation:
 - fringe is a FIFO queue, i.e., new successors go at end



Properties of breadth-first search

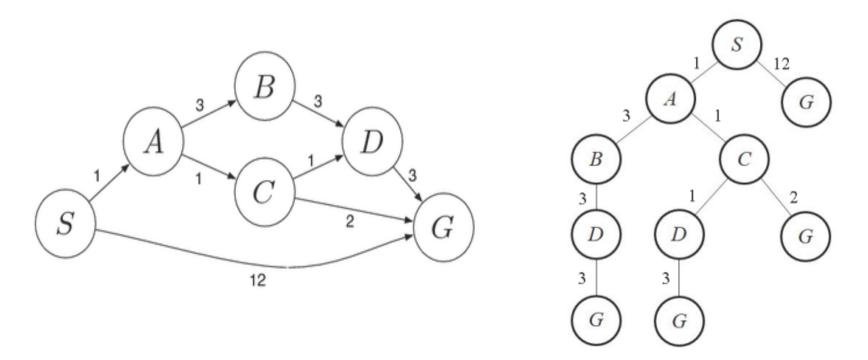
- Complete? Yes (if b is finite)
- \square Time? $1+b+b^2+b^3+...+b^d+b(b^d-1)=O(b^{d+1})$
- \square Space? $O(b^{d+1})$ (keeps every node in memory)
- Optimal? Yes (if cost = 1 per step)

Space is the bigger problem (more than time)

Each state has b successors (branching factor) d is the shallower depth

Uniform-cost search

- □ Expand least-cost unexpanded node
- Implementation:
 - fringe = queue ordered by path cost
- Find the solution with minimum cumulative cost, i.e. an optimal solution



Uniform-cost search (Solution)

```
Initialization: { [ S , 0 ] }

Iteration1: { [ S->A , 1 ] , [ S->G , 12 ] }

Iteration2: { [ S->A->C , 2 ] , [ S->A->B , 4 ] , [ S->G , 12 ] }

Iteration3: { [ S->A->C->D , 3 ] , [ S->A->C->G , 4 ] , [ S->A->B->D , 7 ] , [ S->G , 12 ] }

Iteration 4: { [ S->A->C->D->G , 6 ] , [ S->A->C->G , 4 ] , [ S->A->B->D , 7 ] , [ S->G , 12 ] }

Iteration 5: { [ S->A->C->G , 4 ] , [ S->A->C->D->G , 6 ] , [ S->A->B->D->G , 10 ] , [ S->G , 12 ] }

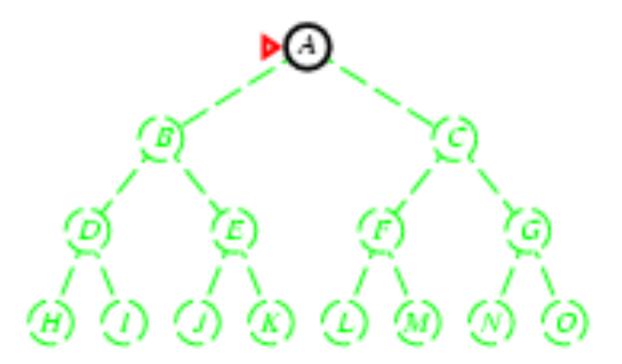
Solution: S->A->C->G.
```

Uniform-cost search

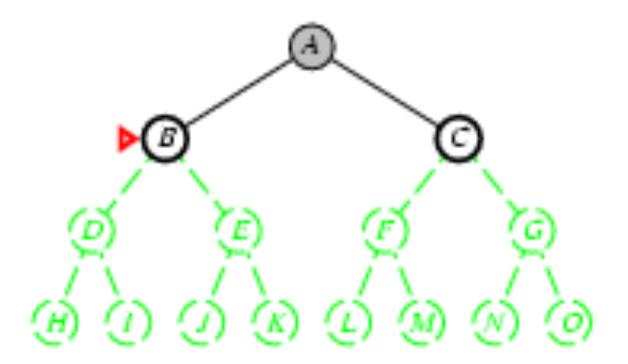
- Complete? Yes, if step cost ≥ ε
- Time? # of nodes with $g \le cost$ of optimal solution, $O(b^{ceiling(C^*/\varepsilon)})$ where C^* is the cost of the optimal solution
- □ Space? # of nodes with $g \le cost$ of optimal solution, $O(b^{ceiling(C^*/ε)})$
- Optimal? Yes nodes expanded in increasing order of g(n)

If all costs are equal → O(bd)

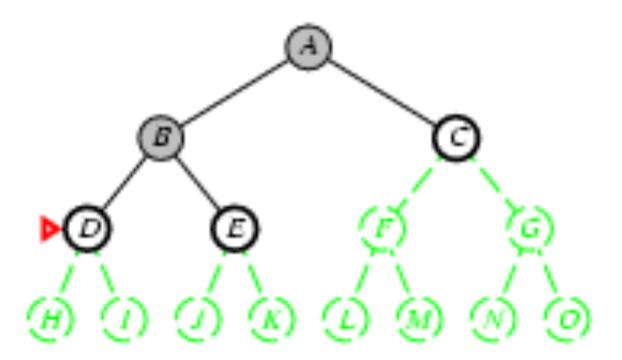
- □ Expand deepest unexpanded node
- Implementation:
 - fringe = LIFO queue, i.e., put successors at front



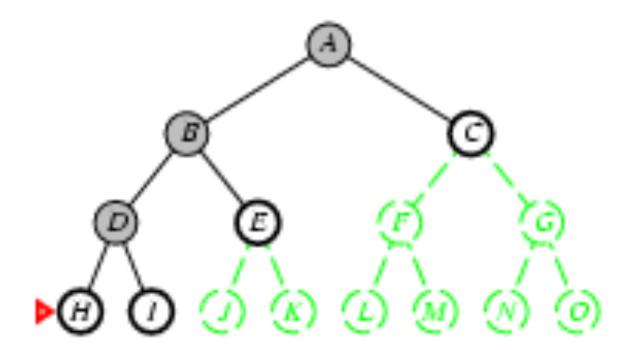
- □ Expand deepest unexpanded node
- Implementation:
 - fringe = LIFO queue, i.e., put successors at front



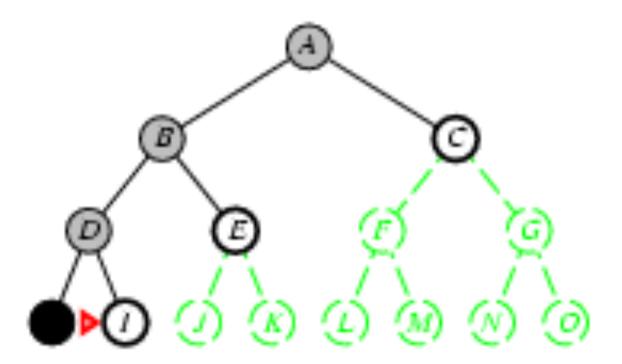
- □ Expand deepest unexpanded node
- Implementation:
 - fringe = LIFO queue, i.e., put successors at front



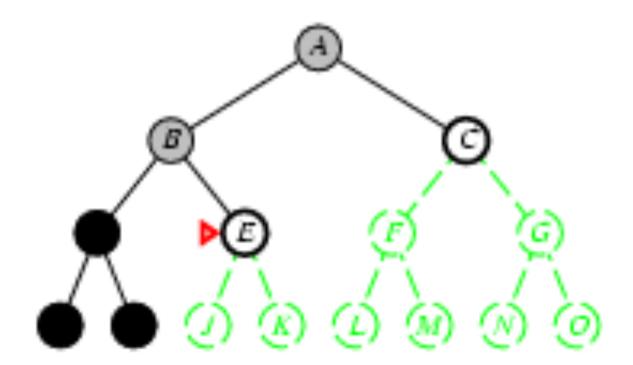
- □ Expand deepest unexpanded node
- Implementation:
 - fringe = LIFO queue, i.e., put successors at front



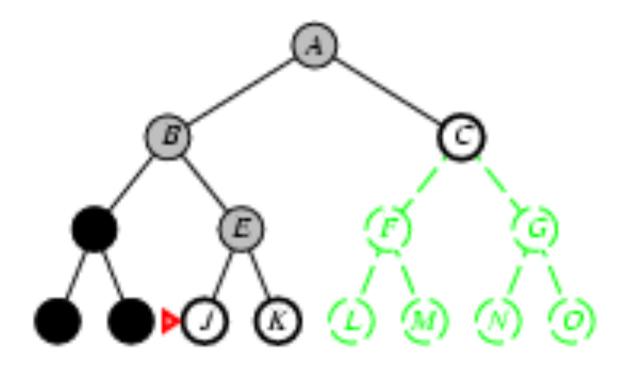
- □ Expand deepest unexpanded node
- Implementation:
 - fringe = LIFO queue, i.e., put successors at front



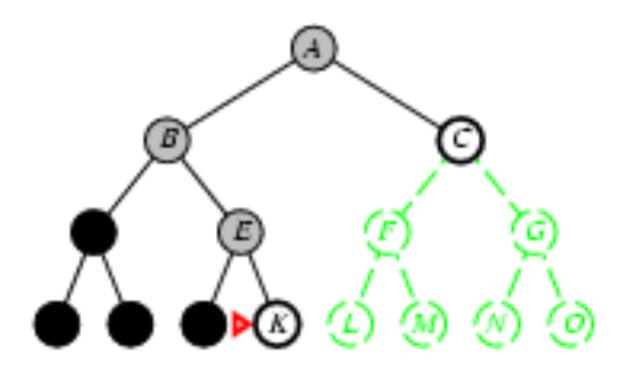
- □ Expand deepest unexpanded node
- Implementation:
 - fringe = LIFO queue, i.e., put successors at front



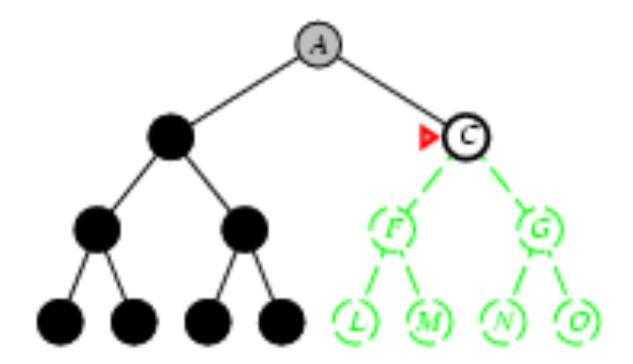
- □ Expand deepest unexpanded node
- Implementation:
 - fringe = LIFO queue, i.e., put successors at front



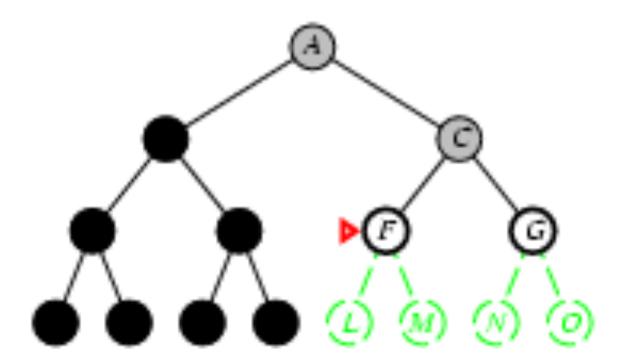
- □ Expand deepest unexpanded node
- Implementation:
 - fringe = LIFO queue, i.e., put successors at front



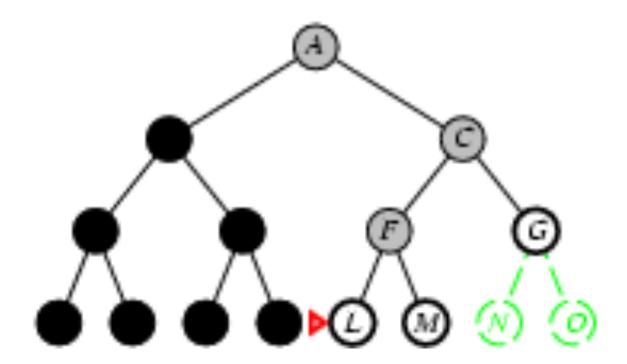
- □ Expand deepest unexpanded node
- Implementation:
 - fringe = LIFO queue, i.e., put successors at front



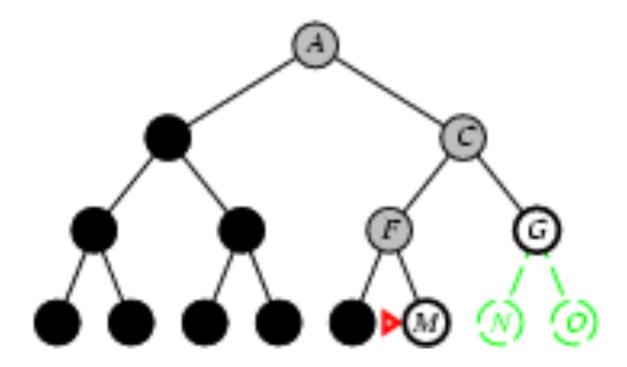
- □ Expand deepest unexpanded node
- Implementation:
 - fringe = LIFO queue, i.e., put successors at front



- □ Expand deepest unexpanded node
- Implementation:
 - fringe = LIFO queue, i.e., put successors at front



- □ Expand deepest unexpanded node
- Implementation:
 - fringe = LIFO queue, i.e., put successors at front



Properties of depth-first search

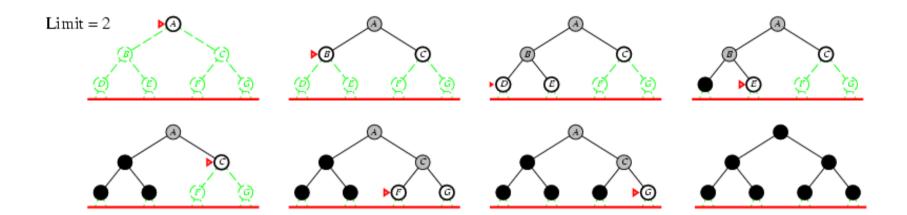
- Complete? No: fails in infinite-depth spaces, spaces with loops
 - Modify to avoid repeated states along path→ complete in finite spaces
- \square Time? $O(b^m)$: terrible if m is much larger than d
 - but if solutions are dense, may be much faster than breadth-first
- □ Space? O(bm)
- Optimal? No

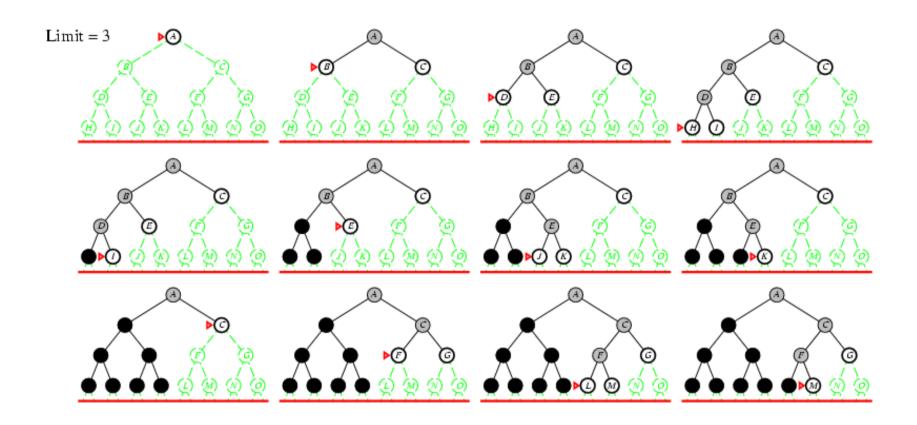
Depth-limited search

- = depth-first search with depth limit L
- ☐ i.e., nodes at depth L have no successor







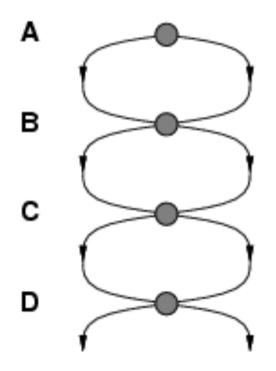


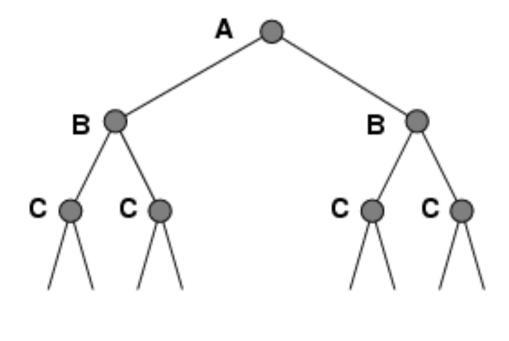
Properties of iterative deepening search

- □ Complete? Yes
- □ Space? O(bd)
- □ Optimal? Yes, if step cost = 1

Repeated states

☐ Failure to detect repeated states can turn a linear problem into an exponential one!





Summary of algorithms

Criterion	Breadth-	Uniform-	Depth-	Depth-	Iterative
	First	Cost	First	Limited	Deepening
Complete?	Yes	Yes	No	No	Yes
Time	$O(b^{d+1})$	$O(b^{\lceil C^*/\epsilon ceil})$	$O(b^m)$	$O(b^l)$	$O(b^d)$
Space	$O(b^{d+1})$	$O(b^{\lceil C^*/\epsilon ceil})$	O(bm)	O(bl)	O(bd)
Optimal?	Yes	Yes	No	No	Yes



Uninformed Search



□ Number of nodes generated in a depth-limited search to depth d with branching factor b:

$$N_{DLS} = b^0 + b^1 + b^2 + \dots + b^{d-2} + b^{d-1} + b^d$$

■ Number of nodes generated in an iterative deepening search to depth d with branching factor b:

$$N_{IDS} = (d+1)b^0 + db^{-1} + (d-1)b^{-2} + ... + 3b^{d-2} + 2b^{d-1} + 1b^d$$

- \Box For b = 10, d = 5,
 - $N_{DLS} = 1 + 10 + 100 + 1,000 + 10,000 + 100,000 = 111,111$
 - $N_{IDS} = 6 + 50 + 400 + 3,000 + 20,000 + 100,000 = 123,456$
- Overhead = (123,456 111,111)/111,111 = 11%