

# Path Planning

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# Outline

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- ☐ **Introduction**
- ☐ Dijkstra
- ☐ A\*
- ☐ A\*PS
- ☐ Theta\*
- ☐ S-Theta\*
- ☐ Conclusions

# Introduction

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- Path-planning problem aims to obtain feasible and optimal (or near to it) routes between two or more points
  - Find optimal (or near to it) paths is not trivial
  - Feasible implies not transverse over obstacles or overcome system limitations
  - Usually parameters: path length, run-time, expanded nodes and number of heading changes
- It is a fundamental task in mobile robots and video games

# Introduction

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- First point to address: the environment

- Local planning vs Long term planning



- Totally observable vs Partially observable



- Discrete vs Continuous



- Dynamic environment? Maybe

- Extra information on the terrain? Maybe

# Introduction

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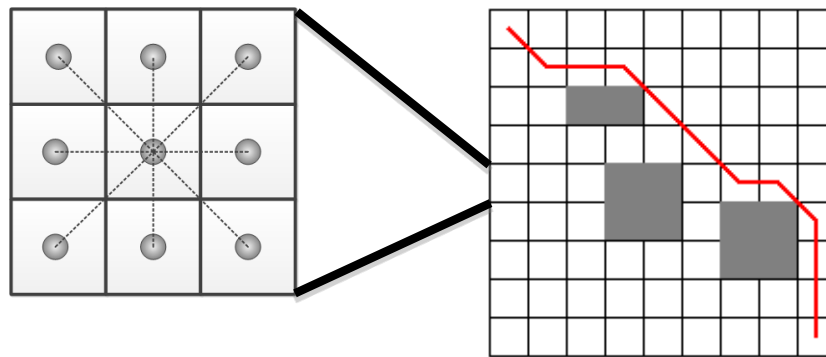
- Long term path-planning is more related to AI
- Easier with fully observable environment
- High effort trying to obtain optimal paths
- In some domains (such as planetary exploration) path-planning and task-planning are highly coupled

# Introduction

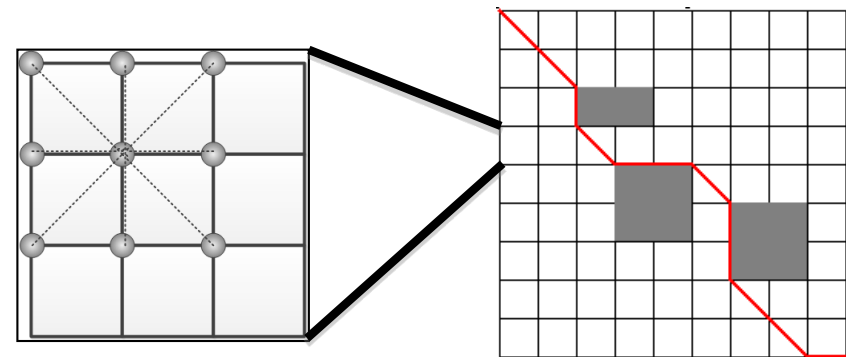
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- ❑ Classical path-planning algorithms are based on A\* heuristic search algorithm
- ❑ Works over 2D grids with blocked and unblocked cells
- ❑ Nodes are (usually) 8-connected with its neighbors
- ❑ Two representations:

Center node



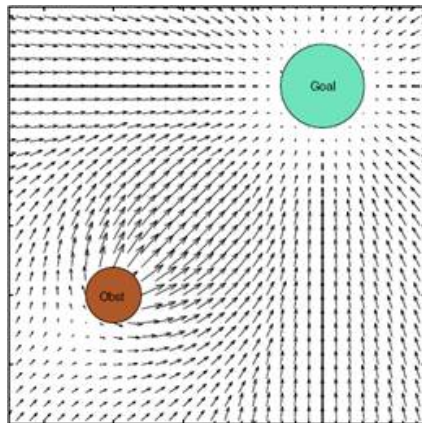
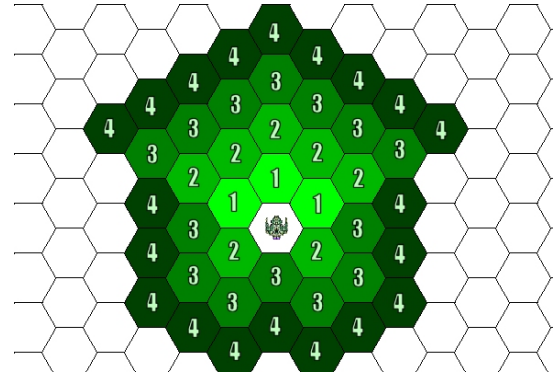
Corner node



# Introduction

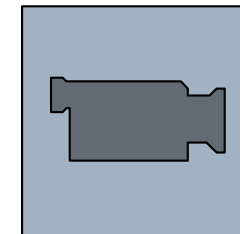
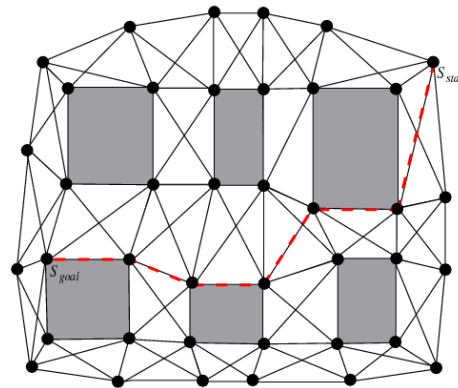
## □ Other representations:

### ■ Hexagonal regions



### ◆ Potential fields

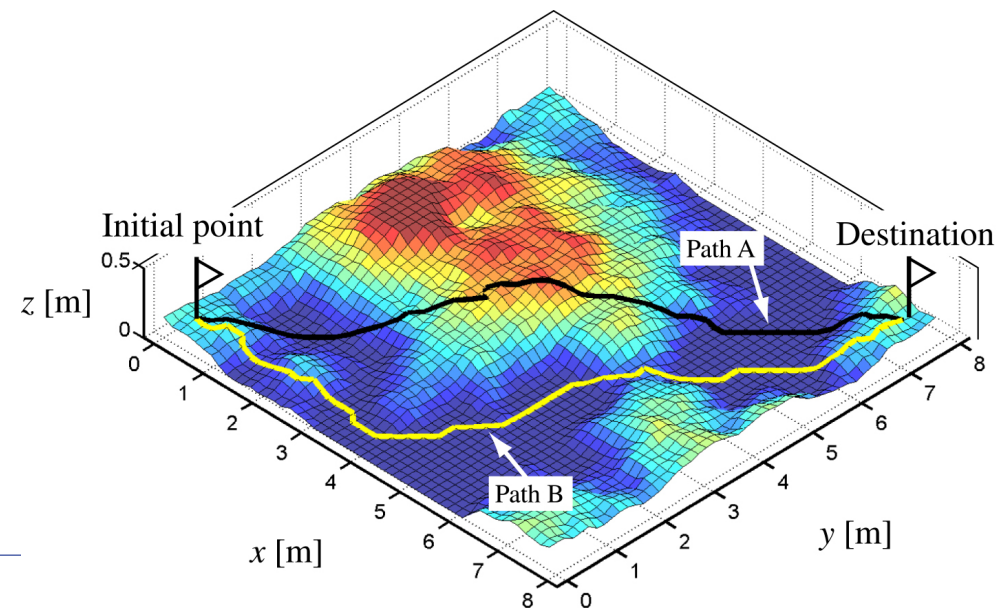
### ■ Visibility graphs



# Introduction

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- For dynamic environments usually a replanning strategy is followed
- Optimize the replanning process?
- What can be taken into consideration into the terrain?
  - Altitude → DEM
  - Hazardous areas





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- ☐ **Dijkstra**
- ☐  $A^*$
- ☐  $A^*PS$
- ☐ Theta\*
- ☐ S-Theta\*
- ☐ Conclusions

# Dijkstra

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- Pick the unvisited vertex with the lowest-distance
- Calculate the distance through it to each unvisited neighbor
- Update the neighbor's distance if smaller
- Mark visited when done with neighbors



# Dijkstra

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```
1  function Dijkstra(Graph, source):
2      dist[source] := 0                // Initializations
3      for each vertex v in Graph:
4          if v ≠ source
5              dist[v] := infinity      // Unknown distance from source to v
6              previous[v] := undefined // Predecessor of v
7          end if
8          PQ.add_with_priority(v, dist[v])
9      end for
10
11
12     while PQ is not empty:           // The main loop
13         u := PQ.extract_min()        // Remove and return best vertex
14         for each neighbor v of u:   // where v has not yet been removed from PQ.
15             alt = dist[u] + length(u, v)
16             if alt < dist[v]        // Relax the edge (u,v)
17                 dist[v] := alt
18                 previous[v] := u
19                 PQ.decrease_priority(v, alt)
20             end if
21         end for
22     end while
23     return previous[]
```

# Outline

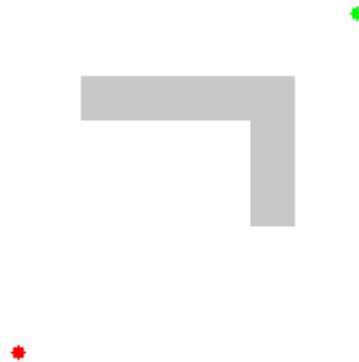
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# A\*

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- A\* makes guided search using two values:
  - Accumulate cost ( $G(t)$ ): cost to reach a node
  - Heuristic ( $H(t)$ ): predicted cost to achieve goal from a node (Euclidian distance, Octal distance)
- Node Evaluation:  $F(t) = G(t) + H(t)$
- A\* is simple, fast and guarantee optimal paths in eight-connected grids
- Artificially restricted to  $45^\circ$  headings



# A\*

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**Algorithm 1** A\* search

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```
1   $G(s) \leftarrow 0$ 
2   $parent(s) \leftarrow s$ 
3   $open \leftarrow \emptyset$ 
4   $open.insert(s, G(s), H(s))$ 
5   $closed \leftarrow \emptyset$ 
6  while  $open \neq \emptyset$  do
7       $p \leftarrow open.pop()$ 
8      if  $p = g$  then
9          return  $path$ 
10     end if
11      $closed.insert(p)$ 
12     for  $t \in neighbours(p)$  do
13         if  $t \notin closed$  then
14             if  $t \notin open$  then
15                  $G(t) \leftarrow \infty$ 
16                  $parent(t) \leftarrow null$ 
17             end if
18              $UpdateVertex(p, t)$ 
19         end if
20     end for
21 end while
22 return  $fail$ 
```

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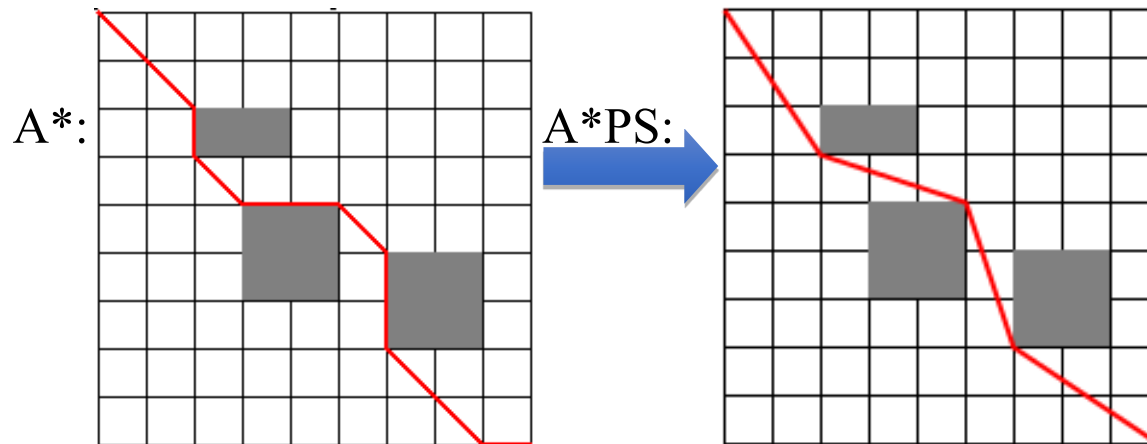
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# A\* Post Processed

- A\* Post Processed (A\*PS) tries to smooth A\* routes by removing intermediate nodes when there is line of sight between 2 no neighbors nodes, that is, there are no obstacles





# Outline

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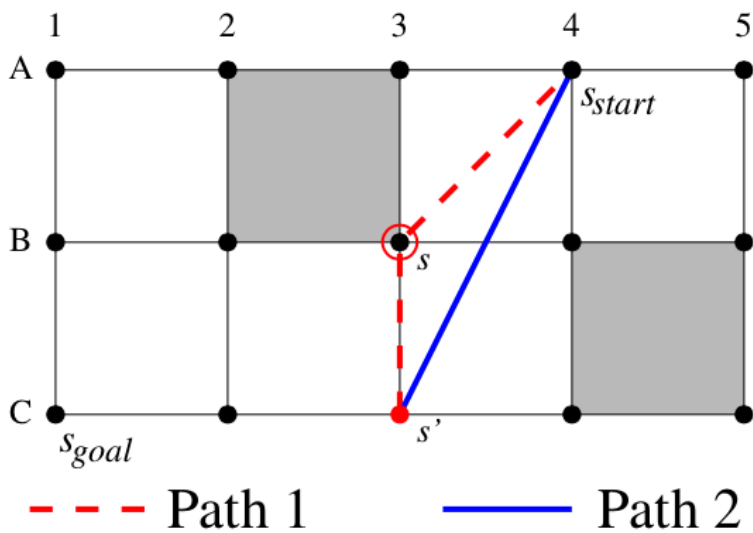
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- ☐ Practice

# Theta\*

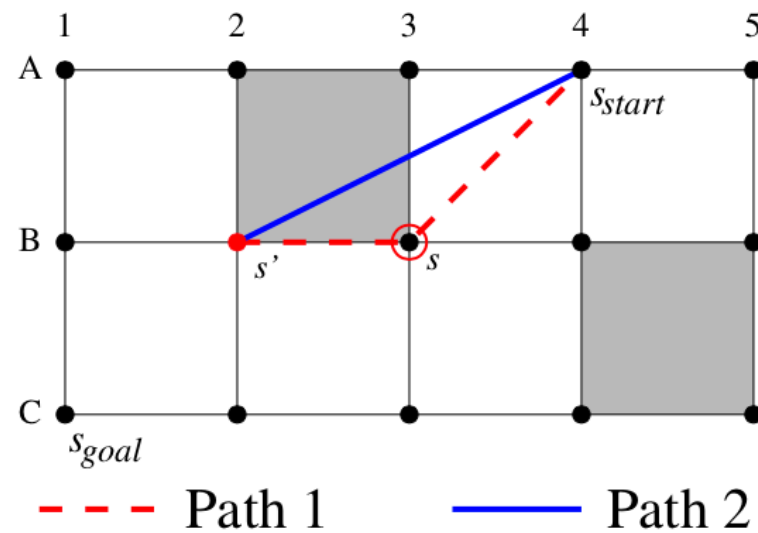
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- ❑ Theta\* is a variation of A\* that integrates the line of sight check during search
- ❑ For this reason Theta\* is not restricted to 45° headings, and gets more realistic paths than A\* without post processing
- ❑ Theta\* is slower than A\* due to line of sight calculation, but paths are shorter and smoothest

# Theta\*



(a) Path 2 is unblocked



(b) Path 2 is blocked

# Theta\*

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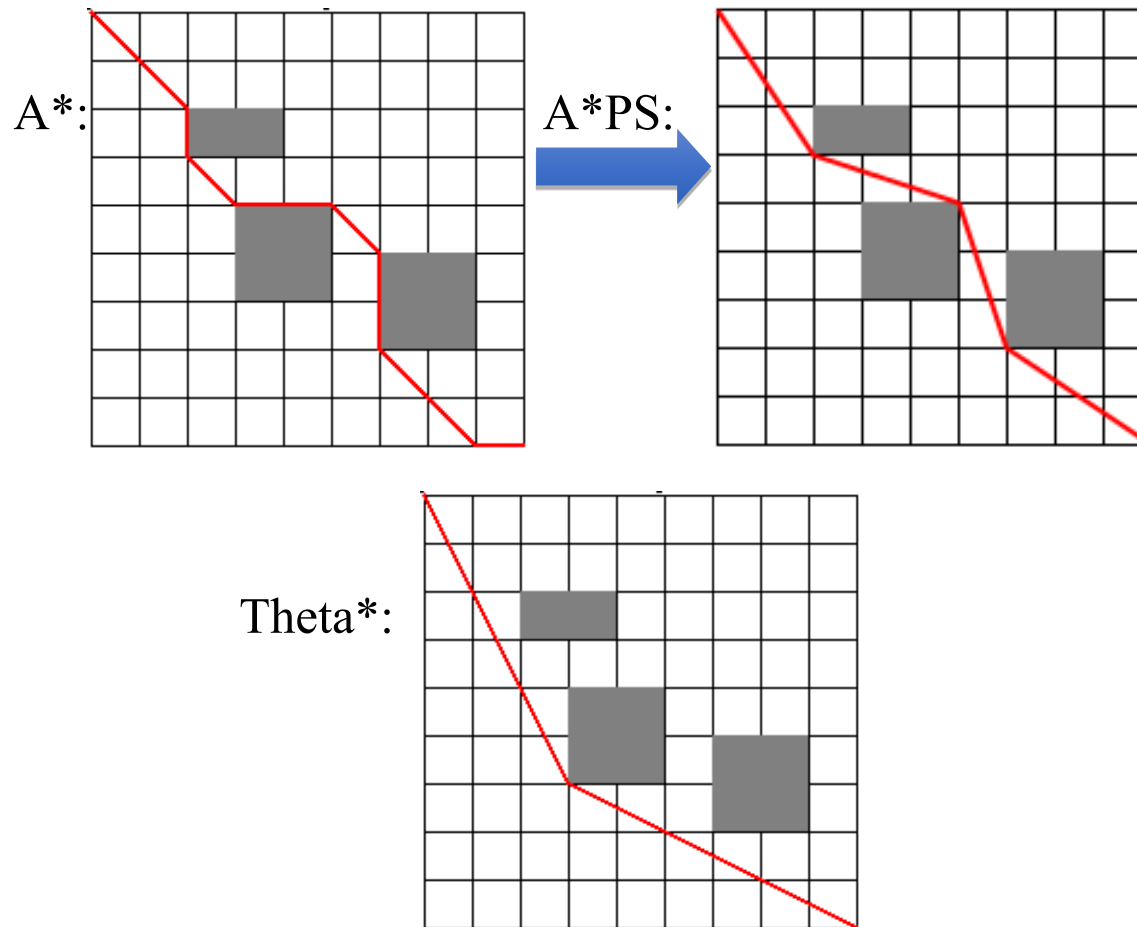
**Algorithm 2** Update vertex function for Basic Theta\*

```
1  UpdateVertex(p, t)
2  if LineOfSight(parent(p), t) then
3      if  $G(\text{parent}(p)) + \text{dist}(\text{parent}(p), t) < G(t)$  then
4           $G(t) \leftarrow G(\text{parent}(p)) + \text{dist}(\text{parent}(p), t)$ 
5           $\text{parent}(t) \leftarrow \text{parent}(p)$ 
6          if  $t \in \text{open}$  then
7              open.remove(t)
8          end if
9          open.insert(t,  $G(t)$ ,  $H(t)$ )
10     end if
11 else
12     if  $G(p) + \text{dist}(p, t) < G(t)$  then
13          $G(t) \leftarrow G(p) + \text{dist}(p, t)$ 
14          $\text{parent}(t) \leftarrow p$ 
15         if  $t \in \text{open}$  then
16             open.remove(t)
17         end if
18         open.insert(t,  $G(t)$ ,  $H(t)$ )
19     end if
20 end if
```

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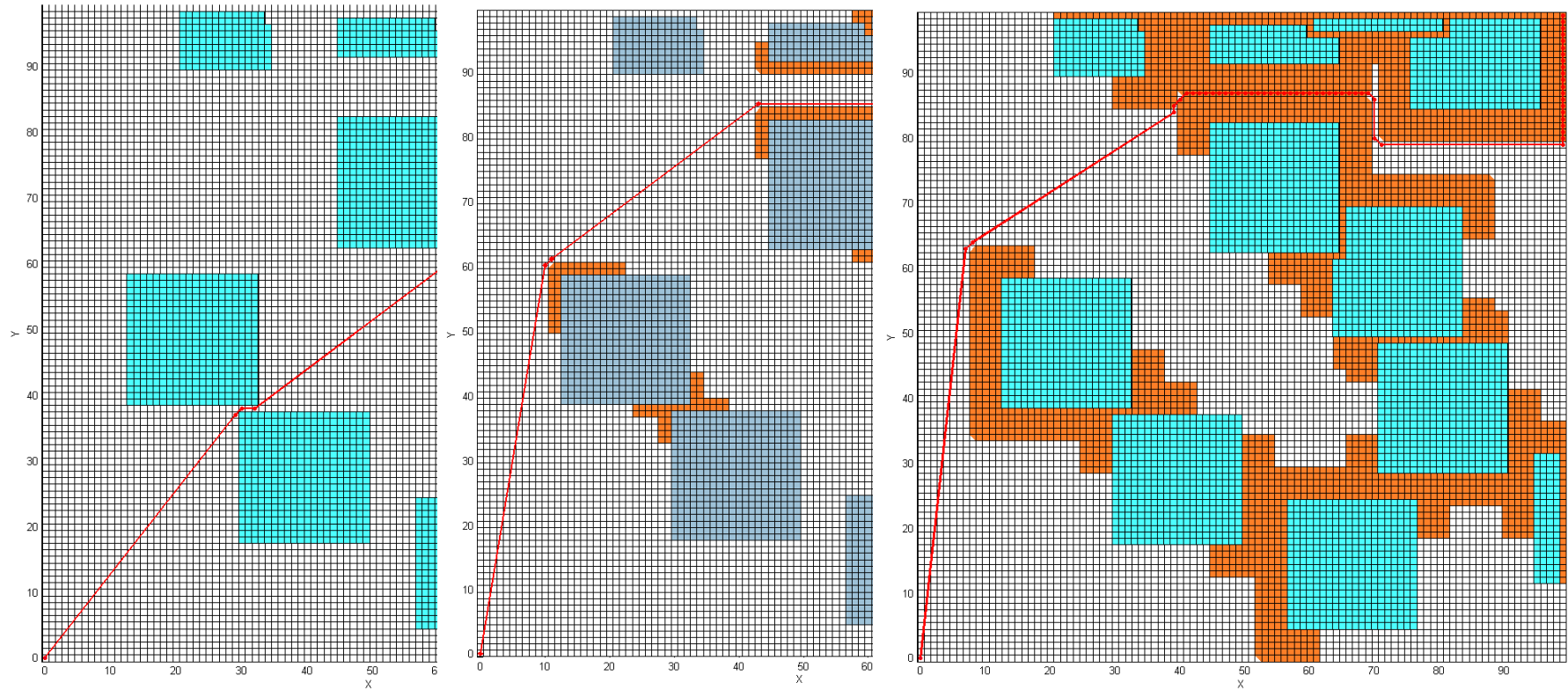
# Theta\*

- A\*PS and Theta\* are any-angle algorithms: not restricted to 45° headings



# Theta\*

- What about crossing the obstacles at the corner?  
→ Safety margin



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# S-Theta\*

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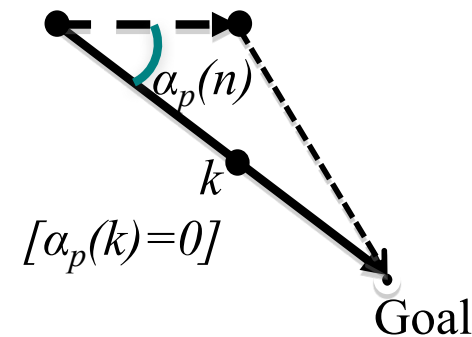
- Theta\* updates a node depending on the distance to reach it, regardless of its orientation
- Adapt Theta\* to take into consideration heading changes during the search process
  - Robotics hardware is usually very limited
  - Rotation cost is greater than the movement straight
- Best path between two points in a free obstacle grid is straight line
- Achieve less heading changes in exchange of a slight degradation of the path length



# S-Theta\*

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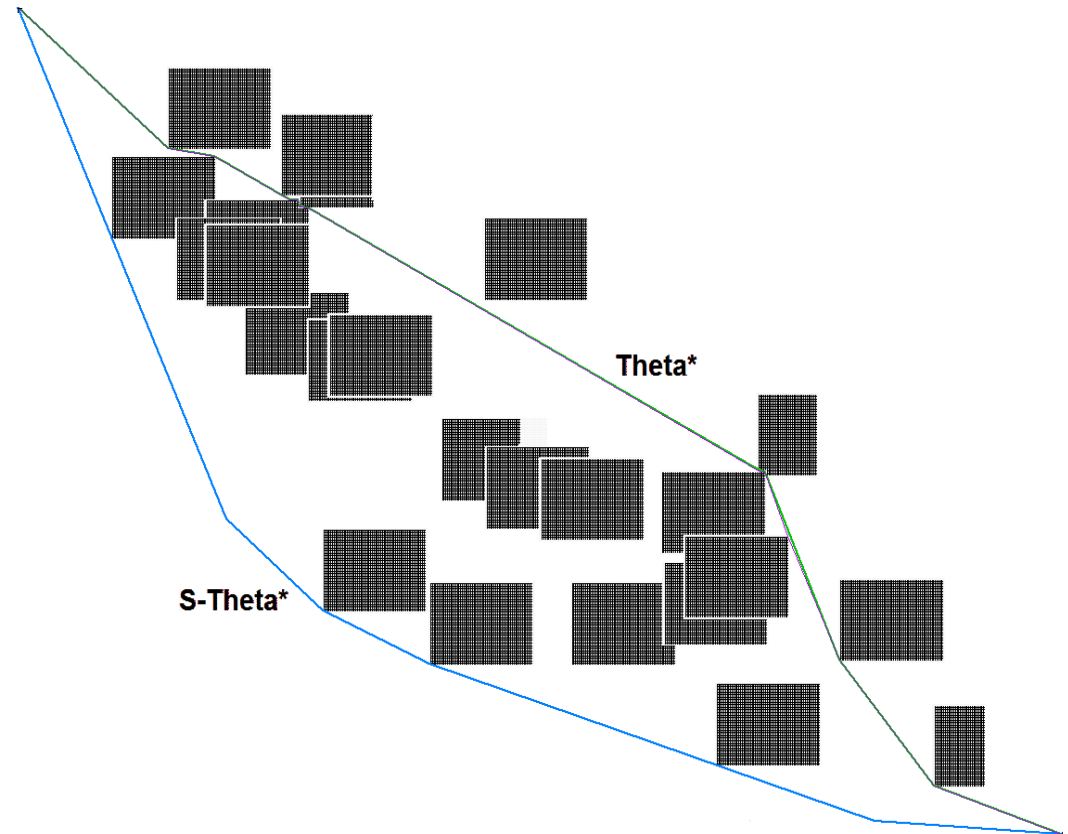
- Include a new term in the cost function:  $\alpha(n)$  that represents the heading change variation to reach a node  $n$  in relationship with the objective and previous nodes
- This term guides the search process to:
  - Smooth heading changes
  - Reduce number of heading changes
- $F(t) = G(t) + H(t) + \alpha(n)$



# S-Theta\*

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- $\alpha(n)$  tries to surround obstacles and return the best path
- Does not expand nodes far from the line
- Need to weight
  - $\alpha(n) = \alpha(n) \times N$
  - $\alpha(n)$  is  $[0^\circ, 180^\circ]$



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# Conclusions

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- ❑ Classical path-planning based on informed search methods provides a good approximation for big observable areas
- ❑ Heuristics are very relevant
- ❑ Representation of the environment is an important point to consider
- ❑ Large number of works, but still possibility to improve
  - Field D\*, Block A\*, Lazy Theta\*...