

**Solution 1.a:**

$$\begin{aligned} \text{Denote } K_1(x, y) &= \varphi_1(x)^T \cdot \varphi_1(y) = [\varphi_1(x)_1, \varphi_1(x)_2, \dots, \varphi_1(x)_N] \begin{bmatrix} \varphi_1(y)_1 \\ \vdots \\ \varphi_1(y)_N \end{bmatrix} \\ &= \varphi_1(x)_1 \cdot \varphi_1(y)_1 + \varphi_1(x)_2 \cdot \varphi_1(y)_2 + \dots + \varphi_1(x)_N \cdot \varphi_1(y)_N = S_1 \end{aligned}$$

$$\begin{aligned} \text{Denote } K_2(x, y) &= \varphi_2(x)^T \cdot \varphi_2(y) = [\varphi_2(x)_1, \varphi_2(x)_2, \dots, \varphi_2(x)_K] \begin{bmatrix} \varphi_2(y)_1 \\ \vdots \\ \varphi_2(y)_K \end{bmatrix} \\ &= \varphi_2(x)_1 \cdot \varphi_2(y)_1 + \varphi_2(x)_2 \cdot \varphi_2(y)_2 + \dots + \varphi_2(x)_K \cdot \varphi_2(y)_K = S_2 \end{aligned}$$

Now we want to show that :

$$K(x, y) = \varphi(x)^T \cdot \varphi(y) = 5K_1(x, y) + 4K_2(x, y) = 5S_1 + 4S_2$$

We define the following mapping –  $\varphi: \mathbb{R}^n \rightarrow \mathbb{R}^{N+K}$  s.t

$$\varphi(x) = [\sqrt{5}\varphi_1(x)_1, \sqrt{5}\varphi_1(x)_2, \dots, \sqrt{5}\varphi_1(x)_N, 2\varphi_2(x)_1, 2\varphi_2(x)_2, \dots, 2\varphi_2(x)_K]$$

when  $\varphi_i(x) = [\varphi_i(x)_1, \varphi_i(x)_2, \dots, \varphi_i(x)_{w_i}]$  and  $i \in \{1, 2\}, w_1 = N, w_2 = K$

Notice that it holds that :

$$\begin{aligned} K(x, y) &= \varphi(x)^T \cdot \varphi(y) \\ &= [\sqrt{5}\varphi_1(x)_1, \sqrt{5}\varphi_1(x)_2, \dots, \sqrt{5}\varphi_1(x)_N, 2\varphi_2(x)_1, 2\varphi_2(x)_2, \dots, 2\varphi_2(x)_K] \cdot \\ &\quad [\sqrt{5}\varphi_1(y)_1, \sqrt{5}\varphi_1(y)_2, \dots, \sqrt{5}\varphi_1(y)_N, 2\varphi_2(y)_1, 2\varphi_2(y)_2, \dots, 2\varphi_2(y)_K]^T \\ &= 5S_1 + 4S_2 = 5K_1(x, y) + 4K_2(x, y) \quad \blacksquare \end{aligned}$$

**Solution 1.b:**

Assume that  $\varphi_1: \mathbb{R}^n \rightarrow \mathbb{R}^m, \varphi_2: \mathbb{R}^n \rightarrow \mathbb{R}^K$

Yes, first, since there exist a linear separator for the data under the mapping  $\varphi_1: \mathbb{R}^n \rightarrow \mathbb{R}^m$  with the weight's vector  $w \in \mathbb{R}^{m+1}$ , it holds that :

$$\text{Sgn} \left( w_0 + \varphi_1(x) \cdot \begin{bmatrix} w_1 \\ \vdots \\ w_m \end{bmatrix} \right) \text{ is a linear classifier.}$$

We want to show that there exist  $w' \in \mathbb{R}^{m+K+1}$  s.t :

$$\text{Sgn} \left( w_0' + \varphi(x) \cdot \begin{bmatrix} w_1' \\ \vdots \\ w_{m+K}' \end{bmatrix} \right) = \text{Sgn} \left( w_0 + \varphi_1(x) \cdot \begin{bmatrix} w_1 \\ \vdots \\ w_m \end{bmatrix} \right) \quad \forall x \in D \text{ and then we could}$$

conclude that  $w_0' + \varphi(x) \cdot \begin{bmatrix} w_1' \\ \vdots \\ w_{m+K}' \end{bmatrix}$  is a linear separator for the data under the

mapping  $\varphi: \mathbb{R}^n \rightarrow \mathbb{R}^{m+K}$  with the weight's vector  $w'$ .

Define  $w' \in \mathbb{R}^{m+K+1}$  in the following order :

$$w' = [w'_0, w'_1, \dots, w'_{m+K+1}] = \left[ w_0, \frac{w_1}{\sqrt{5}}, \dots, \frac{w_m}{\sqrt{5}}, 0_{m+1}, 0_{m+2}, \dots, 0_{m+K} \right]$$

Hence :

$$\begin{aligned} & \text{Sgn} \left( w'_0 + \varphi(x) \cdot \begin{bmatrix} w'_1 \\ \vdots \\ w'_{m+K} \end{bmatrix} \right) \\ &= \text{Sgn} \left( w_0 + \left[ \sqrt{5}\varphi_1(x)_1, \sqrt{5}\varphi_1(x)_2, \dots, \sqrt{5}\varphi_1(x)_N, 2\varphi_2(x)_1, 2\varphi_2(x)_2, \dots, 2\varphi_2(x)_K \right] \cdot \right. \\ & \quad \left. \begin{bmatrix} w_0, \frac{w_1}{\sqrt{5}}, \dots, \frac{w_m}{\sqrt{5}}, 0_{m+1}, 0_{m+2}, \dots, 0_{m+K} \end{bmatrix}^T \right) \\ &= \text{Sgn} (w_0 + w_1\varphi_1(x)_1 + w_2\varphi_1(x)_2 + \dots + w_m\varphi_1(x)_m) \\ &= \text{Sgn} \left( w_0 + \varphi_1(x) \cdot \begin{bmatrix} w_1 \\ \vdots \\ w_m \end{bmatrix} \right) \quad \blacksquare \end{aligned}$$

### Solution 1.c:

We define  $\varphi: S \rightarrow \mathbb{R}^N$  as follows:

$$\varphi(x) = \left[ \underbrace{3, 3, \dots, 3}_{|x|}, \underbrace{0, 0, \dots, 0}_{N-|x|} \right], \forall x \in S.$$

i.e,  $\varphi$  is a vector in  $\mathbb{R}^N$  where the first  $|x|$  entries of it are 3 and the last  $N - |x|$  entries are 0.

By defining  $\varphi$  this way we can show that it satisfies the requirements:

Let  $x, y \in S$  for some  $S = \{1, 2, \dots, N\}$  and w.l.o.g. assume  $|x| < |y|$ . it holds that :

$$\begin{aligned} \varphi(x) \cdot \varphi(y) &= \left[ \underbrace{3, 3, \dots, 3}_{|x|}, \underbrace{0, 0, \dots, 0}_{N-|x|} \right] \cdot \left[ \underbrace{3, 3, \dots, 3}_{|y|}, \underbrace{0, 0, \dots, 0}_{N-|y|} \right]^T \\ &= \sum_i^{|x|} 3 \cdot 3 + \sum_i^{|y|-|x|} 0 \cdot 3 + \sum_i^{N-|y|} 0 \cdot 0 = 9 \min(x, y) \end{aligned}$$

Thus, by defining  $f(x, y) = \min(x, y)$  we get :

$K(x, y) = 9f(x, y) = 9 \min(x, y) = \varphi(x) \cdot \varphi(y)$  and so

$K$  is a valid kernel with respect to  $\varphi$  ■

## Solution 2:

We want to find the maximum possible revenue under the constraint that our budget is 20,000\$. We therefore define:

$B(h, s) = 20h + 170s - 20000$  to be our constraint function, now, we can define the

lagrangian :  $L(h, s, \lambda) = R(h, s) + \lambda \cdot B(h, s) = 200h^{\frac{2}{3}}s^{\frac{1}{3}} + \lambda(20h + 170s - 20000)$ .

We saw that by maximizing such lagrangian we maximize  $R(h, s)$  and maintain our constraint, as the gradient of  $L(h, s, \lambda)$  encapsulates both of our needs when it is equal to  $\vec{0}$ .

This, we now turn to solve  $\nabla L(h, s, \lambda) = \vec{0}$ .

$$\begin{aligned}\frac{\partial L(h, s, \lambda)}{\partial h} &= \frac{2}{3} \cdot 200s^{\frac{1}{3}}h^{-\frac{1}{3}} + 20\lambda = 0 \rightarrow \frac{400}{3} \cdot \left(\frac{s}{h}\right)^{\frac{1}{3}} = -20\lambda \rightarrow \left(\frac{s}{h}\right) = -\frac{3\lambda}{20} \\ \xrightarrow{t=\frac{s}{h}} t^{\frac{1}{3}} &= -\frac{3\lambda}{20} \rightarrow t = -\frac{3\lambda}{20} \cdot t^{\frac{2}{3}} \quad [1]\end{aligned}$$

$$\begin{aligned}\frac{\partial L(h, s, \lambda)}{\partial s} &= \frac{1}{3} \cdot 200 \cdot s^{-\frac{2}{3}} \cdot h^{\frac{2}{3}} + 170\lambda = 0 \rightarrow \frac{200}{3} \cdot \left(\frac{h}{s}\right)^{\frac{2}{3}} = -170\lambda \rightarrow \\ \left(\frac{s}{h}\right)^{-\frac{2}{3}} &= -\frac{51\lambda}{20} \xrightarrow{t=\frac{s}{h}} t^{-\frac{2}{3}} = -\frac{51\lambda}{20} \rightarrow 1 = -\frac{51\lambda}{20} \cdot t^{\frac{2}{3}} \rightarrow \frac{1}{17} = -\frac{3\lambda}{20} \cdot t^{\frac{2}{3}} \quad [2]\end{aligned}$$

From [1], [2] we get :  $t = \frac{s}{h} = \frac{1}{17} \rightarrow h = 17s$ , we also have :

$$\frac{\partial L(h, s, \lambda)}{\partial \lambda} = 20h + 170s - 20000 = 0 \text{ so for } h = 17s \text{ we get}$$

$$340s + 170s = 20000 \rightarrow 510s = 20000 \rightarrow s = \frac{2000}{51} \text{ and so : } h = \frac{2000}{3}$$

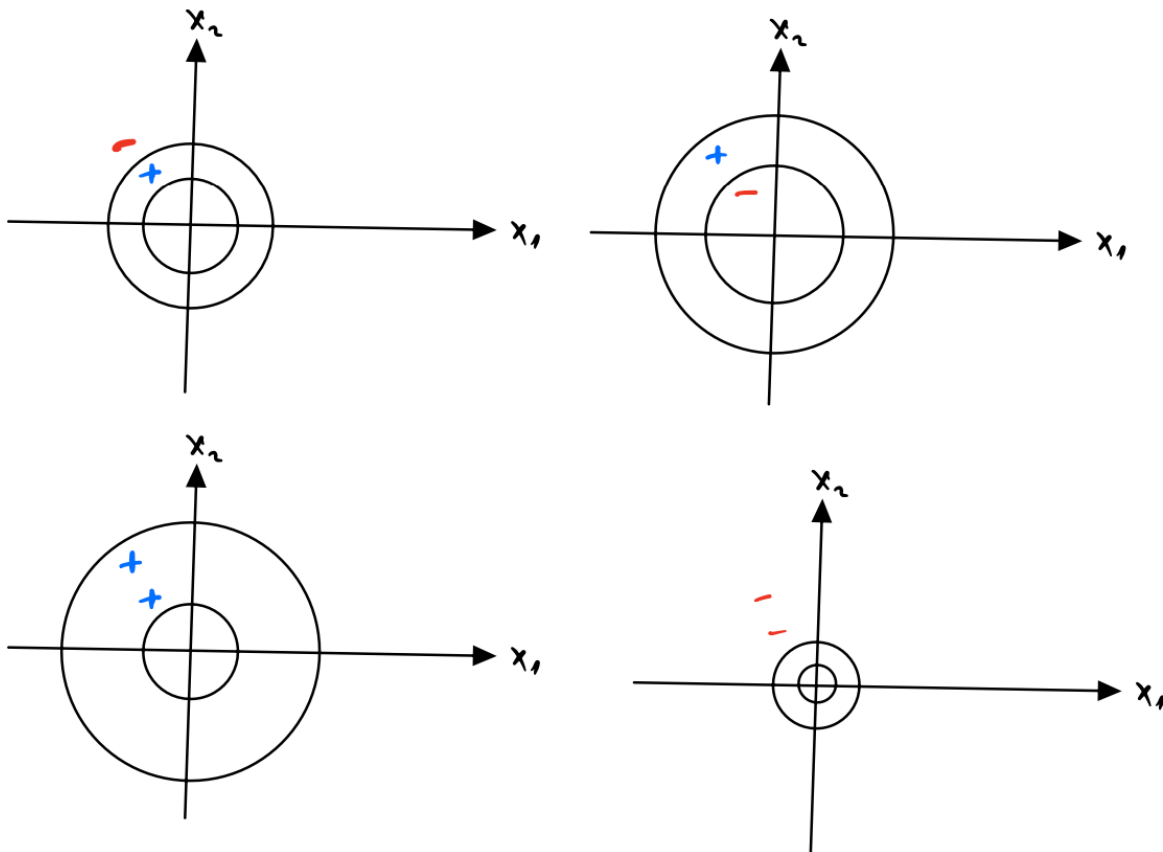
Therefore, for  $(h, s) = \left(\frac{2000}{3}, \frac{2000}{51}\right)$

we get that our revenue is maximal. so, our maximal revenue given that our budget is 20000\$ is :

$$R\left(\frac{2000}{3}, \frac{2000}{51}\right) = 200 \cdot \left(\frac{2000}{3}\right)^{\frac{2}{3}} \cdot \left(\frac{2000}{51}\right)^{\frac{1}{3}} = 51854.81583 \quad \blacksquare$$

### Solution 3.a:

First, we show that  $VC(H) \geq 2$ :



So indeed,  $VC(H) \geq 2$ .

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Now, we proceed to show that  $VC(H) < 3$ .

we note that there are 2 cases of convex hull

for 3 points in  $\mathbb{R}^2$  which are : the 3 points are colinear, or the 3 points form a triangle

for both cases, denote  $r_{x_1}, r_{x_2}, r_{x_3}$  to be the corresponding distances of points  $x_1, x_2, x_3$

from the origin and w.l.o.g. assume that  $r_{x_1} \leq r_{x_2} \leq r_{x_3}$ .

So, for the labeling :  $x_1 = +, x_2 = -, x_3 = +$  . it holds that

$\forall h \in H \exists i \in \{1,2,3\} (h(x_i) \neq y_i)$ , Since our hypothesis space is all origin - centered rings it's impossible to exclude  $x_2$  from within the ring without misclassifying one of the other instances, So,  $VC(H) < 3$ .

Since  $VC(H) \geq 2$  and  $VC(H) < 3$  we deduce that  $VC(H) = 2$ . ■

### Solution 3.b:

We show that  $C$  is PAC – learnable by  $L$  using  $H$ . in order to show that we will prove

that  $\forall 0 < \epsilon < \frac{1}{2}, \forall 0 < \delta < \frac{1}{2}$ , and for all  $c \in C$  and distributions  $\pi$  over  $X$ , the following holds: with data drawn independently according to  $\pi$ ,  $L$  will output, with probability at least  $(1 - \delta)$ , a hypothesis  $h \in H$  such that  $\text{error}_{\pi}(h) \leq \epsilon$  as well as prove that  $L$  operates in time (and sample) complexity that is polynomial in  $\frac{1}{\epsilon}$  and  $\frac{1}{\delta}$ .

Let  $c \in C$ ,  $\epsilon, \delta \in \left(0, \frac{1}{2}\right)$ .

Denote  $r_1^*, r_2^*$  s.t.  $r_1^* \leq r_2^*$  to be the radii of the target origin center annulus.

Assume that our learner is consistent. which means that:

Given the training examples

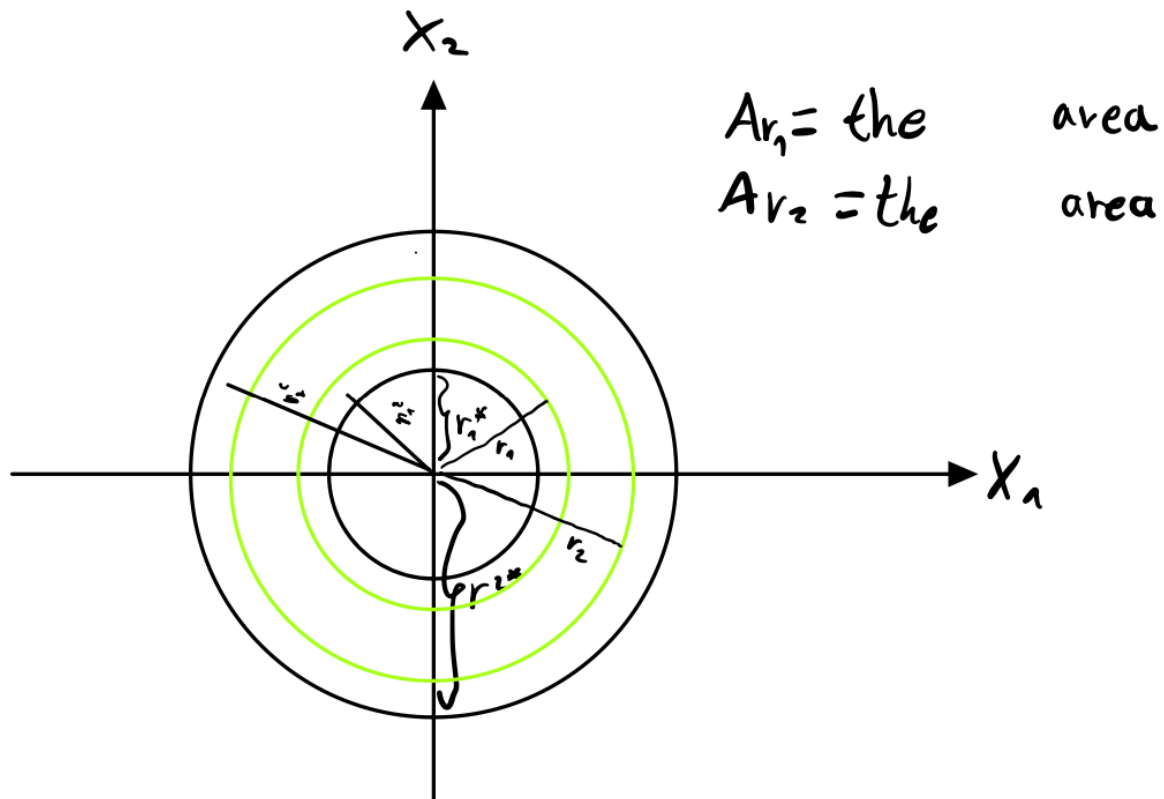
Find positive points with the maximum distance from the origin and minimum distance from the origin,

Define:

$r_1$  = the distance of the positive point with the minimum distance from the origin,

$r_2$  = the distance of the positive point with the maximum distance from the origin.

Output  $h(r_1, r_2)$ .



Consider training data,  $D \in \mathbb{R}^2$ .

Case 1 :

Assume that  $D$  visits each one of the 2 sets  $A_i$  defined above. (recall that  $\pi(A_i) = \frac{\epsilon}{2}$ )  
therefore our error  $r_\pi(x) \leq \epsilon$ .

Otherwise :

Every point in  $D$  never visited at least one of  $A_1, A_2$ .

therefore  $P(\{D \in X^m : L(D), c > \epsilon\}) \leq \sum_{i=1}^2 (P(X - A_i))^m \leq 2 \left(1 - \frac{\epsilon}{2}\right)^m \leq 2e^{\left(-\frac{m\epsilon}{2}\right)}$

Hence :

$$\left(1 - \frac{\epsilon}{2}\right)^m \leq 2e^{\left(-\frac{m\epsilon}{2}\right)} \leq \delta \rightarrow e^{-\frac{m\epsilon}{2} * \ln(e)} \leq e^{\ln\left(\frac{\delta}{2}\right)} \rightarrow -\ln\left(\frac{\delta}{2}\right) \leq \frac{m\epsilon}{2} * \ln(e) \rightarrow \ln\left(\frac{2}{\delta}\right) \leq \frac{m\epsilon}{2}$$

$$\frac{2\ln\left(\frac{2}{\delta}\right)}{\epsilon} \leq m$$

### **Solution 3c:**

In order to get with 95% confidence a hypothesis with at most 5% error we set  $\epsilon = 0.05, \delta = 0.05$ .

Calculating the complexity with the bounds we found in the previous questions:

Sample complexity with the bound we found in 3. b :  $m \geq 147.5$

Sample complexity with the bound for infinite  $H$  :  $m \geq 2992.9$

We found that with the bound of 3. b we have tighter bound as we learned in the lecture that the bound from a direct calculation encapsulates more information about the data since it assumes it comes from a certain distribution whereas the bound for infinite  $H$  does not assume anything about the data we are working on, and it only takes into account the VC dimensions of our  $H$ .

#### Solution 4.a:

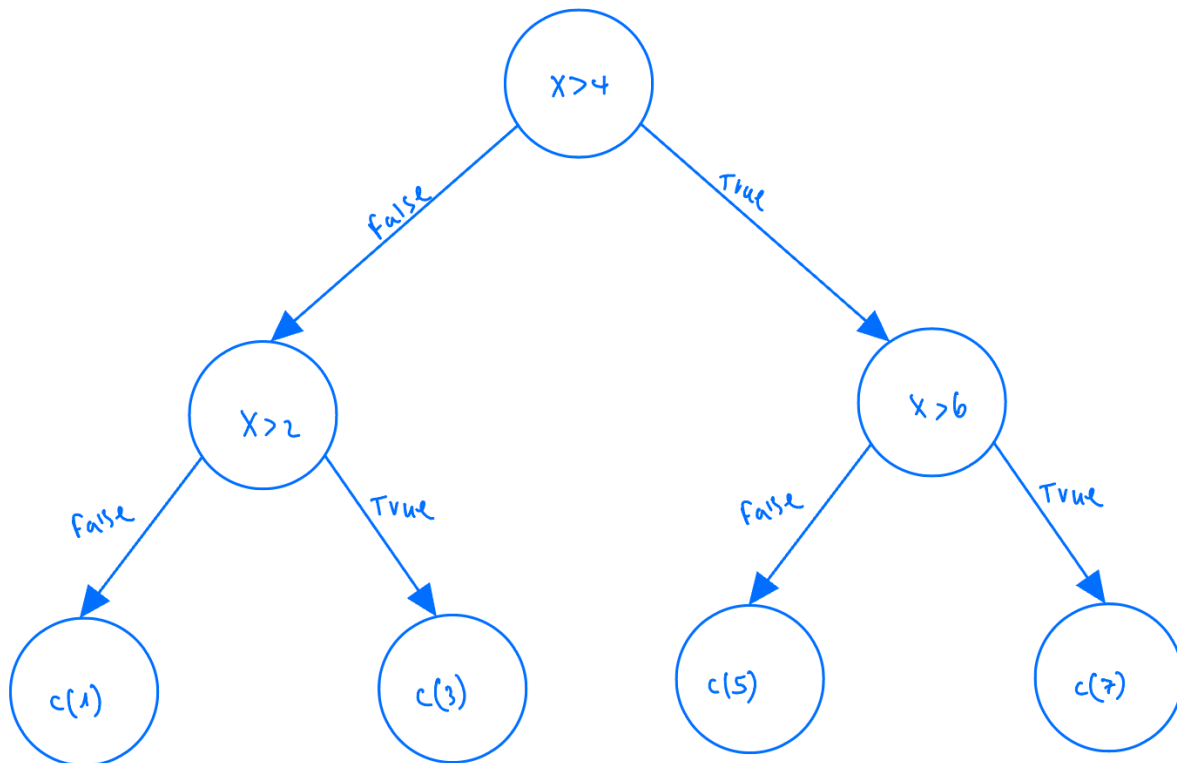
→  $VC(H_3) \geq 4$ :

We choose the following set of points  $\{1, 3, 5, 7\}$ .

There will be a dichotomy for the set of points.

Denote the partition for the dichotomy as  $c(i), \forall i \in \{1, 3, 5, 7\}$

We will show that a binary decision tree with 7 nodes (3 internal nodes and 4 leaf nodes) can indeed shatter this set.



$\forall c \in C$  (concept) we can always choose  $h$  s.t  $h(x_i) = c(x_i), \forall x_i \in \{1, 3, 5, 7\}$ .

←  $VC(H_3) < 5$ :

Let's denote the set of 5 points as  $X = \{x_1, x_2, x_3, x_4, x_5\}$  where  $x_1 < x_2 < x_3 < x_4 < x_5$ .

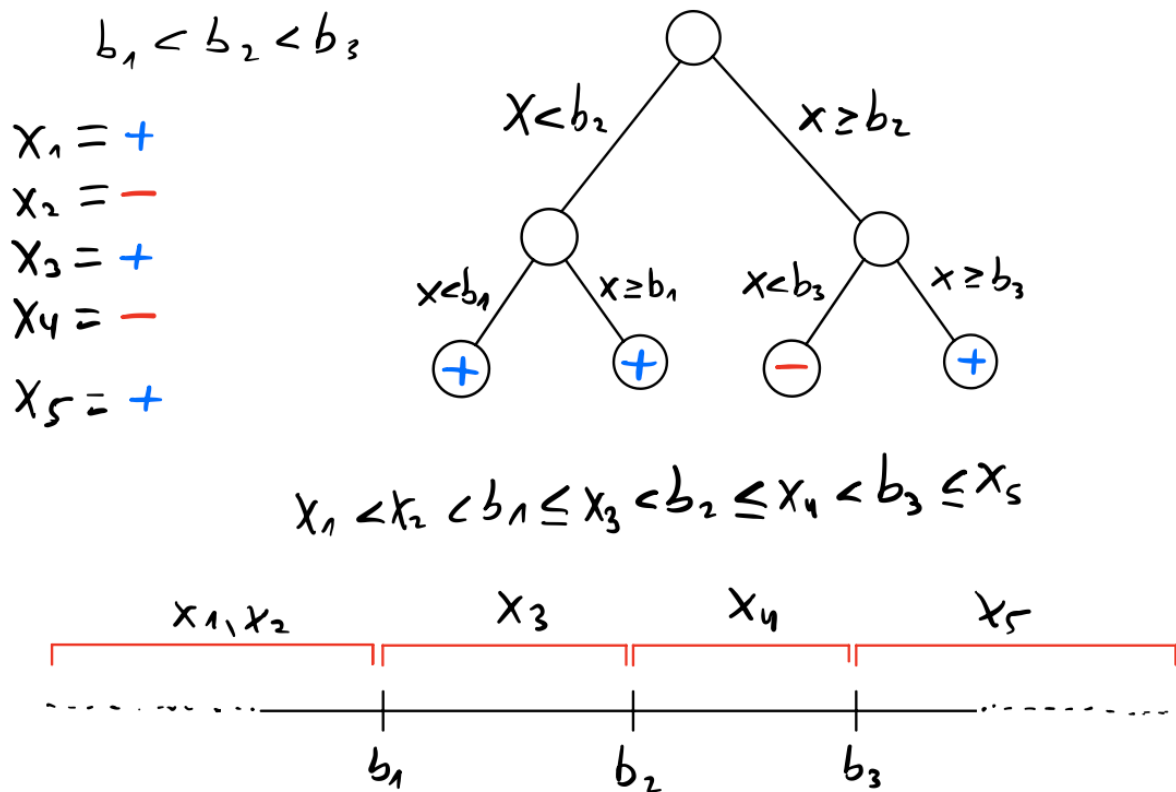
We want to show that there exists a labeling of  $X$  that cannot be realized by a binary decision tree with 7 nodes (3 internal nodes and 4 leaf nodes).

Consider the labeling  $L = \{+, -, +, -, +\}$ .

A binary decision tree with 7 nodes partitions the input space into 4 regions, each

corresponding to a leaf node. Each leaf node can output either + or -  
 Let's denote the decision boundaries of the tree as  $b_1, b_2$  and  $b_3$ , where  $b_1 < b_2 < b_3$ .  
 These boundaries partition the input space into the following 4 regions:

1.  $R_1 = (-\infty, b_1)$
2.  $R_2 = [b_1, b_2)$
3.  $R_3 = [b_2, b_3)$
4.  $R_4 = [b_3, \infty)$



Since we have 5 points and 4 regions, by the Pigeonhole Principle, there must be at least one region that contains at least two points. w.l.o.g., let's say  $x_1, x_2$  fall into the same region, say  $R_1$ .

In the labeling  $L$ ,  $x_1$  and  $x_2$  have different labels. However, all points in the same region get assigned the same label by the decision tree. Therefore, there is no way for the decision tree to realize the labeling  $L$ .



*This shows that a binary decision tree with 7 nodes cannot shatter a set of 5 points, and hence the VC dimension of the hypothesis space of all  $x$ -node decision tree with  $m \leq 3$  is less than 5.*

*Since we previously showed that this hypothesis space can shatter a set of 4 points, we conclude that the VC dimension of the hypothesis space of all " $x$  – node decision tree" with  $m \leq 3$  is exactly 4*

**Solution 4.b:**

$\rightarrow VC(H_m) \geq 2^{m-1}$ :

*We choose the following set of points  $\{1, 2, \dots, 2^{m-1}\}$ .*

*Let be  $S$  be a dichotomy for the set of points.*

*Denote the partition for  $S$  as  $c(i), \forall i \in \{1, 2, \dots, 2^{m-1}\}$*

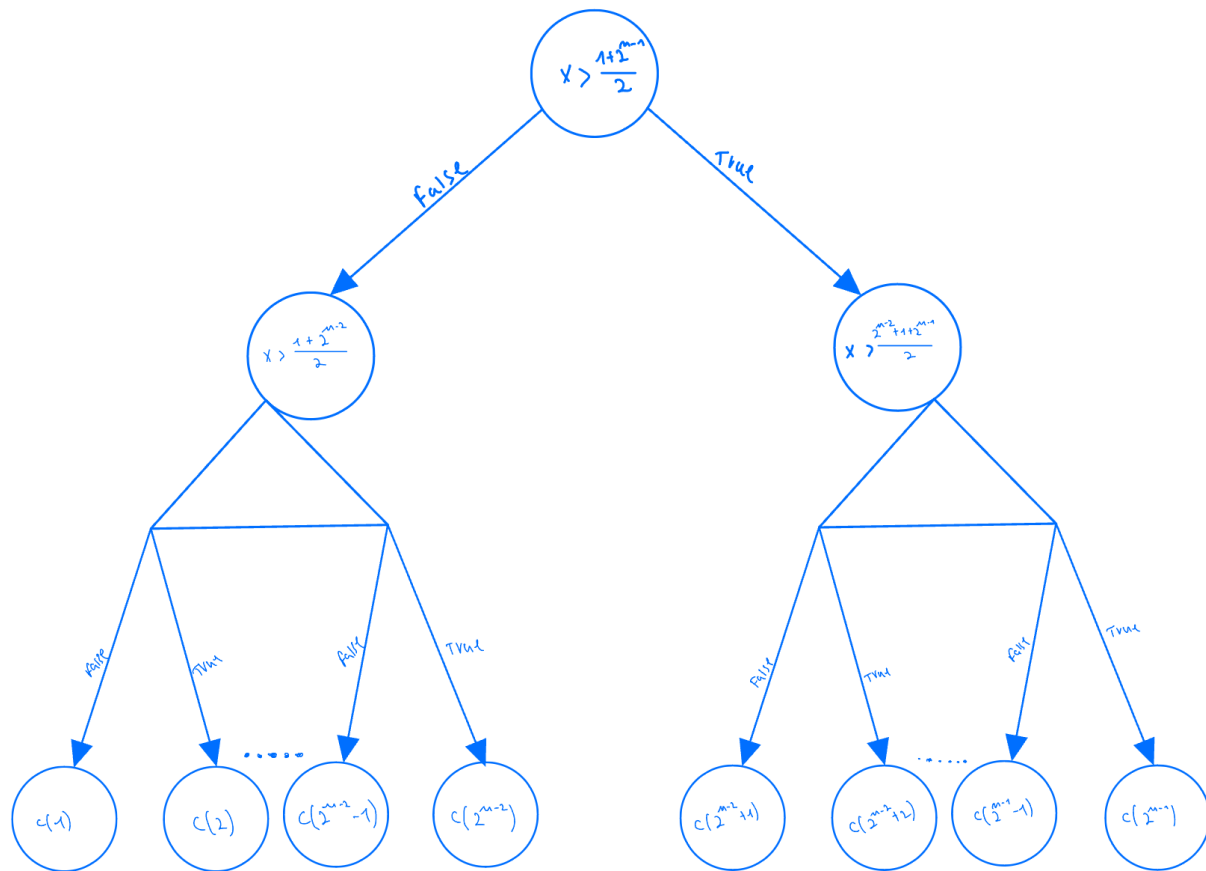
*We will show that a binary decision tree with  $2^n - 1$  nodes can indeed shatter this set.*

*We define the decision boundary for each  $v_i \in V$  (vertices of the decision tree) as follows : Denote  $S_i = \{n_1, n_2, \dots, n_k\}$  and  $B_i$  to be the decision boundary question for each  $v_i$ .*

$\forall v_i \in V, B_i$  defined as  $x > \frac{\min(S_i) + \max(S_i)}{2}$ , if we went to the left child set  $S_i$

*corresponding to  $v_i$  for the left portion of the set  $S_{i-1}$  (old  $S_i$ )*

*and if we went to the right child set  $S_{i+1}$  corresponding to  $v_{i+1}$  for the right portion of the set  $S_{i-1}$ . Therefore we get the following diagram:*



$\forall c \in C$  (concept) we can always choose  $h$  s.t  $h(x_i) = c(x_i), \forall x_i \in \{1, 2, \dots, 2^{m-1}\}$ .

$\leftarrow VC(H_m) < 2^{m-1} + 1$ :

Let's denote the set of  $2^{m-1} + 1$  points as  $X = \{x_1, x_2, \dots, x_{2^{m-1}+1}\}$  where  $x_1 < x_2 < \dots < x_{2^{m-1}+1}$ .

We want to show that there exists a labeling of  $X$  that cannot be realized by a binary decision tree with  $2^m - 1$  nodes.

Consider the labeling  $L = \{+, -, +, -, \dots, (-1)^{2^{m-1}+1}\}$ .

A binary decision tree with  $2^m - 1$  nodes partitions the input space into  $2^{m-1}$  regions, each corresponding to a leaf node. Each leaf node can output either  $+$  or  $-$

Let's denote the decision boundaries of the tree as  $b_1, b_2, \dots, b_{2^{m-1}-1}$ , where  $b_1 < b_2 < \dots < b_{2^{m-1}-1}$ .

These boundaries partition the input space into the following  $2^{m-1}$  regions:

$$\begin{aligned}
R_1 &= (-\infty, b_1) \\
R_2 &= [b_1, b_2) \\
&\vdots \\
R_{2^{m-1}} &= [b_{2^{m-1}-1}, \infty)
\end{aligned}$$

Since we have  $2^{m-1} + 1$  points and  $2^{m-1}$  regions, by the Pigeonhole Principle, there must be at least one region that contains at least two points. w.l.o.g. let's say  $x_1, x_2$  fall into the same region, say  $R_1$ .

In the labeling  $L$ ,  $x_1$  and  $x_2$  have different labels. However, all points in the same region get assigned the same label by the decision tree. Therefore, there is no way for the decision tree to realize the labeling  $L$ .

This shows that a binary decision tree with  $2^m - 1$  nodes cannot shatter a set of  $2^{m-1} + 1$  points, and hence the VC dimension of the hypothesis space of all  $x$ -node decision tree with  $m$  is less than  $2^{m-1} + 1$ .

Since we previously showed that this hypothesis space can shatter a set of  $2^{m-1}$  points, we conclude that the VC dimension of the hypothesis space of all " $x$  - node decision tree" with  $m$  is exactly  $2^{m-1}$ .