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4

5

1. Teoría de números

1.1. Funciones básicas

1.1.1. Función piso y techo

```
lli piso(lli a, lli b){
  if((a >= 0 \&\& b > 0) || (a < 0 \&\& b < 0)){}
    return a / b:
  }else{
    if(a \% b == 0) return a / b:
    else return a / b - 1;
 }
}
lli techo(lli a, lli b){
  if((a >= 0 \&\& b > 0) || (a < 0 \&\& b < 0)){}
    if(a \% b == 0) return a / b;
    else return a / b + 1;
  }else{
    return a / b;
  }
}
```

1.1.2. Exponenciación y multiplicación binaria

```
lli power(lli b, lli e){
    lli ans = 1;
    while(e){
        if(e & 1) ans *= b;
        e >>= 1;
        b *= b;
    }
    return ans;
}

lli multMod(lli a, lli b, lli n){
    lli ans = 0;
    a %= n, b %= n;
    if(abs(b) > abs(a)) swap(a, b);
```

```
if(b < 0){
    a *= -1, b *= -1;
}
while(b){
    if(b & 1) ans = (ans + a) % n;
    b >>= 1;
    a = (a + a) % n;
}
return ans;
}
```

1.1.3. Mínimo común múltiplo y máximo común divisor

```
lli gcd(lli a, lli b){
  lli r:
  while(b != 0) r = a \% b, a = b, b = r;
  return a:
lli lcm(lli a, lli b){
  return b * (a / gcd(a, b));
lli gcd(vector<lli>> & nums){
  lli ans = 0;
  for(lli & num : nums) ans = gcd(ans, num);
  return ans;
}
lli lcm(vector<lli> & nums){
  lli ans = 1:
  for(lli & num : nums) ans = lcm(ans, num);
  return ans:
}
```

1.1.4. Euclides extendido e inverso modular

```
while(r1){
    q = r0 / r1;
    ri = r0 \% r1, r0 = r1, r1 = ri;
    si = s0 - s1 * q, s0 = s1, s1 = si;
    ti = t0 - t1 * q, t0 = t1, t1 = ti;
  s = s0, t = t0;
  return r0;
}
lli modularInverse(lli a, lli m){
  lli r0 = a, r1 = m, ri, s0 = 1, s1 = 0, si;
  while(r1){
    si = s0 - s1 * (r0 / r1), s0 = s1, s1 = si;
   ri = r0 \% r1, r0 = r1, r1 = ri;
  }
  if(r0 < 0) s0 *= -1;
  if(s0 < 0) s0 += m;
  return s0;
}
```

1.1.5. Todos los inversos módulo p

```
//find all inverses (from 1 to p-1) modulo p
vector<lli> allInverses(lli p){
  vector<lli> ans(p);
  ans[1] = 1;
  for(lli i = 2; i < p; ++i)
    ans[i] = p - (p / i) * ans[p % i] % p;
  return ans;
}</pre>
```

1.1.6. Exponenciación binaria modular

```
lli powerMod(lli b, lli e, lli m){
  lli ans = 1;
  b %= m;
  if(e < 0){
    b = modularInverse(b, m);</pre>
```

```
e *= -1;
}
while(e){
  if(e & 1) ans = (ans * b) % m;
  e >>= 1;
  b = (b * b) % m;
}
return ans;
}
```

1.1.7. Teorema chino del residuo

1.1.8. Teorema chino del residuo generalizado

}

7

```
}
while(a0 >= m0) a0 -= m0; while(a0 < 0) a0 += m0;
return {a0, m0};
}</pre>
```

1.1.9. Coeficiente binomial

```
lli ncr(lli n, lli r){
  if(r < 0 || r > n) return 0;
  r = min(r, n - r);
  lli ans = 1;
  for(lli den = 1, num = n; den <= r; den++, num--)
    ans = ans * num / den;
  return ans;
}</pre>
```

1.1.10. Fibonacci

```
//very fast fibonacci
inline void modula(lli & n){
  while(n \ge mod) n -= mod;
}
lli fibo(lli n){
  array<lli, 2> F = {1, 0};
  lli p = 1;
  for(lli v = n; v >>= 1; p <<= 1);
  array<lli, 4> C;
  do{
   int d = (n \& p) != 0;
    C[0] = C[3] = 0;
    C[d] = F[0] * F[0] % mod;
    C[d+1] = (F[0] * F[1] << 1) \% mod;
    C[d+2] = F[1] * F[1] % mod;
    F[0] = C[0] + C[2] + C[3];
    F[1] = C[1] + C[2] + (C[3] << 1);
    modula(F[0]), modula(F[1]);
  }while(p >>= 1);
  return F[1];
```

```
1.2. Cribas
```

1.2.1. Criba de divisores

```
vector<lli> divisorsSum;
vector<vector<int>> divisors;
void divisorsSieve(int n){
   divisorsSum.resize(n + 1, 0);
   divisors.resize(n + 1);
   for(int i = 1; i <= n; ++i){
      for(int j = i; j <= n; j += i){
        divisorsSum[j] += i;
        divisors[j].push_back(i);
      }
   }
}</pre>
```

1.2.2. Criba de primos

```
vector<int> primes;
vector<bool> isPrime;
void primesSieve(int n){
  isPrime.resize(n + 1, true);
  isPrime[0] = isPrime[1] = false;
 primes.push_back(2);
 for(int i = 4; i <= n; i += 2) isPrime[i] = false;</pre>
  int limit = sqrt(n);
 for(int i = 3; i \le n; i += 2){
    if(isPrime[i]){
      primes.push_back(i);
      if(i <= limit)</pre>
        for(int j = i * i; j \le n; j += 2 * i)
          isPrime[j] = false;
   }
 }
}
```

8

1.2.3. Criba de factor primo más pequeño

```
vector<int> lowestPrime;
void lowestPrimeSieve(int n){
  lowestPrime.resize(n + 1, 1);
  lowestPrime[0] = lowestPrime[1] = 0;
  for(int i = 2; i <= n; ++i) lowestPrime[i] = (i & 1 ? i : 2);
  int limit = sqrt(n);
  for(int i = 3; i <= limit; i += 2)
    if(lowestPrime[i] == i)
      for(int j = i * i; j <= n; j += 2 * i)
        if(lowestPrime[j] == j) lowestPrime[j] = i;
}</pre>
```

1.2.4. Criba de factor primo más grande

```
vector<int> greatestPrime;
void greatestPrimeSieve(int n){
  greatestPrime.resize(n + 1, 1);
  greatestPrime[0] = greatestPrime[1] = 0;
  for(int i = 2; i <= n; ++i) greatestPrime[i] = i;
  for(int i = 2; i <= n; i++)
    if(greatestPrime[i] == i)
      for(int j = i; j <= n; j += i)
          greatestPrime[j] = i;
}</pre>
```

1.2.5. Criba de factores primos

```
vector<vector<int>>> primeFactors;
void primeFactorsSieve(lli n){
  primeFactors.resize(n + 1);
  for(int i = 0; i < primes.size(); ++i){
    int p = primes[i];
    for(int j = p; j <= n; j += p)
        primeFactors[j].push_back(p);
  }
}</pre>
```

1.2.6. Criba de la función φ de Euler

```
vector<int> Phi;
void phiSieve(int n){
   Phi.resize(n + 1);
   for(int i = 1; i <= n; ++i) Phi[i] = i;
   for(int i = 2; i <= n; ++i)
      if(Phi[i] == i)
      for(int j = i; j <= n; j += i)
            Phi[j] -= Phi[j] / i;
}</pre>
```

1.2.7. Criba de la función μ

```
vector<int> Mu;
void muSieve(int n){
   Mu.resize(n + 1, -1);
   Mu[0] = 0, Mu[1] = 1;
   for(int i = 2; i <= n; ++i)
      if(Mu[i])
      for(int j = 2*i; j <= n; j += i)
            Mu[j] -= Mu[i];
}</pre>
```

1.2.8. Triángulo de Pascal

9

1.2.9. Segmented sieve

```
vector<int> segmented_sieve(int limit){
  const int L1D_CACHE_SIZE = 32768;
  int raiz = sqrt(limit);
  int segment_size = max(raiz, L1D_CACHE_SIZE);
  int s = 3, n = 3;
  vector<int> primes(1, 2), tmp, next;
  vector<char> sieve(segment_size);
  vector<bool> is_prime(raiz + 1, 1);
  for(int i = 2; i * i <= raiz; i++)
    if(is_prime[i])
      for(int j = i * i; j <= raiz; j += i)
        is_prime[j] = 0;
  for(int low = 0; low <= limit; low += segment_size){</pre>
    fill(sieve.begin(), sieve.end(), 1);
    int high = min(low + segment_size - 1, limit);
    for(; s * s \le high; s += 2){
      if(is_prime[s]){
        tmp.push_back(s);
        next.push_back(s * s - low);
      }
    }
    for(size_t i = 0; i < tmp.size(); i++){</pre>
      int j = next[i];
      for(int k = tmp[i] * 2; j < segment_size; j += k)</pre>
        sieve[j] = 0;
      next[i] = j - segment_size;
    for(; n <= high; n += 2)
      if(sieve[n - low])
        primes.push_back(n);
  }
  return primes;
}
```

1.2.10. Criba de primos lineal

```
vector<int> linearPrimeSieve(int n){
  vector<int> primes;
```

```
vector<bool> isPrime(n+1, true);
for(int i = 2; i <= n; ++i){
   if(isPrime[i])
     primes.push_back(i);
   for(int p : primes){
     int d = i * p;
     if(d > n) break;
     isPrime[d] = false;
     if(i % p == 0) break;
   }
}
return primes;
}
```

1.2.11. Criba lineal para funciones multiplicativas

```
//suppose f(n) is a multiplicative function and
//we want to find f(1), f(2), ..., f(n)
//we have f(pq) = f(p)f(q) if qcd(p, q) = 1
//and\ f(p^a) = q(p, a), where p is prime and a>0
vector<int> generalSieve(int n, function<int(int, int)> g){
  vector<int> f(n+1, 1), cnt(n+1), acum(n+1), primes;
  vector<bool> isPrime(n+1, true);
 for(int i = 2; i \le n; ++i){
    if(isPrime[i]){ //case base: f(p)
     primes.push_back(i);
     f[i] = g(i, 1);
     cnt[i] = 1;
     acum[i] = i;
   for(int p : primes){
     int d = i * p;
     if(d > n) break;
     isPrime[d] = false;
     if(i % p == 0){ //gcd(i, p) != 1
       f[d] = f[i / acum[i]] * g(p, cnt[i] + 1);
        cnt[d] = cnt[i] + 1;
        acum[d] = acum[i] * p;
       break;
     else{ //qcd(i, p) = 1}
```

```
f[d] = f[i] * g(p, 1);
    cnt[d] = 1;
    acum[d] = p;
}
}
return f;
}
```

1.3. Factorización

1.3.1. Factorización de un número

```
vector<pair<lli, int>> factorize(lli n){
  vector<pair<lli, int>> f;
  for(lli p : primes){
    if(p * p > n) break;
    int pot = 0;
    while(n % p == 0){
      pot++;
      n /= p;
    }
    if(pot) f.emplace_back(p, pot);
}
if(n > 1) f.emplace_back(n, 1);
  return f;
}
```

1.3.2. Potencia de un primo que divide a un factorial

```
lli potInFactorial(lli n, lli p){
   lli ans = 0, div = n;
   while(div /= p) ans += div;
   return ans;
}
```

1.3.3. Factorización de un factorial

```
vector<pair<lli, lli>> factorizeFactorial(lli n){
  vector<pair<lli, lli>> f;
  for(lli p : primes){
    if(p > n) break;
    f.emplace_back(p, potInFactorial(n, p));
  }
  return f;
}
```

1.3.4. Factorización usando Pollard-Rho

```
bool isPrimeMillerRabin(lli n){
  if(n < 2) return false;
  if(n == 2) return true:
 lli d = n - 1, s = 0;
 for(; !(d & 1); d >>= 1, ++s);
 for(int i = 0; i < 16; ++i){
    lli a = 1 + rand() % (n - 1);
    lli m = powerMod(a, d, n);
    if (m == 1 \mid \mid m == n - 1) goto exit;
    for(int k = 0; k < s; ++k){
      m = m * m \% n;
      if(m == n - 1) goto exit;
    return false;
    exit:;
  }
  return true;
}
lli getFactor(lli n){
 lli a = 1 + rand() \% (n - 1);
 lli b = 1 + rand() \% (n - 1);
 lli x = 2, y = 2, d = 1;
  while(d == 1){
    x = x * (x + b) \% n + a;
    y = y * (y + b) % n + a;
    y = y * (y + b) \% n + a;
```

```
d = gcd(abs(x - y), n);
  }
  return d;
}
map<lli, int> fact;
void factorizePollardRho(lli n, bool clean = true){
  if(clean) fact.clear();
  while(n > 1 && !isPrimeMillerRabin(n)){
   11i f = n;
   for(; f == n; f = getFactor(n));
   n /= f;
    factorizePollardRho(f, false);
    for(auto & it : fact){
      while(n % it.first == 0){
        n /= it.first;
        ++it.second;
      }
    }
  }
  if(n > 1) ++fact[n];
}
```

1.4. Funciones aritméticas famosas

1.4.1. Función σ

```
//divisor power sum of n
//if pot=0 we get the number of divisors
//if pot=1 we get the sum of divisors
lli sigma(lli n, lli pot){
    lli ans = 1;
    auto f = factorize(n);
    for(auto & factor : f){
        lli p = factor.first;
        int a = factor.second;
        if(pot){
            lli p_pot = power(p, pot);
            ans *= (power(p_pot, a + 1) - 1) / (p_pot - 1);
        }else{
```

```
ans *= a + 1;
}
return ans;
}
```

1.4.2. Función Ω

```
//number of total primes with multiplicity dividing n
int Omega(lli n){
  int ans = 0;
  auto f = factorize(n);
  for(auto & factor : f)
    ans += factor.second;
  return ans;
}
```

1.4.3. Función ω

```
//number of distinct primes dividing n
int omega(lli n){
  int ans = 0;
  auto f = factorize(n);
  for(auto & factor : f)
    ++ans;
  return ans;
}
```

1.4.4. Función φ de Euler

```
//number of coprimes with n less than n
lli phi(lli n){
    lli ans = n;
    auto f = factorize(n);
    for(auto & factor : f)
        ans -= ans / factor.first;
    return ans;
}
```

1.4.5. Función μ

```
//1 if n is square-free with an even number of prime factors
//-1 if n is square-free with an odd number of prime factors
//0 is n has a square prime factor
int mu(lli n){
  int ans = 1;
  auto f = factorize(n);
  for(auto & factor : f){
    if(factor.second > 1) return 0;
    ans *= -1;
  }
  return ans;
}
```

1.5. Orden multiplicativo, raíces primitivas y raíces de la unidad

1.5.1. Función λ de Carmichael

```
//the smallest positive integer k such that for
//every coprime x with n, x^k=1 mod n

lli carmichaelLambda(lli n){
    lli ans = 1;
    auto f = factorize(n);
    for(auto & factor : f){
        lli p = factor.first;
        int a = factor.second;
        lli tmp = power(p, a);
        tmp -= tmp / p;
        if(a <= 2 || p >= 3) ans = lcm(ans, tmp);
        else ans = lcm(ans, tmp >> 1);
    }
    return ans;
}
```

1.5.2. Orden multiplicativo módulo m

```
// the smallest positive integer k such that x^k = 1 mod m
lli multiplicativeOrder(lli x, lli m){
  if(gcd(x, m) != 1) return 0;
  lli order = phi(m);
  auto f = factorize(order);
  for(auto & factor : f){
    lli p = factor.first;
    int a = factor.second;
    order /= power(p, a);
    lli tmp = powerMod(x, order, m);
    while(tmp != 1){
        tmp = powerMod(tmp, p, m);
        order *= p;
    }
  }
  return order;
}
```

1.5.3. Número de raíces primitivas (generadores) módulo m

```
//number of generators modulo m

lli numberOfGenerators(lli m){
   lli phi_m = phi(m);
   lli lambda_m = carmichaelLambda(m);
   if(phi_m == lambda_m) return phi(phi_m);
   else return 0;
}
```

1.5.4. Test individual de raíz primitiva módulo m

```
//test if order(x, m) = phi(m), i.e., x is a generator for Z/mZ
bool testPrimitiveRoot(lli x, lli m){
  if(gcd(x, m) != 1) return false;
  lli order = phi(m);
  auto f = factorize(order);
  for(auto & factor : f){
    lli p = factor.first;
```

```
if(powerMod(x, order / p, m) == 1) return false;
}
return true;
}
```

1.5.5. Test individual de raíz k-ésima de la unidad módulo

m

1.5.6. Encontrar la primera raíz primitiva módulo m

```
lli findFirstGenerator(lli m){
  lli order = phi(m);
  if(order != carmichaelLambda(m)) return -1; //just an
  → optimization, not required
  auto f = factorize(order):
  for(lli x = 1; x < m; x++){
    if(gcd(x, m) != 1) continue;
    bool test = true:
    for(auto & factor : f){
      lli p = factor.first;
      if(powerMod(x, order / p, m) == 1){
       test = false;
        break;
      }
    }
    if(test) return x;
```

```
return -1; //not found }
```

1.5.7. Encontrar la primera raíz k-ésima de la unidad módulo m

```
lli findFirstPrimitiveKthRootUnity(lli k, lli m){
  if(carmichaelLambda(m) % k != 0) return -1; //just an
  → optimization, not required
  auto f = factorize(k);
  for(lli x = 1; x < m; x++){
    if(powerMod(x, k, m) != 1) continue;
    bool test = true:
    for(auto & factor : f){
     lli p = factor.first;
     if(powerMod(x, k / p, m) == 1){
       test = false:
       break;
     }
    }
    if(test) return x;
  return -1; //not found
```

1.5.8. Logaritmo discreto

```
// a^x = b mod m, a and m coprime
pair<lli, lli> discreteLogarithm(lli a, lli b, lli m){
  if(gcd(a, m) != 1) return make_pair(-1, 0); //not found
  lli order = multiplicativeOrder(a, m);
  lli n = sqrt(order) + 1;
  lli a_n = powerMod(a, n, m);
  lli ans = 0;
  unordered_map<lli, lli> firstHalf;
  lli current = a_n;
  for(lli p = 1; p <= n; p++){
    firstHalf[current] = p;
    current = (current * a_n) % m;</pre>
```

```
}
                                                                     if(a < 0) a += p;
  current = b % m;
                                                                     if(a == 0) return 0;
  for(lli q = 0; q \le n; q++){
                                                                     assert(powerMod(a, (p - 1) / 2, p) == 1);
    if(firstHalf.count(current)){
                                                                     if (p \% 4 == 3) return powerMod(a, (p + 1) / 4, p);
      lli p = firstHalf[current];
                                                                    lli s = p - 1;
      lli x = n * p - q;
                                                                     int r = 0;
      return make_pair(x % order, order);
                                                                     while((s & 1) == 0) ++r, s >>= 1;
    }
                                                                    11i n = 2;
    current = (current * a) % m;
                                                                     while(powerMod(n, (p - 1) / 2, p) != p - 1) ++n;
                                                                    lli x = powerMod(a, (s + 1) / 2, p);
  return make_pair(-1, 0); //not found
                                                                    lli b = powerMod(a, s, p);
}
                                                                    lli g = powerMod(n, s, p);
                                                                     while(true){
                                                                      lli t = b;
1.5.9. Raíz k-ésima discreta
                                                                      int m = 0;
                                                                      for(; m < r; ++m){
// x^k = b \mod m, m has at least one generator
                                                                        if(t == 1) break;
vector<lli> discreteRoot(lli k, lli b, lli m){
                                                                        t = t * t % p;
  if(b \% m == 0) return {0};
                                                                      }
  lli g = findFirstGenerator(m);
                                                                      if(m == 0) return x;
  lli power = powerMod(g, k, m);
                                                                      lli gs = powerMod(g, 1 \ll (r - m - 1), p);
  auto y0 = discreteLogarithm(power, b, m);
                                                                      g = gs * gs % p;
  if(y0.first == -1) return {};
                                                                      x = x * gs \% p;
  lli phi_m = phi(m);
                                                                      b = b * g \% p;
  lli d = gcd(k, phi_m);
                                                                      r = m;
  vector<lli> x(d);
                                                                    }
  x[0] = powerMod(g, y0.first, m);
                                                                   }
  lli inc = powerMod(g, phi_m / d, m);
  for(11i i = 1; i < d; i++)
                                                                   1.6. Particiones
    x[i] = x[i - 1] * inc % m;
  sort(x.begin(), x.end());
  return x;
                                                                   1.6.1. Función P (particiones de un entero positivo)
}
                                                                  11i \mod = 1e9 + 7;
1.5.10. Algoritmo de Tonelli-Shanks para raíces cuadradas
                                                                   vector<lli> P;
         módulo p
                                                                   //number of ways to write n as a sum of positive integers
//finds \ x \ such \ that \ x^2 = a \ mod \ p
                                                                   lli partitionsP(int n){
lli sqrtMod(lli a, lli p){
```

if(n < 0) return 0;

if(P[n]) return P[n];

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a %= p;

```
int pos1 = 1, pos2 = 2, inc1 = 4, inc2 = 5;
                                                                        if(j \& 1) return -1;
  lli ans = 0;
                                                                        else return 1;
  for(int k = 1; k \le n; k++){
                                                                      }else{
    lli tmp = (n \ge pos1 ? P[n - pos1] : 0) + (n \ge pos2 ? P[n]
                                                                        return 0;
    \rightarrow - pos2] : 0);
                                                                      }
    if (k \& 1) ans += tmp;
                                                                    }
    else ans -= tmp;
    if(n < pos2) break;</pre>
                                                                    //number of ways to write n as a sum of distinct positive
    pos1 += inc1, pos2 += inc2;
                                                                    \hookrightarrow integers
    inc1 += 3, inc2 += 3;
                                                                    //number of ways to write n as a sum of odd positive integers
  }
                                                                    lli partitionsQ(int n){
  ans %= mod;
                                                                      if(n < 0) return 0;
  if (ans < 0) ans += mod;
                                                                      if(Q[n]) return Q[n];
  return ans;
                                                                      int pos = 1, inc = 3;
}
                                                                      lli ans = 0;
                                                                      int limit = sqrt(n);
void calculateFunctionP(int n){
                                                                      for(int k = 1; k <= limit; k++){</pre>
  P.resize(n + 1);
                                                                        if (k \& 1) ans += Q[n - pos];
  P[0] = 1;
                                                                        else ans -= Q[n - pos];
  for(int i = 1; i <= n; i++)
                                                                        pos += inc;
    P[i] = partitionsP(i);
                                                                        inc += 2;
}
                                                                      }
                                                                      ans <<= 1;
                                                                      ans += s(n);
1.6.2. Función Q (particiones de un entero positivo en dis-
                                                                      ans %= mod;
        tintos sumandos)
                                                                      if (ans < 0) ans += mod;
                                                                      return ans:
                                                                    }
vector<lli>Q;
bool isPerfectSquare(int n){
                                                                    void calculateFunctionQ(int n){
  int r = sqrt(n);
                                                                      Q.resize(n + 1);
                                                                      Q[0] = 1;
  return r * r == n;
                                                                      for(int i = 1; i <= n; i++)
}
                                                                        Q[i] = partitionsQ(i);
                                                                    }
int s(int n){
  int r = 1 + 24 * n;
  if(isPerfectSquare(r)){
                                                                    1.6.3. Número de factorizaciones ordenadas
    int j;
    r = sqrt(r);
                                                                    //number of ordered factorizations of n
    if((r + 1) \% 6 == 0) j = (r + 1) / 6;
                                                                    lli orderedFactorizations(lli n){
    else j = (r - 1) / 6;
```

```
//skip the factorization if you already know the powers
auto fact = factorize(n);
int k = 0, q = 0;
vector<int> powers(fact.size() + 1);
for(auto & f : fact){
  powers[k + 1] = f.second;
  q += f.second;
  ++k;
}
vector<lli> prod(q + 1, 1);
//we need Ncr until the max_power+Omega(n) row
//module if needed
for(int i = 0; i \le q; i++){
 for(int j = 1; j \le k; j++){
    prod[i] = prod[i] * Ncr[powers[j] + i][powers[j]];
  }
}
lli ans = 0;
for(int j = 1; j \le q; j++){
 int alt = 1;
  for(int i = 0; i < j; i++){
    ans = ans + alt * Ncr[j][i] * prod[j - i - 1];
    alt *= -1;
  }
}
return ans;
```

1.6.4. Número de factorizaciones no ordenadas

```
//Number of unordered factorizations of n with
//largest part at most m
//Call unorderedFactorizations(n, n) to get all of them
//Add this to the main to speed up the map:
//mem.reserve(1024); mem.max_load_factor(0.25);
struct HASH{
    size_t operator()(const pair<int,int>&x)const{
        return hash<long long>()(((long long)x.first)^(((long operator)x.second)<<32));
    }</pre>
```

```
};
unordered_map<pair<int, int>, lli, HASH> mem;
lli unorderedFactorizations(int m, int n){
  if(m == 1 && n == 1) return 1;
  if(m == 1) return 0;
  if(n == 1) return 1;
  if(mem.count({m, n})) return mem[{m, n}];
 lli ans = 0;
  int 1 = sqrt(n);
 for(int i = 1; i <= 1; ++i){
    if(n \% i == 0){
      int a = i, b = n / i;
      if(a <= m) ans += unorderedFactorizations(a, b);
      if (a != b && b <= m) ans += unorderedFactorizations(b,
      \rightarrow a);
    }
  }
  return mem[{m, n}] = ans;
}
```

1.7. Otros

1.7.1. Cambio de base

```
string decimalToBaseB(lli n, lli b){
   string ans = "";
   lli d;
   do{
      d = n % b;
      if(0 <= d && d <= 9) ans = (char)(48 + d) + ans;
      else if(10 <= d && d <= 35) ans = (char)(55 + d) + ans;
      n /= b;
   }while(n != 0);
   return ans;
}

lli baseBtoDecimal(const string & n, lli b){
   lli ans = 0;
   for(const char & d : n){
      if(48 <= d && d <= 57) ans = ans * b + (d - 48);</pre>
```

}

```
else if(65 <= d && d <= 90) ans = ans * b + (d - 55);
else if(97 <= d && d <= 122) ans = ans * b + (d - 87);
}
return ans;
}
```

1.7.2. Fracciones continuas

```
//continued fraction of (p+sqrt(n))/q, where p,n,q are positive
\hookrightarrow integers
//returns a vector of terms and the length of the period,
//the periodic part is taken from the right of the array
pair<vector<lli>, int> ContinuedFraction(lli p, lli n, lli q){
  vector<lli> coef;
  lli r = sqrt(n);
  //Skip this if you know that n is not a perfect square
  if(r * r == n){
    lli num = p + r;
    lli den = q;
    lli residue;
    while(den){
      residue = num % den;
      coef.push_back(num / den);
      num = den;
      den = residue;
    }
    return make_pair(coef, 0);
  }
  if((n - p * p) % q != 0){
    n *= q * q;
    p *= q;
    q *= q;
    r = sqrt(n);
  }
  lli a = (r + p) / q;
  coef.push_back(a);
  int period = 0;
  map<pair<lli, lli>, int> pairs;
  while(true){
    p = a * q - p;
```

```
q = (n - p * p) / q;
a = (r + p) / q;
//if p=0 and q=1, we can just ask if q==1 after inserting a
if(pairs.count(make_pair(p, q))){
   period -= pairs[make_pair(p, q)];
   break;
}
coef.push_back(a);
pairs[make_pair(p, q)] = period++;
}
return make_pair(coef, period);
```

1.7.3. Ecuación de Pell

```
//first solution (x, y) to the equation x^2-ny^2=1, n IS NOT a
    → perfect aquare
pair<lli, lli> PellEquation(lli n){
    vector<lli> cf = ContinuedFraction(0, n, 1).first;
    lli num = 0, den = 1;
    int k = cf.size() - 1;
    for(int i = ((k & 1) ? (2 * k - 1) : (k - 1)); i >= 0; i--){
        lli tmp = den;
        int pos = i % k;
        if(pos == 0 && i != 0) pos = k;
        den = num + cf[pos] * den;
        num = tmp;
    }
    return make_pair(den, num);
}
```

1.7.4. Números de Bell

```
//number of ways to partition a set of n elements
//the nth bell number is at Bell[n][0]
vector<vector<int>>> Bell;
void bellSieve(int n){
   Bell.resize(n + 1);
   Bell[0] = {1};
```

```
for(int i = 1; i <= n; ++i){
                                                                       q = p * p;
    Bell[i].resize(i + 1);
                                                                       hi[1] = (hi[p] - temp) * powMod(p, k, Mod) % Mod;
    Bell[i][0] = Bell[i - 1][i - 1];
                                                                       if(hi[1] < 0) hi[1] += Mod;
    for(int j = 1; j <= i; ++j)
                                                                       j = 1 + (p \& 1);
      Bell[i][j] = Bell[i][j-1] + Bell[i-1][j-1];
                                                                       end = (v \le n/q) ? v : n/q;
  }
                                                                       for(i = p + j; i \le 1 + end; i += j){
}
                                                                         if(used[i]) continue;
                                                                         d = i * p;
                                                                         if(d \ll v)
1.7.5. Prime counting function in sublinear time
                                                                           hi[i] = (hi[d] - temp) * powMod(p, k, Mod) % Mod;
                                                                         else
const lli inv_2 = modularInverse(2, Mod);
                                                                           hi[i] = (lo[n/d] - temp) * powMod(p, k, Mod) % Mod;
const lli inv_6 = modularInverse(6, Mod);
                                                                         if(hi[i] < 0) hi[i] += Mod;
const lli inv_30 = modularInverse(30, Mod);
                                                                       if(q \ll v)
lli sum(lli n, int k){
                                                                         for(i = q; i \le end; i += p*j)
  n \%= Mod;
                                                                           used[i] = true;
  if(k == 0) return n:
                                                                       for(i = v; i >= q; i--){
  if(k == 1) return n * (n + 1) % Mod * inv_2 % Mod;
                                                                         lo[i] = (lo[i/p] - temp) * powMod(p, k, Mod) % Mod;
  if(k == 2) return n * (n + 1) % Mod * (2*n + 1) % Mod * inv_6
                                                                         if(lo[i] < 0) lo[i] += Mod;

→ % Mod;

                                                                       }
  if (k == 3) return powMod(n * (n + 1) \% Mod * inv_2 \% Mod, 2,
                                                                     }
  \rightarrow Mod);
                                                                     return hi[1] % Mod;
  if(k == 4) return n * (n + 1) % Mod * (2*n + 1) % Mod *
                                                                   }
  \rightarrow (3*n*(n+1)%Mod -1) % Mod * inv_30 % Mod;
  return 1;
}
//finds the sum of the kth powers of the primes
//less than or equal to n (0 \le k \le 4, add more if you need)
lli SumPrimePi(lli n, int k){
  lli v = sqrt(n), p, temp, q, j, end, i, d;
  vector<lli> lo(v+2), hi(v+2);
  vector<bool> used(v+2);
  for(p = 1; p \le v; p++){
   lo[p] = sum(p, k) - 1;
   hi[p] = sum(n/p, k) - 1;
  }
  for(p = 2; p \leq v; p++){
    if(lo[p] == lo[p-1]) continue;
    temp = lo[p-1];
```

2. Números racionales

2.1. Estructura fraccion

```
struct fraccion{
   ll num, den;
   fraccion(){
       num = 0, den = 1;
   fraccion(ll x, ll y){
       if(y < 0)
           x *= -1, y *=-1;
       ll d = \_gcd(abs(x), abs(y));
       num = x/d, den = y/d;
   fraccion(ll v){
        num = v;
        den = 1;
   fraccion operator+(const fraccion& f) const{
       ll d = \_gcd(den, f.den);
        return fraccion(num*(f.den/d) + f.num*(den/d),
        \rightarrow den*(f.den/d));
   }
   fraccion operator-() const{
        return fraccion(-num, den);
   fraccion operator-(const fraccion& f) const{
       return *this + (-f);
   }
   fraccion operator*(const fraccion& f) const{
        return fraccion(num*f.num, den*f.den);
   }
   fraccion operator/(const fraccion& f) const{
        return fraccion(num*f.den, den*f.num);
   fraccion operator+=(const fraccion& f){
        *this = *this + f;
       return *this;
   fraccion operator = (const fraccion& f){
```

```
*this = *this - f;
    return *this;
fraccion operator++(int xd){
    *this = *this + 1;
    return *this;
fraccion operator--(int xd){
    *this = *this - 1;
    return *this;
fraccion operator*=(const fraccion& f){
    *this = *this * f;
    return *this;
}
fraccion operator/=(const fraccion& f){
    *this = *this / f;
    return *this;
}
bool operator == (const fraccion& f) const{
    11 d = \_gcd(den, f.den);
    return (num*(f.den/d) == (den/d)*f.num);
bool operator!=(const fraccion& f) const{
    11 d = \_gcd(den, f.den);
    return (num*(f.den/d) != (den/d)*f.num);
bool operator >(const fraccion& f) const{
    11 d = \_gcd(den, f.den);
    return (num*(f.den/d) > (den/d)*f.num);
bool operator <(const fraccion& f) const{</pre>
    11 d = \_gcd(den, f.den);
    return (num*(f.den/d) < (den/d)*f.num);
}
bool operator >=(const fraccion& f) const{
    11 d = \_gcd(den, f.den);
    return (num*(f.den/d) >= (den/d)*f.num);
}
bool operator <=(const fraccion& f) const{</pre>
    11 d = \_gcd(den, f.den);
```

```
return (num*(f.den/d) <= (den/d)*f.num);
    }
    fraccion inverso() const{
        return fraccion(den, num);
    }
    fraccion fabs() const{
        fraccion nueva:
        nueva.num = abs(num);
        nueva.den = den;
        return nueva;
    }
    double value() const{
      return (double)num / (double)den;
    string str() const{
        stringstream ss;
        ss << num;
        if(den != 1) ss << "/" << den;
        return ss.str();
    }
};
ostream &operator << (ostream &os, const fraccion & f) {
    return os << f.str();
}
istream &operator>>(istream &is, fraccion & f){
    11 \text{ num} = 0, \text{ den} = 1;
    string str;
    is >> str;
    size_t pos = str.find("/");
    if(pos == string::npos){
        istringstream(str) >> num;
    }else{
        istringstream(str.substr(0, pos)) >> num;
        istringstream(str.substr(pos + 1)) >> den;
    f = fraccion(num, den);
    return is;
}
```

3. Álgebra lineal

3.1. Estructura matrix

```
template <typename T>
struct matrix{
 vector<vector<T>> A;
 int m, n;
 matrix(int m, int n): m(m), n(n){
   A.resize(m, vector<T>(n, 0));
 vector<T> & operator[] (int i){
   return A[i];
  const vector<T> & operator[] (int i) const{
   return A[i];
  static matrix identity(int n){
   matrix<T> id(n, n);
   for(int i = 0; i < n; i++)
     id[i][i] = 1;
   return id;
 matrix operator+(const matrix & B) const{
   assert(m == B.m && n == B.n); //same dimensions
   matrix<T> C(m, n);
   for(int i = 0; i < m; i++)
     for(int j = 0; j < n; j++)
       C[i][j] = A[i][j] + B[i][j];
   return C;
 }
 matrix operator+=(const matrix & M){
    *this = *this + M;
   return *this;
```

```
}
                                                                   }
matrix operator-() const{
                                                                   matrix operator*=(const T & c){
  matrix<T> C(m, n);
                                                                     *this = *this * c;
  for(int i = 0; i < m; i++)
                                                                     return *this;
    for(int j = 0; j < n; j++)
                                                                   }
      C[i][j] = -A[i][j];
  return C;
                                                                   matrix operator^(lli b) const{
}
                                                                     matrix<T> ans = matrix<T>::identity(n);
                                                                     matrix<T> A = *this;
matrix operator-(const matrix & B) const{
                                                                     while(b){
  return *this + (-B);
                                                                       if (b & 1) ans *= A;
}
                                                                       b >>= 1;
                                                                       if(b) A *= A;
                                                                     }
matrix operator-=(const matrix & M){
  *this = *this + (-M);
                                                                     return ans;
                                                                   }
  return *this;
}
                                                                   matrix operator^=(lli n){
                                                                     *this = *this ^ n;
matrix operator*(const matrix & B) const{
  assert(n == B.m); //#columns of 1st matrix = #rows of 2nd
                                                                     return *this;
  \rightarrow matrix
  matrix<T> C(m, B.n);
  for(int i = 0; i < m; i++)
                                                                   bool operator==(const matrix & B) const{
   for(int j = 0; j < B.n; j++)
                                                                     if(m != B.m || n != B.n) return false;
      for(int k = 0; k < n; k++)
                                                                     for(int i = 0; i < m; i++)
        C[i][j] += A[i][k] * B[k][j];
                                                                       for(int j = 0; j < n; j++)
                                                                         if(A[i][j] != B[i][j]) return false;
  return C;
}
                                                                     return true;
                                                                   }
matrix operator*(const T & c) const{
  matrix<T> C(m, n);
                                                                   bool operator!=(const matrix & B) const{
  for(int i = 0; i < m; i++)
                                                                     return !(*this == B);
   for(int j = 0; j < n; j++)
                                                                   }
      C[i][j] = A[i][j] * c;
  return C;
                                                                   void scaleRow(int k, T c){
}
                                                                     for(int j = 0; j < n; j++)
                                                                       A[k][j] *= c;
matrix operator*=(const matrix & M){
                                                                   }
  *this = *this * M;
  return *this;
                                                                   void swapRows(int k, int 1){
```

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```
swap(A[k], A[1]);
}

void addRow(int k, int 1, T c){
  for(int j = 0; j < n; j++)
     A[k][j] += c * A[1][j];
}</pre>
```

3.2. Transpuesta y traza

```
matrix<T> transpose(){
   matrix<T> tr(n, m);
   for(int i = 0; i < m; i++)
      for(int j = 0; j < n; j++)
        tr[j][i] = A[i][j];
   return tr;
}

T trace(){
   T sum = 0;
   for(int i = 0; i < min(m, n); i++)
      sum += A[i][i];
   return sum;
}</pre>
```

3.3. Gauss Jordan

```
//full: true: reduce above and below the diagonal, false:
    reduce only below
//makeOnes: true: make the elements in the diagonal ones,
    false: leave the diagonal unchanged
//For every elemental operation that we apply to the matrix,
//we will call to callback(operation, k, l, value).
//operation 1: multiply row "k" by "value"
//operation 2: swap rows "k" and "l"
//operation 3: add "value" times the row "l" to the row "k"
//It returns the rank of the matrix, and modifies it
int gauss_jordan(bool full = true, bool makeOnes = true,
    function

youngle
```

```
int i = 0, j = 0;
  while(i < m \&\& j < n){
   if(A[i][j] == 0){
      for(int f = i + 1; f < m; f++){
        if(A[f][i] != 0){
          swapRows(i, f);
          if(callback) callback(2, i, f, 0);
          break;
        }
     }
   }
   if(A[i][j] != 0){
     T inv_mult = A[i][j].inverso();
      if(makeOnes && A[i][j] != 1){
        scaleRow(i, inv_mult);
        if(callback) callback(1, i, 0, inv_mult);
      for(int f = (full ? 0 : (i + 1)); f < m; f++){
        if(f != i && A[f][j] != 0){
          T inv_adit = -A[f][j];
          if(!makeOnes) inv_adit *= inv_mult;
          addRow(f, i, inv_adit);
          if(callback) callback(3, f, i, inv_adit);
        }
     }
     i++;
   }
   j++;
 return i;
}
void gaussian_elimination(){
  gauss_jordan(false);
```

3.4. Matriz escalonada por filas y reducida por filas

```
matrix<T> reducedRowEchelonForm(){
  matrix<T> asoc = *this;
```

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```
asoc.gauss_jordan();
return asoc;
}

matrix<T> rowEchelonForm(){
  matrix<T> asoc = *this;
  asoc.gaussian_elimination();
  return asoc;
}
```

3.5. Matriz inversa

```
bool invertible(){
  assert(m == n); //this is defined only for square matrices
 matrix<T> tmp = *this;
  return tmp.gauss_jordan(false) == n;
}
matrix<T> inverse(){
  assert(m == n); //this is defined only for square matrices
  matrix<T> tmp = *this;
  matrix<T> inv = matrix<T>::identity(n);
  auto callback = [&](int op, int a, int b, T e){
   if(op == 1){
     inv.scaleRow(a, e);
   else if(op == 2){
     inv.swapRows(a, b);
   else if(op == 3){
     inv.addRow(a, b, e);
   }
 };
  assert(tmp.gauss_jordan(true, true, callback) == n);
  return inv;
}
```

3.6. Determinante

```
T determinant(){
   assert(m == n); //only square matrices have determinant
   matrix<T> tmp = *this;
   T det = 1;
   auto callback = [&](int op, int a, int b, T e){
      if(op == 1) {
       det /= e;
    }else if(op == 2) {
       det *= -1;
    }
   };
   if(tmp.gauss_jordan(false, true, callback) != n) det = 0;
   return det;
}
```

3.7. Matriz de cofactores y adjunta

```
matrix<T> minor(int x, int y){
  matrix<T> M(m-1, n-1);
 for(int i = 0; i < m-1; ++i)
   for(int j = 0; j < n-1; ++j)
     M[i][j] = A[i < x ? i : i+1][j < y ? j : j+1];
  return M;
}
T cofactor(int x, int y){
  T ans = minor(x, y).determinant();
  if((x + y) \% 2 == 1) ans *= -1;
  return ans;
matrix<T> cofactorMatrix(){
  matrix<T> C(m, n);
 for(int i = 0; i < m; i++)
   for(int j = 0; j < n; j++)
     C[i][j] = cofactor(i, j);
 return C;
}
```

```
matrix<T> adjugate(){
   if(invertible()) return inverse() * determinant();
   return cofactorMatrix().transpose();
}
```

3.8. Factorización PA = LU

```
tuple<matrix<T>, matrix<T>, matrix<T>> PA_LU(){
  matrix<T> U = *this:
  matrix<T> L = matrix<T>::identity(n);
  matrix<T> P = matrix<T>::identity(n);
  auto callback = [&](int op, int a, int b, T e){
    if(op == 2){
     L.swapRows(a, b);
      P.swapRows(a, b);
      L[a][a] = L[b][b] = 1;
      L[a][a + 1] = L[b][b - 1] = 0;
    else if(op == 3){
      L[a][b] = -e;
    }
  };
  U.gauss_jordan(false, false, callback);
  return {P, L, U};
}
```

3.9. Polinomio característico

```
vector<T> characteristicPolynomial(){
   matrix<T> M(n, n);
   vector<T> coef(n + 1);
   matrix<T> I = matrix<T>::identity(n);
   coef[n] = 1;
   for(int i = 1; i <= n; i++){
      M = (*this) * M + I * coef[n - i + 1];
      coef[n - i] = -((*this) * M).trace() / i;
   }
   return coef;
}</pre>
```

3.10. Gram-Schmidt

```
matrix<T> gram_schmidt(){
  //vectors are rows of the matrix (also in the answer)
  //the answer doesn't have the vectors normalized
  matrix<T> B = (*this) * (*this).transpose();
  matrix<T> ans = *this;
  auto callback = [&](int op, int a, int b, T e){
   if(op == 1){
      ans.scaleRow(a, e);
   else if(op == 2){
      ans.swapRows(a, b);
   else if(op == 3){
      ans.addRow(a, b, e);
   }
  }:
  B.gauss_jordan(false, false, callback);
  return ans;
}
```

3.11. Recurrencias lineales

```
//Solves a linear homogeneous recurrence relation of degree
→ "deg" of the form
//F(n) = a(d-1)*F(n-1) + a(d-2)*F(n-2) + \dots + a(1)*F(n-(d-1))
\rightarrow + a(0)*F(n-d)
//with initial values F(0), F(1), ..., F(d-1)
//It finds the nth term of the recurrence, F(n)
//The values of a[0,...,d) are in the array P[]
lli solveRecurrence(lli *P, lli *init, int deg, lli n){
  lli *ans = new lli[deg]();
 lli *R = new lli[2*deg]();
  ans[0] = 1;
  lli p = 1;
  for(lli v = n; v >>= 1; p <<= 1);
    int d = (n \& p) != 0;
    fill(R, R + 2*deg, 0);
    //if deg(mod-1)^2 overflows, just do mod in the
    \rightarrow multiplications
```

```
for(int i = 0; i < deg; i++)
      for(int j = 0; j < deg; j++)
                                                                     for(int j = 0; j < m; ++j){
        R[i + j + d] += ans[i] * ans[j];
                                                                       for(int i = 0; i < n; ++i)
    for(int i = 0; i < 2*deg; ++i) R[i] %= mod;
                                                                         T[i][i] = A[i][i];
    for(int i = deg-1; i >= 0; i--){
                                                                       row[j] = n + j;
      R[i + deg] \% = mod;
                                                                       T[i][n + i] = 1;
      for(int j = 0; j < deg; j++)
                                                                       base[n + j] = 1;
        R[i + j] += R[i + deg] * P[j];
                                                                       T[j][n + m] = b[j];
                                                                     }
    for(int i = 0; i < deg; i++) R[i] \% = mod;
    copy(R, R + deg, ans);
                                                                     for(int i = 0; i < n; ++i)
  }while(p >>= 1);
                                                                       T[m][i] = c[i] * (mini ? 1 : -1);
  lli nValue = 0;
  for(int i = 0; i < deg; i++)
                                                                     while(true){
    nValue += ans[i] * init[i];
                                                                       int p = 0, q = 0;
  return nValue % mod;
                                                                       for(int i = 0; i < n + m; ++i)
}
                                                                         if(T[m][i] <= T[m][p])
                                                                           p = i;
3.12. Simplex
                                                                       for(int j = 0; j < m; ++j)
                                                                         if(T[j][n + m] \le T[q][n + m])
                                                                           q = j;
Parametric Self-Dual Simplex method
Solve a canonical LP:
                                                                       double t = min(T[m][p], T[q][n + m]);
  min or max. c x
  s.t. A x \le b
                                                                       if(t \ge -eps){
    x >= 0
                                                                         vec x(n);
                                                                         for(int i = 0; i < m; ++i)
#include <bits/stdc++.h>
                                                                           if(row[i] < n) x[row[i]] = T[i][n + m];
using namespace std;
                                                                         return {x, T[m][n + m] * (mini ? -1 : 1)}; // optimal
const double eps = 1e-9, oo =
                                                                       }
→ numeric_limits<double>::infinity();
                                                                       if(t < T[q][n + m]){
typedef vector<double> vec;
                                                                         // tight on c -> primal update
typedef vector<vec> mat;
                                                                         for(int j = 0; j < m; ++j)
                                                                           if(T[j][p] >= eps)
pair < vec, double > simplexMethodPD (mat &A, vec &b, vec &c, bool
                                                                             if(T[j][p] * (T[q][n + m] - t) >= T[q][p] * (T[j][n +

    mini = true){
                                                                              \rightarrow m] - t))
  int n = c.size(), m = b.size();
                                                                               q = j;
  mat T(m + 1, vec(n + m + 1));
```

 $if(T[q][p] \le eps)$

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vector<int> base(n + m), row(m);

```
return {vec(n), oo * (mini ? 1 : -1)}; // primal
                                                                        cout << "Numero de restricciones: ";</pre>
         \hookrightarrow infeasible
                                                                        cin >> m;
    }else{
                                                                        cout << "Numero de incognitas: ";</pre>
      // tight on b -> dual update
                                                                        cin >> n;
      for(int i = 0; i < n + m + 1; ++i)
                                                                        mat A(m, \text{vec}(n));
        T[q][i] = -T[q][i];
                                                                        vec b(m), c(n);
                                                                        for(int i = 0; i < m; ++i){
      for(int i = 0; i < n + m; ++i)
                                                                          cout << "Restriccion #" << (i + 1) << ": ";</pre>
        if(T[q][i] >= eps)
                                                                          for(int j = 0; j < n; ++j){
          if(T[q][i] * (T[m][p] - t) >= T[q][p] * (T[m][i] -
                                                                            cin >> A[i][j];
                                                                          }
                                                                          cin >> b[i];
            p = i;
      if(T[q][p] \le eps)
                                                                        cout << "[0]Max o [1]Min?: ";</pre>
        return {vec(n), oo * (mini ? -1 : 1)}; // dual
                                                                        cin >> mini;
         \hookrightarrow infeasible
                                                                        cout << "Coeficientes de " << (mini ? "min" : "max") << " z:</pre>
    }
                                                                        for(int i = 0; i < n; ++i){
    for(int i = 0; i < m + n + 1; ++i)
                                                                          cin >> c[i];
      if(i != p) T[q][i] /= T[q][p];
                                                                        }
                                                                        cout.precision(6);
    T[q][p] = 1; // pivot(q, p)
                                                                        auto ans = simplexMethodPD(A, b, c, mini);
    base[p] = 1;
                                                                        cout << (mini ? "Min" : "Max") << " z = " << ans.second << ",
    base[row[q]] = 0;

→ cuando: \n";

    row[q] = p;
                                                                        for(int i = 0; i < ans.first.size(); ++i){</pre>
                                                                          cout << "x_" << (i + 1) << " = " << ans.first[i] << "\n";</pre>
    for(int j = 0; j < m + 1; ++j){
                                                                        }
      if(j != q){
                                                                        return 0;
        double alpha = T[j][p];
        for(int i = 0; i < n + m + 1; ++i)
          T[j][i] = T[q][i] * alpha;
      }
    }
  }
  return {vec(n), oo};
int main(){
  int m, n;
  bool mini = true;
```

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}

4. FFT

4.1. Declaraciones previas

```
typedef complex<double> comp;
typedef long long int lli;
double PI = acos(-1.0);

int nearestPowerOfTwo(int n){
  int ans = 1;
  while(ans < n) ans <<= 1;
  return ans;
}</pre>
```

4.2. FFT con raíces de la unidad complejas

```
void fft(vector<comp> & X, int inv){
  int n = X.size();
  int len, len2, i, j, k;
  for(i = 1, j = 0; i < n - 1; ++i){
   for (k = n >> 1; (j = k) < k; k >>= 1);
    if(i < j) swap(X[i], X[j]);
  }
  double ang;
  comp t, u, v;
  vector<comp> wlen_pw(n >> 1);
  wlen_pw[0] = 1;
  for(len = 2; len <= n; len <<= 1){
    ang = inv == -1 ? -2 * PI / len : 2 * PI / len;
    len2 = len >> 1;
    comp wlen(cos(ang), sin(ang));
    for(i = 1; i < len2; ++i)
      wlen_pw[i] = wlen_pw[i - 1] * wlen;
    for(i = 0; i < n; i += len)
      for(j = 0; j < len2; ++j){}
        t = X[i + j + len2] * wlen_pw[j];
        X[i + j + len2] = X[i + j] - t;
        X[i + j] += t;
      }
  }
```

```
if(inv == -1)
  for(i = 0; i < n; ++i)
    X[i] /= n;
}</pre>
```

4.3. FFT con raíces de la unidad en $\mathbb{Z}/p\mathbb{Z}$ (NTT)

```
int inverse(int a, int n){
 int r0 = a, r1 = n, ri, s0 = 1, s1 = 0, si;
 while(r1){
   si = s0 - s1 * (r0 / r1), s0 = s1, s1 = si;
   ri = r0 \% r1, r0 = r1, r1 = ri;
 if(s0 < 0) s0 += n;
 return s0;
int p = 7340033;
int root = 5;
int root_1 = inverse(root, p);
int root_pw = 1 << 20;</pre>
void ntt(vector<int> & X, int inv){
 int n = X.size();
 int len, len2, wlen, i, j, k, u, v, w;
 for(i = 1, j = 0; i < n - 1; ++i){
   for (k = n >> 1; (j = k) < k; k >>= 1);
   if(i < j) swap(X[i], X[j]);
 for(len = 2; len <= n; len <<= 1){
   len2 = len >> 1;
   wlen = (inv == -1) ? root_1 : root;
   for(i = len; i < root_pw; i <<= 1)
     wlen = (lli)wlen * wlen % p;
   for(i = 0; i < n; i += len)
     for(j = 0, w = 1; j < len2; ++j){
       u = X[i + j], v = (11i)X[i + j + len2] * w % p;
       X[i + j] = u + v 
       X[i + j + len2] = u - v < 0 ? u - v + p : u - v;
       w = (lli)w * wlen % p;
```

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```
}
if(inv == -1){
  int nrev = inverse(n, p);
  for(i = 0; i < n; ++i)
    X[i] = (lli)X[i] * nrev % p;
}</pre>
```

4.3.1. Otros valores para escoger la raíz y el módulo

| Raíz n-ési- | ω^{-1} | Tamaño | Módulo p |
|-------------------|---------------|--------------|-------------------------------------|
| | ω | | Wodulo p |
| ma de la | | máximo del | |
| unidad (ω) | | arreglo(n) | |
| 15 | 30584 | 2^{14} | $4 \times 2^{14} + 1 = 65537$ |
| 9 | 7282 | 2^{15} | $2 \times 2^{15} + 1 = 65537$ |
| 3 | 21846 | 2^{16} | $1 \times 2^{16} + 1 = 65537$ |
| 8 | 688129 | 2^{17} | $6 \times 2^{17} + 1 = 786433$ |
| 5 | 471860 | 2^{18} | $3 \times 2^{18} + 1 = 786433$ |
| 12 | 3364182 | 2^{19} | $11 \times 2^{19} + 1 = 5767169$ |
| 5 | 4404020 | 2^{20} | $7 \times 2^{20} + 1 = 7340033$ |
| 38 | 21247462 | 2^{21} | $11 \times 2^{21} + 1 = 23068673$ |
| 21 | 49932191 | 2^{22} | $25 \times 2^{22} + 1 = 104857601$ |
| 4 | 125829121 | 2^{23} | $20 \times 2^{23} + 1 = 167772161$ |
| 31 | 128805723 | $3 \ 2^{23}$ | $119 \times 2^{23} + 1 = 998244353$ |
| 2 | 83886081 | 2^{24} | $10 \times 2^{24} + 1 = 167772161$ |
| 17 | 29606852 | 2^{25} | $5 \times 2^{25} + 1 = 167772161$ |
| 30 | 15658735 | 2^{26} | $7 \times 2^{26} + 1 = 469762049$ |
| 137 | 749463956 | 2^{27} | $15 \times 2^{27} + 1 = 2013265921$ |
| | | | |

4.4. Multiplicación de polinomios (convolución lineal)

```
void convolution(vector<comp> & A, vector<comp> & B){
  int sz = A.size() + B.size() - 1;
  int size = nearestPowerOfTwo(sz);
  A.resize(size), B.resize(size);
```

```
fft(A, 1), fft(B, 1);
for(int i = 0; i < size; i++)
    A[i] *= B[i];
fft(A, -1);
    A.resize(sz);
}

void convolution(vector<int> & A, vector<int> & B){
    int sz = A.size() + B.size() - 1;
    int size = nearestPowerOfTwo(sz);
    A.resize(size), B.resize(size);
    ntt(A, 1), ntt(B, 1);
    for(int i = 0; i < size; i++)
        A[i] = (lli)A[i] * B[i] % p;
    ntt(A, -1);
    A.resize(sz);
}</pre>
```

4.5. Aplicaciones

4.5.1. Multiplicación de números enteros grandes

```
string multiplyNumbers(const string & a, const string & b){
 int sgn = 1;
 int pos1 = 0, pos2 = 0;
 while(pos1 < a.size() && (a[pos1] < '1' || a[pos1] > '9')){
   if(a[pos1] == '-') sgn *= -1;
   ++pos1;
 }
 while(pos2 < b.size() && (b[pos2] < '1' || b[pos2] > '9')){
   if(b[pos2] == '-') sgn *= -1;
   ++pos2;
 }
 vector<int> X(a.size() - pos1), Y(b.size() - pos2);
 if(X.empty() || Y.empty()) return "0";
 for(int i = pos1, j = X.size() - 1; i < a.size(); ++i)
   X[j--] = a[i] - '0';
 for(int i = pos2, j = Y.size() - 1; i < b.size(); ++i)</pre>
   Y[j--] = b[i] - '0';
 convolution(X, Y);
```

```
stringstream ss;
  if(sgn == -1) ss << "-";
  int carry = 0;
  for(int i = 0; i < X.size(); ++i){</pre>
    X[i] += carry;
    carry = X[i] / 10;
    X[i] \% = 10;
  }
  while(carry){
    X.push_back(carry % 10);
    carry /= 10;
  }
  for(int i = X.size() - 1; i >= 0; --i)
    ss << X[i];
  return ss.str();
}
```

4.5.2. Inverso de un polinomio

```
vector<int> inversePolynomial(const vector<int> & A){
  vector<int> R(1, inverse(A[0], p));
  //R(x) = 2R(x) - A(x)R(x)^2
  while(R.size() < A.size()){</pre>
    int c = 2 * R.size();
    R.resize(c);
    vector<int> TR = R;
    TR.resize(2 * c);
    vector<int> TF(TR.size());
    for(int i = 0; i < c && i < A.size(); ++i)
      TF[i] = A[i];
    ntt(TR, 1);
    ntt(TF, 1);
    for(int i = 0; i < TR.size(); ++i)
      TR[i] = (lli)TR[i] * TR[i] % p * TF[i] % p;
    ntt(TR, -1);
    for(int i = 0; i < c; ++i){
      R[i] = R[i] + R[i] - TR[i];
      while(R[i] < 0) R[i] += p;
      while(R[i] >= p) R[i] -= p;
    }
```

```
}
  R.resize(A.size());
  return R;
}
```

4.5.3. Raíz cuadrada de un polinomio

```
const int inv2 = inverse(2, p);
vector<int> sqrtPolynomial(const vector<int> & A){
  int r0 = 1; //verify that r0^2 = A[0] \mod p
 vector<int> R(1, r0);
 //R(x) = R(x)/2 + A(x)/(2R(x))
 while(R.size() < A.size()){</pre>
   int c = 2 * R.size();
   R.resize(c);
   vector<int> TF(c):
   for(int i = 0; i < c && i < A.size(); ++i)</pre>
     TF[i] = A[i];
    vector<int> IR = inversePolynomial(R);
    convolution(TF, IR);
   for(int i = 0; i < c; ++i){
     R[i] = R[i] + TF[i];
     if(R[i] >= p) R[i] -= p;
     R[i] = (11i)R[i] * inv2 % p;
   }
 R.resize(A.size());
 return R;
}
```

4.5.4. Cociente y residuo de dos polinomios

```
pair<vector<int>, vector<int>> divide(vector<int> & A,
    vector<int> & B){
    while(!A.empty() && A.back() == 0) A.pop_back();
    while(!B.empty() && B.back() == 0) B.pop_back();
    int n = A.size(), m = B.size();
    if(n < m) return {A, B};</pre>
```

```
if(m == 1)
                                                                    comp w = polar(1.0, PI / n), w1 = w, w2 = 1;
    int inv = inverse(B[0], p);
                                                                    vector<comp> p(m), q(m), b(n);
    for(int i = 0; i < n; ++i)
                                                                    for(int k = 0; k < n; ++k, w2 *= w1, w1 *= w*w){}
      A[i] = (lli)A[i] * inv % p;
                                                                      b[k] = w2;
    return \{A, \{0\}\};
                                                                      p[k] = a[k] * b[k];
                                                                      q[k] = (comp)1 / b[k];
  vector<int> Q = A, Binv = B;
                                                                      if(k) q[m-k] = q[k];
  reverse(Q.begin(), Q.end());
  reverse(Binv.begin(), Binv.end());
                                                                    fft(p, 1), fft(q, 1);
  Binv.resize(n-m+1);
                                                                    for(int i = 0; i < m; i++)
  Binv = inversePolynomial(Binv);
                                                                      p[i] *= q[i];
  Q.resize(n-m+1);
                                                                    fft(p, -1);
  convolution(Q, Binv);
                                                                    for(int k = 0; k < n; ++k)
  Q.resize(n-m+1);
                                                                      a[k] = b[k] * p[k];
                                                                  }
  reverse(Q.begin(), Q.end());
  vector<int> R = B;
  R.resize(m-1);
                                                                  4.6. Convolución de dos vectores reales con solo dos
  vector<int> q = Q;
                                                                         FFT's
  Q.resize(min(n-m+1, m-1));
  convolution(R, Q);
  R.resize(m-1);
                                                                  //A and B are real-valued vectors
  for(int i = 0; i < m-1; ++i){
                                                                  //just do 2 fft's instead of 3
   R[i] = A[i] - R[i];
                                                                  void convolutionTrick(vector<comp> & A, vector<comp> & B){
    while (R[i] < 0) R[i] += p;
                                                                    int sz = A.size() + B.size() - 1;
    while(R[i] >= p) R[i] -= p;
                                                                    int size = nearestPowerOfTwo(sz);
  }
                                                                    vector<comp> C(size);
  while(!R.empty() && R.back() == 0) R.pop_back();
                                                                    comp I(0, 1);
  if(R.empty()) R.push_back(0);
                                                                    for(int i = 0; i < A.size() || i < B.size(); ++i){</pre>
  return {q, R};
                                                                      if(i < A.size()) C[i] += A[i];
}
                                                                      if(i < B.size()) C[i] += I*B[i];
                                                                    fft(C, 1);
4.5.5. DFT con tamaño de vector arbitrario (algoritmo de
                                                                    A.resize(size);
        Bluestein)
                                                                    for(int i = 0, j = 0; i < size; ++i){
                                                                      j = (size-1) & (size-i);
//it evaluates 1, w^2, w^4, ..., w^2(2n-2) on the polynomial
                                                                      A[i] = (conj(C[i]*C[i]) - C[i]*C[i]) * 0.25 * I;
//in this example we do a DFT with arbitrary size
                                                                    fft(A, -1);
```

A.resize(sz);

}

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ESCOM-IPN

void bluestein(vector<comp> & a){

int m = nearestPowerOfTwo(2*n-1);

int n = a.size():

4.7. Convolución con módulo arbitrario

```
//convolution with arbitrary modulo using only 4 fft's
void convolutionMod(vector<int> & A, vector<int> & B, int mod){
  int s = sqrt(mod);
  int sz = A.size() + B.size() - 1:
  int size = nearestPowerOfTwo(sz);
  vector<comp> a(size), b(size);
  for(int i = 0; i < A.size(); ++i)</pre>
    a[i] = comp(A[i] \% s, A[i] / s);
  for(int i = 0; i < B.size(); ++i)</pre>
    b[i] = comp(B[i] \% s, B[i] / s);
  fft(a, 1), fft(b, 1);
  comp I(0, 1);
  vector<comp> c(size), d(size);
  for(int i = 0, j = 0; i < size; ++i){}
    j = (size-1) & (size-i);
    comp e = (a[i] + conj(a[j])) * 0.5;
    comp f = (conj(a[j]) - a[i]) * 0.5 * I;
    comp g = (b[i] + conj(b[j])) * 0.5;
    comp h = (conj(b[j]) - b[i]) * 0.5 * I;
    c[i] = e * g + I * (e * h + f * g);
    d[i] = f * h;
  }
  fft(c, -1), fft(d, -1);
  A.resize(sz);
  for(int i = 0, j = 0; i < sz; ++i){
    j = (size-1) & (size-i);
    lli p0 = (lli)round(real(c[i])) % mod;
    lli p1 = (lli)round(imag(c[i])) % mod;
    lli p2 = (lli)round(real(d[i])) % mod;
    A[i] = p0 + s*(p1 + p2*s \% mod) \% mod;
    while(A[i] >= mod) A[i] -= mod;
    while(A[i] < 0) A[i] += mod;
  }
}
//convolution with arbitrary modulo using CRT
//slower but with no precision errors
void convolutionModCRT(vector<int> & P, vector<int> & Q, int
\rightarrow mod){
```

```
vector < int > A = P, B = P, C = P, D = Q, E = Q;
  int a = 998244353, b = 985661441, c = 754974721;
  p = a, root = 31, root_1 = 128805723, root_pw = 1 << 23;
  convolution(A, D);
  p = b, root = 210, root_1 = 934031556, root_pw = 1 << 22;
  convolution(B, E);
  p = c, root = 362, root_1 = 415027540, root_pw = 1 << 24;
  convolution(C, Q);
  P.resize(A.size());
 for(int i = 0; i < P.size(); ++i){
    int x1 = A[i] \% a;
    if(x1 < 0) x1 += a;
    int x2 = 65710754911 * (B[i] - x1) % b;
    if(x2 < 0) x2 += b;
    int x3 = (41653777411 * (C[i] - x1) % c - x2) * 41180439011

→ % c;

    if(x3 < 0) x3 += c;
   P[i] = x1 + a*(x2 + (11i)x3*b \% mod) \% mod;
    while(P[i] >= mod) P[i] -= mod;
    while(P[i] < 0) P[i] += mod;
 }
}
```

5. Geometría

5.1. Estructura point

```
ld eps = 1e-9, inf = numeric_limits<ld>::max();
bool geq(ld a, ld b){return a-b >= -eps;}
                                                //a >= b
bool leq(ld a, ld b){return b-a >= -eps;}
                                                //a \ll b
bool ge(ld a, ld b){return a-b > eps;}
                                                //a > b
bool le(ld a, ld b){return b-a > eps;}
                                                //a < b
bool eq(ld a, ld b){return abs(a-b) \leq eps;} //a == b
bool neq(ld a, ld b){return abs(a-b) > eps;} //a != b
struct point{
  ld x, y;
  point(): x(0), y(0){}
  point(ld x, ld y): x(x), y(y){}
  point operator+(const point & p) const{return point(x + p.x,
  \rightarrow y + p.y);}
  point operator-(const point & p) const{return point(x - p.x,
  \rightarrow y - p.y);}
  point operator*(const ld & k) const{return point(x * k, y *
  \rightarrow k);}
  point operator/(const ld & k) const{return point(x / k, y /
  \rightarrow k);}
  point operator+=(const point & p){*this = *this + p; return
  → *this;}
  point operator==(const point & p){*this = *this - p; return
  → *this;}
  point operator*=(const ld & p){*this = *this * p; return
  → *this;}
  point operator/=(const ld & p){*this = *this / p; return
  → *this;}
```

```
point rotate(const ld angle) const{
  return point(x * cos(angle) - y * sin(angle), x *

    sin(angle) + y * cos(angle));
point rotate(const ld angle, const point & p){
  return p + ((*this) - p).rotate(angle);
point perpendicular() const{
  return point(-y, x);
}
ld dot(const point & p) const{
  return x * p.x + y * p.y;
}
ld cross(const point & p) const{
  return x * p.y - y * p.x;
ld norm() const{
  return x * x + y * y;
long double length() const{
  return sqrtl(x * x + y * y);
}
point normalize() const{
  return (*this) / length();
}
point projection(const point & p) const{
  return (*this) * p.dot(*this) / dot(*this);
point normal(const point & p) const{
  return p - projection(p);
bool operator==(const point & p) const{
  return eq(x, p.x) && eq(y, p.y);
bool operator!=(const point & p) const{
  return !(*this == p);
```

}

```
}
  bool operator<(const point & p) const{</pre>
    if(eq(x, p.x)) return le(y, p.y);
    return le(x, p.x);
  }
  bool operator>(const point & p) const{
    if(eq(x, p.x)) return ge(y, p.y);
    return ge(x, p.x);
  }
};
istream & operator >> (istream & is, point & P){
  is >> P.x >> P.y;
  return is;
}
ostream & operator << (ostream & os, const point & p) {
  return os << "(" << p.x << ", " << p.y << ")";
}
int sgn(ld x){
  if(ge(x, 0)) return 1;
  if(le(x, 0)) return -1;
  return 0;
}
```

5.2. Líneas y segmentos

5.2.1. Verificar si un punto pertenece a una línea o segmento

5.2.2. Intersección de líneas

```
int intersectLinesInfo(const point & a1, const point & v1,
//line a1+tv1
 //line a2+tv2
 ld det = v1.cross(v2);
 if(eq(det, 0)){
   if(eq((a2 - a1).cross(v1), 0)){
     return -1; //infinity points
   }else{
     return 0; //no points
   }
 }else{
   return 1; //single point
 }
}
point intersectLines(const point & a1, const point & v1, const
→ point & a2, const point & v2){
 //lines a1+tv1, a2+tv2
 //assuming that they intersect
 ld det = v1.cross(v2);
 return a1 + v1 * ((a2 - a1).cross(v2) / det);
}
```

5.2.3. Intersección línea-segmento

5.2.4. Intersección de segmentos

```
int intersectSegmentsInfo(const point & a, const point & b,
\rightarrow const point & c, const point & d){
  //segment ab, segment cd
  point v1 = b - a, v2 = d - c;
  int t = sgn(v1.cross(c - a)), u = sgn(v1.cross(d - a));
  if(t == u){}
    if(t == 0){
      if(pointInSegment(a, b, c) || pointInSegment(a, b, d) ||
      → pointInSegment(c, d, a) || pointInSegment(c, d, b)){
        return -1; //infinity points
      }else{
        return 0; //no point
      }
    }else{
      return 0; //no point
    }
  }else{
    return sgn(v2.cross(a - c)) != sgn(v2.cross(b - c)); //1:

→ single point, 0: no point

  }
}
```

5.2.5. Distancia punto-recta

```
ld distancePointLine(const point & a, const point & v, const

or point & p){
   //line: a + tv, point p
   return abs(v.cross(p - a)) / v.length();
}
```

5.3. Círculos

5.3.1. Distancia punto-círculo

```
ld distancePointCircle(const point & p, const point & c, ld r){
   //point p, center c, radius r
   return max((ld)0, (p - c).length() - r);
}
```

5.3.2. Proyección punto exterior a círculo

5.3.3. Puntos de tangencia de punto exterior

```
pair<point, point> pointsOfTangency(const point & p, const

→ point & c, ld r){
    //point p (outside the circle), center c, radius r
    point v = (p - c).normalize() * r;
    ld theta = acos(r / (p - c).length());
    return {c + v.rotate(-theta), c + v.rotate(theta)};
}
```

5.3.4. Intersección línea-círculo

```
else if(D < 0) return {}; //no intersection
else{ //two points of intersection (chord)
    D = sqrt(D);
    ld t1 = (-B + D) / A;
    ld t2 = (-B - D) / A;
    return {a + v * t1, a + v * t2};
}</pre>
```

5.3.5. Centro y radio a través de tres puntos

5.3.6. Intersección de círculos

```
vector<point> intersectionCircles(const point & c1, ld r1,
\rightarrow const point & c2, ld r2){
  //circle 1 with center c1 and radius r1
  //circle 2 with center c2 and radius r2
  1d A = 2*r1*(c2.y - c1.y);
  1d B = 2*r1*(c2.x - c1.x);
  1d C = (c1 - c2).dot(c1 - c2) + r1*r1 - r2*r2;
  1d D = A*A + B*B - C*C;
  if(eq(D, 0)) return {c1 + point(B, A) * r1 / C};
  else if(le(D, 0)) return {};
  else{
    D = sqrt(D);
    1d cos1 = (B*C + A*D) / (A*A + B*B);
    1d \sin 1 = (A*C - B*D) / (A*A + B*B);
    1d cos2 = (B*C - A*D) / (A*A + B*B);
    1d \sin 2 = (A*C + B*D) / (A*A + B*B);
```

```
return {c1 + point(cos1, sin1) * r1, c1 + point(cos2, sin2)

→ * r1};
}
```

5.3.7. Contención de círculos

```
int circleInsideCircle(const point & c1, ld r1, const point &
\rightarrow c2, ld r2){
 //test if circle 2 is inside circle 1
 //returns "-1" if 2 touches internally 1, "1" if 2 is inside
  \rightarrow 1, "0" if they overlap
 1d 1 = r1 - r2 - (c1 - c2).length();
 return (ge(1, 0) ? 1 : (eq(1, 0) ? -1 : 0));
}
int circleOutsideCircle(const point & c1, ld r1, const point &
\rightarrow c2, ld r2){
 //test if circle 2 is outside circle 1
 //returns "-1" if they touch externally, "1" if 2 is outside
  \rightarrow 1, "0" if they overlap
 ld l = (c1 - c2).length() - (r1 + r2);
 return (ge(1, 0) ? 1 : (eq(1, 0) ? -1 : 0));
int pointInCircle(const point & c, ld r, const point & p){
 //test if point p is inside the circle with center c and
  \rightarrow radius r
 //returns "0" if it's outside, "-1" if it's in the perimeter,

→ "1" if it's inside

 ld l = (p - c).length() - r;
 return (le(1, 0) ? 1 : (eq(1, 0) ? -1 : 0));
5.3.8. Tangentes
```

vector<vector<point>> commonExteriorTangents(const point & c1,

//returns a vector of segments or a single point

 \rightarrow ld r1, const point & c2, ld r2){

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```
for(int i = 0; i < n; ++i){
  if(r1 < r2) return commonExteriorTangents(c2, r2, c1, r1);</pre>
  if(c1 == c2 \&\& abs(r1-r2) < 0) return {};
                                                                        ld si = (b - a).cross(S[i] - a);
  int in = circleInsideCircle(c1, r1, c2, r2);
                                                                        if(eq(si, 0)) continue;
  if(in == 1) return {};
                                                                        point m = getCircle(a, b, S[i]).first;
  else if(in == -1) return {\{c1 + (c2 - c1).normalize() * r1\}\};
                                                                        1d cr = (b - a).cross(m - a);
                                                                        if(le(si, 0)) hi = min(hi, cr);
    pair<point, point> t;
                                                                        else lo = max(lo, cr);
                                                                      }
    if(eq(r1, r2))
      t = \{c1 - (c2 - c1).perpendicular(), c1 + (c2 - c2)\}
                                                                      ld v = (ge(lo, 0) ? lo : le(hi, 0) ? hi : 0);
                                                                      point c = (a + b) / 2 + (b - a).perpendicular() * v / (b - a)
      else
                                                                      \rightarrow a).norm();
      t = pointsOfTangency(c2, c1, r1 - r2);
                                                                      return {c, (a - c).norm()};
    t.first = (t.first - c1).normalize();
                                                                    }
    t.second = (t.second - c1).normalize();
    return {{c1 + t.first * r1, c2 + t.first * r2}, {c1 +
                                                                    pair<point, ld> mec(vector<point> & S, const point & a, int n){
    \rightarrow t.second * r1, c2 + t.second * r2}};
                                                                      random_shuffle(S.begin(), S.begin() + n);
  }
                                                                      point b = S[0], c = (a + b) / 2;
}
                                                                      ld r = (a - c).norm();
                                                                      for(int i = 1; i < n; ++i){
                                                                        if(ge((S[i] - c).norm(), r)){
vector<vector<point>> commonInteriorTangents(const point & c1,
\rightarrow ld r1, const point & c2, ld r2){
                                                                          tie(c, r) = (n == S.size() ? mec(S, S[i], i) : mec2(S, a, a)
  if (c1 == c2 \&\& abs(r1-r2) < 0) return {};
                                                                           \hookrightarrow S[i], i));
  int out = circleOutsideCircle(c1, r1, c2, r2);
                                                                        }
  if(out == 0) return {};
                                                                      }
  else if(out == -1) return {{c1 + (c2 - c1).normalize() *
                                                                      return {c, r};
                                                                    }
  \hookrightarrow r1}};
  else{
    auto t = pointsOfTangency(c2, c1, r1 + r2);
                                                                    pair<point, ld> smallestEnclosingCircle(vector<point> S){
    t.first = (t.first - c1).normalize();
                                                                      assert(!S.empty());
    t.second = (t.second - c1).normalize();
                                                                      auto r = mec(S, S[0], S.size());
    return {{c1 + t.first * r1, c2 - t.first * r2}, {c1 +
                                                                      return {r.first, sqrt(r.second)};
    \rightarrow t.second * r1, c2 - t.second * r2}};
                                                                    }
  }
}
                                                                    5.4. Polígonos
5.3.9. Smallest enclosing circle
                                                                    5.4.1. Perímetro y área de un polígono
pair<point, ld> mec2(vector<point> & S, const point & a, const
                                                                    ld perimeter(vector<point> & P){
→ point & b, int n){
                                                                      int n = P.size();
  ld hi = inf, lo = -hi;
                                                                      ld ans = 0;
```

```
for(int i = 0; i < n; i++){
    ans += (P[i] - P[(i + 1) % n]).length();
}
    return ans;
}

ld area(vector<point> & P){
    int n = P.size();
    ld ans = 0;
    for(int i = 0; i < n; i++){
        ans += P[i].cross(P[(i + 1) % n]);
    }
    return abs(ans / 2);
}</pre>
```

5.4.2. Envolvente convexa (convex hull) de un polígono

```
vector<point> convexHull(vector<point> P){
  sort(P.begin(), P.end());
  vector<point> L, U;
  for(int i = 0; i < P.size(); i++){</pre>
    while(L.size() \geq 2 && leq((L[L.size() - 2] -
    → P[i]).cross(L[L.size() - 1] - P[i]), 0)){
      L.pop_back();
    }
    L.push_back(P[i]);
  for(int i = P.size() - 1; i >= 0; i--){
    while(U.size() \geq 2 && leg((U[U.size() - 2] -
    \rightarrow P[i]).cross(U[U.size() - 1] - P[i]), 0)){
      U.pop_back();
    U.push_back(P[i]);
  }
  L.pop_back();
  U.pop_back();
  L.insert(L.end(), U.begin(), U.end());
  return L;
}
```

5.4.3. Verificar si un punto pertenece al perímetro de un polígono

```
bool pointInPerimeter(vector<point> & P, const point & p){
  int n = P.size();
  for(int i = 0; i < n; i++){
    if(pointInSegment(P[i], P[(i + 1) % n], p)){
      return true;
    }
  }
  return false;
}</pre>
```

5.4.4. Verificar si un punto pertenece a un polígono

5.4.5. Verificar si un punto pertenece a un polígono convexo $O(\log n)$

```
//point in convex polygon in log(n)
//first do preprocess: seg=process(P),
//then for each query call pointInConvexPolygon(seg, p - P[0])
vector<point> process(vector<point> & P){
  int n = P.size();
  rotate(P.begin(), min_element(P.begin(), P.end()), P.end());
```

```
vector<point> seg(n - 1);
                                                                      return false;
  for(int i = 0; i < n - 1; ++i)
    seg[i] = P[i + 1] - P[0];
  return seg;
                                                                    vector<vector<point>> cutPolygon(vector<point> & P, const point
}
                                                                    \rightarrow & a, const point & v){
                                                                      //line a+tv, polygon P
bool pointInConvexPolygon(vector<point> & seg, const point &
                                                                      int n = P.size();
                                                                      if(!lineCutsPolygon(P, a, v)) return {P};
→ p){
  int n = seg.size();
                                                                      int idx = 0:
  if(neq(seg[0].cross(p), 0) \&\& sgn(seg[0].cross(p)) !=
                                                                      vector<vector<point>> ans(2);
                                                                      for(int i = 0; i < n; ++i){
  \rightarrow sgn(seg[0].cross(seg[n - 1])))
    return false;
                                                                        if(intersectLineSegmentInfo(a, v, P[i], P[(i+1)%n])){
                                                                          point p = intersectLines(a, v, P[i], P[(i+1)%n] - P[i]);
  if(neq(seg[n-1].cross(p), 0) \&\& sgn(seg[n-1].cross(p)) !=
  \rightarrow sgn(seg[n - 1].cross(seg[0])))
                                                                          if(P[i] == p) continue;
    return false;
                                                                          ans[idx].push_back(P[i]);
  if(eq(seg[0].cross(p), 0))
                                                                          ans[1-idx].push_back(p);
    return geq(seg[0].length(), p.length());
                                                                          ans[idx].push_back(p);
  int 1 = 0, r = n - 1;
                                                                          idx = 1-idx;
  while (r - 1 > 1) {
                                                                        }else{
    int m = 1 + ((r - 1) >> 1);
                                                                          ans[idx].push_back(P[i]);
    if(geq(seg[m].cross(p), 0)) 1 = m;
                                                                        }
                                                                      }
    else r = m;
  }
                                                                      return ans;
  return eq(abs(seg[1].cross(seg[1 + 1])), abs((p -
                                                                    }
  \rightarrow seg[1]).cross(p - seg[1 + 1])) + abs(p.cross(seg[1])) +
     abs(p.cross(seg[1 + 1])));
                                                                    5.4.7. Centroide de un polígono
}
                                                                    point centroid(vector<point> & P){
5.4.6. Cortar un polígono con una recta
                                                                      point num;
                                                                      1d den = 0;
bool lineCutsPolygon(vector<point> & P, const point & a, const
                                                                      int n = P.size();
\rightarrow point & v){
                                                                      for(int i = 0; i < n; ++i){
                                                                        ld cross = P[i].cross(P[(i + 1) \% n]);
  //line a+tv, polygon P
  int n = P.size();
                                                                        num += (P[i] + P[(i + 1) \% n]) * cross;
  for(int i = 0, first = 0; i \le n; ++i){
                                                                        den += cross;
    int side = sgn(v.cross(P[i\%n]-a));
                                                                      }
    if(!side) continue;
                                                                      return num / (3 * den);
    if(!first) first = side;
                                                                    }
    else if(side != first) return true;
```

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}

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5.4.8. Pares de puntos antipodales

5.4.9. Diámetro y ancho

```
pair<ld, ld> diameterAndWidth(vector<point> & P){
  int n = P.size(), k = 0;
  auto dot = [&](int a, int b){return
  \rightarrow (P[(a+1)\%n]-P[a]).dot(P[(b+1)\%n]-P[b]);};
  auto cross = [&](int a, int b){return
  \rightarrow (P[(a+1)\%n]-P[a]).cross(P[(b+1)\%n]-P[b]);};
  ld diameter = 0:
  ld width = inf:
  while (ge(dot(0, k), 0)) k = (k+1) \% n;
  for(int i = 0; i < n; ++i){
    while (ge(cross(i, k), 0)) k = (k+1) \% n;
    //pair: (i, k)
    diameter = max(diameter, (P[k] - P[i]).length());
    width = min(width, distancePointLine(P[i], P[(i+1)%n] -
    \rightarrow P[i], P[k]));
  }
  return make_pair(diameter, width);
```

5.4.10. Smallest enclosing rectangle

```
pair<1d, 1d> smallestEnclosingRectangle(vector<point> & P){
  int n = P.size();
  auto dot = [&](int a, int b){return
  \rightarrow (P[(a+1)\%n]-P[a]).dot(P[(b+1)\%n]-P[b]);};
  auto cross = [&](int a, int b){return
  \rightarrow (P[(a+1)\%n]-P[a]).cross(P[(b+1)\%n]-P[b]);};
  ld perimeter = inf, area = inf;
 for(int i = 0, j = 0, k = 0, m = 0; i < n; ++i){
    while (ge(dot(i, j), 0)) j = (j+1) \% n;
    if(!i) k = j;
    while (ge(cross(i, k), 0)) k = (k+1) \% n;
    if(!i) m = k;
    while(le(dot(i, m), 0)) m = (m+1) \% n;
    //pairs: (i, k), (j, m)
    point v = P[(i+1)\%n] - P[i];
   ld h = distancePointLine(P[i], v, P[k]);
    ld w = distancePointLine(P[j], v.perpendicular(), P[m]);
    perimeter = min(perimeter, 2 * (h + w));
    area = min(area, h * w);
 return make_pair(area, perimeter);
```

5.5. Par de puntos más cercanos

```
bool comp1(const point & a, const point & b){
   return a.y < b.y;
}

pair<point, point> closestPairOfPoints(vector<point> P){
   sort(P.begin(), P.end(), comp1);
   set<point> S;
   ld ans = inf;
   point p, q;
   int pos = 0;
   for(int i = 0; i < P.size(); ++i){
      while(pos < i && abs(P[i].y - P[pos].y) >= ans){
        S.erase(P[pos++]);
    }
}
```

```
auto lower = S.lower_bound({P[i].x - ans - eps, -inf});
                                                                         return new node({ p });
    auto upper = S.upper_bound({P[i].x + ans + eps, -inf});
                                                                       for(int i = 1; i < r; ++i)
    for(auto it = lower; it != upper; ++it){
                                                                         aux[i].first = (p - aux[i].second).dot(p -
      1d d = (P[i] - *it).length();
                                                                         → aux[i].second);
      if(d < ans)
                                                                       int m = (1 + r) / 2;
                                                                       nth_element(aux.begin() + 1, aux.begin() + m, aux.begin() +
        ans = d;
        p = P[i];
        q = *it;
                                                                       return new node({ p, sqrt(aux[m].first), build(1, m),
      }
                                                                       \rightarrow build(m, r) });
                                                                     }
    S.insert(P[i]);
  }
                                                                     priority_queue<pair<ld, node*>> que;
  return {p, q};
}
                                                                     void k_nn(node *t, point p, int k){
                                                                       if(!t)
                                                                         return;
      Vantage Point Tree (puntos más cercanos a cada
                                                                       ld d = (p - t->p).length();
                                                                       if(que.size() < k)</pre>
      punto)
                                                                         que.push({ d, t });
                                                                       else if(ge(que.top().first, d)){
struct vantage_point_tree{
                                                                         que.pop();
  struct node
                                                                         que.push({ d, t });
    point p;
                                                                       if(!t->1 && !t->r)
   ld th;
                                                                         return;
    node *1, *r;
                                                                       if(le(d, t->th)){
  }*root;
                                                                         k_n(t->1, p, k);
                                                                         if(leq(t->th - d, que.top().first))
  vector<pair<ld, point>> aux;
                                                                           k_nn(t->r, p, k);
                                                                       }else{
  vantage_point_tree(vector<point> &ps){
                                                                         k_nn(t->r, p, k);
    for(int i = 0; i < ps.size(); ++i)</pre>
                                                                         if(leq(d - t->th, que.top().first))
      aux.push_back({ 0, ps[i] });
                                                                           k_nn(t->1, p, k);
    root = build(0, ps.size());
                                                                       }
  }
                                                                     }
  node *build(int 1, int r){
                                                                     vector<point> k_nn(point p, int k){
    if(1 == r)
                                                                       k_nn(root, p, k);
      return 0;
                                                                       vector<point> ans;
    swap(aux[1], aux[1 + rand() \% (r - 1)]);
    point p = aux[1++].second;
                                                                       for(; !que.empty(); que.pop())
```

ans.push_back(que.top().second->p);

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if(1 == r)

```
reverse(ans.begin(), ans.end());
                                                                      QuadEdge* onext = nullptr;
                                                                      bool used = false;
    return ans;
  }
                                                                      QuadEdge* rev() const{return rot->rot;}
};
                                                                      QuadEdge* lnext() const{return rot->rev()->onext->rot;}
                                                                      QuadEdge* oprev() const{return rot->onext->rot;}
                                                                     point dest() const{return rev()->origin;}
      Suma Minkowski
                                                                    }:
vector<point> minkowskiSum(vector<point> A, vector<point> B){
                                                                    QuadEdge* make_edge(const point & from, const point & to){
  int na = (int)A.size(), nb = (int)B.size();
                                                                      QuadEdge* e1 = new QuadEdge;
  if(A.empty() || B.empty()) return {};
                                                                      QuadEdge* e2 = new QuadEdge;
                                                                      QuadEdge* e3 = new QuadEdge;
  rotate(A.begin(), min_element(A.begin(), A.end()), A.end());
                                                                      QuadEdge* e4 = new QuadEdge;
  rotate(B.begin(), min_element(B.begin(), B.end()), B.end());
                                                                      e1->origin = from;
                                                                      e2->origin = to;
  int pa = 0, pb = 0;
                                                                      e3->origin = e4->origin = inf_pt;
  vector<point> M;
                                                                      e1->rot = e3;
                                                                      e2->rot = e4;
  while(pa < na \&\& pb < nb){
                                                                      e3->rot = e2;
    M.push_back(A[pa] + B[pb]);
                                                                      e4->rot = e1;
    ld x = (A[(pa + 1) \% na] - A[pa]).cross(B[(pb + 1) \% nb] -
                                                                      e1->onext = e1;
    \rightarrow B[pb]);
                                                                      e2->onext = e2;
    if(leq(x, 0)) pb++;
                                                                      e3->onext = e4:
    if(geq(x, 0)) pa++;
                                                                      e4->onext = e3;
  }
                                                                     return e1;
                                                                    }
  while(pa < na) M.push_back(A[pa++] + B[0]);</pre>
  while(pb < nb) M.push_back(B[pb++] + A[0]);</pre>
                                                                    void splice(QuadEdge* a, QuadEdge* b){
                                                                      swap(a->onext->rot->onext, b->onext->rot->onext);
  return M;
                                                                      swap(a->onext, b->onext);
}
                                                                    void delete_edge(QuadEdge* e){
      Triangulación de Delaunay
                                                                      splice(e, e->oprev());
                                                                      splice(e->rev(), e->rev()->oprev());
//Delaunay triangulation in O(n \log n)
                                                                      delete e->rot;
const point inf_pt(inf, inf);
                                                                     delete e->rev()->rot;
                                                                     delete e;
struct QuadEdge{
                                                                     delete e->rev();
  point origin;
  QuadEdge* rot = nullptr;
```

```
QuadEdge* connect(QuadEdge* a, QuadEdge* b){
                                                                        QuadEdge *a = make_edge(P[1], P[1 + 1]), *b = make_edge(P[1])
  QuadEdge* e = make_edge(a->dest(), b->origin);
                                                                        \rightarrow + 1], P[r]);
  splice(e, a->lnext());
                                                                        splice(a->rev(), b);
  splice(e->rev(), b);
                                                                        int sg = sgn((P[1 + 1] - P[1]).cross(P[r] - P[1]));
  return e;
                                                                        if(sg == 0)
}
                                                                         return make_pair(a, b->rev());
                                                                        QuadEdge* c = connect(b, a);
                                                                        if(sg == 1)
bool left_of(const point & p, QuadEdge* e){
  return ge((e->origin - p).cross(e->dest() - p), 0);
                                                                          return make_pair(a, b->rev());
}
                                                                          return make_pair(c->rev(), c);
bool right_of(const point & p, QuadEdge* e){
  return le((e->origin - p).cross(e->dest() - p), 0);
                                                                      int mid = (1 + r) / 2;
}
                                                                      QuadEdge *ldo, *ldi, *rdo, *rdi;
                                                                      tie(ldo, ldi) = build_tr(l, mid, P);
ld det3(ld a1, ld a2, ld a3, ld b1, ld b2, ld b3, ld c1, ld c2,
                                                                      tie(rdi, rdo) = build_tr(mid + 1, r, P);
\rightarrow 1d c3) {
                                                                      while(true){
  return a1 * (b2 * c3 - c2 * b3) - a2 * (b1 * c3 - c1 * b3) +
                                                                        if(left_of(rdi->origin, ldi)){
  \rightarrow a3 * (b1 * c2 - c1 * b2);
                                                                         ldi = ldi->lnext();
}
                                                                          continue;
bool in_circle(const point & a, const point & b, const point &
                                                                        if(right_of(ldi->origin, rdi)){
rdi = rdi->rev()->onext;
  1d det = -det3(b.x, b.y, b.norm(), c.x, c.y, c.norm(), d.x,
                                                                          continue;
  \rightarrow d.y, d.norm());
  det += det3(a.x, a.y, a.norm(), c.x, c.y, c.norm(), d.x, d.y,
                                                                        break;
  \rightarrow d.norm()):
  det -= det3(a.x, a.y, a.norm(), b.x, b.y, b.norm(), d.x, d.y,
                                                                      QuadEdge* basel = connect(rdi->rev(), ldi);
                                                                      auto valid = [&basel](QuadEdge* e){return right_of(e->dest(),
  \rightarrow d.norm());
  det += det3(a.x, a.y, a.norm(), b.x, b.y, b.norm(), c.x, c.y,
                                                                      → basel);};
  \rightarrow c.norm());
                                                                      if(ldi->origin == ldo->origin)
  return ge(det, 0);
                                                                        ldo = basel->rev();
}
                                                                      if(rdi->origin == rdo->origin)
                                                                        rdo = basel;
pair<QuadEdge*, QuadEdge*> build_tr(int 1, int r, vector<point>
                                                                      while(true){
                                                                        QuadEdge* lcand = basel->rev()->onext;

→ & P) {

  if(r - 1 + 1 == 2){
                                                                        if(valid(lcand)){
    QuadEdge* res = make_edge(P[1], P[r]);
                                                                          while(in_circle(basel->dest(), basel->origin,
    return make_pair(res, res->rev());
                                                                          → lcand->dest(), lcand->onext->dest())){
                                                                            QuadEdge* t = lcand->onext;
  if(r - 1 + 1 == 3){
                                                                            delete_edge(lcand);
```

```
add();
        lcand = t;
      }
                                                                      P.clear();
    }
                                                                      int kek = 0;
    QuadEdge* rcand = basel->oprev();
                                                                      while(kek < (int)edges.size())</pre>
    if(valid(rcand)){
                                                                        if(!(e = edges[kek++])->used)
      while(in_circle(basel->dest(), basel->origin,
                                                                          add();

→ rcand->dest(), rcand->oprev()->dest())){
                                                                      vector<tuple<point, point, point>> ans;
        QuadEdge* t = rcand->oprev();
                                                                      for(int i = 0; i < (int)P.size(); i += 3){</pre>
        delete_edge(rcand);
                                                                        ans.push_back(make_tuple(P[i], P[i + 1], P[i + 2]));
        rcand = t;
      }
                                                                      return ans;
    }
    if(!valid(lcand) && !valid(rcand))
      break;
    if(!valid(lcand) || (valid(rcand) &&
    → in_circle(lcand->dest(), lcand->origin, rcand->origin,

→ rcand->dest())))
      basel = connect(rcand, basel->rev());
    else
      basel = connect(basel->rev(), lcand->rev());
  }
  return make_pair(ldo, rdo);
}
vector<tuple<point, point, point>> delaunay(vector<point> & P){
  sort(P.begin(), P.end());
  auto res = build_tr(0, (int)P.size() - 1, P);
  QuadEdge* e = res.first;
  vector<QuadEdge*> edges = {e};
  while(le((e->dest() - e->onext->dest()).cross(e->origin -

    e→ e->onext->dest()), 0))

    e = e->onext;
  auto add = [\&P, \&e, \&edges](){
    QuadEdge* curr = e;
    do{
      curr->used = true;
      P.push_back(curr->origin);
      edges.push_back(curr->rev());
      curr = curr->lnext();
    }while(curr != e);
  };
```

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6. Grafos

6.1. Disjoint Set

```
struct disjointSet{
  int N;
  vector<short int> rank;
  vi parent, count;
  disjointSet(int N): N(N), parent(N), count(N), rank(N){}
  void makeSet(int v){
    count[v] = 1;
    parent[v] = v;
  }
  int findSet(int v){
    if(v == parent[v]) return v;
    return parent[v] = findSet(parent[v]);
  }
  void unionSet(int a, int b){
    a = findSet(a), b = findSet(b);
    if(a == b) return:
    if(rank[a] < rank[b]){</pre>
      parent[a] = b;
      count[b] += count[a];
    }else{
      parent[b] = a;
      count[a] += count[b];
      if(rank[a] == rank[b]) ++rank[a];
    }
  }
};
```

6.2. Definiciones

```
struct edge{
  int source, dest, cost;
```

```
edge(): source(0), dest(0), cost(0){}
  edge(int dest, int cost): dest(dest), cost(cost){}
  edge(int source, int dest, int cost): source(source),

→ dest(dest), cost(cost){}
  bool operator==(const edge & b) const{
    return source == b.source && dest == b.dest && cost ==
    → b.cost;
 }
  bool operator<(const edge & b) const{</pre>
    return cost < b.cost;
  bool operator>(const edge & b) const{
    return cost > b.cost;
 }
};
struct path{
  int cost = inf;
  deque<int> vertices;
 int size = 1;
 int prev = -1;
};
struct graph{
  vector<vector<edge>> adjList;
  vector<vb> adjMatrix;
  vector<vi> costMatrix;
  vector<edge> edges;
  int V = 0;
  bool dir = false;
  graph(int n, bool dir): V(n), dir(dir), adjList(n), edges(n),
  → adjMatrix(n, vb(n)), costMatrix(n, vi(n)){
   for(int i = 0; i < n; ++i)
      for(int j = 0; j < n; ++ j)
       costMatrix[i][j] = (i == j ? 0 : inf);
 }
```

```
void add(int source, int dest, int cost){
  adjList[source].emplace_back(source, dest, cost);
  edges.emplace_back(source, dest, cost);
  adjMatrix[source][dest] = true;
  costMatrix[source][dest] = cost;
  if(!dir){
    adjList[dest].emplace_back(dest, source, cost);
    adjMatrix[dest] [source] = true;
    costMatrix[dest] [source] = cost;
  }
}
void buildPaths(vector<path> & paths){
  for(int i = 0; i < V; i++){
    int u = i;
    for(int j = 0; j < paths[i].size; <math>j++){
      paths[i].vertices.push_front(u);
      u = paths[u].prev;
    }
  }
}
```

6.3. DFS genérica

```
void dfs(int u, vi & status, vi & parent){
  status[u] = 1;
  for(edge & current : adjList[u]){
    int v = current.dest;
    if(status[v] == 0){ //not visited
       parent[v] = u;
       dfs(v, status, parent);
  }else if(status[v] == 1){ //explored
    if(v == parent[u]){
       //bidirectional node u<-->v
    }else{
       //back edge u-v
    }
} else if(status[v] == 2){ //visited
       //forward edge u-v
}
```

```
}
status[u] = 2;
}
```

6.4. Dijkstra

```
vector<path> dijkstra(int start){
 priority_queue<edge, vector<edge>, greater<edge>> cola;
  vector<path> paths(V);
  cola.emplace(start, 0);
  paths[start].cost = 0;
  while(!cola.empty()){
    int u = cola.top().dest; cola.pop();
    for(edge & current : adjList[u]){
      int v = current.dest;
      int nuevo = paths[u].cost + current.cost;
      if(nuevo == paths[v].cost && paths[u].size + 1 <</pre>
      → paths[v].size){
        paths[v].prev = u;
        paths[v].size = paths[u].size + 1;
      }else if(nuevo < paths[v].cost){</pre>
        paths[v].prev = u;
        paths[v].size = paths[u].size + 1;
        cola.emplace(v, nuevo);
        paths[v].cost = nuevo;
     }
   }
 buildPaths(paths);
 return paths;
```

6.5. Bellman Ford

```
vector<path> bellmanFord(int start){
  vector<path> paths(V, path());
  vi processed(V);
  vb inQueue(V);
  queue<int> Q;
```

```
paths[start].cost = 0;
                                                                               tmp[i][j] = min(tmp[i][j], tmp[i][k] + tmp[k][j]);
    Q.push(start);
                                                                       return tmp;
    while(!Q.empty()){
      int u = Q.front(); Q.pop(); inQueue[u] = false;
      if(paths[u].cost == inf) continue;
                                                                   6.7. Cerradura transitiva O(V^3)
      ++processed[u];
      if(processed[u] == V){
        cout << "Negative cycle\n";</pre>
                                                                     vector<vb> transitiveClosure(){
                                                                       vector<vb> tmp = adjMatrix;
        return {};
                                                                       for(int k = 0; k < V; ++k)
                                                                         for(int i = 0; i < V; ++i)
      for(edge & current : adjList[u]){
        int v = current.dest;
                                                                           for(int j = 0; j < V; ++j)
                                                                             tmp[i][j] = tmp[i][j] || (tmp[i][k] && tmp[k][j]);
        int nuevo = paths[u].cost + current.cost;
        if(nuevo == paths[v].cost && paths[u].size + 1 <</pre>
                                                                       return tmp;
        → paths[v].size){
                                                                     }
          paths[v].prev = u;
          paths[v].size = paths[u].size + 1;
                                                                   6.8. Cerradura transitiva O(V^2)
        }else if(nuevo < paths[v].cost){</pre>
          if(!inQueue[v]){
                                                                     vector<vb> transitiveClosureDFS(){
            Q.push(v);
                                                                       vector<vb> tmp(V, vb(V));
            inQueue[v] = true;
                                                                       function<void(int, int)> dfs = [&](int start, int u){
                                                                         for(edge & current : adjList[u]){
          paths[v].prev = u;
                                                                           int v = current.dest;
          paths[v].size = paths[u].size + 1;
                                                                           if(!tmp[start][v]){
          paths[v].cost = nuevo;
        }
                                                                             tmp[start][v] = true;
                                                                             dfs(start, v);
      }
                                                                           }
                                                                         }
    buildPaths(paths);
    return paths;
                                                                       };
                                                                       for(int u = 0; u < V; u++)
                                                                         dfs(u, u);
                                                                       return tmp;
6.6. Floyd
  vector<vi> floyd(){
                                                                   6.9. Verificar si el grafo es bipartito
    vector<vi> tmp = costMatrix;
    for(int k = 0; k < V; ++k)
                                                                     bool isBipartite(){
      for(int i = 0; i < V; ++i)
                                                                       vi side(V, -1);
        for(int j = 0; j < V; ++j)
          if(tmp[i][k] != inf && tmp[k][j] != inf)
                                                                       queue<int> q;
```

```
for (int st = 0; st < V; ++st){
    if(side[st] != -1) continue;
    q.push(st);
    side[st] = 0;
    while(!q.empty()){
     int u = q.front();
      q.pop();
      for (edge & current : adjList[u]){
        int v = current.dest;
        if(side[v] == -1) {
          side[v] = side[u] ^ 1;
         q.push(v);
        }else{
          if(side[v] == side[u]) return false;
        }
      }
   }
  }
 return true;
}
```

6.10. Orden topológico

```
vi topologicalSort(){
  int visited = 0;
 vi order, indegree(V);
 for(auto & node : adjList){
   for(edge & current : node){
      int v = current.dest;
      ++indegree[v];
   }
 }
  queue<int> Q;
 for(int i = 0; i < V; ++i){
    if(indegree[i] == 0) Q.push(i);
 while(!Q.empty()){
   int source = Q.front();
   Q.pop();
    order.push_back(source);
```

```
++visited;
for(edge & current : adjList[source]){
   int v = current.dest;
   --indegree[v];
   if(indegree[v] == 0) Q.push(v);
  }
}
if(visited == V) return order;
else return {};
```

6.11. Detectar ciclos

```
bool hasCycle(){
  vi color(V);
  function <bool(int, int) > dfs = [&](int u, int parent){
    color[u] = 1;
   bool ans = false;
   int ret = 0;
   for(edge & current : adjList[u]){
      int v = current.dest;
      if(color[v] == 0)
        ans = dfs(v, u);
      else if(color[v] == 1 && (dir || v != parent || ret++))
        ans = true;
   }
    color[u] = 2;
    return ans;
 };
  for(int u = 0; u < V; ++u)
    if(color[u] == 0 \&\& dfs(u, -1))
      return true;
  return false;
}
```

6.12. Puentes y puntos de articulación

```
pair<vb, vector<edge>> articulationBridges(){
  vi low(V), label(V);
```

}

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```
vb points(V);
  vector<edge> bridges;
  int time = 0;
  function<int(int, int)> dfs = [&](int u, int p){
    label[u] = low[u] = ++time;
    int hijos = 0, ret = 0;
    for(edge & current : adjList[u]){
      int v = current.dest;
      if(v == p && !ret++) continue;
      if(!label[v]){
        ++hijos;
        dfs(v, u);
        if(label[u] <= low[v])</pre>
          points[u] = true;
        if(label[u] < low[v])</pre>
          bridges.push_back(current);
        low[u] = min(low[u], low[v]);
      low[u] = min(low[u], label[v]);
    }
    return hijos;
  };
  for(int u = 0; u < V; ++u)
    if(!label[u])
      points[u] = dfs(u, -1) > 1;
  return make_pair(points, bridges);
}
```

6.13. Componentes fuertemente conexas

```
vector<vi> scc(){
  vi low(V), label(V);
  int time = 0;
  vector<vi> ans;
  stack<int> S;
  function<void(int)> dfs = [&](int u){
    label[u] = low[u] = ++time;
    S.push(u);
    for(edge & current : adjList[u]){
        int v = current.dest;
    }
}
```

```
if(!label[v]) dfs(v);
   low[u] = min(low[u], low[v]);
  if(label[u] == low[u]){
    vi comp;
    while(S.top() != u){
      comp.push_back(S.top());
      low[S.top()] = V + 1;
      S.pop();
    comp.push_back(S.top());
   S.pop();
    ans.push_back(comp);
   low[u] = V + 1;
 }
};
for(int u = 0; u < V; ++u)
  if(!label[u]) dfs(u);
return ans;
```

6.14. Árbol mínimo de expansión (Kruskal)

```
vector<edge> kruskal(){
  sort(edges.begin(), edges.end());
 vector<edge> MST;
 disjointSet DS(V);
 for(int u = 0; u < V; ++u)
   DS.makeSet(u);
 int i = 0:
  while(i < edges.size() && MST.size() < V - 1){</pre>
    edge current = edges[i++];
   int u = current.source, v = current.dest;
   if(DS.findSet(u) != DS.findSet(v)){
     MST.push_back(current);
     DS.unionSet(u, v);
   }
 }
 return MST;
```

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6.15. Máximo emparejamiento bipartito

```
bool tryKuhn(int u, vb & used, vi & left, vi & right){
  if(used[u]) return false;
  used[u] = true;
  for(edge & current : adjList[u]){
    int v = current.dest;
    if(right[v] == -1 || tryKuhn(right[v], used, left,

    right)){
     right[v] = u;
     left[u] = v;
      return true;
    }
  }
  return false;
}
bool augmentingPath(int u, vb & used, vi & left, vi & right){
  used[u] = true;
  for(edge & current : adjList[u]){
    int v = current.dest;
    if(right[v] == -1){
     right[v] = u;
     left[u] = v;
      return true;
    }
  }
  for(edge & current : adjList[u]){
    int v = current.dest;
    if(!used[right[v]] && augmentingPath(right[v], used,
    → left, right)){
     right[v] = u;
     left[u] = v;
      return true;
    }
  }
  return false;
}
//vertices from the left side numbered from 0 to l-1
//vertices from the right side numbered from 0 to r-1
```

```
//graph[u] represents the left side
//qraph[u][v] represents the right side
//we can use tryKuhn() or augmentingPath()
vector<pair<int, int>> maxMatching(int 1, int r){
  vi left(1, -1), right(r, -1);
  vb used(1);
 for(int u = 0; u < 1; ++u){
    tryKuhn(u, used, left, right);
   fill(used.begin(), used.end(), false);
  vector<pair<int, int>> ans;
 for(int u = 0; u < r; ++u){
   if(right[u] != -1){
      ans.emplace_back(right[u], u);
   }
 }
  return ans;
```

6.16. Circuito euleriano

7. Árboles

7.1. Estructura tree

```
struct tree{
  vi parent, level, weight;
 vector<vi> dists, DP;
  int n, root;
  void dfs(int u, graph & G){
    for(edge & curr : G.adjList[u]){
      int v = curr.dest;
      int w = curr.cost;
      if(v != parent[u]){
        parent[v] = u;
        weight[v] = w;
        level[v] = level[u] + 1;
        dfs(v, G);
      }
   }
 }
  tree(int n, int root): n(n), root(root), parent(n), level(n),
  \rightarrow weight(n), dists(n, vi(20)), DP(n, vi(20)){
    parent[root] = root;
  tree(graph & G, int root): n(G.V), root(root), parent(G.V),
  \rightarrow level(G.V), weight(G.V), dists(G.V, vi(20)), DP(G.V,
  \rightarrow vi(20)){
    parent[root] = root;
    dfs(root, G);
 }
  void pre(){
    for(int u = 0; u < n; u++){
      DP[u][0] = parent[u];
      dists[u][0] = weight[u];
    for(int i = 1; (1 << i) <= n; ++i){
```

7.2. k-ésimo ancestro

```
int ancestor(int p, int k){
  int h = level[p] - k;
  if(h < 0) return -1;
  int lg;
  for(lg = 1; (1 << lg) <= level[p]; ++lg);
  lg--;
  for(int i = lg; i >= 0; --i){
    if(level[p] - (1 << i) >= h){
      p = DP[p][i];
    }
  }
  return p;
}
```

7.3. LCA

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```
int lca(int p, int q){
  if(level[p] < level[q]) swap(p, q);
  int lg;
  for(lg = 1; (1 << lg) <= level[p]; ++lg);
  lg--;
  for(int i = lg; i >= 0; --i){
    if(level[p] - (1 << i) >= level[q]){
      p = DP[p][i];
    }
  }
  if(p == q) return p;
  for(int i = lg; i >= 0; --i){
```

```
if(DP[p][i] != -1 && DP[p][i] != DP[q][i]){
    p = DP[p][i];
    q = DP[q][i];
}
return parent[p];
}
```

7.4. Distancia entre dos nodos

```
int dist(int p, int q){
  if(level[p] < level[q]) swap(p, q);</pre>
 int lg;
 for(lg = 1; (1 << lg) <= level[p]; ++lg);
 lg--;
 int sum = 0;
 for(int i = lg; i >= 0; --i){
    if(level[p] - (1 \ll i) >= level[q]){
      sum += dists[p][i];
      p = DP[p][i];
    }
  if(p == q) return sum;
 for(int i = lg; i >= 0; --i){
    if(DP[p][i] != -1 \&\& DP[p][i] != DP[q][i]){
      sum += dists[p][i] + dists[q][i];
      p = DP[p][i];
      q = DP[q][i];
    }
  sum += dists[p][0] + dists[q][0];
  return sum;
}
```

7.5. HLD

7.6. Link Cut

8. Flujos

8.1. Estructura flowEdge

8.2. Estructura flowGraph

```
template<typename T>
struct flowGraph{
 T inf = numeric_limits<T>::max();
 vector<vector<flowEdge<T>*>> adjList;
 vector<int> dist, pos;
 int V;
 flowGraph(int V): V(V), adjList(V), dist(V), pos(V){}
  ~flowGraph(){
   for(int i = 0; i < V; ++i)
     for(int j = 0; j < adjList[i].size(); ++j)</pre>
        delete adjList[i][j];
 }
  void addEdge(int u, int v, T capacity, T cost = 0){
   flowEdge<T> *uv = new flowEdge<T>(v, 0, capacity, cost);
   flowEdge<T> *vu = new flowEdge<T>(u, capacity, capacity,
    \rightarrow -cost);
```

```
uv->res = vu;
vu->res = uv;
adjList[u].push_back(uv);
adjList[v].push_back(vu);
}
```

8.3. Algoritmo de Edmonds-Karp $O(VE^2)$

```
//Maximun Flow using Edmonds-Karp Algorithm O(VE^2)
T edmondsKarp(int s, int t){
 T \max Flow = 0;
  vector<flowEdge<T>*> parent(V);
  while(true){
   fill(parent.begin(), parent.end(), nullptr);
    queue<int> Q;
    Q.push(s);
    while(!Q.empty() && !parent[t]){
      int u = Q.front(); Q.pop();
      for(flowEdge<T> *v : adjList[u]){
        if(!parent[v->dest] && v->capacity > v->flow){
          parent[v->dest] = v;
          Q.push(v->dest);
       }
      }
   }
   if(!parent[t]) break;
   T f = inf;
   for(int u = t; u != s; u = parent[u]->res->dest)
      f = min(f, parent[u]->capacity - parent[u]->flow);
   for(int u = t; u != s; u = parent[u]->res->dest)
      parent[u]->addFlow(f);
    maxFlow += f;
 }
  return maxFlow;
}
```

8.4. Algoritmo de Dinic $O(V^2E)$

```
//Maximun Flow using Dinic Algorithm O(EV^2)
T blockingFlow(int u, int t, T flow){
  if(u == t) return flow;
  for(int &i = pos[u]; i < adjList[u].size(); ++i){</pre>
    flowEdge<T> *v = adjList[u][i];
    if (v\rightarrow capacity > v\rightarrow flow \&\& dist[u] + 1 ==

    dist[v->dest]){
      T fv = blockingFlow(v->dest, t, min(flow, v->capacity -
      \rightarrow v->flow));
      if(fv > 0){
        v->addFlow(fv);
        return fv;
      }
    }
  }
  return 0;
}
T dinic(int s, int t){
  T \max Flow = 0;
  dist[t] = 0;
  while (dist [t] != -1) {
    fill(dist.begin(), dist.end(), -1);
    queue<int> Q;
    Q.push(s);
    dist[s] = 0;
    while(!Q.empty()){
      int u = Q.front(); Q.pop();
      for(flowEdge<T> *v : adjList[u]){
        if(dist[v->dest] == -1 \&\& v->flow != v->capacity){
          dist[v->dest] = dist[u] + 1;
          Q.push(v->dest);
        }
      }
    if(dist[t] != -1){
      T f:
      fill(pos.begin(), pos.end(), 0);
      while(f = blockingFlow(s, t, inf))
        maxFlow += f;
```

```
parent[u]->addFlow(cap[t]);
}
return maxFlow;
return {maxFlow, minCost};
}
```

8.5. Flujo máximo de costo mínimo

```
//Max Flow Min Cost
pair<T, T> maxFlowMinCost(int s, int t){
  vector<bool> inQueue(V);
  vector<T> distance(V), cap(V);
  vector<flowEdge<T>*> parent(V);
  T maxFlow = 0, minCost = 0;
  while(true){
    fill(distance.begin(), distance.end(), inf);
    fill(parent.begin(), parent.end(), nullptr);
    fill(cap.begin(), cap.end(), 0);
    distance[s] = 0;
    cap[s] = inf;
    queue<int> Q;
    Q.push(s);
    while(!Q.empty()){
      int u = Q.front(); Q.pop(); inQueue[u] = 0;
     for(flowEdge<T> *v : adjList[u]){
        if(v->capacity > v->flow && distance[v->dest] >

    distance[u] + v->cost){
          distance[v->dest] = distance[u] + v->cost;
          parent[v->dest] = v;
          cap[v->dest] = min(cap[u], v->capacity - v->flow);
          if(!inQueue[v->dest]){
            Q.push(v->dest);
            inQueue[v->dest] = true;
          }
        }
      }
    if(!parent[t]) break;
    maxFlow += cap[t];
    minCost += cap[t] * distance[t];
    for(int u = t; u != s; u = parent[u]->res->dest)
```

9. Estructuras de datos

9.1. Segment Tree

9.1.1. Minimalistic: Point updates, range queries

```
template<typename T>
struct SegmentTree{
  int N:
  vector<T> ST:
  //build from an array in O(n)
  SegmentTree(int N, vector<T> & arr): N(N){
   ST.resize(N << 1);
   for(int i = 0; i < N; ++i)
     ST[N + i] = arr[i];
   for(int i = N - 1; i > 0; --i)
      ST[i] = ST[i << 1] + ST[i << 1 | 1];
 }
  //single element update in i
  void update(int i, T value){
   ST[i += N] = value; //update the element accordingly
   while(i >>= 1)
     ST[i] = ST[i << 1] + ST[i << 1 | 1];
 }
  //single element update in [l, r]
  void update(int 1, int r, T value){
   1 += N, r += N;
   for(int i = 1; i <= r; ++i)
     ST[i] = value;
   1 >>= 1, r >>= 1;
   while(1 \ge 1){
     for(int i = r; i >= 1; --i)
        ST[i] = ST[i << 1] + ST[i << 1 | 1];
     1 >>= 1, r >>= 1;
   }
 }
```

```
//range query, [l, r]
T query(int l, int r){
  T res = 0;
  for(l += N, r += N; l <= r; l >>= 1, r >>= 1){
    if(l & 1) res += ST[1++];
    if(!(r & 1)) res += ST[r--];
  }
  return res;
}
};
```

9.1.2. Dynamic: Range updates and range queries

```
template<typename T>
struct SegmentTreeDin{
 SegmentTreeDin *left, *right;
 int 1, r;
 T sum, lazy;
 SegmentTreeDin(int start, int end, vector<T> & arr):
  → left(NULL), right(NULL), l(start), r(end), sum(0),
  \rightarrow lazy(0){
   if(1 == r) sum = arr[1];
   else{
     int half = 1 + ((r - 1) >> 1);
     left = new SegmentTreeDin(1, half, arr);
     right = new SegmentTreeDin(half+1, r, arr);
     sum = left->sum + right->sum;
   }
 }
 void propagate(T dif){
   sum += (r - 1 + 1) * dif;
   if(1 != r){
     left->lazy += dif;
     right->lazy += dif;
 }
 T sum_query(int start, int end){
```

```
if(lazy != 0){
      propagate(lazy);
      lazy = 0;
    }
    if(end < 1 | | r < start) return 0:
    if(start <= 1 && r <= end) return sum;
    else return left->sum_query(start, end) +

    right->sum_query(start, end);
  }
  void add_range(int start, int end, T dif){
    if(lazy != 0){
      propagate(lazy);
      lazv = 0;
    }
    if(end < 1 | | r < start) return;
    if(start <= 1 && r <= end) propagate(dif);</pre>
    else{
      left->add_range(start, end, dif);
      right->add_range(start, end, dif);
      sum = left->sum + right->sum;
    }
  }
  void add_pos(int i, T sum){
    add_range(i, i, sum);
  }
};
```

9.1.3. Static: Range updates and range queries

```
template<typename T>
struct SegmentTreeEst{
  int size;
  vector<T> sum, lazy;

void rec(int pos, int 1, int r, vector<T> & arr){
  if(1 == r) sum[pos] = arr[1];
  else{
    int half = 1 + ((r - 1) >> 1);
```

```
rec(2*pos+1, 1, half, arr);
    rec(2*pos+2, half+1, r, arr);
    sum[pos] = sum[2*pos+1] + sum[2*pos+2];
}
SegmentTreeEst(int n, vector<T> & arr): size(n){
  int h = ceil(log2(n));
  sum.resize((1 << (h + 1)) - 1);
 lazy.resize((1 << (h + 1)) - 1);
 rec(0, 0, n - 1, arr);
}
void propagate(int pos, int 1, int r, T dif){
  sum[pos] += (r - 1 + 1) * dif;
  if(1 != r){
   lazy[2*pos+1] += dif;
   lazy[2*pos+2] += dif;
 }
}
T sum_query_rec(int start, int end, int pos, int 1, int r){
  if(lazy[pos] != 0){
   propagate(pos, 1, r, lazy[pos]);
   lazy[pos] = 0;
  if(end < 1 | | r < start) return 0;
  if(start <= 1 && r <= end) return sum[pos];</pre>
  else{
    int half = 1 + ((r - 1) >> 1);
   return sum_query_rec(start, end, 2*pos+1, 1, half) +

    sum_query_rec(start, end, 2*pos+2, half+1, r);
}
T sum_query(int start, int end){
  return sum_query_rec(start, end, 0, 0, size - 1);
}
void add_range_rec(int start, int end, int pos, int 1, int r,
\rightarrow T dif){
```

```
if(lazy[pos] != 0){
                                                                   StPer(int start, int end, T val): left(NULL), right(NULL),
      propagate(pos, 1, r, lazy[pos]);
                                                                   lazy[pos] = 0;
                                                                   StPer(int start, int end, StPer* left, StPer* right):
   }
                                                                   → left(left), right(right), l(start), r(end){
    if(end < 1 | | r < start) return;
                                                                     sum = left->sum + right->sum;
    if(start <= 1 && r <= end) propagate(pos, 1, r, dif);
                                                                   }
    else{
      int half = 1 + ((r - 1) >> 1);
                                                                   T sum_query(int start, int end){
      add_range_rec(start, end, 2*pos+1, 1, half, dif);
                                                                     if(end < 1 | | r < start) return 0;
      add_range_rec(start, end, 2*pos+2, half+1, r, dif);
                                                                     if(start <= 1 && r <= end) return sum;</pre>
      sum[pos] = sum[2*pos+1] + sum[2*pos+2];
                                                                     else return left->sum_query(start, end) +
   }

    right->sum_query(start, end);
 }
                                                                   }
                                                                   StPer* update(int pos, T val){
  void add_range(int start, int end, T dif){
    add_range_rec(start, end, 0, 0, size - 1, dif);
                                                                     if(l == r) return new StPer(l, r, sum + val);
 }
                                                                     int half = 1 + ((r - 1) >> 1);
                                                                     if(pos <= half) return new StPer(1, r, left->update(pos,
  void add_pos(int i, T sum){
                                                                     → val), right);
    add_range(i, i, sum);
                                                                     return new StPer(1, r, left, right->update(pos, val));
  }
                                                                   }
};
                                                                 };
```

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9.1.4. Persistent: Point updates, range queries

9.2. Fenwick Tree

```
template<typename T>
struct FenwickTree{
  int N;
  vector<T> bit;

//build from array in O(n), indexed in O
FenwickTree(int N, vector<T> & arr): N(N){
  bit.resize(N);
  for(int i = 0; i < N; ++i){
   bit[i] += arr[i];
   if((i | (i + 1)) < N)
    bit[i | (i + 1)] += bit[i];
  }
}</pre>
```

```
//single element increment
  void update(int pos, T value){
    while(pos < N){
      bit[pos] += value;
      pos \mid = pos + 1;
    }
  }
  //range\ query,\ [0,\ r]
  T query(int r){
    T res = 0;
    while(r >= 0){
      res += bit[r];
      r = (r \& (r + 1)) - 1;
    }
    return res;
  }
  //range query, [l, r]
  T query(int 1, int r){
    return query(r) - query(1 - 1);
  }
};
```

9.3. SQRT Decomposition

```
struct MOquery{
  int 1, r, index, S;
  bool operator<(const MOquery & q) const{
    int c_o = 1 / S, c_q = q.1 / S;
    if(c_o == c_q)
      return r < q.r;
    return c_o < c_q;
  }
};

template<typename T>
struct SQRT{
  int N, S;
  vector<T> A, B;
```

```
SQRT(int N): N(N){
  this->S = sqrt(N + .0) + 1;
 A.assign(N, 0);
 B.assign(S, 0);
}
void build(vector<T> & arr){
  A = vector<int>(arr.begin(), arr.end());
 for(int i = 0; i < N; ++i) B[i / S] += A[i];
}
//single element update
void update(int pos, T value){
  int k = pos / S;
  A[pos] = value;
 T res = 0;
 for(int i = k * S, end = min(N, (k + 1) * S) - 1; i \le end;
  \rightarrow ++i) res += A[i];
 B[k] = res;
}
//range query, [l, r]
T query(int 1, int r){
 T res = 0;
  int c_1 = 1 / S, c_r = r / S;
  if(c_1 == c_r){
    for(int i = 1; i <= r; ++i) res += A[i];
  }else{
    for(int i = 1, end = (c_1 + 1) * S - 1; i \le end; ++i)
    \rightarrow res += A[i];
    for(int i = c_1 + 1; i \le c_r - 1; ++i) res += B[i];
    for(int i = c_r * S; i <= r; ++i) res += A[i];
  }
  return res;
}
//range queries offline using MO's algorithm
vector<T> MO(vector<MOquery> & queries){
  vector<T> ans(queries.size());
  sort(queries.begin(), queries.end());
```

```
T current = 0;
                                                                        return (right ? right->height : 0) - (left ? left->height :
    int prevL = 0, prevR = -1;
                                                                        \rightarrow 0);
    int i, j;
                                                                     }
    for(const MOquery & q : queries){
      for(i = prevL, j = min(prevR, q.l - 1); i \le j; ++i){
                                                                     AVLNode *maxLeftChild(){
        //remove from the left
                                                                        AVLNode *ret = this;
        current -= A[i];
                                                                       while(ret->left) ret = ret->left;
      }
                                                                       return ret;
      for(i = prevL - 1; i >= q.l; --i){
                                                                     }
        //add to the left
                                                                   };
        current += A[i];
      }
                                                                    template<typename T>
      for(i = max(prevR + 1, q.1); i \le q.r; ++i){
                                                                    struct AVLTree{
        //add to the right
                                                                      AVLNode<T> *root;
        current += A[i];
      }
                                                                      AVLTree(): root(NULL){}
      for(i = prevR; i >= q.r + 1; --i){
        //remove from the right
                                                                      inline int nodeSize(AVLNode<T> *& pos){return pos ?
        current -= A[i];
                                                                      \rightarrow pos->size: 0;}
      }
      prevL = q.1, prevR = q.r;
                                                                      inline int nodeHeight(AVLNode<T> *& pos){return pos ?
      ans[q.index] = current;
                                                                      → pos->height: 0;}
    }
    return ans;
                                                                     inline void update(AVLNode<T> *& pos){
                                                                        if(!pos) return;
                                                                       pos->height = 1 + max(nodeHeight(pos->left),
};
                                                                        → nodeHeight(pos->right));
                                                                       pos->size = 1 + nodeSize(pos->left) + nodeSize(pos->right);
9.4. AVL Tree
template<typename T>
                                                                      int size(){return nodeSize(root);}
struct AVLNode{
  AVLNode<T> *left, *right;
                                                                      void leftRotate(AVLNode<T> *& x){
  short int height;
                                                                        AVLNode<T> *y = x->right, *t = y->left;
  int size;
                                                                       y->left = x, x->right = t;
  T value;
                                                                       update(x), update(y);
                                                                       x = y;
  AVLNode(T value = 0): left(NULL), right(NULL), value(value),
  \rightarrow height(1), size(1){}
                                                                     void rightRotate(AVLNode<T> *& y){
  inline short int balance(){
                                                                       AVLNode<T> *x = y->left, *t = x->right;
```

```
x->right = y, y->left = t;
                                                                       if(!pos->left) pos = pos->right;
  update(y), update(x);
                                                                       else if(!pos->right) pos = pos->left;
  y = x;
                                                                       else{
}
                                                                         pos->value = pos->right->maxLeftChild()->value;
                                                                         erase(pos->right, pos->value);
void updateBalance(AVLNode<T> *& pos){
                                                                       }
                                                                     }
  if(!pos) return;
  short int bal = pos->balance();
                                                                     update(pos), updateBalance(pos);
  if(bal > 1){
    if(pos->right->balance() < 0) rightRotate(pos->right);
    leftRotate(pos);
                                                                   void insert(T value){insert(root, value);}
  else if(bal < -1){
    if(pos->left->balance() > 0) leftRotate(pos->left);
                                                                   void erase(T value){erase(root, value);}
    rightRotate(pos);
  }
                                                                   void updateVal(T old, T New){
}
                                                                     if(search(old))
                                                                       erase(old), insert(New);
void insert(AVLNode<T> *&pos, T & value){
                                                                   }
  if(pos){
    value < pos->value ? insert(pos->left, value) :
                                                                   T kth(int i){

    insert(pos->right, value);

                                                                     assert(0 <= i && i < nodeSize(root));</pre>
    update(pos), updateBalance(pos);
                                                                     AVLNode<T> *pos = root;
  }else{
                                                                     while(i != nodeSize(pos->left)){
                                                                       if(i < nodeSize(pos->left)){
    pos = new AVLNode<T>(value);
  }
                                                                         pos = pos->left;
}
                                                                       }else{
                                                                         i -= nodeSize(pos->left) + 1;
                                                                         pos = pos->right;
AVLNode<T> *search(T & value){
  AVLNode<T> *pos = root;
  while(pos){
                                                                     }
    if(value == pos->value) break;
                                                                     return pos->value;
    pos = (value < pos->value ? pos->left : pos->right);
  }
                                                                   int lessThan(T & x){
  return pos;
}
                                                                     int ans = 0;
                                                                     AVLNode<T> *pos = root;
void erase(AVLNode<T> *&pos, T & value){
                                                                     while(pos){
                                                                       if(x > pos->value){
  if(!pos) return;
  if(value < pos->value) erase(pos->left, value);
                                                                         ans += nodeSize(pos->left) + 1;
  else if(value > pos->value) erase(pos->right, value);
                                                                         pos = pos->right;
  else{
                                                                       }else{
```

```
pos = pos->left;
    }
  }
  return ans;
}
int lessThanOrEqual(T & x){
  int ans = 0;
  AVLNode<T> *pos = root;
  while(pos){
   if(x < pos->value){
      pos = pos->left;
    }else{
      ans += nodeSize(pos->left) + 1;
      pos = pos->right;
    }
  }
  return ans;
}
int greaterThan(T & x){
  int ans = 0;
  AVLNode<T> *pos = root;
  while(pos){
   if(x < pos->value){
      ans += nodeSize(pos->right) + 1;
      pos = pos->left;
    }else{
      pos = pos->right;
    }
  }
  return ans;
}
int greaterThanOrEqual(T & x){
  int ans = 0;
  AVLNode<T> *pos = root;
  while(pos){
    if(x > pos->value){
      pos = pos->right;
    }else{
```

```
ans += nodeSize(pos->right) + 1;
       pos = pos->left;
   }
   return ans;
  }
  int equalTo(T & x){
    return lessThanOrEqual(x) - lessThan(x);
  }
  void build(AVLNode<T> *& pos, vector<T> & arr, int i, int j){
   if(i > j) return;
    int m = i + ((j - i) >> 1);
   pos = new AVLNode<T>(arr[m]);
   build(pos->left, arr, i, m - 1);
   build(pos->right, arr, m + 1, j);
   update(pos);
  }
  void build(vector<T> & arr){
    build(root, arr, 0, (int)arr.size() - 1);
  }
  void output(AVLNode<T> *pos, vector<T> & arr, int & i){
   if(pos){
      output(pos->left, arr, i);
      arr[++i] = pos->value;
      output(pos->right, arr, i);
   }
 }
  void output(vector<T> & arr){
   int i = -1;
    output(root, arr, i);
 }
};
```

9.5. Treap

```
template<typename T>
struct TreapNode{
  TreapNode<T> *left, *right;
  T value;
  int key, size;
  //fields for queries
  bool rev;
  T sum, add;
  TreapNode(T value = 0): value(value), key(rand()), size(1),
  → left(NULL), right(NULL), sum(value), add(0), rev(false){}
};
template<typename T>
struct Treap{
  TreapNode<T> *root;
  Treap(): root(NULL) {}
  inline int nodeSize(TreapNode<T>* t){return t ? t->size: 0;}
  inline T nodeSum(TreapNode<T>* t){return t ? t->sum : 0;}
  inline void update(TreapNode<T>* &t){
    if(!t) return;
    t->size = 1 + nodeSize(t->left) + nodeSize(t->right);
    t->sum = t->value; //reset node fields
    push(t->left), push(t->right); //push changes to child
    \rightarrow nodes
    t->sum = t->value + nodeSum(t->left) + nodeSum(t->right);
    \rightarrow //combine(left,t,t), combine(t,right,t)
  }
  int size(){return nodeSize(root);}
  void merge(TreapNode<T>* &t, TreapNode<T>* t1, TreapNode<T>*
  if(!t1) t = t2;
```

```
else if(!t2) t = t1;
  else if(t1->key > t2->key)
    merge(t1->right, t1->right, t2), t = t1;
  else
    merge(t2\rightarrow left, t1, t2\rightarrow left), t = t2;
 update(t);
}
void split(TreapNode<T>* t, T & x, TreapNode<T>* &t1,
→ TreapNode<T>* &t2){
 if(!t)
    return void(t1 = t2 = NULL);
  if(x < t->value)
    split(t->left, x, t1, t->left), t2 = t;
  else
    split(t->right, x, t->right, t2), t1 = t;
 update(t);
}
void insert(TreapNode<T>* &t, TreapNode<T>* x){
  if(!t) t = x;
  else if(x->key > t->key)
    split(t, x->value, x->left, x->right), t = x;
  else
    insert(x->value < t->value ? t->left : t->right, x);
  update(t);
}
TreapNode<T>* search(T & x){
  TreapNode<T> *t = root;
  while(t){
   if(x == t->value) break;
    t = (x < t->value ? t->left : t->right);
  }
  return t;
}
void erase(TreapNode<T>* &t, T & x){
  if(!t) return;
  if(t->value == x)
    merge(t, t->left, t->right);
```

```
else
                                                                   //OPERATIONS FOR IMPLICIT TREAP
    erase(x < t->value ? t->left : t->right, x);
  update(t);
                                                                   inline void push(TreapNode<T>* t){
}
                                                                     if(!t) return;
                                                                     //add in range example
void insert(T & x){insert(root, new TreapNode<T>(x));}
                                                                     if(t->add){
                                                                       t->value += t->add;
void erase(T & x){erase(root, x);}
                                                                       t->sum += t->add * nodeSize(t);
                                                                       if(t->left) t->left->add += t->add;
void updateVal(T & old, T & New){
                                                                       if(t->right) t->right->add += t->add;
  if(search(old))
                                                                       t->add = 0;
    erase(old), insert(New);
}
                                                                     //reverse range example
                                                                     if(t->rev){
T kth(int i){
                                                                       swap(t->left, t->right);
  assert(0 <= i && i < nodeSize(root));</pre>
                                                                       if(t->left) t->left->rev ^= true;
  TreapNode<T> *t = root;
                                                                       if(t->right) t->right->rev ^= true;
  while(i != nodeSize(t->left)){
                                                                       t->rev = false;
    if(i < nodeSize(t->left)){
                                                                     }
                                                                   }
      t = t->left;
    }else{
      i -= nodeSize(t->left) + 1;
                                                                   void split2(TreapNode<T>* t, int i, TreapNode<T>* &t1,
      t = t->right;
                                                                   → TreapNode<T>* &t2){
    }
                                                                     if(!t)
                                                                       return void(t1 = t2 = NULL);
  return t->value;
                                                                     push(t);
}
                                                                     int curr = nodeSize(t->left);
                                                                     if(i <= curr)</pre>
int lessThan(T & x){
                                                                       split2(t->left, i, t1, t->left), t2 = t;
  int ans = 0;
                                                                     else
  TreapNode<T> *t = root;
                                                                       split2(t->right, i - curr - 1, t->right, t2), t1 = t;
  while(t){
                                                                     update(t);
                                                                   }
    if(x > t->value){
      ans += nodeSize(t->left) + 1;
      t = t->right;
                                                                   inline int aleatorio(){
    }else{
                                                                     return (rand() << 15) + rand();
      t = t->left;
                                                                   }
    }
  }
                                                                   void merge2(TreapNode<T>* &t, TreapNode<T>* t1, TreapNode<T>*
  return ans;
                                                                   }
                                                                     push(t1), push(t2);
```

```
if(!t1) t = t2;
                                                                    T nth(TreapNode<T>* t, int i){
  else if(!t2) t = t1;
                                                                      push(t);
  else if(aleatorio() % (nodeSize(t1) + nodeSize(t2)) <</pre>
                                                                      assert(0 <= i && i < nodeSize(t));</pre>
  \rightarrow nodeSize(t1))
                                                                      int curr = nodeSize(t->left);
    merge2(t1->right, t1->right, t2), t = t1;
                                                                      if(i == curr)
                                                                        return t->value;
    merge2(t2->left, t1, t2->left), t = t2;
                                                                      else if(i < curr)</pre>
  update(t);
                                                                        return nth(t->left, i);
}
                                                                      else
                                                                        return nth(t->right, i - curr - 1);
                                                                    }
//insert the element "x" at position "i"
void insert_at(T & x, int i){
  if(i > nodeSize(root)) return;
                                                                    //update value of element at position "i" with "x"
  TreapNode<T> *t1 = NULL, *t2 = NULL;
                                                                    void update_at(T & x, int i){update_at(root, x, i);}
  split2(root, i, t1, t2);
  merge2(root, t1, new TreapNode<T>(x));
                                                                    //ith element
                                                                    T nth(int i){return nth(root, i);}
  merge2(root, root, t2);
}
                                                                    //add "val" in [l, r]
//delete element at position "i"
                                                                    void add_update(T & val, int l, int r){
void erase at(int i){
                                                                      TreapNode<T> *t1 = NULL, *t2 = NULL, *t3 = NULL;
  if(i >= nodeSize(root)) return;
                                                                      split2(root, 1, t1, t2);
  TreapNode<T> *t1 = NULL, *t2 = NULL, *t3 = NULL;
                                                                      split2(t2, r - 1 + 1, t2, t3);
  split2(root, i, t1, t2);
                                                                      t2->add += val;
  split2(t2, 1, t2, t3);
                                                                      merge2(root, t1, t2);
  merge2(root, t1, t3);
                                                                      merge2(root, root, t3);
}
                                                                    }
                                                                    //reverse [l, r]
void update_at(TreapNode<T>* t, T & x, int i){
  push(t);
                                                                    void reverse_update(int 1, int r){
  assert(0 <= i && i < nodeSize(t));</pre>
                                                                      TreapNode<T> *t1 = NULL, *t2 = NULL, *t3 = NULL;
  int curr = nodeSize(t->left);
                                                                      split2(root, 1, t1, t2);
  if(i == curr)
                                                                      split2(t2, r - 1 + 1, t2, t3);
                                                                      t2->rev ^= true;
    t->value = x;
  else if(i < curr)</pre>
                                                                      merge2(root, t1, t2);
    update_at(t->left, x, i);
                                                                      merge2(root, root, t3);
    update_at(t->right, x, i - curr - 1);
  update(t);
                                                                    //rotate [l, r] k times to the right
}
                                                                    void rotate_update(int k, int l, int r){
```

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```
TreapNode<T> *t1 = NULL, *t2 = NULL, *t3 = NULL, *t4 =
                                                                     vector<int> logs;
    → NULL;
                                                                    int K, N;
    split2(root, 1, t1, t2);
    split2(t2, r - 1 + 1, t2, t3);
                                                                     SparseTable(vector<T> & arr){
    k %= nodeSize(t2);
                                                                      N = arr.size():
    split2(t2, nodeSize(t2) - k, t2, t4);
                                                                      K = log2(N) + 2;
    merge2(root, t1, t4);
                                                                      ST.assign(K + 1, vector < T > (N));
    merge2(root, root, t2);
                                                                      logs.assign(N + 1, 0);
    merge2(root, root, t3);
                                                                      for(int i = 2; i \le N; ++i)
  }
                                                                        logs[i] = logs[i >> 1] + 1;
                                                                      for(int i = 0; i < N; ++i)
  //sum query in [l, r]
                                                                        ST[0][i] = arr[i];
  T sum_query(int 1, int r){
                                                                      for(int j = 1; j \le K; ++j)
    TreapNode<T> *t1 = NULL, *t2 = NULL, *t3 = NULL;
                                                                        for(int i = 0; i + (1 << j) <= N; ++i)
                                                                          ST[j][i] = min(ST[j-1][i], ST[j-1][i+(1 << (j-1)[i])
    split2(root, 1, t1, t2);
    split2(t2, r - 1 + 1, t2, t3);
                                                                           → 1))]); //put the function accordingly
                                                                    }
    T ans = nodeSum(t2);
    merge2(root, t1, t2);
    merge2(root, root, t3);
                                                                    T sum(int 1, int r){ //non-idempotent functions
                                                                      T ans = 0;
    return ans;
  }
                                                                      for(int j = K; j \ge 0; --j){
                                                                        if((1 << j) <= r - 1 + 1){
  void inorder(TreapNode<T>* t){
                                                                          ans += ST[j][1];
    if(!t) return;
                                                                          1 += 1 << j;
    push(t);
                                                                        }
                                                                      }
    inorder(t->left);
    cout << t->value << " ";
                                                                      return ans;
                                                                    }
    inorder(t->right);
  }
                                                                    T minimal(int 1, int r){ //idempotent functions
  void inorder(){inorder(root);}
                                                                       int j = logs[r - l + 1];
};
                                                                      return min(ST[j][1], ST[j][r - (1 << j) + 1]);
                                                                    }
                                                                  };
     Sparse table
                                                                  9.7. Disjoint
9.6.1. Normal
```

//build on $O(n \log n)$, queries in O(1) for any operation

template<typename T>

struct DisjointSparseTable{

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template<typename T>

vector<vector<T>> ST;

struct SparseTable{

```
vector<vector<T>> left, right;
                                                                      vector<int> pref; //just use this if you want sums
  int K, N;
                                                                      //queries indexed in base 1, complexity for all queries:
  DisjointSparseTable(vector<T> & arr){
                                                                      \hookrightarrow O(log(max_element))
    N = arr.size():
                                                                      //build from [from, to) with non-negative values in range [x,
    K = log2(N) + 2;
    left.assign(K + 1, vector<T>(N));
                                                                      //you can use vector iterators or array pointers
                                                                      WaveletTree(vector<int>::iterator from, vector<int>::iterator
    right.assign(K + 1, vector<T>(N));
    for(int j = 0; (1 << j) <= N; ++j){
                                                                      \rightarrow to, int x, int y): lo(x), hi(y){
      int mask = (1 << j) - 1;</pre>
                                                                        if(from >= to) return;
      T acum = 0; //neutral element of your operation
                                                                        int m = (lo + hi) / 2;
      for(int i = 0; i < N; ++i){
                                                                        auto f = [m](int x){return x <= m;};
        acum += arr[i]; //your operation
                                                                        freq.reserve(to - from + 1);
        left[i][i] = acum;
                                                                        freq.push_back(0);
        if((i & mask) == mask) acum = 0; //neutral element of
                                                                        pref.reserve(to - from + 1);
                                                                        pref.push_back(0);

→ your operation

      }
                                                                        for(auto it = from; it != to; ++it){
      acum = 0; //neutral element of your operation
                                                                          freq.push_back(freq.back() + f(*it));
      for(int i = N-1; i >= 0; --i){
                                                                          pref.push_back(pref.back() + *it);
        acum += arr[i]; //your operation
        right[j][i] = acum;
                                                                        if(hi != lo){
        if((i & mask) == 0) acum = 0; //neutral element of your
                                                                          auto pivot = stable_partition(from, to, f);
        \hookrightarrow operation
                                                                          left = new WaveletTree(from, pivot, lo, m);
      }
                                                                          right = new WaveletTree(pivot, to, m + 1, hi);
    }
                                                                        }
  }
                                                                      }
  T query(int 1, int r){
                                                                      //kth element in [l, r]
    if(l == r) return left[0][1];
                                                                      int kth(int 1, int r, int k){
    int i = 31 - __builtin_clz(l^r);
                                                                        if(1 > r) return 0;
    return left[i][r] + right[i][l]; //your operation
                                                                        if(lo == hi) return lo;
  }
                                                                        int lb = freq[l - 1], rb = freq[r];
};
                                                                        int inLeft = rb - lb;
                                                                        if(k <= inLeft) return left->kth(lb + 1, rb, k);
                                                                        else return right->kth(l - lb, r - rb, k - inLeft);
      Wavelet Tree
9.8.
                                                                      }
struct WaveletTree{
                                                                      //number of elements less than or equal to k in [l, r]
  int lo, hi;
                                                                      int lessThanOrEqual(int 1, int r, int k){
  WaveletTree *left, *right;
                                                                        if(l > r \mid \mid k < lo) return 0;
  vector<int> freq;
                                                                        if(hi \leq k) return r - 1 + 1;
```

```
int lb = freq[l - 1], rb = freq[r];
    return left->lessThanOrEqual(lb + 1, rb, k) +
    → right->lessThanOrEqual(1 - lb, r - rb, k);
  }
  //number of elements equal to k in [l, r]
  int equalTo(int 1, int r, int k){
    if(1 > r \mid \mid k < lo \mid \mid k > hi) return 0;
    if(lo == hi) return r - l + 1;
    int lb = freq[1 - 1], rb = freq[r];
    int m = (lo + hi) / 2;
    if(k <= m) return left->equalTo(lb + 1, rb, k);
    else return right->equalTo(1 - lb, r - rb, k);
  }
  //sum of elements less than or equal to k in [l, r]
  int sum(int 1, int r, int k){
    if(1 > r \mid \mid k < lo) return 0;
    if(hi <= k) return pref[r] - pref[l - 1];</pre>
    int lb = freq[l - 1], rb = freq[r];
    return left->sum(lb + 1, rb, k) + right->sum(l - lb, r -
    \hookrightarrow rb, k);
  }
};
```

9.9. Ordered Set C++

```
if(t == 0){ //insert}
      conj.insert(n);
   }else if(t == 1){ //search
      if(conj.find(n) != conj.end()) cout << "Found\n";</pre>
      else cout << "Not found\n";
   }else if(t == 2){ //delete
      conj.erase(n);
   }else if(t == 3){ //update
      cin >> m;
      if(conj.find(n) != conj.end()){
        conj.erase(n);
        conj.insert(n);
     }
   }else if(t == 4){ //lower bound
      cout << conj.order_of_key(n) << "\n";</pre>
   }else if(t == 5){ //qet nth element
      auto pos = conj.find_by_order(n);
      if(pos != conj.end()) cout << *pos << "\n";</pre>
      else cout << "-1\n";
   }
 }
 return 0;
}
```

9.10. Splay Tree

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9.11. Red Black Tree

10. Cadenas

10.1. Trie

```
struct Node{
    bool isWord = false;
  map<char, Node*> letters;
};
struct Trie{
  Node* root;
  Trie(){
    root = new Node();
  }
  inline bool exists(Node * actual, const char & c){
    return actual->letters.find(c) != actual->letters.end();
  }
  void InsertWord(const string& word){
    Node* current = root;
    for(auto & c : word){
      if(!exists(current, c))
        current->letters[c] = new Node();
      current = current->letters[c];
    current->isWord = true;
  }
  bool FindWord(const string& word){
    Node* current = root;
    for(auto & c : word){
      if(!exists(current, c))
        return false;
      current = current->letters[c];
    }
    return current->isWord;
  }
  void printRec(Node * actual, string acum){
```

```
if(actual->isWord){
      cout << acum << "\n";
    for(auto & next : actual->letters)
      printRec(next.second, acum + next.first);
  }
  void printWords(const string & prefix){
   Node * actual = root;
   for(auto & c : prefix){
      if(!exists(actual, c)) return;
      actual = actual->letters[c];
    printRec(actual, prefix);
 }
};
10.2. KMP
struct kmp{
  vector<int> aux;
  string pattern;
  kmp(string pattern){
    this->pattern = pattern;
    aux.resize(pattern.size());
    int i = 1, j = 0;
    while(i < pattern.size()){</pre>
      if(pattern[i] == pattern[j])
        aux[i++] = ++j;
      else{
        if(j == 0) aux[i++] = 0;
        else j = aux[j - 1];
      }
   }
  vector<int> search(string & text){
    vector<int> ans;
```

int i = 0, j = 0;

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```
while(i < text.size() && j < pattern.size()){
    if(text[i] == pattern[j]){
        ++i, ++j;
        if(j == pattern.size()){
            ans.push_back(i - j);
            j = aux[j - 1];
        }
    }else{
        if(j == 0) ++i;
        else j = aux[j - 1];
    }
    return ans;
}</pre>
```

10.3. Aho-Corasick

```
const int M = 26;
struct node{
  vector<int> child;
  int p = -1;
  char c = 0;
  int suffixLink = -1, endLink = -1;
  int id = -1;
  node(int p = -1, char c = 0) : p(p), c(c){
    child.resize(M, −1);
 }
};
struct AhoCorasick{
  vector<node> t;
  vector<int> lenghts;
  int wordCount = 0;
  AhoCorasick(){
    t.emplace_back();
  }
```

```
void add(const string & s){
  int u = 0;
 for(char c : s){
    if(t[u].child[c-'a'] == -1){
      t[u].child[c-'a'] = t.size();
      t.emplace_back(u, c);
   u = t[u].child[c-'a'];
  t[u].id = wordCount++;
 lenghts.push_back(s.size());
void link(int u){
  if(u == 0){
    t[u].suffixLink = 0;
   t[u].endLink = 0;
   return;
  }
  if(t[u].p == 0){
   t[u].suffixLink = 0;
   if(t[u].id != -1) t[u].endLink = u;
    else t[u].endLink = t[t[u].suffixLink].endLink;
   return;
  int v = t[t[u].p].suffixLink;
  char c = t[u].c;
  while(true){
    if(t[v].child[c-'a'] != -1){
     t[u].suffixLink = t[v].child[c-'a'];
     break;
   }
    if(v == 0){
     t[u].suffixLink = 0;
     break;
    v = t[v].suffixLink;
  if(t[u].id != -1) t[u].endLink = u;
  else t[u].endLink = t[t[u].suffixLink].endLink;
}
```

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```
void build(){
    queue<int> Q;
    Q.push(0);
    while(!Q.empty()){
     int u = Q.front(); Q.pop();
      link(u);
     for(int v = 0; v < M; ++v)
        if(t[u].child[v] != -1)
          Q.push(t[u].child[v]);
    }
  }
  int match(const string & text){
    int u = 0;
    int ans = 0;
    for(int j = 0; j < text.size(); ++j){</pre>
      int i = text[j] - 'a';
      while(true){
        if(t[u].child[i] != -1){
          u = t[u].child[i];
          break;
        }
        if(u == 0) break;
        u = t[u].suffixLink;
      }
      int v = u;
      while(true){
        v = t[v].endLink;
        if(v == 0) break;
        ++ans;
        int idx = j + 1 - lenghts[t[v].id];
        cout << "Found word #" << t[v].id << " at position " <<</pre>
        \rightarrow idx << "\n";
        v = t[v].suffixLink;
      }
    }
    return ans;
  }
};
```

10.4. Rabin-Karp

10.5. Suffix Array

10.6. Función Z

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11. Varios

11.1. Lectura y escritura de __int128

```
//cout for __int128
ostream & operator << (ostream & os, const __int128 & value) {
  char buffer[64];
  char *pos = end(buffer) - 1;
  *pos = ' \setminus 0';
  __int128 tmp = value < 0 ? -value : value;
  do{
    --pos;
    *pos = tmp \% 10 + ^{'}0';
    tmp /= 10;
  }while(tmp != 0);
  if(value < 0){
    --pos;
    *pos = '-';
  return os << pos;
//cin for __int128
istream &operator>>(istream &is, __int128 & value){
  char buffer[64];
  is >> buffer;
  char *pos = begin(buffer);
  int sgn = 1;
  value = 0;
  if(*pos == '-'){
    sgn = -1;
    ++pos;
  }else if(*pos == '+'){
    ++pos;
  }
  while(*pos != '\0'){
    value = (value << 3) + (value << 1) + (*pos - '0');
    ++pos;
  }
  value *= sgn;
  return is;
```

}

11.2. Longest Common Subsequence (LCS)

```
int lcs(string & a, string & b){
  int m = a.size(), n = b.size();
  vector<vector<int>> aux(m + 1, vector<int>(n + 1));
  for(int i = 1; i <= m; ++i){
    for(int j = 1; j <= n; ++j){
      if(a[i - 1] == b[j - 1])
        aux[i][j] = 1 + aux[i - 1][j - 1];
      else
        aux[i][j] = max(aux[i - 1][j], aux[i][j - 1]);
    }
}
return aux[m][n];
}</pre>
```

11.3. Longest Increasing Subsequence (LIS)

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11.4. Levenshtein Distance

11.5. Día de la semana

```
//0:saturday, 1:sunday, ..., 6:friday
int dayOfWeek(int d, int m, lli y){
  if(m == 1 || m == 2){
    m += 12;
    --y;
  }
  int k = y % 100;
  lli j = y / 100;
  return (d + 13*(m+1)/5 + k + k/4 + j/4 + 5*j) % 7;
}
```

11.6. 2SAT

```
struct satisfiability_twosat{
  int n;
  vector<vector<int>> imp;

satisfiability_twosat(int n) : n(n), imp(2 * n) {}

void add_edge(int u, int v){imp[u].push_back(v);}
```

```
int neg(int u){return (n << 1) - u - 1;}</pre>
  void implication(int u, int v){
    add_edge(u, v);
    add_edge(neg(v), neg(u));
  vector<bool> solve(){
    int size = 2 * n;
    vector<int> S, B, I(size);
    function < void(int) > dfs = [&](int u){
      B.push_back(I[u] = S.size());
      S.push_back(u);
      for(int v : imp[u])
        if(!I[v]) dfs(v);
        else while (I[v] < B.back()) B.pop_back();</pre>
      if(I[u] == B.back())
        for(B.pop_back(), ++size; I[u] < S.size();</pre>

    S.pop_back())

          I[S.back()] = size;
    };
    for(int u = 0; u < 2 * n; ++u)
      if(!I[u]) dfs(u);
    vector<bool> values(n);
    for(int u = 0; u < n; ++u)
      if(I[u] == I[neg(u)]) return {};
      else values[u] = I[u] < I[neg(u)];</pre>
    return values;
 }
};
```

11.7. Código Gray

```
//gray code
int gray(int n){
  return n ^ (n >> 1);
}

//inverse gray code
int inv_gray(int g){
  int n = 0;
  while(g){
    n ^= g;
    g >>= 1;
  }
  return n;
}
```

11.8. Contar número de unos en binario en un rango