



ASIS CTF FINALS - BIT GAME

Steliyan Syarov

GENERAL INFO:

- Released in the last 1 / 3 of the CTF
- Reverse **crypto** challenge
- Description:
 - A game for grown ups, just play with bits to find out the sought after secret.
 - 2 files given: **single_bits.py + output.txt**
- Research + Math!

```

import random
from Crypto.Util.number import *
from flag import flag

def gen_rand(nbit, l):
    R = []
    while True:
        r = random.randint(0, nbit-1)
        if r not in R:
            R.append(r)
            if len(R) == l:
                break
    R.sort()
    rbit = '1'
    for i in range(l-1):
        rbit += (R[i+1] - R[i] - 1) * '0' + '1'
    rbit += (nbit - R[-1] - 1) * '0'
    return int(rbit, 2)

def genkey(p, l):
    n = len(bin(p)[2:])
    f, skey = gen_rand(n, l), gen_rand(n, l)
    pkey = f * inverse(skey, p) % p
    return (p, pkey), skey

def encrypt(msg, pkey):
    p, g = pkey
    msg, enc, n = bytes_to_long(msg), [], len(bin(p)[2:])
    for b in bin(msg)[2:]:
        s, t = gen_rand(n, l), gen_rand(n, l)
        c = (s * g + t) % p
        if b == '0':
            enc.append((c, t))
        else:
            enc.append((p-c, t))
    return enc

```

WHAT WAS GIVEN...

```

p = 862718293348820473429344482784628181556388621521298319395315527974911
l = 5

```

```

pkey, skey = genkey(p, l)
enc = encrypt(flag, pkey)
H = pkey[1] ** 2 % p

```

```

print('H =', H)
print('enc =', enc)

```

```

enc=[
    (693452030319839722726551032492658894345570855091353510812978453643631,
    374147273652096600518280302026002167242456968463360),
    (68855961238876671161727227684383103592686215888250125242406607896646,
    1684996666696915012012216278925205994248717536439095543946901716996),
    (378940288706735756349711787806658208666776697748239130653541467467677,
    46768052395269458121440307166481563409072781262848),
    (96192682537252933190719942697876279568205071028966045728552811342752,
    1427269470777442821119947625424370788877729792),
    (326046970830869347127025860081671929137127663636883986249275995969665,
    187072209578355573854590213492815704776701722494976),
    (340905918585524940327083299326143363870527182200486260808937793514394,
    48264630071912636341289319625352763693625718863839232),
    (580732406123392374746863117598833633024192305365332270959719531036156,
    6628619432568334886610593630907295735406448532595523987702272),
    ...

```

WHAT DOES THE CODE DO ?

```
import random
from Crypto.Util.number import *
from flag import flag
```

```
def gen_rand(nbit, l):
    R = []
    while True:
        r = random.randint(0, nbit-1)
        if r not in R:
            R.append(r)
            if len(R) == l:
                break
    R.sort()
    rbit = '1'
    for i in range(l-1):
        rbit = (R[i+1] - R[i] - 1) * '0' + '1'
    rbit += (nbit - R[-1] - 1) * '0'
    return int(rbit, 2)
```

```
def genkey(p, l):
    n = len(bin(p)[2:])
    f, skey = gen_rand(n, l), gen_rand(n, l)
    pkey = f * inverse(skey, p) % p
    return (p, pkey), skey
```

```
def encrypt(msg, pkey):
    p, g = pkey
    msg, enc, n = bytes_to_long(msg), [], len(bin(p)[2:])
    for b in bin(msg)[2:]:
        s, t = gen_rand(n, l), gen_rand(n, l)
        c = (s * g + t) % p
        if b == '0':
            enc.append((c, t))
        else:
            enc.append((p-c, t))
    return enc
```

```
p = 862718293348826473429344482784628181556388621521298319395315527974911
l = 5
```

```
pkey, skey = genkey(p, l)
enc = encrypt(flag, pkey)
H = pkey[1] ** 2 % p
```

```
print('H =', H)
print('enc =', enc)
```

```
def gen_rand(nbit, l):
    R = []
    while True:
        r = random.randint(0, nbit-1)
        if r not in R:
            R.append(r)
            if len(R) == l:
                break
    R.sort()
    rbit = '1'
    for i in range(l-1):
        rbit += (R[i+1] - R[i] - 1) * '0' + '1'
    rbit += (nbit - R[-1] - 1) * '0'
    return int(rbit, 2)
```

[illegible]

WHAT DOES THE CODE DO ?

```
import random
from Crypto.Util.number import *
from flag import flag

def gen_rand(nbit, l):
    R = []
    while True:
        r = random.randint(0, nbit-1)
        if r not in R:
            R.append(r)
            if len(R) == l:
                break
    R.sort()
    rbit = '1'
    for i in range(l-1):
        rbit += (R[i+1] - R[i] - 1) * '0' + '1'
    rbit += (nbit - R[-1] - 1) * '0'
    return int(rbit, 2)

def genkey(p, l):
    n = len(bin(p)[2:])
    f, skey = gen_rand(n, l), gen_rand(n, l)
    pkey = f * inverse(skey, p) % p
    return (p, pkey), skey

def encrypt(msg, pkey):
    p, g = pkey
    msg, enc, n = bytes_to_long(msg), [], len(bin(p)[2:])
    for b in bin(msg)[2:]:
        s, t = gen_rand(n, l), gen_rand(n, l)
        c = (s * g + t) % p
        if b == '0':
            enc.append((c, t))
        else:
            enc.append((p-c, t))
    return enc

p = 862718293348820473429344482784628181556388621521298319395315527974911
l = 5

pkey, skey = genkey(p, l)
enc = encrypt(flag, pkey)
H = pkey[l] ** 2 % p

print('H =', H)
print('enc =', enc)
```

```
def genkey(p, l):
    n = len(bin(p)[2:])      # n = 229
    f, skey = gen_rand(n, l), gen_rand(n, l)
    pkey = f * inverse(skey, p) % p
    return (p, pkey), skey
```

WHAT DOES THE CODE DO ?

```
import random
from Crypto.Util.number import *
from flag import flag

def gen_rand(nbit, l):
    R = []
    while True:
        r = random.randint(0, nbit-1)
        if r not in R:
            R.append(r)
            if len(R) == l:
                break
    R.sort()
    rbit = '1'
    for i in range(l-1):
        rbit += (R[i+1] - R[i] - 1) * '0' + '1'
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    return int(rbit, 2)

def genkey(p, l):
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    pkey = f * inverse(skey, p) % p
    return (p, pkey), skey

def encrypt(msg, pkey):
    p, g = pkey
    msg, enc, n = bytes_to_long(msg), [], len(bin(p)[2:])
    for b in bin(msg)[2:]:
        s, t = gen_rand(n, l), gen_rand(n, l)
        c = (s * g + t) % p
        if b == '0':
            enc.append((c, t))
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            enc.append((p-c, t))
    return enc

p = 862718293348820473429344482784628181556388621521298319395315527974911
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pkey, skey = genkey(p, l)
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H = pkey[l] ** 2 % p

print('H =', H)
print('enc =', enc)
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```
def encrypt(msg, pkey):
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```

WHAT DOES THE CODE DO ?

```
import random
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    R.sort()
    rbit = '1'
    for i in range(l-1):
        rbit += (R[i+1] - R[i] - 1) * '0' + '1'
    rbit += (nbit - R[-1] - 1) * '0'
    return int(rbit, 2)

def genkey(p, l):
    n = len(bin(p)[2:])
    f, skey = gen_rand(n, l), gen_rand(n, l)
    pkey = f * inverse(skey, p) % p
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def encrypt(msg, pkey):
    p, g = pkey
    msg, enc, n = bytes_to_long(msg), [], len(bin(p)[2:])
    for b in bin(msg)[2:]:
        s, t = gen_rand(n, l), gen_rand(n, l)
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            enc.append((p-c, t))
    return enc
```

```
p = 862718293348820473429344482784628181556388621521298319395315527974911
l = 5

pkey, skey = genkey(p, l)
enc = encrypt(flag, pkey)
H = pkey[l] ** 2 % p

print('H =', H)
print('enc =', enc)
```

```
p = 862718293348820473429344482784628181556388621521298319395315527974911
l = 5
```

```
pkey, skey = genkey(p, l)
enc = encrypt(flag, pkey)
H = pkey[l] ** 2 % p
```

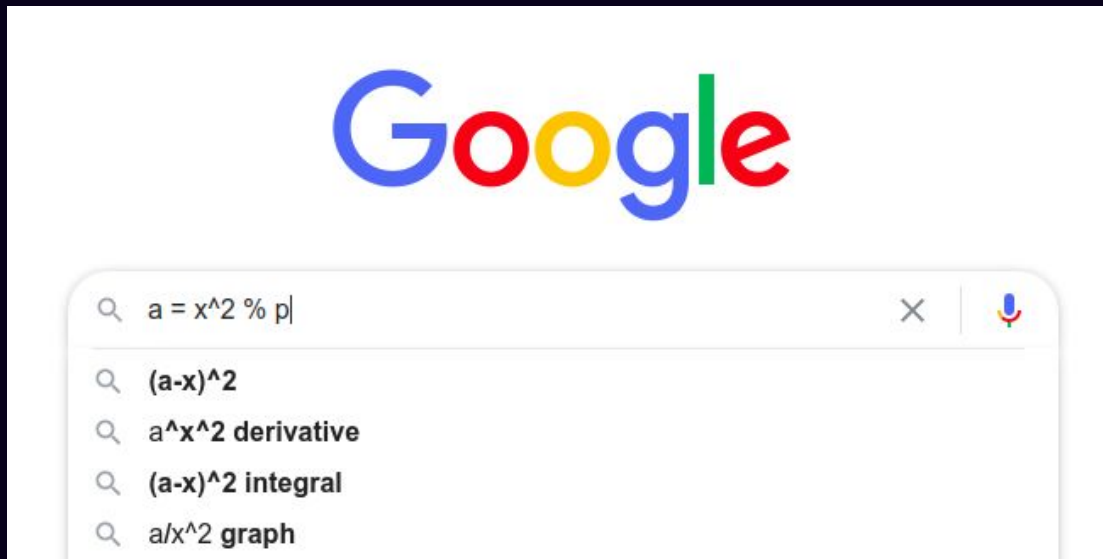
```
print('H =', H)
print('enc =', enc)
```

WHERE DO WE BEGIN ?

$H = \text{pkey}[1]^{** 2 \% p}$

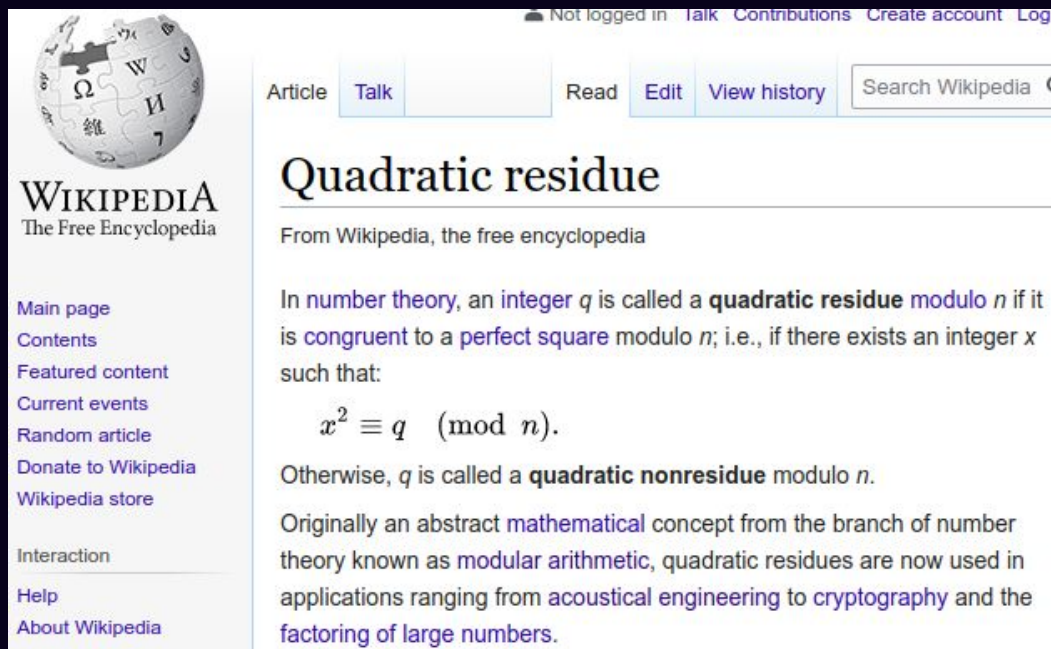
$H = 381704527450191606347421195235742637659723827441243208291869156144963$

$p = 862718293348820473429344482784628181556388621521298319395315527974911$



WHAT IS A QUADRATIC RESIDUE ?

$$H = \text{pkey}[1] ** 2 \% p \quad \Rightarrow \quad \text{pkey}^2 \equiv H \pmod{p}$$



The screenshot shows the Wikipedia page for "Quadratic residue". The page layout includes a sidebar on the left with links like "Main page", "Contents", and "Featured content". The main content area has a title "Quadratic residue" and a subtitle "From Wikipedia, the free encyclopedia". The text explains that in number theory, an integer q is a quadratic residue modulo n if it is congruent to a perfect square modulo n . It includes the equation $x^2 \equiv q \pmod{n}$ and mentions that otherwise, q is a quadratic nonresidue modulo n . The page also notes that quadratic residues are used in applications like acoustical engineering, cryptography, and factoring of large numbers.

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Quadratic residue

From Wikipedia, the free encyclopedia

In [number theory](#), an [integer](#) q is called a **quadratic residue modulo** n if it is [congruent](#) to a [perfect square](#) modulo n ; i.e., if there exists an integer x such that:

$$x^2 \equiv q \pmod{n}.$$

Otherwise, q is called a **quadratic nonresidue** modulo n .

Originally an abstract [mathematical](#) concept from the branch of number theory known as [modular arithmetic](#), quadratic residues are now used in applications ranging from [acoustical engineering](#) to [cryptography](#) and the [factoring of large numbers](#).

HOW TO SOLVE A QUADRATIC RESIDUE ?

Our p is not a prime number !

Factorization of p to factors p_i +

Find solutions to each factor: $pkey^2 \equiv H \pmod{p_i}$ +

Chinese remainder theorem to combine the solutions

```
p = 862718293348820473429344482784628181556388621521298319395315527974911
factors = [1504073, 20492753, 59833457464970183, 467795120187583723534280000348743236593]
```

HOW TO FIND SOLUTIONS FOR THE FACTORS ?

- 1) If a is a quadratic residue $(\bmod p)$ and $p \equiv 3 \pmod{4}$ then $a^{(p+1)/4}$ is a solution to $x^2 \equiv a \pmod{p}$.
- 2) In all other cases there is no analogous formula and one may use the Tonelli-Shanks algorithm.

TONELLI-SHANKS ALGORITHM

$x^2 \equiv H \pmod{p_i}$, where p_i are prime numbers, factors of p

```
p0 = 1504073:  
399666  
1104407
```

```
p1 = 20492753:  
7111848  
13380905
```

```
p2 = 59833457464970183:  
34240854883018057  
25592602581952126
```

```
p3 = 467795120187583723534280000348743236593:  
308269479959806774875048102517512730884  
15952564022776948659231897831230505709
```

```
def legendre(a, p):  
    return pow(a, (p - 1) // 2, p)  
  
def tonelli(n, p):  
    assert legendre(n, p) == 1, "not a square (mod p)"  
    q = p - 1  
    s = 0  
    while q % 2 == 0:  
        q //= 2  
        s += 1  
    if s == 1:  
        return pow(n, (p + 1) // 4, p)  
    for z in range(2, p):  
        if p - 1 == legendre(z, p):  
            break  
    c = pow(z, q, p)  
    r = pow(n, (q + 1) // 2, p)  
    t = pow(n, q, p)  
    m = s  
    t2 = 0  
    while (t - 1) % p != 0:  
        t2 = (t * t) % p  
        for i in range(1, m):  
            if (t2 - 1) % p == 0:  
                break  
            t2 = (t2 * t2) % p  
        b = pow(c, 1 << (m - i - 1), p)  
        r = (r * b) % p  
        c = (b * b) % p  
        t = (t * c) % p  
        m = i  
    return r
```

WHAT IS THE CHINESE REMAINDER THEOREM ?

Theorem statement [\[edit \]](#)

Let n_1, \dots, n_k be integers greater than 1, which are often called *moduli* or *divisors*. Let us denote by N the product of the n_i .

The Chinese remainder theorem asserts that if the n_i are *pairwise coprime*, and if a_1, \dots, a_k are integers such that $0 \leq a_i < n_i$ for every i , then there is one and only one integer x , such that $0 \leq x < N$ and the remainder of the *Euclidean division* of x by n_i is a_i for every i .

This may be restated as follows in term of *congruences*: If the n_i are pairwise coprime, and if a_1, \dots, a_k are any integers, then there exists an integer x such that

$$\begin{aligned} x &\equiv a_1 \pmod{n_1} \\ &\vdots \\ x &\equiv a_k \pmod{n_k}, \end{aligned}$$

and any two such x are congruent modulo N .^[12]

Looking for a solution of:

$$H \equiv pkey^2 \pmod{p}$$

Factors p_i must be pairwise *coprime*! $\Rightarrow \gcd(p_i, p_j) = 1$

HAVING TROUBLE APPLYING THE THEOREM...

```
from functools import reduce
def chinese_remainder(n, a):
    sum = 0
    prod = reduce(lambda a, b: a*b, n)
    for n_i, a_i in zip(n, a):
        p = prod // n_i
        sum += a_i * mul_inv(p, n_i) * p
    return sum % prod

def mul_inv(a, b):
    b0 = b
    x0, x1 = 0, 1
    if b == 1: return 1
    while a > 1:
        q = a // b
        a, b = b, a%b
        x0, x1 = x1 - q * x0, x0
    if x1 < 0: x1 += b0
    return x1
```

```
if __name__ == '__main__':
    pi = [1504073, 20492753, 59833457464970183, 467795120187583723534280000348743236593]

    r0 = [1104407, 13380905, 25592602581952126, 159525640227776948659231897831230505709]
    r1 = [1104407, 13380905, 25592602581952126, 308269479959806774875048102517512730884]
    r2 = [1104407, 13380905, 34240854883018057, 159525640227776948659231897831230505709]
    r3 = [1104407, 13380905, 34240854883018057, 308269479959806774875048102517512730884]
    r4 = [1104407, 7111848, 25592602581952126, 159525640227776948659231897831230505709]
    r5 = [1104407, 7111848, 25592602581952126, 308269479959806774875048102517512730884]
    r6 = [1104407, 7111848, 34240854883018057, 159525640227776948659231897831230505709]
    r7 = [1104407, 7111848, 34240854883018057, 308269479959806774875048102517512730884]
    r8 = [399666, 13380905, 25592602581952126, 159525640227776948659231897831230505709]
    r9 = [399666, 13380905, 25592602581952126, 308269479959806774875048102517512730884]
    r10 = [399666, 13380905, 34240854883018057, 159525640227776948659231897831230505709]
    r11 = [399666, 13380905, 34240854883018057, 308269479959806774875048102517512730884]
    r12 = [399666, 7111848, 25592602581952126, 159525640227776948659231897831230505709]
    r13 = [399666, 7111848, 25592602581952126, 308269479959806774875048102517512730884]
    r14 = [399666, 7111848, 34240854883018057, 159525640227776948659231897831230505709]
    r15 = [399666, 7111848, 34240854883018057, 308269479959806774875048102517512730884]

    rs = [r0, r1, r2, r3, r4, r5, r6, r7, r8, r9, r10, r11, r12, r13, r14, r15]
    for r in rs:
        print(chinese_remainder(pi, r))
```


RESULTS OF THE CHINESE REMAINDER THEOREM:

YESSS!

```
possible_pkeys = [  
739258514585797449032297222084055811470702691611125687711372074996687,  
460372941321907147479217537764873459531020265638440613169538863817034,  
785232755191570522054580268914288666635203565395845945113513087262967,  
506347181927680220501500584595106314695521139423160870571679876083314,  
86101158007865811093157347685756838828674540402077828886977707662969,  
582126006814767809378493792538386036347062978048093214327944496483316,  
44267527335610710524512040903173061894857656284200226876603191954338,  
628100247420540882400776839368618891511563851832813471730085508749596,  
234618045928279591028567643416009290044824769688484847665230019225315,  
818450766013209762904832441881455119661530965237098092518712336020573,  
280592286534052664050850690246242145209325643473205105067371031491595,  
1706713270162362497771005927059793269643217500520030525537820311942,  
356371111421140252927843898189521866860867482098137448823635651891597,  
77485538157249951374764213870339514921185056125452374281802440711944,  
402345352026913325950126945019754722025368355882857706225776664157877,  
123459778763023024397047260700572370085685929910172631683943452978224]
```

NEXT STEP AFTER FINDING THE PUBLIC KEY ...

```
import random
from Crypto.Util.number import *
from flag import flag

def gen_rand(nbit, l):
    R = []
    while True:
        r = random.randint(0, nbit-1)
        if r not in R:
            R.append(r)
            if len(R) == l:
                break
    R.sort()
    rbit = '1'
    for i in range(l-1):
        rbit += (R[i+1] - R[i] - 1) * '0' + '1'
    rbit += (nbit - R[-1] - 1) * '0'
    return int(rbit, 2)

def genkey(p, l):
    n = len(bin(p)[2:])
    f, skey = gen_rand(n, l), gen_rand(n, l)
    pkey = f * inverse(skey, p) % p
    return (p, pkey), skey
```

```
def encrypt(msg, pkey):
    p, g = pkey
    msg, enc, n = bytes_to_long(msg), [], len(bin(p)[2:])
    for b in bin(msg)[2:]:
        s, t = gen_rand(n, l), gen_rand(n, l)
        c = (s * g + t) % p
        if b == '0':
            enc.append((c, t))
        else:
            enc.append((p-c, t))
    return enc
```

```
p = 862718293348828473429344482784628181556388621521298319395315527974911
l = 5
```

```
pkey, skey = genkey(p, l)
enc = encrypt(flag, pkey)
H = pkey[l] ** 2 % p
```

```
print('H =', H)
print('enc =', enc)
```

... reversing $c = (s * g + t) \% p$

```
def decrypt(msg, pkey):
    p, g = pkey
    msg, enc, n = bytes_to_long(msg), [], len(bin(p)[2:])
    for b in bin(msg)[2:]:
        s, t = gen_rand(n, l), gen_rand(n, l)
        c = (s * g + t) % p
        if b == '0':
            enc.append((c, t))
        else:
            enc.append((p-c, t))
    return enc
```

```
enc=[
(693452030319839722726551032492658894345570855091353510812978453643631,
374147273652096600518280302026002167242456968463360),
(688559612388766711617272277684383103592686215888250125242406607896646,
1684996666696915012012216278925205994248717536439095543946901716996),
```


BRUTE FORCING 'S'?

```
c = (s * pkey + t) % p
(c - t) ≡ s * pkey (mod p)
(c - t) * pkey ≡ s * pkey^2 (mod p)
(c - t) * pkey ≡ s * H (mod p)    because H ≡ pkey^2 (mod p)
```

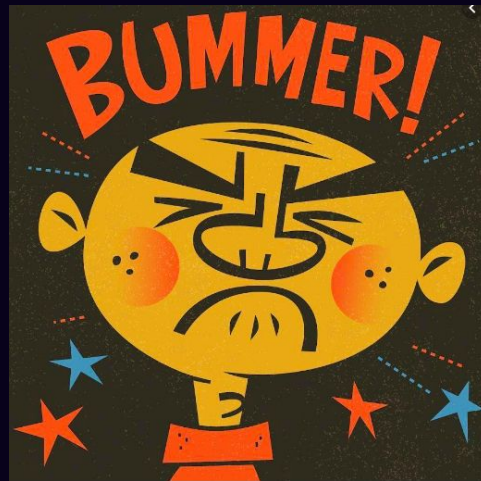
Rule:

$a \equiv b \pmod{p}$

$a = b + p * k$

```
s * H = (c - t) * pkey + p * k
s = ((c - t) * pkey + p * k) / H
```

```
def find_s():
    for k in range(100000000):
        s = ((c - t) * pkey + p * k) / H
        if(s == int(round(s))):
            print(s)
```



THE **PROPER** WAY TO REVERSE THE EXPRESSION !

```
c = (s * pkey + t) % p
(c - t) ≡ s * pkey (mod p)
(c - t) * pkey(-1) ≡ s * pkey * pkey(-1) (mod p)    pkey(-1) is modular inverse of pkey mod p
(c - t) * pkey(-1) ≡ s * 1 (mod p)

=> s = (c - t) * inverse(pkey, p) % p
```

Modular inverse:

$$ax \equiv 1 \pmod{m}$$

$$aa^{-1} \equiv 1 \pmod{m}$$

Note: modular inverse of a number **n mod m** exists **only** if **gcd(n, m) = 1**

The library I used "Crypto.Util.number" doesn't check this and might produce wrong answers. Better alternatives are: "gmpy2" and "sage".

FINAL EXPLOIT:

```
#!/usr/bin/env python3
enc = [...] # Given in the output.txt

# Calculated via the "Tonelli-Shanks algorithm" and the "Chinese remainder theorem"
possible_pkeys = [
739258514585797449032297222084055811470702691611125687711372074996687,
...
123459778763023024397047260700572370085685929910172631683943452978224]

H = 381704527450191606347421195235742637659723827441243208291869156144963
p = 862718293348820473429344482784628181556388621521298319395315527974911 # p = 2^229 - 1

from Crypto.Util.number import *

flag = ''

# Check all possible public keys, because only one of them is used for the encryption
for pkey in possible_pkeys:

    # For each flag try to reverse its bits
    for i in range(len(enc)):
        c = enc[i][0]
        t = enc[i][1]

        s = ((c - t) * inverse(pkey, p)) % p
        x = bin(s)[2:].count('1')
        if (x == 5):
            flag += '0'
            continue

        s = (((p - c) - t) * inverse(pkey, p)) % p
        x = bin(s)[2:].count('1')
        if (x == 5):
            flag += '1'
            continue

    # If we got a meaningful flag, reverse it.
    if (flag != ''):
        print(long_to_bytes(int(flag,2)))
        flag = ''
```

ASIS{T0y_3XampL3_w1tH__m3rSEnN3__nUmB3r5}

COUNTERMEASURES.

1. Choosing p to be prime, so that it can't be factored into smaller numbers.
2. Using a better random number generator.

THANK YOU

FOR YOUR ATTENTION !