

# Disjoint Sets

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# DisjointSets Data Structure

- ◆ Data structure for maintaining a partition of a set into disjoint subsets (data structure sometimes called *Partition* rather than DisjointSets)
- ◆ General features
  - *Data:*
    - ◆ Universe  $U$  – the base set that is being partitioned (this set is never altered)
    - ◆ Collection  $C = \{X_1, X_2, \dots, X_n\}$  of subsets of the universe – the subsets are disjoint and their union is  $U$  (these subsets are modified when the data structure is used – size of  $C$  shrinks because of repeated union operations)
  - *Operations:*
    - ◆  $\text{find}(x)$  – returns the subset  $X_i$  to which  $x$  belongs
    - ◆  $\text{union}(A, B)$  – replaces the subsets  $A, B$  in  $C$  with  $A \cup B$ .

# Array based Implementation of DisjointSets as Trees.

- ◆ The elements of each set  $X$  in the collection  $C$  are represented by nodes in a tree  $T_x$ ; the set  $X$  itself is referenced by its root  $r_x$ .
- ◆  $\text{find}(x)$  returns the root of the tree to which  $x$  belongs
- ◆  $\text{union}(x,y)$  joins the tree  $x$  belongs to to the tree  $y$  belongs to by pointing root of one to the root of the other.

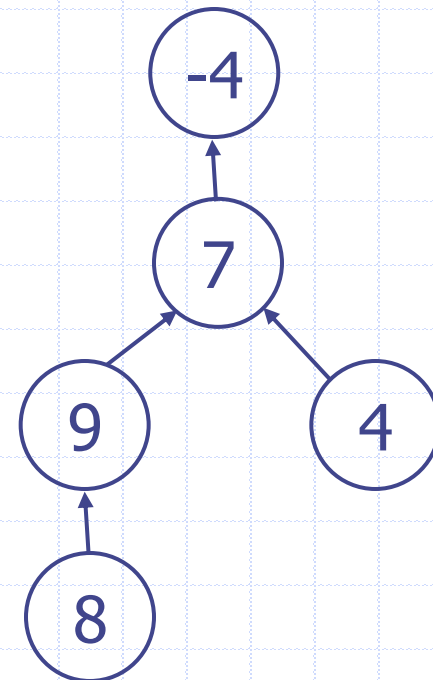
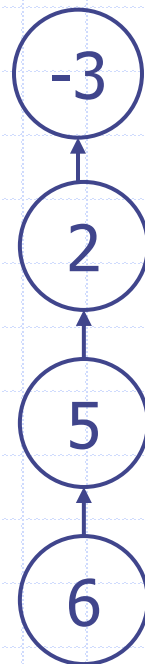
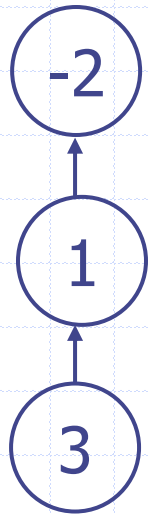
# Array Implementation

1	2	3	4	5	6	7	8	9
-2	-3	1	7	2	5	-4	9	7

$U = \{1, 2, 3, 4, 5, 6, 7, 8, 9\}$

$C = \{\{1, 3\}, \{2, 5, 6\}, \{4, 7, 8, 9\}\}$

Tree representations:



# Example

$$U = \{1, 2, 3, 4, 5, 6, 7, 8, 9\}$$

$$C = \{\{1, 3\}, \{2, 5, 6\}, \{4, 7, 8, 9\}\}$$

$$\text{Find}(1) = \text{Find}(3) = 1. \text{ (The root)}$$

$\text{Arr}[1] = -2$ . Here 2 denote number of items. Also, the negative sign indicates end of the tree.

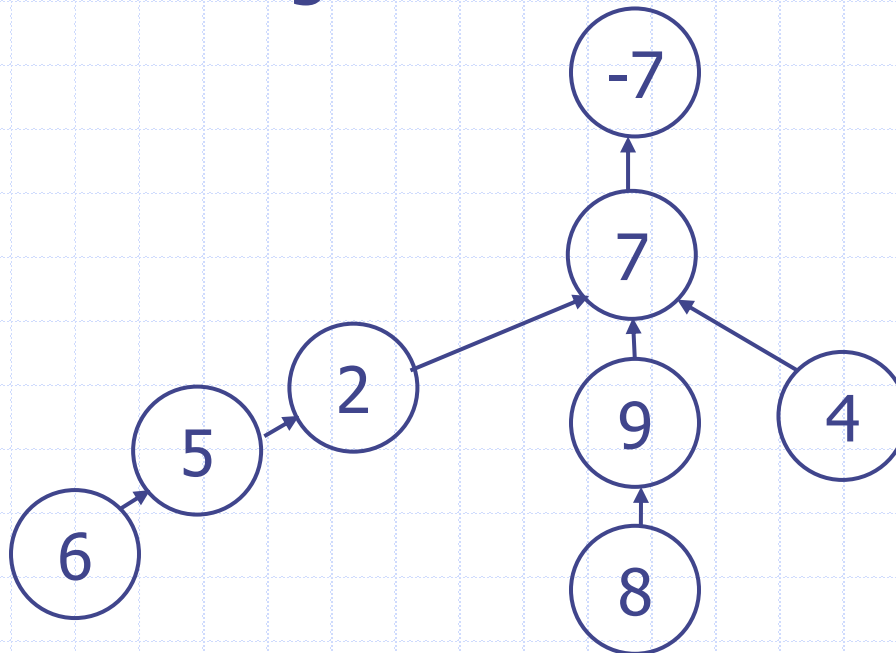
$$\text{Find}(2) = \text{Find}(5) = \text{Find}(6) = 2$$

$$\text{Find}(4) = \text{Find}(7) = \text{Find}(8) = \text{Find}(9) = 7$$

1	2	3	4	5	6	7	8	9
-2	7	1	7	2	5	-7	9	7

Union(6, 8)

- Make the "root" of the smaller tree point to root of larger tree.
- Change the value of number of items



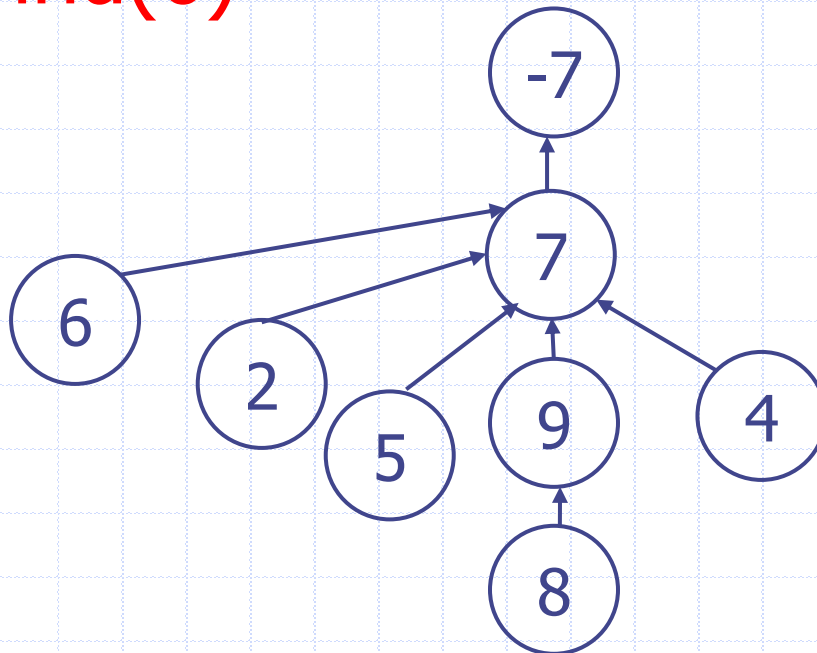
Find(5) is 7

1	2	3	4	5	6	7	8	9
-2	7	1	7	7	7	-7	9	7

## Find with path compression

All elements from the “findItem” to the root of the tree will point to the root of the tree.

Find(6)



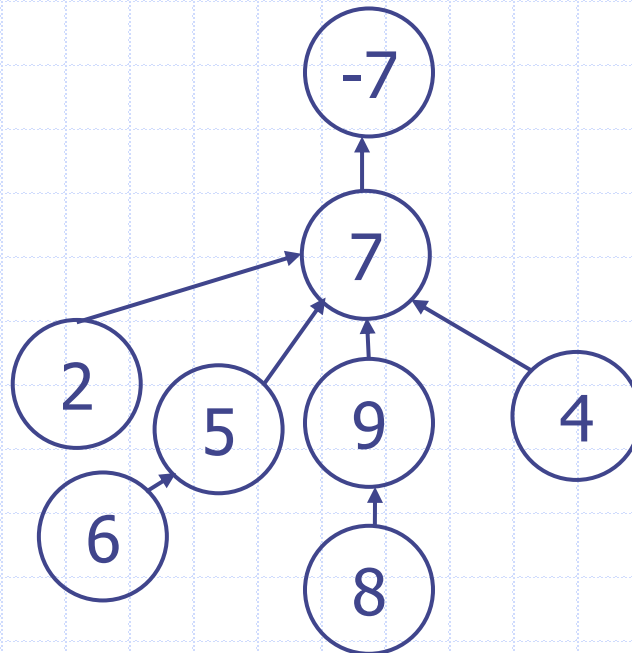
Note that 6 and 5  
pointing to 7.

1	2	3	4	5	6	7	8	9
-2	7	1	7	7	5	-7	9	7

## Find with path compression

All elements from the "find item" to the root of the tree will point to the root of the tree.

**Find(5)** (Instead of Find(6))



Note that 6 still points to 5 since 6 was "before" 5.



# Summary

Thus, the disjoint sets are implemented using an 1-dimensional array. With path compression, we can almost achieve  $O(1)$  time complexity for union and find operations.

In this example, index 0 is not used. If elements in the sets are 0, 1, ..., then we can use index location 0 as well.