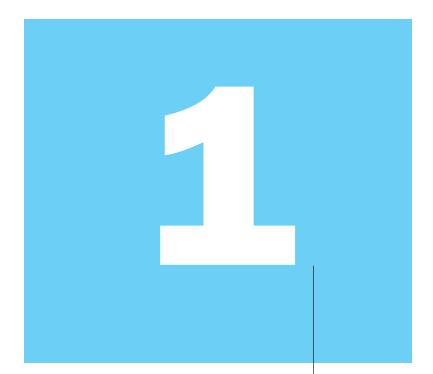
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# **Solutions**

**1.1 Personal computer:** Computer that emphasizes delivery of good performance to a single user at low cost and usually executes third-party software.

**Server:** Computer used for large workloads and usually accessed via a network.

**Embedded computer:** Computer designed to run one application or one set of related applications and integrated into a single system.

#### 1.2

- a. Performance via Pipelining
- **b.** Dependability via Redundancy
- c. Performance via Prediction
- d. Make the Common Case Fast
- e. Hierarchy of Memories
- f. Performance via Parallelism
- g. Use Abstraction to Simplify Design
- **1.3** The program is compiled into an assembly language program, which is itself assembled into a machine-language program.

## 1.4

- **a.**  $1280 \times 1024$  pixels = 1,310,720 pixels => 1,310,720  $\times$  3 = 3,932,160 bytes/frame.
- **b.**  $3,932,160 \text{ bytes} \times (8 \text{ bits/byte}) / 100E6 \text{ bits/second} = 0.31 \text{ seconds}$

# 1.5

Desktop Processor	Year	Tech	Max. Clock Speed (GHz)	Integer IPC/core	Cores	Max. DRAM Bandwidth (GB/s)	SP Floating Point (Gflop/s)	L3 cache (MiB)
Westmere i7-620	2100	32	3.33	4	2	17.1	107	4
lvy Bridge i7-3770K	2013	22	3.90	6	4	25.6	250	8
Broadwell i7-6700K	2015	14	4.20	8	4	34.1	269	8
Kaby Lake i7-7700K	2017	14	4.50	8	4	38.4	288	8
Coffee Lake i7-9700K	2019	14	4.90	8	8	42.7	627	12
Imp./year	·	20%	4%	7%	15%	10%	19%	12%
Doubles every		4 years	18 years	10 years	5 years	7 years	4 years	6 years

## 1.6

a. performance of P1 (instructions/sec) = 
$$3 \times 10^9/1.5 = 2 \times 10^9$$
  
performance of P2 (instructions/sec) =  $2.5 \times 10^9/1.0 = 2.5 \times 10^9$   
performance of P3 (instructions/sec) =  $4 \times 10^9/2.2 = 1.8 \times 10^9$ 

**b.** cycles(P1) = 
$$10 \times 3 \times 10^9 = 30 \times 10^9$$
 s  
cycles(P2) =  $10 \times 2.5 \times 10^9 = 25 \times 10^9$  s  
cycles(P3) =  $10 \times 4 \times 10^9 = 40 \times 10^9$  s

**c.** No. instructions(P1) = 
$$30 \times 10^9 / 1.5 = 20 \times 10^9$$

No. instructions(P2) = 
$$25 \times 10^9/1 = 25 \times 10^9$$

No. instructions(P3) = 
$$40 \times 10^9/2.2 = 18.18 \times 10^9$$

$$CPI_{new} = CPI_{old} \times 1.2$$
, then  $CPI(P1) = 1.8$ ,  $CPI(P2) = 1.2$ ,  $CPI(P3) = 2.6$ 

$$f$$
 = No. instr. × CPI/time, then

$$f(P1) = 20 \times 10^9 \times 1.8/7 = 5.14 \,\text{GHz}$$

$$f(P2) = 25 \times 10^9 \times 1.2/7 = 4.28 \,\text{GHz}$$

$$f(P3) = 18.18 \times 10^9 \times 2.6/7 = 6.75 \,\text{GHz}$$

# 1.7

**a.** Class A:  $10^5$  instr. Class B:  $2 \times 10^5$  instr. Class C:  $5 \times 10^5$  instr. Class D:  $2 \times 10^5$  instr.

Time = No. instr.  $\times$  CPI/clock rate

Total time P1 = 
$$(10^5 + 2 \times 10^5 \times 2 + 5 \times 10^5 \times 3 + 2 \times 10^5 \times 3)/(2.5 \times 10^9) = 10.4 \times 10^{-4} \text{ s}$$

Total time P2 =  $(10^5 \times 2 + 2 \times 10^5 \times 2 + 5 \times 10^5 \times 2 + 2 \times 10^5 \times 2)/(3 \times 10^9) = 6.66 \times 10^{-4} \text{ s}$ 

$$CPI(P1) = 10.4 \times 10^{-4} \times 2.5 \times 10^{9}/10^{6} = 2.6$$

$$CPI(P2) = 6.66 \times 10^{-4} \times 3 \times 10^{9} / 10^{6} = 2.0$$

**b.** clock cycles(P1) =  $10^5 \times 1 + 2 \times 10^5 \times 2 + 5 \times 10^5 \times 3 + 2 \times 10^5 \times 3 = 26 \times 10^5$ clock cycles(P2) =  $10^5 \times 2 + 2 \times 10^5 \times 2 + 5 \times 10^5 \times 2 + 2 \times 10^5 \times 2 = 20 \times 10^5$ 

**S-4** 

# 1.8

**a.** CPI = 
$$T_{exec} \times f/No.$$
 instr.

Compiler A CPI = 1.1

Compiler B CPI = 1.25

**b.** 
$$f_B/f_A = (\text{No. instr.}(B) \times \text{CPI}(B))/(\text{No. instr.}(A) \times \text{CPI}(A)) = 1.37$$

**c.** 
$$T_A/T_{new} = 1.67$$

$$T_{\rm R}/T_{\rm new} = 2.27$$

# 1.9

**1.9.1** 
$$C = 2 \times DP/(V^2 \times F)$$

Pentium 4: C = 3.2E - 8F

Core i5 Ivy Bridge: C = 2.9E-8F

**1.9.2** Pentium 4: 10/100 = 10%

Core i5 Ivy Bridge: 30/70 = 42.9%

**1.9.3** 
$$(S_{new} + D_{new})/(S_{old} + D_{old}) = 0.90$$

$$D_{new} = C \times V_{new} \times 2 \times F$$

$$S_{old} = V_{old} \times I$$

$$S_{new} = V_{new} \times I$$

Therefore:

$$V_{\text{new}} = [D_{\text{new}}/(C \times F)]1/2$$

$$D_{new} = 0.90 \times (S_{old} + D_{old}) - S_{new}$$

$$S_{new} = V_{new} \times (S_{old}/V_{old})$$

Pentium 4:

$$S_{new} = V_{new} \times (10/1.25) = V_{new} \times 8$$

$$D_{new} = 0.90 \times 100 - V_{new} \times 8 = 90 - V_{new} \times 8$$

$$V_{new} = [(90 - V_{new} \times 8)/(3.2E8 \times 3.6E9)]^{1/2}$$

$$V_{new} = 0.85 V$$

Core i5:

$$S_{pew} = V_{pew} \times (30/0.9) = V_{pew} \times 33.3$$

$$D_{new} = 0.90 \times 70 - V_{new} \times 33.3 = 63 - V_{new} \times 33.3$$

$$V_{\text{new}} = [(63 - V_{\text{new}} \times 33.3)/(2.9E8 \times 3.4E9)]^{1/2}$$

$$V_{new} = 0.64 V$$

## Chapter 1 Solutions

**S-6** 

р	# arith inst.	# L/S inst.	# branch inst.	cycles	ex. time	speedup
1	2.56E9	1.28E9	2.56E8	1.92E10	9.60	1.00
2	1.83E9	9.14E8	2.56E8	1.41E10	7.04	1.36
4	9.14E8	4.57E8	2.56E8	7.68E9	3.84	2.50
8	4.57E8	2.29E8	2.56E8	4.48E9	2.24	4.29

# 1.10

#### 1.10.1

р	ex. time
1	41.0
2	29.3
4	14.6
8	7.33

1.10.2

1.10.3 3

#### 1.11

**1.11.1** die area<sub>15cm</sub> = wafer area/dies per wafer = 
$$\pi \times 7.5^2/84 = 2.10 \text{ cm}^2$$
  
yield<sub>15cm</sub> =  $1/(1 + (0.020 \times 2.10/2))^2 = 0.9593$   
die area<sub>20cm</sub> = wafer area/dies per wafer =  $\pi \times 10^2/100 = 3.14 \text{ cm}^2$   
yield<sub>20cm</sub> =  $1/(1 + (0.031 \times 3.14/2))^2 = 0.9093$ 

1.11.2 
$$\operatorname{cost/die}_{15\text{cm}} = 12/(84 \times 0.9593) = 0.1489$$
  
 $\operatorname{cost/die}_{20\text{cm}} = 15/(100 \times 0.9093) = 0.1650$ 

1.11.3 die area<sub>15cm</sub> = wafer area/dies per wafer = 
$$\pi \times 7.5^2/(84 \times 1.1) = 1.91 \text{ cm}^2$$
  
yield<sub>15cm</sub> =  $1/(1 + (0.020 \times 1.15 \times 1.91/2))^2 = 0.9575$   
die area<sub>20cm</sub> = wafer area/dies per wafer =  $\pi \times 10^2/(100 \times 1.1) = 2.86 \text{ cm}^2$   
yield<sub>20cm</sub> =  $1/(1 + (0.03 \times 1.15 \times 2.86/2))^2 = 0.9082$ 

**1.11.4** defects per area<sub>0.92</sub> = 
$$(1-y^{.5})/(y^{.5} \times \text{die\_area/2}) = (1 - 0.92^{.5})/(0.92^{.5} \times 2/2) = 0.043 \text{ defects/cm}^2$$

defects per area<sub>0.95</sub> = 
$$(1-y^{.5})/(y^{.5} \times \text{die\_area/2}) = (1 - 0.95^{.5})/(0.95^{.5} \times 2/2) = 0.026 \text{ defects/cm}^2$$

## 1.12

**1.12.1** CPI = clock rate  $\times$  CPU time/instr. count

clock rate = 1/cycle time = 3 GHz

 $CPI(bzip2) = 3 \times 10^9 \times 750/(2389 \times 10^9) = 0.94$ 

**1.12.2** SPEC ratio = ref. time/execution time

SPEC ratio(bzip2) = 9650/750 = 12.86

1.12.3 CPU time = No. instr.  $\times$  CPI/clock rate

If CPI and clock rate do not change, the CPU time increase is equal to the increase in the number of instructions, that is 10%.

**1.12.4** CPU time(before) = No. instr.  $\times$  CPI/clock rate

CPU time(after) =  $1.1 \times \text{No. instr.} \times 1.05 \times \text{CPI/clock}$  rate

CPU time(after)/CPU time(before) =  $1.1 \times 1.05 = 1.155$ . Thus, CPU time is increased by 15.5%.

**1.12.5** SPECratio = reference time/CPU time

SPECratio(after)/SPECratio(before) = CPU time(before)/CPU time(after) = 1/1.1555 = 0.86. The SPECratio is decreased by 14%.

**1.12.6** CPI = (CPU time  $\times$  clock rate)/No. instr.

$$CPI = 700 \times 4 \times 10^{9}/(0.85 \times 2389 \times 10^{9}) = 1.37$$

1.12.7 Clock rate ratio =  $4 \,\text{GHz}/3 \,\text{GHz} = 1.33$ 

They are different because, although the number of instructions has been reduced by 15%, the CPU time has been reduced by a lower percentage.

- **1.12.8** 700/750 = 0.933. CPU time reduction: 6.7%
- **1.12.9** No. instr. = CPU time  $\times$  clock rate/CPI

No. instr. = 
$$960 \times 0.9 \times 4 \times 10^9 / 1.61 = 2146 \times 10^9$$

**1.12.10** Clock rate = No. instr.  $\times$  CPI/CPU time.

Clock rate 
$$_{\text{new}}$$
 = No. instr. × CPI/0.9 × CPU time = 1/0.9 clock rate  $_{\text{old}}$  = 4.44 GHz

**1.12.11** Clock rate = No. instr.  $\times$  CPI/CPU time.

Clock rate  $_{new}$  = No. instr. × 0.85 × CPI/0.80 CPU time = 0.85/0.80, clock rate  $_{old}$  = 3.18 GHz

## 1.13

**1.13.1** 
$$T(P1) = 5 \times 10^9 \times 0.9/(4 \times 10^9) = 1.125 s$$

$$T(P2) = 10^9 \times 0.75/(3 \times 10^9) = 0.25 s$$

clock rate(P1) > clock rate(P2), performance(P1) < performance(P2)

**1.13.2**  $T(P1) = No. instr. \times CPI/clock rate$ 

T(P1) = 2.2531021 s

 $T(P2) 5 N \times 0.75/(3 \times 10^9)$ , then  $N = 9 \times 10^8$ 

1.13.3 MIPS = Clock rate  $\times$  10<sup>-6</sup>/CPI

MIPS(P1) = 
$$4 \times 10^9 \times 10^{-6}/0.9 = 4.44 \times 10^3$$

MIPS(P2) = 
$$3 \times 10^9 \times 10^{-6} / 0.75 = 4.0 \times 10^3$$

MIPS(P1) > MIPS(P2), performance(P1) < performance(P2) (from 11a)

**1.13.4** MFLOPS = No. FP operations  $\times 10^{-6}$ /T

$$MFLOPS(P1) = .4 \times 5E9 \times 1E-6/1.125 = 1.78E3$$

$$MFLOPS(P2) = .4 \times 1E9 \times 1E-6/.25 = 1.60E3$$

MFLOPS(P1) > MFLOPS(P2), performance(P1) < performance(P2) (from 11a)

# 1.14

**1.14.1** 
$$T_{fp} = 70 \times 0.8 = 56 \text{ s. } T_{new} = 56 + 85 + 55 + 40 = 236 \text{ s. Reduction: } 5.6\%$$

**1.14.2** 
$$T_{\text{new}} = 250 \times 0.8 = 200 \text{ s}, T_{\text{fp}} + T_{\text{l/s}} + T_{\text{branch}} = 165 \text{ s}, T_{\text{int}} = 35 \text{ s}. \text{ Reduction time INT: } 58.8\%$$

**1.14.3** 
$$T_{new} = 250 \times 0.8 = 200 \text{ s}, T_{fp} + T_{int} + T_{1/s} = 210 \text{ s}. \text{ NO}$$

# 1.15

**1.15.1** Clock cycles =  $\mathrm{CPI}_{\mathrm{fp}} \times \mathrm{No.}$  FP instr. +  $\mathrm{CPI}_{\mathrm{int}} \times \mathrm{No.}$  INT instr. +  $\mathrm{CPI}_{\mathrm{l/s}} \times \mathrm{No.}$  L/S instr. +  $\mathrm{CPI}_{\mathrm{branch}} \times \mathrm{No.}$  branch instr.

 $T_{CPIJ}$  = clock cycles/clock rate = clock cycles/2 × 10<sup>9</sup>

clock cycles = 
$$512 \times 10^6$$
;  $T_{CPU} = 0.256$  s

To have the number of clock cycles by improving the CPI of FP instructions:

 $\begin{aligned} & \text{CPI}_{\text{improved fp}} \times \text{No. FP instr.} + \text{CPI}_{\text{int}} \times \text{No. INT instr.} + \text{CPI}_{\text{l/s}} \times \text{No. L/S instr.} + \\ & \text{CPI}_{\text{branch}} \times \text{No. branch instr.} = \text{clock cycles/2} \end{aligned}$ 

 $\begin{aligned} & \text{CPI}_{\text{improved fp}} = \left( \text{clock cycles/2} - \left( \text{CPI}_{\text{int}} \times \text{No. INT instr.} + \text{CPI}_{\text{l/s}} \times \text{No. L/S} \right. \\ & \text{instr.} + \text{CPI}_{\text{branch}} \times \text{No. branch instr.} \right) / \text{No. FP instr.} \end{aligned}$ 

$$CPI_{improved fp} = (256 - 462)/50 < 0 = = > not possible$$

**1.15.2** Using the clock cycle data from a.

To have the number of clock cycles improving the CPI of L/S instructions:

$$\begin{aligned} & \text{CPI}_{\text{fp}} \times \text{No. FP instr.} + \text{CPI}_{\text{int}} \times \text{No. INT instr.} + \text{CPI}_{\text{improved l/s}} \times \text{No. L/S instr.} \\ & + \text{CPI}_{\text{branch}} \times \text{No. branch instr.} = \text{clock cycles/2} \end{aligned}$$

 $\begin{aligned} & \text{CPI}_{_{\text{improved }I/s}} = (\text{clock cycles/2} - (\text{CPI}_{_{\text{fp}}} \times \text{No. FP instr.} + \text{CPI}_{_{\text{int}}} \times \text{No. INT} \\ & \text{instr.} + \text{CPI}_{_{\text{branch}}} \times \text{No. branch instr.})) / \text{No. L/S instr.} \end{aligned}$ 

$$CPI_{improved \, l/s} = (256 - 198)/80 = 0.725$$

**1.15.3** Clock cycles =  $CPI_{fp} \times No. FP instr. + CPI_{int} \times No. INT instr. + <math>CPI_{l/s} \times No. L/S instr. + CPI_{branch} \times No. branch instr.$ 

 $T_{CPU} = clock \ cycles/clock \ rate = clock \ cycles/2 \times 10^9$ 

$$\text{CPI}_{\text{int}} = 0.6 \times 1 = 0.6; \text{CPI}_{\text{fp}} = 0.6 \times 1 = 0.6; \text{CPI}_{\text{l/s}} = 0.7 \times 4 = 2.8; \text{CPI}_{\text{branch}} = 0.7 \times 2 = 1.4$$

$$T_{_{\mathrm{CPU}}}$$
 (before improv.) = 0.256 s;  $T_{_{\mathrm{CPU}}}$  (after improv.) = 0.171 s

# 1.16

processors	exec. time/ processor	time w/overhead	speedup	actual speedup/ideal speedup
1	100			
2	50	54	100/54 = 1.85	1.85/2 = .93
4	25	29	100/29 = 3.44	3.44/4 = 0.86
8	12.5	16.5	100/16.5 = 6.06	6.06/8 = 0.75
16	6.25	10.25	100/10.25 = 9.76	9.76/16 = 0.61