# DTU Course 02156 Logical Systems and Logic Programming (2021)

Week	Date	Main Topics (Prolog Programming in All Lessons)
35 #01	31/8	Course Prerequisites & Tutorial on Logical Systems and Logic Programming
36 #02	7/9	Chapter 1 - Introduction (Prolog Note)
37 #03	14/9	Chapter 2 - Propositional Logic: Formulas, Models, Tableaux
38 #04	21/9	Chapter 3 - Propositional Logic: Deductive Systems
39 #05	28/9	"Isabelle" - Propositional Logic: Sequent Calculus Verifier (SeCaV)
40 #06	5/10	Chapter 4 - Propositional Logic: Resolution
41 #07	12/10	Chapter 7 - First-Order Logic: Formulas, Models, Tableaux
42		(Autumn Vacation)
43 #08	26/10	Chapter 8 - First-Order Logic: Deductive Systems
44 #09	2/11	"Isabelle" - First-Order Logic: Sequent Calculus Verifier (SeCaV)
45 #10	9/11	Chapter 9 - First-Order Logic: Terms and Normal Forms
46 #11	16/11	Chapter 10 - First-Order Logic: Resolution
47 #12	23/11	Chapter 11 - First-Order Logic: Logic Programming
48 #13	30/11	Chapter 12 - First-Order Logic: Undecidability and Model Theory & Course Evaluation

Responsible: Associate Professor Jørgen Villadsen <jovi@dtu.dk>

## **Assignments & Exam**

## MUST BE SOLVED INDIVIDUALLY

Assignment-1 Deadline Sunday 26/9 (Available Wednesday 15/9)

Assignment-2 Deadline Sunday 10/10 (Available Wednesday 29/9)

Assignment-3 Deadline Sunday 31/10 (Available Wednesday 13/10)

Assignment-4 Deadline Sunday 14/11 (Available Wednesday 3/11)

Assignment-5 Deadline Thursday 2/12 (Available Wednesday 17/11)

Written Exam Tuesday 14/12 (2 Hours / No Computer / All Notes Allowed)

The mandatory assignments and the written exam are evaluated as a whole – even if you do well in the mandatory assignments then you still must do decent in the written exam in order to pass the course!

A TEACHER MUST IMMEDIATELY REPORT ANY SUSPICION OF CHEATING TO THE STUDY ADMINISTRATION FOR FURTHER ACTIONS



### **Sequent Calculus Verifier**

SeCaV formalizes first-order logic with constants and functions

SeCaV verifies one-sided sequent calculus proofs

SeCaV uses the Isabelle proof assistant

SeCaV is a tool for teaching logic

Jørgen Villadsen

Asta Halkjær From

Alexander Birch Jensen

**Anders Schlichtkrull** 

#### Agenda — Week #5

Motivation — Logic — Is it important...???

Isabelle Proof Assistant

Sequent Calculus Verifier (SeCaV)

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\* Predicate logic and naive set theory as first symbolic logics

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Alan Turing 1912-1954

⋆ Computability – Homosexual – Ate apple with cyanide – Suicide?

FLoC is a large research event in computer science (every 4 years)

1996 New York, USA
1999 Trento, Italy
2002 Copenhagen
2006 Seattle, USA
2010 Edinburgh, UK
2014 Vienna, Austria

Logicians: Informatics, Mathematics, Linguistics, Philosophy, etc.

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Oxford, UK, 6-19 July 2018

http://www.floc2018.org/

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More than 2000 participants...

#### Isabelle Proof Assistant

Isabelle is a generic proof assistant

http://isabelle.in.tum.de/

The Archive of Formal Proofs is a collection of proof libraries, examples, and larger scientific developments, mechanically checked in the theorem prover Isabelle

http://www.isa-afp.org/

The world's first operating-system kernel with an end-to-end proof of implementation correctness and security enforcement is now open source

http://sel4.systems/

480000 lines of Isabelle source files = 9 Gutenberg bibles

Tableaux are used...

theory Scratch imports Main begin

theorem "p \<or> \<not>p" by blast

end

#### Wikipedia

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Each conditional tautology is inferred from other conditional tautologies on earlier lines in a formal argument according to rules and procedures of inference, giving a better approximation to the style of natural deduction used by mathematicians than David Hilbert's earlier style of formal logic where every line was an unconditional tautology.

Tableaux turned "tree upside down and signs reversed":

$$\begin{array}{lll}
\neg((p\lor q)\to (q\lor p)) \\
p\lor q, \neg(q\lor p) \\
\hline
P\lor q, \neg q, \neg p \\
\hline
L R & 2. \vdash \neg q, q, p & Axiom \\
p, \neg q, \neg p & 3. \vdash \neg(p\lor q), q, p & \beta\lor, 1, 2 \\
\times & 4. \vdash \neg(p\lor q), (q\lor p) & \alpha\lor, 3 \\
R & 5. \vdash (p\lor q)\to (q\lor p) & \alpha\to, 4 \\
q, \neg q, \neg p & \times
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Hence 
$$\vdash (A \lor B) \to (B \lor A)$$

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**Soundness and Completeness Theorem:**  $\models A$  **iff**  $\vdash A$ 

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### **Soundness and Completeness Theorem:** $\models A$ **iff** $\vdash A$

Note: Use implicit set formation for conclusions in the system  $\mathcal G$