PART I

Task Control Block Linked List

Following the previous homework (HW1), please add some code to the μ C/OS-II scheduler in the kernel level to observe the operations of the task control block (TCB) and TCB linked list.

The screenshot results. (10%)

```
D:\ntust\EmbeddingOS\RTOS_M11107309_PA1\RTOS_M
         created, Thread ID 17316
OSTick
Task[63] created, TCB address b2ffa0
-----After TCB[63] being linked-----
Previous TCB point to address
                                   0
                              b2ffa0
Current TCB point to address
        TCB point to address
Next
                                   0
Task[ 1] created, TCB address b2fffc
-----After TCB[ 1] being linked-----
Previous TCB point to address
                                   0
Current TCB point to address
                              b2fffc
        TCB point to address
                              b2ffa0
Next
Task[ 2] created, TCB address b30058
-----After TCB[ 2] being linked-----
Previous TCB point to address
Current TCB point to address
                              b30058
Next
        TCB point to address
                              b2fffc
TCB addr
                                 Next TCB addr
Task
       Prev TCB addr
                 0
                        b30058
                                      b2fffc
2
                        b2fffc
1
            b30058
                                      b2ffa0
63
            b2fffc
                        b2ffa0
                                           0
```

A report that describes your implementation (please attach the screenshot of the code and MARK the modified part). (10%)

(如未標示,整段即為全新 code)

os_cpu_c.c

建立 task 時會進行 trace 並輸出訊息,將其更改為題目要求的 TCB address。

os_core.c

OSTCBInit()

```
OSTCBList = ptcb;

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printf("-----After TCB[%2d] being linked-----\n", prio);

2454

printf("Previous TCB point to address\t%6x\n", ptcb->OSTCBPrev);

printf("Current TCB point to address\t%6x\n", ptcb);

printf("Next TCB point to address\t%6x\n\n", ptcb->OSTCBNext);

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printf("Next TCB point to address\t%6x\n\n", ptcb->OSTCBNext);

/// PA#1 part1

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OSRdyGrp |= ptcb->OSTCBBitY; /* Make task ready to run

OSRdyTbl[ptcb->OSTCBY] |= ptcb->OSTCBBitX;
```

OSStart()

任務建立好後在 OSStart,呼叫 ptcb,輸出 TCB list,利用 ptcb->OSTCBNext 與while,歷遍 TCB list。

PART II

RM Scheduler Implementation

To implement the Rate Monotonic (RM) scheduler for periodic tasks and observe the scheduling behaviors.

```
Periodic Task Set = \{\tau_{ID} \text{ (ID, arrival time, execution time, period)}\}\

Example Task Set 1 = \{\tau_{I} \text{ (1, 1, 2, 4), } \tau_{I} \text{ (2, 0, 4, 10)}\}\

Example Task Set 2 = \{\tau_{I} \text{ (1, 3, 4, 14), } \tau_{I} \text{ (2, 0, 2, 8), } \tau_{I} \text{ (3, 0, 4, 10), } \tau_{I} \text{ (4, 24, 2, 12)}\}\

Example Task Set 3 = \{\tau_{I} \text{ (1, 2, 2, 10), } \tau_{I} \text{ (2, 1, 1, 5), } \tau_{I} \text{ (3, 0, 8, 15)}\}\
```

The correctness of schedule results of examples. Note the testing task set might not be the same as the given example task set. (40%)

Task Set 1

ick	Event	CreateTask ID	NextTask ID	ResponseTime	#of ContextSwitch	PreemptionTime	OSTimeDly
	Preemption	task(2)(0)	task(1)(0)				
	Completion	task(1)(0)	task(2)(0)				
5	Preemption	task(2)(0)	task(1)(1)				
	Completion	task(1)(1)	task(2)(0)				
8	Completion	task(2)(0)	task(63)	8			
	Preemption	task(63)	task(1)(2)				
1	Completion	task(1)(2)	task(2)(1)				
	Preemption	task(2)(1)	task(1)(3)				
	Completion	task(1)(3)	task(2)(1)				
7	Completion	task(2)(1)	task(1)(4)			2/t	
)	Completion	task(1)(4)	task(63)			0	
9	Preemption	task(63)	task(2)(2)				
	Preemption	task(2)(2)	task(1)(5)				
	Completion	task(1)(5)	task(2)(2)				
	Preemption	task(2)(2)	task(1)(6)				
	Completion	task(1)(6)	task(2)(2)				
3	Completion	task(2)(2)	task(63)				
9	Preemption	task(63)	task(1)(7)				

Task Set 2

Tick	Event	CreateTask ID	NextTask ID	ResponseTime	#of ContextSwitch	PreemptionTime	OSTimeDly
	Completion	task(2)(0)	task(3)(0)				
	Completion	task(3)(0)	task(4)(0)				
	Completion	task(4)(0)	task(1)(0)				
	Preemption	task(1)(0)	task(2)(1)				
10	Completion	task(2)(1)	task(3)(1)				
14	Completion	task(3)(1)	task(1)(0)				
16	Completion	task(1)(0)	task(2)(2)	16			
18	Completion	task(2)(2)	task(1)(1)				
20	Preemption	task(1)(1)	task(3)(2)				
24	Completion	task(3)(2)	task(2)(3)				
26	Completion	task(2)(3)	task(4)(1)				
28	Completion	task(4)(1)	task(1)(1)				
30	Completion	task(1)(1)	task(3)(3)	13			

Task Set 3

Tick	Event	CreateTask ID	NextTask ID	ResponseTime	#of ContextSwitch	PreemptionTime	OSTimeDly
1	Preemption	task(3)(0)	task(2)(0)				
2	Completion	task(2)(0)	task(1)(0)				4
4	Completion	task(1)(0)	task(3)(0)				8
6	Preemption	task(3)(0)	task(2)(1)				
7	Completion	task(2)(1)	task(3)(0)				4
11	Preemption	task(3)(0)	task(2)(2)				
12	Completion	task(2)(2)	task(1)(1)				4
14	Completion	task(1)(1)	task(3)(0)				8
15	Completion	task(3)(0)	task(3)(1)				0
16	Preemption	task(3)(1)	task(2)(3)				
17	Completion	task(2)(3)	task(3)(1)				4
21	Preemption	task(3)(1)	task(2)(4)				
22	Completion	task(2)(4)	task(1)(2)				4
24	Completion	task(1)(2)	task(3)(1)				8
26	Preemption	task(3)(1)	task(2)(5)				
27	Completion	task(2)(5)	task(3)(1)				4
28	Completion	task(3)(1)	task(63)	13			2
30	Preemption	task(63)	task(3)(2)				

A report that describes your implementation (please attach the screenshot of the code and MARK the modified part). (40%)

(如未標示,整段即為全新 code)

ucos_ii.h

global variable

建立全域變數,用於後面計算 response time。

Main function

```
int i, j, key;
       task_para_set tmp_TaskParameter;
       for (j = 1; j < TASK_NUMBER; j++)</pre>
          key = TaskParameter[j].TaskPeriodic;
          tmp_TaskParameter = TaskParameter[j];
          while (i >= 0 && TaskParameter[i].TaskPeriodic > key)
          TaskParameter[i + 1] = tmp_TaskParameter;
      /// create n tasks
for (n = 0;n < TASK_NUMBER;n++)</pre>
          OSTaskCreateExt(task1,
              &TaskParameter[n],
               &Task_STK[n][TASK_STACKSIZE - 1],
               TaskParameter[n].TaskID,
               &Task_STK[n][0],
               TASK_STACKSIZE,
               &TaskParameter[n],
               (OS_TASK_OPT_STK_CHK | OS_TASK_OPT_STK_CLR)); /// opt
      /// PA#1 par
```

依照題目要求,需要能夠根據輸入的任務數量,建立任務,所以利用 for loop 進行 n 個 task 的建立。

因為要進行 RM scheduling, 所以必須將 task 的 priority 根據 period 進行排列(成反比),因此針對輸入的 period 將任務重新排列。

Task function

根據給定的 arrive time, execution time, period 進行運作的 task funtion。execution 期間卡在 while 迴圈中不斷檢查執行時間是否結束,結束後則使用 OSTimeDly()進入 waiting 狀態,直到下個周期開始。

OSStart()

根據題目要求,輸出標題。

variable

建立變數,等會用於計算 response time、preemtion time、job number 與 CtxSw number。

OSTimeTick()

在 Time tick 中,設定 task 的啟動時間。

OS SchedNew()

```
if ((Output_err = fopen_s(&Output_fp, "./Output.txt", "a")) != 0)
      printf("can't open Output.txt");
/// Next task and current task are same.
if (OSPrioCur = OSPrioHighRdy)
      if (OSTCBCur->self_continue == 1)
            printf("%2d\t", OSTimeGet());
fprintf(Output_fp, "%2d\t", OSTimeGet());
            printf(" Completion\t");
fprintf(Output_fp, " Completion\t");
            printf(" task(%2d)(%2d)", OSTCEPrioTb1[OSPrioCur]->OSTCEId, job_number[OSPrioCur]);
fprintf(Output_fp, " task(%2d)(%2d)", OSTCEPrioTb1[OSPrioCur]->OSTCEId, job_number[OSPrioCur]);
            job_number[OSPrioCur]++;
if (OSPrioHighRdy == 63)
                   \begin{array}{lll} printf(" & task(\%2d) & & \\ tt", & OSPrioHighRdy); \\ fprintf(Output\_fp, " & task(\%2d) & \\ tt", & OSPrioHighRdy); \\ \end{array} 
                   \begin{array}{lll} printf("& task(\%2d)(\%2d)\t\t", OSTCEPrioTbl[OSPrioHighRdy]-SOSTCEId, job_number[OSPrioHighRdy]); \\ fprintf(Output_fp, "& task(\%2d)(\%2d)\t\t", OSTCEPrioTbl[OSPrioHighRdy]-SOSTCEId, job_number[OSPrioHighRdy]); \\ \end{array} 
            /// ResponseTime
printf("%4d\t\t", (OSTimeGet() - task_start_time[OSPrioCur]));
fprintf(Output_fp, "%4d\t\t", (OSTimeGet() - task_start_time[OSPrioCur]));
/// # of ContextSwitch
            /// task_preempt_count[OSPrioCur]++; /// same task
printf("%7d\t\t", task_preempt_count[OSPrioCur]);
fprintf(Output_fp, "%7d\t\t", task_preempt_count[OSPrioCur]);
            printf("%4d\t", task_preempt_time_acc[OSPrioCur]);
            fprintf(Output_fp, "%4d\t", task_preempt_time_acc[OSPrioCur]);
// DelayTime: Period - PersonseTime
            // Delaylime: Ferrou - Responselime
printf("Middu", (TaskParameter(OSPrioCur - 1].TaskPeriodic - (OSTimeGet() - task_start_time[OSPrioCur])));
fprintf(Output_fp, "%11d\n", (TaskParameter(OSPrioCur - 1].TaskPeriodic - (OSTimeGet() - task_start_time[OSPrioCur])));
            task_preempt_time[OSPrioCur] = 0;
task_preempt_count[OSPrioCur] = 0;
            task_preempt_time_acc[OSPrioCur] = 0;
            if (task preempt time[OSPrioHighRdv] != 0) /// if the next task has been preempted
                  task_preempt_time_acc[OSPrioHighRdy] = task_preempt_time_acc[OSPrioHighRdy] + (OSTimeGet() - task_preempt_time[OSPrioHighRdy]);
            task_start_time[OSTCBCur->OSTCBPrio] = OSTimeGet();
else if ((OSPrioCur != OSPrioHighRdy && OSTimeGet() != 0))
     if (OSPrioCur = 0)
            task_start_time[OSPrioHighRdy] = OSTimeGet();
            printf("%2d\t", OSTimeGet());
fprintf(Output_fp, "%2d\t", OSTimeGet());
            if ((OSRdyTb1[OSPrioCur >> 3] & OSMapTb1[OSPrioCur & 0x07]) = 0)
                 printf(" Completion\t");
fprintf(Output_fp, " Completion\t");
                        printf(" task(%2d)(%2d)", OSTCEPrioTb1[OSPrioCur]->OSTCEId, job_number[OSPrioCur]);
fprintf(Output_fp, " task(%2d)(%2d)", OSTCEPrioTb1[OSPrioCur]->OSTCEId, job_number[OSPrioCur]);
                 if (OSPrioHighRdv == 63)
```

```
task(%2d)(%2d)\t\t", OSTCEPrioTb1[OSPrioHighRdy]->OSTCEId, job_number[OSPrioHighRdy]);
tput_fp, "task(%2d)(%2d)\t\t", OSTCEPrioTb1[OSPrioHighRdy]->OSTCEId, job_number[OSPrioHighRdy]);
             printf("
  /// wesponserime
printf("Mad'\t", (OSTimeGet() - task_start_time[OSPrioCur]));
fprintf(Output_fp, "%4d\t\t", (OSTimeGet() - task_start_time[OSPrioCur]));
  for Contextoricum
for your count(OSPrioCur)++;
printf("%7d\t\t", task_preempt_count(OSPrioCur));
fprintf(Output_fp, "%7d\t\t", task_preempt_count(OSPrioCur));
  printf("%4d\t", task_preempt_time_acc[OSPrioCur]);
fprintf(Output_fp, "%4d\t", task_preempt_time_acc[OSPrioCur]);
  // Delayline: Feriod - Responseline
print("Wildum", (TaskParameter(OSPrioCur - 1].TaskPeriodic - (OSTimeGet() - task_start_time[OSPrioCur])));
fprintf(Output_fp, "Wild\n", (TaskParameter(OSPrioCur - 1].TaskPeriodic - (OSTimeGet() - task_start_time[OSPrioCur])));
  task_preempt_count[OSPrioCur] = 0;
task_preempt_time_acc[OSPrioCur] = 0;
  if (task_preempt_time[OSPrioHighRdy] != 0)
           task_preempt_time_acc[OSPrioHighRdy] = task_preempt_time_acc[OSPrioHighRdy] + (OSTimeGet() - task_preempt_time[OSPrioHighRdy]
task_preempt_count[OSPrioHighRdy]++;
job_number[OSPrioCur]++;
if (OSRdyTb1[y] = Ou) {
   OSRdyGrp &= (OS_PRIO)-OSTCBCur->OSTCBBitY;
                      OSTCBCur->OSTCBDly = ticks;
OS_TRACE_TASK_DLY(ticks);
                     OS_EXIT_CRITICAL();
           OSTCBCur->already_delay = 1;
           printf(" Completion\t");
fprintf(Output_fp, " Completion\t");
           /// CurrentTask
printf(" task(%2d)(%2d)", OSTCEPrioTb1[OSPrioCur]->OSTCBId, job_number[OSPrioCur]);
fprintf(Output_fp, " task(%2d)(%2d)", OSTCEPrioTb1[OSPrioCur]->OSTCBId, job_number[OSPrioCur]);
           if (OSPrioHighRdv == 63)
                     printf(" task(%2d)(%2d)\t\t", OSTCEPrioTb[OSPrioHighRdy]->OSTCEId, job_number[OSPrioHighRdy]);
fprintf(Output_fp, " task(%2d)(%2d)\t\t", OSTCEPrioTb[OSPrioHighRdy]->OSTCEId, job_number[OSPrioHighRdy]);
             printf(^{\$4d}\t^{*}, (OSTimeGet() - task\_start\_time[OSPrioCur])); \\ fprintf(Output\_fp, ^{\$4d}\t^{*}, (OSTimeGet() - task\_start\_time[OSPrioCur])); \\ 
           for Confeatorition
for interference in the form of the form o
```

```
printf("%11d\n", (TaskParameter[OSPrioCur - 1].TaskPeriodic - (OSTimeGet() - task_start_time[OSPrioCur])));
fprintf(Output_fp, "%11d\n", (TaskParameter[OSPrioCur - 1].TaskPeriodic - (OSTimeGet() - task_start_time[OSPrioCur])));
                        task_preempt_time[OSPrioCur] = 0;
task_preempt_count[OSPrioCur] = 0;
task_preempt_time_acc[OSPrioCur] = 0;
                              task_preempt_time_acc[OSPrioHighRdy] = task_preempt_time_acc[OSPrioHighRdy] + (OSTimeGet() - task_preempt_time[OSPrioHighRdy]
                       task_preempt_count[OSPrioHighRdy]++;
                        job_number[OSPrioCur]++
                        printf(" Preemption\t");
fprintf(Output_fp, " Preemption\t");
                             printf(" task(%2d)(%2d)", OSTCEPrioTbl[OSPrioCur]->OSTCEId, job_number[OSPrioCur]);
fprintf(Output_fp, " task(%2d)(%2d)", OSTCEPrioTbl[OSPrioCur]->OSTCEId, job_number[OSPrioCur]);
                        if (OSPrioHighRdy = 63)
                               \begin{array}{ll} printf(" & task(\%2d) & \n", OSPrioHighRdy); \\ fprintf(Output\_fp, " & task(\%2d) & \n", OSPrioHighRdy); \end{array} 
                              printf(" task(%2d)(%2d)\n", OSTCEPrioTb1[OSPrioHighRdy]->OSTCEId, job_number[OSPrioHighRdy]);
fprintf(Output_fp, " task(%2d)(%2d)\n", OSTCEPrioTb1[OSPrioHighRdy]->OSTCEId, job_number[OSPrioHighRdy]);
                       task_preempt_time[OSPrioCur] = OSTimeGet(); /// if the next task has been preempted
task_preempt_count[OSPrioCur]++;
                       /// task_preempt_time[OSPrioHighRdy]
task_preempt_count[OSPrioHighRdy]++;
fclose(Output_fp);
```

透過 OSRdyTbl 去確認 task 是否 ready 並進入 waiting,來區分任務為 completion與 preemption,其中加入 OSPrioCur與 OSPrioHighRdy,可以協助判斷 task 的切換。

資料	判斷、記錄方式			
Event	Task a 切換至 Task b 時,若 Task a 在 OSRdyTbl 上仍處於 ready 狀態(==1),則此情況為 preemtpion;反之則為 completion。			
Response Time	「Completion 的時間點」減掉「task 被設為 ready 的時間點(記錄於 task_start_time)」,即為 response time。			
# of ContextSwitch	記錄於 task_preempt_count,每當 task 切換則+1,completion輸出後歸零。			
Preemption Time	發生 preemtpion 時,記錄「task 被 preempt 的時間點」, 與「後續返回到該 task 的時間點」之差距,即為 preemption time。再 task complete 之前有多次 preemption 則進行累加。			

Report

這次的 PA1 主要為 RM Scheduling 的實作,首先透過熟悉 TCB 的運作過程,再到如何計算 Response Time 等等工作時間,過程挫折感滿重的,需要將紙本的知識應用於實作上,常常因為不夠清楚,導致一直踩坑。

雖然沒有做出完整的功能,但卻更加熟悉 RM 的設計流程,非常有趣,也透過這次作業知道了生活上有許事情要安排,就好像這次作業一樣,要如何評估是非常重要的,才不會因為期中考與作業同時來的時候,導致 Miss Deadline。