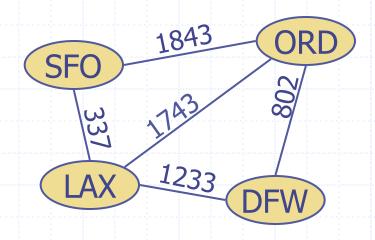
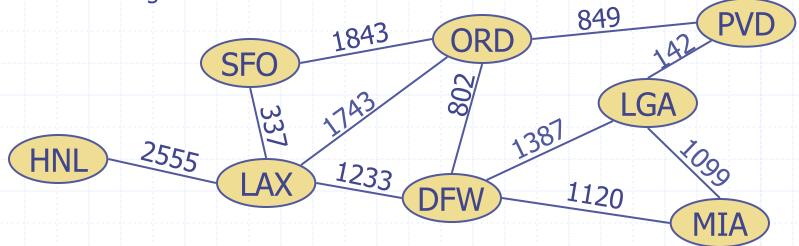
Presentation for use with the textbook Data Structures and Algorithms in Java, 6th edition, by M. T. Goodrich, R. Tamassia, and M. H. Goldwasser, Wiley, 2014

Graphs



Graphs

- \Box A graph is a pair (V, E), where
 - V is a set of nodes, called vertices
 - *E* is a collection of pairs of vertices, called edges
 - Vertices and edges are positions and store elements
- Example:
 - A vertex represents an airport and stores the three-letter airport code
 - An edge represents a flight route between two airports and stores the mileage of the route



Edge Types

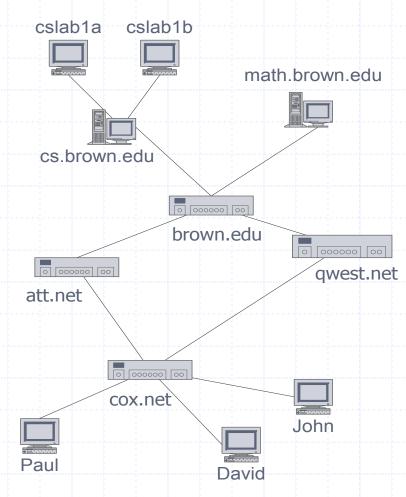
- Directed edge
 - ordered pair of vertices (u,v)
 - first vertex u is the origin
 - second vertex v is the destination
 - e.g., a flight
- Undirected edge
 - unordered pair of vertices (u,v)
 - e.g., a flight route
- Directed graph
 - all the edges are directed
 - e.g., route network
- Undirected graph
 - all the edges are undirected
 - e.g., flight network





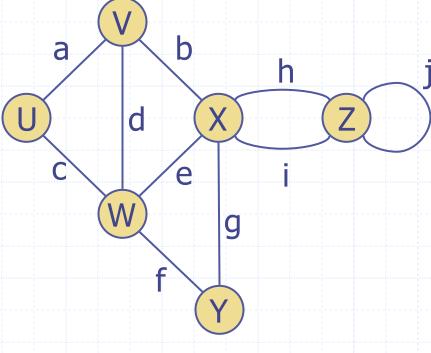
Applications

- Electronic circuits
 - Printed circuit board
 - Integrated circuit
- Transportation networks
 - Highway network
 - Flight network
- Computer networks
 - Local area network
 - Internet
 - Web
- Databases
 - Entity-relationship diagram



Terminology

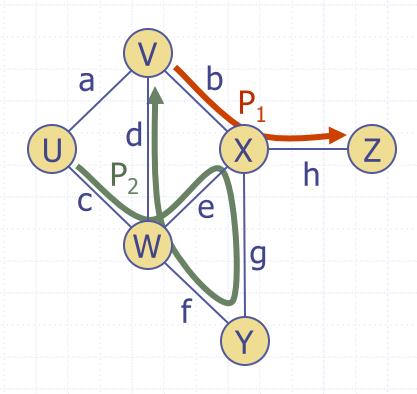
- End vertices (or endpoints) of an edge
 - U and V are the endpoints of a
- Edges incident on a vertex
 - a, d, and b are incident on V
- Adjacent vertices
 - U and V are adjacent
- Degree of a vertex
 - X has degree 5
- Parallel edges
 - h and i are parallel edges
- Self-loop
 - j is a self-loop



Terminology (cont.)

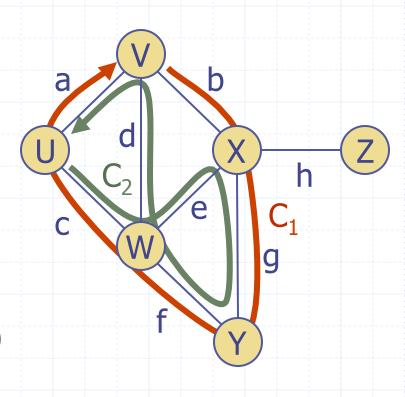
Path

- sequence of alternating vertices and edges
- begins with a vertex
- ends with a vertex
- each edge is preceded and followed by its endpoints
- Simple path
 - path such that all its vertices and edges are distinct
- Examples
 - $P_1 = (V,b,X,h,Z)$ is a simple path
 - P₂=(U,c,W,e,X,g,Y,f,W,d,V) is a path that is not simple



Terminology (cont.)

- Cycle
 - circular sequence of alternating vertices and edges
 - each edge is preceded and followed by its endpoints
- Simple cycle
 - cycle such that all its vertices and edges are distinct
- Examples
 - C₁=(V,b,X,g,Y,f,W,c,U,a,≼) is a simple cycle
 - C₂=(U,c,W,e,X,g,Y,f,W,d,V,a,≼) is a cycle that is not simple



Properties

Property 1

 $\sum_{v} \deg(v) = 2m$

Proof: each edge is counted twice

Property 2

In an undirected graph with no self-loops and no multiple edges

$$m \le n (n-1)/2$$

Proof: each vertex has degree at most (n-1)

What is the bound for a directed graph?

Notation

m

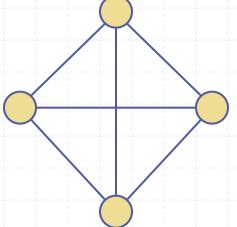
number of vertices number of edges deg(v) degree of vertex v

Example

$$= n = 4$$

$$\mathbf{m} = 6$$

$$\bullet \deg(v) = 3$$



Vertices and Edges

- A graph is a collection of vertices and edges.
- We model the abstraction as a combination of three data types: Vertex, Edge, and Graph.
- A Vertex is a lightweight object that stores an arbitrary element provided by the user (e.g., an airport code)
 - We assume it supports a method, element(), to retrieve the stored element.
- An **Edge** stores an associated object (e.g., a flight number, travel distance, cost), retrieved with the element() method.

Graph ADT

numVertices(): Returns the number of vertices of the graph.

vertices(): Returns an iteration of all the vertices of the graph.

numEdges(): Returns the number of edges of the graph.

edges(): Returns an iteration of all the edges of the graph.

getEdge(u, v): Returns the edge from vertex u to vertex v, if one exists; otherwise return null. For an undirected graph, there is no difference between getEdge(u, v) and getEdge(v, u).

endVertices(e): Returns an array containing the two endpoint vertices of edge e. If the graph is directed, the first vertex is the origin and the second is the destination.

opposite(v, e): For edge e incident to vertex v, returns the other vertex of the edge; an error occurs if e is not incident to v.

outDegree(v): Returns the number of outgoing edges from vertex v.

in Degree(v): Returns the number of incoming edges to vertex v. For an undirected graph, this returns the same value as does out Degree(v).

outgoingEdges(v): Returns an iteration of all outgoing edges from vertex v.

incomingEdges(v): Returns an iteration of all incoming edges to vertex v. For an undirected graph, this returns the same collection as does outgoingEdges(v).

insertVertex(x): Creates and returns a new Vertex storing element x.

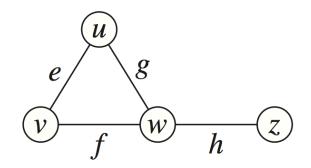
insertEdge(u, v, x): Creates and returns a new Edge from vertex u to vertex v, storing element x; an error occurs if there already exists an edge from u to v.

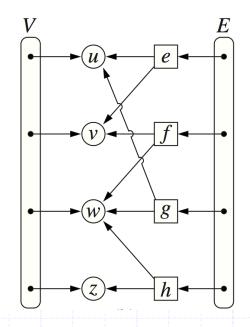
removeVertex(v): Removes vertex v and all its incident edges from the graph.

removeEdge(e): Removes edge e from the graph.

Edge List Structure

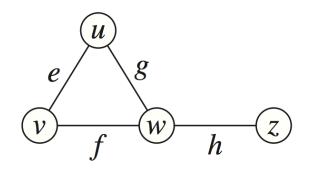
- Vertex object
 - element
 - reference to position in vertex sequence
- Edge object
 - element
 - origin vertex object
 - destination vertex object
 - reference to position in edge sequence
- Vertex sequence
 - sequence of vertex objects
- Edge sequence
 - sequence of edge objects

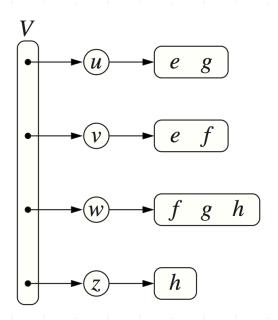




Adjacency List Structure

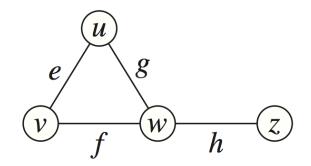
- Incidence sequence for each vertex
 - sequence of references to edge objects of incident edges
- Augmented edge objects
 - references to associated positions in incidence sequences of end vertices





Adjacency Matrix Structure

- Edge list structure
- Augmented vertex objects
 - Integer key (index) associated with vertex
- 2D-array adjacency array
 - Reference to edge object for adjacent vertices
 - Null for non nonadjacent vertices
- The "old fashioned"version just has 0 for no edge and 1 for edge



0	1	2	3
	e	g	
e		f	
g	f		h
		h	
	e	e e	$egin{array}{ c c c c c c c c c c c c c c c c c c c$

Performance

 n vertices, m edges no parallel edges no self-loops 	Edge List	Adjacency List	Adjacency Matrix
Space	n+m	n + m	n^2
incidentEdges(v)	m	deg(v)	n
areAdjacent (v, w)	m	$\min(\deg(v), \deg(w))$	1
insertVertex(o)	1	1	n^2
insertEdge(v, w, o)	1	1	1
removeVertex(v)	m	deg(v)	n^2
removeEdge(e)	1	1	1