

# Lógica Computacional

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FCT UNL

Aula Prática 7

Semântica da Lógica de Primeira Ordem.

## 1. Interpretação de termos e fórmulas.

Considere a assinatura  $\Sigma = (SF, SP)$  onde:

- $SF_0 = \{zero\}$ ,  $SF_1 = \{suc\}$ ,  $SF_2 = \{\oplus, \otimes\}$  e;
- $SP_1 = \{Par, Impar\}$ ,  $SP_2 = \{Eq, Meq\}$ .

Considere também a estrutura de interpretação  $\mathbf{Nat} = (\mathbb{N}_0, I)$ , sendo:

- $\underline{zero}_I = 0$ ;
- $\underline{suc}_I : \mathbb{N}_0 \rightarrow \mathbb{N}_0$  tal que  $\underline{suc}_I(n) = n + 1$ ;
- $\underline{\oplus}_I : \mathbb{N}_0^2 \rightarrow \mathbb{N}_0$  tal que  $\underline{\oplus}_I(n, m) = n + m$ ;
- $\underline{\otimes}_I : \mathbb{N}_0^2 \rightarrow \mathbb{N}_0$  tal que  $\underline{\otimes}_I(n, m) = n \times m$ ;
- $\underline{Par}_I : \mathbb{N}_0 \rightarrow \{0, 1\}$  tal que  $\underline{Par}_I(n) = 1$  sse  $n$  é par;
- $\underline{Impar}_I : \mathbb{N}_0 \rightarrow \{0, 1\}$  tal que  $\underline{Impar}_I(n) = 1$  sse  $n$  é ímpar;
- $\underline{Eq} : \mathbb{N}_0^2 \rightarrow \{0, 1\}$  tal que  $\underline{Eq}_I(n, m) = 1$  sse  $n = m$ ;
- $\underline{Meq}_I : \mathbb{N}_0^2 \rightarrow \{0, 1\}$  tal que  $\underline{Meq}_I(n, m) = 1$  sse  $n \leq m$ .

Assuma a atribuição  $\rho : X \rightarrow \mathbb{N}_0$  tal que  $\rho(n) = 3$  e  $\rho(m) = 2$ .

(a) Determine a interpretação dos seguintes termos em  $\mathbf{Nat}$ .

- i.  $\llbracket zero \rrbracket_{\mathbf{Nat}}^\rho$
- ii.  $\llbracket n \rrbracket_{\mathbf{Nat}}^\rho$
- iii.  $\llbracket suc(n) \rrbracket_{\mathbf{Nat}}^\rho$
- iv.  $\llbracket \oplus(suc(zero), m) \rrbracket_{\mathbf{Nat}}^\rho$
- v.  $\llbracket \otimes(\oplus(m, suc(n)), \oplus(suc(zero), m)) \rrbracket_{\mathbf{Nat}}^\rho$

(b) Determine se são verdadeiras as afirmações seguintes.

- i.  $\mathbf{Nat}, \rho \models Meq(zero, n)$ ;
- ii.  $\mathbf{Nat}, \rho \models Meq(m, \oplus(suc(zero), m))$ ;
- iii.  $\mathbf{Nat}, \rho \models Eq(n, m) \wedge Meq(n, suc(n))$ ;
- iv.  $\mathbf{Nat}, \rho \models Par(n) \rightarrow Impar(n)$ ;
- v.  $\mathbf{Nat}, \rho \models \exists n Impar(suc(n))$ ;
- vi.  $\mathbf{Nat}, \rho \models \exists n Eq(suc(n), zero)$ ;
- vii.  $\mathbf{Nat}, \rho \models \forall n Eq(suc(n), m)$ ;
- viii.  $\mathbf{Nat}, \rho \models \forall n \neg Eq(suc(n), zero)$ .

2. Consequência semântica

Verifique se são verdadeiras as seguintes afirmações.

- (a)  $\{P(y)\} \models \forall x P(x)$
- (b)  $\{\exists x P(x)\} \models \forall x P(x)$
- (c)  $\{\exists x P\} \models \forall x P$
- (d)  $\{\forall x P(x) \rightarrow \forall x Q(x)\} \models \forall x (P(x) \rightarrow Q(x))$
- (e)  $\{\forall x (P(x) \vee Q(x))\} \models \forall x P(x) \vee \forall x Q(x)$
- (f)  $\{\exists x P(x) \wedge \exists x Q(x)\} \models \exists x (P(x) \wedge Q(x))$
- (g)  $\{\exists x (P(x) \wedge Q(x))\} \models \exists x P(x) \wedge \exists x Q(x)$
- (h)  $\{\forall x (\neg(P(x) \wedge Q(x)))\} \models \forall x \neg P(x) \wedge \forall x \neg Q(x)$
- (i)  $\{\forall x (P(x) \rightarrow \neg Q(x)), \neg Q(a)\} \models P(a)$
- (j)  $\{\forall x (P(x) \rightarrow \neg Q(x)), P(a)\} \models \neg Q(a)$
- (k)  $\{\forall x (P(x) \rightarrow Q(x)), \forall x (Q(x) \rightarrow R(x))\} \models \forall x (P(x) \rightarrow R(x))$
- (l)  $\{\exists x (\varphi \wedge \psi)\} \models \exists x \varphi \wedge \exists x \psi$
- (m)  $\{\forall x \varphi \wedge \forall x \psi\} \models \forall x (\varphi \wedge \psi)$
- (n)  $\{\exists x \varphi \vee \exists x \psi\} \models \exists x (\varphi \vee \psi)$
- (o)  $\{\exists x \neg \varphi\} \models \neg \forall x \varphi$
- (p)  $\{\forall x \neg \varphi\} \models \neg \exists x \varphi$

Se  $x \notin \text{VL}(\psi)$ :

- (q)  $\{\forall x (\varphi \wedge \psi)\} \models \forall x \varphi \wedge \psi$
- (r)  $\{\forall x (\varphi \vee \psi)\} \models \forall x \varphi \vee \psi$
- (s)  $\{\exists x (\varphi \wedge \psi)\} \models \exists x \varphi \wedge \psi$
- (t)  $\{\exists x (\varphi \vee \psi)\} \models \exists x \varphi \vee \psi$