

# HW3 Question 1 Proof

CMSC 423 Fall 2014

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- (A) First we show that the MaxDist quantity is non-increasing as more centers are added in farthest-first traversal. Let  $X^{k-1}$  be the set of centers selected in the first  $k - 1$  iterations of farthest-first traversal, and  $x_k$  be the center selected in the  $k$ -th step, then  $\text{MaxDist}(\text{Data}, X) \leq \text{MaxDist}(\text{Data}, X^{k-1})$ .

$$\text{MaxDist}(\text{Data}, X) = \max_{u \in \text{Data}} \min_{x \in X} d(u, x) \quad (1)$$

$$= \max_{u \in \text{Data}} \min \left[ \min_{x \in X^{k-1}} d(u, x), d(u, x_k) \right] \quad (2)$$

$$\leq \max_{u \in \text{Data}} \min_{x \in X^{k-1}} d(u, x) \quad (3)$$

$$= \text{MaxDist}(\text{Data}, X^{k-1}) \quad (4)$$

From this it follows that  $\text{MaxDist}(\text{Data}, X) \leq \text{MaxDist}(\text{Data}, X^t)$  for all  $t=[1, \dots, k-1]$ .

- (B) Using (A), we show that  $d(x_i, x_j) \geq \text{MaxDist}(\text{Data}, X)$  for all  $x_i, x_j \in X$ .

Assume  $i < j$ , that is,  $x_i$  was chosen as a center in an earlier iteration than  $x_j$ . Then,

$$d(x_i, x_j) \geq \min_{x_t: t=[1, \dots, j-1]} d(x_t, x_j) = \text{MaxDist}(\text{Data}, X^{j-1}) \geq \text{MaxDist}(\text{Data}, X)$$

where the last inequality follows from (A).

- (C) Let  $u \in \text{Data}$  be such that  $d(u, X) = \text{MaxDist}(\text{Data}, X)$ , then by definition  $d(u, x) \geq \text{MaxDist}(\text{Data}, X)$  for all  $x \in X$ .
- (D) From (B) and (C), it follows that  $X \cup \{u\}$  is a set of  $k + 1$  points in Data with distance between every pair of points greater than or equal to  $\text{MaxDist}(\text{Data}, X)$ .
- (E) The optimal  $k$  clustering must include two of the points defined in (D) in one of its clusters. Therefore,  $\text{MaxDist}(\text{Data}, X_{\text{opt}}) \geq \frac{\text{MaxDist}(\text{Data}, X)}{2}$ . Which proves the result.