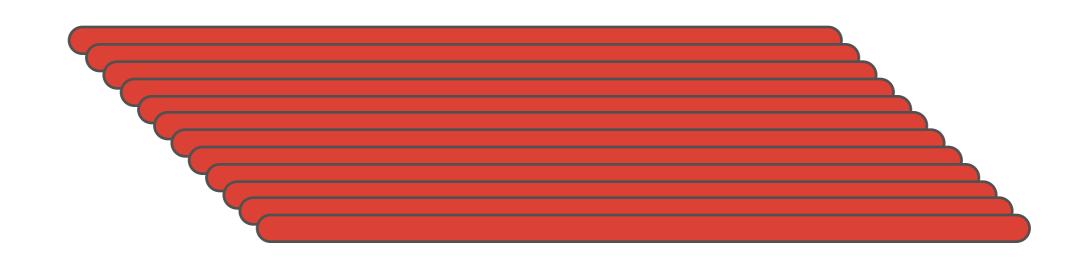
Introduction to genome assembly

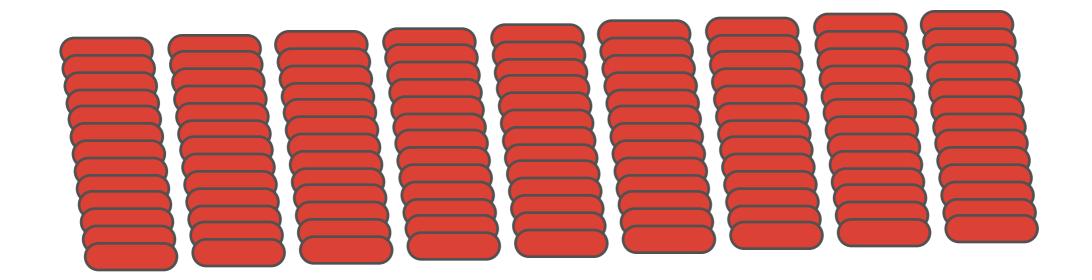
CMSC423
Many slides courtesy of Ben Langmead

SEC-GEN SEQUENCING





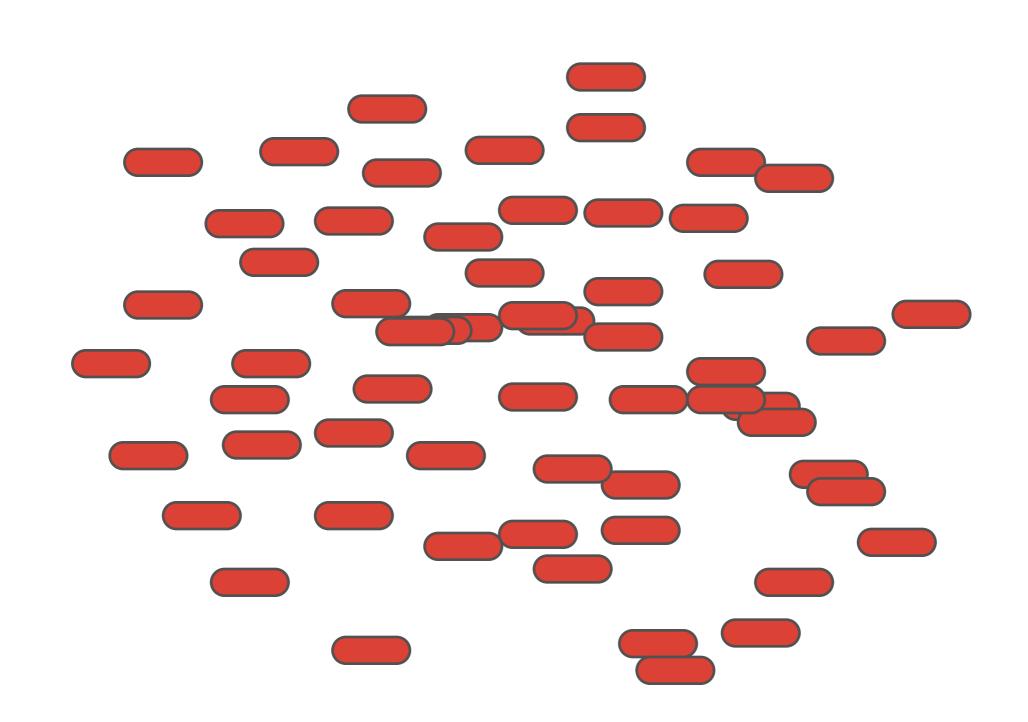
SEC-GEN SEQUENCING



Fragmentation is random, i.e., not equal-sized (but hard to draw)



SEC-GEN SEQUENCING



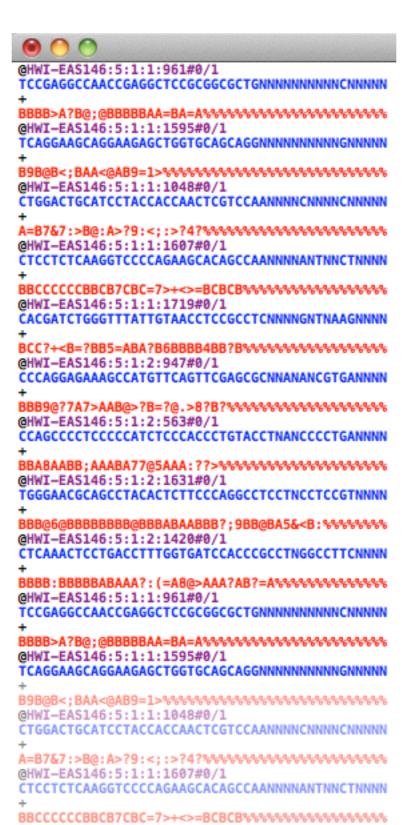


SECOND-GENERATION SEQUENCING

- "Ultra high throughput" DNA sequencing
 - 6 gigabases / day vs.
 - 3 gigabases / 13 years (human genome project, more or less)
 - 200 bp long reads



From reads to evidence



From reads to evidence



2. Comparative

Sequence-wise, individuals of a species are nearly identical

Well curated, annotated "reference" genomes exist





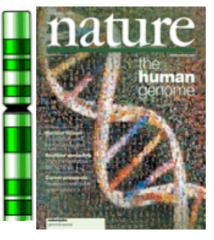




D. melanogaster, Science, 2000

H. sapiens, Nature, 2000 and Science, 2000

M. musculus, Nature, 2002



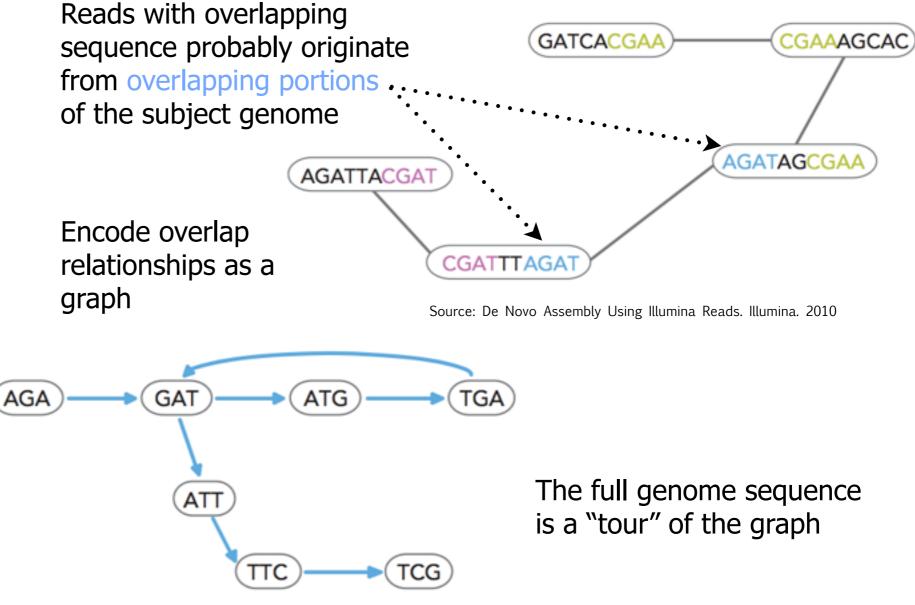
Idea: "Map" reads to their point of origin with respect to a reference, then study differences

From reads to evidence



I. de novo

Assume nothing! - let reads tell us everything



Source: De Novo Assembly Using Illumina Reads. Illumina. 2010 http://www.illumina.com/Documents/products/technotes/technote_denovo_assembly.pdf

What we'll cover

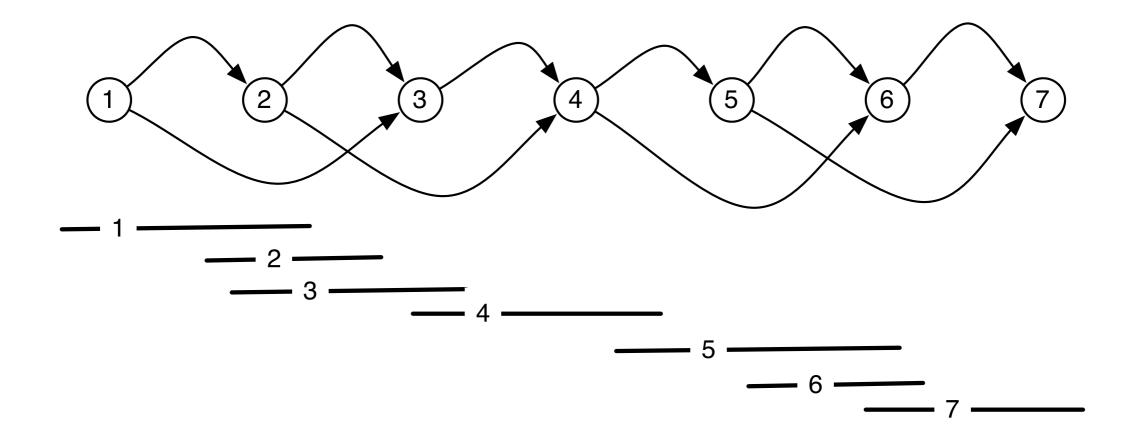
- Genome assembly as graph problems
 - Two representations:
 - Overlap graph
 - How much sequencing required for assembly
 - DeBruijn graph
- How to get assemblies from solutions to graph problems

Overlap Graph

Overlap graph:

Nodes = reads

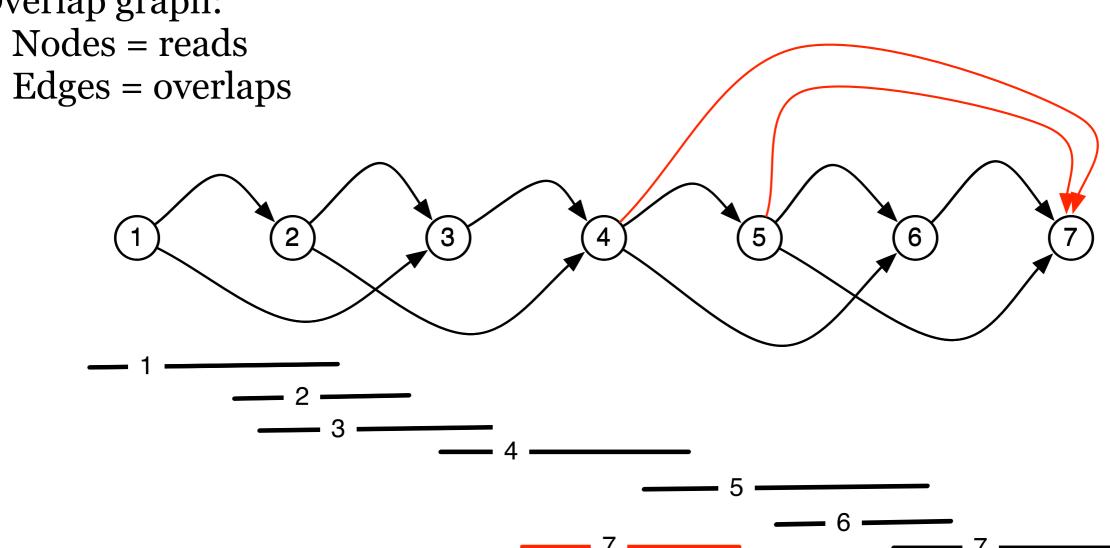
Edges = overlaps



Given overlap graph, how can we find a good candidate assembly?

Overlap Graph

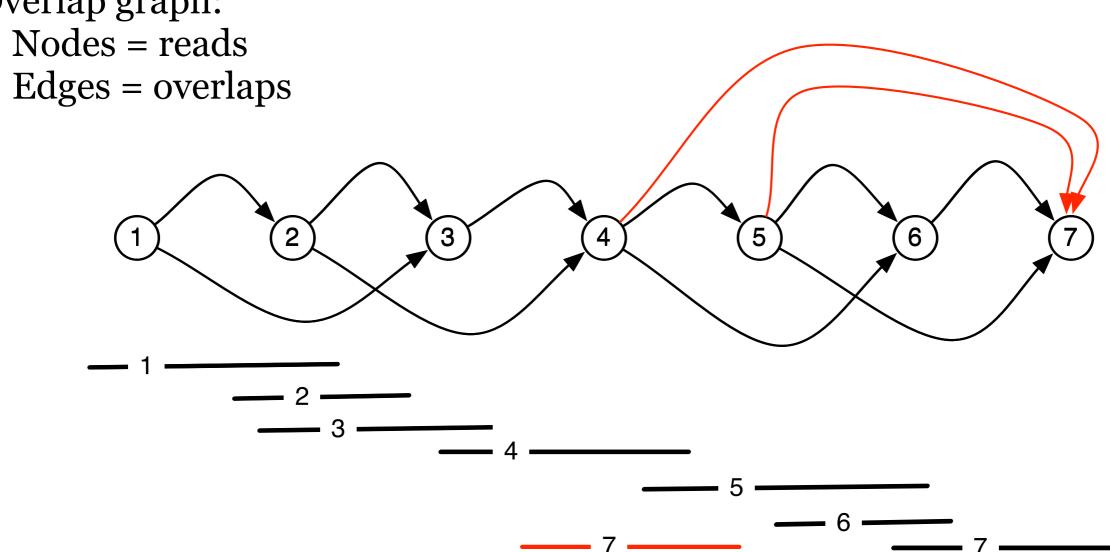
Overlap graph:



Given overlap graph, how can we find a good candidate assembly?

Overlap Graph

Overlap graph:



Given overlap graph, how can we find a good candidate assembly?

Hamiltonian Path (aka Traveling Salesman Path): visit every node in the graph exactly once.

Hamiltonian Path

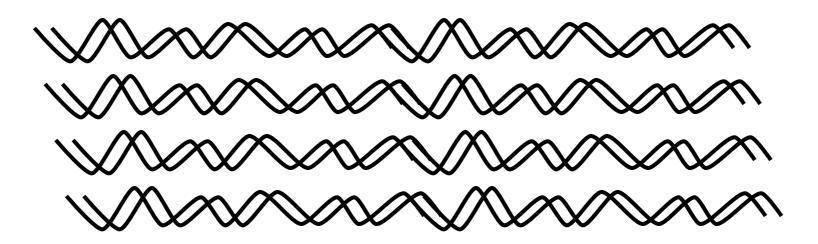
- Motivation: Every read must be used in exactly one place in the genome.
- Hamiltonian Path is NP-hard.
- Though good solvers exist, they can't operate on the millions of reads from a sequencing project.
- Solution: greedy walk along the graph.



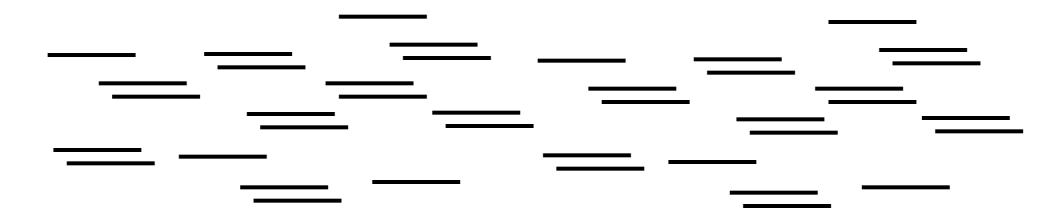
Optimal Hamiltonian path of 24,978 cities in Sweden (Applegate et al, 2004, www.tsp.gatech.edu/sweden/index.html).

Shotgun Sequencing

Many copies of the DNA



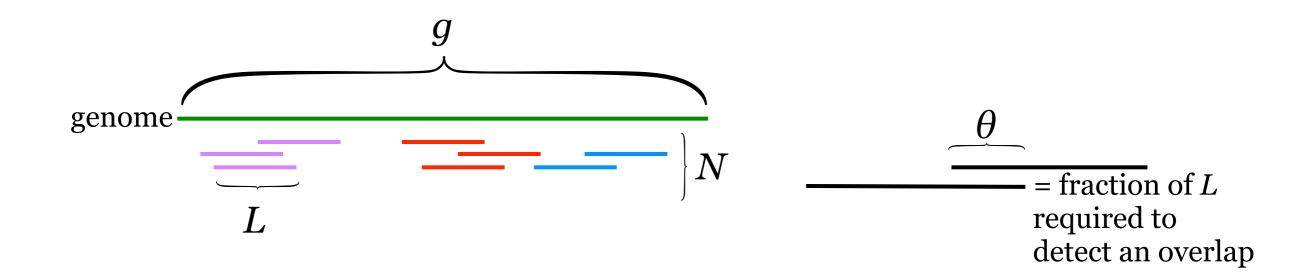
Shear it, randomly breaking them into many small pieces, read ends of each:



Assemble into original genome:

Lander-Waterman Statistics

How many reads to we need to be sure we cover the whole genome?



An *island* is a contiguous group of reads that are connected by overlaps of length $\geq \theta L$. (Various colors above)

Want: Expression for expected # of islands given N, g, L, θ .

Expected # of Islands

 $\lambda := N/g = \text{probability a read starts at a given position}$ (assuming random sampling)

Pr(*k* reads start in an interval of length *x*)

x trials, want k "successes," small probability λ of success Expected # of successes = λx

Poisson approximation to binomial distribution:

$$\Pr(k \text{ reads in length } x) = e^{-\lambda x} \frac{(\lambda x)^k}{k!}$$

Expected # of islands = $N \times Pr(\text{read is at rightmost end of island})$

$$= N \times \text{Pr(o reads start in } (1-\theta)L)$$

$$= Ne^{-\lambda(1-\theta)L} \frac{(\lambda(1-\theta)L)^0}{0!}$$

$$= Ne^{-\lambda(1-\theta)L}$$

$$= Ne^{-\lambda(1-\theta)L}$$

$$= Ne^{-(1-\theta)LN/g} \leftarrow LN/g \text{ is called the coverage } c.$$

Expected # of Islands, 2

Rewrite to depend more directly on the things we can control: c and θ

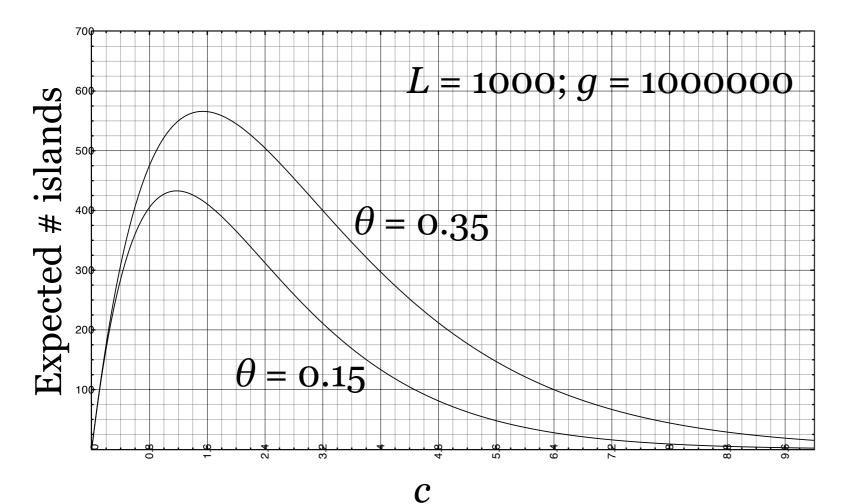
Expected # of islands =
$$Ne^{-(1-\theta)LN/g}$$

$$= Ne^{-(1-\theta)c}$$

$$= \frac{L/g}{L/g} N e^{-(1-\theta)c}$$

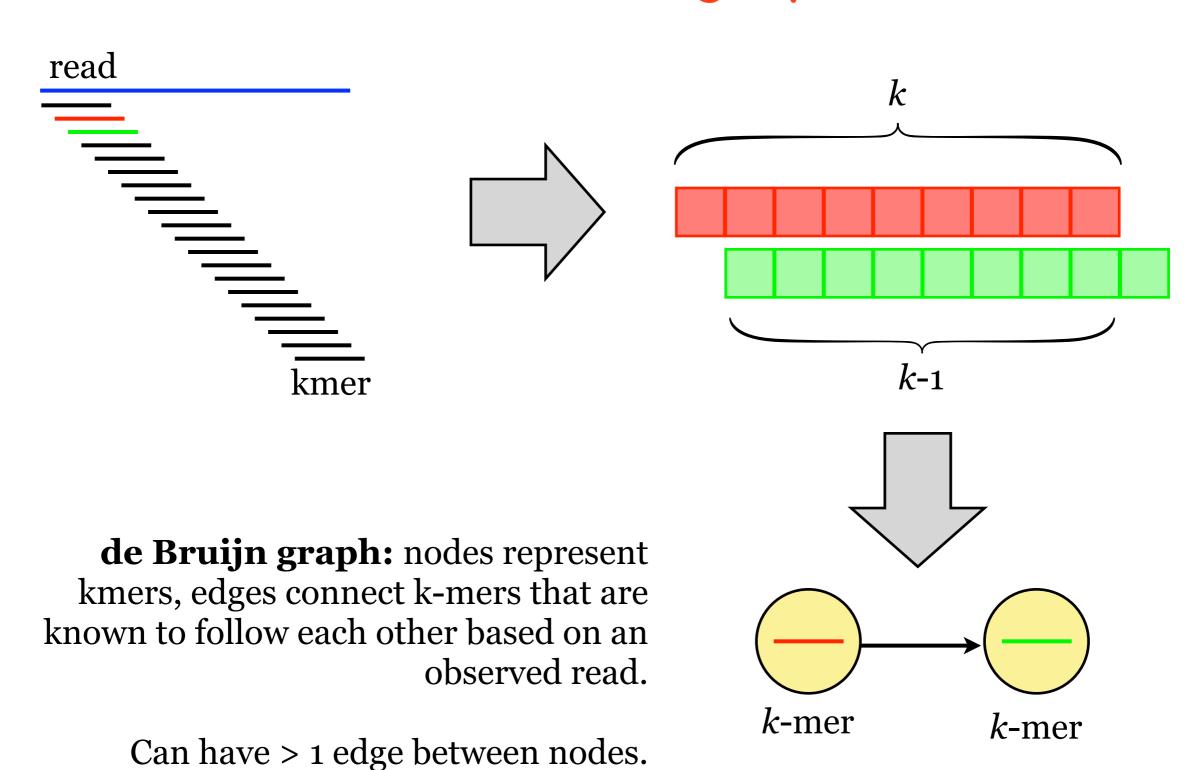
$$= \frac{g}{L} c e^{-(1-\theta)c}$$

$$= \frac{g}{L}ce^{-(1-\theta)c}$$

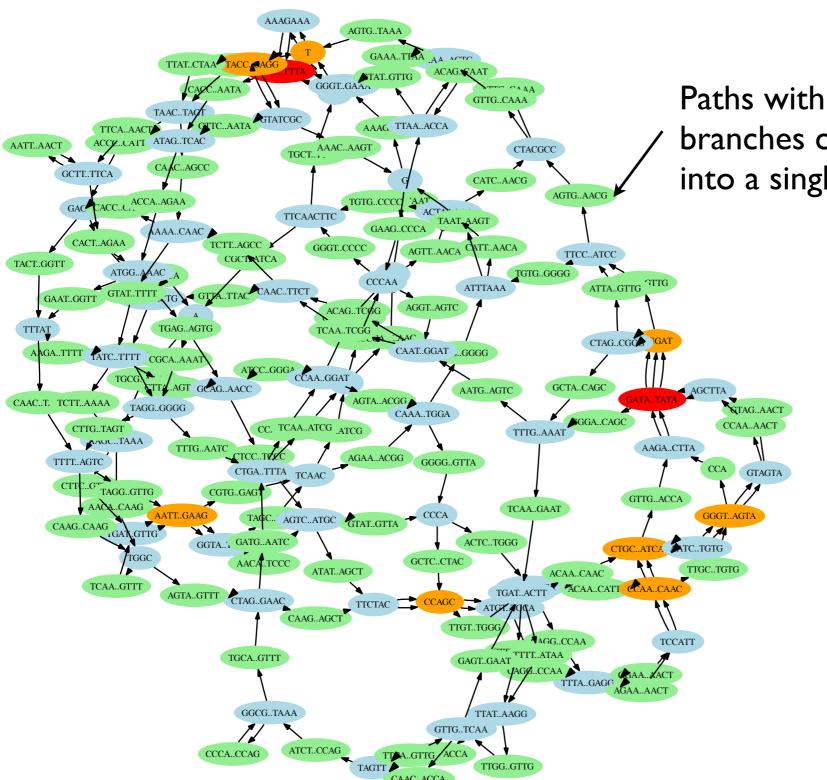


Assembly via Eulerian Path

de Bruijn graph



Example bacterial de Bruijn graph

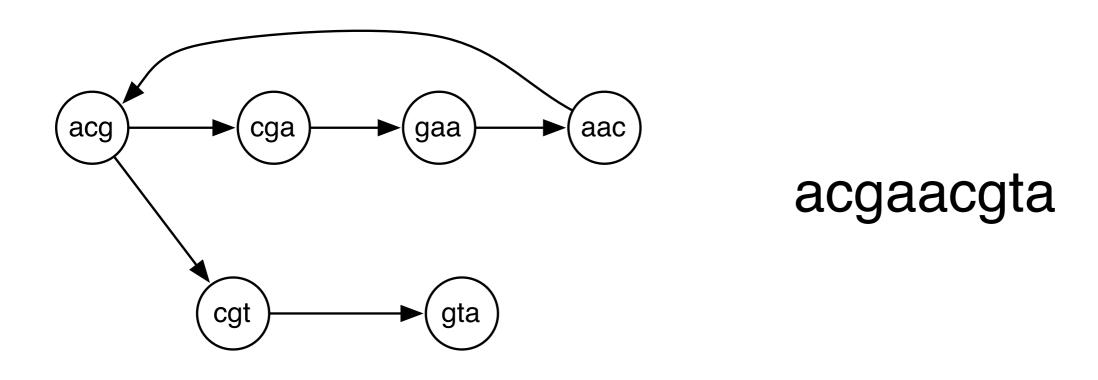


Paths with no branches compressed into a single node

Eulerian path = use every edge exactly once.

With perfect data, the genome can be reconstructed by some Eulerian path through this graph

Assembly via Eulerian Path



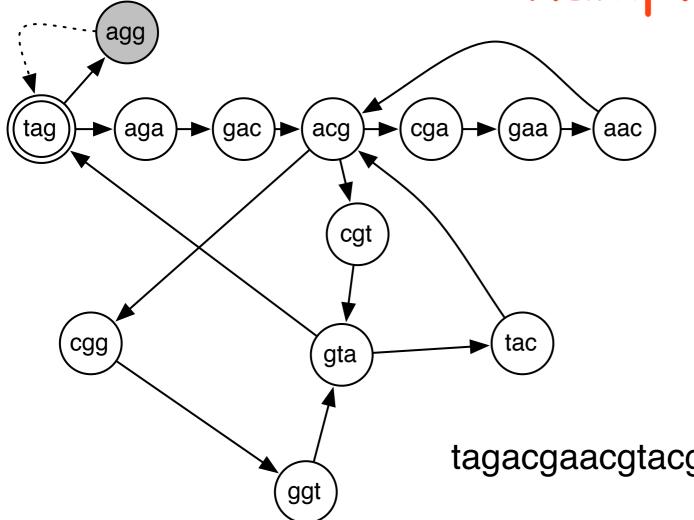
Let dG(s) be the de Bruijn graph of string s. Then s corresponds to some Eulerian path in dG(s).

A directed graph has an Eulerian path if and only if:

- One node has one more edge leaving it than entering
- One node has one more edge entering than leaving
- •All other nodes have the same number of edges entering and leaving

How can we find such a path?

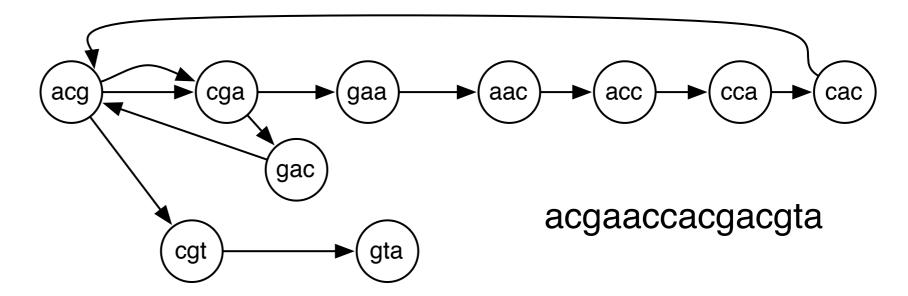
Examples



A directed graph has an Eulerian cycle if and only if:

•All nodes have the same number of edges entering and leaving

tagacgaacgtacggtagg



Eulerian Path Algorithm

Connect node with out-degree < in-degree to node with out-degree < in-degree. So that we will have an Eulerian cycle.

Why will you return to *u*?

*How can find such

a node quickly?

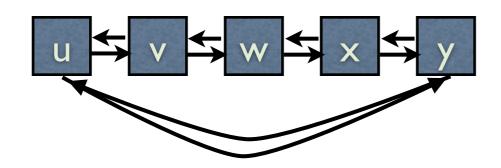
Walk from some arbitrary node u until you return to u, creating a doubly liked list of the path you visit.

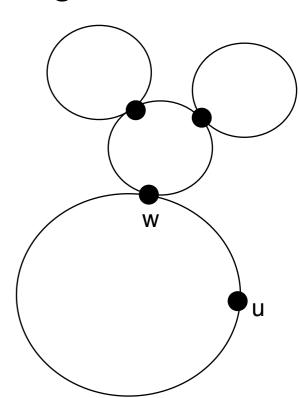
Repeat until all edges used:

•Start from some node w on the current tour with unused edges*.

•Walk along unused edges until you return to w, inserting the visited nodes

after w into the current tour list.





Eulerian Path Algorithm

Connect node with out-degree < in-degree to node with out-degree < in-degree. So that we will have an Eulerian cycle.

Why will you return to *u*?

*How can find such

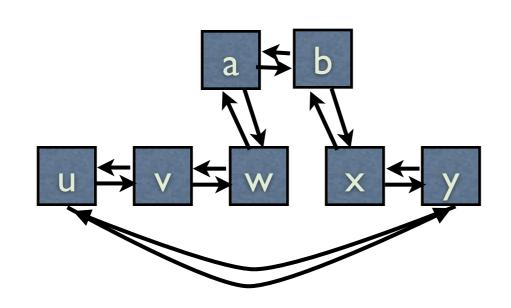
Walk from some arbitrary node u until you return to u, creating a doubly liked list of the path you visit.

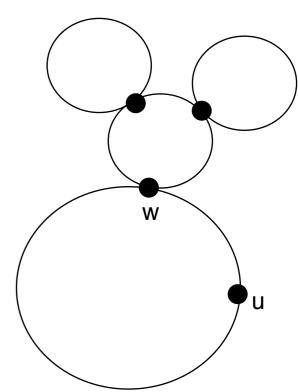
Repeat until all edges used:

a node quickly? •Start from some node w on the current tour with unused edges*.

•Walk along unused edges until you return to w, inserting the visited nodes

after w into the current tour list.



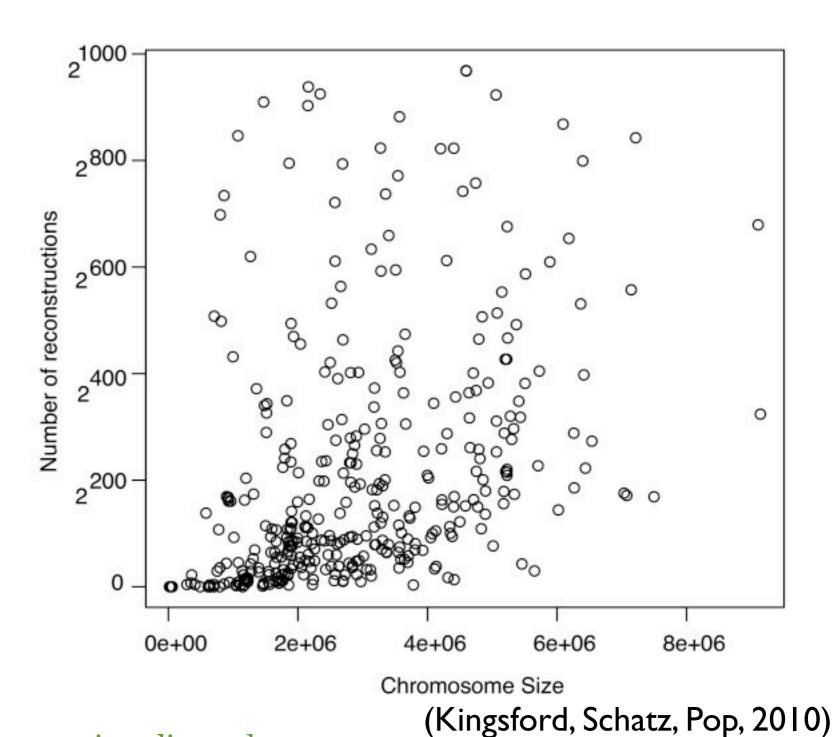


The Problem with Eulerian Paths

There are typically an astronomical number of possible Eulerian tours with perfect data.

Adding back constraints to limit # of tours leads to a NP-hard problem.

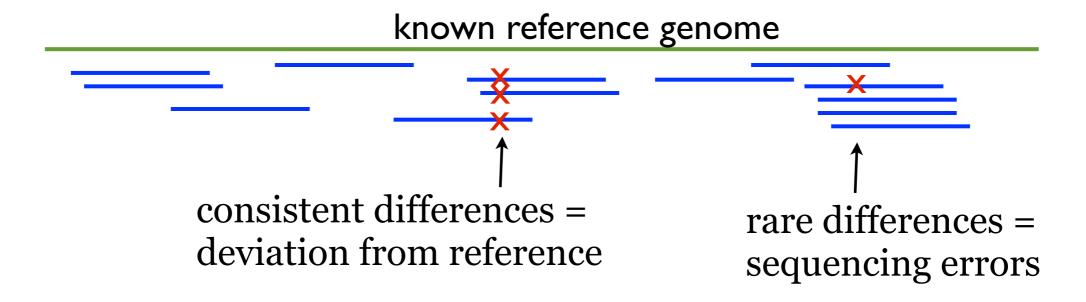
With imperfect data, there are usually NO Eulerian tours.



Aside: counting # of Eulerian tours in a directed graph is easy, but in an undirected graph is #P-complete (hard).

Comparative Assembly

Align reads to known genome:



Can use much lower coverage (e.g. 4X coverage instead of 20-30X for *de novo* assembly).

Aligning a large # of short sequences to one large sequence is an important special case of sequence alignment.

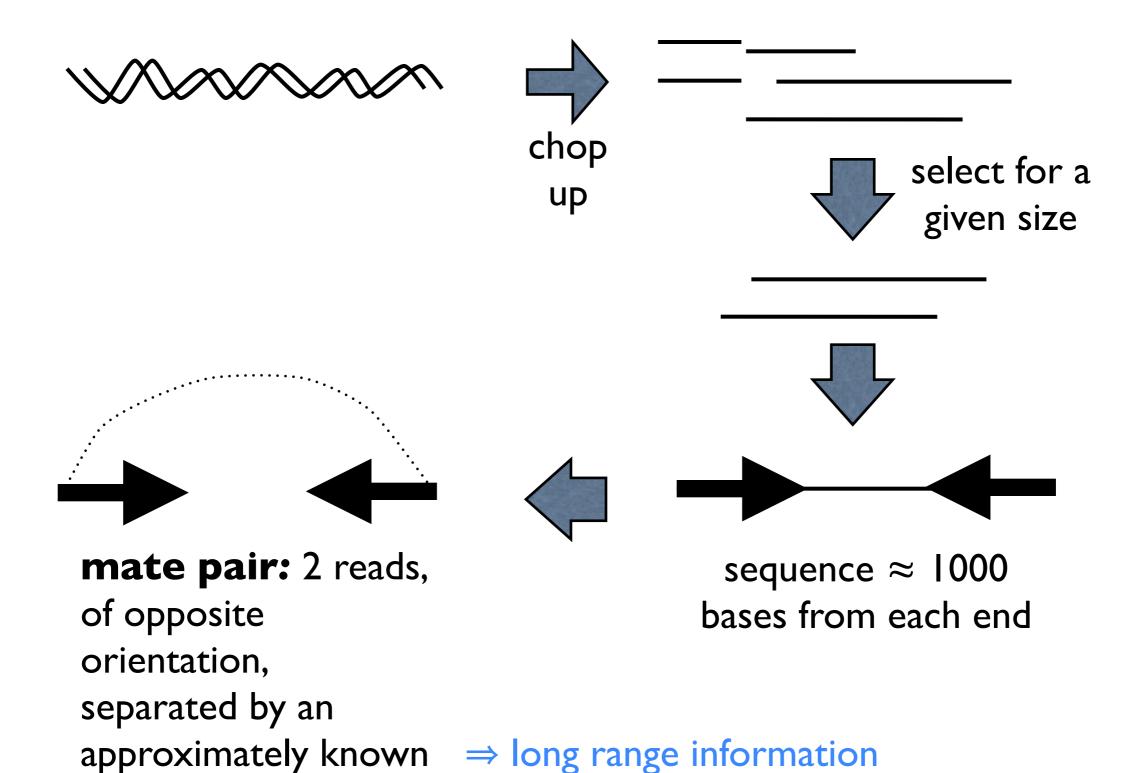
"1000" Genomes Project

find variants
that occur in >
1% of the
population:
sequence
≈2500 genomes
at 4X coverage,
align them to
reference.

1000 Genomes Samples						6-Sept-11		
Population	Status	Available to research community (dates approx)	DNA sequenced from blood	Offspring samples from trios	First set	Second set	Third set	Total
Utah residents (CEPH) with Northern and Western European ancestry (CEU)	Available	Available	no	yes	100			100
Toscani in Italia (TSI)	Available	Available	no	no	100			100
British from England and Scotland (GBR)	Available	Available	no	no	96	4		100
Finnish from Finland (FIN)	Available	Available	no	no	100			100
Iberian populations in Spain (IBS)	Available to project	Available	no	yes	30	70		100
Total European ancestry					426	74		500
Han Chinese in Beijing, China (CHB)	Available	Available	no	no	100			100
Japanese in Toyko, Japan (JPT)	Available	Available	no	no	100			100
Han Chinese South (CHS)	Available	Available	most	yes	100			100
Chinese Dai in Xishuangbanna (CDX)	Available to project	Oct-Dec 2011	some	no		100		100
Kinh in Ho Chi Minh City, Vietnam (KHV)	Available to project	Oct-Dec 2011	yes	some		100		100
Chinese in Denver, Colorado (CHD) (pilot 3 only)	Available	Available	no	no				0
TOTAL East Asian ancestry					300	200		500
Yoruba in Ibadan, Nigeria (YRI)	Available	Available	no	yes	100			100
Luhya in Webuye, Kenya (LWK)	Available	Available	no	no	100			100
Gambian in Western Division, The Gambia	Collecting samples	Mar-May 2012	no	yes			100	100

http://www.1000genomes.org/about#ProjectSamples

Mate Pairs



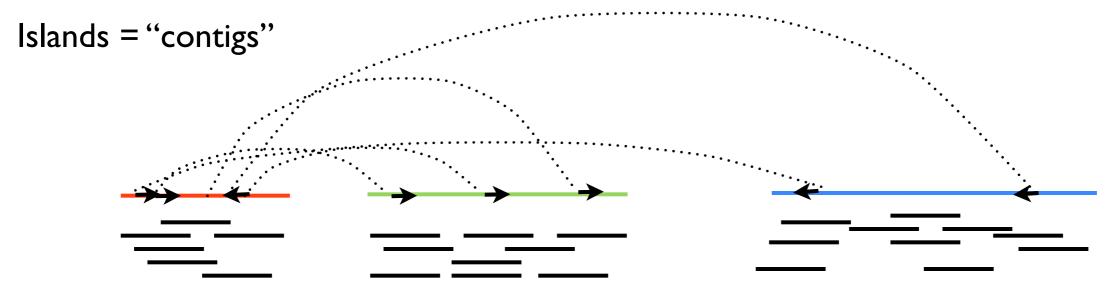
distance

Scaffolding

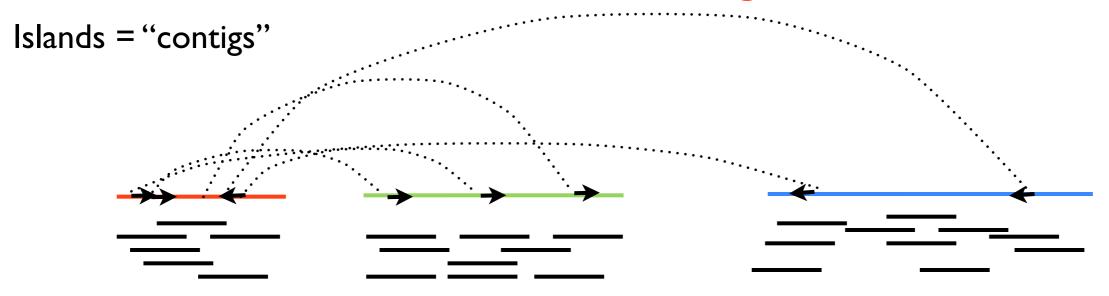
Islands = "contigs"

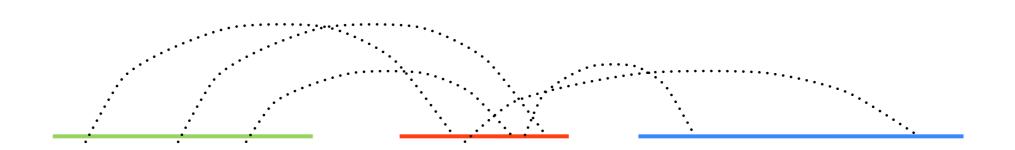


Scaffolding



Scaffolding





Summary

- Sanger sequencing reads DNA via synthesis; 800-1000bp.
- Assembly Paradigms:
 - Shortest Common Superstring (NP-hard; sensitive to repeats)
 - Hamiltonian cycle in overlap graph (NP-hard)
 - Eulerian cycle in de Bruijn graph (polynomial in basic form, but large # of solutions)
- Overlap alignment can be computed with slight variant of sequence alignment DP.
 - K-mer hashing technique avoids all pairs overlap alignment

Hard vs. Easy

- Eulerian path visit every edge once (easy)
- Hamiltonian path visit every node once (hard)
- Shortest common supersequence (easy)
- Shortest common superstring (hard)
- Counting Eulerian tours in directed graphs (easy)
- Counting Eulerian tours in undirected graphs (hard)
- Aligning 2 sequences (easy)
- Aligning k > 2 sequences (hard)
- Shortest path (easy)
- Longest path (hard)

References

- http://www.cbcb.umd.edu/research/assembly_primer
- http://www.ncbi.nlm.nih.gov/pmc/articles/PMC2874646/
- http://www.math.ucsd.edu/~gptesler/186/slides/shotgun_f13handout.pdf
- http://www.biomedcentral.com/1471-2105/11/21/abstract