

Introduction

Synchronous Message Eychange

Implementing RISC-V in SME

Fetching instruction and increment Instruction decode and execution

Branching Simple datapat

Instructions

Further work

Implementing RISC-V in Synchronous Message Exchange

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eScience lunch talk October 2, 2019



Overview

Introduction

Motivation

Synchronous

Message Exchange

Implementin RISC-V in SME

Fetching instruct and increment Instruction decoder and execution Memory access Branching Simple datapath Instructions

Further wor

2 Implementing RISC-V in SME
Single Cycle RISC-V
Fetching instruction and increment
Instruction decode and execution
Memory access
Branching
Simple datapath
Instructions
Control

3 Further work



Motivation

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Synchronous Message Exchan

Implementing RISC-V in SME

Fetching instruction and increment Instruction decode and execution

Branching Simple datapatl

Instructions Control

Further work

- Most measurement instruments are based on Intel x86 CPU
- Limits the bandwidth at which data collection is possible
- Limits the possibility for custom solutions



Motivation

Introduction

Synchronous Message Exchange

Implementing RISC-V in SME

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Memory access
Branching
Simple datapath

- RISC-V especially well suited embedding in scientific instruments
- Fully customizable = faster and more power efficient
- Could be implemented in a FPGA
- But FPGA programming is complicated





Synchronous Message Exchange

Synchronous Message Exchange

- Write FPGA applications within .Net framework¹
- Enjoy productivity features of modern language
- Automatic VHDI conversion
- Based on Communicating Sequential Processes (CSP)²

 $^{^{}m 1}$ Building Hardware from C# models, Kenneth Skovhede and Brian Vinter

 $^{^{2}\}mathsf{Communicating}$ sequential processes, Charles Antony Richard Hoare



Synchronous Message Exchange

Introduction

Motivation

Synchronous Message Exchange

Implementin RISC-V in SME

Single Cycle RISC-N Fetching instruction and increment Instruction decode and execution

Branching
Simple datapath

Further wor

- Processes are isolated no memory sharing
- Communicates with each other through channels
- Globally synchronous
- Perfect for implementing RISC-V!



Figure: One to one

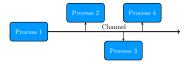


Figure: One to many



Single Cycle RISC-V

Introduction

Synchronous Message Exchang

Implementing RISC-V in SMF

Single Cycle RISC-V

Fetching instruction and increment

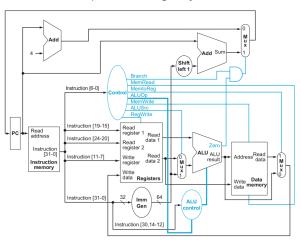
Instruction decode and execution

Memory access Branching

Simple datapa

- ..

Full datapath for single cycle RISC-V





Fetching instruction and increment

Introduction

Synchronous Message Exchang

Implementing RISC-V in SMF

Single Cycle RISC-V Fetching instruction and increment

Instruction decod

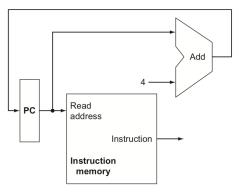
and execution

Branching

Simple datapat Instructions

- ..

We need memory for instructions and a way to remember current instruction





Instruction decode and execution

Introduction

Motivation

Motivation
Synchronous
Message Exchang

Implementing RISC-V in SME

Single Cycle RISC-V

Instruction decode and execution

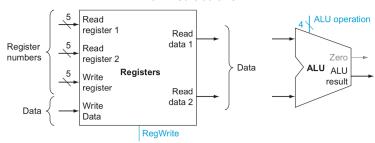
and execution

Simple datapath

Instructions

E....bloom ...oul

Storage for data we currently work with and unit for execution for instructions





Instruction decode and execution

Introduction

Motivation

Synchronous Message Exchange

Implementing RISC-V in SME

Fetching instruction

Instruction decode and execution

Memory access
Branching
Simple datapath

Instructions

Further wor

Conventions according to RISC-V spec sheet

Register	ABI Name	Description	Saver
x0	zero	Hard-wired zero	_
x1	ra	Return address	Caller
x2	sp	Stack pointer	Callee
x3	gp	Global pointer	_
x4	tp	Thread pointer	_
х5	t0	Temporary/alternate link register	Caller
x6-7	t1-2	Temporaries	Caller
x8	s0/fp	Saved register/frame pointer	Callee
x9	s1	Saved register	Callee
x10-11	a0-1	Function arguments/return values	Caller
x12-17	a2-7	Function arguments	Caller
x18-27	s2-11	Saved registers	Callee
x28-31	t3-6	Temporaries	Caller
f0-7	ft0-7	FP temporaries	Caller
f8-9	fs0-1	FP saved registers	Callee
f10-11	fa0-1	FP arguments/return values	Caller
f12-17	fa2-7	FP arguments	Caller
f18-27	fs2-11	FP saved registers	Callee
f28-31	ft8-11	FP temporaries	Caller

Taken from The RISC-V Instruction Set Manual - Volume I: User-Level ISA



Memory access

Introduction

Synchronous Message Exchange

Implementing RISC-V in SME

Fetching instruction and increment

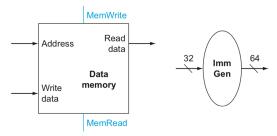
Memory access

Simple datapat

Instructio

Further wor

We need memory unit for storing data and immediate generation unit for calculating memory address





Branching

Introduction

Motivation

Synchronous Message Exchange

Implementing RISC-V in SME

Fetching instructionand increment
Instruction decode

and execution

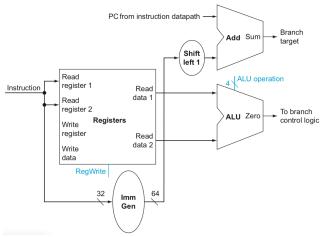
Memory access

Branching

Simple datapate

Further we

Additional unit needed for branching instructions





Simple datapath

Introduction Motivation

Synchronous Message Exchange

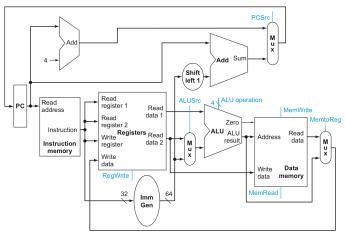
Implementing RISC-V in SME

Fetching instruction and increment
Instruction decode and execution

and execution Memory access

Simple datapath

Wiring it up we have our simple datapath





Instructions

Introduction Motivation

Synchronous Message Exchang

Implementing RISC-V in SME

Fetching instruction and increment
Instruction decoder and execution
Memory access

Simple datapar

Further wor

Instruction Formats

31	25	24	20	19	15	14	11	7 6		0
funct7		rs2		rs1		funct3	rd		opcode	
	imm[11:0]			rs1		funct3	rd		opcode	
imm[11:5]	imm[5]	imm[4:0]		rs1		funct3	rd		opcode	
imm[11:5		rs2		rs1		imm[4:0]	rd		opcode	
imm[12]10	:5]	rs2		rs1		imm[4:1 11]	rd		opcode	
imm[31:12]						rd		opcode		
imm[20 10:1 11 19:12]						rd		opcode		

^{*} This is a special case of the RV64I I-type format used by slli, srli and srai instructions where the lower 6 bits (imm[5] and imm[4:0]) are used to determine the shift amount (shamt). If slliw, srliw and sraiw are used it should generate an error if imm[5] \(\neq 0 \)

RV64I Base Instructions

Name	Fmt	Opcode	Funct3	Funct7/	Assembly	Description (in C)
				imm[11:5]		
Add	R	0110011	000	0000000	add rd, rs1, rs2	rd = rs1 + rs2
Subtract	R	0110011	000	0100000	sub rd, rs1, rs2	rd = rs1 - rs2
AND	R	0110011	111	0000000	and rd, rs1, rs2	rd = rs1 & rs2
OR	R	0110011	110	0000000	or rd, rs1, rs2	$rd = rs1 \mid rs2$
XOR	R	0110011	100	0000000	xor rd, rs1, rs2	$rd = rs1 \hat{r} rs2$



Simple datapath

Introduction Motivation

Synchronous Message Exchange

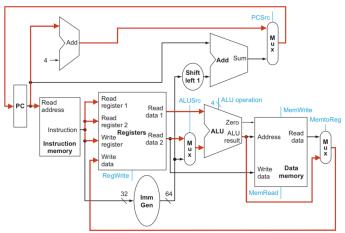
Implementing RISC-V in SME

Fetching instructio and increment Instruction decode and execution Memory access

Branching
Simple datapat
Instructions

Further wor

The R-Type instruction datapath would look like





Control

Introduction Motivation

Synchronous Message Exchang

Implementing RISC-V in SME

Fetching instructio and increment Instruction decode and execution

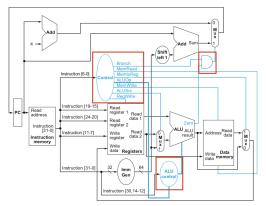
Memory access

Simple datapat

Control

Further we

We need control unit to determine path for instructions. Main Control and ALU control are separated to simplify complexity of system.



Taken from Computer organization and design RISC V



Further work

Introducti

Motivation

Synchronous Message Exchan

Implementin RISC-V in

Single Cycle RISC-Fetching instructio

Instruction decod

and execution

Memory access

Branching Simple datapat

Instructions

Further work

- Implement rest of base instructions
- Pipeline the datapath
- Test on a FPGA