## Risc V Reference Card

#### **Instruction Formats**

31		$^{25}$	24	20	19		15	14	12	11	7	6	0	
	funct7		rs2			rs1		func	t3		rd	opcode		R-type
	imm[11:0]				rs1 funct3			$_{\mathrm{rd}}$	opcode		I-type			
	imm[11:6]		imm[5:0]			rs1		func	t3		$^{\mathrm{rd}}$	opcode		$I$ -type $^*$
	imm[11:5]		rs2			rs1		func	:t3	j	imm[4:0]	opcode		S-type
	imm[12 10:5]		rs2			rs1		func	t3	in	nm[4:1 11]	opcode		B-type
	imm[31:12]										$_{\mathrm{rd}}$	opcode		U-type
imm[20 10:1 11 19:12]											rd	opcode		$_{ m J-type}$

<sup>\*</sup> This is a special case of the RV64I I-type format used by slli, srli and srai instructions where the lower 6 bits in the immediate are used to determine the shift amount (shamt). If slliw, srliw and sraiw are used it should generate an error if  $imm[6] \neq 0$ 

### **RV64I Base Instructions**

Name	Fmt	Opcode	Funct3	Funct7/ imm[11:5]	Assembly	Description (in C)
Add	R	0110011	000	0000000	add rd, rs1, rs2	rd = rs1 + rs2
Subtract	R	0110011	000	0100000	sub rd, rs1, rs2	rd = rs1 - rs2
AND	R	0110011	111	0000000	and rd, rs1, rs2	rd = rs1 & rs2
OR	R	0110011	110	0000000	or rd, rs1, rs2	$rd = rs1 \mid rs2$
XOR	R	0110011	100	0000000	xor rd, rs1, rs2	$rd = rs1 \hat{r} rs2$
Shift Left Logical	R	0110011	001	0000000	sll rd, rs1, rs2	$rd = rs1 \ll rs2$
Set Less Than	R	0110011	010	0000000	slt rd, rs1, rs2	rd = (rs1 < rs2)?1:0
Set Less Than (U)*	R	0110011	011	0000000	sltu rd, rs1, rs2	rd = (rs1 < rs2)?1:0
Shift Right Logical	R	0110011	101	0000000	srl rd, rs1, rs2	$rd = rs1 \gg rs2$
Shift Right Arithmetic <sup>†</sup>	R	0110011	101	0100000	sra rd, rs1, rs2	$rd = rs1 \gg rs2$
Add Word	R	0111011	000	0000000	addw rd, rs1, rs2	rd = rs1 + rs2
Subtract Word	R	0111011	000	0100000	subw rd, rs1, rs2	rd = rs1 + rs2 rd = rs1 - rs2
Shift Left Logical Word	R	0111011	001	0000000	sllw rd, rs1, rs2	$rd = rs1 \ll rs2$
Shift Right Logical Word	R	0111011	101	0000000	srlw rd, rs1, rs2	$rd = rs1 \gg rs2$
Shift Right Arithmetic Word <sup>†</sup>	R	0111011	101	0100000	sraw rd, rs1, rs2	$rd = rs1 \gg rs2$
Add Immediate	I	0010011	000	0100000	addi rd, rs1, imm	rd = rs1 # is2 $rd = rs1 + imm$
AND Immediate	I	0010011	111		and rd, rs1, imm	rd = rs1 & imm
OR Immediate	I	0010011	110		or rd, rs1, imm	rd = rs1 & mm
XOR Immediate	Ī	0010011	100		xor rd, rs1, imm	rd = rs1 $nmm$
Shift Left Logical Immediate	Ī	0010011	001	0000000	slli rd, rs1, shamt	$rd = rs1 \ll shamt$
Shift Right Logical Immediate	Ī	0010011	101	0000000	srli rd, rs1, shamt	$rd = rs1 \gg shamt$
Shift Right Arithmetic Immediate <sup>†</sup>	Ī	0010011	101	0100000	srai rd, rs1, shamt	$rd = rs1 \gg shamt$
Set Less Than Immediate	I	0010011	010	0100000	slti rd, rs1, imm	rd = 181 % shant rd = (rs1 < imm)?1:0
Set Less Than Immediate (U)*	I	0010011	011		sltiu rd, rs1, imm	rd = (rs1 < imm) : 1:0 rd = (rs1 < imm) : 1:0
Add Immediate Word	I	0010011	000		addiw rd, rs1, imm	rd = (rs1 < mm) : 1:0 rd = rs1 + imm
Shift Left Logical Immediate Word	I	0011011	000	0000000	slliw rd, rs1, shamt	rd = rs1 + mm $rd = rs1 \ll shamt$
Shift Right Logical Immediate Word	I	0011011	101	0000000	srliw rd, rs1, shamt	$rd = rs1 \ll shamt$ $rd = rs1 \gg shamt$
Shift Right Arithmetic Imm Word <sup>†</sup>	I	0011011	101	0100000	· '	
Load Byte	Ī	0000011	000	0100000	sraiw rd, rs1, shamt	$rd = rs1 \gg shamt$ rd = M[rs1+imm][0:7]
Load Byte Load Half	I	0000011	000		lb rd, rs1, imm lh rd, rs1, imm	rd = M[rs1+imm][0:7] rd = M[rs1+imm][0:15]
Load Word	I	0000011	010		lw rd, rs1, imm	rd = M[rs1+imm][0:15] rd = M[rs1+imm][0:31]
Load Word Load Doubleword	I	0000011	010		ld rd, rs1, imm	$  \mathbf{rd} = \mathbf{M}[\mathbf{rs1} + \mathbf{mm}][0.31]$ $  \mathbf{rd} = \mathbf{M}[\mathbf{rs1} + \mathbf{mm}][0.63]$
	I	ł				1 1 1
Load Byte (U)*		0000011	100		lbu rd, rs1, imm	rd = M[rs1+imm][0:7]
Load Half (U)*	I	0000011	101		lhu rd, rs1, imm	rd = M[rs1+imm][0:15]
Load Word (U) <sup>*</sup>	I	0000011	110		lwu rd, rs1, imm	rd = M[rs1+imm][0:31]
Store Byte	S	0100011	000		sb rs1, rs2, imm	M[rs1+imm][0:7] = rs2[0:7]
Store Half	S	0100011	001		sh rs1, rs2, imm	M[rs1+imm][0:15] = rs2[0:15]
Store Word	S	0100011	010		sw rs1, rs2, imm	M[rs1+imm][0:31] = rs2[0:31]
Store Doubleword	S	0100011	011		sd rs1, rs2, imm	M[rs1+imm][0:63] = rs2[0:63]
Branch If Equal	В	1100011	000		beq rs1, rs2, imm	if(rs1 == rs2) PC += imm
Branch Not Equal	В	1100011	001		bne rs1, rs2, imm	if(rs1 != rs2) PC += imm
Branch Less Than	В	1100011	100		blt rs1, rs2, imm	if(rs1 < rs2) PC += imm
Branch Greater Than Or Equal	В	1100011	101		bge rs1, rs2, imm	$if(rs1 \ge rs2) PC += imm$
Branch Less Than (U)*	В	1100011	110		bltu rs1, rs2, imm	if(rs1 < rs2) PC += imm
Branch Greater Than Or Equal $(U)^*$	В	1100011	111		bgeu rs1, rs2, imm	$if(rs1 \ge rs2) PC += imm$
Load Upper Immediate	U	0110111			lui rd, imm	$rd = imm \ll 12$
Add Upper Immediate To PC	U	0010111			auipc rd, imm	$rd = PC + (imm \ll 12)$
Jump And Link	J	1101111			jal rd, imm	rd = PC + 4; $PC += imm$
Jump And Link Register	I	1100111	000		jalr rd, rs1, imm	rd = PC + 4; $PC = rs1 + imm$

<sup>\*</sup>Assumes values are unsigned integers and zero extends  $\dagger$  Fills in with sign bit during right shift and msb (most significant bit) extends

### **RV64M Standard Extension Instructions**

Name	Fmt	Opcode	Funct3	Funct7	Assembly	Description (in C)
Multiply	R	0110011	000	0000001	mul rd, rs1, rs2	$rd = (rs1 \cdot rs2)[63:0]$
Multiply Upper Half	R	0110011	001	0000001	mulh rd, rs1, rs2	$rd = (rs1 \cdot rs2)[127:64]$
Multiply Upper Half Sign/Unsigned <sup>†</sup>	R	0110011	010	0000001	mulhsu rd, rs1, rs2	$rd = (rs1 \cdot rs2)[127:64]$
Multiply Upper Half (U)*	R	0110011	011	0000001	mulhu rd, rs1, rs2	$rd = (rs1 \cdot rs2)[127:64]$
Divide	R	0110011	100	0000001	div rd, rs1, rs2	rd = rs1 / rs2
Divide (U)*	R	0110011	101	0000001	divu rd, rs1, rs2	rd = rs1 / rs2
Remainder	R	0110011	110	0000001	rem rd, rs1, rs2	rd = rs1 % rs2
Remainder (U)*	R	0110011	111	0000001	remu rd, rs1, rs2	rd = rs1 % rs2
Multiply Word	R	0111011	000	0000001	mulw rd, rs1, rs2	$rd = (rs1 \cdot rs2)[63:0]$
Divide Word	R	0111011	100	0000001	divw rd, rs1, rs2	rd = rs1 / rs2
Divide Word (U)*	R	0111011	101	0000001	divuw rd, rs1, rs2	rd = rs1 / rs2
Remainder Word	R	0111011	110	0000001	remw rd, rs1, rs2	rd = rs1 % rs2
Remainder Word (U)*	R	0111011	111	0000001	remuw rd, rs1, rs2	rd = rs1 % rs2

<sup>\*</sup>Assumes values are unsigned integers and zero extends  $^\dagger$  Multiply with one operand signed and the other unsigned

# Bibliography

- [1] C. A. R. Hoare. Communicating sequential processes. Communications of the ACM, 21(8):666–677, 1978. ISSN 1557-7317.
- [2] B. Vinter and K. Skovhede. Synchronous message exchange for hardware designs.  $Communicating\ Process\ Architectures$ , pages 201–212, 2014.