Integer arithmetic - Multiplication

Computer Systems 5th Lecture, Sep 10 2017

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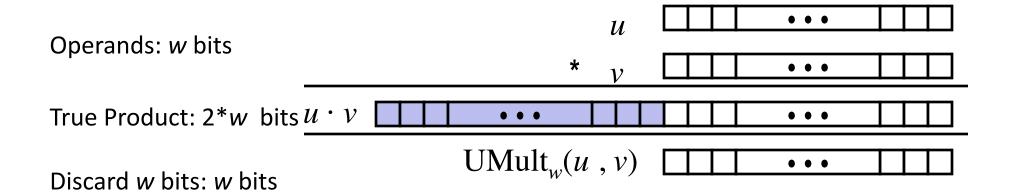
Based on slides by:

Randal E. Bryant and David R. O'Hallaron

Multiplication

- Goal: Computing Product of w-bit numbers x, y
 - Either signed or unsigned
- But, exact results can be bigger than w bits
 - Unsigned: up to 2w bits
 - Result range: $0 \le x * y \le (2^w 1)^2 = 2^{2w} 2^{w+1} + 1$
 - Two's complement min (negative): Up to 2w-1 bits
 - Result range: $x * y \ge (-2^{w-1})*(2^{w-1}-1) = -2^{2w-2} + 2^{w-1}$
 - Two's complement max (positive): Up to 2w bits, but only for $(TMin_w)^2$
 - Result range: $x * y \le (-2^{w-1})^2 = 2^{2w-2}$
- So, maintaining exact results...
 - would need to keep expanding word size with each product computed
 - is done in software, if needed
 - e.g., by "arbitrary precision" arithmetic packages

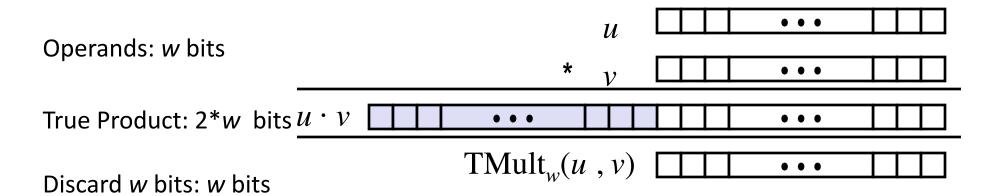
Unsigned Multiplication in C



- Standard Multiplication Function
 - Ignores high order w bits
- **Implements Modular Arithmetic**

$$UMult_{w}(u, v) = u \cdot v \mod 2^{w}$$

Signed Multiplication in C



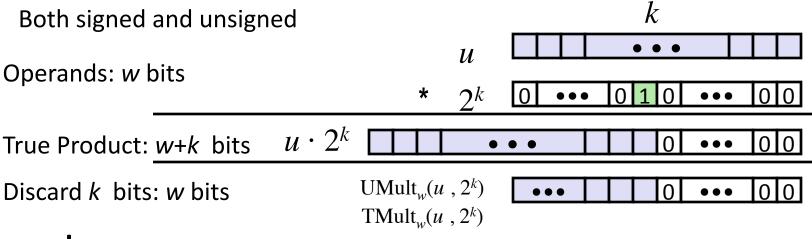
Standard Multiplication Function

- Ignores high order w bits
- Some of which are different for signed vs. unsigned multiplication
- Lower bits are the same

Power-of-2 Multiply with Shift

Operation

- **u** << **k** gives **u** * **2**^k
- Both signed and unsigned



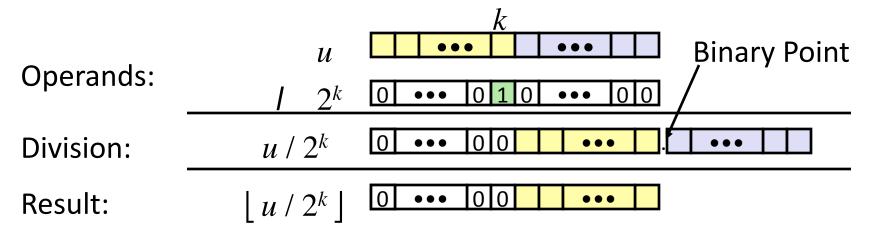
Examples

$$u << 5$$
 - $u << 3$ == $u * 24$

- Most machines shift and add faster than multiply
 - Compiler generates this code automatically

Unsigned Power-of-2 Divide with Shift

- Quotient of Unsigned by Power of 2
 - $\mathbf{u} \gg \mathbf{k}$ gives $\left[\mathbf{u} / 2^{k}\right]$
 - Uses logical shift



	Division	Computed	Hex	Binary
x	15213	15213	3B 6D	00111011 01101101
x >> 1	7606.5	7606	1D B6	00011101 10110110
x >> 4	950.8125	950	03 B6	00000011 10110110
x >> 8	59.4257813	59	00 3B	00000000 00111011

Unsigned Division C

Operands: w bits

True Product: 2*w bits

Discard w bits: w bits

- Standard Multiplication Function
 - Ignores high order w bits
- **Implements Modular Arithmetic**

$$UMult_{w}(u, v) = u \cdot v \mod 2^{w}$$

Today: Integer arithmetic and floating point

- Recap
- Integer arithmetic
 - Addition, negation, multiplication, shifting
 - Summary
- Floating Points

Arithmetic: Basic Rules

Addition:

- Unsigned/signed: Normal addition followed by truncate, same operation on bit level
- Unsigned: addition mod 2^w
 - Mathematical addition + possible subtraction of 2^w
- Signed: modified addition mod 2^w (result in proper range)
 - Mathematical addition + possible addition or subtraction of 2^w
- Fast execution (usually 1 clock cycle)

Multiplication:

- Unsigned/signed: Normal multiplication followed by truncate, same operation on bit level
- Unsigned: multiplication mod 2^w
- Signed: modified multiplication mod 2^w (result in proper range)
- Slower execution (usually multiple clock cycles)

Why Should I Use Unsigned?

- Don't use without understanding implications
 - Easy to make mistakes

```
unsigned i;
for (i = cnt-2; i >= 0; i--)
  a[i] += a[i+1];
```

Can be very subtle

```
#define DELTA sizeof(int)
int i;
for (i = CNT; i-DELTA >= 0; i-= DELTA)
. . .
```

Counting Down with Unsigned

Proper way to use unsigned as loop index

```
unsigned i;
for (i = cnt-2; i < cnt; i--)
  a[i] += a[i+1];</pre>
```

- See Robert Seacord, Secure Coding in C and C++
 - C Standard guarantees that unsigned addition will behave like modular arithmetic
 - $0-1 \rightarrow UMax$
- Even better

```
size_t i;
for (i = cnt-2; i < cnt; i--)
   a[i] += a[i+1];</pre>
```

- Data type size t defined as unsigned value with length = word size
- Code will work even if cnt = UMax
- What if cnt is signed and < 0?</p>

Why Should I Use Unsigned? (cont.)

- Do Use When Performing Modular Arithmetic
 - Multiprecision arithmetic
- *Do* Use When Using Bits to Represent Sets
 - Logical right shift, no sign extension