

# Models of Lambek calculus with subexponentials

Daniel Rogozin<sup>1,2</sup>

<sup>1</sup>Lomonosov Moscow State University

<sup>2</sup>Serokell OÜ

**Abstract**

## 1 The Lambek Calculus with subexponentials

**Definition 1.** A subexponential signature is an ordered quintuple:

$$\Sigma = \langle \mathcal{I}, \leq, \mathcal{W}, \mathcal{C}, \mathcal{E} \rangle,$$

where  $I = \{s_1, \dots, s_n\}$ ,  $\mathcal{I}, \leq$  is a preorder.  $\mathcal{W}, \mathcal{C}, \mathcal{E}$  are upwardly closed subsets of  $I$  and  $\mathcal{W} \cap \mathcal{C} \subseteq \mathcal{E}$ .

**Definition 2.**

$$\mathcal{F}_\Sigma ::= Tp \mid (\mathcal{F}_\Sigma / \mathcal{F}_\Sigma) \mid (\mathcal{F}_\Sigma \backslash \mathcal{F}_\Sigma) \mid (\mathcal{F}_\Sigma \bullet \mathcal{F}_\Sigma) \mid (\mathcal{F}_\Sigma \vee \mathcal{F}_\Sigma) \mid (\mathcal{F}_\Sigma \wedge \mathcal{F}_\Sigma) \mid (!_s \mathcal{F}_\Sigma)_{s \in \Sigma}$$

**Definition 3.** Noncommutative linear logic with subexponentials ( $SMALC_\Sigma$ ), where  $\Sigma$  is a subexponential signature.

$$\begin{array}{c} \overline{A \rightarrow A} \text{ }^{ax} \\ \frac{\Gamma \rightarrow A \quad \Delta, B, \Theta \rightarrow C}{\Delta, \Gamma, A \backslash B, \Theta \rightarrow C} \backslash \rightarrow \quad \frac{A, \Pi \rightarrow B}{\Pi \rightarrow A \backslash B} \rightarrow \backslash \\ \frac{\Gamma \rightarrow A \quad \Delta, B, \Theta \rightarrow C}{\Delta, B / A, \Gamma, \Theta \rightarrow C} / \rightarrow \quad \frac{\Pi, A \rightarrow B}{\Pi \rightarrow B / A} \rightarrow / \\ \frac{\Gamma, A, B, \Delta \rightarrow C}{\Gamma, A \bullet B, \Delta \rightarrow C} \bullet \rightarrow \quad \frac{\Gamma \rightarrow A \quad \Delta \rightarrow B}{\Gamma, \Delta \rightarrow A \bullet B} \rightarrow \bullet \\ \frac{\Gamma, A_i, \Delta \rightarrow B}{\Gamma, A_1 \& A_2, \Delta \rightarrow B} \&, i = 1, 2 \rightarrow \quad \frac{\Gamma \rightarrow A \quad \Gamma \rightarrow B}{\Gamma \rightarrow A \& B} \rightarrow \& \\ \frac{\Gamma, A, \Delta \rightarrow C \quad \Gamma, B, \Delta \rightarrow C}{\Gamma, A \vee B, \Delta \rightarrow C} \vee \rightarrow \quad \frac{\Gamma \rightarrow A_i}{\Gamma \rightarrow A_1 \vee A_2} \rightarrow \vee, i = 1, 2 \\ \frac{\Gamma, \Delta \rightarrow A}{\Gamma, 1, \Delta \rightarrow A} 1 \rightarrow \quad \overline{\rightarrow 1} \rightarrow 1 \\ \frac{\Gamma, A, \Delta \rightarrow C}{\Gamma, !^s A, \Delta \rightarrow C} ! \rightarrow \quad \frac{!^{s_1} A_1, \dots, !^{s_n} A_n \rightarrow A}{!^{s_1} A_1, \dots, !^{s_n} A_n \rightarrow !^s A} \rightarrow !, \forall j, s_j \geq s \end{array}$$

Structural rules:

$$\begin{array}{c}
\frac{\Gamma, !^s A, \Delta, !^s A, \Theta \rightarrow B}{\Gamma, !^s A, \Delta, \Theta \rightarrow B} \mathbf{ncontr}_1, s \in C \qquad \frac{\Gamma, !^s A, \Delta, !^s A, \Theta \rightarrow B}{\Gamma, \Delta, !^s A, \Theta \rightarrow B} \mathbf{ncontr}_2, s \in C \\
\\
\frac{\Gamma, \Delta, !^s A, \Theta \rightarrow B}{\Gamma, !^s A, \Delta, \Theta \rightarrow A} \mathbf{ex}_1, s \in E \qquad \frac{\Gamma, !^s A, \Delta, \Theta \rightarrow B}{\Gamma, \Delta, !^s A, \Theta \rightarrow A} \mathbf{ex}_2, s \in E \\
\\
\frac{\Gamma, \Delta \rightarrow B}{\Gamma, !^s A, \Delta \rightarrow B} \mathbf{weak}_1, s \in C \qquad \frac{\Gamma \rightarrow A \quad \Pi, A, \Delta \rightarrow B}{\Gamma, \Pi, \Delta \rightarrow B} \mathbf{cut}
\end{array}$$

**Theorem 1.**

1. Cut-rule is admissible;
2.  $SMALC_\Sigma$  is undecidable, if  $C \neq \emptyset$ ;
3. If  $C$  is empty, then the decidability problem of  $SMALC_\Sigma$  belongs to  $PSPACE$ .

## 2 Semantics

**Definition 4.** *Quantale*

A quantale is a triple  $\mathcal{Q} = \langle A, \bigvee, \cdot \rangle$ , such that  $\langle A, \bigvee \rangle$  is a complete lattice and  $\langle A, \cdot \rangle$  is a semigroup, such that for all indexing set  $I$ :

1.  $a \cdot \bigvee_{i \in I} b_i = \bigvee_{i \in I} (a \cdot b_i)$ ;
2.  $\bigvee_{i \in I} a_i \cdot b = \bigvee_{i \in I} (a_i \cdot b)$

A quantale is called unital, if  $\langle A, \cdot \rangle$  is a monoid.

Some example of quantales:

- Let  $A$  be a semigroup (monoid), then  $\langle \mathcal{P}(A), \cdot, \subseteq \rangle$  is a free (unital) quantale.
- Let  $R$  be a ring and  $Sub(R)$  be a set of additive subgroups of  $R$ . We define  $A \cdot B$  as an additive subgroup generated by finite sums of products  $ab$  and order is defined by inclusion.
- Any locale is a quantale with  $\cdot = \wedge$ .

It is easy to see, that any (unital) quantale is a residual (monoid) semigroup. We define divisions as follows:

1.  $a \backslash b = \bigvee \{c \mid a \cdot c \leq b\}$
2.  $b / a = \bigvee \{c \mid c \cdot a \leq b\}$

**Definition 5.** Let  $\mathcal{Q}_1, \mathcal{Q}_2$  be quantales. A quantale homomorphism is a map  $f : \mathcal{Q}_1 \rightarrow \mathcal{Q}_2$ , such that:

1. for all  $a, b \in \mathcal{Q}_1$ ,  $f(a \cdot b) = f(a) \cdot f(b)$ ;
2. for all indexing set  $I$ ,  $f(\bigvee_{i \in I} a_i) = \bigvee_{i \in I} f(a_i)$ .

If  $\mathcal{Q}_1, \mathcal{Q}_2$  are unital quantales, then a unital homomorphism is a quantale homomorphism, such that  $f(\varepsilon) = \varepsilon$ .

**Definition 6.**

Let  $\mathcal{Q} = \langle A, \bigvee, \cdot \rangle$  be a quantale.  $\mathcal{S} \subseteq \mathcal{Q}$  is said to be a subquantale, if  $\mathcal{S}$  is closed under multiplication and sups.

There occurs the following simple statement:

**Proposition 1.**

Let  $\mathcal{Q}_1, \mathcal{Q}_2$  be quantales and  $\mathcal{S} \subseteq \mathcal{Q}_1$  is a subquantale of  $\mathcal{Q}_1$ .

Then, if  $f : \mathcal{Q}_1 \rightarrow \mathcal{Q}_2$  is a quantale homomorphism, then  $f(\mathcal{S}) \subseteq \mathcal{Q}_2$  is a subquantale of  $\mathcal{Q}_2$ .

In other words, a homomorphic image of subquantale is a subquantale.

*Proof.*

It is clearly that  $f(\mathcal{S}) \subseteq \mathcal{Q}_2$  is a submonoid of  $\mathcal{Q}_2$ . Let  $a_i \in \mathcal{S}$  for each  $i \in I$ , so  $\bigvee_{i \in I} a_i \in \mathcal{S}$ , but  $f(a_i) \in f(\mathcal{S})$  for any  $i \in I$ , so  $f(\bigvee_{i \in I} a_i) = \bigvee_{i \in I} (f(a_i)) \in f(\mathcal{S})$ , so  $f(\mathcal{S})$  is closed under joins, so  $f(\mathcal{S})$  is a subquantale of  $\mathcal{Q}_2$   $\square$

**Definition 7.**

Let  $\mathcal{Q} = \langle A, \bigvee, \cdot \rangle$  be a quantale. The center of a quantale is the subquantale  $\mathcal{Z}(\mathcal{Q}) = \{a \in A \mid \forall b \in A, a \cdot b = b \cdot a\}$

**Definition 8.**

An open modality (or quantic conucleus) on quantale  $\mathcal{Q}$  is a map  $\Box : \mathcal{Q} \rightarrow \mathcal{Q}$ , such that

1.  $\Box x \leq x$ ;
2.  $\Box x = \Box \Box x$ ;
3.  $x \leq y \Rightarrow \Box x \leq \Box y$ ;
4.  $\Box x \cdot \Box y = \Box(\Box x \cdot \Box y)$ .

For unital quantale, we require that  $\Box e = e$ .

Note that, we may replace the last condition on equivalent condition  $\Box(x) \cdot \Box(y) \leq \Box(x \cdot y)$ .

**Definition 9.**

We define a partial order on open modalities on  $\mathcal{Q}$  as  $\Box_1 \leq \Box_2 \Leftrightarrow \forall a \in \mathcal{Q}, \Box_1(a) \leq \Box_2(a)$ .

**Lemma 1.**  $\Box_1 a_1 \cdot \Box_2 a_2 \leq \Box(\Box_1 a_1 \cdot \Box_2 a_2)$ , where  $\Box_i \leq \Box, i = 1, 2$ .

*Proof.*

$$\begin{aligned} \Box_1 a_1 \cdot \Box_2 a_2 &\leq \\ \Box_1(\Box_1 a_1) \cdot \Box_2(\Box_2 a_2) &\leq \\ \Box(\Box_1 a_1) \cdot \Box(\Box_2 a_2) &\leq \\ \Box(\Box_1(a_1) \cdot \Box_2(a_2)) &\leq \end{aligned}$$

$\square$

**Definition 10.**

1. An open modality is called central, if for all  $a, b \in \mathcal{Q}, \Box a \cdot b = b \cdot \Box a$ .
2. An open modality is called weak square-increasing, if for all  $a, b \in \mathcal{Q}, \Box a \cdot b \leq \Box a \cdot b \cdot \Box a$  and  $b \cdot \Box a \leq \Box a \cdot b \cdot \Box a$ .

3. An open modality is called *unital*, if  $\forall a \in Q, \Box a \leq e$ .

**Lemma 2.**

Let  $\Box$  be an open modality on some unital quantale  $\mathcal{Q} = \langle A, \bigvee, \cdot, e \rangle$ . Then, if  $\Box$  is unital and weak idempotent, then  $\Box$  is central.

*Proof.*

$$\begin{aligned}
& b \cdot \Box a \leq \\
& \quad \text{Right weak square-increase} \\
& \Box a \cdot b \cdot \Box a \leq \\
& \quad \text{Unitality} \\
& \Box a \cdot b \cdot e \leq \\
& \quad \text{Identity} \\
& \Box a \cdot b \leq \\
& \quad \text{Left weak square-increase} \\
& \Box a \cdot b \cdot \Box a \leq \\
& \quad \text{Unitality} \\
& e \cdot b \cdot \Box a \leq \\
& \quad \text{Identity} \\
& b \cdot \Box a
\end{aligned}$$

Hence,  $b \cdot \Box a = \Box a \cdot b$ , so for all  $a \in \mathcal{Q}, \Box a \in \mathcal{Z}(\mathcal{Q})$ .

□

**Proposition 2.**

Let  $\mathcal{Q}$  be a quantale and  $S \subseteq \mathcal{Q}$  a subquantale, then  $\Box : \mathcal{Q} \rightarrow \mathcal{Q}$ , such that  $\Box a = \bigvee \{s \in S \mid s \leq a\}$ , is an open modality.

*Proof.* See Rosenthal.

□

**Proposition 3.**

Let  $\mathcal{Q}$  be a quantale and  $S_1 \subseteq S_2 \subseteq \mathcal{Q}$ .

Then  $\Box_{S_1}(a) \leq \Box_{S_2}(a)$ .

*Proof.* Immediately.

□

**Proposition 4.**

Let  $\mathcal{Q}$  be a quantale and  $S \subseteq \mathcal{Q}$  a subquantale, then the following operations are open modalities:

1.  $\Box_z(a) = \bigvee \{s \in S \mid s \leq a, s \in \mathcal{Z}(\mathcal{Q})\};$
2.  $\Box_1(a) = \bigvee \{s \in S \mid s \leq a, s \leq 1\};$
3.  $\Box_{idem}(a) = \bigvee \{s \in S \mid s \leq a, \forall b \in \mathcal{Q}, b \cdot s \vee s \cdot b \leq s \cdot b \cdot s\};$
4. Operations with combinations of conditions above;

*Proof.* Immediately.

□

**Proposition 5.**

1.  $\forall a \in \mathcal{Q}, \Box_{1, idem}(a) \leq \Box_z(a).$
2.  $\forall a \in \mathcal{Q}, \Box_{z, 1, idem} = \Box_{1, idem}(a)$

*Proof.* Follows from Lemma 3. □

**Definition 11.** *Interpretation of subexponential signature*

Let  $\Sigma = \langle I, \leq, W, C, E \rangle$  be a subexponential signature, where  $|I| = n$  and  $\square_{\mathcal{Q}}$  is a category of open modalities on a quantale  $\mathcal{Q}$ . Subexponential interpretation is a contravariant functor  $\sigma : I \rightarrow \square_{\mathcal{Q}}$  defined as follows:

$$\sigma(s_i) = \begin{cases} \square_i : \mathcal{Q} \rightarrow \mathcal{Q}, \text{ s.t. } \forall a \in \mathcal{Q}, \square_i(a) = \{s \in S_i \mid s \leq a\}, \\ \quad \text{if } s_i \notin W \cap C \cap E \\ \square_i : \mathcal{Q} \rightarrow \mathcal{Q}, \text{ s.t. } \forall a \in \mathcal{Q}, \square_i(a) = \{s \in S_i \mid s \leq a, s \leq \mathbb{1}\}, \\ \quad \text{if } s_i \in W \\ \square_i : \mathcal{Q} \rightarrow \mathcal{Q}, \text{ s.t. } \forall a \in \mathcal{Q}, \square_i(a) = \{s \in S_i \mid s \leq a, s \in \mathcal{Z}(\mathcal{Q})\}, \\ \quad \text{if } s_i \in E \\ \square_i : \mathcal{Q} \rightarrow \mathcal{Q}, \text{ s.t. } \forall a \in \mathcal{Q}, \square_i(a) = \{s \in S_i \mid s \leq a, \forall b, b \cdot s \vee s \cdot b \leq s \cdot b \cdot s\}, \\ \quad \text{if } s_i \in E \\ \text{otherwise, if } s_i \text{ belongs to some intersection of subsets, then we combine the relevant conditions} \end{cases}$$

**Definition 12.** Let  $\mathcal{Q}$  be an unital quantale,  $f : Tp \rightarrow \mathcal{Q}$  a valuation and  $\sigma : I \rightarrow \square_{\mathcal{Q}}$  a subexponential interpretation, then interpretation is defined inductively:

$$\begin{aligned} \llbracket p_i \rrbracket &= f(p_i) \\ \llbracket \mathbb{1} \rrbracket &= e \\ \llbracket A \bullet B \rrbracket &= \llbracket A \rrbracket \cdot \llbracket B \rrbracket \\ \llbracket A \setminus B \rrbracket &= \llbracket A \rrbracket \setminus \llbracket B \rrbracket \\ \llbracket A / B \rrbracket &= \llbracket A \rrbracket / \llbracket B \rrbracket \\ \llbracket A \& B \rrbracket &= \llbracket A \rrbracket \wedge \llbracket B \rrbracket \\ \llbracket A \vee B \rrbracket &= \llbracket A \rrbracket \vee \llbracket B \rrbracket \\ \llbracket !_{s_i} A \rrbracket &= \sigma(s_i) \llbracket A \rrbracket \end{aligned}$$

**Definition 13.**  $\Gamma \models A \Leftrightarrow \forall f, \forall \sigma, \llbracket \Gamma \rrbracket \leq \llbracket A \rrbracket$

**Theorem 2.**  $\Gamma \rightarrow A \Rightarrow \llbracket \Gamma \rrbracket \leq \llbracket A \rrbracket$

*Proof.* We consider the promotion case, the rest modal cases are immediatly shown.

Let  $!_{s_1} A_1, \dots, !_{s_n} A_n \rightarrow A$  and  $\forall i, s \leq s_i$ .

Then  $\forall a \in \mathcal{Q}, \sigma(s_i)(a) \leq \sigma(s)(a)$ .

By IH,  $\sigma(s_1) \llbracket A_1 \rrbracket \cdots \sigma(s_n) \llbracket A_n \rrbracket \leq \llbracket A \rrbracket$ .

Thus,  $\sigma(s)(\sigma(s_1) \llbracket A_1 \rrbracket \cdots \sigma(s_n) \llbracket A_n \rrbracket) \leq \sigma(s)(\llbracket A \rrbracket)$ .

By Lemma 5,  $\sigma(s_1) \llbracket A_1 \rrbracket \cdots \sigma(s_n) \llbracket A_n \rrbracket \leq \sigma(s)(\sigma(s_1) \llbracket A_1 \rrbracket \cdots \sigma(s_n) \llbracket A_n \rrbracket)$ .

So,  $\sigma(s_1) \llbracket A_1 \rrbracket \cdots \sigma(s_n) \llbracket A_n \rrbracket \leq \sigma(s)(\llbracket A \rrbracket)$ . □

### 3 Quantale completeness

**Definition 14.**

Let  $\mathcal{F} \subseteq Fm$ , an ideal is a subset  $\mathcal{I} \subseteq \mathcal{F}$ , such that:

- If  $B \in \mathcal{I}$  and  $A \rightarrow B$ , then  $A \in \mathcal{I}$ ;
- If  $A, B \in \mathcal{I}$ , then  $A \vee B \in \mathcal{I}$ .

**Definition 15.**

Let  $S \subseteq \mathcal{F} \subseteq Fm$ , then  $\bigvee S = \bigcap \{\mathcal{I} \subseteq \mathcal{F} \mid S \subseteq \mathcal{I}\}$

The following conditions hold similarly to [?]:

**Lemma 3.**

1.  $\bigvee S$  is an ideal;
2.  $A \subseteq Fm$ , then  $\{B \mid B \rightarrow A\} = \bigvee \{A\}$ ;
3.  $\bigvee \{A\} \subseteq \bigvee \{B\}$  iff  $A \rightarrow B$ ;
4. Let  $\mathcal{Q} = \{\bigvee S \mid S \subseteq Fm\}$  and  $\bigvee \mathcal{A} \cdot \bigvee \mathcal{B} = \bigvee \{A \bullet B \mid A \in \mathcal{A}, B \in \mathcal{B}\}$ . Then  $\langle \mathcal{Q}, \subseteq, \cdot, \bigvee \mathbf{1} \rangle$  is a unital quantale.

We extend this construction for polymodal case as follows:

**Lemma 4.** Let  $!_s \in I$  and  $A \in \mathcal{F}_\Sigma$ , then  $\Box_s(\bigvee \{A\}) = \bigvee \{!_s B \mid !_s B \rightarrow A\}$  is a quantic conucleus.

*Proof.*

See Yetter. □

**Lemma 5.** Let  $A \in \mathcal{F}_\Sigma$ , then  $\Box_s \bigvee \{A\} = \bigvee \{!_s A\}$ , for each  $s \in \mathcal{I}$ .

*Proof.* Let  $A \in Fm$  and  $s \in \mathcal{I}$ .

Let  $!_s B \in \Box_s \bigvee \{A\}$ , then  $!_s B \rightarrow A$ , then  $!_s B \rightarrow !_s A$  by promotion. So,  $!_s B \in \bigvee \{!_s A\}$ .

Let  $C \in \bigvee \{!_s A\}$ , then  $C \rightarrow !_s A$ , so  $!_s C \rightarrow !_s A$  by dereliction, but  $!_s A \rightarrow A$ , hence  $!_s C \rightarrow A$  by cut. So,  $!_s C \in \Box_s \bigvee \{A\}$ . □

**Lemma 6.**

Let  $i, j \in I$  and  $i \leq j$ , then for all  $A \in \mathcal{F}_\Sigma$ ,  $\Box_j(\bigvee \{A\}) \subseteq \Box_i(\bigvee \{A\})$ .

*Proof.*

Let  $i, j \in I$  and  $i \leq j$ , then for all  $A \in \mathcal{F}_\Sigma$ ,  $!_j A \rightarrow !_i A$  by promotion. Then  $\bigvee \{!_j A\} \subseteq \bigvee \{!_i A\}$ , so  $\Box_j(\bigvee \{A\}) \subseteq \Box_i(\bigvee \{A\})$ . □

**Lemma 7.**

For all  $A \in \mathcal{F}_\Sigma$ ,

1. Let  $s \in W$ , then  $\Box_s \{A\} \subseteq \{\mathbf{1}\}$ ;
2. Let  $s \in E$ , then  $\Box_s(\bigvee \{A\}) \cdot \bigvee \{B\} = \bigvee \{B\} \cdot \Box_s(\bigvee \{A\})$ .
3. Let  $s \in C$ , then  $(\Box_s \bigvee A \cdot \bigvee B) \cup (\bigvee B \cdot \Box_s \bigvee A) \subseteq \Box_s \bigvee A \cdot \bigvee B \cdot \Box_s \bigvee A$ , for all  $B \subseteq Fm$ .

*Proof.*

1. Follows from  $!_s A \rightarrow \mathbf{1}$ , so  $s \in W$ ;
2. Follows from  $!_s A \bullet B \leftrightarrow B \bullet !_s A$ ;
3. Follows from  $!_s A \bullet B \rightarrow !_s A \bullet B \bullet !_s A$  and similarly for  $B \bullet !_s A$ .

□

**Definition 16.**

Let  $\mathcal{Q}$  be a syntactic quantale as proposed above and  $\mathcal{I} = \langle I, \leq, W, C, E \rangle$  be a subexponential signature.

We define a map  $\Box : \mathcal{I} \rightarrow \text{Mod}_{\mathcal{Q}}$  as follows:

$$\Box(i)(\bigvee\{A\}) = \{!_i B \mid !_i B \rightarrow A\}.$$

**Lemma 8.**  $\Box$  is a subexponential interpretation.

*Proof.* Follows from lemmas above. □

**Lemma 9.**

Let  $\mathcal{Q}$  be a quantale constructed above and  $(\Box_{s_i})_{s_i \in \Sigma}$  be a family of quantic conuclei on  $\mathcal{Q}$ . Then there exist a model  $\langle \mathcal{Q}, \llbracket \cdot \rrbracket \rangle$ , such that  $\llbracket A \rrbracket = \bigvee\{A\}$ ,  $A \in \text{Fm}$ .

*Proof.*

We define an interpretation as follows:

1.  $\llbracket p_i \rrbracket = \bigvee\{p_i\}$
2.  $\llbracket \mathbf{1} \rrbracket = \bigvee\{\mathbf{1}\}$
3.  $\llbracket A \bullet B \rrbracket = \bigvee\{A \bullet B\}$
4.  $\llbracket A/B \rrbracket = \bigvee\{A/B\}$
5.  $\llbracket B \setminus A \rrbracket = \bigvee\{B \setminus A\}$
6.  $\llbracket A \& B \rrbracket = \bigvee\{A \& B\}$
7.  $\llbracket A \vee B \rrbracket = \bigvee\{A \vee B\}$
8.  $\llbracket !_s A \rrbracket = \Box(s)(\bigvee\{A\}) = \bigvee\{!_s A\}.$

□

**Theorem 3.**  $\Gamma \models A \Rightarrow \Gamma \rightarrow A$ .

*Proof.* Follows from lemmas above. □

## 4 Relational semantics

**Definition 17.**

Let  $A$  be a set. Then relational quantale on  $A$  is a triple  $\mathcal{Q} = \langle \mathcal{A}, \bigvee, \mathcal{I} \rangle$  where  $\mathcal{A} \subseteq 2^{A \times A}$ :

- $\langle \mathcal{A}, \bigvee, \subseteq \rangle$  is a complete semi-lattice;
- Multiplication is defined as  $R \circ S = \{ \langle a, c \rangle \mid \exists b \in A, \langle a, b \rangle \in R \text{ and } \langle b, c \rangle \in S \}$
- $\langle \mathcal{A}, \circ, \mathcal{I} \rangle$  is a monoid;
- For each indexing set  $J$ ,  $R \circ \bigvee_{i \in I} S_i = \bigvee_{i \in I} (R \circ S_i)$  and  $\bigvee_{i \in I} R_i \circ S = \bigvee_{i \in I} (R_i \circ S)$ .

**Theorem 4.**

Let  $\mathcal{Q} = \langle A, \leq, \cdot, \bigvee \rangle$  be a unital quantale and  $\mathcal{S}$  is a subquantale of  $\mathcal{Q}$ .

Then  $\langle \mathcal{Q}, \Box_{\mathcal{S}} \rangle$  is isomorphic to some relational quantale of  $A$  with some quantic conucleus  $\hat{\Box}$ .

*Proof.*

Consider a relational quantale proposed by Brown and Gurr.

This quantale is 4-tuple  $\theta(\mathcal{Q}) = \langle \mathcal{R}, \subseteq, \circ, \bigvee \rangle$  defined as follows:

1.  $\theta(a) = \{ \langle b, c \rangle \mid b \leq a \cdot c \};$
2.  $\theta(a \cdot b) = \theta(a) \circ \theta(b);$
3.  $\theta(\bigvee_{i \in I} a_i) = \bigvee_{i \in I} \theta(a_i);$
4.  $\theta(\varepsilon) = \{ \langle b, c \rangle \mid b \cdot \varepsilon \leq c \} = \{ \langle b, c \rangle \mid b \leq c \}$

Let  $\mathcal{S} \subseteq \mathcal{Q}$ , so  $\Box_{\mathcal{S}} a := \bigvee \{ s \mid s \in \mathcal{S}, s \leq a \}$  is quantic conucleus.

So,  $\theta(\mathcal{S}) \subseteq \theta(\mathcal{Q})$  is a subquantale of  $\theta(\mathcal{Q})$ .

Let us define  $\hat{\Box}\theta(a) := \bigvee \{ \theta(s) \mid \theta(s) \in \theta(\mathcal{S}), \theta(s) \subseteq \theta(a) \}$ , so

$$\begin{aligned} \theta(\Box_{\mathcal{S}} a) &= \{ \langle p, q \rangle \mid p \leq \Box_{\mathcal{S}} a \cdot q \} = \\ &= \{ \langle p, q \rangle \mid p \leq \bigvee \{ s \mid s \in \mathcal{S}, s \leq a \} \cdot q \} = \end{aligned}$$

Homomorphism

$$\theta(\bigvee_{s \in \mathcal{S}, s \leq a} s) =$$

Homomorphism preserves sups

$$\bigvee_{s \in \mathcal{S}, s \leq a} \theta(s) =$$

Unfolding

$$\bigvee \{ \theta(s) \mid s \in \mathcal{S}, s \leq a \} =$$

Unfolding

$$\bigvee \{ \theta(s) \mid \theta(s) \in \theta(\mathcal{S}), \theta(s) \subseteq \theta(a) \} = \hat{\Box}\theta(a)$$

So,  $\hat{\Box}\theta(a) = \theta(\Box_{\mathcal{S}} a)$ .

□

## 5 Syntactic concept lattices

**Definition 18.** Let  $\mathcal{L}$  be a finite alphabet and  $L \subseteq \mathcal{L}^*$  be a language.

We define maps  $[\cdot]^\triangleright : \mathcal{P}(\mathcal{L}^*) \rightarrow \mathcal{P}(\mathcal{L}^* \times \mathcal{L}^*)$  and  $[\cdot]^\triangleleft : \mathcal{P}(\mathcal{L}^* \times \mathcal{L}^*) \rightarrow \mathcal{P}(\mathcal{L}^*)$  as follows:

1.  $M \subseteq \mathcal{L}^*, M^\triangleright = \{ (x, y) \mid \forall w \in M, xwy \in L \};$
2.  $C \subseteq \mathcal{L}^* \times \mathcal{L}^*, C^\triangleleft = \{ w \mid \forall (x, y) \in C, xwy \in L \}$

Note that compositions  $[\cdot]^\triangleleft^\triangleright$  and  $[\cdot]^\triangleright^\triangleleft$  form closure operators, by the way  $[\cdot]^\triangleleft$  and  $[\cdot]^\triangleright$  are connected via contravariant Galois connection.

**Definition 19.** A syntactic concept is a pair  $\langle S, C \rangle$ , where  $S \subseteq \mathcal{L}^*$  and  $C \subseteq \mathcal{L}^* \times \mathcal{L}^*$ , such that  $S^\triangleright = C$  and  $C^\triangleleft = S$ .