

**Definition 1.** *Monoidal comonad*

A monoidal comonad on some monoidal category  $\mathcal{C}$  is a triple  $\langle \mathcal{F}, \epsilon, \delta \rangle$ , where  $\mathcal{F}$  is a monoidal endofunctor and  $\epsilon : \mathcal{F} \Rightarrow \text{Id}_{\mathcal{C}}$  (counit) and  $\delta : \mathcal{F} \Rightarrow \mathcal{F}^2$  (comultiplication), such that the following diagrams commute:

$$\begin{array}{ccc}
 \mathcal{F}A \otimes \mathcal{F}B & \xrightarrow{\phi_{A,B}} & \mathcal{F}(A \otimes B) \\
 \delta_A \otimes \delta_B \downarrow & & \searrow \delta_{A \otimes B} \\
 \mathcal{F}\mathcal{F}A \otimes \mathcal{F}\mathcal{F}B & \xrightarrow{\phi_{\mathcal{F}A, \mathcal{F}B}} & \mathcal{F}(\mathcal{F}A \otimes \mathcal{F}B) \\
 & \nearrow \mathcal{F}(\phi_{A,B}) & \\
 & \mathcal{F}\mathcal{F}(A \otimes B) & 
 \end{array}
 \qquad
 \begin{array}{ccc}
 \mathcal{F}A \otimes \mathcal{F}B & \xrightarrow{\phi_{A,B}} & \mathcal{F}(A \otimes B) \\
 \epsilon_A \otimes \epsilon_B \searrow & & \swarrow \epsilon_{A \otimes B} \\
 & A \otimes B & 
 \end{array}$$
  

$$\begin{array}{ccc}
 \mathbb{1} & \xrightarrow{\phi} & \mathcal{F}\mathbb{1} \\
 \phi \downarrow & & \downarrow \delta_{\mathbb{1}} \\
 \mathcal{F}\mathbb{1} & \xrightarrow{\mathcal{F}(\phi)} & \mathcal{F}\mathcal{F}\mathbb{1} \\
 & & \downarrow \delta_{\mathcal{F}\mathbb{1}} \\
 & & \mathcal{F}\mathcal{F}\mathcal{F}\mathbb{1}
 \end{array}$$
  

$$\begin{array}{ccc}
 \mathbb{1} & \xrightarrow{id_{\mathbb{1}}} & \mathbb{1} \\
 \phi \searrow & & \swarrow \epsilon_{\mathbb{1}} \\
 & \mathcal{F}\mathbb{1} & 
 \end{array}$$

**Definition 2.** *Biclosed monoidal category*

Let  $\mathcal{C}$  be a monoidal category. Biclosed monoidal category is a monoidal category with the following additional data:

1. Bifunctors  $\_ \multimap \_ , \_ \multimap \_ : \mathcal{C}^{op} \times \mathcal{C} \rightarrow \mathcal{C}$ ;
2. Natural isomorphism  $\mathbf{curry}_{A,B,C} : \text{Hom}(A \otimes B, C) \cong (B, A \multimap C)$ ;
3. Natural isomorphism  $\mathbf{curry}'_{A,B,C} : \text{Hom}(A \otimes B, C) \cong (A, C \multimap B)$ ;
4. For each  $A, B \in \text{Ob}_{\mathcal{C}}$ , there are exist arrows  $ev_{A,B} : A \otimes (A \multimap B) \rightarrow B$  and  $ev'_{A,B} : (B \multimap A) \otimes A \rightarrow B$ , such that for all  $f : A \otimes C \rightarrow B$ :
  - (a)  $\Lambda_l \circ (id_A \otimes \mathbf{curry}(f)) = f$ ;
  - (b)  $\Lambda_r \circ (\mathbf{curry}'(f) \otimes id_A) = f$

**Definition 3.** Let  $F$  be endofunctor and  $A \in \text{Ob}_{\mathcal{C}}$ , then a coalgebra of  $F$  is a tuple  $\langle A, \theta \rangle$ , where  $\theta : A \rightarrow FA$ .

Given coalgebras  $\langle A, \theta \rangle$  and  $\langle B, \psi \rangle$ , a homomorphism is a morphism  $f : A \rightarrow B$ , s.t. the diagram below commutes:

$$\begin{array}{ccc}
 A & \xrightarrow{\theta} & FA \\
 f \downarrow & & \downarrow Ff \\
 B & \xrightarrow{\psi} & FB
 \end{array}$$

that is,  $Ff \circ \theta = \psi \circ f$

**Definition 4.** *Subexponential model structure*

Let  $\Sigma = \langle I, \leq, W, C, E \rangle$  be a subexponential signature and  $\mathcal{C}$  be a biclosed monoidal category, then a subexponential model structure is  $\langle \mathcal{C}, \{\mathcal{F}_s\}_{s \in I} \rangle$  with the following additional data:

- for all  $s \in I$ ,  $\mathcal{F}_s$  is a monoidal comonad;
- if  $s \in W$ , then for all  $A \in \text{Ob}(\mathcal{C})$ , there exists a morphism  $w_{As} : F_s A \rightarrow \mathbb{1}$ ;
- if  $s \in C$ , then for all  $A \in \text{Ob}(\mathcal{C})$ , there exists morphisms  $w_{Al} : F_s A \otimes A \otimes F_s A \rightarrow F_s A \otimes B$  and  $w_{Ar} : F_s A \otimes A \otimes F_s A \rightarrow B \otimes F_s A$ ;
- if  $s \in E$ , then for all  $A \in \text{Ob}(\mathcal{C})$ , there is an isomorphism,  $e_A : F_s A \otimes B \cong B \otimes F_s A$ ;
- if  $s_1 \in W$ ,  $s_2 \in I$  and  $s_1 \leq s_2$ , then there is a morphism  $w_{As_2} : F_{s_2} A \rightarrow \mathbb{1}$  for all  $A \in \text{Ob}(\mathcal{C})$  and ditto for  $E$  and  $C$ ;
- Let  $\bigotimes_{s \in J, i=0}^n F_s A$ , where  $J \subset I$ , and  $s' \in I$ , s.t.  $s \geq s'$  for all  $s \in J'$ ; Then there exists morphism a morphism  $\theta_{\bigotimes_{s \in J, i=1}^n F_{s_j} A_i} : \bigotimes_{s \in J, i=0}^n F_s A \rightarrow F_{s'}(\bigotimes_{s \in J, i=0}^n F_s A)$ , such that  $\langle \bigotimes_{s \in J, i=1}^n F_{s_j} A_i, \theta_{\bigotimes_{s \in J, i=1}^n F_{s_j} A_i} \rangle$  is a coalgebra on  $F_s$ .

**Definition 5.** Let  $\langle \mathcal{C}, \{\mathcal{F}_s\}_{s \in I} \rangle$  be a subexponential model structure for subexponential signature  $\Sigma = \langle I, \leq, W, C, E \rangle$ . Let  $v : Tp \rightarrow \text{Ob}(\mathcal{C})$  be a valuation map. Then the interpretation function  $\llbracket \cdot \rrbracket$  is defined as follows:

- (1)  $\llbracket \mathbb{1} \rrbracket = \mathbb{1}$
- (2)  $\llbracket A \setminus B \rrbracket = \llbracket A \rrbracket \multimap \llbracket B \rrbracket$
- (3)  $\llbracket A / B \rrbracket = \llbracket A \rrbracket \multimap \llbracket B \rrbracket$
- (4)  $\llbracket A \bullet B \rrbracket = \llbracket A \rrbracket \otimes \llbracket B \rrbracket$
- (5)  $\llbracket !_s A \rrbracket = F_s \llbracket A \rrbracket$

**Theorem 1.** *The following statements are equivalent:*

- $SMLC_\Sigma + (\text{cut}) \vdash \Gamma \Rightarrow A$
- $SMLC_\Sigma \vdash \Gamma \Rightarrow A$
- $\exists f, f : \llbracket \Gamma \rrbracket \rightarrow \llbracket A \rrbracket$

*Proof.*

(1)  $\Rightarrow$  (2): cut elimination.

- (2)  $\Rightarrow$  (3): Soundness:

$$\begin{array}{c}
 \overline{id_A : A \rightarrow A} \\
 \\
 \frac{f : \Gamma \rightarrow A \quad g : \Delta \otimes B \otimes \Theta \rightarrow C}{g \circ (id_\Delta \otimes (ev_{A, B_l} \circ (f \otimes id_{A \multimap B})) \otimes id_\Theta) : \Delta \otimes (\Gamma \otimes A \multimap B) \otimes \Theta \rightarrow C} \\
 \\
 \frac{f : A \otimes \Pi \rightarrow B}{\Lambda_l(f) : \Pi \rightarrow A \multimap B} \\
 \\
 \frac{f : \Gamma \rightarrow A \quad g : \Delta \otimes B \otimes \Theta \rightarrow C}{g \circ (id_\Delta \otimes (ev_{A, B_l} \circ (id_{B \multimap A} \otimes f)) \otimes id_\Theta) : \Delta \otimes (B \multimap A \otimes \Gamma) \otimes \Theta \rightarrow C}
 \end{array}$$

$$\begin{array}{c}
\frac{f : \Pi \otimes A \rightarrow B}{\Lambda_r(f) : \Pi \rightarrow B \multimap A} \\
\\
\frac{f : \Gamma \otimes A \otimes B \otimes \Delta \rightarrow C}{f \circ (\alpha_{\Gamma, A, B} \otimes id_{\Delta}) : \Gamma \otimes (A \otimes B) \otimes \Delta \rightarrow C} \\
\\
\frac{f : \Gamma \rightarrow A \quad g : \Delta \rightarrow B}{f \otimes g : \Gamma \otimes \Delta \rightarrow A \otimes B} \\
\\
\frac{f : \Gamma \otimes A_i \otimes \Delta \rightarrow B}{f \circ (id_{\Gamma} \otimes \pi_i id_{\Delta}) : \Gamma \otimes (A_1 \times A_2) \otimes \Delta \rightarrow B} \\
\\
\frac{f : \Gamma \rightarrow A \quad g : \Gamma \rightarrow B}{\langle f, g \rangle : \Gamma \rightarrow A \times B} \\
\\
\frac{f : \Gamma \otimes A \otimes \Delta \rightarrow C \quad g : \Gamma \otimes B \otimes \Delta \rightarrow C}{id_{\Gamma} \otimes [f, g] \otimes id_{\Delta} : \Gamma \otimes (A + B) \otimes \Delta \rightarrow C} \\
\\
\frac{}{id_{\mathbb{1}} : \mathbb{1} \rightarrow \mathbb{1}} \\
\\
\frac{f : \Gamma \otimes \Delta \rightarrow A}{f \circ (\rho_{\Gamma} \otimes id_{\Delta}) : (\Gamma \otimes \mathbb{1}) \otimes \Delta \rightarrow A} \\
\\
\frac{f : \Gamma \otimes A \otimes \Delta \rightarrow B}{f \circ (id_{\Gamma} \otimes \delta_s^A \otimes id_{\Delta}) : \Gamma \otimes F_s A \otimes \Delta \rightarrow B} \\
\\
\frac{f : F_{s_1} A_1 \otimes \cdots \otimes F_{s_n} A_n \rightarrow B}{F_s(f) : F_s(F_{s_1} A_1 \otimes \cdots \otimes F_{s_n} A_n) \rightarrow F_s B} \\
\\
\frac{F_s(f) \circ \theta_{\otimes_{s \in J, i=1}^n F_{s_j} A_i} : F_{s_1} A_1 \otimes \cdots \otimes F_{s_n} A_n \rightarrow F_s B}{F_s(f) \circ \theta_{\otimes_{s \in J, i=1}^n F_{s_j} A_i} : F_{s_1} A_1 \otimes \cdots \otimes F_{s_n} A_n \rightarrow F_s B} \\
\\
\frac{f : \Gamma \otimes \Delta \rightarrow A}{f \circ (\rho_{\Gamma} \otimes id_{\Delta}) : (\Gamma \otimes \mathbb{1}) \otimes \Delta \rightarrow A} \\
\\
\frac{f \circ (\rho_{\Gamma} \otimes id_{\Delta}) \circ (id_{\Gamma} \otimes w_{As}) \otimes id_{\Delta} : (\Gamma \otimes F_s A) \otimes \Delta \rightarrow A}{f : \Gamma \otimes (F_s A \otimes B \otimes F_s A) \otimes \Delta \rightarrow C} \\
\\
\frac{f : \Gamma \otimes (F_s A \otimes B \otimes F_s A) \otimes \Delta \rightarrow C}{f \circ (id_{\Gamma} \otimes c_{As}^l \otimes id_{\Delta}) : \Gamma \otimes (F_s A \otimes B) \otimes \Delta \rightarrow C} \\
\\
\frac{f : \Gamma \otimes (F_s A \otimes B \otimes F_s A) \otimes \Delta \rightarrow C}{(id_{\Gamma} \otimes c_{As}^r \otimes id_{\Delta}) \circ f : \Gamma \otimes (B \otimes F_s A) \otimes \Delta \rightarrow C} \\
\\
\frac{f : \Gamma \otimes (\Delta \otimes F_s A) \otimes \Theta \rightarrow B}{(id_{\Gamma} \otimes (id_{\Delta} \otimes e_{As}) \otimes id_{\Theta}) \circ f : \Gamma \otimes (F_s A \otimes \Delta) \otimes \Theta \rightarrow B} \\
\\
\frac{f : \Gamma \otimes (F_s A \otimes \Delta) \otimes \Theta \rightarrow B}{(id_{\Gamma} \otimes (id_{\Delta} \otimes e_{As}^{-1}) \otimes id_{\Theta}) \circ f : \Gamma \otimes (\Delta \otimes F_s A) \otimes \Theta \rightarrow B}
\end{array}$$

- Completeness:

**Definition 6.**

□

## 1 Concrete model