On R-models

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Definition 1.

R-model is a pair $\mathcal{M} = \langle W, R, v \rangle$, where R is a transitive relation on W and $v : Tp \to 2^R$ is a valuation, such that:

- 1. $\mathcal{M}, w \Vdash p_i \Leftrightarrow w \in v(p_i)$;
- 2. $\mathcal{M}, \langle a, b \rangle \Vdash A \bullet B \Leftrightarrow \text{there exists } c \in W, \mathcal{M}, a \Vdash A \text{ and } \mathcal{M}, b \Vdash B;$
- 3. $\mathcal{M}, \langle a, b \rangle \Vdash A \backslash B \Leftrightarrow \text{for all } c \in R^{-1}(a), \ \mathcal{M}, \langle c, a \rangle \Vdash A \text{ implies } \mathcal{M}, \langle c, b \rangle \Vdash B;$
- 4. $\mathcal{M}, \langle a, b \rangle \Vdash B/A \Leftrightarrow \text{for all } c \in R(a), \mathcal{M}, \langle a, c \rangle \Vdash A \text{ implies } \mathcal{M}, \langle b, c \rangle \Vdash B;$
- 5. $\mathcal{M}, \langle a, b \rangle \Vdash \Gamma \to A \Leftrightarrow \mathcal{M}, \langle a, b \rangle \Vdash \Gamma \text{ implies } \mathcal{M}, \langle a, b \rangle \Vdash A$

where $\mathcal{M}, \langle a, b \rangle \Vdash \Gamma$ denotes $\mathcal{M}, \langle a, b \rangle \Vdash A_1 \bullet \cdots \bullet A_n$, or, equivalently, there exist c_1, \ldots, c_{n-1} , such that $\mathcal{M}, \langle a, c_1 \rangle \Vdash A_1, \mathcal{M}, \langle c_1, c_2 \rangle \Vdash A_2, \ldots, \mathcal{M}, \langle c_{n-1}, b \rangle \Vdash A_n$ implies that $\mathcal{M}, \langle a, b \rangle \Vdash B$.

Theorem 1. Let \mathbb{F} is a R-frame, then $\mathbb{F} \models L$.

Definition 2.

Let \mathcal{F}_1 , \mathcal{F}_2 be R-frames and $\mathcal{M}_1 = \langle \mathcal{F}_1, v_1 \rangle$, $\mathcal{M}_2 = \langle \mathcal{F}_2, v_2 \rangle$ be R-models. A map $f: W_1 \to W_2$ is said to be a R-frame p-morphism if the following conditions hold:

- 1. f is onto;
- 2. $\forall a, b \in W_1(aR_1b \Rightarrow f(a)R_2f(b))$ (monotonicity);
- 3. $\forall d \in W_1, c \in W_2, f(d)R_2c \Rightarrow \exists c' \in W_1, f(c') = c \& dR_1c' \text{ (left lift property)};$
- 4. $\forall d \in W_1, c \in W_2, cR_2f(d) \Rightarrow \exists c' \in W_1, f(c') = c \& c'R_1d \text{ (right lift property)};$

A map $f: \mathcal{F}_1 \to \mathcal{F}_2$ is R-model p-moprhism, iff:

• $\mathcal{M}_1, \langle a, b \rangle \Vdash p_i \Leftrightarrow \mathcal{M}_2, \langle f(a), f(b) \rangle \Vdash p_i$

Lemma 1. Let $f: \mathcal{M}_1 \twoheadrightarrow \mathcal{M}_2$, then $\mathcal{M}_1, \langle a, b \rangle \Vdash A \Leftrightarrow \mathcal{M}_2, \langle f(a), f(b) \rangle \Vdash A$, for all $a, b \in W_1$ and for all $A \in Fm$

Proof.

 $1. \Rightarrow$

- (a) Basic case: follows from the definition.
- (b) Let $A = B \bullet C$ and $\mathcal{M}_1, \langle a, b \rangle \Vdash B \bullet C$, then there exists $c \in W_1$, such that $\mathcal{M}_1, \langle a, c \rangle \Vdash B$ and $\mathcal{M}_1, \langle c, b \rangle \Vdash C$.

Then, aR_1c and cR_1b , so $f(a)R_2f(c)$ and $f(c)R_2f(b)$.

Thus, by IH, $\mathcal{M}_2, \langle f(a), f(c) \rangle \Vdash B$ and $\mathcal{M}_2, \langle f(c), f(b) \rangle \Vdash C$, so $\mathcal{M}_2, \langle f(a), f(b) \rangle \Vdash B \bullet C$.

(c) Let $A = B \setminus C$ and $\mathcal{M}_1, \langle a, b \rangle \Vdash B \setminus C$. Let $c \in W_2$, such that $cR_2f(a)$. Then, by left lift property, there exist $d \in W_1$, such that f(d) = c and dR_1a .

Thus, $\mathcal{M}_1, \langle d, a \rangle \Vdash A$ implies $\mathcal{M}_1, \langle d, b \rangle \Vdash B$.

By IH, $\mathcal{M}_2, \langle c, f(a) \rangle \Vdash A$ implies $\mathcal{M}_2, \langle c, f(b) \rangle \Vdash B$, then $\mathcal{M}_2, \langle f(a), f(b) \rangle \Vdash A \backslash B$.

- (d) Similarly to (c), but via right lift property.
- (e) Let $\mathcal{M}_1, \langle a, b \rangle \Vdash A_1, \dots, A_n \to A$.

Thus, there exist $c_1, \ldots, c_{n-1} \in W_1$, such that $\mathcal{M}_1, \langle a, c_1 \rangle \Vdash A_1, \ldots, \mathcal{M}_1, \langle c_{n-2}, c_{n-1} \rangle \Vdash A_n$ implies that $\mathcal{M}_1, \langle c_{n-1}, b \rangle \Vdash B$.

By IH, $\mathcal{M}_2, \langle f(a), f(c_1) \rangle \Vdash A_1, \dots, \mathcal{M}_1, \langle f(c_{n-1}), f(b) \rangle \Vdash A_n$.

So, $\mathcal{M}_2, \langle f(a), f(b) \rangle \Vdash A_1 \bullet \cdots \bullet A_n$.

Thus, $\mathcal{M}_2, \langle f(a), f(b) \rangle \Vdash B$.

2. ←

- (a) Basic case: follows from the definition.
- (b) Let $A = B \bullet C$. Let $\mathcal{M}_2, \langle f(a), f(b) \rangle \Vdash B \bullet C$. Then there exists $c \in W_2$, such that $\mathcal{M}_2, \langle f(a), c \rangle \Vdash B$ and $\mathcal{M}_2, \langle c, f(b) \rangle \Vdash C$. So far as f is surjection, then there exists $d \in W_1$, such that c = f(d), then $\mathcal{M}_2, \langle f(a), f(d) \rangle \Vdash B$ and $\mathcal{M}_2, \langle f(d), f(b) \rangle \Vdash C$, and, by IH, $\mathcal{M}_1, \langle a, d \rangle \Vdash B$ and $\mathcal{M}_1, \langle d, b \rangle \Vdash C$, then $\mathcal{M}_1, \langle a, b \rangle \Vdash B \bullet C$.
- (c) Let $A = B \setminus C$ and $\mathcal{M}_2, \langle f(a), f(b) \rangle \Vdash B \setminus C$. Let $c \in W_1$ and cR_1a , then $f(c)R_1f(a)$ by monotonicity, so $\mathcal{M}_2, \langle f(c), f(a) \rangle \Vdash A$ implies $\mathcal{M}_2, \langle f(c), f(b) \rangle \Vdash B$. By IH, $\mathcal{M}_1, \langle c, a \rangle \Vdash A$ implies $\mathcal{M}_1, \langle c, b \rangle \Vdash B$. Thus, $\mathcal{M}_1, \langle c, a \rangle \Vdash A \setminus B$.
- (d) Similarly to (c).
- (e) Let $\mathcal{M}_2, \langle f(a), f(b) \rangle \Vdash \Gamma \to A$, so by case (b) and IH, $\mathcal{M}_2, \langle a, b \rangle \Gamma$, so $\mathcal{M}_\infty, \langle a, b \rangle A$.

Lemma 2.

- 1. Let \mathcal{M}_1 and \mathcal{M}_2 be R-models and $\mathcal{M}_1 \twoheadrightarrow \mathcal{M}_2$. Then $\mathcal{M}_1 \models A$ iff $\mathcal{M}_2 \models A$.
- 2. Let \mathcal{F}_1 and \mathcal{F}_2 be R-frames and $\mathcal{F}_1 \twoheadrightarrow \mathcal{F}_2$, then $\mathcal{F}_1 \models A$ implies $\mathcal{F}_2 \models A$.
- 3. $\mathcal{F}_1 \cong \mathcal{F}_2$, then $Log(\mathcal{F}_1) = Log(\mathcal{F}_2)$.

Proof. Standardly.