

Quantale model of Lambek calculus with subexponentials

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1 Calculus

Definition 1. A subexponential signature is an ordered quintuple:

$$\Sigma = \langle I, \leq, W, C, E \rangle,$$

where $I = \{s_1, \dots, s_n\}$, $\langle I, \leq \rangle$ is a preorder. W, C, E are subsets of I and $W \cup C \subseteq E$.

Definition 2. Noncommutative linear logic with subexponentials ($SMALC_\Sigma$), where Σ is a subexponential signature.

$$\begin{array}{c}
 \overline{A \Rightarrow A} \text{ }^{ax} \\
 \\
 \frac{\Gamma \Rightarrow A \quad \Delta, B, \Theta \Rightarrow C}{\Delta, \Gamma, A \backslash B, \Theta \Rightarrow C} \backslash \rightarrow \qquad \frac{A, \Pi \Rightarrow B}{\Pi \Rightarrow A \backslash B} \rightarrow \backslash \\
 \\
 \frac{\Gamma \Rightarrow A \quad \Delta, B, \Theta \Rightarrow C}{\Delta, B / A, \Gamma, \Theta \Rightarrow C} / \rightarrow \qquad \frac{\Pi, A \Rightarrow B}{\Pi \Rightarrow B / A} \rightarrow / \\
 \\
 \frac{\Gamma, A, B, \Delta \Rightarrow C}{\Gamma, A \bullet B, \Delta \Rightarrow C} \bullet \rightarrow \qquad \frac{\Gamma \Rightarrow A \quad \Delta \Rightarrow B}{\Gamma, \Delta \Rightarrow A \bullet B} \rightarrow \bullet \\
 \\
 \frac{\Gamma, A_i, \Delta \Rightarrow B}{\Gamma, A_1 \& A_2, \Delta \Rightarrow B} \&, i = 1, 2 \rightarrow \qquad \frac{\Gamma \Rightarrow A \quad \Gamma \Rightarrow B}{\Gamma \Rightarrow A \& B} \rightarrow \& \\
 \\
 \frac{\Gamma, A, \Delta \Rightarrow C \quad \Gamma, B, \Delta \Rightarrow C}{\Gamma, A \vee B, \Delta \Rightarrow C} \vee \rightarrow \qquad \frac{\Gamma \Rightarrow A_i}{\Gamma \Rightarrow A_1 \vee A_2} \rightarrow \vee, i = 1, 2 \\
 \\
 \frac{\Gamma, \Delta \Rightarrow A}{\Gamma, \mathbf{1}, \Delta \Rightarrow A} \mathbf{1} \rightarrow \qquad \overline{\Rightarrow \mathbf{1}} \rightarrow \mathbf{1} \\
 \\
 \frac{\Gamma, A, \Delta \Rightarrow C}{\Gamma, !^s A, \Delta \Rightarrow C} ! \rightarrow \qquad \frac{!^{s_1} A_1, \dots, !^{s_n} A_n \Rightarrow A}{!^{s_1} A_1, \dots, !^{s_n} A_n \Rightarrow !^s A} \rightarrow !, \forall j, s_j \geq s \\
 \\
 \frac{\Gamma, \Delta \Rightarrow B}{\Gamma, !^s A, \Delta \Rightarrow B} \text{weak}_!, s \in C
 \end{array}$$

$$\frac{\Gamma, !^s A, \Delta, !^s A, \Theta \Rightarrow B}{\Gamma, !^s A, \Delta, \Theta \Rightarrow B} \text{ ncontr}_1, s \in C$$

$$\frac{\Gamma, !^s A, \Delta, !^s A, \Theta \Rightarrow B}{\Gamma, \Delta, !^s A, \Theta \Rightarrow B} \text{ ncontr}_2, s \in C$$

$$\frac{\Gamma, \Delta, !^s A, \Theta \Rightarrow B}{\Gamma, !^s A, \Delta, \Theta \Rightarrow A} \text{ ex}_1, s \in E$$

$$\frac{\Gamma, !^s A, \Delta, \Theta \Rightarrow B}{\Gamma, \Delta, !^s A, \Theta \Rightarrow A} \text{ ex}_1, s \in E$$

Lemma 1. *Let $A \Leftrightarrow B$, then $C[p_i := A] \Leftrightarrow C[p_i := B]$*

Proof. By induction on C . □

Lemma 2. • $!_{s_i} \Gamma \rightarrow A$ iff $!_{s_i} \Gamma \rightarrow !_{s_i} A$.

• $!_{s_i} A \leftrightarrow !_{s_i} (!_{s_i} A)$

Proof.

1. $!_{s_i} \Gamma \rightarrow A$ iff $!_{s_i} \Gamma \rightarrow !_{s_i} A$;

$$\frac{!_{s_i} \Gamma \rightarrow A}{!_{s_i} \Gamma \rightarrow !_{s_i} A} \rightarrow_{s_i}$$

$$\frac{!_{s_i} \Gamma \rightarrow !_{s_i} A \quad \frac{A \rightarrow A}{!_{s_i} A \rightarrow A} !_{s_i} \rightarrow}{!_{s_i} \Gamma \rightarrow A} \text{ cut}$$

2. $!_{s_i} A \leftrightarrow !_{s_i} !_{s_i} A$

$$\frac{\frac{A \rightarrow A}{!_{s_i} A \rightarrow A}}{!_{s_i} !_{s_i} A \rightarrow !_{s_i} A} !_{s_i} \rightarrow$$

□

2 Semantics

Definition 3. *Quantale*

A quantale is a triple $\langle A, \vee, \cdot \rangle$, such that $\langle A, \vee \rangle$ is a complete lattice and $\langle A, \cdot \rangle$ is a semigroup. A quantale is called unital, if $\langle A, \cdot \rangle$ is a monoid.

It is easy to see, that any (unital) quantale is a residual (monoid) semigroup. We define divisions as follows:

1. $a \backslash b = \bigvee \{c \mid a \cdot c \leq b\}$

$$2. b/a = \bigvee \{c \mid c \cdot a \leq b\}$$

Definition 4.

Let $\langle A, \bigvee, \cdot \rangle$ be a quantale. The center of a quantale is the set $Z(Q) = \{a \in Q \mid \forall b \in Q, a \cdot b = b \cdot a\}$

Definition 5. An open modality (or quantic conucleus) on quantale Q is a map $I : Q \rightarrow Q$, such that

1. $I(x) \leq x$;
2. $I(x) = I(I(x))$;
3. $x \leq y \Rightarrow I(x) \leq I(y)$;
4. $I(x) \cdot I(y) = I(I(x) \cdot I(y))$.

Definition 6. We define a partial order on open modalities on Q as $I_1 \leq I_2 \Leftrightarrow \forall a \in Q, I_1(a) \leq I_2(a)$.

Lemma 3. Let Q be a quantale and \mathcal{I} be a set of all open modalities on Q . Then \mathcal{I} for a locally small category.

Proof. $\langle \mathcal{I}, \leq \rangle$ form a partial order, so $\langle \mathcal{I}, \leq \rangle$ is a locally small category. \square

Lemma 4.

Let $\langle A, \bigvee, \cdot \rangle$ be a quantale and $I : Q \rightarrow Q$ is an open modality on Q , then $I(x) \cdot I(y) \leq I(x \cdot y)$.

Proof.

$I(x) \cdot I(y) \leq x \cdot y$, then $I(I(x) \cdot I(y)) \leq I(x \cdot y)$, but $I(x) \cdot I(y) \leq I(I(x) \cdot I(y))$. Thus, $I(x) \cdot I(y) \leq I(x \cdot y)$. \square

Definition 7. An open modality is called central, if $\forall a, b \in Q, I(a) \cdot b = b \cdot I(a)$.

Definition 8. An open modality is called weak idempotent, if $\forall a, b \in Q, I(a) \cdot b \leq I(a) \cdot b \cdot I(a)$ and $b \cdot I(a) \leq I(a) \cdot b \cdot I(a)$.

Definition 9. An open modality is called unital, if $\forall a \in Q, I(a) \leq e$.

Lemma 5. Let I be an interior on some unital quantale $\langle Q, \bigvee, \cdot, e \rangle$. Then, if I is unital and weak idempotent, then I is central.

Proof.

$$\begin{aligned}
& b \cdot I(a) \leq \\
& \quad \text{Right weak idempotence} \\
& I(a) \cdot b \cdot I(a) \leq \\
& \quad \text{Unitality} \\
& I(a) \cdot b \cdot I(e) \leq \\
& \quad \text{Identity} \\
& I(a) \cdot b \leq \\
& \quad \text{Left weak idempotence} \\
& I(a) \cdot b \cdot I(a) \leq \\
& \quad \text{Unitality} \\
& e \cdot b \cdot I(a) \leq \\
& \quad \text{Identity} \\
& b \cdot I(a)
\end{aligned}$$

Hence, $b \cdot I(a) = I(a) \cdot b$

\square

Proposition 1.

Let Q be a quantale and $S \subseteq Q$ a subquantale, then $I : Q \rightarrow Q$, such that $I(a) = \bigvee \{s \in S \mid x \leq a\}$, is an open modality. Moreover, $\{x \in Q \mid I(x) = x\} = S$.

Proof. See □

Proposition 2.

Let Q be a quantale and $S_1, S_2 \subseteq Q$, such that $S_1 \subseteq S_2$.
Then $I_1(a) \leq I_2(a)$.

Proof.

Let $a \in Q$, so $\{s \in S_1 \mid s \leq a\} \subseteq \{s \in S_2 \mid s \leq a\}$, so $\bigvee \{s \in S_1 \mid s \leq a\} \subseteq \bigvee \{s \in S_2 \mid s \leq a\}$.
Thus, $I_1(a) \leq I_2(a)$. □

Proposition 3.

Let Q be a quantale and $S \subseteq Q$ a subquantale, then the following operations are open modalities:

1. $I_z(a) = \bigvee \{s \in S \mid s \leq a, s \in Z(Q)\};$
2. $I_{\mathbb{1}}(a) = \bigvee \{s \in S \mid s \leq a, s \leq \mathbb{1}\};$
3. $I_{idem}(a) = \bigvee \{s \in S \mid s \leq a, \forall b \in Q, b \cdot s \vee s \cdot b \leq s \cdot b \cdot s\};$
4. $I_{z, \mathbb{1}}, I_{z, idem}, I_{\mathbb{1}, idem}, I_{z, \mathbb{1}, idem}.$

Proof. Immediately. □

Proposition 4.

1. $\forall a \in Q, I_{\mathbb{1}, idem}(a) \leq I_z(a).$
2. $\forall a \in Q, I_{z, \mathbb{1}, idem} = I_{\mathbb{1}, idem}(a)$

Proof. Follows from Lemma 3. □

Proposition 5.

1. $I_z(a) \vee I_{\mathbb{1}}(a) \vee I_{idem}(a) \leq I(a)$
2. $I_{z, \mathbb{1}, idem} \leq I_{z, \mathbb{1}}(a) \wedge I_{z, idem}(a)$

Lemma 6. $\forall a \in Q, I_1(a) \leq I_2(I_1(a))$, if $I_1(a) \leq I_2(a)$.

Proof. $I_1(a) \leq I_1(I_1(a)) \leq I_2(I_1(a))$ □

Lemma 7. $I_1(a_1) \cdot I_2(a_2) \leq I'(I_1(a_1) \cdot I_2(a_2))$, where $I_i \leq I', i = 1, 2$.

Proof.

$$\begin{aligned} I_1(a_1) \cdot I_2(a_2) &\leq \\ I_1(I_1(a_1)) \cdot I_2(I_2(a_2)) &\leq \\ I'(I_1(a_1)) \cdot I'(I_2(a_2)) &\leq \\ I'(I_1(a_1) \cdot I_2(a_2)) & \end{aligned}$$
□

Definition 10. *Interpretation of subexponential signature*

Let $\Sigma = \langle I, \leq, W, C, E \rangle$ be a subexponential signature, where $|I| = n$ and \mathcal{Q} is a category of open modalities on Q . Subexponential interpretation is a contravariant functor $\sigma : I \rightarrow \mathcal{Q}$ defined as follows:

$$\sigma(s_i) = \begin{cases} \square_i : Q \rightarrow Q, \text{ s.t. } \forall a \in Q, \square_i(a) = \{s \in S_i \mid s \leq a\}, \\ \quad \text{if } s_i \notin W \cap C \cap E \\ \square_i : Q \rightarrow Q, \text{ s.t. } \forall a \in Q, \square_i(a) = \{s \in S_i \mid s \leq a, \leq 1\}, \\ \quad \text{if } s_i \in W \\ \square_i : Q \rightarrow Q, \text{ s.t. } \forall a \in Q, \square_i(a) = \{s \in S_i \mid s \leq a, \in Z(Q)\}, \\ \quad \text{if } s_i \in E \\ \square_i : Q \rightarrow Q, \text{ s.t. } \forall a \in Q, \square_i(a) = \{s \in S_i \mid s \leq a, \forall b, b \cdot s \vee s \cdot b \leq s \cdot b \cdot s\}, \\ \quad \text{if } s_i \in E \\ \text{otherwise, if } s_i \text{ belongs to some intersection of subsets, then we combine the relevant conditions} \end{cases}$$

Definition 11. Let Q be a quantale, $f : Tp \rightarrow Q$ a valuation and $\sigma : I \rightarrow \mathcal{S}$ a subexponential interpretation, then interpretation is defined inductively:

$$\begin{aligned} \llbracket p_i \rrbracket &= f(p_i) \\ \llbracket 1 \rrbracket &= e \\ \llbracket A \bullet B \rrbracket &= \llbracket A \rrbracket \cdot \llbracket B \rrbracket \\ \llbracket A \setminus B \rrbracket &= \llbracket A \rrbracket \setminus \llbracket B \rrbracket \\ \llbracket A/B \rrbracket &= \llbracket A \rrbracket / \llbracket B \rrbracket \\ \llbracket A \& B \rrbracket &= \llbracket A \rrbracket \wedge \llbracket B \rrbracket \\ \llbracket A \vee B \rrbracket &= \llbracket A \rrbracket \vee \llbracket B \rrbracket \\ \llbracket !_{s_i} A \rrbracket &= \sigma(s_i) \llbracket A \rrbracket \end{aligned}$$

Definition 12. $\Gamma \models A \Leftrightarrow \forall f \forall \sigma, \llbracket \Gamma \rrbracket \leq \llbracket A \rrbracket$

Theorem 1. $\Gamma \rightarrow A \Rightarrow \llbracket \Gamma \rrbracket \leq \llbracket A \rrbracket$

Proof. We consider cases with modal rules.

1. Let $!_{s_1} A_1, \dots, !_{s_n} A_n \rightarrow A$ and $\forall i, s \leq s_i$.
Then $\forall a \in Q, \sigma(s_i)(a) \leq \sigma(s)(a)$.
By IH, $\sigma(s_1) \llbracket A_1 \rrbracket \cdot \dots \cdot \sigma(s_n) \llbracket A_n \rrbracket \leq \llbracket A \rrbracket$.
Thus, $\sigma(s)(\sigma(s_1) \llbracket A_1 \rrbracket \cdot \dots \cdot \sigma(s_n) \llbracket A_n \rrbracket) \leq \sigma(s)(\llbracket A \rrbracket)$.
By Lemma 5, $\sigma(s_1) \llbracket A_1 \rrbracket \cdot \dots \cdot \sigma(s_n) \llbracket A_n \rrbracket \leq \sigma(s)(\sigma(s_1) \llbracket A_1 \rrbracket \cdot \dots \cdot \sigma(s_n) \llbracket A_n \rrbracket)$.
So, $\sigma(s_1) \llbracket A_1 \rrbracket \cdot \dots \cdot \sigma(s_n) \llbracket A_n \rrbracket \leq \sigma(s)(\llbracket A \rrbracket)$.
2. Let $\Gamma, A, \Delta \rightarrow B$.
By IH, $\llbracket \Gamma \rrbracket \cdot \llbracket A \rrbracket \cdot \llbracket \Delta \rrbracket \leq \llbracket B \rrbracket$.
By the definition, $\sigma(s_i)(\llbracket A \rrbracket) \leq \llbracket A \rrbracket$.
So, $\llbracket \Gamma \rrbracket \cdot \sigma(s_i)(\llbracket A \rrbracket) \cdot \llbracket \Delta \rrbracket \leq \llbracket B \rrbracket$.
3. Let $\Gamma, \Delta \rightarrow B$, $A \in Fm$, and $s_i \in W$.
So, $\llbracket \Gamma \rrbracket \cdot \llbracket \Delta \rrbracket \leq \llbracket B \rrbracket$, then $\llbracket \Gamma \rrbracket \cdot e \cdot \llbracket \Delta \rrbracket \leq \llbracket B \rrbracket$, where $e \in Q$ is unit.
By the definition of unital open modality, $\sigma(s_i)(\llbracket A \rrbracket) \leq e$.
Thus, $\llbracket \Gamma \rrbracket \cdot \sigma(s_i)(\llbracket A \rrbracket) \cdot \llbracket \Delta \rrbracket \leq \llbracket B \rrbracket$.

4. Let $\Gamma, !_{s_i} A, \Delta, !_{s_i} A, \Pi \rightarrow B$ and $s_i \in C$.

By IH, $\llbracket \Gamma \rrbracket \cdot \sigma(s_i)(\llbracket A \rrbracket) \cdot \llbracket \Delta \rrbracket \cdot \sigma(s_i)(\llbracket A \rrbracket) \cdot \llbracket \Pi \rrbracket \leq \llbracket B \rrbracket$.

By the definition, $\sigma(s_i)(\llbracket A \rrbracket) \cdot \llbracket \Delta \rrbracket \leq \sigma(s_i)(\llbracket A \rrbracket) \cdot \llbracket \Delta \rrbracket \cdot \sigma(s_i)(\llbracket A \rrbracket)$.

Then $\llbracket \Gamma \rrbracket \cdot \sigma(s_i)(\llbracket A \rrbracket) \cdot \llbracket \Delta \rrbracket \cdot \llbracket \Pi \rrbracket \leq \llbracket B \rrbracket$

5. Let $\Gamma, !_{s_i} A, \Delta, \Pi \rightarrow B$ and $s_i \in E$, so $\sigma(s_i)(a) \in Z(Q)$ for all $a \in Q$ by the definition.

By IH, $\llbracket \Gamma \rrbracket \cdot \sigma(s_i)(\llbracket A \rrbracket) \cdot \llbracket \Delta \rrbracket \cdot \llbracket \Pi \rrbracket \leq \llbracket B \rrbracket$

Hence, $\llbracket \Gamma \rrbracket \cdot \llbracket \Delta \rrbracket \cdot \sigma(s_i)(\llbracket A \rrbracket) \cdot \llbracket \Pi \rrbracket \leq \llbracket B \rrbracket$.

□

3 Quantale completeness

Definition 13.

Let $\mathcal{F} \subseteq Fm$, an ideal is a subset $\mathcal{I} \subseteq \mathcal{F}$, such that:

- If $B \in \mathcal{I}$ and $A \rightarrow B$, then $A \in \mathcal{I}$;
- If $A, B \in \mathcal{I}$, then $A \vee B \in \mathcal{I}$.

Definition 14.

Let $S \subseteq \mathcal{F} \subseteq Fm$, then $\bigvee S = \bigcap \{ \mathcal{I} \subseteq \mathcal{F} \mid S \subseteq \mathcal{I} \}$

Proposition 6. $\bigvee S$ is an ideal.

Lemma 8. $A \subseteq Fm$, then $\{B \mid B \rightarrow A'\} = \bigvee A$.

Proof.

Let $A \subseteq Fm$. Then $\{B \mid B \rightarrow A', A' \in A\} \subseteq \bigvee A$, so far as $A' \rightarrow A'$ by axiom.

On the other hand, $\{B \mid B \rightarrow A', A' \in A\}$ is an ideal, hence, $A \subseteq \{B \mid B \rightarrow A', A' \in A\}$. □

Lemma 9. $\bigvee A \subseteq \bigvee B$ iff $\forall A' \in A, \forall B' \in B, A' \rightarrow B'$.

Proof. Let $\bigvee A \subseteq \bigvee B$, then $\{C \mid C \rightarrow A', A' \in A\} \subseteq \{D \mid D \rightarrow B', B' \in B\}$.

Thus, for all $A' \in A$, $A' \in \{C \mid C \rightarrow A', A' \in A\}$, then $A' \in \{D \mid D \rightarrow B', B' \in B\}$, hence $A' \rightarrow B'$, for all $B' \in B$.

On the other hand, let $A' \rightarrow B'$ for all $A' \in A$, $B' \in B$ and $C \in \bigvee A$.

Thus, $C \rightarrow A'$, then $C \rightarrow B'$ by cut, so $C \in \bigvee B$.

□

Lemma 10. Let $\mathcal{Q} = \{\bigvee S \mid S \subseteq Fm\}$ and $\bigvee \mathcal{A} \cdot \bigvee \mathcal{B} = \{A \bullet B \mid A \in \mathcal{A}, B \in \mathcal{B}\}$. Then $\langle \mathcal{Q}, \subseteq, \cdot, \bigvee 1 \rangle$ is a quantale.

Proof. See □

Lemma 11. Let $!_s \in I$, $I \notin W \cap E \cap C$, then $\Box_s(\bigvee A) = \bigvee \{B \mid B \rightarrow \bigvee !_s A', A' \in A\}$ is a quantic conucleus.

Proof.

1. $\Box_s(\bigvee A) \subseteq \bigvee A$;

Let $B \in \Box_s(\bigvee A)$, then for all $A' \in A$, $B \rightarrow !_s A'$, but $!_s A' \rightarrow A'$, then $B \rightarrow A'$, so $B \in \bigvee A$.

2. $\Box_s(\Box_s(\bigvee A)) = \bigvee \Box_s(\bigvee A)$;
 $\Box_s(\Box_s(\bigvee A)) =$
 $\{B \mid B \rightarrow \bigvee !_s !_s A', A' \in A\} =$, that follows from equivalence $!_s !_s B \leftrightarrow !_s B$.
 $\{B \mid B \rightarrow \bigvee !_s A', A' \in A\}$
3. $\bigvee A \subseteq \bigvee B \Rightarrow \Box_s(\bigvee A) \subseteq \Box_s(\bigvee B)$;
Follows from admissibility of K-rule for all $s \in I$.
4. $\Box_s \bigvee A \cdot \Box_s \bigvee B = \Box_s(\Box_s \bigvee A \cdot \Box_s \bigvee B)$.
 $\Box_s \bigvee A \cdot \Box_s \bigvee B =$
 $\bigvee \{C \bullet D \mid C \bullet D \rightarrow !_s A' \bullet !_s B'\} =$
 $\bigvee \{C \bullet D \mid C \bullet D \rightarrow !_s (!_s A' \bullet !_s B')\} =$
 $\Box_s(\Box_s \bigvee A \cdot \Box_s \bigvee B)$

□

Lemma 12.

1. Let $s \in W$, then for all $A \subseteq Fm$, $\mathbf{1} \in \Box_s(\bigvee A)$;
2. Let $s \in E$, then $\Box_s(\bigvee A) \cdot \bigvee B = \bigvee B \cdot \Box_s(\bigvee A)$.
3. Let $s \in C$, then $(\Box_s \bigvee A \cdot \bigvee B) \cup (\bigvee B \cdot \Box_s \bigvee A) \subseteq \Box_s \bigvee A \cdot \bigvee B \cdot \Box_s \bigvee A$, for all $B \subseteq Fm$.

Proof. 1. Let $s \in W$, then for all $A \subseteq Fm$, $\Box_s(\bigvee A) = \{!_s B \mid !_s B \rightarrow A', A' \in A\}$. But, $!_s B \rightarrow \mathbf{1}$, hence, $\mathbf{1} \in \Box_s(\bigvee A)$, so far as $\Box_s(\bigvee A)$ is an ideal.

2.
 $\Box_s(\bigvee A) \cdot \bigvee B =$
 $\bigvee \{!_s C \bullet D \mid !_s C \bullet D \rightarrow A' \bullet B', A' \in A, B' \in B\} =$
 $\bigvee \{D \bullet !_s C \mid D \bullet !_s C \rightarrow A' \bullet B', A' \in A, B' \in B\} =$
 $\bigvee B \cdot \Box_s(\bigvee A)$

3.
 $\Box_s \bigvee A \cdot \bigvee B = \bigvee \{!_s C \bullet D \mid !_s C \bullet D \rightarrow A' \bullet B'\}$. $!_s C \bullet D \rightarrow !_s C \bullet D \bullet !_s C$, hence $\Box_s \bigvee A \cdot \bigvee B \subseteq$
 $\Box_s \bigvee A \cdot \bigvee B \cdot \Box_s \bigvee A$.
Similarly with $\bigvee B \cdot \Box_s \bigvee A$.

□

Lemma 13.

Let $i, j \in I$ and $i \leq j$, then for all $A \subseteq Fm$, $\Box_j(\bigvee A) \subseteq \Box_i(\bigvee A)$.

Proof. Let $i, j \in I$ and $i \leq j$. Let $B \in \Box_j(\bigvee A)$, then $\forall A', B \rightarrow !_j A'$.

But $!_j A \rightarrow !_i A$. Then $B \rightarrow !_i A$ by hence. So, $B \in \Box_i(\bigvee A)$.

□

Definition 15. Let Q be a syntactic quantale as proposed above and $\mathcal{I} = \langle I, \leq, W, C, E \rangle$ be a subexponential signature.

We define a map $\Box : \mathcal{I} \rightarrow \text{Mod}_Q$ as follows:

$$\Box(i)(\bigvee A) = \{B \mid B \rightarrow !_i A\}.$$

Lemma 14. \square is a subexponential interpretation.

Proof. Follows from lemmas 11 and 12. \square

Lemma 15.

Let Q be a quantale constructed above and $\square_1, \dots, \square_n$ be a family of quantic conuclei on Q . Then there exist a model $\langle Q, \llbracket \cdot \rrbracket \rangle$, such that $\llbracket A \rrbracket = \bigvee \{A\}$, $A \in Fm$.

Proof.

We define an interpretation as follows:

1. $\llbracket p_i \rrbracket = \bigvee \{p_i\}$
2. $\llbracket 1 \rrbracket = \bigvee \{1\}$
3. $\llbracket A \bullet B \rrbracket = \bigvee \{A \bullet B\}$
4. $\llbracket A/B \rrbracket = \bigvee \{A/B\}$
5. $\llbracket B \setminus A \rrbracket = \bigvee \{B \setminus A\}$
6. $\llbracket A \& B \rrbracket = \bigvee \{A \& B\}$
7. $\llbracket A \vee B \rrbracket = \bigvee \{A \vee B\}$
8. $\llbracket !_s A \rrbracket = \square(s)(\bigvee A) = \{B \mid B \rightarrow !_s A\} = \bigvee \{ !_s A \}$

\square

Theorem 2. $\Gamma \models A \Rightarrow \Gamma \rightarrow A$.

Proof. Follows from lemmas 9, 12, 13, 14. \square