

Definition 1. *Monoidal comonad*

A monoidal comonad on some monoidal category \mathcal{C} is a triple $\langle \mathcal{F}, \epsilon, \delta \rangle$, where \mathcal{F} is a monoidal endofunctor and $\epsilon : \mathcal{F} \Rightarrow \text{Id}_{\mathcal{C}}$ (counit) and $\delta : \mathcal{F} \Rightarrow \mathcal{F}^2$ (comultiplication), such that the following diagrams commute:

$$\begin{array}{ccc}
 \mathcal{F}A \otimes \mathcal{F}B & \xrightarrow{\phi_{A,B}} & \mathcal{F}(A \otimes B) \\
 \downarrow \delta_A \otimes \delta_B & & \searrow \delta_{A \otimes B} \\
 \mathcal{F}\mathcal{F}A \otimes \mathcal{F}\mathcal{F}B & \xrightarrow{\phi_{\mathcal{F}A, \mathcal{F}B}} & \mathcal{F}(\mathcal{F}A \otimes \mathcal{F}B) \\
 & \nearrow \mathcal{F}(\phi_{A,B}) & \\
 & \mathcal{F}\mathcal{F}(A \otimes B) &
 \end{array}
 \qquad
 \begin{array}{ccc}
 \mathcal{F}A \otimes \mathcal{F}B & \xrightarrow{\phi_{A,B}} & \mathcal{F}(A \otimes B) \\
 \searrow \epsilon_A \otimes \epsilon_B & & \swarrow \epsilon_{A \otimes B} \\
 & A \otimes B &
 \end{array}$$

$$\begin{array}{ccc}
 \mathbb{1} & \xrightarrow{\phi} & \mathcal{F}\mathbb{1} \\
 \phi \downarrow & & \downarrow \delta_{\mathbb{1}} \\
 \mathcal{F}\mathbb{1} & \xrightarrow{\mathcal{F}(\phi)} & \mathcal{F}\mathcal{F}\mathbb{1} \\
 & & \downarrow \delta_{\mathcal{F}\mathbb{1}} \\
 & & \mathcal{F}\mathcal{F}\mathcal{F}\mathbb{1}
 \end{array}$$

$$\begin{array}{ccc}
 \mathbb{1} & \xrightarrow{id_{\mathbb{1}}} & \mathbb{1} \\
 \phi \searrow & & \swarrow \epsilon_{\mathbb{1}} \\
 & \mathcal{F}\mathbb{1} &
 \end{array}$$

Definition 2. *Biclosed monoidal category*

Let \mathcal{C} be a monoidal category. Biclosed monoidal category is a monoidal category with the following additional data:

1. Bifunctors $_ \multimap _ , _ \multimap _ : \mathcal{C}^{op} \times \mathcal{C} \rightarrow \mathcal{C}$;
2. Natural isomorphism $\mathbf{curry}_{A,B,C} : \text{Hom}(A \otimes B, C) \cong (B, A \multimap C)$;
3. Natural isomorphism $\mathbf{curry}'_{A,B,C} : \text{Hom}(A \otimes B, C) \cong (A, C \multimap B)$;
4. For each $A, B \in \text{Ob}_{\mathcal{C}}$, there are exist arrows $ev_{A,B} : A \otimes (A \multimap B) \rightarrow B$ and $ev'_{A,B} : (B \multimap A) \otimes A \rightarrow B$, such that for all $f : A \otimes C \rightarrow B$:
 - (a) $\Lambda_l \circ (id_A \otimes \mathbf{curry}(f)) = f$;
 - (b) $\Lambda_r \circ (\mathbf{curry}'(f) \otimes id_A) = f$

Definition 3. Let F be endofunctor and $A \in \text{Ob}_{\mathcal{C}}$, then a coalgebra of F is a tuple $\langle A, \theta \rangle$, where $\theta : A \rightarrow FA$.

Given coalgebras $\langle A, \theta \rangle$ and $\langle B, \psi \rangle$, a homomorphism is a morphism $f : A \rightarrow B$, s.t. the diagram below commutes:

$$\begin{array}{ccc}
 A & \xrightarrow{\theta} & FA \\
 f \downarrow & & \downarrow Ff \\
 B & \xrightarrow{\psi} & FB
 \end{array}$$

that is, $Ff \circ \theta = \psi \circ f$

Definition 4. *Subexponential model structure*

Let $\Sigma = \langle I, \leq, W, C, E \rangle$ be a subexponential signature and \mathcal{C} be a biclosed monoidal category, then a subexponential model structure is $\langle \mathcal{C}, \{\mathcal{F}_s\}_{s \in I} \rangle$ with the following additional data:

- for all $s \in I$, \mathcal{F}_s is a monoidal comonad;
- if $s \in W$, then for all $A \in \text{Ob}(\mathcal{C})$, there exists a morphism $w_{As} : F_s A \rightarrow \mathbb{1}$;
- if $s \in C$, then for all $A \in \text{Ob}(\mathcal{C})$, there exists morphisms $w_{Al} : F_s A \otimes A \otimes F_s A \rightarrow F_s A \otimes B$ and $w_{Ar} : F_s A \otimes A \otimes F_s A \rightarrow B \otimes F_s A$;
- if $s \in E$, then for all $A \in \text{Ob}(\mathcal{C})$, there is an isomorphism, $e_A : F_s A \otimes B \cong B \otimes F_s A$;
- if $s_1 \in W$, $s_2 \in I$ and $s_1 \leq s_2$, then there is a morphism $w_{As_2} : F_{s_2} A \rightarrow \mathbb{1}$ for all $A \in \text{Ob}(\mathcal{C})$ and ditto for E and C ;
- Let $\bigotimes_{s \in J, i=0}^n F_s A$, where $J \subset I$, and $s' \in I$, s.t. $s \geq s'$ for all $s \in J'$; Then there exists morphism a morphism $\theta_{\bigotimes_{s \in J, i=1}^n F_{s_j} A_i} : \bigotimes_{s \in J, i=0}^n F_s A \rightarrow F_{s'}(\bigotimes_{s \in J, i=0}^n F_s A)$, such that $\langle \bigotimes_{s \in J, i=1}^n F_{s_j} A_i, \theta_{\bigotimes_{s \in J, i=1}^n F_{s_j} A_i} \rangle$ is a coalgebra on F_s .

Definition 5. Let $\langle \mathcal{C}, \{\mathcal{F}_s\}_{s \in I} \rangle$ be a subexponential model structure for subexponential signature $\Sigma = \langle I, \leq, W, C, E \rangle$. Let $v : Tp \rightarrow \text{Ob}(\mathcal{C})$ be a valuation map. Then the interpretation function $\llbracket \cdot \rrbracket$ is defined as follows:

- (1) $\llbracket \mathbb{1} \rrbracket = \mathbb{1}$
- (2) $\llbracket A \setminus B \rrbracket = \llbracket A \rrbracket \multimap \llbracket B \rrbracket$
- (3) $\llbracket A / B \rrbracket = \llbracket A \rrbracket \multimap \llbracket B \rrbracket$
- (4) $\llbracket A \bullet B \rrbracket = \llbracket A \rrbracket \otimes \llbracket B \rrbracket$
- (5) $\llbracket !_s A \rrbracket = F_s \llbracket A \rrbracket$

Theorem 1. *The following statements are equivalent:*

- $SMLC_\Sigma + (\text{cut}) \vdash \Gamma \Rightarrow A$
- $SMLC_\Sigma \vdash \Gamma \Rightarrow A$
- $\exists f, f : \llbracket \Gamma \rrbracket \rightarrow \llbracket A \rrbracket$

Proof.

(1) \Rightarrow (2): cut elimination.

- (2) \Rightarrow (3): Soundness:

$$\begin{array}{c}
 \overline{id_A : A \rightarrow A} \\
 \\
 \frac{f : \Gamma \rightarrow A \quad g : \Delta \otimes B \otimes \Theta \rightarrow C}{g \circ (id_\Delta \otimes (ev_{A, B_l} \circ (f \otimes id_{A \multimap B})) \otimes id_\Theta) : \Delta \otimes (\Gamma \otimes A \multimap B) \otimes \Theta \rightarrow C} \\
 \\
 \frac{f : A \otimes \Pi \rightarrow B}{\Lambda_l(f) : \Pi \rightarrow A \multimap B} \\
 \\
 \frac{f : \Gamma \rightarrow A \quad g : \Delta \otimes B \otimes \Theta \rightarrow C}{g \circ (id_\Delta \otimes (ev_{A, B_l} \circ (id_{B \multimap A} \otimes f)) \otimes id_\Theta) : \Delta \otimes (B \multimap A \otimes \Gamma) \otimes \Theta \rightarrow C}
 \end{array}$$

$$\begin{array}{c}
\frac{f : \Pi \otimes A \rightarrow B}{\Lambda_r(f) : \Pi \rightarrow B \multimap A} \\
\\
\frac{f : \Gamma \otimes A \otimes B \otimes \Delta \rightarrow C}{f \circ (\alpha_{\Gamma, A, B} \otimes id_{\Delta}) : \Gamma \otimes (A \otimes B) \otimes \Delta \rightarrow C} \\
\\
\frac{f : \Gamma \rightarrow A \quad g : \Delta \rightarrow B}{f \otimes g : \Gamma \otimes \Delta \rightarrow A \otimes B} \\
\\
\frac{f : \Gamma \otimes A_i \otimes \Delta \rightarrow B}{f \circ (id_{\Gamma} \otimes \pi_i id_{\Delta}) : \Gamma \otimes (A_1 \times A_2) \otimes \Delta \rightarrow B} \\
\\
\frac{f : \Gamma \rightarrow A \quad g : \Gamma \rightarrow B}{\langle f, g \rangle : \Gamma \rightarrow A \times B} \\
\\
\frac{f : \Gamma \otimes A \otimes \Delta \rightarrow C \quad g : \Gamma \otimes B \otimes \Delta \rightarrow C}{id_{\Gamma} \otimes [f, g] \otimes id_{\Delta} : \Gamma \otimes (A + B) \otimes \Delta \rightarrow C} \\
\\
\frac{}{id_{\mathbb{1}} : \mathbb{1} \rightarrow \mathbb{1}} \\
\\
\frac{f : \Gamma \otimes \Delta \rightarrow A}{f \circ (\rho_{\Gamma} \otimes id_{\Delta}) : (\Gamma \otimes \mathbb{1}) \otimes \Delta \rightarrow A} \\
\\
\frac{f : \Gamma \otimes A \otimes \Delta \rightarrow B}{f \circ (id_{\Gamma} \otimes \delta_s^A \otimes id_{\Delta}) : \Gamma \otimes F_s A \otimes \Delta \rightarrow B} \\
\\
\frac{f : F_{s_1} A_1 \otimes \cdots \otimes F_{s_n} A_n \rightarrow B}{F_s(f) : F_s(F_{s_1} A_1 \otimes \cdots \otimes F_{s_n} A_n) \rightarrow F_s B} \\
\\
\frac{F_s(f) \circ \theta_{\otimes_{s \in J, i=1}^n F_{s_j} A_i} : F_{s_1} A_1 \otimes \cdots \otimes F_{s_n} A_n \rightarrow F_s B}{F_s(f) \circ \theta_{\otimes_{s \in J, i=1}^n F_{s_j} A_i} : F_{s_1} A_1 \otimes \cdots \otimes F_{s_n} A_n \rightarrow F_s B} \\
\\
\frac{f : \Gamma \otimes \Delta \rightarrow A}{f \circ (\rho_{\Gamma} \otimes id_{\Delta}) : (\Gamma \otimes \mathbb{1}) \otimes \Delta \rightarrow A} \\
\\
\frac{f \circ (\rho_{\Gamma} \otimes id_{\Delta}) \circ (id_{\Gamma} \otimes w_{As}) \otimes id_{\Delta} : (\Gamma \otimes F_s A) \otimes \Delta \rightarrow A}{f \circ (\rho_{\Gamma} \otimes id_{\Delta}) \circ (id_{\Gamma} \otimes w_{As}) \otimes id_{\Delta} : (\Gamma \otimes F_s A) \otimes \Delta \rightarrow A} \\
\\
\frac{f : \Gamma \otimes (F_s A \otimes B \otimes F_s A) \otimes \Delta \rightarrow C}{f \circ (id_{\Gamma} \otimes c_{As}^l \otimes id_{\Delta}) : \Gamma \otimes (F_s A \otimes B) \otimes \Delta \rightarrow C} \\
\\
\frac{f : \Gamma \otimes (F_s A \otimes B \otimes F_s A) \otimes \Delta \rightarrow C}{(id_{\Gamma} \otimes c_{As}^r \otimes id_{\Delta}) \circ f : \Gamma \otimes (B \otimes F_s A) \otimes \Delta \rightarrow C} \\
\\
\frac{f : \Gamma \otimes (\Delta \otimes F_s A) \otimes \Theta \rightarrow B}{(id_{\Gamma} \otimes (id_{\Delta} \otimes e_{As}) \otimes id_{\Theta}) \circ f : \Gamma \otimes (F_s A \otimes \Delta) \otimes \Theta \rightarrow B} \\
\\
\frac{f : \Gamma \otimes (F_s A \otimes \Delta) \otimes \Theta \rightarrow B}{(id_{\Gamma} \otimes (id_{\Delta} \otimes e_{As}^{-1}) \otimes id_{\Theta}) \circ f : \Gamma \otimes (\Delta \otimes F_s A) \otimes \Theta \rightarrow B}
\end{array}$$

- Completeness:

Definition 6.

□

1 Concrete model