

# Categorical model of noncommutative linear logic with subexponentials

**Definition 1.** A subexponential signature is an ordered quintuple:

$$\Sigma = \langle I, \leq, W, C, E \rangle,$$

where  $I = \{s_1, \dots, s_n\}$ ,  $\langle I, \leq \rangle$  is a preorder.  $W, C, E$  are subsets of  $I$  and  $W \cup C \subseteq E$ .

**Definition 2.** Noncommutative linear logic with subexponentials ( $SMALC_\Sigma$ ), where  $\Sigma$  is a subexponential signature.

$$\begin{array}{c}
\overline{A \Rightarrow A} \text{ }^{ax} \\
\\
\frac{\Gamma \Rightarrow A \quad \Delta, B, \Theta \Rightarrow C}{\Delta, \Gamma, A \backslash B, \Theta \Rightarrow C} \backslash \rightarrow \qquad \frac{A, \Pi \Rightarrow B}{\Pi \Rightarrow A \backslash B} \rightarrow \backslash \\
\\
\frac{\Gamma \Rightarrow A \quad \Delta, B, \Theta \Rightarrow C}{\Delta, B / A, \Gamma, \Theta \Rightarrow C} / \rightarrow \qquad \frac{\Pi, A \Rightarrow B}{\Pi \Rightarrow B / A} \rightarrow / \\
\\
\frac{\Gamma, A, B, \Delta \Rightarrow C}{\Gamma, A \bullet B, \Delta \Rightarrow C} \bullet \rightarrow \qquad \frac{\Gamma \Rightarrow A \quad \Delta \Rightarrow B}{\Gamma, \Delta \Rightarrow A \bullet B} \rightarrow \bullet \\
\\
\frac{\Gamma, A_i, \Delta \Rightarrow B}{\Gamma, A_1 \& A_2, \Delta \Rightarrow B} \&, i = 1, 2 \rightarrow \qquad \frac{\Gamma \Rightarrow A \quad \Gamma \Rightarrow B}{\Gamma \Rightarrow A \& B} \rightarrow \& \\
\\
\frac{\Gamma, A, \Delta \Rightarrow C \quad \Gamma, B, \Delta \Rightarrow C}{\Gamma, A \vee B, \Delta \Rightarrow C} \vee \rightarrow \qquad \frac{\Gamma \Rightarrow A_i}{\Gamma \Rightarrow A_1 \vee A_2} \rightarrow \vee, i = 1, 2 \\
\\
\frac{\Gamma, \Delta \Rightarrow A}{\Gamma, \mathbf{1}, \Delta \Rightarrow A} \mathbf{1} \rightarrow \qquad \frac{}{\Rightarrow \mathbf{1}} \rightarrow \mathbf{1} \\
\\
\frac{\Gamma, A, \Delta \Rightarrow C}{\Gamma, !^s A, \Delta \Rightarrow C} ! \rightarrow \qquad \frac{!^{s_1} A_1, \dots, !^{s_n} A_n \Rightarrow A}{!^{s_1} A_1, \dots, !^{s_n} A_n \Rightarrow !^s A} \rightarrow !, \forall j, s_j \geq s \\
\\
\frac{\Gamma, \Delta \Rightarrow B}{\Gamma, !^s A, \Delta \Rightarrow B} \text{weak}_!, s \in C \\
\\
\frac{\Gamma, !^s A, \Delta, !^s A, \Theta \Rightarrow B}{\Gamma, !^s A, \Delta, \Theta \Rightarrow B} \text{ncontr}_1, s \in C \\
\\
\frac{\Gamma, !^s A, \Delta, !^s A, \Theta \Rightarrow B}{\Gamma, \Delta, !^s A, \Theta \Rightarrow B} \text{ncontr}_2, s \in C
\end{array}$$

$$\frac{\Gamma, \Delta, !^s A, \Theta \Rightarrow B}{\Gamma, !^s A, \Delta, \Theta \Rightarrow A} \mathbf{ex}_1, s \in E$$

$$\frac{\Gamma, !^s A, \Delta, \Theta \Rightarrow B}{\Gamma, \Delta, !^s A, \Theta \Rightarrow A} \mathbf{ex}_1, s \in E$$

**Definition 3.** *Monoidal comonad*

A monoidal comonad on some monoidal category  $\mathcal{C}$  is a triple  $\langle \mathcal{F}, \epsilon, \delta \rangle$ , where  $\mathcal{F}$  is a monoidal endofunctor and  $\epsilon : \mathcal{F} \Rightarrow Id_{\mathcal{C}}$  (counit) and  $\epsilon : \mathcal{F} \Rightarrow \mathcal{F}^2$  (comultiplication), such that the following diagrams commute:

$$\begin{array}{ccc} \mathcal{F}A \otimes \mathcal{F}B & \xrightarrow{\phi_{A,B}} & \mathcal{F}(A \otimes B) \\ \downarrow \delta_A \otimes \delta_B & & \searrow \delta_{A \otimes B} \\ \mathcal{F}\mathcal{F}A \otimes \mathcal{F}\mathcal{F}B & \xrightarrow{\phi_{\mathcal{F}A, \mathcal{F}B}} & \mathcal{F}\mathcal{F}(A \otimes B) \\ & \nearrow \mathcal{F}(\phi_{A,B}) & \end{array} \quad \begin{array}{ccc} \mathcal{F}A \otimes \mathcal{F}B & \xrightarrow{\phi_{A,B}} & \mathcal{F}(A \otimes B) \\ \searrow \epsilon_A \otimes \epsilon_B & & \swarrow \epsilon_{A \otimes B} \\ & A \otimes B & \end{array}$$
  

$$\begin{array}{ccc} \mathbb{1} & \xrightarrow{\phi} & \mathcal{F}\mathbb{1} \\ \phi \downarrow & & \downarrow \delta_{\mathbb{1}} \\ \mathcal{F}\mathbb{1} & \xrightarrow{\mathcal{F}(\phi)} & \mathcal{F}\mathcal{F}\mathbb{1} \end{array}$$
  

$$\begin{array}{ccc} \mathbb{1} & \xrightarrow{id_{\mathbb{1}}} & \mathbb{1} \\ \phi \searrow & & \swarrow \epsilon_{\mathbb{1}} \\ & \mathcal{F}\mathbb{1} & \end{array}$$

**Definition 4.** *Biclosed monoidal category*

Let  $\mathcal{C}$  be a monoidal category. Biclosed monoidal category is a monoidal category with the following additional data:

1. Bifunctors  $\_ \multimap \_, \_ \multimap \_ : \mathcal{C}^{op} \times \mathcal{C} \rightarrow \mathcal{C}$ ;
2. Natural isomorphism  $\mathbf{curry}_{A,B,C} : Hom(A \otimes B, C) \cong (B, A \multimap C)$ ;
3. Natural isomorphism  $\mathbf{curry}'_{A,B,C} : Hom(A \otimes B, C) \cong (A, C \multimap B)$ ;
4. For each  $A, B \in Ob_{\mathcal{C}}$ , there are exist arrows  $ev_{A,B} : A \otimes (A \Rightarrow B) \rightarrow B$  and  $ev'_{A,B} : (B \Leftarrow A) \otimes A \rightarrow B$ , such that for all  $f : A \otimes C \rightarrow B$ :

$$(a) \ \Lambda_l \circ (id_A \otimes \mathbf{curry}(f)) = f;$$

$$(b) \ \Lambda_r \circ (\mathbf{curry}'(f) \otimes id_A) = f$$

**Definition 5.** Let  $F$  be endofunctor and  $A \in Ob_{\mathcal{C}}$ , then a coalgebra of  $F$  is a tuple  $\langle A, \theta \rangle$ , where  $\theta : A \rightarrow FA$ .

Given coalgebras  $\langle A, \theta \rangle$  and  $\langle A, \psi \rangle$ , a homomorphism is a morphism  $f : A \rightarrow B$ , s.t. the diagram below commutes:

$$\begin{array}{ccc}
A & \xrightarrow{\theta} & FA \\
f \downarrow & & \downarrow Ff \\
B & \xrightarrow{\psi} & FB
\end{array}$$

that is,  $Ff \circ \theta = \psi \circ f$

**Definition 6.** *Subexponential model structure*

Let  $\Sigma = \langle I, \leq, W, C, E \rangle$  be a subexponential model structure and  $\mathcal{C}$  be a biclosed monoidal category, then a subexponential model structure is  $\langle \mathcal{C}, \{\mathcal{F}_s\}_{s \in I} \rangle$  with the following additional data:

- for all  $s \in I$ ,  $\mathcal{F}_s$  is a monoidal comonad;
- if  $s \in W$ , then for all  $A \in \text{Ob}(\mathcal{C})$ , there exists a morphism  $w_{As} : F_s A \rightarrow \mathbb{1}$ ;
- if  $s \in C$ , then for all  $A \in \text{Ob}(\mathcal{C})$ , there exists morphisms  $w_{Al} : F_s A \otimes B \rightarrow F_s A \otimes A \otimes F_s A$  and  $w_{Ar} : B \otimes F_s A \rightarrow F_s A \otimes A \otimes F_s A$ ;
- if  $s \in E$ , then for all  $A \in \text{Ob}(\mathcal{C})$ , there is an isomorphism,  $e_A : F_s A \otimes B \cong B \otimes F_s A$ ;
- if  $s_1 \in W$ ,  $s_2 \in I$  and  $s_1 \leq s_2$ , then there is a morphism  $w_{As_2} : F_{s_2} A \rightarrow \mathbb{1}$  for all  $A \in \text{Ob}(\mathcal{C})$  and ditto for  $E$  and  $C$ ;
- Let  $\bigotimes_{s \in J, i=0}^n F_s A$ , where  $J \subset I$ , and  $s' \in I$ , s.t.  $s \geq s'$  for all  $s \in J$ ; Then there exists a morphism  $\theta : \bigotimes_{s \in J, i=0}^n F_s A \rightarrow F_{s'}(\bigotimes_{s \in J, i=0}^n F_s A)$ , such that  $\langle \bigotimes_{s \in J, i=0}^n F_s A, \theta \rangle$  is a coalgebra on  $F_{s'}$ .

**Theorem 1.** *The following statements are equivalent:*

- $SMLC_\Sigma + (\text{cut}) \vdash \Gamma \Rightarrow A$
- $SMLC_\Sigma \vdash \Gamma \Rightarrow A$
- $\exists f, f : \llbracket \Gamma \rrbracket \rightarrow \llbracket A \rrbracket$

*Proof.* • (1)  $\Rightarrow$  (2): cut elimination.

- (2)  $\Rightarrow$  (3): Soundness:

$$\begin{array}{c}
\overline{id_A : A \rightarrow A} \\
\\
\frac{f : \Gamma \rightarrow A \quad g : \Delta \otimes B \otimes \Theta \rightarrow C}{g \circ (id_\Delta \otimes (ev_{A, B_l} \circ (f \otimes id_{A \multimap B})) \otimes id_\Theta) : \Delta \otimes (\Gamma \otimes A \multimap B) \otimes \Theta \rightarrow C} \\
\\
\frac{f : A \otimes \Pi \rightarrow B}{\Lambda_l(f) : \Pi \rightarrow A \multimap B} \\
\\
\frac{f : \Gamma \rightarrow A \quad g : \Delta \otimes B \otimes \Theta \rightarrow C}{g \circ (id_\Delta \otimes (ev_{A, B_l} \circ (id_{B \multimap A} \otimes f)) \otimes id_\Theta) : \Delta \otimes (B \multimap A \otimes \Gamma) \otimes \Theta \rightarrow C}
\end{array}$$

$$\frac{f : \Pi \otimes A \rightarrow B}{\Lambda_r(f) : \Pi \rightarrow B \multimap A}$$

$$\frac{f : \Gamma \otimes A \otimes B \otimes \Delta \rightarrow C}{f \circ (\alpha_{\Gamma, A, B} \otimes id_{\Delta}) : \Gamma \otimes (A \otimes B) \otimes \Delta \rightarrow C}$$

$$\frac{f : \Gamma \rightarrow A \quad g : \Delta \rightarrow B}{f \otimes g : \Gamma \otimes \Delta \rightarrow A \otimes B}$$

$$\overline{id_{\mathbb{1}} : \mathbb{1} \rightarrow \mathbb{1}}$$

$$\frac{f : \Gamma \otimes \Delta \rightarrow A}{f \circ (\rho_{\Gamma} \otimes id_{\Delta}) : (\Gamma \otimes \mathbb{1}) \otimes \Delta \rightarrow A}$$

- Completeness:

□

## 1 Concrete model

**Definition 7.** *Quantale*