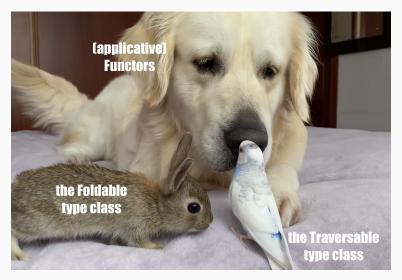
# Functional programming, Seminar No. 5

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# **Today**

# We will study



# **Functor**

#### Motivation

· Let us have a look at the following functions:

```
map :: (a -> b) -> [a] -> [b]
map _ [] = []
map f (x : xs) = f x : map f xs

mapMaybe :: (a -> b) -> Maybe a -> Maybe b
mapMaybe _ Nothing = Nothing
mapMaybe f (Just x) = Just (f x)
```

- These function are similar. Here we have a unary function that we carry through a parametrised type.
- We generalise that with the type class Functor.

#### **Functor: the definition and instances**

```
class Functor (f :: * -> *) where
  fmap :: (a -> b) -> (f a -> f b)
instance Functor Maybe where
  fmap _ Nothing = Nothing
  fmap f (Just x) = Just (f x)
instance Functor [] where
  map [] = []
 map f (x : xs) = f x : map f xs
```

#### The full definition of Functor

```
class Functor (f :: * -> *) where
  fmap
           :: (a -> b) -> f a -> f b
  (<\$) :: a -> f b -> f a
  (<$)
         = fmap . const
infixl 4 <$>, <$
(<\$>) :: Functor f => (a -> b) -> f a -> f b
(<\$>) = fmap
void :: Functor f \Rightarrow f a \rightarrow f ()
void x = () < x
```

# An example of a Functor instance

```
import Data.Functor
data Tree a = Leaf a | Node (Tree a) a (Tree a)
  deriving Show
instance Functor Tree where
  fmap f (Leaf a) = Leaf (f a)
  fmap f (Node ls a rs) = Node (fmap f ls) (f a) (fmap f rs)
left = Node (Leaf 2) 3 (Leaf 5)
right = Node (Leaf 5) 7 (Leaf 11)
tree = Node left 13 right
treeWord = (\x -> \text{show } x ++ \text{show } x) < > \text{tree}
voidTree = void tree
constTree = "Anna" <$ treeWord</pre>
                                                                  5/30
```

### The example of a Functor instance. The DeriveFunctor

One may derive the Functor instance automatically for some data types.

```
{-# LANGUAGE DeriveFunctor #-}
import Data.Functor

data Tree a = Leaf a | Node (Tree a) a (Tree a)
    deriving (Show, Functor)
```

The derived instance in this example is equivalent to our version.

```
treeWord = (\x -> show x ++ show x) <$> tree
voidTree = void tree
constTree = "Anna" <$ treeWord</pre>
```

# Functor instances for two-parametric data types

Let us take a look at the Functor for type constructors that have kind \* -> \* -> \* instance Functor ((,) a) where fmap f (x,y) = (x, f y)instance Functor ((->) r) where fmap = (.)instance Functor (Either a) where  $fmap \_ (Left x) = Left x$ fmap f (Right y) = Right (f y)

### The Functor laws

Any Functor instance has to obey the following axioms:

```
fmap id fx = fx

fmap (f . g) fx = (fmap f . fmap g) fx
```

# The Functor laws. Example

Let us check that the list data type is really a functor.

```
fmap id [] = map id [] = []
fmap id (x : xs) =
  id x : fmap id xs =
  x : fmap id xs = -- IH
  x : xs
fmap (f . g) [] = []
fmap (f . g) (x : xs) =
  (f \cdot g) \times fmap (f \cdot g) \times fmap = --IH
  (f \cdot g) \times (fmap f \cdot fmap g) \times =
  f (g x) : fmap f (fmap g xs)
```

**Applicative functors** 

#### **Motivation**

It is clear that we would like to have something like fmap for functions of an arbitrary arity:

```
fmap2
    :: (a -> b -> c)
    -> f a -> f b -> f c

fmap3
    :: (a -> b -> c -> d)
    -> f a -> f b -> f c -> f d

fmap4
    :: (a -> b -> c -> d -> e)
    -> f a -> f b -> f c -> f d -> e
```

# The Applicative class

```
class Functor f => Applicative f where
    {-# MINIMAL pure, ((<**>) | liftA2) #-}
    pure :: a -> f a

    (<**>) :: f (a -> b) -> f a -> f b
    (<**>) = liftA2 id -- the same as liftA2 ($)

liftA2 :: (a -> b -> c) -> f a -> f b -> f c
liftA2 f x = (<**) (fmap f x)</pre>
```

# The Applicative class. A couple of examples

```
instance Applicative Maybe where
  pure = Just
  Nothing <*> _ = Nothing
  _ <*> Nothing = Nothing
  Just f <*> Just x = Just (f x)

instance Applicative [] where
  pure x = [x]
  fs <*> fx = [ f x | f <- fs, x <- xs]</pre>
```

# The Applicative laws

# The Applicative class. Another Applicative instance for lists

- The list data type might have an alternative Applicative instance
- Recall the function zipWith:

```
zipWith :: (a -> b -> c) -> [a] -> [b] -> [c]
zipWith _ [] _ = []
zipWith _ _ [] = []
zipWith f (x:xs) (y:ys) = f x y : zipWith f xs ys
```

- The signature of zipWith corresponds to the signature of liftA2
- On the other hand, we cannot have two instances for the same data type

# The Applicative class. Another Applicative instance for lists

```
newtype ZipList a = ZipList { getZipList :: [a] }
instance Functor ZipList where
  fmap f = ZipList . getZipList . fmap f
instance Applicative ZipList where
  liftA2 f (ZipList xs) (ZipList ys) =
   ZipList (zipWith f xs ys)
  zipF <*> zipX = liftA2 ($)
  pure = ???
```

How to implement pure and preverse the applicative laws? If we define pure as below, that would break all those axioms.

```
pure x = ZipList [x]
```

# The Applicative class. Another Applicative instance for lists

· Here is the proper instance:

```
instance Applicative ZipList where
liftA2 f (ZipList xs) (ZipList ys) =
   ZipList (zipWith f xs ys)
zipF <*> zipX = liftA2 ($)
pure a = ZipList $ iterate x
   where iterate x = x : iterate x
```

# Monoid

### The Monoid definition

```
class Semigroup a where
  (<>) :: a -> a -> a
    -- a binary associative operation
class Semigroup a => Monoid a where
  mempty :: a
    -- a neutral element
  mappend :: a -> a -> a
 mappend = (<>)
  mconcat :: [a] -> a
  mconcat = foldr mappend mempty
```

The operation in a semigroup should associative and mempty is a neutral element

#### The Monoid instances

```
instance Semigroup [a] where
  (<>) = (++)

instance Monoid [a] where
  mempty = []
```

It is clear that (++) is an associative operation and [] is neutral.

#### **Numbers and Booleans as monoids**

```
newtype Sum a = Sum { getSum :: a }
  deriving (Show, Eq, Ord)

instance Num a => Semigroup (Sum a) where
  Sum a <> Sum b = Sum (a + b)

instance Num a => Monoid (Sum a) where
  mempty = Sum 0
```

One has the Monoid instance for any numerical type with the product as a binary operation. For that one needs to introduce the following new type:

```
newtype Sum a = Sum { getSum :: a }
deriving (Show, Eq, Ord)
```

### **Numbers and Booleans as monoids**

```
newtype All = All { getAll :: Bool }
  deriving (Show, Eq, Show)

instance Semigroup All where
  All a <> All b = All (a && b)

instance Monoid All where
  mempty = All True
```

The same for disjunction.

# Back to the Applicative class

The Monoid type class allows one to have the Applicative instance for tuples as follows:

```
instance Monoid a => Applicative ((,) a) where
  pure x = (mempty x, x)
  (a, f) <*> (b, x) = (a <> b, f x)
```

# Foldable

#### **Motivation**

· Before we took a look at such fold functions as foldr

```
foldr :: (a -> b -> b) -> b -> [a] -> b
foldr _ ini [] = []
foldr f ini (x : xs) = f x (foldr f ini xs)
```

- One may generalise the idea of folding and consider a broader class of foldable data structures
- Before that, we introduce the type class called Foldable

# The Foldable type class

```
class Foldable t where
  {-# MINIMAL foldMap | foldr #-}
  fold :: Monoid m => t m -> m
  fold = foldMap id
  foldMap :: Monoid m => (a -> m) -> t a -> m
  foldMap f = foldr (mappend . f) mempty
  foldr :: (a \rightarrow b \rightarrow b) \rightarrow b \rightarrow t a \rightarrow b
  foldr f z t = appEndo (foldMap (Endo . f) t) z
  foldl :: (b \rightarrow a \rightarrow b) \rightarrow b \rightarrow t a \rightarrow b
     foldl f z t =
       appEndo (getDual (foldMap (Dual . Endo . flip f) t)) z
```

# Useful functions for foldable data types

Here we provide type signatures only:

```
toList :: Foldable t => t a -> [a]
null :: Foldable => t a -> Bool
length :: Foldable t => t a -> Int
elem :: (Eq a, Foldable t) => a -> t a -> Bool
maximum :: (Ord a, Foldable t) => t a -> a
sum, product :: (Num a, Foldable t) => t a -> a
```

#### Foldable instances.

```
instance Foldable [] where
  elem = List.elem
  foldl = List.foldl
  foldr = List.foldr
  length = List.length
  maximum = List.maximum
  product = List.product
```

# Foldable instances. Other examples

```
instance Foldable (Either a) where
   foldMap _ (Left _) = mempty
   foldMap f (Right y) = f y
   foldr _z (Left _) = z
   foldr f z (Right v) = f v z
   length (Left _) = 0
   length (Right _) = 1
   null
                     = isLeft
instance Foldable ((,) a) where
   foldMap f (_, y) = f y
   foldr f z (_, y) = f y z
```

**Traversable** 

# The motivating example

```
dist :: Applicative f => [f a] -> f [a]
dist [] = pure []
dist (x : xs) = liftA2 (:) x (dist xs)

> dist (Just <$> [1,2,4])
Just [1,2,4]
> dist [Just 1, Nothing]
Nothing
> getZipList $ dist $ map ZipList [[1,2,3], [4,5,6], [7,8,9]]
[[1,4,7],[2,5,8],[3,6,9]]
```

#### The Traversable definition

According to the documentation, Traversable describes "functors representing data structures that can be traversed from left to right".

```
class (Functor t, Foldable t) => Traversable t where
  traverse :: Applicative f => (a -> f b) -> t a -> f (t b)
  traverse f = sequenceA . fmap f

sequenceA :: Applicative f => t (f a) -> f (t a)
  sequenceA = traverse id
  {-# MINIMAL traverse / sequenceA #-}
```

#### The Traversable instances

```
instance Travesable Maybe where
  traverse _ Nothing = Nothing
  traverse f (Just x) = Just <$> f x

instace Traversable [] where
  traverse _ g = foldr consF (pure [])
  where
  consF x ys = liftA2 (:) (g x) ys
```

# Summary

# **Summary**

#### Today we

• introduced such type classes as Functor, Applicative, Monoid, Foldable, and Traversable

### **Summary**

#### Today we

• introduced such type classes as Functor, Applicative, Monoid, Foldable, and Traversable

Next time, we will

· study monads!