

Stuff related to tensor products of modal algebras

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1 Preliminaries

The basic definitions according to [1].

Definition 1. An n -modal algebra is an algebra $\mathcal{B} = \langle B, \vee, \neg, \perp, \Diamond_1, \dots, \Diamond_n \rangle$, where $\mathcal{B} = \langle B, \vee, \neg, \perp \rangle$ is a Boolean algebra and for each $i = 1, \dots, n$ \Diamond_i is an order-preserving unary function obeying:

1. $\Diamond_i p \vee \Diamond_i q = \Diamond_i(p \vee q)$
2. $\Diamond_i \perp = \perp$

Definition 2. An n -normal modal logic is a set of formulas containing Boolean tautologies, formulas $\Diamond_i p \vee \Diamond_i q \leftrightarrow \Diamond_i(p \vee q)$ and $\Diamond_i \perp \leftrightarrow \perp$ for $i = 1, \dots, n$, and is closed under **MP**, **Sub**, and the monotonicity rule: from $\varphi \rightarrow \psi$ infer $\Diamond_i \varphi \rightarrow \Diamond_i \psi$.

Definition 3.

1. A Kripke n -frame is a structure $\mathcal{F} = \langle W, R_1, \dots, R_n \rangle$, where $R_i \subseteq W \times W$ for each $i = 1, \dots, n$.
2. Let \mathcal{F} be a Kripke frame. A Kripke model over \mathcal{F} is a tuple $\mathcal{M} = \langle \mathcal{F}, \vartheta \rangle$, where $\vartheta : PV \rightarrow 2^W$ is a valuation. Modalised formulas have the following semantics:

$$\mathcal{M}, x \models \Diamond_i \varphi \Leftrightarrow \exists y \in R_i(x) \mathcal{M}, y \models \varphi.$$

A formula φ is true in \mathcal{M} iff $\|\varphi\|_{\mathcal{M}} = W$. φ is valid in $\mathcal{F} = \langle W, R_1, \dots, R_n \rangle$ iff $\|\varphi\|_{\mathcal{M}} = W$ for every model \mathcal{M} over \mathcal{F} .

3. The complex algebra of a Kripke frame $\mathcal{F} = \langle W, R_1, \dots, R_n \rangle$ is a complete modal algebra defined as $\mathcal{F}^+ = \langle \mathcal{P}(W), \cup, \neg, \emptyset, R_1^{-1}, \dots, R_n^{-1} \rangle$.
4. A general n -frame is a structure $\mathcal{F} = \langle W, R_1, \dots, R_n, \mathcal{A} \rangle$, where \mathcal{A} is a subalgebra of $\langle W, R_1, \dots, R_n \rangle^+$.

Definition 4. Let \mathcal{F} be an n -frame, then $\text{Log}(\mathcal{F}) = \{\varphi \mid \mathcal{F} \models \varphi\}$. If \mathbb{F} is a class of n -frames, then $\text{Log}(\mathbb{F}) = \bigcap_{\mathcal{F} \in \mathbb{F}} \text{Log}(\mathcal{F})$. We use the same notation for logics of general frames and modal algebras.

We discuss the background on products of modal logics and products of Kripke frames [8] [9].

Definition 5. Let $\mathcal{F}_1 = \langle W_1, R_1^1, \dots, R_1^n \rangle$ be an n -frame and $\mathcal{F}_2 = \langle W_2, R_2^1, \dots, R_2^m \rangle$. The product frame of \mathcal{F}_1 and \mathcal{F}_2 is an $n + m$ -frame of the form

$$\mathcal{F}_1 \times \mathcal{F}_2 = \langle W_1 \times W_2, R_h^1, \dots, R_h^n, R_v^1, \dots, R_v^m \rangle$$

such that for all $u_1, u_2 \in W_1$ and for all $v_1, v_2 \in W_2$,

$$\begin{aligned} \langle u_1, v_1 \rangle R_h^i \langle u_2, v_2 \rangle &\text{ iff } u_1 R_1^i u_2 \text{ and } v_1 = v_2 \text{ for } 1 \leq i \leq n. \\ \langle u_1, v_1 \rangle R_v^j \langle u_2, v_2 \rangle &\text{ iff } u_1 = u_2 \text{ and } u_1 R_2^j u_2 \text{ for } 1 \leq j \leq m. \end{aligned}$$

Let \mathcal{L}_1 be a normal n -modal logic and \mathcal{L}_2 a normal m -modal logic, the product of \mathcal{L}_1 and \mathcal{L}_2 is defined as

$$\mathcal{L}_1 \times \mathcal{L}_2 = \text{Log}(\text{Frames}(\mathcal{L}_1) \times \text{Frames}(\mathcal{L}_2))$$

Proposition 1.

1. Let $\mathcal{L}_1, \mathcal{L}_2$ be modal logics, then $\mathcal{L}_1 * \mathcal{L}_2 \subseteq \mathcal{L}_1 \times \mathcal{L}_2$.
2. Let $\mathcal{L}_1, \mathcal{L}_2$ be Kripke complete modal logics, then $\mathcal{L}_1 \times \mathcal{L}_2 = \text{Log}(\text{Frames}_r(\mathcal{L}_1) \times \text{Frames}_r(\mathcal{L}_2))$, where $\text{Frames}_r(\mathcal{L}_i) = \{\mathcal{F} \in \text{Frames}(\mathcal{L}_i) \mid \mathcal{F} \text{ is rooted}\}$ for $i = 1, 2$.

1.1 Axiomatising products

The following properties hold for a product frame having the form $\mathcal{F} = \langle W, R_1, \dots, R_n \rangle$:

1. commutativity 1
2. commutativity 2
3. confluence

2 Tensor products of modal algebras. The basic definitions and results

Definition 6. A commutative associative algebra

Definition 7. Tensor product of them

Definition 8. Tensor product of modal algebras

Definition 9. Tensor product of general frames

TODO: describe related underlying results

3 Note on incomplete modal logics

TODO: consider the system containing **GL**, McKinsey, seriality, and linearity.

4 Solution of Problem 1

5 Solution of Problem 4

Definition 10. A finitely axiomatisable normal modal logic

References

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