

Varieties of representable relation algebra reducts

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1 Question

Let τ be a subsignature of operations expressible in the signature of RA . $\mathbf{R}(\tau)$ is a variety iff ??.
Investigate discriminators for representable relation algebra reducts.

2 Varieties and discriminators

Standardly, a class of algebras is called *variety*, if it can be determined by some equational theory, or, equivalently, it is closed under homomorphic images, subalgebras and direct products. Given a class of algebras \mathcal{K} , $\mathbf{V}(\mathcal{K})$ is a variety generated by \mathcal{K} or, equivalently, $\mathbf{HSP}(\mathcal{K})$, the closure of \mathcal{K} under homomorphic images, subalgebras and direct products.

Let $\{A_i \mid i \in I\}$ be an indexed family of algebras, then a subalgebra $A \subseteq \prod_{i \in I} A_i$ is a subdirect product if $\pi_i(A) = A_i$. An embedding $\alpha : A \rightarrow \prod_{i \in I} A_i$ is subdirect if $\alpha(A)$ is subdirect product. An algebra is subdirectly irreducible if for every subdirect embedding $\alpha : A \rightarrow \prod_{i \in I} A_i$ there exists $i \in I$ such that $\pi_i \circ \alpha : A \rightarrow A_i$ is an isomorphism.

An equivalence relation θ on an algebra A is called *congruence*, if θ respects any operation. $\mathbf{Con}(A)$ is the set of all congruences on A . An algebra is called *simple*, if $\mathbf{Con}(A/\theta) = \{\Delta, \nabla\}$, where Δ and ∇ are trivial congruences. One can obtain a simple algebra by factorising it through the maximal congruence [SB81, Theorem 8].

One can equivalently define subdirectly irreducible algebras using congruences as follows. an algebra is subdirectly irreducible iff either A is trivial or there exists a minimal congruence in $\mathbf{Con}(A) \setminus \{\Delta\}$.

Recall that a Stone space is a compact Hausdorff zero-dimensional topological space. A subdirect product $A \subseteq \prod_{x \in X} A_x$ over a Stone space X if

1. for all $a, b \in A$ $\{x \in X \mid a(x) = b(x)\}$ is clopen.
2. for all $a, b \in A$ and for all clopen $Y \subseteq X$ $a \upharpoonright_Y \cup a \upharpoonright_{X \setminus Y} \in A$.

A variety \mathcal{V} is *arithmetical*, if it is congruence-permutable and congruence-distributive, or, equivalently, there exists a ternary term p such that:

$$\mathcal{V} \models p(x, y, x) \approx p(x, y, y) \approx p(y, y, x) \approx x \quad (1)$$

A ternary term $t(x, y, z)$ for an algebra A if, for all $a, b, c \in A$:

$$t(a, b, c) = \begin{cases} c & \text{if } a = b \\ a & \text{otherwise} \end{cases} \quad (2)$$

A variety \mathcal{V} is called discriminator if there exists a class \mathcal{K} such that $\mathcal{V} = \mathbf{V}(\mathcal{K})$ and there exists a term $t(x, y, z)$, which is a discriminator term for every member of \mathcal{K} . It is known that if

an algebra A has a discriminator term, then A is simple [SB81, Lemma 9.2]. Moreover, we have the following property of discriminator terms, see [SB81, Theorem 9.4].

Theorem 1. *Let $t(x, y, z)$ be a discriminator term for every member of a class \mathcal{K} :*

1. $\mathbf{V}(\mathcal{K})$ is an arithmetical variety.
2. Every indecomposable member of $\mathbf{V}(\mathcal{K})$ is simple.
3. Simple algebras are precisely members of $\text{ISP}_U(\mathcal{K}_+)$.
4. Every member of $\mathbf{V}(\mathcal{K})$ is isomorphic to a Boolean product of simple algebras.

3 BAOs, relation algebras and their reducts

3.1 Discriminator varieties of BAOs

Let B be a Boolean algebra, an operator is an n -ary function $f : A^n \rightarrow A$ such that, for all $x_1, \dots, x_n, x, y \in B$:

- $f(x_1, \dots, x + y, \dots, x_n) = f(x_1, \dots, x, \dots, x_n) + f(x_1, \dots, y, \dots, x_n)$
- $f(x_1, \dots, 0, \dots, x_n) = 0$

A Boolean algebra with operators is an algebra $M = (B, (f_i)_{i \in I})$, where each f_i is an operator.

In the case of BAOs, one can define discriminator simpler, as a unary term $d(x)$ such that, for all $a \in M$, where M is a BAO:

$$d(x) = \begin{cases} 0 & \text{if } x = 0 \\ 1 & \text{otherwise} \end{cases} \quad (3)$$

One can characterise discriminator varieties as follows, see [AGM⁺98, Lemma 2.1]:

Theorem 2. *Let \mathcal{V} be a variety of BAOs and $d(x)$ a unary term, then the following are equivalent:*

1. d is a discriminator variety.
2. The following equations are valid in \mathcal{V} :
 - (a) $x \leq d(x)$
 - (b) $d(d(x)) \leq d(x)$
 - (c) $d(-d(x)) \leq -d(x)$
 - (d) $f(x_0, \dots, x_{n-1}) \leq d(x_i)$ for all $n > 0$ and for every operator f of M

3.2 Relation algebras and their reducts

In this subsection, we consider relation algebras, a kind of BAOs.

Definition 1.

A relation algebra is an algebra $\mathcal{R} = (R, 0, 1, +, -, ;, \smile, \mathbf{1}')$ such that $(R, 0, 1, +, -)$ is a Boolean algebra and the following hold:

1. $(R, ;, \mathbf{1}')$ is a monoid

2. $(a + b); c = (a; c) + (b; c)$
3. $a^{\smile\smile} = a$
4. $(a + b)^{\smile} = a^{\smile} + b^{\smile}$
5. $(a; b)^{\smile} = b^{\smile}; a^{\smile}$
6. $a^{\smile}; (-a; b) \leq -b$

where $a \leq b$ iff $a + b = b$. **RA** denotes the class of all relation algebras.

Definition 2. A proper relation algebra is an algebra $\mathcal{R} = (R, \emptyset, W, \cup, -, |, \smile, \mathbf{1})$ such that $R \subseteq \mathcal{P}(W)$, where $W \subseteq X \times X$ is an equivalence relation; $|$ is relation composition, \smile is relation converse, **Id** is a diagonal subset of W , that is:

1. $a|b = \{(x, z) \mid \exists y (x, y) \in a \ \& \ (y, z) \in b\}$
2. $a^{\smile} = \{(x, y) \mid (y, x) \in a\}$
3. **Id** = $\{(x, y) \mid x = y\}$

The class of all proper relation algebras is denoted as **PRA**. **Rs** is the class of all relation set algebras, proper relation algebra with a diagonal subrelation as an identity. **RRA** is the class of all representable relation algebras, that is, the closure of **PRA** under isomorphic copies.

3.3 Varieties

3.3.1 Boolean algebras with residuated operators

Let \mathcal{A}_0 be a Boolean algebra, a unary operation on f is called *residuated* if there exists a *residual* operation g such that for all $a, b \in \mathcal{A}_0$:

$$f(a) \leq b \text{ iff } a \leq f(b)$$

Equivalently, f is residuated if there exists a conjugate operation h such that for all $a, b \in \mathcal{A}_0$:

$$f(a) \cdot b = 0 \text{ iff } a \wedge h(b) = 0$$

f is *self-conjugate* whenever f is equal to its conjugate operation.

Theorem 3. Let f be a residuated operator on \mathcal{A}_0 , then

1. f is normal and completely additive,
2. f and h are conjugate operations on \mathcal{A}_0 iff they are normal and the following holds for all $a, b \in \mathcal{A}_0$

$$f(a) \cdot b \leq f(a \cdot h(b)) \text{ and } a \cdot h(b) \leq h(f(a) \cdot b)$$

Two n -ary operators f and h are conjugate in the i -th argument if $f_{\bar{a},i}$ is conjugate to $g_{\bar{a},i}$. Let \mathcal{A} be a BAO, \mathcal{A} is called a *Boolean algebra with residuated operators* if for each operator f of \mathcal{A} and for each $i < ar(f)$ there exists an $ar(f)$ -ry term t , which is conjugate to f in the i -th argument. **RA** is an example, where \smile is self-conjugate and composition has conjugate terms $x; y^{\smile}$ and $x^{\smile}; y$.

An r -algebra is a Boolean algebra with three residuated binary operations \bullet , \triangleright and \triangleleft , where \triangleright and \triangleleft are right and left conjugates to \bullet , that is, the following conditions are equivalent:

1. $(x \bullet y) \cdot z = 0$
2. $(x \triangleright z) \cdot y = 0$
3. $(z \triangleleft y) \cdot x = 0$

3.3.2 Residuated monoids

3.3.3 Positive relation algebras

3.3.4 Domain algebras

3.4 Non-varieties

3.5 Unknown

4 Decidability aspects

4.1 Current results

4.2 Problems

References

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