# Characterising representable positive relation algebras via Priestley duality

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# 1 Distributive lattice representation and Priestley duality

Given a bounded distributive lattice  $\mathcal{L}$ , a proper subset  $F \subset \mathcal{L}$  is said to be a *filter* if it is upward closed and closer under finite infima. A filter F is *prime* if  $a + b \in F$  implies either  $a \in F$  or  $b \in F$ . The spectrum of  $\mathcal{L}$ , denoted as  $\text{Spec}(\mathcal{L})$ , is the set of all prime filters.

A filter is *complete* if whenever  $\Pi T$  exists for  $T \subseteq F$ , then  $\Pi T \in F$ . A filter is *completely prime* if whenever  $\Sigma T$  exists for  $T \subseteq F$ , then there exists  $t \in T$  such that  $t \in F$ . The dual definitions are for ideals.

**Proposition 1.** Let  $h: \mathcal{L} \to \mathcal{R}$  be a representation, then then

$$h^{-1}[x] = \{a \in \mathcal{L} \mid x \in h(a)\} \in \operatorname{Spec}(\mathcal{L})$$

Recall that a *Priestley space* is a triple  $\mathcal{X} = (X, \tau, \leq)$  such that  $(X, \tau)$  is a compact topological space,  $(X, \leq)$  is a bounded poset such that if  $x \leq y$ , then there exists a clopen U such that  $x \in U$  and  $y \notin U$ . Given a bounded distributive lattice  $\mathcal{L}$ , define the map  $\phi : \mathcal{L} \to 2^{\operatorname{Spec}(\mathcal{L})}$  such that

$$\phi: a \mapsto \{F \in \operatorname{Spec}(\mathcal{L}) \mid a \in F\}$$

**Fact 1.** 1. The sets  $\phi(a)$  and  $-\phi(a)$  form the subbasis of the topology  $\tau$  on Spec( $\mathcal{L}$ ).

2.  $(\operatorname{Spec}(\mathcal{L}), \tau, \subseteq)$  is a Priestley space.

Given a Priestley space  $\mathcal{X} = (X, \tau, \leq)$ , the set  $ClOp(\mathcal{X})$  consists of all clopens of  $\mathcal{X}$ . The structure  $(ClOp(\mathcal{X}), \cap, \cup, \emptyset, X)$  is a distributive lattice.

**Fact 2.** Let  $\mathcal{L}$  be a distributive lattice and let  $\mathcal{X}$  be a Priestley space:

- 1.  $\mathcal{L} \hookrightarrow \mathcal{L}^+ = (2^{\operatorname{Spec}(\mathcal{L})}, \cap, \cup, \varnothing, \operatorname{Spec}(\mathcal{L})),$
- 2.  $\mathcal{L} \cong \mathrm{ClOp}(\mathrm{Spec}(\mathcal{L})),$
- 3.  $\mathcal{X} \cong \operatorname{Spec}(\operatorname{ClOp}(\mathcal{X})),$
- 4. The categories of Priestley spaces and bounded distributive lattices are dually equivalent.

## 2 Representatiting positive relation algebras

**Definition 1.** A positive relation algebra is a algebra  $\mathcal{R} = (R, \cdot, +, \cdot, \cdot, \cdot, 0, 1, 1')$  such that

- 1.  $(R, \cdot, +, 0, 1)$  is a bounded distributive lattice,
- 2.  $(R,;,\mathbf{1}')$  is a monoid,
- 3. for all  $a, b, c \in R$

(a) 
$$a; (b+c) = a; b+a; c,$$

(b) 
$$a^{\smile} = a$$
,

(c) 
$$(a+b)^{\smile} = a^{\smile} + b^{\smile}$$
,

$$(d) (a;b)^{\smile} = b^{\smile}; a^{\smile},$$

(e) 
$$(a;b) \cdot c^{\smile} = 0 \leftrightarrow (b;c) \cdot c^{\smile} = 0.$$

A positive relation algebra  $\mathcal{R}$  is representable if there exists a one-to-one function  $h: \mathcal{R} \to 2^{X \times X}$  over the base set  $X \neq \emptyset$  such that:

- $f(a \cdot b) = f(a) \cap f(b)$ ,
- f(a + b) = f(a) + f(b),
- $f(0) = \emptyset$ ,
- $f(1) = \bigcup_{a \in \mathcal{R}} f(a),$
- $f(\mathbf{1}') = \Delta_X$ ,
- $f(a;b) = \{(x,z) \mid \exists y \in X ((x,y) \in f(a) \& (y,z) \in f(b))\} = f(a)|f(b),$
- $f(a) = \{(y, x) \mid (x, y) \in f(a)\}.$

## 3 Spectral spaces for positive relation algebras

A PRA-space is a structure  $(X, \tau, \leq, R, I, E)$  where  $X = (X, \tau, \leq)$  is a Priestley space and  $R \subseteq X^3$ ,  $I \subseteq X^2$  and  $E \subseteq X$  such that:

- For all  $x, y, z, w \in X$  there exists  $u \in X$  such that R(x, y, u) and R(u, z, w) iff there exists  $v \in X$  such that R(y, z, v) and R(x, v, w),
- If  $A, B \subseteq X$  are upward closed, so is  $R[A, B, \_]$ ,
- I(A) is upward closed clopen whenever A is upward closed clopen,
- I(x) is closed for each  $x \in X$ ,
- For all  $x, y, z \in X$ ,  $x \leq y$  and I(x, z) imply I(y, z),
- For all  $x, y \in X$  there exists  $z \in X$  such that x = y iff I(z, y) and I(x, z),
- For all  $x, y, z \in X$  there exists  $u \in X$  such that I(u, z) and R(x, y, u) iff there exist  $u, w \in X$  such that R(v, w, z), I(y, v) and I(x, w).

- For all  $x, y, u, v \in X$ , R(u, v, y) and I(x, u) implies R(x, y, v),
- E is upward closed clopen such that for each clopen  $A \subseteq X$  one has

$$R[A, E, \_] = R[E, A, \_] = A$$

**Lemma 1.** Let  $\mathcal{R}$  be a positive relation algebra, then

1.

2.

## 4 Complete representability

## 4.1 Completely representable distibutive lattices

Let  $\mathcal{L}$  be a bounded distributive lattice, then a set  $S \subseteq 2^{\mathcal{L}}$  is said to be distinguishing if for every  $a, b \in \mathcal{L}$  such that  $a \neq b$  there exists  $s \in S$  such that either  $a \in s$  and  $b \notin b$  or vice versa.

**Theorem 1.** Let  $\mathcal{L}$  be a bounded distributive lattice, then

- 1.  $\mathcal{L}$  is completely representable iff  $\mathcal{L}$  has a distinguishing set of complete, completely prime filters,
- 2.  $(\mathcal{L}_+)^+$  is completely representable.

TODO: read [EH12]

### 4.2 Completely representable positive relation algebras

#### 5 The main result

For that we need such model theoretic notions as saturation and types, see [Hod93, Section 6.3].

**Definition 2.** Let  $\mathcal{M}$  be a first-order structure of a signature L and  $S \subseteq \mathcal{M}$ . Let L(S) be an extension of L with copies of elements from S as additional constants. We assume that Cnst(L) and S are disjoint.

- 1. Let  $n < \omega$ , an n-type over S is a set  $\mathcal{T}$  of L(S) formulas  $A(\overline{x})$ , where  $\overline{x}$  is a fixed n-tuple of elements from S. Notation:  $\mathcal{T}(\overline{x})$ . A type is an n-type for some  $n < \omega$ .
- 2. An n-type  $\mathcal{T}(\overline{x})$  is realised in  $\mathcal{M}$ , if there exists  $\overline{m} \in \mathcal{M}^n$  such that  $\mathcal{M} \models A(\overline{m})$  for every  $A \in \mathcal{T}(\overline{x})$ .  $\mathcal{M}$  omits  $\mathcal{T}(\overline{x})$ , if  $\mathcal{T}(\overline{x})$  is not realised in  $\mathcal{M}$ .
- 3.  $\mathcal{T}(\overline{x})$  is finitely satisfied in  $\mathcal{M}$ , if every finite subtype  $\mathcal{T}_0(\overline{x}) \subseteq \mathcal{T}(\overline{x})$  is realised in  $\mathcal{M}$ . We can reformulate that as  $\mathcal{M} \models \exists \overline{a} \bigwedge_{A \in \mathcal{T}_0} A(\overline{a})$ .
- 4. Let T be a theory, then a type  $\mathcal{T}$  over the empty set of constants is T-consistent, if there exists a model  $\mathcal{M} \models T$  such that  $\mathcal{T}$  is finitely satisfied in  $\mathcal{M}$ .
- 5. Let  $\kappa$  be a cardinal, then  $\mathcal{M}$  is  $\kappa$ -saturated, if for every  $S \subseteq \mathcal{M}$  with  $|S| < \kappa$  every finitely satisfied 1-type  $\mathcal{T}$  is realised in  $\mathcal{M}$ .

By default, a saturated model is an  $\omega$ -saturated model for us.

The useful facts, they are from [CK90] and [Hod93]:

**Fact 3.** Let  $\mathcal{M}$  be an FO-structue and  $\kappa$  a cardinal, then:

- 1.  $\mathcal{M}$  is  $\kappa$ -saturated iff every finitely satisfiable  $\alpha$ -type (an arbitrary  $\alpha \leq \kappa$ ) with fewer than  $\kappa$  parameters is realised in  $\mathcal{M}$ .
- 2. If  $\mathcal{M}$  is  $\kappa$ -saturated, then  $\mathcal{M}$  is  $\lambda$ -saturated for every  $\lambda < \kappa$ .
- 3. Every consistent theory has a  $\kappa$ -saturated model and every model has an elementary  $\kappa$ -saturated extension.
- 4. Let  $(\mathcal{M}_i)_{i<\omega}$  a family of structures of the (at most) countable signature and D a non-principal ultrafilter over  $\omega$ , then  $\Pi_D \mathcal{M}_i$  is  $\omega_1$ -saturated.

Let  $\mathcal{A}$  be a positive relation algebra, define the first-order relational language of the form

$$\mathcal{L}(\mathcal{A}) = (=, \{R_a^2\}_{a \in \mathcal{A}})$$

The  $\mathcal{L}(\mathcal{A})$ -theory  $T_{\mathcal{A}}$  consists of the following statements:

- $\sigma_1 = \forall x \forall y (\mathbf{1}'(x,y) \leftrightarrow (x=y))$
- $\sigma_+(R, S, T) = \forall x \forall y (R(x, y) \leftrightarrow S(x, y) \lor T(x, y))$
- $\sigma(R, S, T) = \forall x \forall y (R(x, y) \leftrightarrow S(x, y) \land T(x, y))$
- $\sigma_1(R, S, T) = \forall x \forall y (R(x, y) \leftrightarrow \exists z (S(x, z) \land T(z, y)))$
- $\sigma_{\smile}(R,S) = \forall x \forall y (R(x,y) \leftrightarrow S(y,x))$
- $\sigma_{\neq 0} = \exists x \exists y R(x,y)$  for any  $R_a$  such that  $a \neq 0$
- $\sigma_0 = \neg \exists x \exists y 0(x,y)$
- $\sigma_1 = \forall x \forall y (R(x,y) \to \mathbf{1}(x,y))$

**Proposition 2.**  $T_{\mathcal{A}}$  is satisfiable whenever  $\mathcal{A}$  is representable.

**Theorem 2.** Let A be a positive relation algebra, then R is representable iff  $(R_+)^+$  is completely representable.

Theorem 3. RPRA is a canonical variety.

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