

KTH Royal Institute of Technology

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Contest (1)

```
template.cpp
#include <bits/stdc++.h>
#define all(x) begin(x), end(x)
using namespace std;

using ll = long long;
int main() {
   cin.tie(0)->sync_with_stdio(0);
   cin.exceptions(cin.failbit);
}
```

10 lines

Data structures (2)

BIT.h

Description: Query [l, r] sums, and point updates. kth() returns the smallest index i s.t. query(0, i) >= k**Time:** $\mathcal{O}(\log n)$ for all ops.

ne: $O(\log n)$ for all ops. 33f78c, 22 lines

```
template <typename T>
struct BIT {
  vector<T> s;
  int n:
  BIT(int n) : s(n + 1), n(n) {}
  void update(int i, T v) {
   for (i++; i <= n; i += i & -i) s[i] += v;
  T query(int i) {
   T ans = 0;
    for (i++; i; i -= i & -i) ans += s[i];
    return ans;
  T query(int 1, int r) { return query(r) - query(l - 1); }
  int kth(T k) { // returns n if k > sum of tree
   if (k <= 0) return -1;
   int i = 0;
    for (int pw = 1 << __lg(n); pw; pw >>= 1)
     if (i + pw \le n \&\& s[i + pw] \le k) k -= s[i += pw];
    return i:
};
```

```
KDBIT.h
```

```
Description: k-dimensional BIT. BIT<int, N, M> gives an N \times M BIT.
Query: bit.query(x1, x2, y1, y2) Update: bit.update(x, y, delta)
Time: O(\log^k n) Status: Tested
template <class T, int... Ns>
struct BIT {
 T val = 0;
  void update(T v) { val += v; }
 T query() { return val; }
template <class T, int N, int... Ns>
struct BIT<T, N, Ns...> {
 BIT<T, Ns...> bit[N + 1];
  // map < int, BIT < T, Ns... >> bit;
  // if the memory use is too high
  template <class... Args>
  void update(int i, Args... args) {
    for (i++; i <= N; i += i & -i) bit[i].update(args...);</pre>
  template <class... Args>
  T query(int i, Args... args) {
    for (i++; i; i -= i & -i) ans += bit[i].query(args...);
    return ans:
  template <class... Args,
            enable_if_t<(sizeof...(Args) ==
                          2 * sizeof...(Ns))>* = nullptr>
  T query (int 1, int r, Args... args) {
    return query(r, args...) - query(l - 1, args...);
};
DSU.h
Description: Maintains union of disjoint sets
Time: \mathcal{O}(\alpha(N))
                                                       c22586, 14 lines
struct DSU {
  vector<int> s:
  DSU(int n) : s(n, -1) {}
  int find(int i) { return s[i] < 0 ? i : s[i] = find(s[i]); }</pre>
  bool join(int a, int b) {
    a = find(a), b = find(b);
    if (a == b) return false;
    if (s[a] > s[b]) swap(a, b);
    s[a] += s[b], s[b] = a;
    return true;
  int size(int i) { return -s[find(i)]; }
  bool same(int a, int b) { return find(a) == find(b); }
RMQ.h
Description: Constant time subarray min/max queries for a fixed array
Time: O(nlogn) initialization and O(1) queries. Status: Tested 536eac, 15 lines
template <typename T, class Compare = less<T>>
struct RMQ {
  vector<vector<T>> t;
  Compare cmp;
  RMQ(vector<T>& a) : t(\underline{lg(a.size())} + 1, a) {
    int n = a.size(), lg = __lg(n);
    for (int k = 1, len = 1; k <= lq; k++, len <<= 1)</pre>
      for (int i = 0; i + 2 * len - 1 < n; i++)
        t[k][i] = min(t[k-1][i], t[k-1][i+len], cmp);
 T query(int a, int b) {
    int k = __lg(b - a + 1), len = 1 << k;</pre>
```

return min(t[k][a], t[k][b - len + 1], cmp);

```
,
```

```
Splay.h
```

else

Description: An implicit balanced BST. You only need to change update() and prop().

If used for link-cut tree, code everything up to splay(). Time: amortized $O(\log n)$ for all operations

```
struct node {
 node *ch[2] = \{0\}, *p = 0;
 int cnt = 1, val;
 node (int val, node * 1 = 0, node * r = 0)
   : ch{1, r}, val(val) {}
int cnt(node* x) { return x ? x->cnt : 0; }
int dir(node* p, node* x) { return p && p->ch[0] != x; }
void setLink(node* p, node* x, int d) {
 if (p) p \rightarrow ch[d] = x;
 if (x) x->p = p;
node* update(node* x) {
 if (!x) return 0;
 x - cnt = 1 + cnt(x - ch[0]) + cnt(x - ch[1]);
  setLink(x, x->ch[0], 0);
 setLink(x, x->ch[1], 1);
 return x;
void prop(node* x) {
 if (!x) return;
 // update(x); // needed if prop() can change subtree sizes
void rotate(node* x, int d) {
 if (!x || !x->ch[d]) return;
 node *y = x - > ch[d], *z = x - > p;
  setLink(x, y->ch[d ^ 1], d);
  setLink(y, x, d^{1});
 setLink(z, y, dir(z, x));
 update(x);
 update(y);
node* splay(node* x) {
  while (x && x->p) {
    node *y = x->p, *z = y->p;
    // prop(z), prop(y), prop(x); // needed for LCT
    int dy = dir(y, x), dz = dir(z, y);
    if (!z)
      rotate(y, dy);
    else if (dy == dz)
      rotate(z, dz), rotate(y, dy);
      rotate(y, dy), rotate(z, dz);
  return x;
// the returned node becomes the new root, update the root
// pointer!
node* nodeAt(node* x, int pos) {
 if (!x) return 0;
 while (prop(x), cnt(x->ch[0]) != pos)
    if (pos < cnt(x->ch[0]))
     x = x -> ch[0];
```

```
pos -= cnt(x->ch[0]) + 1, x = x->ch[1];
  return splay(x);
node* merge(node* 1, node* r) {
 if (!1 || !r) return 1 ?: r;
 1 = nodeAt(1, cnt(1) - 1);
  setLink(l, r, 1);
  return update(1);
// first is everything < pos, second is >= pos
pair<node*, node*> split(node* t, int pos) {
  if (pos <= 0 || !t) return {0, t};</pre>
  if (pos > cnt(t)) return {t, 0};
  node *1 = nodeAt(t, pos - 1), *r = 1->ch[1];
  if (r) 1 \rightarrow ch[1] = r \rightarrow p = 0;
  return {update(1), update(r)};
// insert a new node between pos-1 and pos
node* insert(node* t, int pos, int val) {
  auto [1, r] = split(t, pos);
  return update (new node (val, l, r));
// apply lambda to all nodes in an inorder traversal
template <class F>
void each(node* x, F f) {
 if (x) \text{ prop}(x), \text{ each}(x->\text{ch}[0], f), f(x), \text{ each}(x->\text{ch}[1], f);
```

Geometry (3)

$\underline{\text{Graphs}}$ (4)

Time: $\mathcal{O}(|V| + |E|)$

cont.clear();

return val[j] = low;

template <class G, class F>

ncomps++;

SCCTarjan.h

Description: Finds strongly connected components of a directed graph. Visits/indexes SCCs in reverse topological order.

Usage: scc(graph) returns an array that has the ID of each node's SCC. $scc(graph, [\&](vector < int > \& v) { ... })$ calls the lambda on each SCC, and returns the same array.

```
namespace SCCTarjan {
 vector<int> val, comp, z, cont;
 int Time, ncomps;
  template <class G, class F>
  int dfs(int j, G& g, F& f) {
   int low = val[j] = ++Time, x;
    z.push_back(j);
   for (auto e : q[j])
     if (comp[e] < 0) low = min(low, val[e] ?: dfs(e, g, f));</pre>
    if (low == val[j]) {
     do {
       x = z.back();
        z.pop back();
       comp[x] = ncomps;
       cont.push_back(x);
      } while (x != j);
      f(cont);
```

```
vector<int> scc(G& g, F f) {
    int n = q.size();
    val.assign(n, 0);
    comp.assign(n, -1);
    Time = ncomps = 0;
    for (int i = 0; i < n; i++)</pre>
     if (comp[i] < 0) dfs(i, g, f);</pre>
    return comp;
  template <class G> // convenience function w/o lambda
 vector<int> scc(G& g) {
    return scc(g, [](auto& v) {});
} // namespace SCCTarjan
SCCKosaraju.h
Description: Finds strongly connected components of a directed graph.
Visits/indexes SCCs in topological order.
Usage: scc(graph) returns an array that has the ID
of each node's SCC.
Time: \mathcal{O}(|V| + |E|)
                                                      9b78e7, 29 lines
namespace SCCKosaraju {
 vector<vector<int>> adj, radj;
 vector<int> todo, comp;
  vector<bool> vis;
  void dfs1(int x) {
   vis[x] = 1;
    for (int y : adj[x])
      if (!vis[y]) dfs1(y);
    todo.push_back(x);
  void dfs2(int x, int i) {
   comp[x] = i;
    for (int v : radj[x])
      if (comp[y] == -1) dfs2(y, i);
 vector<int> scc(vector<vector<int>>& _adj) {
    adj = \_adj;
    int time = 0, n = adj.size();
    comp.resize(n, -1), radj.resize(n), vis.resize(n);
    for (int x = 0; x < n; x++)
      for (int y : adj[x]) radj[y].push_back(x);
    for (int x = 0; x < n; x++)
      if (!vis[x]) dfs1(x);
    reverse(todo.begin(), todo.end());
```

LCA.h

for (int x : todo)

}; // namespace SCCKosaraju

return comp;

358d18, 37 lines

Description: Answers lowest common ancestor queries on a rooted tree using RMQ. Works with both directed and undirected adjacency lists.

Time: $\mathcal{O}(1)$ queries with an $\mathcal{O}(n \log n)$ precomp (RMQ).

if (comp[x] == -1) dfs2(x, time++);

```
"../../content/data-structures/RMQ.h" 7e028d, 20 lines
struct LCA {
  int T = 0;
  vector<int> pre, path, times;
  RMQ<int> rmq;
  LCA (vector<vector<int>>& adj, int root = 0)
      : pre(adj.size()), rmq((dfs(root, -1, adj), times)) {}
  void dfs(int u, int p, vector<vector<int>>& adj) {
    pre[u] = T++;
  for (int v : adj[u])
      if (v != p) {
      path.push_back(u), times.push_back(pre[u]);
      dfs(v, u, adj);
    }
}
```

```
}
int lca(int u, int v) {
  if (u == v) return u;
  tie(u, v) = minmax(pre[u], pre[v]);
  return path[rmq.query(u, v - 1)];
};
```

Mathematics (5)

Fraction.h

Description: Struct for representing fractions/rationals. All ops are $O(\log N)$ due to GCD in constructor. Uses cross multiplication alde34, 27 lines

```
template <typename T>
struct O
 Ta, b;
 Q(T p, T q = 1) {
   T g = gcd(p, q);
   a = p / q;
   b = q / g;
   if (b < 0) a = -a, b = -b;
 T gcd(T x, T y) const { return __gcd(x, y); }
 Q operator+(const Q& o) const {
    return {a * o.b + o.a * b, b * o.b};
 O operator-(const O& o) const
    return *this + O(-o.a, o.b);
 Q operator*(const Q& o) const { return {a * o.a, b * o.b}; }
 Q operator/(const Q& o) const { return *this * Q(o.b, o.a); }
 Q recip() const { return {b, a}; }
 int signum() const { return (a > 0) - (a < 0); }</pre>
 return a * o.b < o.a * b;
 friend ostream& operator<<(ostream& cout, const Q& o) {</pre>
   return cout << o.a << "/" << o.b;</pre>
};
```

FractionOverflow.h

Description: Safer struct for representing fractions/rationals. Comparison is 100% overflow safe; other ops are safer but can still overflow. All ops are $O(\log N)$.

```
template <typename T>
struct 00 H
  T a, b;
  QO(T p, T q = 1) {
    T g = gcd(p, q);
    a = p / g;
    b = q / g;
    if (b < 0) a = -a, b = -b;
  T gcd(T x, T y) const { return __gcd(x, y); }
  QO operator+(const QO& o) const {
    T g = gcd(b, o.b), bb = b / g, obb = o.b / g;
    return {a * obb + o.a * bb, o.b * obb};
  OO operator-(const OO& O) const {
    return *this + QO(-o.a, o.b);
  QO operator*(const QO& o) const {
    T g1 = gcd(a, o.b), g2 = gcd(o.a, b);
    return { (a / g1) * (o.a / g2), (b / g2) * (o.b / g1) };
```

```
QO operator/(const QO& o) const {
    return *this * QO(o.b, o.a);
  QO recip() const { return {b, a}; }
  int signum() const { return (a > 0) - (a < 0); }</pre>
  static bool lessThan(T a, T b, T x, T y) {
   if (a / b != x / y) return a / b < x / y;</pre>
   if (x % y == 0) return false;
   if (a % b == 0) return true;
    return lessThan(y, x % y, b, a % b);
  bool operator<(const QO& o) const {
    if (this->signum() != o.signum() || a == 0) return a < o.a;</pre>
    if (a < 0)
      return lessThan(abs(o.a), o.b, abs(a), b);
    else
      return lessThan(a, b, o.a, o.b);
  friend ostream& operator<<(ostream& cout, const QO& o) {</pre>
    return cout << o.a << "/" << o.b;</pre>
};
```

PrimeSieve.h

Description: Prime sieve for generating all primes up to a certain limit. isprime[i] is true iff i is a prime.

Time: $\lim_{n\to\infty} 100'000'000 \approx 0.8 \text{ s.}$ Runs 30% faster if only odd indices are stored.

dc4f55, 14 lines

```
const int MAX PR = 5'000'000;
bitset<MAX_PR> isprime;
vector<int> primeSieve(int lim) {
  isprime.set();
  isprime[0] = isprime[1] = 0;
  for (int i = 4; i < lim; i += 2) isprime[i] = 0;</pre>
  for (int i = 3; i * i < lim; i += 2)</pre>
    if (isprime[i])
      for (int j = i * i; j < lim; j += i * 2) isprime[j] = 0;</pre>
  vector<int> pr;
  for (int i = 2; i < lim; i++)
   if (isprime[i]) pr.push_back(i);
  return pr;
```

PrimeSieveFast.h

Description: Prime sieve for generating all primes smaller than LIM.

Time: LIM=1e9 $\approx 1.5s$

```
a1933d, 23 lines
const int LIM = 1e8;
bitset<LIM> isPrime:
vector<int> primeSieve() {
  const int S = round(sqrt(LIM)), R = LIM / 2;
  vector < int > pr = {2}, sieve(S + 1);
  pr.reserve(int(LIM / log(LIM) * 1.1));
  vector<pair<int, int>> cp;
  for (int i = 3; i <= S; i += 2)
   if (!sieve[i]) {
      cp.push_back(\{i, i * i / 2\});
      for (int j = i * i; j <= S; j += 2 * i) sieve[j] = 1;</pre>
  for (int L = 1; L <= R; L += S) {</pre>
    array<bool, S> block{};
    for (auto& [p, idx] : cp)
      for (int i = idx; i < S + L; idx = (i += p))</pre>
       block[i - L] = 1;
    for (int i = 0; i < min(S, R - L); i++)
      if (!block[i]) pr.push_back((L + i) * 2 + 1);
```

```
for (int i : pr) isPrime[i] = 1;
 return pr;
Miscellaneous (6)
NDimensional Vector.h
                                                    3c0f61, 12 lines
template <int D, typename T>
struct Vec : public vector<Vec<D - 1, T>> {
 static_assert(D >= 1,
                "Vector dimension must be greater than zero!");
 template <typename... Args>
 Vec(int n = 0, Args... args)
   : vector<Vec<D - 1, T>>(n, Vec<D - 1, T>(args...)) {}
template <typename T>
struct Vec<1, T> : public vector<T> {
 Vec(int n = 0, const T& val = T()) : vector<T>(n, val) {}
Submasks.h
                                                    35424b, 3 lines
for (int mask = 0; mask < (1 << n); mask++)
 for (int sub = mask; sub; sub = (sub - 1) & mask)
// do thing
Strings
ZValues.h
                                                    151ee3, 10 lines
vector<int> zValues(string& s) {
 int n = ( int )s.length();
 vector<int> z(n);
 for (int i = 1, l = 0, r = 0; i < n; ++i) {
   if (i \le r) z[i] = min(r - i + 1, z[i - 1]);
   while (i + z[i] < n \&\& s[z[i]] == s[i + z[i]]) ++z[i];
   if (i + z[i] - 1 > r) 1 = i, r = i + z[i] - 1;
 return z;
```