Mark: 5/20 (total score: 7.4/29)



+287/1/52+

Computer Security Foundations 11/11/2022

**Duration: 1H** 

LEIC Midterm

This is a multiple-choice test that will be corrected automatically. Please follow these rules:

- · Mark your answer using only blue and black pen. No pencil or light-coloured pens.
- Check only inside each box and be generous on ink. Erased boxes will not be detected automatically.
- Only the boxes matter for the automatic correction. You may underline text or take notes on the sides.

The test is marked for 20 points. There are 30 questions in total, each with 4 options. Each question is worth 20/30 points. Students can check one or two choices per question. The scoring for each question is as follows:

- One checked correct answer (100%).
- One checked incorrect answer (-20%).
- No checked answers (0%).

- Two checked answers, one correct (50%).
- Two checked answers, none correct (-20%).
- More than two checked answers (-20%).

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code your 9-digit upYYYYXXXXX student number horizontally on the left, and replicate it below. Write also your first and last name below.

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## Group 1 Introduction (4 questions)

Question 1.1 & Pick the incorrect statement or indicate that all are correct.

- A security policy describes how security mechanisms are used to realize a security model.
- A trust model describes which assumptions we can rely on to build security.

X All are correct.

A security goal is stated in terms of avoiding loss of value to an asset.

0/1

	Question 1.2 \$\infty\$ Which of the following is not a common reason for an attacker to compromise a server that answers to requests from a large number of clients/users?
-0.2/1	<ul> <li>As part of a supply-chain attack.</li> <li>☐ Geo-political and strategic motivation.</li> <li>☐ To cause a data breach.</li> </ul>
	Question 1.3 & Which of the following concepts is not associated with the risk management approach to security?
-0.2/1	<ul><li>☐ Cost/benefit analysis.</li><li>☐ Threat analysis.</li><li>☐ Mitigation.</li></ul>
	Question 1.4  Which of the following options is not usually categorised as an exploit?
1/1	<ul> <li>☐ JavaScript performing heap spraying.</li> <li>☐ Text filled into edit box that contains JavaScript.</li> <li>☑ Overflow in an integer computation.</li> <li>☐ Message that triggers an array access out of bounds.</li> </ul>
	Group 2 Software Security (10 questions)
	Question 2.1 $\clubsuit$ Suppose you intend to use Return Oriented Programming to execute the instructions of functions $f_1$ , $f_2$ , $f_3$ , $f_4$ , $f_5$ (in some order) and you have placed their addresses in the stack as illustrated below. Which of the following functions could take parameters and the ROP strategy would still work without any changes to the part of the stack shown in the figure?
1/1	All functions could take parameters. High address $&f_3$ Function $f_3$ .  It is not possible to pass any parameters.  Function $f_2$ .  Low address $&f_3$ $&f_1$ $&f_5$ $&f_4$ $&f_4$
	Question 2.2  Certain format strings can make a program crash (with very high probability). Which of the following printf commands is very likely to crash the program?
0/1	<pre>int n=1; printf("%n%f",&amp;n); int n=0; printf("%f%d",n); int n=0; printf("%s%x",&amp;n); int n=1; printf("%s%d",n);</pre>
	Question 2.3 A canary built from string terminating bytes offers better security than a random canary because:
1/1	<ul> <li>It makes it more difficult to overwrite the stack.</li> <li>☐ It is harder to guess than a random canary.</li> <li>☐ They can be generated much more efficiently.</li> <li>☐ Random bytes can be all 0.</li> </ul>

		Question 2.4 A Recall the buffer overflow lab that you have performed, where a buffer in a local variable receives an untested input. For the exploitation strategy where the injected code (shellcode) is placed at the end of the input bytes, you should:				
		X In the place of the vulnerable function's return address, write another address pointing to a higher place in the stack.				
0.5/1		None of the other options.				
0.5/1	->	In the place of the vulnerable function's return address, write another address pointing to a lower place in the stack.				
		In the place of the vulnerable function's return address, write the address of the buffer.				
		Question 2.5 - One can bypass Address Space Layout Randomization (ASLR) protections by:				
	$\rightarrow$	All of the other choices.				
1/1		Jumping to system code using Return Oriented Programming.				
		Trial and error.				
	Extracting addresses using other vulnerabilities.					
		0				
		Question 2.6 Consider the code below, which introduced a critical vulnerability in openSSH. An attacker that controls the value of nresp can cause a memory management problem because:				
-0.2/1		It can force the if statement not to execute. nresp = packet_get_int();  It can lead to not allocation enough memory. if (nresp > 0) {     response = xmalloc(nresp*sizeof(char*));     for (i = 0; i < nresp; i++)         response[i] = packet_get_string(NULL); }  None of the other choices.				
		Question 2.7 $\clubsuit$ Recall what you have studied about the classical configuration of the stack region managed by a function $f$ that is called by a function $g$ . Indicate which of the following options is correct for the positioning of the following data pieces: ① local variables of $f$ , ② frame pointer of $g$ , ③ parameters passed by $g$ , ④ return address back to $g$ .				
		$\square$ $(4) \Rightarrow A, (3) \Rightarrow B, (1) \Rightarrow C, (2) \Rightarrow D$ High address				
1/1		$(3) \Rightarrow A, (4) \Rightarrow B, (2) \Rightarrow C, (1) \Rightarrow D$				
		$\square \square \rightarrow \Lambda, \square \rightarrow B, \square \rightarrow C, \bigoplus \rightarrow D$				
		Question 2.8 Recall the buffer overflow lab. You explored a buffer overflow vulnerability by creating a malicious input that overwrites the return address of the vulnerable function with an address within your buffer where you stored <i>shellcode</i> to be executed. Which security technique had to be disabled during compilation to allow you to overwrite the return address of the vulnerable function?				
0.04:		None of the other choices Stack Canary				
-0.2/1		Data Execution Prevention Address Space Layout Randomisation				

	Question 2.9 . Information-flow analysis	sis encompasses a class of program analysis technically the illustrates an example of:	ques				
0.2/1	Control-flow integrity.  Dynamic taint analysis.	<pre>int printf(untainted char *fmt,); tainted char *fgets();</pre>					
0.271	<ul><li>☐ Constant-time analysis.</li><li>☒ Static taint analysis.</li></ul>	<pre>tainted char *name = fgets(,network_fd); printf(name); // FAIL: tainted not untainted</pre>					
	Group 3 Systems Security	(10 questions)					
	Question 3.1 & Remember the environr	nent variable and setuid program lab. The Linux permission bits with each entry. For a file with the					
)/1		o, others) in the permission bits, has the executable	e bit				
	only by the owner or by a user belongi Only by a user belonging to the group.	The World of the Control of World Control of the Co					
		system is a particular distribution built on top of Lin	nux.				
	It implements a form of Mandatory Admissions.	ccess Control, so that no application can change its	per-				
)/1	There is no root user, to prevent applic	cations from escalating privileges.					
<i>)</i> // I	It isolates different applications by registering each application with a different UNIX user.						
	It enforces Manifest Permissions per application, to restrict its system capabilities.						
		the address space of the Linux kernel is not complet to kernel mapping. Which of the following is correct					
	Part of each user process's memory is a of the address translation tables.	mapped into the kernel memory space, to reduce the	size				
)/1	Part of the kernel memory space is map up system calls.	Part of the kernel memory space is mapped into the memory space of each user process, to speed					
	Since the discovery of speculative execution attacks such as Meltdown and Spectre, Linux no longer uses kernel mapping, as it violates the principle of separation of privilege.						
	Part of the kernel memory space is map the number of cache misses.	ped into the memory space of each user process, to rec	duce				
	Question 3.4 . Which of the above is no classes?	t a systems security principle that we have studied in	ı the				
. –	Separation of privilege	Closed design	- Land oral				
).5/1	Compromise recording	Economy of mechanism	tracks which				

Question 3.5 & Which fundamental mechanism does a UNIX operating system have in place to enforce isolation? System call interposition, by preventing different user processes to communicate via system calls. Software fault isolation, by ensuring that the control-flow of user processes never accesses invalid virtual memory regions. 0.5/1Software fault isolation, by virtualising the address space of different user processes, monitoring all virtual address translations. System call interposition, by virtualising the address space of different user processes, monitoring all virtual address translations. Question 3.6 \( \text{Many modern operating systems require system support for an hardware component called a Trusted Platform Module (TPM). Which is the main rationale for that requirement? To store user's cookies when browsing the web. To guarantee that the user only installs official software. -0.2/1To protect the user login process. To prevent malicious bootloaders from compromising the operating system's boot process. Question 3.7 . The UNIX filesystem can be seen as an instance of: Access control lists, since only the owner of each file (or root) may read, write or execute it. Attribute-based access control, since users may change their group membership without the need to change file permissions. 0.5/1Capability lists, since each file enumerates the permissions for all the users who may read, write or execute it. Role-based access control, since each file has an access control list for three fixed roles (owner, group, others) and the association between users and groups may be modified independently. Question 3.8 In the environment variable and setuid program lab, we have experimented with environment variables and setuid programs. Which of the following sentences is not true? If the shell detects that it is being run under a setuid process, it may drop its privilege. When a process forks a child process, it passes on its environment, excluding some critical variables if the fork is a system call. 0.5/1The system function allows calling a shell function within a program with the program's environment. When the shell forks a child process, it passes on its environment, excluding some critical variables if it has a different effective user id.

Ouestion 3.9 . Under Linux, Docker crucially relies on seccomp-bpf to confine containers. Select the incorrect choice or mark that all choices are correct. All choices are correct. If an attacker has access to certain system calls inside a container, it can exploit Docker to acquire root in the host OS. 0/1Using the default Docker seccomp-bpf filters, an attacker that acquires root in the container does not directly acquire root in the host OS. Users can manually configure seccomp-bpf filters to block various system calls inside the con-Question 3.10 . Virtual machines are commonly used as a security mechanism. Which of the following sentences is not true? An advantage of virtual machines is that software may be transparently executed as in a nonvirtualised operating system. A disadvantage of virtual machines is that malware may exploit the virtualisation layer to remain -0.2/1undetected. A disadvantage of virtual machines is that malware may detect that is is being virtualised. An advantage of virtual machines is that acquiring root in the virtual machine does not grant root in a host operating system. Web Security (6 questions) Group 4 Which sentence best characterises a Cross-Site Request Forgery (CSRF) attack? Question 4.1 & "When a malicious origin can send requests to a server in another origin ...": ... and measure side-channels such as re-... and read its response. sponse time. -0.2/1... where the user is already logged in. ... and trigger a side-effect on the server. Question 4.2 . Cookies may be used for various purposes. Name the incorrect one below. Tracking user activity. Personalising user experience. 1/1 Managing user sessions. Encrypting user communication. Question 4.3 & Consider the following SQL query, written in some server-side library, that is vulnerable to SQL injection. Which of the following statements is true? uName = getRequestString("username"); uPass = getRequestString("userpassword"); sql = 'SELECT\_\*\_FROM\_Users\_WHERE\_Name\_="' + uName + '"\_AND\_Pass\_="' + uPass + '"' A SQL injection attack that bypasses authentication is only possible because the clause for the Name field appears before the clause for the Pass field. A SQL injection attack will only be able to read existing data from the USERS table. 0.5/1A SQL injection attack would not be possible if the clauses for the Name and Pass fields did not enclose strings with double quotes ("). A SQL injection attack may be able to delete the USERS table.

		Question 4.4 A classical protection against Cross-S adopt a Content Security Policy (CSP). Which of the following		
0/1		<ul> <li>DOM-based XSS attacks cannot be prevented with a when processing a user request.</li> <li>Stored XSS attacks cannot be prevented with a CSP stored in the server.</li> <li>Preventing reflected XSS attacks requires both CSP XSS attacks using inline scripts are blocked by a decoupled of the company of</li></ul>	because the malicion and SRI (Subresour fault CSP.	ous payloads are alread
0/1	_	①=Frames,②=DOM,③=Images,④=Sub-frames ① ①=Pages,②=HTML,③=JavaScript,④=Popups ② ①=DOM,②=Cookies,③=iFrames,④=Fetch 【 ①=Pages,②=Cookies,③=HTTP,④=Frames	Systems Security Processes Files Sockets Sub-processes	Web Security ① ② ③ ④ ④
0/1	-3	Question 4.6 ♣ According the Same-Origin Policy (SO HTML code in a page from origin A can send a PO JavaScript code in a frame from origin A can exchan JavaScript code in a page from origin A can inspect HTML code in a page from origin A can send and refrom origin B.	ST request to a page nge data with a fram the HTML code of	e from origin B. ne from origin B. a frame from origin B.