

Chapel: Base Language



Goals of this Talk



- Help you understand code in subsequent slide decks
- Give you the basic skills to program in Chapel
- Provide a survey of Chapel's base language features
- Impart an appreciation for the base language design

Note: There is more in this slide deck than we will be able to cover, so consider it a reference and overview





"Hello World" in Chapel: Two Versions

Fast prototyping

```
writeln("Hello, world!");
```

"Production-grade"

```
module Hello {
  proc main() {
    writeln("Hello, world!");
  }
}
```



Characteristics of Chapel



Design points

- Identifying parallelism & locality is user's job, not compiler's
- No unexpected compiler-inserted array temporaries
- No pointers and limited opportunities for aliasing
- Intentionally not an extension of an existing language

Chapel Influences



C, Modula: basic syntax

ZPL, HPF: data parallelism, index sets, distributed arrays

CRAY MTA C/Fortran: task parallelism, synchronization

CLU (see also Ruby, Python, C#): iterators

Scala (see also ML, Matlab, Perl, Python, C#): type inference

Java, C#: OOP, type safety

C++: generic programming/templates (but with a different syntax)



Outline



- Introductory Notes
- Elementary Concepts
 - Lexical structure
 - Types, variables, and constants
 - Operators and Assignments
 - Compound Statements
 - Input and output
- Data Types and Control Flow
- Program Structure



Lexical Structure



Comments

```
/* standard
    C style
    multi-line */
// standard C++ style single-line
```

- Identifiers:
 - Composed of A-Z, a-z, _, \$, 0-9
 - Cannot start with 0-9
- Case-sensitive







Туре	Description	Default Value	Currently-Supported Bit Widths	Default Bit Width
bool	logical value	false	8, 16, 32, 64	impl. dep.
int	signed integer	0	8, 16, 32, 64	64
uint	unsigned integer	0	8, 16, 32, 64	64
real	real floating point	0.0	32, 64	64
imag	imaginary floating point	0.0i	32, 64	64
complex	complex floating points	0.0 + 0.0i	64, 128	128
string	character string	\\ //	N/A	N/A

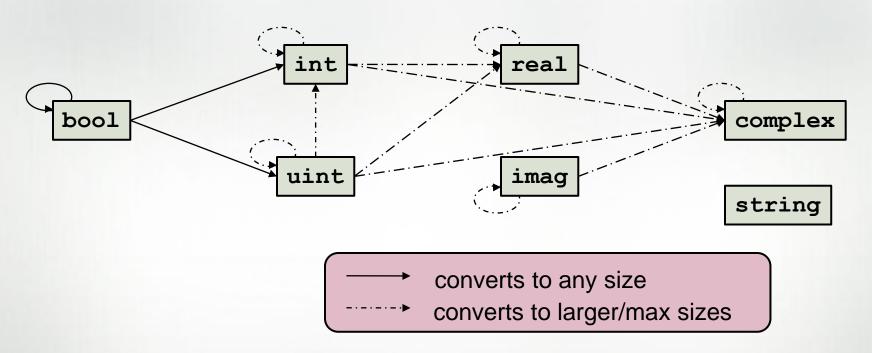
Syntax

```
primitive-type:
  type-name [( bit-width )]
```





Implicit Type Conversions (Coercions)



- Notes:
 - reals do not implicitly convert to ints as in C
 - ints and uints don't interconvert as handily as in C



Type Aliases and Casts



Basic Syntax

```
type-alias-declaration:
   type identifier = type-expr;

cast-expr:
   expr : type-expr
```

- Semantics
 - type aliases are simply symbolic names for types
 - casts are supported between any primitive types
- Examples

```
type elementType = complex(64);

5:int(8) // store value as int(8) rather than int
"54":int // convert string to an int
249:elementType // convert int to complex(64)
```





Variables, Constants, and Parameters

Basic syntax

```
declaration:
  var  identifier [: type] [= init-expr];
  const identifier [: type] [= init-expr];
  param identifier [: type] [= init-expr];
```

Semantics

- var/const: execution-time variable/constant
- param: compile-time constant
- No *init-expr* \Rightarrow initial value is the type's default
- No *type* \Rightarrow type is taken from *init-expr*



Config Declarations



Syntax

```
config-declaration:
  config type-alias-declaration
  config declaration
```

- Semantics
 - Like normal, but supports command-line overrides
 - Must be declared at module/file scope
- Examples

```
config param intSize = 32;
config type elementType = real(32);
config const epsilon = 0.01:elementType;
config var start = 1:int(intSize);
```

```
% chpl myProgram.chpl -sintSize=64 -selementType=real
% a.out --start=2 --epsilon=0.00001
```





Basic Operators and Precedence

Operator	Description	Associativity	Overloadable
:	cast	left	no
**	exponentiation	right	yes
! ~	logical and bitwise negation	right	yes
* / %	multiplication, division and modulus	left	yes
unary + -	positive identity and negation	right	yes
+ -	addition and subtraction	left	yes
<< >>	shift left and shift right	left	yes
<= >= < >	ordered comparison	left	yes
== !=	equality comparison	left	yes
&	bitwise/logical and	left	yes
^	bitwise/logical xor	left	yes
1	bitwise/logical or	left	yes
& &	short-circuiting logical and	left	via isTrue
11	short-circuiting logical or	left	via isTrue







Kind	Description
=	simple assignment
+= -= *= /= %= **= &= = ^= &&= = <<= >>=	compound assignment (e.g., $x += y$; is equivalent to $x = x + y$;)
<=>	swap assignment

Note: assignments are only supported at the statement level







Syntax

```
compound-stmt:
{ stmt-list }
```

- Semantics
 - As in C, permits a series of statements to be used in place of a single statement
- Example

```
writeln("Starting a compound statement");
x += 1;
writeln("Ending the compound statement");
}
```



Console Input/Output



Output

- write(expr-list): writes the argument expressions
- writeln(...) variant: writes a linefeed after the arguments

Input

- read(expr-list): reads values into the argument expressions
- read(type-list): reads values of given types, returns as tuple
- readln(...) variant: same, but reads through next linefeed

• Example:

```
var first, last: string;
write("what is your name? ");
read(first);
last = read(string);
writeln("Hi ", first, " ", last);
```

What is your name?

Chapel User

Hi Chapel User

I/O to files and strings also supported



Outline



- Introductory Notes
- Elementary Concepts
- Data Types and Control Flow
 - Tuples
 - Ranges
 - Arrays
 - For loops
 - Other control flow
- Program Structure



Tuples



Syntax

```
heterogeneous-tuple-type:
  ( type, type-list )
homogenous-tuple-type:
  param-int-expr * type

tuple-expr:
  ( expr, expr-list )
```

Purpose

- supports lightweight grouping of values

 (e.g., when passing or returning procedure arguments)
- multidimensional arrays use tuple indices

```
var coord: (int, int, int) = (1, 2, 3);
var coordCopy: 3*int = coord;
var (i1, i2, i3) = coord;
var triple: (int, string, real) = (7, "eight", 9.0);
```



Range Values



Syntax

```
range-expr:
[low] .. [high]
```

Semantics

Regular sequence of integers

```
low <= high: low, low+1, low+2, ..., high
```

low > high: degenerate (an empty range)

low or high unspecified: unbounded in that direction



Range Operators



Syntax

Semantics

- by: strides range; negative stride ⇒ start from high
- #: selects initial count elements of range
- align: specifies the alignment of a strided range
- [] or (): intersects the two ranges

```
1..6 by 2 // 1, 3, 5
1..6 by -1 // 6, 5, 4, ..., 1
1..6 #4 // 1, 2, 3, 4
1..6[3..] // 3, 4, 5, 6
```

```
1.. by 2 // 1, 3, 5, ...
1.. by 2 #3 // 1, 3, 5
1.. by 2 align 2 // 2, 4, ...
1.. #3 by 2 // 1, 3
0..#n // 0, ..., n-1
```



Array Types



Syntax

```
array-type:
[ index-set-expr ] elt-type
```

- Semantics
 - Stores an element of elt-type for each index
 - Array values expressed using square brackets

Examples

```
var A: [1..3] int = [5, 3, 9], // 3-element array of ints
B: [1..3, 1..5] real, // 2D array of reals
C: [1..3][1..5] real; // array of arrays of reals
```

Much more on arrays in data parallelism section later...



For Loops



Syntax

```
for-loop:
   for index-expr in iteratable-expr { stmt-list }
```

- Semantics
 - Executes loop body serially, once per loop iteration
 - Declares new variables for identifiers in index-expr
 - type and const-ness determined by iteratable-expr
 - iteratable-expr could be a range, array, or iterator
- Examples

```
var A: [1..3] string = [" DO", " RE", " MI"];

for i in 1..3 { write(A(i)); } // DO RE MI
for a in A { a += "LA"; } write(A); // DOLA RELA MILA
```



Zipper Iteration



Syntax

```
zipper-for-loop:
  for index-expr in zip( iteratable-exprs ) { stmt-list }
```

- Semantics
 - Zipper iteration is over all yielded indices pair-wise
- Example

```
var A: [0..9] real;
for (a,i,j) in zip(A, 1..10, 2..20 by 2) do
    a = j + i/10.0;
writeln(A);
```

2.1 4.2 6.3 8.4 10.5 12.6 14.7 16.8 18.9 21.0







Conditional statements

```
if cond { computeA(); } else { computeB(); }
```

While loops

```
while cond {
  compute();
}
```

```
do {
   compute();
} while cond;
```

Select statements

```
select key {
  when value1 { compute1(); }
  when value2 { compute2(); }
  otherwise { compute3(); }
}
```

Note: Chapel also has expression-level conditionals and for loops





Control Flow: Braces vs. Keywords

Most control flow supports keyword-based forms for single-statement versions

Conditional statements

```
if cond then computeA(); else computeB();
```

While loops

```
while cond do
  compute();
```

For loops

```
for indices in iteratable-expr do
  compute();
```

Select statements

```
select key {
  when value1 do compute1();
  when value2 do compute2();
  otherwise do compute3();
}
```



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- Program Structure
 - Procedures and iterators
 - Modules and main()
 - Records and classes
 - Generics
 - Other basic language features



Procedures, by example



Example to compute the area of a circle

```
proc area(radius: real): real {
  return 3.14 * radius**2;
}
writeln(area(2.0)); // 12.56
```

```
proc area(radius) {
  return 3.14 * radius**2;
}
```

Argument and return types can be omitted

Example of argument default values, naming



Iterators



- Iterator: a procedure that generates values/variables
 - Used to drive loops or populate data structures
 - Like a procedure, but yields values back to invocation site
 - Control flow logically continues from that point

```
iter fibonacci(n) {
  var current = 0,
    next = 1;
  for 1..n {
    yield current;
    current += next;
    current <=> next;
}
```

```
for f in fibonacci(7) do
    writeln(f);

0
1
2
3
5
8
```



Argument and Return Intents



- Arguments can optionally be given intents
 - (blank): varies with type; follows principle of least surprise
 - most types: const
 - arrays, domains, sync vars: passed by reference
 - const: disallows modification of the formal
 - in: copies actual into formal at start; permits modifications
 - out: copies formal into actual at procedure return
 - inout: does both of the above
 - ref: pass by reference
 - param/type: formal must be a param/type (evaluated at compile-time)
- Return types can also have intents
 - (blank)/const: cannot be modified (without assigning to a variable)
 - var: permits modification back at the callsite
 - type: returns a type (evaluted at compile-time)
 - param: returns a param value (evaluated at compile-time)



Modules



Syntax

```
module-def:
  module identifier { code }

module-use:
  use module-identifier;
```

Semantics

- all Chapel code is stored in modules
- use-ing a module makes its symbols visible in that scope
- module-level statements are executed at program startup
 - typically used to initialize the module
- for convenience, a file containing code outside of a module declaration creates a module with the file's name



Program Entry Point: main()



- Semantics
 - Chapel programs start by:
 - initializing all modules
 - executing main(), if it exists

```
M1.chpl:
use M2;
writeln("Initializing M1");
proc main() { writeln("Running M1"); }
```

```
M2.chpl:
module M2 {
  writeln("Initializing M2");
}
```

```
% chpl M1.chpl M2.chpl
% ./a.out
Initializing M2
Initializing M1
Running M1
```



Revisiting "Hello World"



Fast prototyping

```
hello.chpl
writeln("Hello, world!"); ==
```

"Production-grade"

```
module HelloWorld {
   proc main() {
      writeln("Hello, world!");
   }
}
```

```
module hello {
  writeln("Hello, world!");
}
```

Module-level code is executed during module initialization

main() executed when program begins running



Records and Classes



- Chapel's struct/object types
 - Contain variable definitions (fields)
 - Contain procedure & iterator definitions (methods)
 - Records: value-based (e.g., assignment copies fields)
 - Classes: reference-based (e.g., assignment aliases object)
 - Record : Class :: C++ struct : Java class

```
record circle {
  var radius: real;
  proc area() {
    return pi*radius**2;
  }
}
```



Records and Classes



- Chapel's struct/object types
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 - Records: value-based (e.g., assignment copies fields)
 - Classes: reference-based (e.g., assignment aliases object)
 - Record : Class :: C++ struct : Java class

```
class circle {
  var radius: real;
  proc area() {
    return pi*radius**2;
  }
}
```

```
var c1, c2: circle;
c1 = new c1(radius=1.0);
c2 = c1; // aliases c1's circle
c1.radius = 5.0;
writeln(c2.radius); // 5.0
delete c1; // users delete classes
```







Methods without arguments need not use parenthesis

```
proc circle.circumference {
  return 2* pi * radius;
}
writeln(c1.area(), " ", c1.circumference);
```

Methods can be defined for any type

```
proc int.square() {
   return this**2;
}
writeln(5.square());
```



Generic Procedures



Generic procedures can be defined using type and param arguments:

```
proc foo(type t, x: t) { ... }
proc bar(param bitWidth, x: int(bitWidth)) { ... }
```

Or by simply omitting an argument type (or type part):

```
proc goo(x, y) { ... }
proc sort(A: []) { ... }
```

Generic procedures are instantiated for each unique argument signature:



Generic Objects



Generic objects can be defined using type and param fields:

```
class Table { param size: int; var data: size*int; }
class Matrix { type eltType; ... }
```

Or by simply eliding a field type (or type part):

```
record Triple { var x, y, z; }
```

Generic objects are instantiated for each unique type signature:

```
// instantiates Table, storing data as a 10-tuple
var myT: Table(10);
// instantiates Triple as x:int, y:int, z:real
var my3: Triple(int, int, real) = new Triple(1, 2, 3.0);
```



Other Base Language Features not covered today



- Enumerated types
- Unions
- Type select statements, argument type queries
- Parenthesis-less functions/methods
- Procedure dispatch constraints ("where" clauses)
- Compile-time features for meta-programming
 - type/param procedures
 - folded conditionals
 - unrolled for loops
 - user-defined compile-time warnings and errors



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Status: Base Language Features

- Most features working well
- Performance is currently suboptimal in some cases
- Some semantic checks are incomplete
 - e.g., constness-checking for members, arrays
- Error messages could use improvement at times
- OOP features are limited in certain respects
 - generic classes w/ subclassing, user constructors
- Memory for strings is currently leaked



Future Directions



- Error handling/Exceptions
- Fixed-length strings
- Concepts/Interfaces (joint work with CU Boulder)
- Improved namespace control
 - private fields/methods in classes and records
 - module symbol privacy, filtering, renaming
- Interoperability with other languages (joint with LLNL)



Questions?



- Introductory Notes
 - Characteristics
 - Influences
- Elementary Concepts
 - Lexical structure
 - Types, variables, and constants
 - Operators and assignments
 - Compound Statements
 - Input and output

- Data Types and Control Flow
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