

Hierarchical Locales: Exposing the Node Architecture in Chapel

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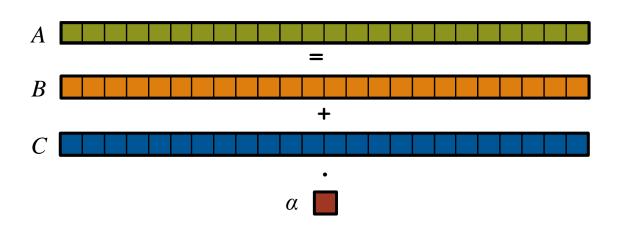
STREAM Triad: a trivial parallel computation

CRAY

Given: *m*-element vectors *A*, *B*, *C*

Compute: $\forall i \in 1..m, A_i = B_i + \alpha \cdot C_i$

In pictures:





STREAM Triad: MPI



```
#include <hpcc.h>
```

```
static int VectorSize;
static double *a, *b, *c;
int HPCC StarStream(HPCC Params *params) {
  int myRank, commSize;
  int rv, errCount;
 MPI Comm comm = MPI COMM WORLD;
 MPI Comm size ( comm, &commSize );
 MPI Comm rank ( comm, &myRank );
 rv = HPCC Stream( params, 0 == myRank);
 MPI Reduce ( &rv, &errCount, 1, MPI INT, MPI SUM,
   0, comm );
  return errCount;
int HPCC Stream(HPCC Params *params, int doIO) {
  register int j;
 double scalar;
 VectorSize = HPCC LocalVectorSize( params, 3,
   sizeof(double), 0 );
  a = HPCC XMALLOC( double, VectorSize );
 b = HPCC XMALLOC( double, VectorSize );
  c = HPCC XMALLOC( double, VectorSize );
```

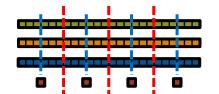


```
if (!a || !b || !c) {
  if (c) HPCC free(c);
  if (b) HPCC free(b);
  if (a) HPCC free(a);
  if (doIO) {
    fprintf( outFile, "Failed to allocate memory
  (%d).\n", VectorSize );
    fclose( outFile );
  return 1;
for (j=0; j<VectorSize; j++) {</pre>
 b[i] = 2.0;
  c[j] = 0.0;
scalar = 3.0;
for (j=0; j<VectorSize; j++)</pre>
  a[i] = b[i] + scalar*c[i];
HPCC free(c);
HPCC free (b);
```

HPCC free(a);



STREAM Triad: MPI+OpenMP





MPI + OpenMP

```
#include <hpcc.h>
                                                        if (!a || !b || !c) {
#ifdef OPENMP
                                                          if (c) HPCC free(c);
#include <omp.h>
                                                          if (b) HPCC free(b);
#endif
                                                          if (a) HPCC free(a);
                                                          if (doIO) {
static int VectorSize;
static double *a, *b, *c;
                                                            fprintf( outFile, "Failed to allocate memory
                                                          (%d).\n", VectorSize );
int HPCC StarStream(HPCC Params *params) {
                                                            fclose( outFile );
  int myRank, commSize;
  int rv, errCount;
                                                          return 1;
  MPI Comm comm = MPI COMM WORLD;
 MPI Comm size ( comm, &commSize );
                                                      #ifdef OPENMP
  MPI Comm rank ( comm, &myRank );
                                                      #pragma omp parallel for
                                                      #endif
  rv = HPCC Stream( params, 0 == myRank);
                                                        for (j=0; j<VectorSize; j++) {</pre>
  MPI Reduce ( &rv, &errCount, 1, MPI INT, MPI SUM,
                                                         b[j] = 2.0;
   0, comm );
                                                          c[j] = 0.0;
  return errCount;
                                                        scalar = 3.0;
int HPCC Stream(HPCC Params *params, int doIO) {
                                                      #ifdef OPENMP
  register int j;
                                                      #pragma omp parallel for
  double scalar;
                                                      #endif
                                                        for (j=0; j<VectorSize; j++)</pre>
  VectorSize = HPCC LocalVectorSize( params, 3,
                                                          a[i] = b[i] + scalar*c[i];
   sizeof(double), 0 );
                                                        HPCC free(c);
  a = HPCC XMALLOC( double, VectorSize );
                                                        HPCC free (b);
  b = HPCC XMALLOC( double, VectorSize );
                                                        HPCC free(a);
  c = HPCC XMALLOC( double, VectorSize );
```



STREAM Triad: MPI+OpenMP vs. CUDA

MPI + OpenMP

```
#ifdef OPENMP
#include <omp.h>
#endif
static int VectorSize;
static double *a, *b, *c;
int HPCC StarStream(HPCC Params *params) {
 int myRank, commSize;
 int rv, errCount;
 MPI Comm comm = MPI COMM WORLD;
 MPI Comm size ( comm, &commSize );
 MPI Comm rank ( comm, &myRank );
 rv = HPCC Stream( params, 0 == myRank);
 MPI Reduce ( &rv, &errCount, 1, MPI INT, MPI SUM, 0, comm );
  return errCount;
int HPCC Stream (HPCC Params *params, int doIO) {
 register int j;
 double scalar;
 VectorSize = HPCC LocalVectorSize( params, 3, sizeof(double), 0 );
 a = HPCC XMALLOC( double, VectorSize );
 b = HPCC XMALLOC( double, VectorSize );
 c = HPCC XMALLOC( double, VectorSize );
 if (!a || !b || !c) {
   if (c) HPCC free(c);
   if (b) HPCC free(b);
   if (a) HPCC_free(a);
      fprintf( outFile, "Failed to allocate memory (%d).\n", VectorSize );
     fclose ( outFile );
   return 1;
#ifdef OPENMP
#pragma omp parallel for
 for (j=0; j<VectorSize; j++) {
  b[j] = 2.0;
   c[j] = 0.0;
 scalar = 3.0;
#ifdef OPENMP
#pragma omp parallel for
#endif
 for (j=0; j<VectorSize; j++)
   a[j] = b[j] + scalar*c[j];
 HPCC free(c);
 HPCC free (b);
 HPCC free(a);
  return 0;
```

CUDA

```
#define N
                2000000
int main() {
 float *d a, *d b, *d c;
  float scalar:
  cudaMalloc((void**)&d a, sizeof(float)*N);
 cudaMalloc((void**)&d b, sizeof(float)*N);
  cudaMalloc((void**)&d c, sizeof(float)*N);
  dim3 dimBlock(128);
  dim3 dimGrid(N/dimBlock.x );
  if ( N % dimBlock.x != 0 ) dimGrid
 set array<<<dimGrid,dimBlock>>>(d b, .5f, N);
 set array<<<dimGrid,dimBlock>>>(d c, .5f, N);
  scalar=3.0f:
 STREAM Triad<<<dimGrid,dimBlock>>>(d b, d c, d_a, scalar, N);
  cudaThreadSynchronize();
  cudaFree(d a);
  cudaFree(d b);
 cudaFree(d c);
 global void set array(float *a, float value, int len) {
  int idx = threadIdx.x + blockIdx.x * blockDim.x;
  if (idx < len) a[idx] = value;
 global void STREAM Triad( float *a, float *b, float *c,
                              float scalar, int len) {
  int idx = threadIdx.x + blockIdx.x * blockDim.x;
  if (idx < len) c[idx] = a[idx] + scalar*b[idx];
```

STREAM Triad: Chapel

```
#ifdef OPENMP
#include <omp.h>
#endif
                     config const m = 1000,
                                       alpha = 3.0;
                                                                                       the special
                     const ProblemSpace = {1..m} dmapped
MPI Comm rank ( comm, &myRank );
                                                                                       sauce
                    var A, B, C: [ProblemSpace] real;
                    B = 2.0;
 double scalar;
 VectorSize = HPCC LocalVectorSize(
                    C = 3.0;
 a = HPCC XMALLOC ( double, VectorSi
                    A = B + alpha * C;
                                                                                   c, da, scalar, N);
 -------
```

<u>Philosophy:</u> Good language design can tease details of locality and parallelism away from an algorithm, permitting the compiler, runtime, applied scientist, and HPC expert to each focus on their strengths.



#pragi

Outline

- ✓ Motivation
- Chapel Background
- Exposing the node architecture: locale models
- Project Status and Next Steps



What is Chapel?



- An emerging parallel programming language
 - Design and development led by Cray Inc.
 - in collaboration with academia, labs, industry
 - Initiated under the DARPA HPCS program
- Overall goal: Improve programmer productivity
 - Improve the programmability of parallel computers
 - Match or beat the performance of current programming models
 - Support better portability than current programming models
 - Improve the robustness of parallel codes
- A work-in-progress



Chapel's Implementation

- Being developed as open source at SourceForge
- Licensed as BSD software
- Target Architectures:
 - multicore desktops and laptops
 - Cray architectures
 - commodity clusters
 - systems from other vendors
 - in-progress: CPU+accelerator hybrids, ...



Chapel Design Goal

A = B + alpha * C;

Have it perform like this...

```
#ifdef OPENMP
#include <omp.h>
#endif
static int VectorSize;
static double *a, *b, *c;
int HPCC StarStream (HPCC Params *params) {
 int myRank, commSize;
 int rv, errCount;
 MPI Comm comm = MPI COMM WORLD;
 MPI Comm size ( comm, &commSize );
 MPI Comm rank ( comm, &myRank );
 rv = HPCC_Stream( params, 0 == myRank);
 MPI Reduce ( &rv, &errCount, 1, MPI INT, MPI SUM, 0, comm );
 return errCount;
int HPCC Stream(HPCC_Params *params, int doIO) {
 register int j;
 double scalar;
 VectorSize = HPCC LocalVectorSize( params, 3, sizeof(double), 0 );
 a = HPCC XMALLOC( double, VectorSize );
 b = HPCC XMALLOC ( double, VectorSize );
 c = HPCC XMALLOC( double, VectorSize );
 if (!a || !b || !c) {
   if (c) HPCC free(c);
   if (b) HPCC free(b);
   if (a) HPCC free(a);
     fprintf( outFile, "Failed to allocate memory (%d).\n", VectorSize );
      fclose(outFile);
```

...on any architecture available

```
#pragma omp parallel for
#endif
for (j=0; j<VectorSize; j++) {
    b(j] = 2.0;
    c(j] = 0.0;
}
scalar = 3.0;
#ifdef _OPENMP
#pragma omp parallel for
#endif
for (j=0; j<VectorSize; j++)
    a(j] = b(j]+scalar*c(j);
...</pre>
```



How can this be achieved?



Don't expose low-level mechanisms

Chapel provides syntax to express locality and parallelism, which are portable to any architecture

Don't expect the compiler to do everything

Chapel's locality and parallelism mechanisms are implemented in Chapel modules not our compiler

Don't keep the implementation too close

Chapel is an open-source project with contributors from around the world



Outline

- ✓ Motivation
- √ Chapel Background
- > Exposing the node architecture: locale models
- Project Status and Next Steps



Exposing the node architecture: Locale models



Locale models

- The locale type
- Locality/affinity control
- Task parallelism
- Example: The NUMA locale model
- STREAM Triad revisited
 - Domain maps
 - Leader-follower iterators



The Locale Model



 Every Chapel program employs a locale model to describe the system architecture

Locale models define:

- Abstract node architecture (the *locale* type)
- Memory allocation
- Task scheduling policies
- Locality/affinity control

Locale models are written in Chapel

- Class and procedural interface for the compiler
- Utilizes runtime interfaces (e.g., communication, tasking/threading)
- Programmer role: Architecture specialist



The locale type



The locale type

- an abstract unit of architecture
 - Can have memory, processing units, etc.
- can be nested (hierarchical)
- typically a compute node (multicore processor or SMP) or your laptop
- Chapel programs run on one or more locales



Defining Locales



Specify # of locales when running Chapel programs

```
% a.out --numLocales=8 % a.out -nl 8
```

Chapel provides built-in locale variables

```
config const numLocales: int = ...;
const Locales: [0..#numLocales] locale = ...;

Locales    Lo L1 L2 L3 L4 L5 L6 L7
```

User's main() begins executing on locale 0



Locale Operations



Locale methods support queries about the target system:

```
proc locale.numCores { ... }
proc locale.id { ... }
proc locale.name { ... }
```

On statements support placement of computations:

```
writeln("on locale 0");

on Locales[1] do
  writeln("now on locale 1");

writeln("on locale 0 again");
```



Locale Operations



Locales can be nested:

```
proc locale.getChild(...) { ... }
proc locale.getChildCount { ... }
```

- Children (sublocales) are locales
- On statements support placement of computations on sublocales:

```
writeln("on locale 0");

on here.getChild(1) do
   writeln("now on locale 0, sublocale 1");

writeln("on locale 0 again");
```



On statements: Rewriting



Conceptually, the Chapel compiler translates:

```
on here.getChild(1) do
  writeln("now on locale 0, sublocale 1");
```

into:

```
proc on_fn() {
  writeln("now on locale 0, sublocale 1");
}

chpl_executeOn(here.getChild(1), on_fn, ...);
```

chpl_executeOn() is defined by the Locale Model



Task parallelism

- Tasks are the basic unit of parallelism
- Tasks can be created by the following Chapel statements:
 - begin
 - cobegin
 - coforall



Task parallelism: begin statements



```
// create a fire-and-forget task for a statement
begin writeln("hello world");
writeln("good bye");
```

Possible outputs:

```
hello world
good bye
```

```
good bye
hello world
```





```
// create a task per child statement
cobegin {
  producer(1);
  producer(2);
  consumer(1);
} // implicit join of the three tasks here
```



Task parallelism: coforall loops



```
// create a task per iteration
coforall t in 0..#numTasks {
  writeln("Hello from task ", t, " of ", numTasks);
} // implicit join of the numTasks tasks here
writeln("All tasks done");
```

Sample output:

```
Hello from task 2 of 4
Hello from task 0 of 4
Hello from task 3 of 4
Hello from task 1 of 4
All tasks done
```



Begin statement: Rewriting



Conceptually, the Chapel compiler translates:

```
begin writeln("hello world");
```

into:

```
proc begin_fn() {
  writeln("hello world");
}
chpl_taskListAddBegin(begin_fn, ...);
```

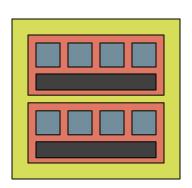
chpl_taskListAddBegin() is defined by the Locale Model



Example: The NUMA locale model

CRAY

- The first interesting Locale Model prototype
- n NUMA domains per locale (sublocale)
 - m cores per NUMA domain
- Memory divided between NUMA domains
 - shared by both NUMA domains



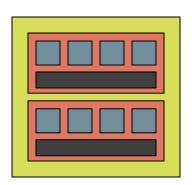


Example: The NUMA locale model



```
class LocaleModel: AbstractLocaleModel {
   var numSublocales: int;
   var childLocales: [0..#numSublocales] NumaDomain;
   ...
   proc getChild(i) return childLocales[i];
   ...
}
```

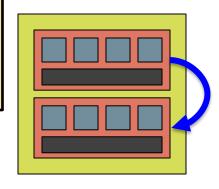
```
class NumaDomain: AbstractLocaleModel {
    ...
    proc getChild(i) return nil;
    ...
}
```





Example: chpl_executeOn()

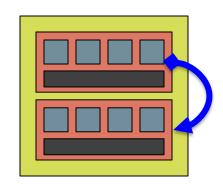
```
// no off-node case, locale ids simplified
proc chpl executeOn(loc, fn, args) {
  if (loc == here) {
    // run directly on this numa domain
    chpl ftable call(fn, args);
  } else {
    // move to a different numa domain
    var orig = here.sublocid;
    chpl task setSubloc(loc.sublocid);
    chpl ftable call(fn, args);
    chpl task setSubloc(orig);
```





Example: chpl_taskListAddBegin()

```
// Variant 1:
// Start the task on the specified sublocale
// Let the tasking layer decide what to do with "any"
proc chpl_taskListAddBegin(loc, fn, args) {
   chpl_task_addToTaskList(fn, args, loc.sublocid, ...);
}
```





Example: chpl_taskListAddBegin()

```
Variant 1.
 / Variant 2:
// Start the task on the specified sublocale
// Schedule "any" in a round robin fashion
proc chpl taskListAddBegin(loc, fn, args) {
  const sublocid
      = if loc.sublocid != subloc any then
          loc.sublocid; // use specified subloc
        else
          loc.nextSubloc.fetchAdd(1) % loc.numSublocales;
                        // round robin using atomic count
  chpl task addToTaskList(fn, args, sublocid, ...);
```



Example: chpl_taskListAddBegin()

```
Variant 1.
// Variant 3:
// Start the task on the specified sublocale
// Schedule "any" based on load
proc chpl taskListAddBegin(loc, fn, args) {
  const sublocid
    = if loc.sublocid != subloc any then
        loc.sublocid; // use specified subloc
      else
       getBestSubloc(loc);
  chpl task addToTaskList(fn, args, sublocid, ...);
proc getBestSubloc(loc) {
  const (, sublocid)
          = minloc reduce (loc.numTasks(), 0..#numSubLocs);
  return sublocid;
```



Back to STREAM Triad

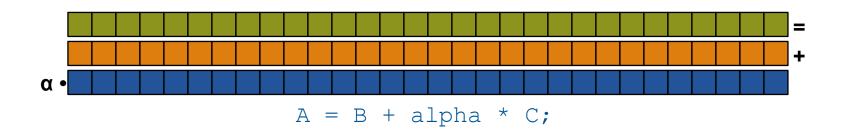




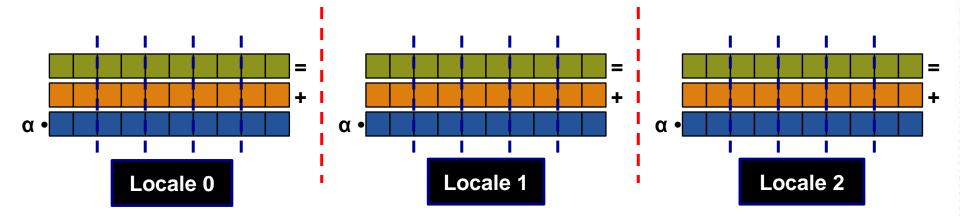
The Special Sauce: Domain maps



Domain maps are "recipes" that instruct the compiler how to map the global view of a computation...



...to the target locales' memory and processors:





Domain maps in a nutshell



Domain maps specify

- data allocation and layout
- parallel iteration
- low level implementation of array and domain operations
 - e.g., slicing, reallocating, reshaping

Domain maps are written in Chapel

- Class and procedural interface for the compiler
- Utilizes lower level concepts like coforall and on statements
- Programmer role: HPC specialist



Back to STREAM Triad

a = b + alpha * c;

- How many tasks?
- Where are they executed?
- How are the iterations assigned to tasks?

A: Chapel's *leader-follower* iterators are designed to give programmers full control over such decisions



Leader-Follower Iterators: Definition

- Chapel defines all forall loops in terms of leaderfollower iterators:
 - leader iterators: create parallelism, assign iterations to tasks
 - follower iterators: serially execute work generated by leader

Given...

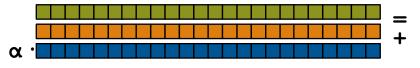
- ... A is defined to be the *leader*
- ...A, B, and C are all defined to be followers



Leader-follower iterators: Rewriting

Conceptually, the Chapel compiler translates:

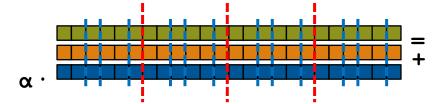
```
forall (a,b,c) in zip(A,B,C) do
  a = b + alpha * c;
```



Leader

into:

```
coforall subloc in sublocales do on subloc do
  coforall tid in here.numCores {
    for (a,b,c) in zip(A,B,C) {
        a = b + alpha * c;
    }
    Followers
}
```





Outline

CRAY

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Implementation Status -- Version 1.8.0 (Oct 2013)



Overall Status:

- Most features work at a functional level
 - some features need to be improved or re-implemented (e.g., OOP)
- Many performance optimizations remain
 - particularly for distributed memory (multi-locale) execution



Next Steps



- Evolve from Prototype- to Production-grade
 - Add/Improve missing features
 - Performance optimizations
- Target more complex compute node types
 - e.g., CPU+GPU, Intel Phi, ...
- Continue to grow the user and developer communities
 - Work toward transitioning Chapel from Cray-controlled to community-governed



Chapel Design Goal Revisited

Write code like this...

Have it perform like this...

```
#ifdef OPENMP
#include <omp.h>
#endif
static int VectorSize;
static double *a, *b, *c;
int HPCC StarStream (HPCC Params *params) {
 int myRank, commSize;
 int rv, errCount;
 MPI Comm comm = MPI COMM WORLD;
 MPI Comm size ( comm, &commSize );
 MPI Comm rank ( comm, &myRank );
 rv = HPCC_Stream( params, 0 == myRank);
 MPI Reduce ( &rv, &errCount, 1, MPI INT, MPI SUM, 0, comm );
 return errCount;
int HPCC Stream(HPCC_Params *params, int doIO) {
 register int j;
 double scalar;
 VectorSize = HPCC LocalVectorSize( params, 3, sizeof(double), 0 );
 a = HPCC XMALLOC( double, VectorSize );
 b = HPCC XMALLOC ( double, VectorSize );
 c = HPCC XMALLOC( double, VectorSize );
  if (!a || !b || !c) {
   if (c) HPCC free(c);
   if (b) HPCC free(b);
   if (a) HPCC free(a);
     fprintf( outFile, "Failed to allocate memory (%d).\n", VectorSize );
      fclose(outFile);
```

...on any architecture available

```
#pragma omp parallel for
#endif
  for (j=0; j<VectorSize; j++) {
    b[j] = 2.0;
    c[j] = 0.0;
}
scalar = 3.0;
#ifdef _OPENMP
#pragma omp parallel for
#endif
  for (j=0; j<VectorSize; j++)
    a[j] = b[j]+scalar*c[j];
...</pre>
```



Chapel Design Goal Revisted



√ High-level programmability

- Applications programmers expression parallelism and locality
- Chapel's locale models and domain maps can insulate programmers from details of the architecture

~ Performance

- Designed to enable good performance, results yet to be seen
- In cases where we've optimized, we've seen very promising results, e.g.,
 - Dynamic iterators that match the performance of OpenMP [Chamberlain, et al., PGAS 2011]
 - Communication aggregation implemented via the domain map interface [Sanz, et al., SBAC-PAD 2012]



Chapel at SC13

- Emerging Technologies Booth (all week)
 - Booth #3547: staffed by Chapel team members; poster and handouts
- Talk (Tues @ 3:20): Hierarchical Locales: Exposing the Node Architecture in Chapel
 - KISTI booth (#3713): Sung-Eun Choi (Cray Inc.)
- Poster (Tues @ 5:15): Towards Co-Evolution of Auto-Tuning and Parallel Languages
 - Posters Session: Ray Chen (University of Maryland)
- Chapel Lightning Talks BoF (Wed @ 12:15)
 - 5-minute talks on education, MPI-3, Big Data, Autotuning, Futures, MiniMD
- Talk (Wed @ 4:30): Chapel, an Emerging Parallel Language
 - HPC Impact Theatre (booth #3947): Brad Chamberlain (Cray Inc.)
- **Happy Hour** (Wed @ 5pm): 4th annual Chapel Users Group (CHUG) Happy Hour
 - Pi Bar (just across the street at 1400 Welton St): open to public, dutch treat
- HPC Education (Thus @ 1:30pm): High-Level Parallel Programming Using Chapel
 - David Bunde (Knox College) and Kyle Burke (Colby College)

