



Chapel: Background



Sustained Performance Milestones



1 GF - 1988: Cray Y-MP; 8 Processors

• Static finite element analysis





1 TF - 1998: Cray T3E; 1,024 Processors

Modeling of metallic magnet atoms

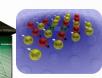




1 PF - 2008: Cray XT5; 150,000 Processors

• Superconductive materials





1 EF - ~2018: Cray ____; ~10,000,000 Processors

• TBD



Sustained Performance Milestones



1 GF - 1988: Cray Y-MP; 8 Processors

- Static finite element analysis
- Fortran77 + Cray autotasking + vectorization





1 TF – 1998: Cray T3E; 1,024 Processors

- Modeling of metallic magnet atoms
- Fortran + MPI (?)





1 PF – 2008: Cray XT5; 150,000 Processors

- Superconductive materials
- C++/Fortran + MPI + vectorization





- 1 EF ~2018: Cray ____; ~10,000,000 Processors
- TBD
- TBD: C/C++/Fortran + MPI + CUDA/OpenCL/OpenMP/??? or ???







Why Do HPC Programming Models Change?

HPC has traditionally given users...

...low-level, control-centric programming models

...ones that are closely tied to the underlying hardware

benefits: lots of control; decent generality; easy to implement

downsides: lots of user-managed detail; brittle to changes





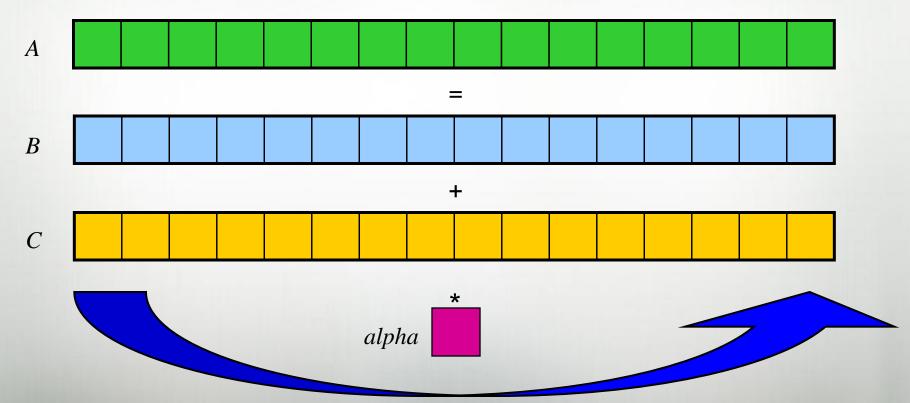
Introduction to STREAM Triad



Given: m-element vectors A, B, C

Compute: $\forall i \in 1..m, A_i = B_i + \alpha \cdot C_i$

Pictorially:







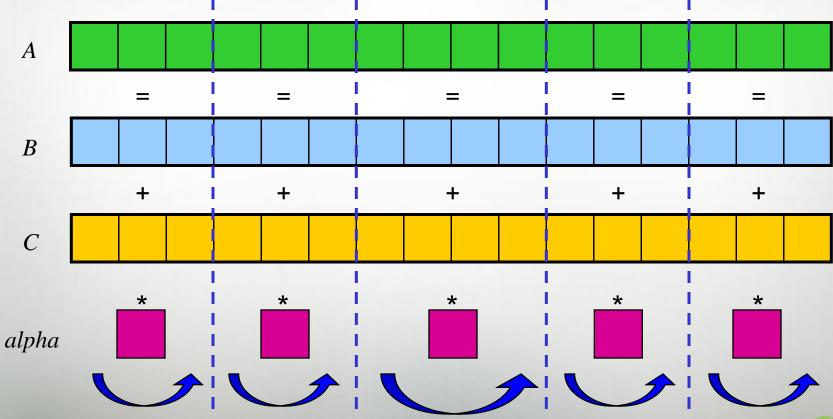
Introduction to STREAM Triad



Given: *m*-element vectors *A*, *B*, *C*

Compute: $\forall i \in 1..m, A_i = B_i + \alpha \cdot C_i$

Pictorially (in parallel):









```
MPI
```

```
#include <hpcc.h>
static int VectorSize;
static double *a, *b, *c;
int HPCC StarStream(HPCC Params *params) {
  int myRank, commSize;
 int rv, errCount;
 MPI Comm comm = MPI COMM WORLD;
 MPI Comm size ( comm, &commSize );
 MPI Comm rank ( comm, &myRank );
  rv = HPCC Stream( params, 0 == myRank);
 MPI Reduce ( &rv, &errCount, 1, MPI INT, MPI SUM,
   0, comm );
  return errCount;
int HPCC Stream(HPCC Params *params, int doIO) {
  register int j;
  double scalar;
 VectorSize = HPCC LocalVectorSize( params, 3,
   sizeof(double), 0 );
  a = HPCC XMALLOC( double, VectorSize );
  b = HPCC XMALLOC( double, VectorSize );
  c = HPCC XMALLOC( double, VectorSize );
```

```
if (!a || !b || !c) {
  if (c) HPCC free(c);
  if (b) HPCC free(b);
  if (a) HPCC free(a);
  if (doIO) {
    fprintf( outFile, "Failed to allocate memory
  (%d).\n", VectorSize );
    fclose( outFile );
  return 1;
for (j=0; j<VectorSize; j++) {</pre>
 b[j] = 2.0;
  c[i] = 0.0;
scalar = 3.0;
for (j=0; j<VectorSize; j++)</pre>
  a[i] = b[i] + scalar*c[i];
HPCC free(c);
HPCC free(b);
HPCC free(a);
return 0;
```







MPI + OpenMP #include <hpcc.h> #ifdef OPENMP #include <omp.h> #endif static int VectorSize; static double *a, *b, *c; int HPCC StarStream(HPCC Params *params) { int myRank, commSize; int rv, errCount; MPI Comm comm = MPI COMM WORLD; MPI Comm size (comm, &commSize); MPI Comm rank (comm, &myRank); rv = HPCC Stream(params, 0 == myRank); MPI Reduce (&rv, &errCount, 1, MPI INT, MPI SUM, 0, comm); return errCount; int HPCC Stream(HPCC Params *params, int doIO) { register int j; double scalar; VectorSize = HPCC LocalVectorSize(params, 3, sizeof(double), 0); a = HPCC XMALLOC(double, VectorSize); b = HPCC XMALLOC(double, VectorSize); c = HPCC XMALLOC(double, VectorSize);

```
if (!a || !b || !c) {
    if (c) HPCC free(c);
    if (b) HPCC free(b);
    if (a) HPCC free(a);
    if (doIO) {
      fprintf( outFile, "Failed to allocate memory
    (%d).\n", VectorSize);
      fclose( outFile );
    return 1;
#ifdef OPENMP
#pragma omp parallel for
#endif
  for (j=0; j<VectorSize; j++) {</pre>
   b[j] = 2.0;
    c[j] = 0.0;
  scalar = 3.0;
#ifdef OPENMP
#pragma omp parallel for
#endif
  for (j=0; j<VectorSize; j++)</pre>
    a[i] = b[i] + scalar*c[i];
  HPCC free(c);
  HPCC free (b);
  HPCC free(a);
  return 0;
```







MPI + OpenMP

```
#include <hpcc.h>
#ifdef OPENMP
#include <omp.h>
#endif
static int VectorSize;
static double *a, *b, *c;
int HPCC StarStream(HPCC Params *params)
 int myRank, commSize;
 int rv, errCount;
 MPI Comm comm = MPI COMM WORLD;
 MPI Comm size ( comm, &commSize );
 MPI_Comm_rank( comm, &myRank);
 rv = HPCC Stream( params, 0 == myRank);
 MPI Reduce ( &rv, &errCount, 1, MPI INT, MPI SUM, 0, comm );
 return errCount;
int HPCC_Stream(HPCC_Params *params, int doIO) {
 register int j;
 double scalar:
 VectorSize = HPCC LocalVectorSize( params, 3, sizeof(double), 0 );
 a = HPCC XMALLOC( double, VectorSize );
 b = HPCC XMALLOC( double, VectorSize );
 c = HPCC XMALLOC( double, VectorSize );
 if (!a || !b || !c) {
   if (c) HPCC free(c);
   if (b) HPCC free(b);
   if (a) HPCC_free(a);
   if (doIO) {
     fprintf( outFile, "Failed to allocate memory (%d).\n", VectorSize );
     fclose( outFile );
   return 1;
#ifdef OPENMP
#pragma omp parallel for
 for (j=0; j<VectorSize; j++) {
  b[j] = 2.0;
   c[j] = 0.0;
 scalar = 3.0;
#ifdef OPENMP
#pragma omp parallel for
 for (j=0; j<VectorSize; j++)
  a[j] = b[j] + scalar*c[j];
 HPCC free(c);
 HPCC free (b);
 HPCC free(a);
```

return 0;

CUDA

```
#define N
                2000000
int main() {
  float *d a, *d b, *d c;
  float scalar;
  cudaMalloc((void**)&d a, sizeof(float)*N);
  cudaMalloc((void**)&d b, sizeof(float)*N);
  cudaMalloc((void**)&d c, sizeof(float)*N);
  dim3 dimBlock(128);
  dim3 dimGrid(N/dimBlock.x );
  if ( N % dimBlock.x != 0 ) dimGrid.x+=1;
  set array<<<dimGrid,dimBlock>>>(d b, .5f, N);
  set array<<<dimGrid,dimBlock>>>(d c, .5f, N);
  scalar=3.0f;
  STREAM Triad<<<dimGrid,dimBlock>>>(d b, d c, d a, scalar, N);
  cudaThreadSynchronize();
  cudaFree(d a);
  cudaFree (d b);
  cudaFree(d c);
  global void set array(float *a, float value, int len) {
  int idx = threadIdx.x + blockIdx.x * blockDim.x;
  if (idx < len) a [idx] = value;
 global void STREAM Triad(float *a, float *b, float *c,
                              float scalar, int len) {
  int idx = threadIdx.x + blockIdx.x * blockDim.x;
  if (idx < len) c[idx] = a[idx]+scalar*b[idx];</pre>
```





```
MPI + OpenMP
                                                                                                   CUDA
#include <hpcc.h>
#ifdef OPENMP
                                                              #define N
                                                                                  2000000
#include <omp.h>
#endif
                                                             int main() {
static int VectorSize;
static double *a, *b, *c;
                                                                float *d a, *d b, *d c;
                                                                float scalar;
int HPCC StarStream (HPCC Params *params) {
 int myRank, commSize;
 int rv, errCount;
 MPI Comm comm = MPI COMM WORLD;
                                                                   Chapel
 MPI Comm size ( comm, &commSize );
 MPI Comm rank ( comm, &myRank );
                                                                                                                          the special
 rv = HPCC Stream( params, 0 == myRar
 MPI Reduce ( &rv, &errCount, 1, MPI
                              config const m = 1000,
                                                                                                                          sauce
                                                         alpha = 3.0;
int HPCC_Stream(HPCC_Params *params,
 register int j;
 double scalar;
                                                                                                                      N);
                              const ProbSpace = [1..m] dmapped
 VectorSize = HPCC LocalVectorSize(
                                                                                                                      N);
 a = HPCC XMALLOC( double, VectorSize
 b = HPCC XMALLOC ( double, VectorSize
 c = HPCC XMALLOC( double, VectorSize
                              var A, B, C: [ProbSpace] real;
 if (!a || !b || !c) {
                                                                                                                      c, da, scalar, N);
  if (c) HPCC_free(c);
  if (b) HPCC free(b);
  if (a) HPCC_free(a);
  if (doIO) {
    fprintf ( outFile, "Failed to al
    fclose( outFile );
   return 1;
#ifdef OPENMP
                              A = B + alpha * C;
#pragma omp parallel for
                                                                                                                     value, int len) {
 for (j=0; j<VectorSize; j++) {
                                                                int idx = threadIdx.x + blockIdx.x * blockDim.x;
  b[j] = 2.0;
  c[j] = 0.0;
                                                                if (idx < len) a [idx] = value;
 scalar = 3.0;
#ifdef OPENMP
                                                               global void STREAM Triad(float *a, float *b, float *c,
#pragma omp parallel for
                                                                                                    float scalar, int len) {
 for (j=0; j<VectorSize; j++)
  a[j] = b[j] + scalar*c[j];
                                                                int idx = threadIdx.x + blockIdx.x * blockDim.x;
 HPCC free(c);
                                                                if (idx < len) c[idx] = a[idx]+scalar*b[idx];
 HPCC free (b);
 HPCC free(a);
return 0;
```



Why Do HPC Programming Models Change?

HPC has traditionally given users...

...low-level, control-centric programming models

...ones that are closely tied to the underlying hardware

benefits: lots of control; decent generality; easy to implement downsides: lots of user-managed detail; brittle to changes

one characterization of Chapel's goals:

- Raise the level of abstraction to insulate parallel algorithms from underlying hardware when possible/practical
- Yet permit control over such details using appropriate abstraction and separation of concerns



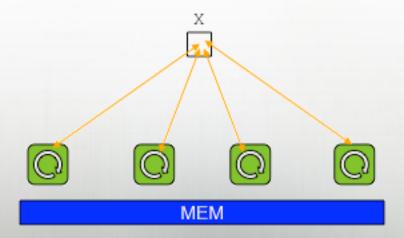


Shared Memory Programming Models



e.g., OpenMP, pthreads

- + support dynamic, fine-grain parallelism
- + considered simpler, more like traditional programming
 - "if you want to access something, simply name it"
- no support for expressing locality/affinity; limits scalability
- bugs can be subtle, difficult to track down (race conditions)
- tend to require complex memory consistency models



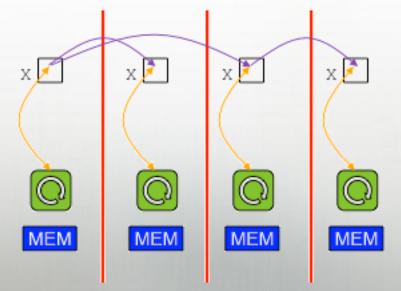




Distributed Memory Programming Models

e.g., MPI

- + a more constrained model; can only access local data
- + run on most large-scale parallel platforms
 - and for many of them, can achieve near-optimal performance
- + are relatively easy to implement
- + can serve as a strong foundation for higher-level models
- + users are able to get real work done with them





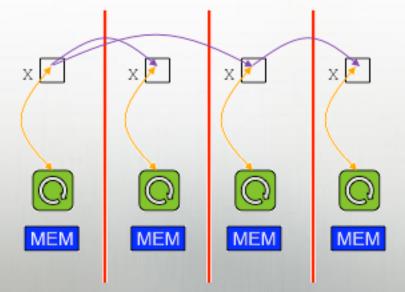




Distributed Memory Programming Models

e.g., MPI

- communication must be used to get copies of remote data
 - and tends to reveal too much about how to transfer data, not simply what
- only supports "cooperating executable"-level parallelism
- couples data transfer and synchronization
- has frustrating classes of bugs of its own
 - e.g., mismatches between sends/recvs, buffer overflows, etc.





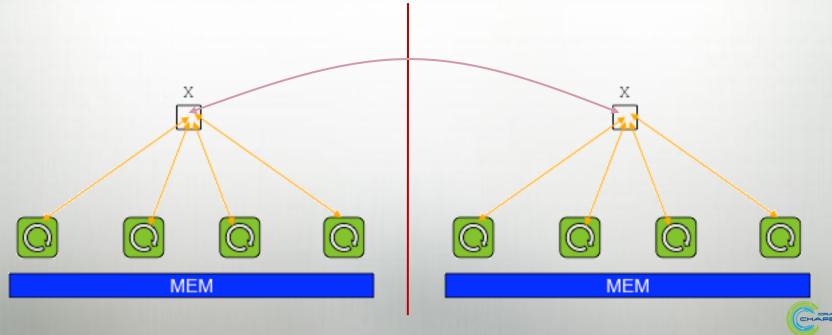


Hybrid Programming Models



e.g., MPI+OpenMP, MPI+pthreads, MPI+CUDA, ...

- + support a division of labor: each handles what it does best
- + permit overheads to be amortized across processor cores
- require multiple distinct notations to express a single logical parallel algorithm, each with its own distinct semantics



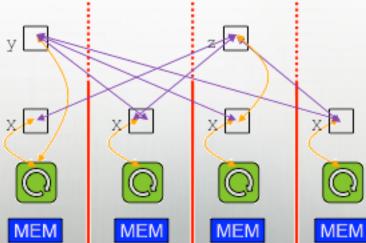




PGAS (Partitioned Global Address Space) Models

e.g., Co-Array Fortran (CAF), Unified Parallel C (UPC)

- + support a shared namespace, like shared-memory
- + support a strong sense of ownership and locality
 - each variable is stored in a particular memory segment
 - tasks can access any visible variable, local or remote
 - local variables are cheaper to access than remote ones
- + implicit communication eases user burden; permits compiler use best mechanisms available





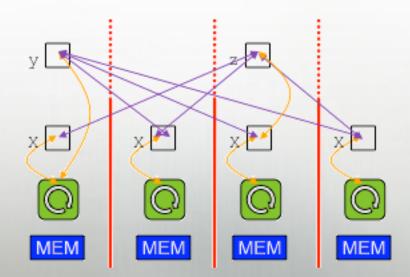




PGAS (Partitioned Global Address Space) Models

e.g., Co-Array Fortran (CAF), Unified Parallel C (UPC)

- retain many of the downsides of shared-memory
 - error cases, memory consistency models
- restricted to SPMD programming and execution models
- data structures not as flexible/rich as one might like



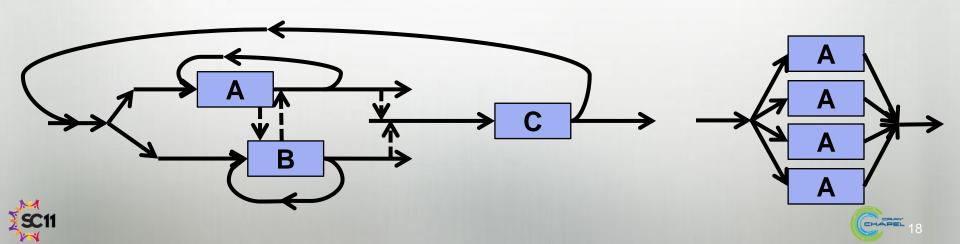






Chapel: A Next-Generation PGAS Language

- a PGAS language, but non-traditional:
 - more general/dynamic/multithreaded parallelism
 - concepts for composable data and task parallelism
 - distinct concepts for locality vs. parallelism
 - e.g., locale type represents architectural locality
 - productivity features
 - type inference, iterator functions, rich array types, OOP, ...



Chapel's Origins



• HPCS: High Productivity Computing Systems #PCS



- Overall goal: Raise high-end user productivity by 10x
 Productivity = Performance + Programmability + Portability + Robustness
- Phase II: Cray, IBM, Sun (July 2003 June 2006)
 - Goal: Propose new productive system architectures
 - Each vendor created a new programming language
 - Cray: Chapel
 - IBM: X10
 - Sun: Fortress
- Phase III: Cray, IBM (July 2006)
 - Goal: Develop the systems proposed in phase II
 - Each vendor implemented a compiler for their language
 - Sun also continued their Fortress effort without HPCS funding





Chapel's Productivity Goals



- Vastly improve programmability over current languages
 - Writing parallel programs
 - Reading, modifying, porting, tuning, maintaining them
- Support performance at least as good as MPI
 - Competitive with MPI on generic clusters
 - Better than MPI on more capable architectures
- Improve portability over current languages
 - As ubiquitous as MPI but more abstract
 - More portable than OpenMP, UPC, and CAF are thought to be
- Improve robustness via improved semantics
 - Eliminate common error cases
 - Provide better abstractions to help avoid other errors





Outline



- Chapel's Context
- Chapel's Motivating Themes
 - 1. General parallel programming
 - 2. Global-view abstractions
 - 3. Multiresolution design
 - 4. Control over locality/affinity
 - 5. Reduce gap between mainstream & HPC languages







With a unified set of concepts...

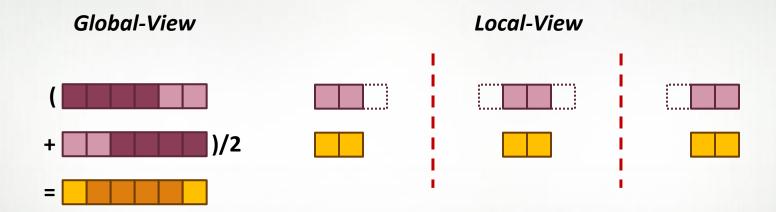
- ...express any parallelism desired in a user's program
 - Styles: data-parallel, task-parallel, concurrency, nested, ...
 - Levels: model, function, loop, statement, expression
- ...target all parallelism available in the hardware
 - Systems: multicore desktops, clusters, HPC systems, ...
 - Levels: machines, nodes, cores, instructions







In pictures: "Apply a 3-Point Stencil to a vector"

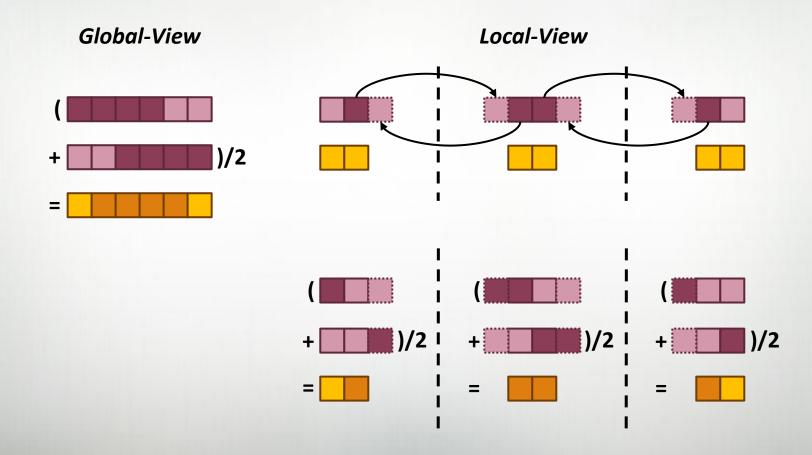








In pictures: "Apply a 3-Point Stencil to a vector"







2) Global-View Abstractions



In code: "Apply a 3-Point Stencil to a vector"

Global-View

```
proc main() {
    var n = 1000;
    var A, B: [1..n] real;

    forall i in 2..n-1 do
        B[i] = (A[i-1] + A[i+1])/2;
}
```

Local-View (SPMD)

```
proc main() {
 var n = 1000;
  var p = numProcs(),
      me = myProc(),
      myN = n/p
  var A, B: [0..myN+1] real;
  if (me < p-1) {
    send(me+1, A[myN]);
    recv (me+1, A[myN+1]);
  if (me > 0) {
    send (me-1, A[1]);
    recv (me-1, A[0]);
  forall i in 1..myN do
    B[i] = (A[i-1] + A[i+1])/2;
```

Bug: Refers to uninitialized values at ends of A



2) Global-View Abstractions



In code: "Apply a 3-Point Stencil to a vector"

Global-View

```
proc main() {
   var n = 1000;
   var A, B: [1..n] real;

   forall i in 2..n-1 do
       B[i] = (A[i-1] + A[i+1])/2;
}
```

Communication becomes geometrically more complex for higher-dimensional arrays

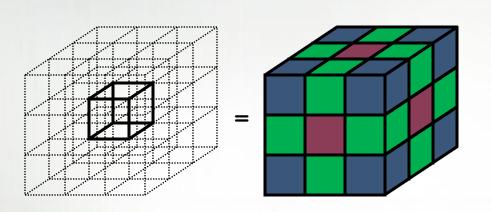
Local-View (SPMD)

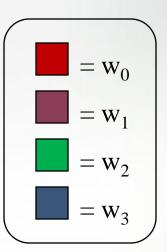
```
proc main (/
                     Assumes p divides n
  var n = 1000;
  var p = numProcs(),
      me = myProc(),
      myN = n/p
      mvLo = 1,
      myHi = myN;
  var A, B: [0..myN+1] real;
  if (me < p-1) {
    send(me+1, A[myN]);
    recv (me+1, A[myN+1]);
  } else
    myHi = myN-1;
  if (me > 0) {
    send (me-1, A[1]);
    recv (me-1, A[0]);
  } else
    myLo = 2;
  forall i in myLo..myHi do
    B[i] = (A[i-1] + A[i+1])/2;
```

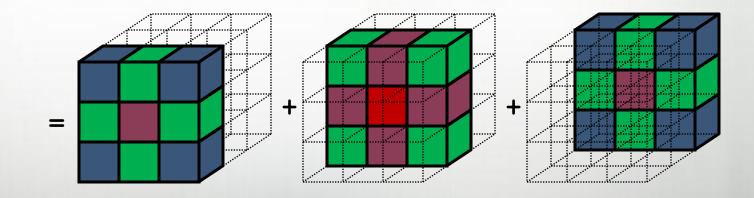


2) rprj3 Stencil from NAS MG















2) rprj3 Stencil from NAS MG in Fortran + MPI

```
subroutine rprj3( r,mlk,m2k,m3k,s,mlj,m2j,m3j,k)
 implicit none
 double precision x1(m), y1(m), x2,y2
 else
d3 = 1
        if( debug_vec(0) .ge. 1 ) then
  call rep_nrm(s,mlj,m2j,m3j,' rprj3',k-1)
endif
 subroutine norm2u3(r,n1,n2,n3,rnm2,rnmu,nx,ny,nx)
 integer nl, n2, n3, nx, ny, nz
double precision rnm2, rnmu, r(n1,n2,n3)
double precision s, a, ss
integer i3, i2, i1, ierr
 dn = 1.0d0*nx*nv*nz
> mpr.mnu = ss
call mpi allreduce(s, ss, 1, dp_type,
> mpi_sum,mpi_comm_world,ierr)
s = ss
rnm2=sqrt(s / dn)
 subroutine rep nrm(u.nl.n2.n3.title.kk)
 integer nl, n2, n3, kk
double precision u(n1,n2,n3)
character*8 title
```

```
call norm2u3(u,nl,n2,n3,rmm2,rmmu,nx(kk),ny(kk),nr(kk))
if (ms.eq. root) then
    write(*,7)kk,title,rmm2,rmmu
    format('Level',12,'in',a8,':norms=',D21.14,D21.14)
    andif
 implicit none
if( .not. dead(kk) ) then
  do axis = 1, 3
   if( nprocs .ne. 1) then
  call zero3(u,n1,n2,n3)
endif
return
end
  subroutine ready( axis, dir, k )
  integer axis, dir, k
integer buff_id,buff_len,i,ierr
do i=1,nm2
buff(i,buff_id) = 0.0D0
  msg_id(axis,dir,1) = msg_type(axis,dir) +1000*me
call mpi irecv(buff(l,buff_id), buff_len,

> dp_type, nbr(axis, -dir,k), neg_type(axis,dir),

> pain_comm_world, msg_id(axis,dir,l), ierr)

return
end
  subroutine give3( axis, dir, u, n1, n2, n3, k )
 implicit none
 integer axis, dir, n1, n2, n3, k, ierr double precision u( n1, n2, n3 )
 integer i3, i2, i1, buff_len,buff_id
          do i3=2,n3-1
do i2=2,n2-1
buff_len = buff_len + 1
buff[buff_len,buff_id] = u(2, i2,i3)
enddo
          call mpi_send(
  buff(1, buff_id), buff_len,dp_type,
    nbr(axis,dir,k), mag_type(axis,dir),
  mpi_comm_world,ierr)
                    buff_len = buff_len + 1
buff(buff_len, buff_id) = u(nl-1, i2,i3)
                     buff_len = buff_len + 1
```

```
buff(buff_len, buff_id ) = u(i1, 2,i3)
enddo
         call mpi_send(
   buff(I, buff_id), buff_len,dp_type,
   nbr(axis,dir,k), msg_type(axis,dir),
   mpi_comm_world,ierr)
     else if ( dir .eq. +1 ) then
        call mpi_send(
  buff(1, buff_id), buff_len,dp_type,
  nbr(axis, dir, k), msg_type(axis,dir),
  mpi_comm_world, ierr)
         call mpi_send(
   buff(1, buff id), buff len,dp_type,
   nbr(axis, dīr, k), msg_type(axis,dir),
   mpi_comm_world, ierr)
     else if (dir .eg. +1 ) then
                    buff len = buff len + 1
buff(buff len, buff id) = u(i1,i2,n3-1)
         call mpi_send(
   buff(I, buff id), buff_len,dp_type,
   nbr(axis,dIr,k), msg_type(axis,dir),
   mpi_comm_world,ierr)
subroutine take3 (axis. dir. u. nl. n2. n3)
integer axis, dir, n1, n2, n3
double precision u( n1, n2, n3 )
         do i3=2.n3-1
         id=2,n2-1
indx = indx + 1
u(n1,i2,i3) = buff(indx, buff_id)
enddo
enddo
    else if ( dir .eq. +1 ) then
         do i3=2,n3-1

do i2=2,n2-1

indx = indx + 1

u(1,i2,i3) = buff(indx, buff_id)
              do il=1,n1
  indx = indx + 1
  u(il.n2.i3) = buff(indx.buff id)
     else if ( dir .eq. +1 ) then
         do i3=2.n3-1
                   indx = indx + 1
u(i1,1,i3) = buff(indx, buff_id)
if(axis .eq. 3)then
```

```
u(i1,i2,n3) = buff(indx, buff id)
  subroutine commlp(axis, u, n1, n2, n3, kk)
 implicit none
  integer i3, i2, i1, buff_len,buff_id
integer i, kk, indx
 dir = -1
 buff_id = 3 + dir
buff len = nm2
  do i=1,nm2
buff(i,buff_id) = 0.000
 buff id = 3 + dir
  do i=1,nm2
buff(i,buff_id) = 0.000
 if(axis.eq. 1)then
do i3=2,n3-1
do i2=2,n3-1
bufflen= bufflen+ 1
buff(bufflen, buffid) = u(n1-1, i2,i3)
        do il=1,nl
buff_len = buff_len + 1
buff(buff_len, buff_id) = u(il,n2-1,i3)
if( axis .eq. 3 ) then
do i2-1,n2
do .n1
buff_len buff_len + 1
buff_len, buff_id ) = u( i1,i2,n3-1)
enddo
  buff id = 2 + dir
if (axis .eq. 1 ) then
do i3=2,n3-1
do i2=2,n2-1
buff len = buff_len + 1
buff(buff_len,buff_id) = u(2, i2,i3)
enddo
  if(axis .eq. 3)then
          do il=1,n2
do il=1,n1
buff len = buff len + 1
buff(buff_len, buff_id) = u(il,i2,2)
 do i=1,nm2
  buff(i,4) = buff(i,3)
  buff(i,2) = buff(i,1)
enddo
```

```
buff_id = 3 + dir
if( axis .eq. 1 )then
do i3=2,n3-1
do i2=2,n2-1
indx = indx + 1
u(n1,12,13) = buff(indx, buff_id)
enddo
enddo
          indx = indx + 1
  u(i1,n2,i3) = buff(indx, buff_id )
enddo
if( axis .eq. 3 ) then

do i2=1,n2

do i1=1,n1

indx = indx + 1

u(i1,i2,n3) = buff(indx, buff_id)
if( axis .eq. 1 )then
do i3=2,n3-1
do i2=2,n2-1
indx = indx + 1
u(1,i2,i3) = buff(indx, buff_id)
 if(axis .eq. 2)then
do i3=2,n3-1
          do il=1,nl
  indx = indx + 1
  u(il,1,i3) = buff(indx, buff_id)
if( axis .eq. 3 ) then
    do i2=1,n2
    do i1=1,n1
    indx = indx + 1
        u(i1,i2,1) = buff(indx, buff_id)
enddo
  integer n1,n2,n3,i1,i2,i3,i,ierr
double precision z(n1,n2,n3)
  double precision z(
integer ml, m2, m3
   do i=0,nprocs-1
if( me .eq. i )then
write(*,*)'id = ', me
           write(*,*)' id = ', me

do i3=1,m3

do i1=1,m1

write(*,6)(z(i1,i2,i3),i2=1,m2)

enddo

write(*,*)' - - - - - '
             write(*,*)' '
format(15f6.3)
call mpi_barrier(mpi_comm_world,ierr)
 subroutine zero3(z,n1,n2,n3)
 integer n1, n2, n3
double precision z(n1,n2,n3)
integer i1, i2, i3
do i3=1,n3
do i2=1,n2
do i1=1,n1
z(i1,i2,i3)=0.000
enddo
```





2) rprj3 Stencil from NAS MG in Chapel

Our previous work in ZPL demonstrated that such compact codes can result in better performance than Fortran + MPI while also supporting more flexibility at runtime*.

^{*}e.g., the Fortran + MPI rprj3 code shown previously not only assumes p divides n, it also assumes that p and n are specified at compile-time and powers of two.





2) Classifying Current Programming Models

	System	Data Model	Control Model
Communication Libraries	MPI/MPI-2	Local-View	Local-View
	SHMEM, ARMCI, GASNet	Local-View	SPMD
Shared Memory	OpenMP, Pthreads	Global-View (trivially)	Global-View (trivially)
PGAS Languages	Co-Array Fortran	Local-View	SPMD
	UPC	Global-View	SPMD
	Titanium	Local-View	SPMD
PGAS Libraries	Global Arrays	Global-View	SPMD







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HPCS Languages	Chapel	Global-View	Global-View
	X10 (IBM)	Global-View	Global-View
	Fortress (Sun)	Global-View	Global-View







2) Global-View Programming: A Final Note

 A language may support both global- and local-view programming — in particular, Chapel does

```
proc main() {
  coforall loc in Locales do
    on loc do
       MySPMDProgram(loc.id, Locales.numElements);
}
proc MySPMDProgram(me, p) {
    ...
}
```

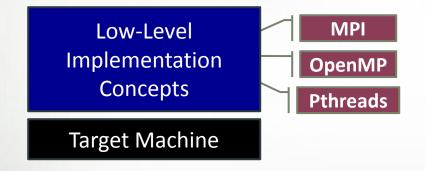






3) Multiresolution Language Design: Motivation





Target Machine

"Why is everything so difficult?"
"Why don't my programs port trivially?"

"Why don't I have more control?"



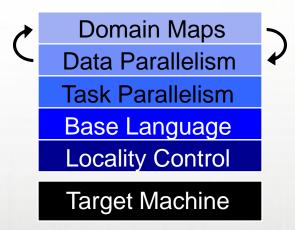






Multiresolution Design: Support multiple tiers of features

- higher levels for programmability, productivity
- lower levels for performance, control
- build the higher-level concepts in terms of the lower
 Chapel language concepts



separate concerns appropriately for clean design





4) Control over Locality/Affinity



Consider:

- Scalable architectures package memory near processors
- Remote accesses take longer than local accesses

Therefore:

- Placement of data relative to computation affects scalability
- Give programmers control of data and task placement

Note:

- As core counts grow, locality will matter more on desktops
- GPUs and accelerators already expose node-level locality







5) Reduce Gap Between HPC & Mainstream Languages

Consider:

- Students graduate with training in Java, Matlab, Perl, Python
- Yet HPC programming is dominated by Fortran, C/C++, MPI

We'd like to narrow this gulf with Chapel:

- to leverage advances in modern language design
- to better utilize the skills of the entry-level workforce...
- ...while not ostracizing the traditional HPC programmer
 - e.g., support object-oriented programming, but make it optional





Questions?



- Chapel's Context
- Chapel's Motivating Themes
 - 1. General parallel programming
 - 2. Global-view abstractions
 - 3. Multiresolution design
 - 4. Control over locality/affinity
 - 5. Reduce gap between mainstream & HPC languages



