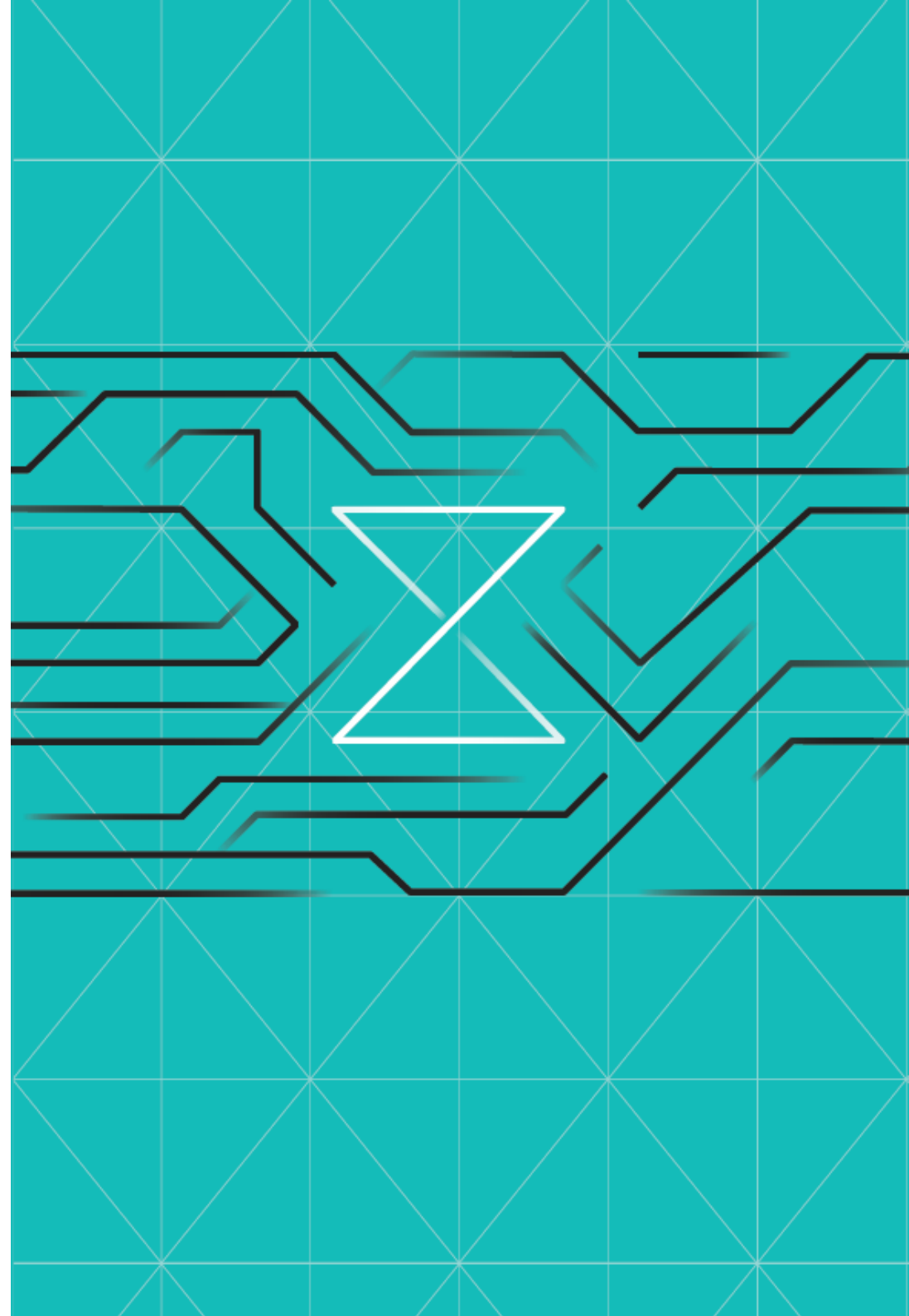


ASM2

Developers Assemble!

- [1 BACKGROUND](#)
- [2 PREPARATION](#)
- [3 THE CHALLENGE](#)
- [4 USE TSO TEST FACILITY](#)
- [5 MODIFY AND RE-COMPILE](#)
- [6 VALIDATION AND NEXT STEPS](#)



1 BACKGROUND

This challenge takes you a bit deeper into IBM Z Assembler Language.

Some large companies need programmers, developers and architects that understand IBM Z Assembler language.

These technical positions pay well and lead to promotions in the technical ranks. Familiarity with IBM Z Assembler on a resume is a major “attention-getter” for securing an interview. You can say you have familiarity with IBM Z Assembler as a result of completing these challenges.

Do not let the technical details in the challenge scare you away. By the end of this challenge you will realize you are capable of beginning to understand the technical details of how IBM Z computer software works.

2 PREPARATION

A relatively simple assembler program will be used and explained.

Each computer architecture has machine instructions unique to the architecture. All computer languages supported by the architecture must be translated into the unique machine instructions of the underlying computer architecture. Each computer architecture has an assembly language which includes mnemonics that are assembled into machine instructions understood by the computer. Compilers and interpreters translate supported computer languages into the unique machine instructions understood by the hosting computer.

Computer processing memory is used to load and store the machine instructions and data for processing. The operating system keeps track of processing memory locations using addresses where some of the memory is free, some of the memory has machine instructions, and some of the memory has data.

Higher level languages such as C/C++, Java, COBOL, etc. were created to make programming the computer easier by hiding the complexity of the underlying machine instructions, addressable memory, and registers. Registers are at the top of the memory hierarchy, and provide the fastest way to access data.

IBM Z mainframe computer architecture, z/Architecture, is what is commonly known as Complex Instruction Set Computing (CISC); intel x86 and compatible processors are also CISC architectures.

The assembler code that follows is explained in greater detail using the assembler compiled output.

```

6  *-----* 00060000
7  *      register equates      * 00070000
8  *-----* 00080000
9  R0      EQU   0              register 0      00090000
10 BASEREG EQU   12             base register   00100000
11 SAVEREG EQU   13             save area register 00110000
12 RETREG  EQU   14             caller's return address 00120000
13 ENTRYREG EQU   15            entry address   00130000
14 RETCODE EQU   15            return code     00140000
15      EJECT                    00150000
16 *-----* 00160000
17 *      standard entry setup, save area chaining, establish * 00170000
18 *      base register and addressability      * 00180000
19 *-----* 00190000
20      USING ASMPGM,ENTRYREG      establish addressability 00200000
21      B      SETUP              branch around eyecatcher 00210000
22      DC     CL8'ASMPGM'         program name   00220000
23      DC     CL8'&SYSDATE'       program assembled date 00230000
24      STM     RETREG,BASEREG,12(SAVEREG) save caller's registers 00240000
25      BALR    BASEREG,R0         establish base register 00250000
26      DROP    ENTRYREG          drop initial base register 00260000
27      USING   *,BASEREG         establish addressability 00270000
28      LA      ENTRYREG,SAVEAREA  point to this program save area 00280000
29      ST      SAVEREG,4(,ENTRYREG) save address of caller 00290000
30      ST      ENTRYREG,8(SAVEREG) save address of this program 00300000
31      LR      SAVEREG,ENTRYREG   point to this program savearea 00310000
32      EJECT                    00320000

```

The section shown above is responsible for “standard linkage”, a technique for passing data and processing control from a caller program to a called program.

Read the comments to the right of the assembler mnemonics.

```

33 *-----* 00330000
34 *      program body      * 00340000
35 *-----* 00350000
36      L      2,=C'Begin'      00360001
37      LA     2,=C'Begin'      00370001
38      LOOPINIT DS    0H        00380000
39      SR     2,2              00390000
40      L      2,=F'4'          00400001
41      L      3,=F'1'          00410000
42      LOOP   DS    0H        00420000
43      A      3,=F'1'          00430001
44      BCT    2,LOOP           00440001
45      STOP1  LH     3,HALFCON  00450001
46      STOP2  A      3,FULLCON  00460001
47      STOP3  ST     3,HEXCON   00470001
48      EJECT                    00471001

```

Next comes the logic to be executed, the program body.

```

61  *      storage and constant definitions.                * 00590000
62  *      print output definition.                          * 00600000
63  *-----* 00610000
64  SAVEAREA DC 18F'-1'                                register save area 00710000
65  FULLCON  DC F'-1'                                    00711001
66  HEXCON   DC XL4'9ABC'                                00712001
67  HALFCON  DC H'32'                                    00713001
68  |         | END ASMPGM                                00720000

```

Generally, the bottom of an Assembler source code contains the definition of program constants, and other structures for handling data in processing memory storage.

```

69  //GO      EXEC PGM=ASMPGM                            00730000
70  //STEPLIB DD DSN=&SYSUID..LOAD,DISP=SHR              00740000
71  //PRINT   DD SYSOUT=*                                00750002
72

```

And finally, the job-step that runs the program created by the Assemble and Link steps

3 THE CHALLENGE

Compile assembler program

- copy 'ZXP.PUBLIC.SOURCE(ASMPGM)' to your SOURCE dataset.
- submit the program with right-click from the DATA SETS view in VSCode, and select Submit Job.
- Review the assembler compilation details by checking the job output called ASMPGM:SYSPRINT

65	Loc	Object Code	Addr1	Addr2	Stmt	Source Statement	HLASM R6.0	2023/02/14 06.17
66	000000		00000	000ED	2	ASMPGM CSECT		00050000
67					3	*-----*		00060000
68					4	* register equates		* 00070000
69					5	*-----*		00080000
70			00000		6	R0 EQU 0 register 0		00090000
71			0000C		7	BASEREG EQU 12 base register		00100000
72			0000D		8	SAVEREG EQU 13 save area register		00110000
73			0000E		9	RETREG EQU 14 caller's return address		00120000
74			0000F		10	ENTRYREG EQU 15 entry address		00130000
75	i		0000F		11	RETCODE EQU 15 return code		00140000

Things to look for in the compiler listings:

The column headings in the program listing:

- Loc - abbreviation for location, program location column
- Object Code - machine instructions column. commonly called op codes
- Addr1 - memory address location column
- Addr2 - memory address location column
- Stmt - line numbers associated with each line of code

- Source Statement - assembler code labels, mnemonics, and mnemonic operands

The program code itself:

- Stmt 2 is the first assembler code line where:
Loc 00000 represents the beginning location
Addr1 00000 Addr2 000FD represents the starting and ending location of the program in hexadecimal - effectively the program length in bytes - x'FD' or 253 decimal
ASMPGM is a user chosen label
CSECT is an assembler directive declaring a Control Section
- Stmt 3, 4, 5 are comments - identified by the character in column 1 being an asterisk "*"
 - Stmt 6 to 11 show user selected labels equated (EQU) to computer register numbers - names to use in the program instead of just the register numbers.

Following Stmt 11 is assembler code frequently reused for ["standard linkage"](#) - storing the location and register contents of the calling program to return control to when this program is finished.

79		13	*-----*	00160000
80		14	* standard entry setup, save area chaining, establish	* 00170000
81		15	* base register and addressability	* 00180000
82		16	*-----*	00190000
83	R:F 00000	17	USING ASMPGM,ENTRYREG	establish addressability 00200000
84	000000 47F0 F014 00014	18	B SETUP	branch around eyecatcher 00210000
85	000004 C1E2D4D7C7D44040	19	DC CL8'ASMPGM'	program name 00220000
86		20	DC CL8'&SYSDATE'	program assembled date 00230000
87	00000C F0F261F1F461F2F3	+	DC CL8'02/14/23'	program assembled date 00230000
88	000014 90EC D00C 0000C	21	SETUP STM RETREG,BASEREG,12(SAVEREG)	save caller's registers 00240000
89	000018 05C0	22	BALR BASEREG,R0	establish base register 00250000
90		23	DROP ENTRYREG	drop initial base register 00260000
91	R:C 0001A	24	USING *,BASEREG	establish addressability 00270000
92	00001A 41F0 C06E 00088	25	LA ENTRYREG,SAVEAREA	point to this program save area 00280000
93	00001E 50D0 F004 00004	26	ST SAVEREG,4(,ENTRYREG)	save address of caller 00290000
94	000022 50F0 D008 00008	27	ST ENTRYREG,8(,SAVEREG)	save address of this program 00300000
95	i000026 18DF	28	LR SAVEREG,ENTRYREG	point to this program savearea 00310000

- Stmt 18 includes assembler mnemonic **B**, a branch instruction.

The operand is SETUP, a program label that is a quick way to resolve addressable location in memory.

- To the left of Stmt 18 you can see an operation code ("OPCODE") that begins with 47, which is the opcode for the mnemonic **B**.
- Stmt 18 Addr2 is 00014. Loc 00014, Stmt 22, is label SETUP with more assembler code to execute.

While the explanation below is elaborate, it is just preparation for you to see these changes happen at breakpoints during program execution later in the challenge.

The program body includes -

Statement	OPCODE	Mnemonic	Description
34	x'58'	L	Load register 2 with characters 'Begin'
35	x'41'	LA	Load Address register 2 with address of 'Begin' address location. You will see the register content difference between the 2 instructions later. You will might notice the Addr2 is 000F8 memory address location for 'Begin', =C'Begin' resulted in assembler assigning a memory address location for the characters "Begin"
37	x'1B'	SR	subtract register 2 from itself
38	x'58'	L	load 4 into register 2

ASM2|230808-2334

Statement	OPCODE	Mnemonic	Description
39	x'58'	L	load 1 into register 3
42	x'5A'	A	add 1 to register 3
43	x'46'	BCT	decrement register 2 by 1 and branch to the label loop until register 2 is zero
45	x'48'	LH	load halfword into register with content of memory location assigned to HALFCON label
46	x'5A'	A	add memory location assigned to FULLCON label into register 3
47	x'50'	ST	store register 3 in memory location of HEXCON

ASM2/230809-2334

```

99          30 *-----* 00330000
100          31 *      program body * 00340000
101          32 *-----* 00350000
102      000028 5820 C0CE      000E8 33      L      2,=C'Begin' 00360001
103      00002C 4120 C0CE      000E8 34      LA      2,=C'Begin' 00370001
104      000030          35 LOOPINIT DS      0H 00380000
105      000030 1B22          36      SR      2,2 00390000
106      000032 5820 C0C6      000E0 37      L      2,=F'4' 00400001
107      000036 5830 C0CA      000E4 38      L      3,=F'1' 00410000
108      00003A          39 LOOP      DS      0H 00420000
109      00003A 5A30 C0CA      000E4 40      A      3,=F'1' 00430001
110      00003E 4620 C020      0003A 41      BCT      2,LOOP 00440001
111      000042 4830 C0BE      000D8 42 STOP1 LH      3,HALFCON 00450001
112      000046 5A30 C0B6      000D0 43 STOP2 A      3,FULLCON 00460001
113      i00004A 5030 C0BA      000D4 44 STOP3 ST      3,HEXCON 00470001

```

At Stmt 56 you will see 'WTO':

- Left is blank - no op code, no operands, and no addresses
- WTO, abbreviation for Write To Operator, is an assembler macro
- The text following WTO is written to the system log

Below, you can see commonly reused code for 'standard linkage', restoring the caller program registers. This is how this program returns control to the program that called it.

```

117          46 *-----* 00480000
118          47 *      standard exit - restore caller's registers and * 00490000
119          48 *      return to caller * 00500000
120          49 *-----* 00510000
121      00004E          50 EXIT DS      0H halfword boundary alignment 00520000
122      00004E 58D0 D004      00004 51      L      SAVEREG,4(,SAVEREG) restore caller's save area addr 00530000
123      000052 58E0 D00C      0000C 52      L      RETREG,12(,SAVEREG) restore return address register 00540000
124      000056 980C D014      00014 53      LM      R0,BASEREG,20(SAVEREG) restore all regs. except reg15 00550000
125          54      WTO      'Giving control back to system' 00551001

```

Next up: Stmt 58 to 64 with a '+' following the statement numbers is the WTO macro expanded code

Stmt 65 with op code beginning with x'07' associated with mnemonic BR, branch

The address location represented by label RETREG is register 14 which contains the address location of the caller program.

So, the execution result is a branch to an address location of the caller program.

126	00005A 0700		56+	CNOP	0,4		01-WTO
127	00005C A715 0013	00082	57+	BRAS	1,IHB0001A	BRANCH AROUND MESSAGE	@LCC 01-WTO
128	000060 0021		58+	DC	AL2(33)	TEXT LENGTH	@YA17152 01-WTO
129	000062 0000		59+	DC	B'0000000000000000'	MCSFLAGS	01-WTO
130	000064 C789A58995874083		60+	DC	C'Giving control back to system'		X01-WTO
131	00006C 9695A39996934082		+			MESSAGE TEXT	@L6C
132	000082		61+	IHB0001A DS	0H		01-WTO
133	000082 0A23		62+	SVC	35	ISSUE SVC 35	@L6A 01-WTO
134	i000084 07FE		63	BR	RETREG	return to caller	00560000

And finally there are some defined constants with assembly-assigned addressable memory locations

The opcode area is used as content at the addressable memory location when the program starts execution.

A few observations -

- Label SAVEAREA at location 98 into the program is 18 fullwords in length
- Label FULLCON at location E0 into the program is 1 fullword in length
- Label HEXCON at location E4 into the program is a hexadecimal length of 4.
- Label HALFCON at location E8 into the program is a halfword where of decimal value H'32' but inspection of the stored value in memory is x'20'

- Program body included =F'4', =F'1', and =C'Begin', where the = resulted in the assembly to machine code assigning addressable locations for the values at the end of the executable program, F0, F4, and F8, respectively.

```
138          65 *-----* 00580000
139          66 *      storage and constant definitions.      * 00590000
140          67 *      print output definition.                  * 00600000
141          68 *-----* 00610000
142      000086 0000
143      000088 FFFFFFFFFFFFFFFF      69 SAVEAREA DC 18F'-1'      register save area      00710000
144      0000D0 FFFFFFFF      70 FULLCON DC F'-1'      00711001
145      0000D4 00009ABC      71 HEXCON DC XL4'9ABC'      00712001
146      0000D8 0020      72 HALFCON DC H'32'      00713001
147      000000      73      END ASMPGM      00720000
148      0000E0 00000004      74      =F'4'
149      0000E4 00000001      75      =F'1'
150      0000E8 C285878995      76      =C'Begin'
```

Fullword	32 bits - 4 bytes - of storage
Halfword	16 bits - 2 bytes - of storage

Other terms that might be useful: [Glossary of z/OS Terms and Abbreviations](#)

Use a calculator to convert between hexadecimal and decimal (if you need it)

4 USE TSO TEST FACILITY

The TSO TEST facility is used to execute a program with the added capability to stop the program execution at chosen address locations, to inspect changes in registers, and make changes in other address locations in the program.

The TSO Ready prompt is required to use the TSO TEST facility - if you are using ISPF, you will need to exit that environment, and get to TSO READY.

In this challenge, you can execute all the commands through the Zowe CLI in the same way as the REXX1 challenge.

To start things off, create an active TSO connection with Zowe CLI and capture the Address-Space Key (ASKEY) -

```
export ASKEY='zowe tso start address-space --sko`  
echo $ASKEY
```

If you are using a Windows Command shell, then the following will work better:

```
for /f "delims=" %i in ('zowe tso start address-space --sko -a FB3') do set ASKEY=%i  
echo %ASKEY%
```

You should see something similar to

```
Z#####-153-aacwaaat
```

where Z##### is your Z userid.

Once you have your TSO session key, you need to prepare TSO to recognise “short-hand” dataset names - for this challenge, it means that the TEST program, which will try and use a **LOAD** dataset, will use *your* LOAD dataset.

`zowe tso send address-space $ASKEY --data "PROFILE PREFIX(Z#####)"` (make sure to replace Z##### with your Z-userid!)

Send `zowe tso send address-space $ASKEY --data "PROFILE"` to see the resulting profile setting.

Now tell TSO to start the TEST environment with your assembled program

`zowe tso send address-space $ASKEY --data "test (asmpgm)"` - this will look in the LOAD dataset

For Windows Command Shell - `zowe tso send address-space %ASKEY% --data "test (asmpgm)"` - this will look in the LOAD dataset

Run through the following sets of commands by sending them as data to the TSO session - for example, for a command **LISTPSW**, use:

`zowe tso send address-space $ASKEY --data "LISTPSW"`

For Windows Command Shell - `zowe tso send address-space %ASKEY% --data "LISTPSW"`

(You should find it easy by pressing the cursor-up key and overtyping the last command)

NOTE: *if your zowe address-space times out between commands, you will see a message like :*

IZUG1126E: z/OSMF cannot correlate the request for key "Z#####-153-aacwaaat" with an active z/OS application session.

if that happens, you just need redo the command to start a new address-space, redo the command to start the TEST of asmpgm, then send the last `at` command that you got to, followed by `go`; that should take you back to where you were ...

- `LISTPSW`

Enter the above command to list the current address location of your ASMPGM under the **INSTR ADDR** column

- `l 0r:15r`

List the 'current' content of the 16 registers, 0-15. Observe register 15 (15r) is the address of ASMPGM which will be executed shortly. Register 14 is the address of TSO Ready environment which called TEST facility. When TEST facility is ended, control is returned to the address in register 14

- `l 15r% length(240)`

List the content starting at address location based on register 15. Observe the content is machine instructions and data from the compile output.

For example, the first byte, x'47' machine instruction (OPCODE) was created by the branch, B mnemonic at Stmt 18

- `l 15r% length(240) c`

List the content starting at address location based on register 15 in character format. Observe the literal text embedded in the assembler program. The first instruction in the program is to branch around the literal "eye catchers" embedded in the executable program

- `l +0 length(240)`

List the content starting 'relative' address location 0. Observe content is the same as listing 'absolute' address location starting at 11F00.

The program executable area can be referenced by either 'relative' or 'absolute' addressing

- `at +14 (l 0r:15r)`

at +14 is Stmt 22 - SETUP label and STM assembler mnemonic Stopping at +14 means the instruction at +14 is NOT executed until proceeding past +14 location

Enter the above command to stop program execution at 'relative' address location 14, then list the contents of all the registers (using `l 0r:15r`)

- `at +28 (l 0r:15r)` at +28 is Stmt 34 - L 2,=C'Begin' to load register 2 with content of character string 'Begin'.

Enter the above command to stop program execution at 'relative' address location 28, then list the contents of all the registers

- `at +2C (l 2r)`

at +2C is Stmt 35 - LA 2,=C'Begin' to load register 2 with address location of character string 'Begin'. Enter the above command to stop program execution at 'relative' address location 2C

- list content of register 2 which now contains the result of L 2,=C'Begin'

- `at +30 (l 2r;l +E8 length(5) c)`

at +30 is Stmt 37 - subtract register 2 from itself, SR 2,2, filling register 2 with zeroes Enter the above command to stop program execution at 'relative' address location 30, then

- list content of register 2 containing result of LA 2,=C'Begin'

Observe register 2 contains an 'absolute' address location, then the 'absolute' address location is listed. The content of the 'absolute' address location is displayed

- `at +32 (l 2r)`

at +32 is Stmt 38 - load register 2 with the number 4, L 2,=F'4'

Enter the above command to stop program execution at 'relative' address location 32, then - register 2 should now contain zeros as a result of the previously executed SR 2,2

- `at +36 (l 2r)`

at +36 is Stmt 39 - load register 3 with the number 1, L 3,=F'1'

Enter the above command to stop program execution at 'relative' address location 36, then check register 2. Observe register 2 changed from 0000000 to 0000004 as the result of L 2,=F'4'

- `at +3A (l 2r:3r)`

at +3A is both Stmt 41 and Stmt 42 where:

Stmt 41 has label LOOP at the +3A Defined Storage location, LOOP DS 0H

Stmt 42 results in add 1 to register 3, A 3,=F'1'

Enter the above command to stop program execution at 'relative' address location 3A

l 2r:3r

list content of registers 2 and 3 to that registers 2 equals 00000004 and register 3 equals 00000001 for the first time through the loop.

The program execution is designed to branch back to the +3A location, therefore, setting a stopping point at +3A will be executed again

- at +3E (l 2r:3r)

at +3E is Stmt 43 - branch on count, BCT 2,LOOP

Enter the above command to stop program execution at 'relative' address location 3E.

Each BCT execution results in decrement of register 2 by 1

The result of BCT 2,LOOP is to branch to label LOOP at +3A until register 2 is zero

The next stop is at +3A, storage address location of label LOOP adding 1 to register 3 until register 2 is zero

- Enter the following commands to be executed when register 3 is zero, terminating the loop

at +42

at +46

at +4A

at +84

- list content of registers 2 and 3, then branch to LOOP

Once register 2 is zero, then execution will stop at +42

You will need to enter `go` numerous times to proceed with program execution

When execution stops at +42, then

- `l 2r:3r` Register 2 and 3 show the final results of the loop Register 2 contains 00000000 as a result of BCT count decrement during each loop execution Register 3 contains 00000005 as a result of adding 1 during each loop execution

- `l +d8`

Location +D8 has the content of HALFCON HALFCON is defined as a halfword, 'H', therefore the first 4 bytes, 0020 represents the halfword

- `go` to continue execution

When execution stops at +46

```
l 3r Stmt 45 -      STOP1      LH 3,HALFCON      was just executed
```

Register 3 is a full word, therefore load halfword, 'LH', resulted in register 3 containing 00000020

```
l +D0
```

Content of value FULLCON is displayed, 'FFFFFFFF'

However, the assembler Define Constant statement is DC F'-1'

Why? FULLCON is an example of Two's Complement

Two's Complement Explained

For the purpose of computer arithmetic operation, FULLCON is -1

- `go`
- When execution stops at +4A

Stmt 46 - STOP2 A 3,FULLCON was just executed

l 3r

Register 3, 0000001F, is the result of adding x'20' to x'-1'

Before continuing with program execution, list value of address location for HEXCON (+D4)

l +D4

Observe address location currently shows 00009ABC

- `go`
- When execution stops at +84,

Stmt 47 - STOP3 ST 3,HEXCON was just executed ST 3,HEXCON stores the content of register 3 into address location HEXCON

l 3r

Register 3 still contains the result of the previous add operation unaltered by ST op code

l +D4

HEXCON address location was written over using the ST, Store, op code with content of Register 3

- `go`

The last `go` resulted in ASMPGM program termination returning control to caller TS0 TEST facility

Enter `end` to terminate TS0 TEST facility, returning control to TS0

5 MODIFY AND RE-COMPILE

Copy your SOURCE member ASMPGM to ASM2PGM:

1. Update the JCL to ensure output produces a load module called **ASM2PGM**
2. Use register 6 for all operations where register 2 was previously used
3. Use register 7 for all operations where register 3 was previously used
4. Initialize the register used to sum each add operation in the loop with a zero value
5. Execute the loop 10 times adding 5 to the register being used to sum each add operation
6. Submit and verify results via the Job output.

Remember that the register used to store the sum of the arithmetic add operation is a hexadecimal value.

6 VALIDATION AND NEXT STEPS

Submit the **CHKASM2** jcl located in **ZXP.PUBLIC.JCL** and review completion code for **0000**

Once you have completed the updates correctly and received credit, maybe you would like to get a bit more detail about how the IBM System 390 instruction set works, how data is represented for computation, and larger-scale applications are built.

Check out the [The z/Architecture Assembler language course series](#) which has just launched on the IBM Training platform.

Unit 1 - Numbering systems	▼
Unit 2 - z/Architecture concepts	▼
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