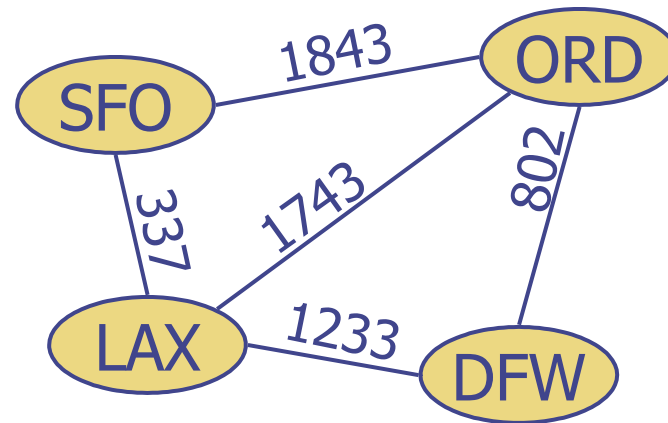


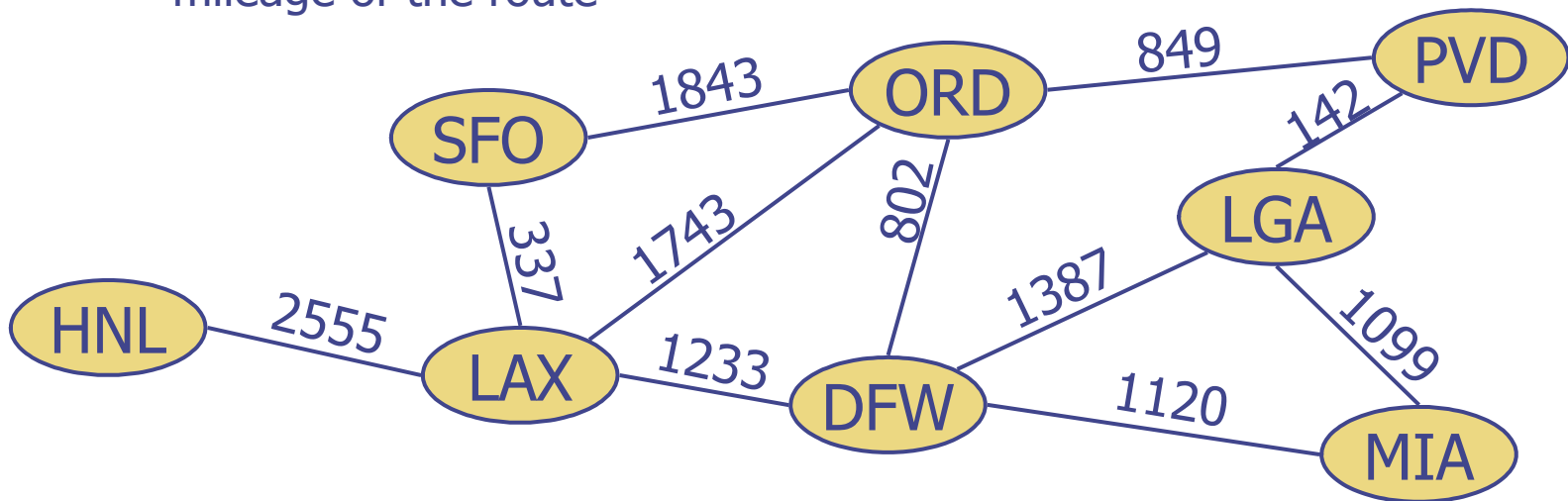
Presentation for use with the textbook **Data Structures and Algorithms in Java, 6<sup>th</sup> edition**, by M. T. Goodrich, R. Tamassia, and M. H. Goldwasser, Wiley, 2014

# Graphs



# Graphs

- A graph is a pair  $(V, E)$ , where
  - $V$  is a set of nodes, called **vertices**
  - $E$  is a collection of pairs of vertices, called **edges**
  - Vertices and edges are positions and store elements
- Example:
  - A vertex represents an airport and stores the three-letter airport code
  - An edge represents a flight route between two airports and stores the mileage of the route



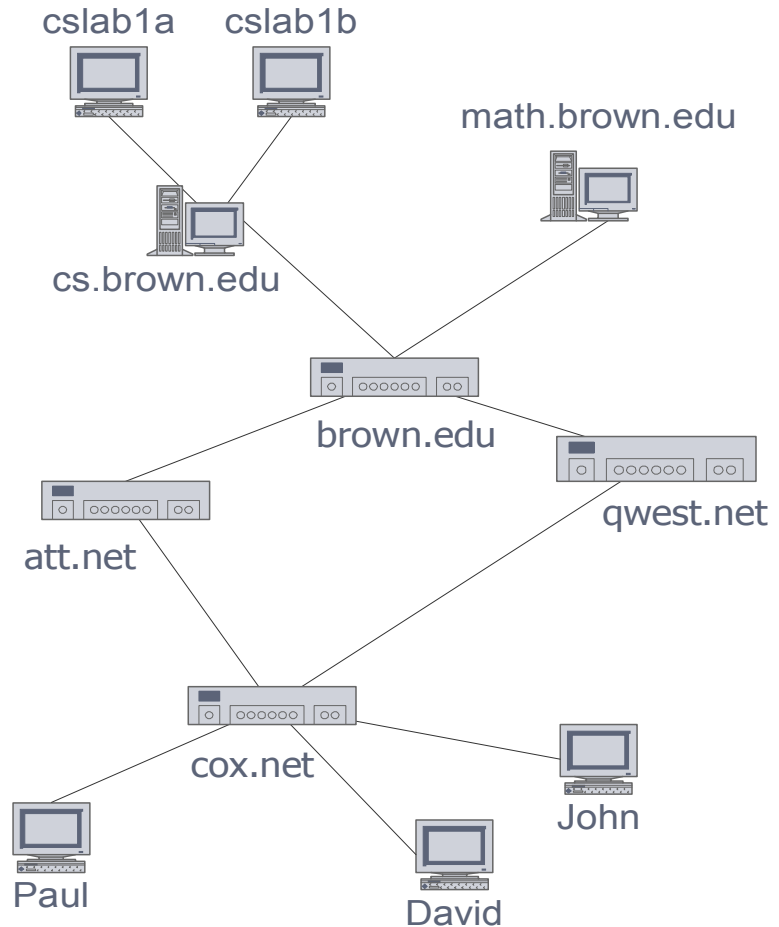
# Edge Types

- Directed edge
  - ordered pair of vertices  $(u, v)$
  - first vertex  $u$  is the origin
  - second vertex  $v$  is the destination
  - e.g., a flight
- Undirected edge
  - unordered pair of vertices  $(u, v)$
  - e.g., a flight route
- Directed graph
  - all the edges are directed
  - e.g., route network
- Undirected graph
  - all the edges are undirected
  - e.g., flight network



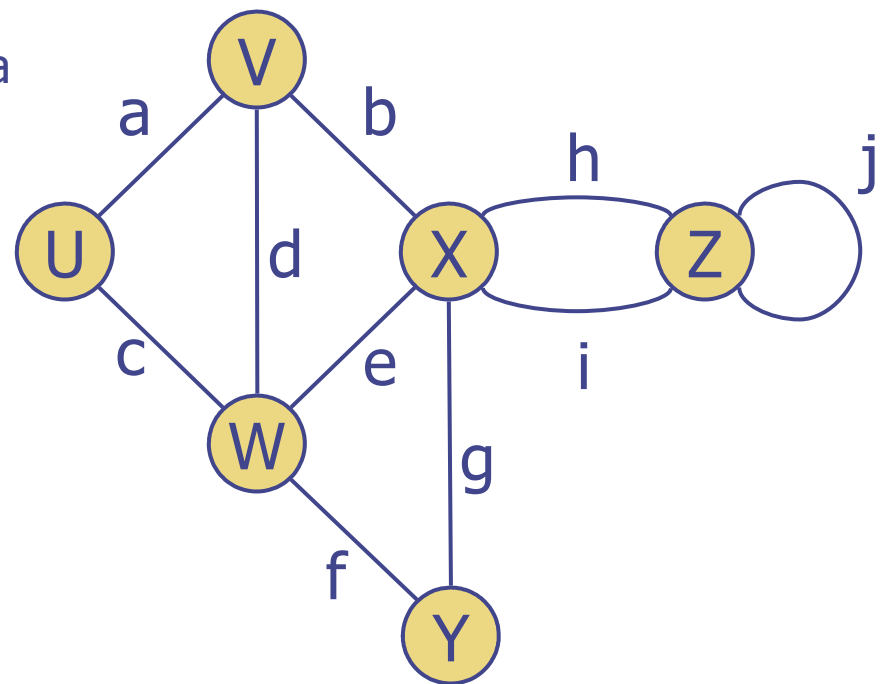
# Applications

- ❑ Electronic circuits
  - Printed circuit board
  - Integrated circuit
- ❑ Transportation networks
  - Highway network
  - Flight network
- ❑ Computer networks
  - Local area network
  - Internet
  - Web
- ❑ Databases
  - Entity-relationship diagram



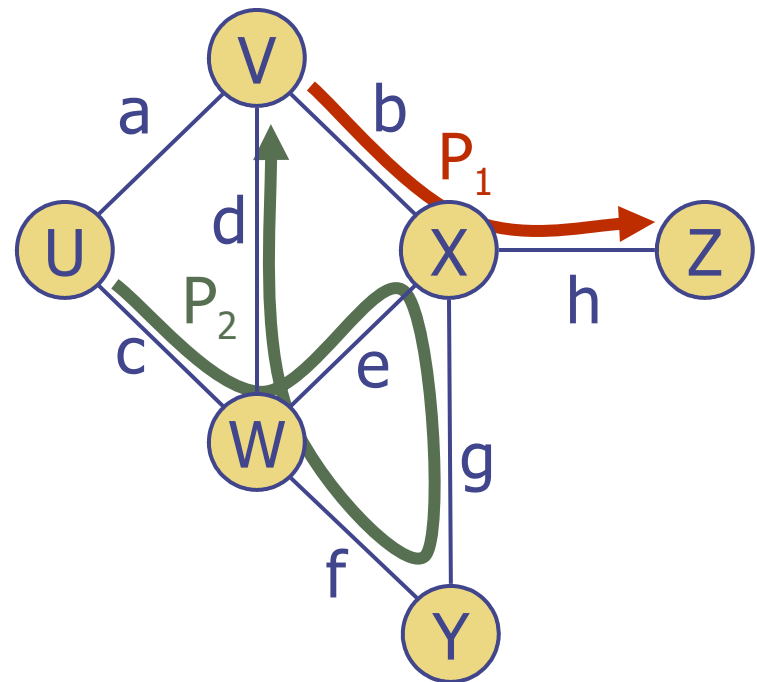
# Terminology

- End vertices (or endpoints) of an edge
  - U and V are the endpoints of a
- Edges incident on a vertex
  - a, d, and b are incident on V
- Adjacent vertices
  - U and V are adjacent
- Degree of a vertex
  - X has degree 5
- Parallel edges
  - h and i are parallel edges
- Self-loop
  - j is a self-loop



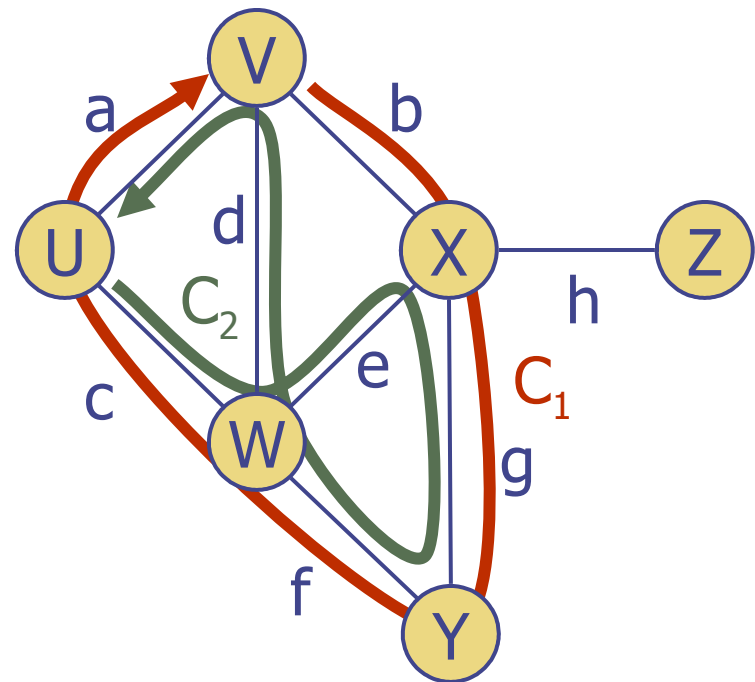
# Terminology (cont.)

- Path
  - sequence of alternating vertices and edges
  - begins with a vertex
  - ends with a vertex
  - each edge is preceded and followed by its endpoints
- Simple path
  - path such that all its vertices and edges are distinct
- Examples
  - $P_1 = (V, b, X, h, Z)$  is a simple path
  - $P_2 = (U, c, W, e, X, g, Y, f, W, d, V)$  is a path that is not simple



# Terminology (cont.)

- Cycle
  - circular sequence of alternating vertices and edges
  - each edge is preceded and followed by its endpoints
- Simple cycle
  - cycle such that all its vertices and edges are distinct
- Examples
  - $C_1 = (V, b, X, g, Y, f, W, c, U, a, \downarrow)$  is a simple cycle
  - $C_2 = (U, c, W, e, X, g, Y, f, W, d, V, a, \downarrow)$  is a cycle that is not simple



# Properties

## Property 1

$$\sum_v \deg(v) = 2m$$

Proof: each edge is  
counted twice

## Property 2

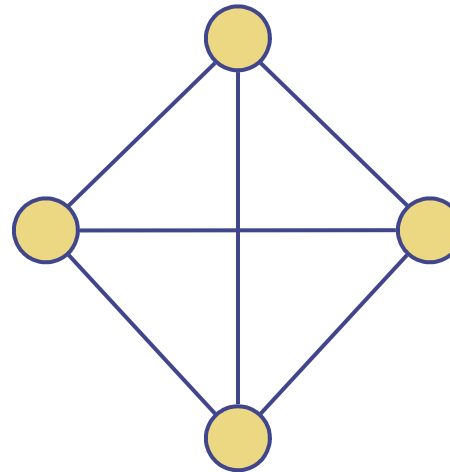
In an undirected graph  
with no self-loops and  
no multiple edges

$$m \leq n(n-1)/2$$

Proof: each vertex has  
degree at most  $(n-1)$

## Notation

$n$	number of vertices
$m$	number of edges
$\deg(v)$	degree of vertex $v$



## Example

- $n = 4$
- $m = 6$
- $\deg(v) = 3$

What is the bound for a  
directed graph?



# Vertices and Edges

- A **graph** is a collection of **vertices** and **edges**.
- We model the abstraction as a combination of three data types: Vertex, Edge, and Graph.
- A **Vertex** is a lightweight object that stores an arbitrary element provided by the user (e.g., an airport code)
  - We assume it supports a method, `element()`, to retrieve the stored element.
- An **Edge** stores an associated object (e.g., a flight number, travel distance, cost), retrieved with the `element( )` method.

# Graph ADT

`numVertices()`: Returns the number of vertices of the graph.

`vertices()`: Returns an iteration of all the vertices of the graph.

`numEdges()`: Returns the number of edges of the graph.

`edges()`: Returns an iteration of all the edges of the graph.

`getEdge( $u, v$ )`: Returns the edge from vertex  $u$  to vertex  $v$ , if one exists; otherwise return null. For an undirected graph, there is no difference between `getEdge( $u, v$ )` and `getEdge( $v, u$ )`.

`endVertices( $e$ )`: Returns an array containing the two endpoint vertices of edge  $e$ . If the graph is directed, the first vertex is the origin and the second is the destination.

`opposite( $v, e$ )`: For edge  $e$  incident to vertex  $v$ , returns the other vertex of the edge; an error occurs if  $e$  is not incident to  $v$ .

`outDegree( $v$ )`: Returns the number of outgoing edges from vertex  $v$ .

`inDegree( $v$ )`: Returns the number of incoming edges to vertex  $v$ . For an undirected graph, this returns the same value as does `outDegree( $v$ )`.

`outgoingEdges( $v$ )`: Returns an iteration of all outgoing edges from vertex  $v$ .

`incomingEdges( $v$ )`: Returns an iteration of all incoming edges to vertex  $v$ . For an undirected graph, this returns the same collection as does `outgoingEdges( $v$ )`.

`insertVertex( $x$ )`: Creates and returns a new Vertex storing element  $x$ .

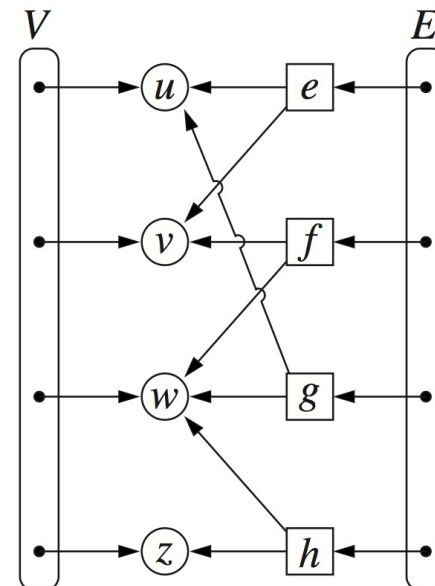
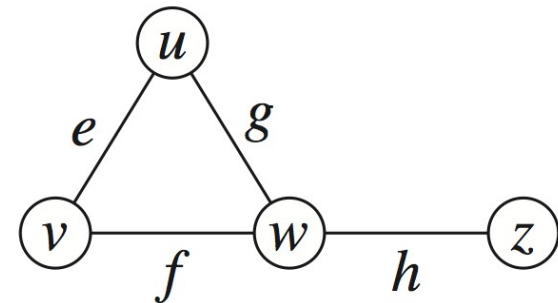
`insertEdge( $u, v, x$ )`: Creates and returns a new Edge from vertex  $u$  to vertex  $v$ , storing element  $x$ ; an error occurs if there already exists an edge from  $u$  to  $v$ .

`removeVertex( $v$ )`: Removes vertex  $v$  and all its incident edges from the graph.

`removeEdge( $e$ )`: Removes edge  $e$  from the graph.

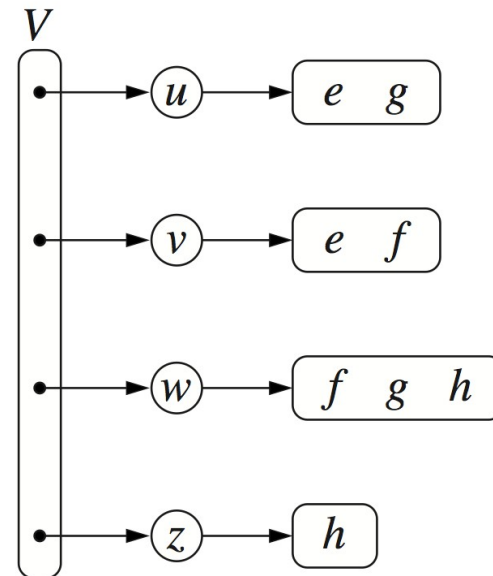
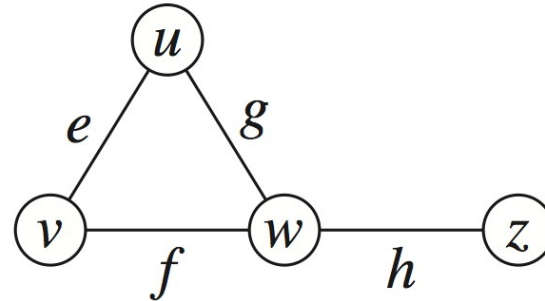
# Edge List Structure

- Vertex object
  - element
  - reference to position in vertex sequence
- Edge object
  - element
  - origin vertex object
  - destination vertex object
  - reference to position in edge sequence
- Vertex sequence
  - sequence of vertex objects
- Edge sequence
  - sequence of edge objects



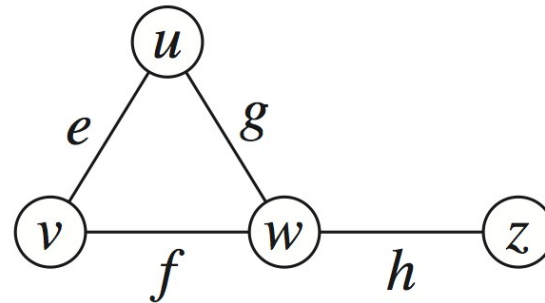
# Adjacency List Structure

- Incidence sequence for each vertex
  - sequence of references to edge objects of incident edges
- Augmented edge objects
  - references to associated positions in incidence sequences of end vertices



# Adjacency Matrix Structure

- ❑ Edge list structure
- ❑ Augmented vertex objects
  - Integer key (index) associated with vertex
- ❑ 2D-array adjacency array
  - Reference to edge object for adjacent vertices
  - Null for non adjacent vertices
- ❑ The “old fashioned” version just has 0 for no edge and 1 for edge



		0	1	2	3
$u$	→ 0		$e$	$g$	
$v$	→ 1	$e$		$f$	
$w$	→ 2	$g$	$f$		$h$
$z$	→ 3			$h$	

# Performance

<ul style="list-style-type: none"> <li>▪ <math>n</math> vertices, <math>m</math> edges</li> <li>▪ no parallel edges</li> <li>▪ no self-loops</li> </ul>	Edge List	Adjacency List	Adjacency Matrix
Space	$n + m$	$n + m$	$n^2$
<b>incidentEdges</b> ( $v$ )	$m$	$\text{deg}(v)$	$n$
<b>areAdjacent</b> ( $v, w$ )	$m$	$\min(\text{deg}(v), \text{deg}(w))$	1
<b>insertVertex</b> ( $o$ )	1	1	$n^2$
<b>insertEdge</b> ( $v, w, o$ )	1	1	1
<b>removeVertex</b> ( $v$ )	$m$	$\text{deg}(v)$	$n^2$
<b>removeEdge</b> ( $e$ )	1	1	1