Lab 5-2

1. Reference Code:

- Lab 4 and Lab 5-1
- The Bounded-Buffer Problem in the Lecture slides.

2. Overall Constraints:

For mutual exclusion locks (**Data Synchronization**): *spinlock* or *mutex*For synchronization of activities between threads (**Process Synchronization**): *semaphore* or *condition variable*

- **3. Lab Assignment**: Submission date on the moodle (Mar. 9, Mon. 11:59 pm), Online Grading during lecture sessions or Offline grading using the submitted files.
 - **Specific Constraints** (like the textbook):
 - 1) **n** buffers, each can hold one item (A global array: *char arr[n]*)
 - 2) Mutex *mtx* initialized to the value 1
 - 3) Semaphore *full* initialized to the value 0
 - 4) Semaphore *empty* initialized to the value n
 - In the main thread: create two threads (one is **producer thread** the other is **consumer thread**).
 - But, explicitly **consumer thread MUST** be created firstly (e.g., give 1 second sleep between the creation of the two threads).
 - The **consumer** thread prints the buffer values and the number of buffer size after the operation (e.g., [AA] [2], [A] [1] or [] [0] etc. **NOTE**: from 0 to buffer size), and the **producer** thread prints also the number of buffer size after the operation (e.g., A 1, AA 2 or AAA 3 etc. **NOTE**: from 1 to buffer size).
 - But, the bounded-buffer problem MUST be synchronized (Process Sync) with the
 protection of the critical section problem (Data Sync) in the shared array (i.e., the *arr*global array) by operating the buffer array and printing.
 - Constraint1: MUST use a mutex and two counting semaphores.
 - Constraint2: Semaphores MUST be implemented outside the critical section.

• Assumption:

For convenience, can use $\mathbf{n} = 10$, (but **MUST** be scalable up to any \mathbf{n} buffers). For convenience, create and test only 20 operations, (but **MUST** be scalable up to any N operations).

- You **MUST** implement the synchronization techniques **inside the created threads** (not the main thread) except initialization.
- Lastly, the expected **Output** could be like below. (**NOTE**: the results depend on the remainder section or speed. Please test the speed variations of the two threads using the **usleep() API**.

Expected Output 1: Producing speed is faster than consuming speed.

jurn@ZBookG4:~/work/cs332\$./lab5 con thread pro thread **A** 1 [0] **A** 1 AA 2 **AAA 3** AAAA 4 AAAAA 5 AAAAAA 6 [AAAAA] [5] AAAAAA 6 AAAAAAA 7 AAAAAAA 8 AAAAAAAA 9 **AAAAAAAAA** 10 [AAAAAAAA] [9] **AAAAAAAAA 10** [AAAAAAAA] [9] **AAAAAAAAA** 10 [AAAAAAAA] [9] AAAAAAAAA 10

Expected Output 1: Consuming speed faster than producing speed. (Always, the consumer will wait the first produced item)

```
jurn@ZBookG4:~/work/cs332$ ./lab5_2
con thread
pro thread
A 1
[] [0]
A 1
[] [0]
A 1
```

[] [0] A 1

[] [0]

A 1

[] [0]

A 1

[] [0]

A 1

[] [0]

A 1

[] [0] A 1

[] [0]

A 1

[] [0]