Lab 5-1

1. Reference Code:

- Lab 4 ex
- The Bounded-Buffer Problem in the Lecture slides.

2. Overall Constraints:

For mutual exclusion locks (**Data Synchronization**): *spinlock* or *mutex*For synchronization of activities between threads (**Process Synchronization**): *semaphore* or *condition variable*

3. Lab Assignment: Submission date on the moodle (Mar. 6, Fri. 1 pm), Online Grading during lecture sessions or Offline grading using the submitted files.

- Specific Constraints:
 - 1) **n** buffers, each can hold one item (A global array: *char arr[n]*)
 - 2) **Mutex mtx** initialized to the value 1
 - 3) A **condition variable** *full* initialization
 - 4) A **condition variable** *empty* initialization
- In the main thread: create two threads (one is **producer thread** the other is **consumer thread**).
 - But, explicitly **consumer thread MUST** be created firstly (e.g., give 1 second sleep between the creation of the two threads).
- The **consumer** thread prints the buffer values and the number of buffer size after the operation (e.g., [AA] [2], [A] [1] or [] [0] etc. **NOTE**: from 0 to buffer size-1), and the **producer** thread prints also the number of buffer size after the operation (e.g., A 1, AA 2 or AAA 3 etc. **NOTE**: from 1 to buffer size).
 - But, the bounded-buffer problem MUST be synchronized (Process Sync) with the
 protection of the critical section problem (Data Sync) in the shared array (i.e., the *arr*global array) by operating the buffer array and printing.
 - Constraint1: MUST use a mutex and two condition variables.
 - Constraint2: Semaphores MUST be implemented outside the critical section.

• Assumption:

For convenience, can use $\mathbf{n} = 10$, (but **MUST** be scalable up to any \mathbf{n} buffers). For convenience, create and test only 20 operations, (but **MUST** be scalable up to any N operations).

- You **MUST** implement the synchronization techniques **inside the created threads** (not the main thread) except initialization.
- Lastly, the expected **Output** could be like below. (**NOTE**: the results depend on the remainder section or speed. Please test the speed variations of the two threads using the **usleep() API** or **for loop**.

Expected Output 1: Producing speed is faster than consuming speed.

```
jurn@ZBookG4:~/work/cs332$ ./lab5 1
con thread
pro thread
A 1
[0]
A 1
AA 2
AAA 3
AAAA 4
AAAAA 5
AAAAAA 6
[AAAAA] [5]
AAAAAA 6
AAAAAAA 7
AAAAAAAA 8
AAAAAAAA 9
AAAAAAAAA 10
[AAAAAAAA] [9]
AAAAAAAAA 10
[AAAAAAAA] [9]
AAAAAAAAA 10
[AAAAAAAA] [9]
AAAAAAAAA 10
```

Expected Output 1: Consuming speed is faster than producing speed. (Always, the consumer will wait the first produced item)

```
jurn@ZBookG4:~/work/cs332$ ./lab5_2
con thread
pro thread
A 1
[] [0]
A 1
[] [0]
```

A 1

[] [0]

A 1

[] [0] A 1

[] [0]

A 1

[] [0]

A 1

[] [0]

A 1

[] [0]

A 1

[] [0] A 1

[] [0]