APIC

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APIC ("Advanced Programmable Interrupt Controller") is the updated Intel standard for the older PIC. It is used in multiprocessor systems and is an integral part of all recent Intel (and compatible) processors. The APIC is used for sophisticated interrupt redirection, and for sending interrupts between processors. These things weren't possible using the older PIC specification.

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Detection

The CPUID.01h:EDX[bit 9] flag specifies whether a CPU has a built-in local APIC. You can find all of the APICs on a system (both local and IO APICs) by parsing the MADT.

Local APIC and IO-APIC

In an APIC-based system, each CPU is made of a "core" and a "local APIC". The local APIC is responsible for handling cpu-specific interrupt configuration. Among other things, it contains the *Local Vector Table (LVT)* that translates events such as "internal clock" and other "local" interrupt sources into a interrupt vector (e.g. LocalINT1 pin could be raising an NMI exception by storing "2" in the corresponding entry of the LVT).

More information about the local APIC can be found in Chapter 10 of the Intel System Programming Guide, Vol 3A Part 1 (https://www.intel.com/content/dam/www/public/us/en/documents/manuals/64-ia-32-architectures-software-developer-vol-3a-part-1-manual.pdf).

In addition, there is an I/O APIC (e.g. intel 82093AA) that is part of the chipset and provides multi-processor interrupt management, incorporating both static and dynamic symmetric interrupt distribution across all processors. In systems with multiple I/O subsystems, each subsystem can have its own set of interrupts.

Each interrupt pin is individually programmable as either edge or level triggered. The interrupt vector and interrupt steering information can be specified per interrupt. An indirect register accessing scheme optimizes the memory space needed to access the I/O APIC's internal registers. To increase system flexibility when assigning memory space usage, the I/O APIC's two-register memory space is relocatable, but defaults to 0xFEC00000.

The Intel standards for the APIC can be found on the Intel site under the category "Multiprocessor Specification", or simply this PDF file

(http://web.archive.org/web/20070112195752/http://developer.intel.com/design/pentium/datashts/24201606.pdf)

Inter-Processor Interrupts

Inter-Processor Interrupts (IPIs) are generated by a local APIC and can be used as basic signaling for scheduling coordination, multi-processor bootstrapping, etc.

Local APIC configuration

The local APIC is enabled at boot-time and can be disabled by clearing bit 11 of the IA32_APIC_BASE Model Specific Register (MSR) (see example below, this only works on CPUs with family >5, as the Pentium does not have such MSR). The CPU then receives its interrupts directly from a 8259-compatible PIC. The Intel Software Developer's Manual, however states that, once you have disabled the local APIC through IA32_APIC_BASE you can't enable it anymore until a complete reset. The I/O APIC can also be configured to run in legacy mode so that it emulates an 8259 device.

The local APIC's registers are memory-mapped in physical page FEE00xxx (as seen in table 8-1 of Intel P4 SPG). This address is the same for each local APIC that exists in a configuration, meaning you are only able to directly access the registers of the local APIC of the core that your code is currently executing on. Note that there is a MSR that specifies the actual APIC base (only available on CPUs with family >5). The MADT contains the local APIC base and on 64-bit systems it may also contain a field specifying a 64-bit base address override which you ought to use instead. You can choose to leave the Local APIC base just where you find it, or to move it at your pleasure. **Note:** I don't think you can move it any further than the 4th Gb.

To enable the Local APIC to receive interrupts it is necessary to configure the "Spurious Interrupt Vector Register". The correct value for this field is the IRQ number that you want to map the spurious interrupts to within the lowest 8 bits, and the 8th bit set to 1 to actually enable the APIC (see the specification for more details). You should choose an interrupt number that has its lowest 4 bits set and is above 32 (as you might guess); easiest is to use 0xFF. This is important on some older processors because the lowest 4 bits for this value must be set to 1 on these.

Disable the 8259 PIC properly. This is nearly as important as setting up the APIC. You do this in two steps: masking all interrupts and remapping the IRQs. Masking all interrupts disables them in the PIC. Remapping is what you probably already did when you used the PIC: you want interrupt requests to start at 32 instead of 0 to avoid conflicts with the exceptions. You should then avoid using these interrupt vectors for other purposes. This is necessary because even though you masked all interrupts on the PIC, it could still give out spurious interrupts which will then be misinterpreted from your kernel as exceptions.

Here are some code examples on setting up the APIC:

```
#define IA32_APIC_BASE_MSR 0x1B
#define IA32_APIC_BASE_MSR_BSP 0x100 // Processor is a BSP
#define IA32_APIC_BASE_MSR_ENABLE 0x800

/** returns a 'true' value if the CPU supports APIC
  * and if the local APIC hasn't been disabled in MSRs
  * note that this requires CPUID to be supported.
  */
bool check_apic() {
    uint32_t eax, edx;
    cpuid(1, &eax, &edx);
    return edx & CPUID_FEAT_EDX_APIC;
}

/* Set the physical address for local APIC registers */
void cpu_set_apic_base(uintptr_t apic) {
    uint32_t edx = 0;
    uint32_t eax = (apic & 0xfffff000) | IA32_APIC_BASE_MSR_ENABLE;
```

```
#ifdef PHYSICAL MEMORY EXTENSION
   edx = (apic >> 32) \& 0x0f:
#endif
   cpuSetMSR(IA32 APIC BASE MSR, eax, edx);
/**
 * Get the physical address of the APIC registers page
 * make sure you map it to virtual memory ;)
uintptr t cpu get apic base() {
   uint32 t eax, edx;
   cpuGetMSR(IA32_APIC_BASE_MSR, &eax, &edx);
#ifdef __PHYSICAL_MEMORY_EXTENSION_
  return (eax & 0xffffff000) ((edx & 0x0f) << 32);
  return (eax & 0xfffff000):
#endif
void enable apic() {
    /* Hardware enable the Local APIC if it wasn't enabled */
    cpu set apic base(cpu get apic base());
    /* Set the Spurious Interrupt Vector Register bit 8 to start receiving interrupts */
   write_reg(0xF0, ReadRegister(0xF0) \mid 0x100);
```

Local APIC and x86 SMM Attacks

The APIC was introduced to the core Intel processor architecture skeleton in Intel's 82489DX discrete chip (https://4donline.ihs.com/images/VipMasterIC/IC/INTL/INTLD047/INTLD047-2-1259.pdf? hkey=EF798316E3902B6ED9A73243A3159BB0) in a similar time period as System Management Mode was introduced to operating systems. In original architecture, the APIC could not be mapped to memory, and it wasn't until later changes that it became mappable.

As System Management Mode's memory (SMRAM) is given a protected range of memory (which can vary from system to system), it is possible to map the APIC memory location into the SMRAM. The result of this is that SMM memory is pushed outside its protected range and exposed to lesser-privileged permission rings. Using this method, attackers can leverage their permissions using System Management Mode, which is protected from all rings above -2.

In newer generation Intel processors (starting with the Intel Atom (https://en.wikipedia.org/wiki/Intel_Atom) in 2013), this has been taken into account. An undocumented check is performed against the System Management Range Registers when the APIC is relocated to memory. This check ensures that the APIC does not overlap with the SMRAM. **However**, this relies on the SMRR to be configured correctly. Otherwise, this mitigation will not work properly and attackers will still be able to user the attack.

Local APIC registers

The local APIC registers are memory mapped to an address that can be found in the MP/MADT tables. Make sure you map these to virtual memory if you are using paging. Each register is 32 bits long, and expects to be written and read as a 32 bit integer. Although each register is 4 bytes, they are all aligned on a 16 byte boundary.

List of local APIC registers (TODO: Add descriptions for all registers):

| Offset | Register name | Read/Write permissions | |
|-------------|-------------------|------------------------|--|
| 000h - 010h | Reserved | | |
| 020h | LAPIC ID Register | Read/Write | |
| | | | |

| 03011 | Date Version Register | rteda orny |
|-------------|---|------------|
| 040h - 070h | Reserved | |
| 080h | Task Priority Register (TPR) | Read/Write |
| 090h | Arbitration Priority Register (APR) | Read only |
| 0A0h | Processor Priority Register (PPR) | Read only |
| 0B0h | EOI register | Write only |
| 0C0h | Remote Read Register (RRD) | Read only |
| 0D0h | Logical Destination Register | Read/Write |
| 0E0h | Destination Format Register | Read/Write |
| 0F0h | Spurious Interrupt Vector Register | Read/Write |
| 100h - 170h | In-Service Register (ISR) | Read only |
| 180h - 1F0h | Trigger Mode Register (TMR) | Read only |
| 200h - 270h | Interrupt Request Register (IRR) | Read only |
| 280h | Error Status Register | Read only |
| 290h - 2E0h | Reserved | |
| 2F0h | LVT Corrected Machine Check Interrupt (CMCI) Register | Read/Write |
| 300h - 310h | Interrupt Command Register (ICR) | Read/Write |
| 320h | LVT Timer Register | Read/Write |
| 330h | LVT Thermal Sensor Register | Read/Write |
| 340h | LVT Performance Monitoring Counters Register | Read/Write |
| 350h | LVT LINTO Register | Read/Write |
| 360h | LVT LINT1 Register | Read/Write |
| 370h | LVT Error Register | Read/Write |
| 380h | Initial Count Register (for Timer) | Read/Write |
| 390h | Current Count Register (for Timer) | Read only |
| 3A0h - 3D0h | Reserved | |

Read only

Read/Write

EOI Register

3E0h

3F0h

030h

LAPIC Version Register

Write to the register with offset 0xB0 using the value 0 to signal an end of interrupt. A non-zero value causes a general protection fault.

Local Vector Table Registers

Reserved

There are some special interrupts that the processor and LAPIC can generate themselves. While external interrupts are configured in the I/O APIC, these interrupts must be configured using registers in the LAPIC. The most interesting registers are: 0x320 = 1apic timer, 0x350 = 1apic timer, 0x360 = 1apic timer

Register format:

| Bits 0-7 | The vector number |
|--------------------------------|--------------------------------|
| Bits 8-11 (reserved for timer) | 100b if NMI |
| Bit 12 | Set if interrupt pending. |
| Bit 13 (reserved for timer) | Polarity, set is low triggered |
| | |

Divide Configuration Register (for Timer)

| Bit 14 (reserved for timer) | Remote IRR |
|-----------------------------|--------------------------------------|
| Bit 15 (reserved for timer) | trigger mode, set is level triggered |
| Bit 16 | Set to mask |
| Bits 17-31 | Reserved |

The vector number, or starting page number for SIPIs

Spurious Interrupt Vector Register

The offset is 0xF0. The low byte contains the number of the spurious interrupt. As noted above, you should probably set this to 0xFF. To enable the APIC, set bit 8 (or 0x100) of this register. If bit 12 is set then EOI messages will not be broadcast. All the other bits are currently reserved.

Interrupt Command Register

Bits

0 - 7

The interrupt command register is made of two 32-bit registers; one at 0x300 and the other at 0x310. It is used for sending interrupts to different processors. The interrupt is issued when 0x300 is written to, but not when 0x310 is written to. Thus, to send an interrupt command one should first write to 0x310, then to 0x300. At 0x310 there is one field at bits 24-27, which is local APIC ID of the target processor (for a physical destination mode). Here is how 0x300 is structured:

| Bits 8- 10 | The destination mode. 0 is normal, 1 is lowest priority, 2 is SMI, 4 is NMI, 5 can be INIT or INIT level deassert, 6 is a SIPI. |
|-------------------|---|
| Bit 11 | The destination mode. Clear for a physical destination, or set for a logical destination. If the bit is clear, then the destination field in 0x310 is treated normally. |
| Bit 12 | Delivery status. Cleared when the interrupt has been accepted by the target. You should usually wait until this bit clears after sending an interrupt. |
| Bit 13 | Reserved |
| Bit 14 | Clear for INIT level de-assert, otherwise set. |
| Bit 15 | Set for INIT level de-assert, otherwise clear. |
| Bits 18- 19 | Destination type. If this is > 0 then the destination field in 0x310 is ignored. 1 will always send the interrupt to itself, 2 will send it to all processors, and 3 will send it to all processors aside from the current one. It is best to avoid using modes 1, 2 and 3, and stick with 0. |
| Bits 20- | Reserved |

IO APIC Configuration

The IO APIC uses two registers for most of its operation - an address register at IOAPICBASE+0 and a data register at IOAPICBASE+0x10. All accesses must be done on 4 byte boundaries. The address register uses the bottom 8 bits for register select. Here is some example code that illustrates this:

```
uint32_t cpuReadIoApic(void *ioapicaddr, uint32_t reg)
{
   uint32_t volatile *ioapic = (uint32_t volatile *)ioapicaddr;
   ioapic[0] = (reg & 0xff);
   return ioapic[4];
}
```

```
void cpuWriteIoApic(void *ioapicaddr, uint32_t reg, uint32_t value)
{
  uint32_t volatile *ioapic = (uint32_t volatile *)ioapicaddr;
  ioapic[0] = (reg & 0xff);
  ioapic[4] = value;
}
```

Note the use of the volatile keyword. This prevents a compiler like Visual C from reordering or optimizing away the memory accesses, which would be a Bad Thing $^{\text{m}}$. The volatile keyword is put before the '*' sign. It means that the *value pointed to* is volatile, not the pointer itself.

IO APIC Registers

Bits

0-7

Bit

Bit

15

Bit

16

Bits

Using the methods described above, the following registers are accessible.

Interrupt vector. Allowed values are from 0x10 to 0xFE.

Trigger mode. 0 = Edge sensitive, 1 = Level sensitive.

| 0x00 | Get/set the IO APIC's id in bits 24-2/. All other bits are reserved. |
|----------------------|--|
| I () X () I | Get the version in bits 0-7. Get the maximum amount of redirection entries in bits 16-23. All other bits are reserved. Read only. |
| 0x02 | Get the arbitration priority in bits 24-27. All other bits are reserved. Read only. |
| | Contains a list of redirection entries. They can be read from and written to. Each entries uses two addresses, e.g. 0x12 and 0x13. |

Here is what a redirection entry looks like.

| | Type of delivery mode. $0 = Normal$, $1 = Low priority$, $2 = System management interrupt$, $4 = Non maskable interrupt$, $5 = INIT$, $7 = External$. All others are reserved. |
|-----|---|
| | Destination mode. Affects how the destination field is read, 0 is physical mode, 1 is logical. If the Destination Mode of this entry is Physical Mode, bits 56-59 contain an APIC ID. |
| Rit | |

Bit 12 Set if this interrupt is going to be sent, but the APIC is busy. Read only.

Polarity of the interrupt. 0 = High is active, 1 = Low is active.

Bit Used for level triggered interrupts only to show if a local APIC has received the interrupt (= 1), or has sent an EOI (= 0). Read only.

an EOI (= 0). Read only.

Interrupt mask. Stops the interrupt from reaching the processor if set.

For more information, check out chapter 3 of the I/O APIC datasheet

(http://web.archive.org/web/20161130153145/http://download.intel.com/design/chipsets/datashts/29056601.pdf)

The redirection table allows you to choose which external interrupts are sent to which processors and with which interrupt vectors. When choosing the processors you should consider: spreading out the workload between the processors, avoiding processors in a low-power state, and avoiding throttled processors. When choosing the

interrupt vectors you should remember that interrupts 0x00 to 0x1F are reserved for internal processor exceptions, the interrupts you remapped the PIC to may receive spurious interrupts, that 0xFF is probably where you put the APIC spurious interrupt, and that the upper 4 bits of an interrupt vector indicate its priority.

Logical Destination Mode

Logical destination mode uses an 8-bit logical APIC ID, contained in the LDR (logical destination register, unique to each APIC). All APICs compare their local ID to the destination code sent with the interrupt. This allows to target a group of processors by programming them with the same logical APIC ID.

The LDR is formatted as follows

| Bits 0-23 | Reserved. | | |
|--------------|------------------|---|---|
| Bits 24- | Flat model | Bitmap of target processors (bit identifies single processor; supports a maximum of 8 local APIC units) | |
| | Cluster model | Bits 24-27 | Local APIC address (identifies the specific processor in a group) |
| | | Bits 28-31 | Cluster address (identifies a group of processors) |

The DFR (destination format register) specifies Flat or Cluster model and is structured as follows

| Bits 0-27 | Reserved. |
|------------|---|
| Bits 28-31 | Model (1111b = Flat model, 0000b = Cluster model) |

'Don't use cluster mode addressing, especially "hierarchical cluster mode". AFAIK it was intended for large NUMA systems, where there's a "node controller" for each NUMA domain that forwarded interrupts to CPUs within that NUMA domain (with a seperate APIC bus for each NUMA domain). Unless your chipset has these "node controllers" (or "cluster managers" as Intel calls them) it won't work, and no modern computers have them (AFAIK there are only a few obscure Pentium III/P6 NUMA systems that ever did). You want to use "flat model" for normal SMP and for most NUMA systems (including AMD's).' (Brendan (http://forum.osdev.org/viewtopic.php? f=1&t=14808&start=17))

More info can be found in "Pentium Processor System Architecture. Chapter 15: The APIC" (https://books.google.nl/books?id=TVzjEZg1--YC&printsec=frontcover)

SIPI Sequence

The INIT-SIPI sequence is done by the operating system after the firmware already gives the first INIT signal. This example code from the [Silcos Kernel http://wiki.osdev.org/Silcos_Kernel]

// NOTE: ICRLow and ICRHigh are types in the Silcos kernel. If your code uses direct bit // manipulations you must replace some code with bit operations.

See Also

Articles

- 8259 PIC
- IOAPIC
- APIC timer

Threads

- APIC timer (http://www.osdev.org/phpBB2/viewtopic.php?t=10686)
- Mapping the I/O APIC (http://www.osdev.org/phpBB2/viewtopic.php?t=11529)
- Brendan gives some general info on the APIC and implementing it. (http://www.osdev.org/phpBB2/viewtopic.php?p=107868#107868)

External Links

- original I/O APIC specification/datasheet (http://www.intel.com/design/chipsets/datashts/290566.htm)
- updated I/O APIC specification/datasheet (http://developer.intel.com/design/chipsets/specupdt/290710.htm)
- Volume 3A:System Programming Guide, Part 1,manuals has a chapter on the APIC (http://www.intel.com/products/processor/manuals/)
- Volume 3A:System Programming Guide, Chapter 10.4 for further reading about the LAPIC (http://www.intel.com/products/processor/manuals/)
- Advanced Programmable Interrupt Controller by Mike Rieker (http://web.archive.org/web/20140308064246/http://www.osdever.net/tutorials/pdf/apic.pdf)
- "The Importance of Implementing APIC-Based Interrupt Subsystems on Uniprocessor PCs". Microsoft. 18 September 2010
 - (http://web.archive.org/web/20100918084750/http://www.microsoft.com/whdc/archive/apic.mspx)
- Pentium Processor System Architecture (https://books.google.nl/books?id=TVzjEZg1--YC&printsec=frontcover)

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